

# A Note on the Two Type Procurement Auction Equilibrium

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The goal of this note is extend the baseline model discussed in Bolotny and Vasserman (2019) to allow for two-dimensional types.

We begin by defining the primitives of the model in order of importance for the exercise.

Bidders indexed by  $i$  are described types consisting of two parameters:

- A “cost” parameter:  $\alpha^i$ 
  - This captures the relative cost-efficiency of bidder  $i$ .
- A “risk” parameter:  $\gamma^i$ 
  - This captures the extent to which bidder  $i$  dislike risks.

The role of  $\gamma$  is mediated by a “utility” function that models how much a bidder values a given bet. In particular, we will use:

$$u(\pi) = 1 - \exp(-\gamma\pi)$$

To solidify what this means, suppose that a bidder faces a bet: win \$100 with probability 0.5 or \$0 with probability 0.5. If the bidder has  $\gamma = 0.01$ , then the expectation of his utility for the bet is:

$$\begin{aligned} 0.5 \cdot u(100) + 0.5 \cdot u(0) &= 0.5 \cdot (1 - \exp(-\frac{100}{100})) + 0.5 \cdot (1 - \exp(-\frac{0}{100})) \\ &= 0.5 \cdot (1 - \exp(-1)) + 0.5 \cdot (0) \approx 0.316 \end{aligned}$$

By contrast, if he were able to get \$50 for sure, he would value this at

$$u(50) = 1 - \exp(-\frac{50}{100}) \approx 0.393$$

so you can see that the bidder “dislikes” risk. In general, the higher that  $\gamma$  is, the larger the discrepancy between the value of a “bet” and a “sure thing” of equal expected value.

The rest of the parameters in the model explain the way that a bidder of type  $(\alpha^i, \gamma^i)$  values a given bid that he is considering to submit to the auction.

This is described in the problem definition section just below.

For each of  $t = 1, \dots, T$  materials (items) that procurement project will require, there is:

- $q_t^e$ :
  - The DOT engineer’s estimate of the quantity of item  $t$  that will be needed for the project
- $q_t^b$ :
  - The bidder’s estimated quantity of item  $t$  that will be used
- $\sigma_t^2$ 
  - The variance of the bidder’s estimate of  $q_t^b$
- $c_t$ :
  - The market unit rate for item  $t$

Note that I often use “DOT” to refer to the government org that runs the auction – DOT stands for Department of Transportation.

## The Problem Definition

We are interested in solving for the “equilibrium bidding” function that maps each possible type  $(\tilde{\alpha}, \tilde{\gamma})$  to a “score”  $s(\tilde{\alpha}, \tilde{\gamma})$ .

We can think the “score” function as a map:

$$s : [\underline{\alpha}, \bar{\alpha}] \times [\underline{\gamma}, \bar{\gamma}] \rightarrow [\underline{S}, \bar{S}]$$

where each of  $[\underline{\alpha}, \bar{\alpha}]$ ,  $[\underline{\gamma}, \bar{\gamma}]$  and  $[\underline{S}, \bar{S}]$  are subsets of  $\mathbb{R}_+$ .

In order for a prospective “score” function to be an equilibrium, it needs maximize the “expected utility” from participating in the auction for each possible bidder type  $(\tilde{\alpha}, \tilde{\gamma})$  — let’s call this  $EU(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))$ .

This is the product of the (expected) utility that the bidder would get from completing the project if he wins (given his bid) times the probability that he wins.

$$EU(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma})) = E[u(\pi(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))) \mid \text{win}] \times \Pr(\text{win} \mid s(\tilde{\alpha}, \tilde{\gamma}))$$

Let’s break this down. The rules of the auction are that the bidder with the **lowest** score wins the auction. Thus, we can write

$$\Pr(\text{win} \mid s(\tilde{\alpha}, \tilde{\gamma})) = \Pr(s(\tilde{\alpha}, \tilde{\gamma}) < \text{all other bidders' scores})$$

Let’s assume there are only two bidders to make things simple (so that “all other bidders” is just the one other bidder). Our construction requires that **every** type of bidder uses the same score function  $s(\cdot, \cdot)$ .

Thus the probability that  $s(\tilde{\alpha}, \tilde{\gamma})$  is lower than the other bidder’s score is the probability that the other bidder  $j$  drew a type  $(\alpha^j, \gamma^j)$  such that  $s(\tilde{\alpha}, \tilde{\gamma}) < s(\alpha^j, \gamma^j)$ .

How might we find this distribution? We will assume (as a primitive of the model – not something that we compute) that each of  $\tilde{\alpha}$  and  $\tilde{\gamma}$  are drawn independently according to some probability distributions. For instance a simple case would have them distributed uniformly on their domains:

$$\alpha \sim U[\underline{\alpha}, \bar{\alpha}] \text{ and } \gamma \sim U[\underline{\gamma}, \bar{\gamma}].$$

Let’s write  $f_\alpha(\tilde{\alpha})$  and  $F_\alpha(\tilde{\alpha})$  for the pdf and cdf of the  $\alpha$  distribution evaluated at  $\tilde{\alpha}$ , respectively, and  $f_\gamma(\tilde{\gamma})$  and  $F_\gamma(\tilde{\gamma})$  as the pdf and cdf for  $\gamma$ .

So (plugging in) for a proposed score function  $s(\cdot)$ , the expected utility of participation can be written:

$$EU(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma})) = E[u(\pi(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))) \mid \text{win}] \times \int_{\underline{\alpha}}^{\bar{\alpha}} \int_{\underline{\gamma}}^{\bar{\gamma}} [\mathbf{1}\{s(\tilde{\alpha}, \tilde{\gamma}) < s(\alpha^j, \gamma^j)\}] f(\alpha^j) f(\gamma^j) d\alpha^j d\gamma^j$$

Note that the probability of winning critically depends on the choice of the scoring function. For a given scoring function, the probability of winning is easy to compute.

Now the second part – the expected utility of profits conditional on winning. The math for this is the same as in the simpler case detailed in the paper, so I won’t replicate it here — instead I’ll present the formula and explain how it relates to the overall problem.

$$E[u(\pi(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))) \mid \text{win}] = \left( 1 - \exp \left( -\tilde{\gamma} \sum_{t=1}^T q_t^b (b_t^*(s(\tilde{\alpha}, \tilde{\gamma})) - \tilde{\alpha} c_t) - \frac{\tilde{\gamma} \sigma_t^2}{2} (b_t^*(s(\tilde{\alpha}, \tilde{\gamma})) - \tilde{\alpha} c_t)^2 \right) \right)$$

This formula includes a new object: the unit bid  $b_t^*(s(\tilde{\alpha}, \tilde{\gamma}))$ . What is this?

Although only the score determines who wins the auction, in reality, bidders don't actually submit "scores" directly – instead they submit a unit bid  $b_t$  for each item  $t$  in the procurement project. The "score" for a bidder submitting the vector  $\{b_1, \dots, b_T\}$  of bids, is then computed by multiplying by the DOT's estimate of the quantity of each item that will be needed and summing:

$$\text{score} = \sum_t b_t q_t^e.$$

But a result from the paper is that we don't have to think too hard about how to choose the unit bids in equilibrium — it is sufficient to solve for an equilibrium **score** function as we've discussed before, under the assumption that for any score (and  $(\tilde{\alpha}, \tilde{\gamma})$  type), bidders will deterministically choose the unit bids  $b_t^*(s(\tilde{\alpha}, \tilde{\gamma}))$  to maximize  $E[u(\pi(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))) \mid \text{win}]$ .

That is:

$$b_t^*(s(\tilde{\alpha}, \tilde{\gamma})) = \arg \max_{\{b_t\}} \left[ 1 - \exp \left( -\tilde{\gamma} \sum_{t=1}^T q_t^b (b_t - \tilde{\alpha} c_t) - \frac{\tilde{\gamma} \sigma_t^2}{2} (b_t - \tilde{\alpha} c_t)^2 \right) \right]$$

$$\text{s.t. } \sum_{t=1}^T b_t q_t^e = s(\tilde{\alpha}, \tilde{\gamma})$$

and  $b_t \geq 0$  for each  $t$

Note: this really this should be  $b_t^*(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))$  but I'm suppressing the other parameters since they're implied.

The optimization program above can be rewritten as a fairly standard constrained quadratic program:

$$b_t^*(s(\tilde{\alpha}, \tilde{\gamma})) = \arg \max_{\{b_t\}} \left[ \sum_{t=1}^T \underbrace{q_t^b (b_t - \tilde{\alpha} c_t)}_{\text{Expectation of Profits}} \underbrace{- \frac{\tilde{\gamma} \sigma_t^2}{2} (b_t - \tilde{\alpha} c_t)^2}_{\text{Variance of Profits}} \right]$$

$$\text{s.t. } \sum_{t=1}^T b_t q_t^e = s(\tilde{\alpha}, \tilde{\gamma})$$

and  $b_t \geq 0$  for each  $t$

We need a numerical solver to solve this in general because there is no closed form way to know which non-negativity constraints will bind (e.g. which bids will be 0 at the optimum.) However, once we solve this, we know that the solution for any non-zero bid will follow:

$$b_t^*(s(\tilde{\alpha}, \tilde{\gamma})) = \tilde{\alpha} \cdot c_t + \frac{1}{\tilde{\gamma}} \cdot \frac{q_t^b}{\sigma_t^2} + \frac{q_t^e}{\sigma_t^2 \sum_{p: b_p^* > 0} \left[ \frac{(q_p^e)^2}{\sigma_p^2} \right]} \left( s - \sum_{p: b_p^* > 0} \left[ \tilde{\alpha} \cdot c_p q_p^e + \frac{1}{\tilde{\gamma}} \cdot \frac{q_p^b}{\sigma_p^2} q_p^e \right] \right)$$

If all of the item bids are above 0 at the optimum, then the summations in the formula above are over all items  $t = 1, \dots, T$ . Otherwise, the summations are over the items that have unit bids above 0 at the optimum.

We can find the partial derivative with respect to  $s$ ,  $\frac{\partial b_t(s)}{\partial s}$ , from this equation:

$$\frac{\partial b_t(\tilde{s})}{\partial s} = \frac{q_t^e}{\sigma_t^2 \sum_{p: b_p^* > 0} \left[ \frac{(q_p^e)^2}{\sigma_p^2} \right]} \quad \text{if } b_t(\tilde{s}) > 0$$

and  $\frac{\partial b_t(\tilde{s})}{\partial s} = 0$  if  $b_t(\tilde{s}) = 0$ .

Going back to the expected utility for participating in an auction, let's go back to a more high level expression:

$$EU(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma})) = \underbrace{E[u(\pi(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))) \mid \text{win}]}_{\text{Value of winning}} \times \underbrace{\Pr(s(\tilde{\alpha}, \tilde{\gamma}) < s(\alpha^j, \gamma^j))}_{\text{Probability of winning}}$$

and to simplify notation, let's write this as:

$$EU(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma})) \equiv V(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma})) \times P(s(\tilde{\alpha}, \tilde{\gamma})).$$

We would like to find the function  $s(\cdot, \cdot)$  such every bidder type  $(\tilde{\alpha}, \tilde{\gamma})$  will be maximally happy with his bid  $s(\tilde{\alpha}, \tilde{\gamma})$  — meaning that he could not profit by bidding a different score if his opponent is using  $s(\cdot, \cdot)$  to determine *her* bid.

A quick inspection of the  $EU(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))$  function shows that it is in fact concave in  $s$  (as a value), and so a sufficient condition for the optimality of  $s(\cdot, \cdot)$  is that the first order condition of  $EU(\cdot)$  holds:

$$\frac{\partial EU(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))}{\partial \tilde{s}} = 0.$$

This first order condition is what will define the differential equation that we need to solve in order to find a function  $s(\cdot, \cdot)$  that will satisfy the conditions to be an “equilibrium”.

Taking the derivative from our expression above:

$$\frac{\partial EU(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))}{\partial \tilde{s}} = \frac{\partial}{\partial \tilde{s}} V(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma})) \times P(s(\tilde{\alpha}, \tilde{\gamma})) + V(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma})) \times \frac{\partial}{\partial \tilde{s}} P(s(\tilde{\alpha}, \tilde{\gamma})) = 0.$$

Note that we can compute  $\frac{\partial}{\partial \tilde{s}} V(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))$  fairly easily — this is exactly the same as in the one dimensional type case:

Let's simplify notation one more time and write:

$$CE(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma})) = \sum_{t=1}^T q_t^b (b_t^*(s(\tilde{\alpha}, \tilde{\gamma})) - \tilde{\alpha} c_t) - \frac{\tilde{\gamma} \sigma_t^2}{2} (b_t^*(s(\tilde{\alpha}, \tilde{\gamma})) - \tilde{\alpha} c_t)^2$$

so that

$$V(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma})) = 1 - \exp[-\tilde{\gamma} CE(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))]$$

$$\frac{\partial}{\partial \tilde{s}} V(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma})) = \tilde{\gamma} \frac{\partial}{\partial \tilde{s}} CE(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma})) \times \exp[-\tilde{\gamma} CE(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))],$$

where

$$\frac{\partial}{\partial \tilde{s}} CE(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma})) = \sum_{t=1}^T \left[ \frac{\partial b_t^*(s(\tilde{\alpha}, \tilde{\gamma}))}{\partial s} (q_t^b - \tilde{\gamma} \sigma_t^2 (b_t^*(s(\tilde{\alpha}, \tilde{\gamma})) - \tilde{\alpha} c_t)) \right].$$

As noted before, we solve for  $b_t^*(s(\tilde{\alpha}, \tilde{\gamma}))$  and  $\frac{\partial b_t^*(s(\tilde{\alpha}, \tilde{\gamma}))}{\partial s}$  numerically, and so  $\frac{\partial}{\partial \tilde{s}} CE(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))$  and consequently  $\frac{\partial}{\partial \tilde{s}} V(\tilde{\alpha}, \tilde{\gamma}, s(\tilde{\alpha}, \tilde{\gamma}))$  are well defined and computable for any function  $s(\cdot, \cdot)$ .

The challenge is in figuring out how to figure out  $P(s(\tilde{\alpha}, \tilde{\gamma}))$ . For a given guess of  $s(\cdot, \cdot)$  this is easily done — just integrate over the distributions of  $\alpha$  and  $\gamma$ . But this is where my knowledge of implementing differential equation solvers is sparse.