

Ans to Ques no 1 (a)

As the award is given in a constant time. The whole award given-time complexity would be $O(n)$.

Ans to the Ques no 1 (b)

The asymptotic time complexity of the following function is:

$$O(1) * O(n) * O(n) * (O(1) + O(1))$$

$$\Rightarrow O(n^2)$$

Ans to the Ques no 1 (c)

Following function is:

$$T(n) = 625T\left(\frac{n}{5}\right) + n^3$$

As per Master's Theorem,

$$T(n) = aT\left(\frac{n}{b}\right) + cn^k$$

here,

$$a = 625$$

$$b = 5$$

$$k = 3$$

$$c = 1$$

$$\therefore b^k < a$$

$$\Rightarrow 5^3 < 625$$

$$\Rightarrow 125 < 625$$

if $b^k < a$, the time complexity would be $O(n^{\log_b a})$

$$\therefore T_e = O(n^{\log_5 625})$$

$$= O(n^4)$$

Ans

Ans to the Ques no 2 (a)

Four

Time complexity of the following function is:

$$O(\log_7 n) * O(n) * \Theta(O(n) + O(n))$$

$$\Rightarrow O(n^2 \log_7 n)$$

Ans the Ques no 2 (b)

Given recurrence relation is:

$$T(n) = T(n/2) + T(n/4) + n$$

Substitution method,

$$T(n) = T(n/2) + T(n/4) + n$$

$$= [T(n/2^2) + T(n/2^3) + n/2] + n$$

$$= [T(n/2^3) + T(n/2^4) + n/2^2] + n/2 + n$$

$$= [T(n/2^4) + T(n/2^5) + n/2^3] + n/2^2 + n/2 + n$$

$$= T(n/2^k) + T(n/2^{k-1}) + \dots + T(n/2^{k-2}) + n/2^{k-3} + n/2^{k-4} + \dots + n$$

Assume,

$$n = 2^k$$

$$n/2^k = 1$$

$$n = 2^k$$

$$k = \log_2 n$$

\therefore Time complexity of the given recurrence relation is $O(\log n)$

Ans

Sorting 1:

a) Since there is no requirement about stability and it should need to do it in place method that's why I choose to bubble sort.

In bubble sort we will move left to right swapping adjacent elements as needed. Each pass moves the next largest element into its final position.

15	8	11	112	33	39	88	41
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↓

8	15	11	112	33	39	88	41
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↓

8	11	15	112	33	39	88	41
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↓

8	11	15	33	112	39	88	41
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↓

8	11	15	33	39	112	88	41
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↓

8	11	15	33	39	88	112	41
---	----	----	----	----	----	-----	----

↓

8	11	15	33	39	88	41	112
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↓

8	11	15	33	39	41	88	112
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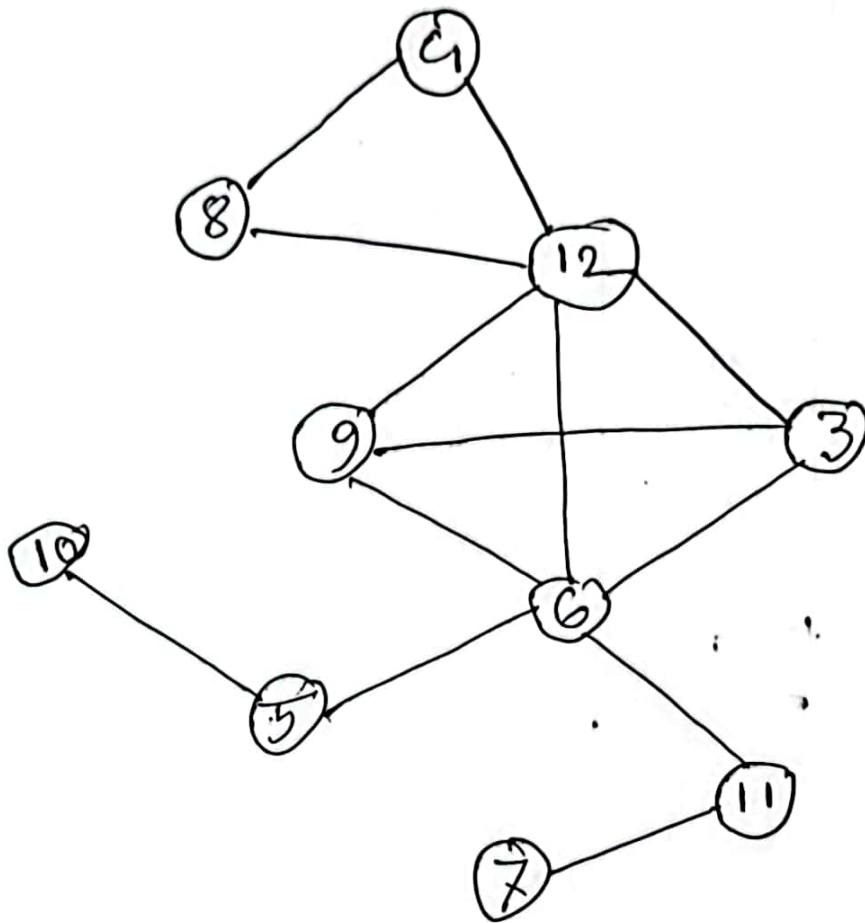
In bubble sort it iterates each elem and compare with the next one: if $a[i] > a[i+1]$ then it'll swap otherwise $a[i] == a[i+1]$.

- b) The best case occurs when the array is already sorted. In a sorted list there is 0 number of swapping and $(n-1)$ numbers of comparison. Hence the best case of time complexity of bubble sort is $O(n)$.

Ans.

Graph 2:

a)



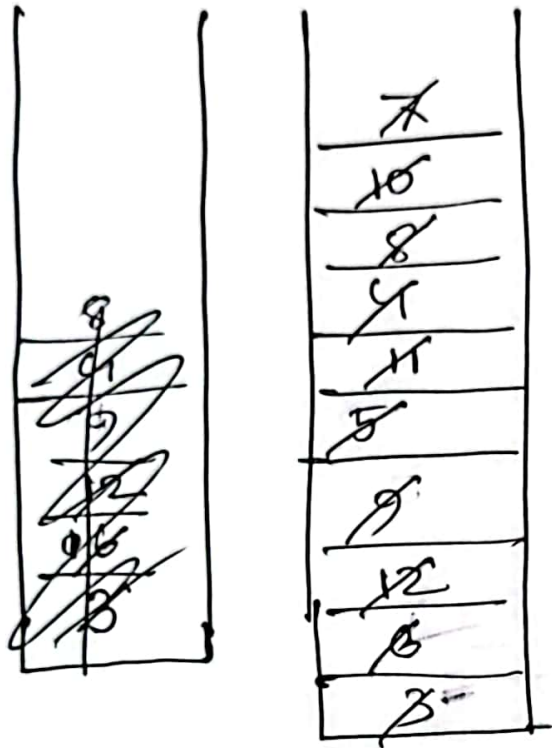
b) Yes bill is right. there are at least 4 triangles.

we can see from the graph.

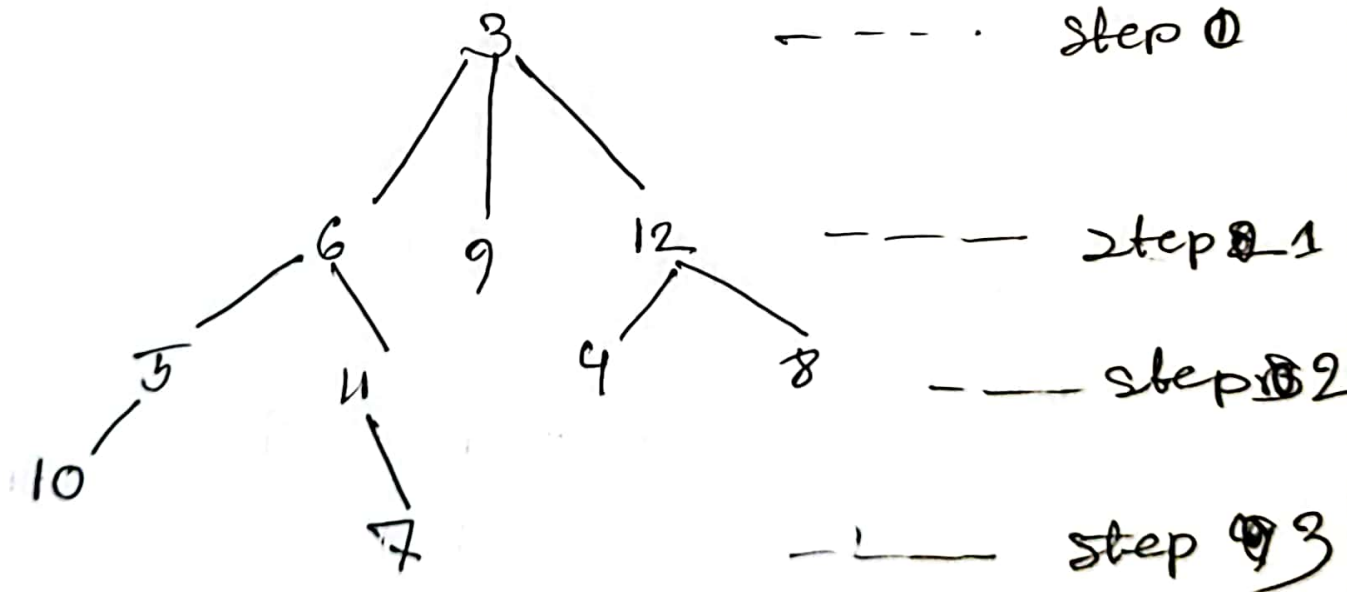
edges $[1, 8, 12]$, $[12, 9, 3]$, $[3, 9, 6]$, $[12, 9, 6]$

make a connected and make triangles.

c) From the graph



visited node



for node $[6, 9, 12]$ shortest distance is 1 from node 3 and for $[6, 11, 4, 8]$ shortest distance is 2 from node 3 and for $[10, 7]$ the shortest distance is 3 from node 3

d) In this graph we saw that we have 13 edges. From no 3 to 12 vertex we have $\frac{n(n-1)}{2}$ number of edges if we keep the graph simple. Now, for 3 to 12 number of nodes the total edges will be $= \frac{10(9)}{2}$
 $= 45$
edges

\therefore Number of edges we need more

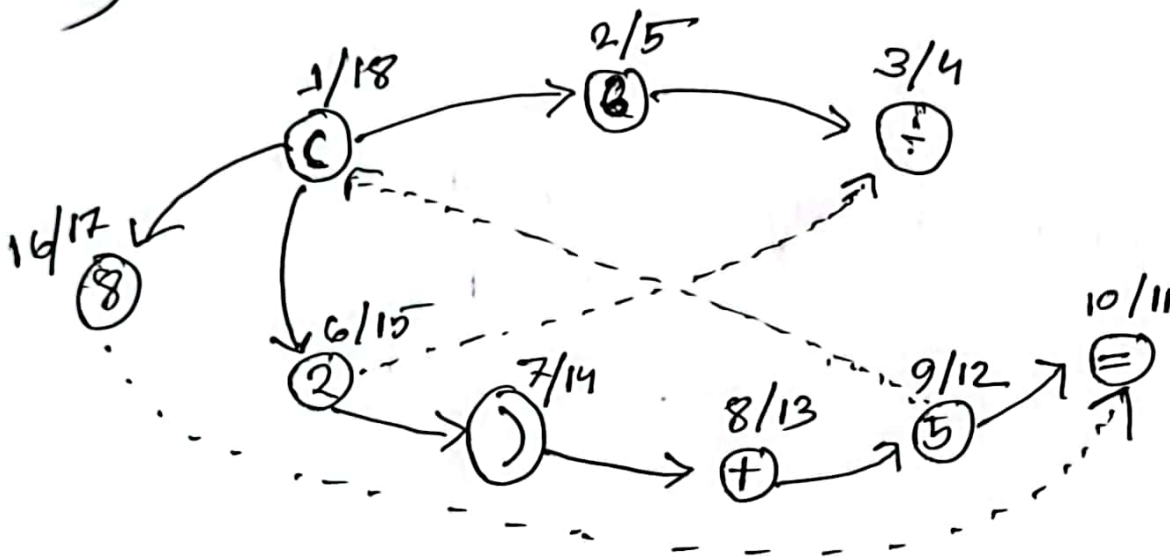
$$= (45 - 13)$$

$$= 32 \text{ edges} \quad \underline{\text{Ans}}$$

Graph 3

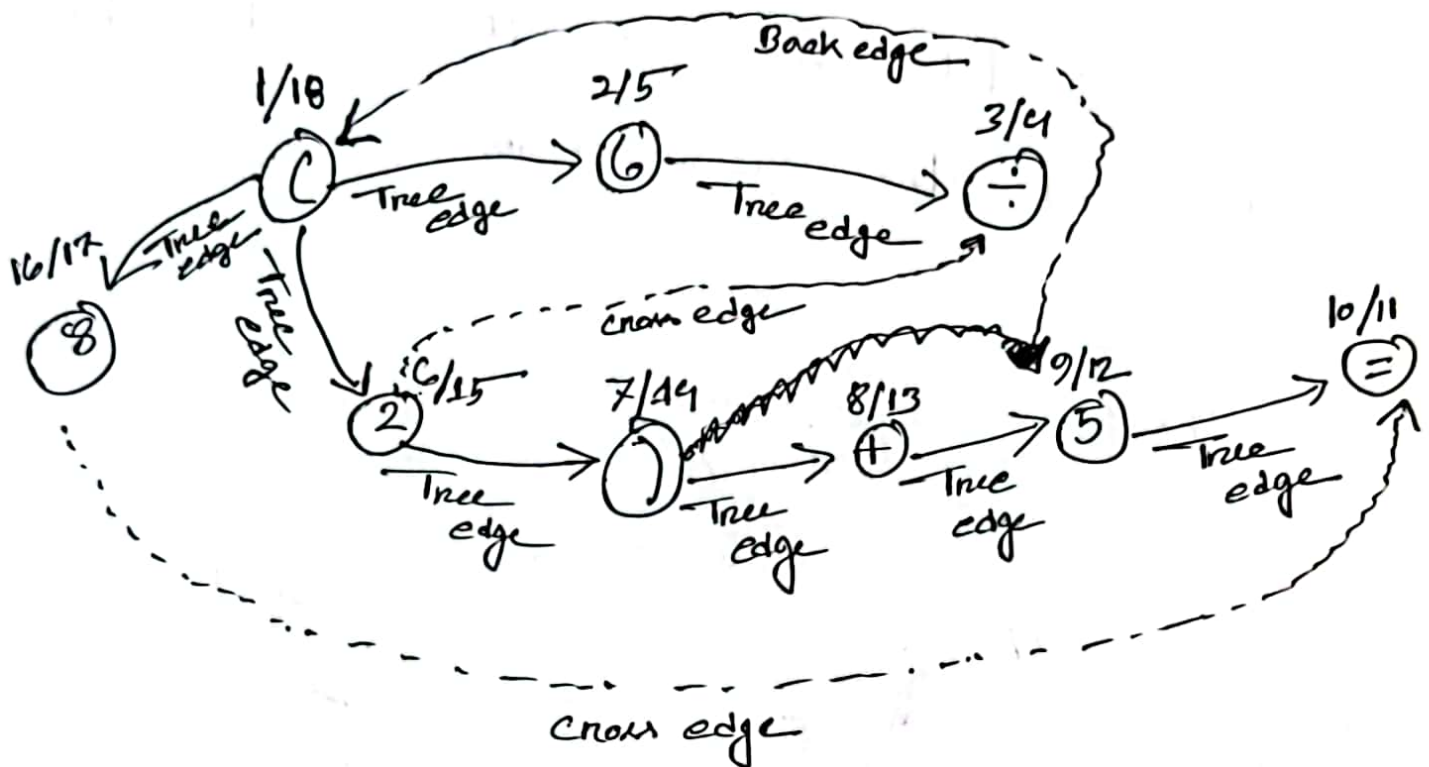
a) There are 8 edges in the DFS tree of the equation.

b) ~~10~~



After running a DFS from the equation and making a DFS tree we can say that Bill was right — the equation is achievable by running a DFS.

c) From the DFS Tree we can see that

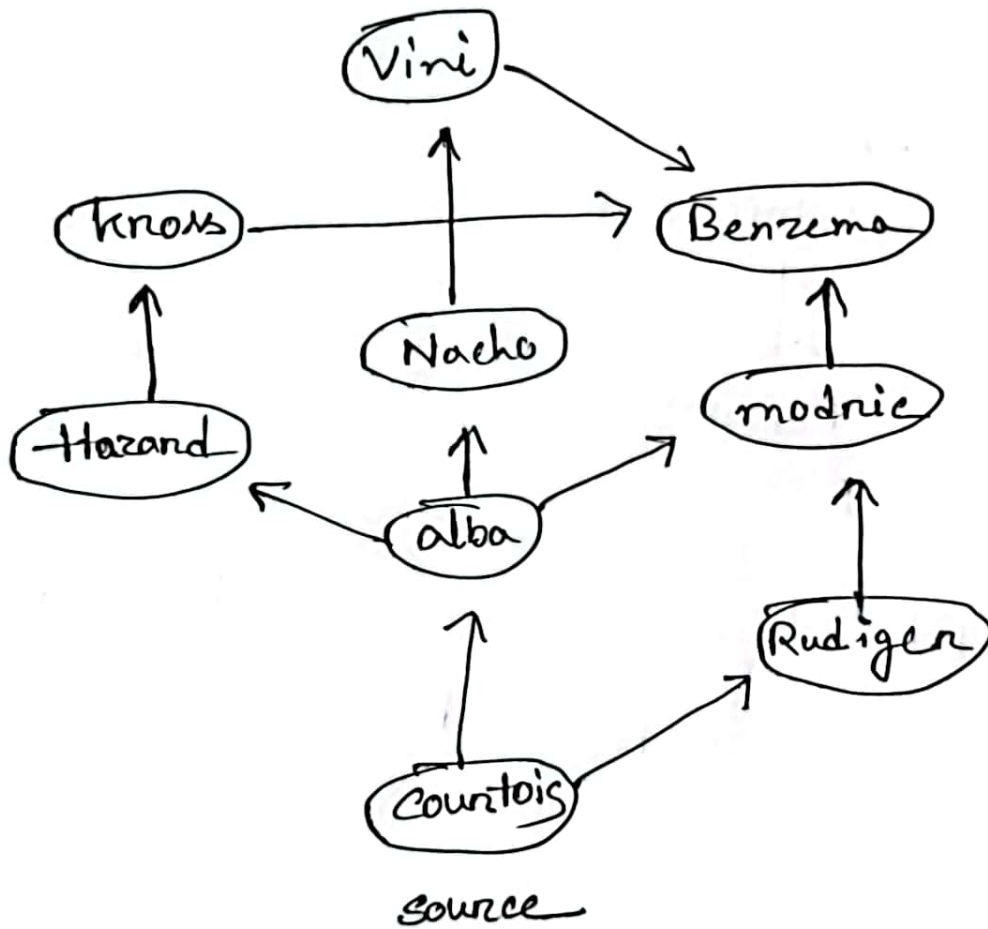


From DFS tree we can see that there are 2 cross edges from node "8" to node "=" and node "2" to node "÷", and there are one back edge from node "5" to node "C" and rest of all means total 8 tree edges.

Graph 1:

a) we will use BFS algorithm to know the shortest path.

Here is the BFS Graph

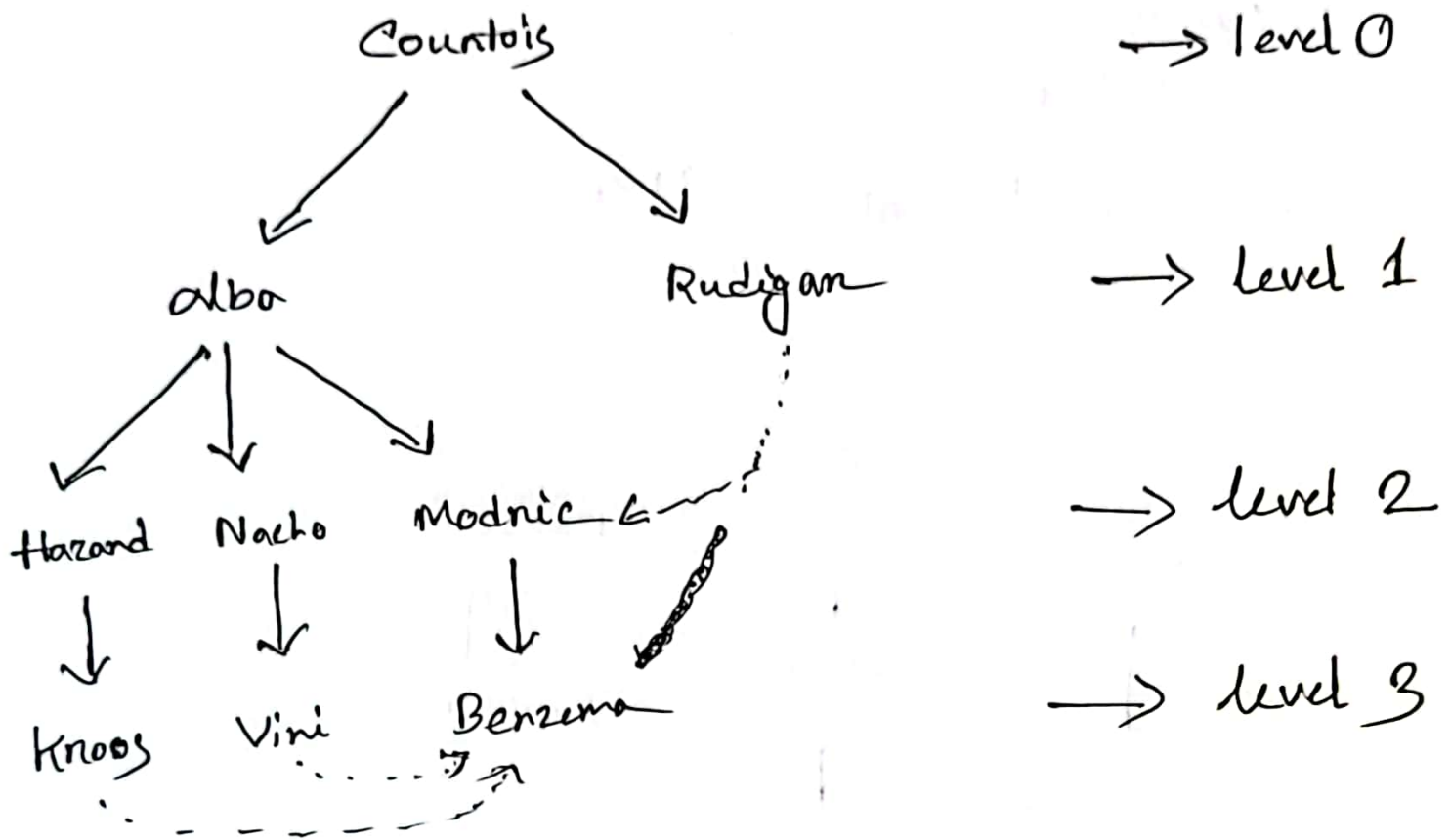


New drawn BFS tree:

Q

Courtois	alba	Rudiger	Hazard	Nacho	modric	knows	Vini	Benzema	
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BFS Tree:



From the BFS tree we can see that minimum 3 passes need to reach the ball from Courtois to Benzema.

b) From the BFS tree we know that there are 3 minimum passes need to reach the ball from Courtois to Benzema. The players who are required for this minimum passes shown below.

Courtois
↓
Rudigan
↓
~~Benzema~~
Modric
↓
Benzema

or

Courtois
↓
Alba
↓
Modric
↓
Benzema

Sorting 2

a) we will do the bubble sort algorithm in the list, as the time complexity of Bubble sort is $O(n^2)$.

steps are shown below.

23	2	19	3	7	11	5	13
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↓

2	23	19	3	7	11	5	13
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↓

2	19	23	3	7	11	5	13
---	----	----	---	---	----	---	----

↓

2	19	3	23	7	11	5	13
---	----	---	----	---	----	---	----

↓

2	19 3	19	23	7	11	5	13
---	-----------------	----	----	---	----	---	----

↓

2	3	19	7	23	11	5	13
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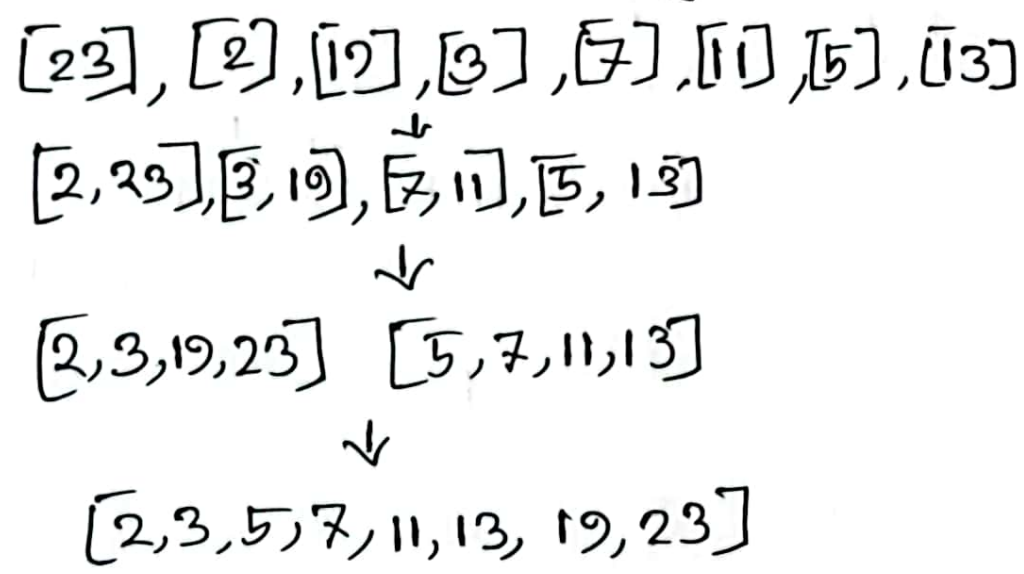
↓

2	3	19	7	23	23	5	13
---	---	----	---	---------------	----	---	----

⋮ ⋮ ⋮ ⋮ ⋮ ⋮ ⋮

2	3	5	7	11	13	19	23
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b) To sort the given array in linear time, we can use merge function. as we know the time complexity of merge function is $O(n)$. First we can divide the array into individual subarrays, then we will merge adjacent subarrays pairwise, creating new subarrays.



c) we will use binary search to add 15 into the sorted list.

steps are shown below :

① $\overset{L}{[2, 3, 5, 7, 11, 13, 19, 23]}^R$

left = 0

Right = 7

mid = 3

$7 < 15$

(ii) $\text{left} = \text{middle} + 1 = 4$
 $\text{Right} = 7$
 $\text{mid} = 5$

$$[2, 3, 5, 7, \overset{L}{11}, \overset{m}{13}, \overset{R}{19}, 23]$$

$$13 < 15$$

(iii) $\text{left} = \text{middle} + 1 = 6$
 $\text{Right} = 7$
 $\text{mid} = 6$

$$[2, 3, 5, 7, 11, \overset{L/m}{13}, \overset{R}{19}, 23]$$

$$19 > 15$$

(iv) $\text{Right} = \text{mid} - 1 = 5$
 $\text{left} = 6$
 $\text{mid} = 6$

$$[2, 3, 5, 7, 11, \overset{R}{13}, \overset{L/m}{19}, 23]$$

Now insert the value into index 6

$$[2, 3, 5, 7, 11, 13, 15, 19, 23]$$

The updated sorted array is $[2, 3, 5, 7, 11, 13, 15, 19, 23]$
 with the element 15 inserted at the correct position.

Sorting 3

(10)

21201240
Tom

- ⑥ Jimmy's strategy of modifying bubble sort and insertion sort algorithm to find the first five largest and five lowest numbers in just 5 iteration, this approach is not reliable because:
- ① Bubble and insertion sort are designed to sort the whole entire array not to finding the largest or smallest elements.
 - ② Only in 5 iteration the sorting will not complete fully. Consequently, the rightmost and leftmost 5 numbers may not actually represent the largest or smallest elements respectively.
 - ③ The largest elements may not necessarily occupy the rightmost position after a limited number of iterations and the smallest elements may not occupy the leftmost positions. The order of elements can be arbitrary within the partially

sorted subarray.

- (iv) After 5 iteration we can see the
five right most number will be $[12, 15, 6, 10, 18]$
and left most number will be $[2, 3, 4, 7, 9]$

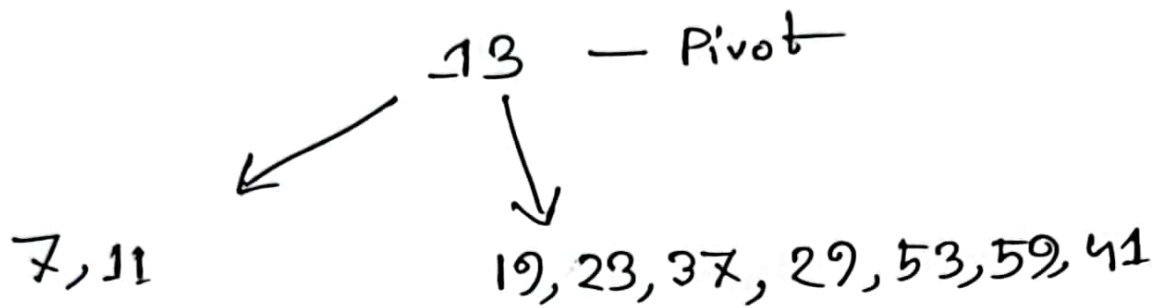
so, as we can see from this, Jimmy's strategy
of modifying bubble sort and insertion sort
is not reliable.

- (v) In the given list: $[13, 11, 19, 7, 23, 37, 29,$
 $53, 59, 41]$

in this list the first element is 13.
for 13 to be placed in its correct
position, it should be greater than
11, 7 and ~~less~~ smaller than (23, 37, 29, 53
59, 41).

\therefore The pivot before the first partition
was 13.

(d)

$$[13, 11, 19, 7, 23, 37, 29, 53, 59, 41]$$


After partitioning using 13 as the pivot
the list will look like:

$$[7, 11, 13, 19, 23, 37, 29, 53, 59, 41]$$

Ans