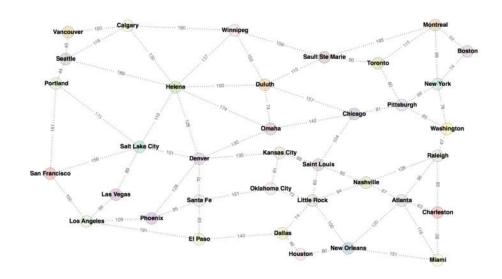
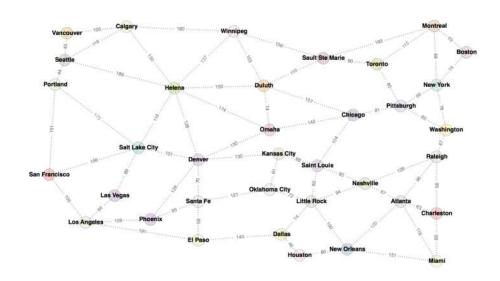


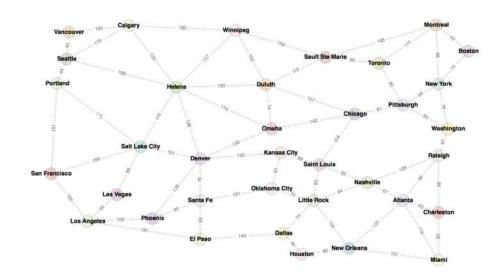
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- Modify BFS: Prioritize by cost not depth \rightarrow **Expand node** n with the lowest path cost g(n)

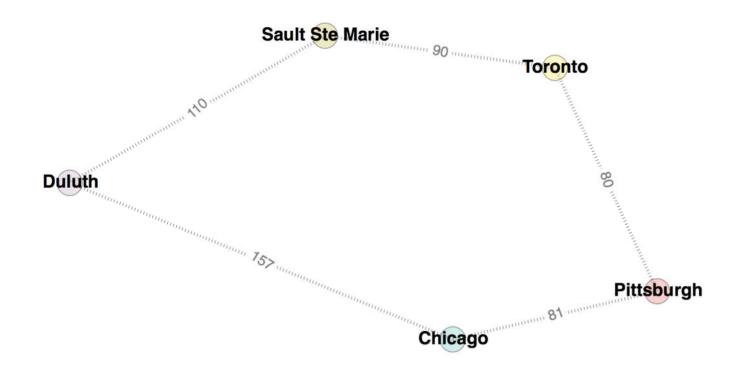


- The arcs in the search graph may have weights (different cost attached). How to leverage this information?
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- We want the cheapest not shallowest solution.
- Modify BFS: Prioritize by cost not depth \rightarrow **Expand node** n with the lowest path cost g(n)
- Explores increasing costs.

UCS algorithm

```
function Uniform-Cost-Search(initialState, goalTest)
returns Success or Failure: /* Cost f(n) = g(n) */
frontier = Heap.new(initialState)
explored = Set.new()
while not frontier.isEmpty():
     state = frontier.deleteMin()
     explored.add(state)
     if goalTest(state):
          return Success(state)
     for neighbor in state.neighbors():
          if neighbor not in frontier \cup explored:
                frontier.insert(neighbor)
          else if neighbor in frontier:
                frontier.decreaseKey(neighbor)
```

return FAILURE



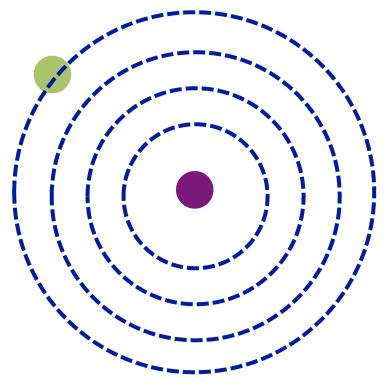
Go from Chicago to Sault Ste Marie. Using BFS, we would find Chicago-Duluth-Sault Ste Marie. However, using UCS, we would find Chicago-Pittsburgh-Toronto-Sault Ste Marie, which is actually the shortest path!

• Complete Yes, if solution has a finite cost.

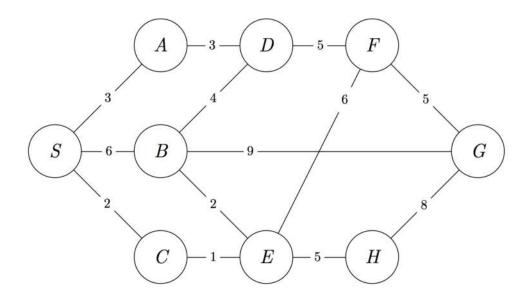
Time

- Suppose C^* : cost of the optimal solution
- Every action costs at least ϵ (bound on the cost)
- The effective depth is roughly C^*/ϵ (how deep the *cheapest* solution could be).
- $-O(b^{C^*/\epsilon})$
- Space # of nodes with $g \leq \text{cost of optimal solution, } O(b^{C^*/\epsilon})$
- Optimal Yes
- Implementation: fringe = queue ordered by path cost g(n), lowest first = Heap!

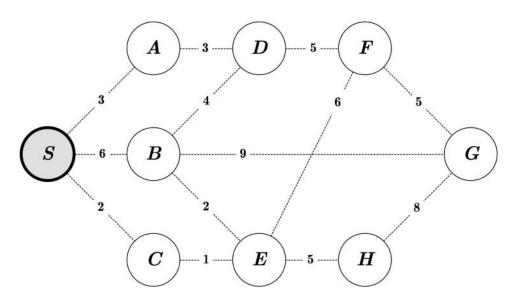
While complete and optimal, UCS explores the space in every direction because no information is provided about the goal!



Exercise



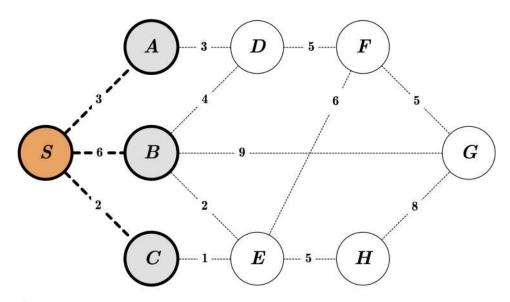
Question: What is the order of visits of the nodes and the path returned by BFS, DFS and UCS?



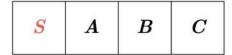
Queue:

 $oldsymbol{S}$

Order of Visit:

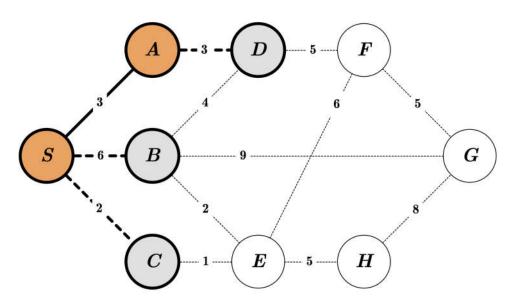


Queue:

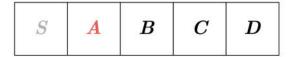


Order of Visit:

 \boldsymbol{S}

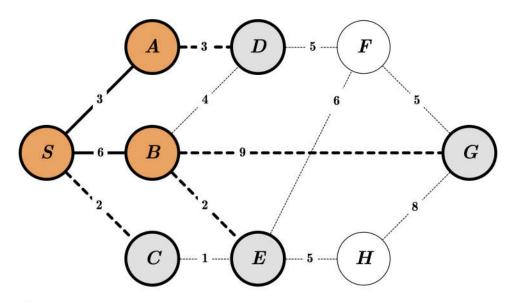


Queue:

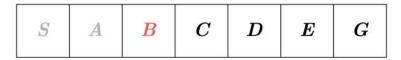


Order of Visit:

S A

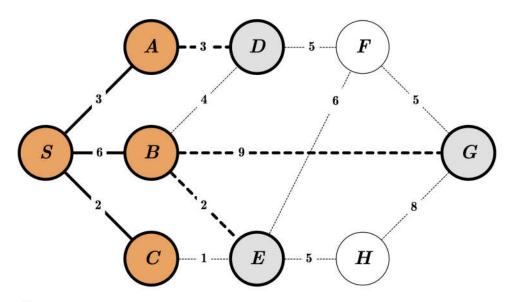


Queue:

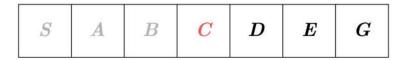


Order of Visit:

 $S \qquad A \qquad B$

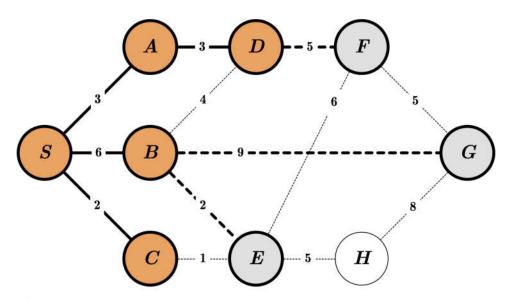


Queue:



Order of Visit:

S A B C

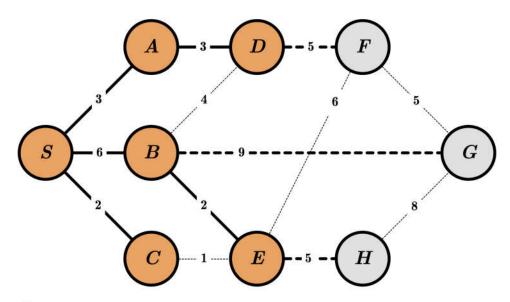


Queue:



Order of Visit:

S A B C D

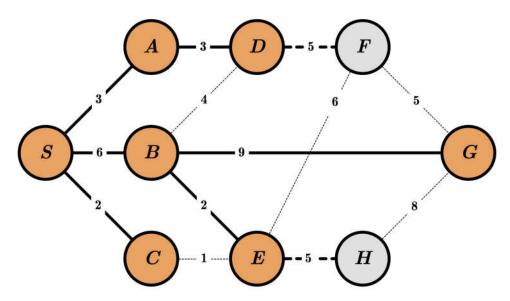


Queue:



Order of Visit:

S A B C D E

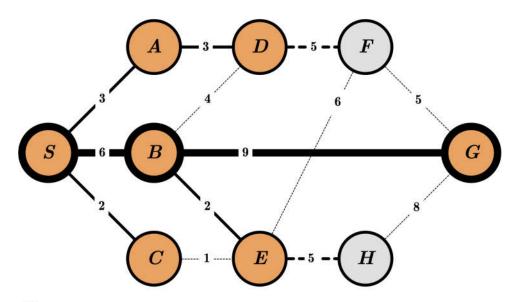


Queue:



Order of Visit:

S A B C D E G

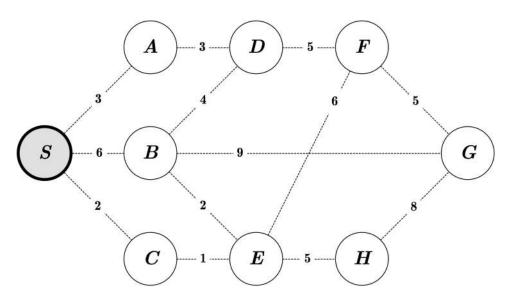


Queue:



Order of Visit:

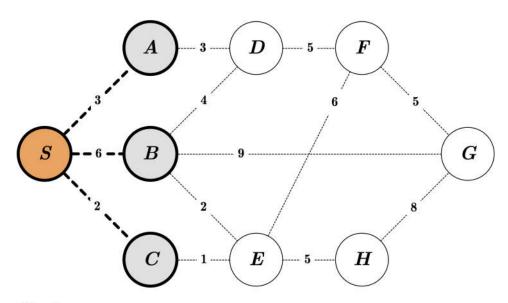
S A B C D E G



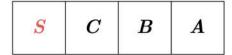
Stack:

 $oldsymbol{S}$

Order of Visit:

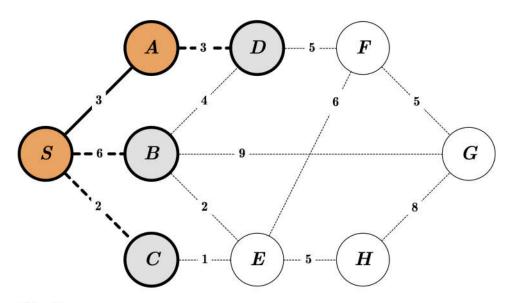


Stack:

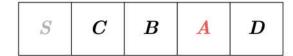


Order of Visit:

 \boldsymbol{S}

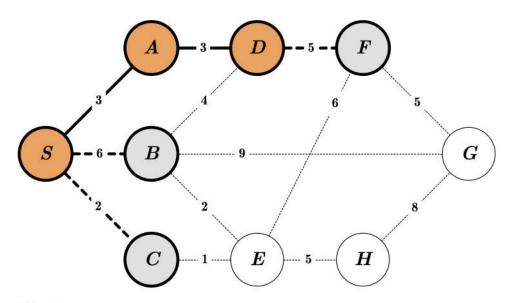


Stack:



Order of Visit:

S A

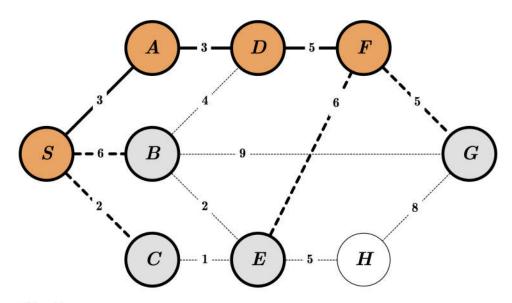


Stack:



Order of Visit:

 $S \quad A \quad D$

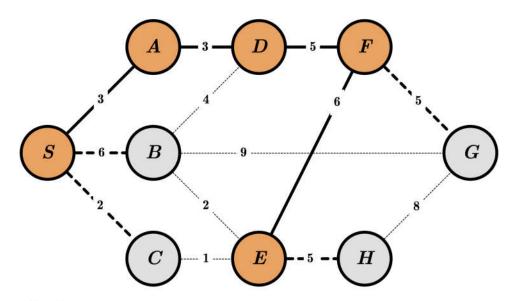


Stack:



Order of Visit:

S A D F

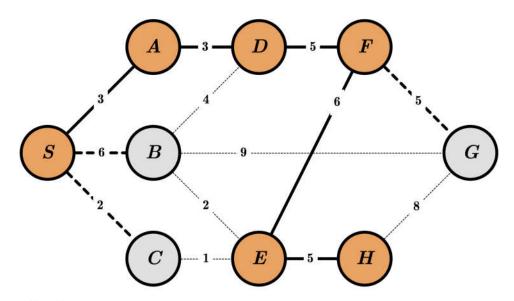


Stack:



Order of Visit:

S A D F E

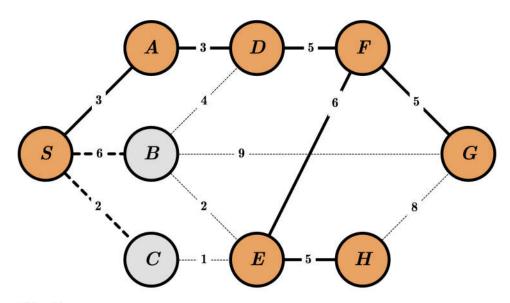


Stack:



Order of Visit:

S A D F E H

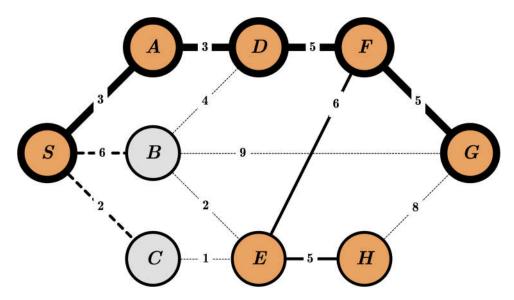


Stack:



Order of Visit:

S A D F E H G

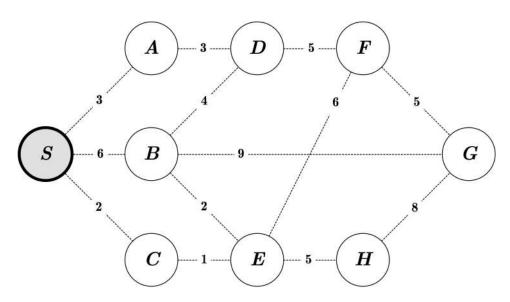


Stack:



Order of Visit:

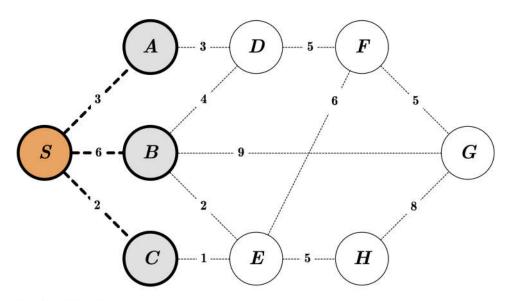
S A D F E H G



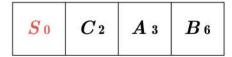
Priority Queue:

 $oldsymbol{S}$ 0

Order of Visit:

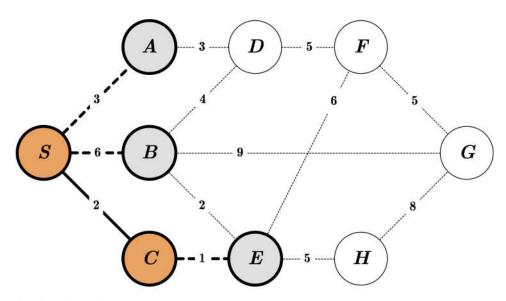


Priority Queue:

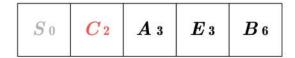


Order of Visit:

 \boldsymbol{S}

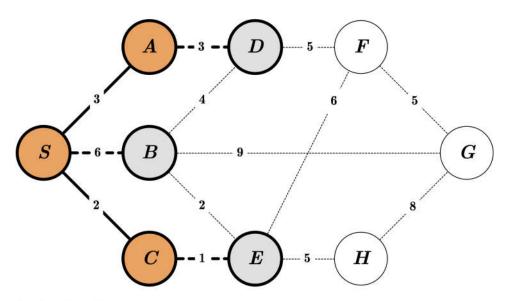


Priority Queue:



Order of Visit:

S C

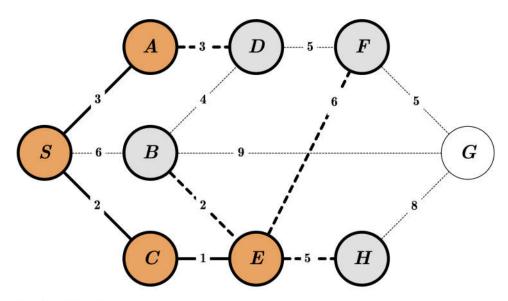


Priority Queue:

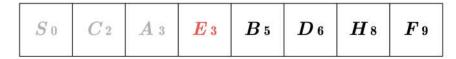


Order of Visit:

 $S \quad C \quad A$

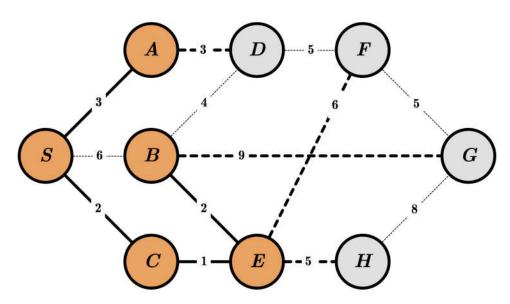


Priority Queue:



Order of Visit:

 $S \quad C \quad A \quad E$

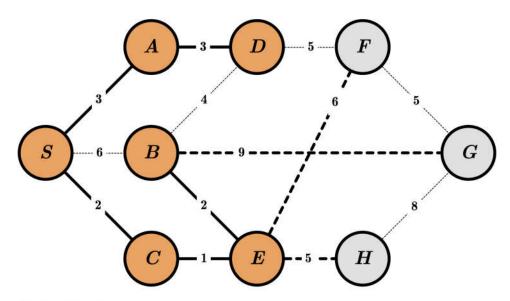


Priority Queue:



Order of Visit:

S C A E B

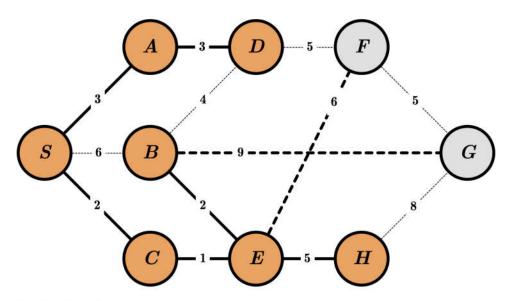


Priority Queue:

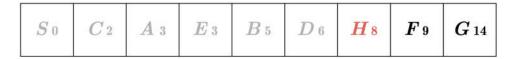


Order of Visit:

S C A E B D

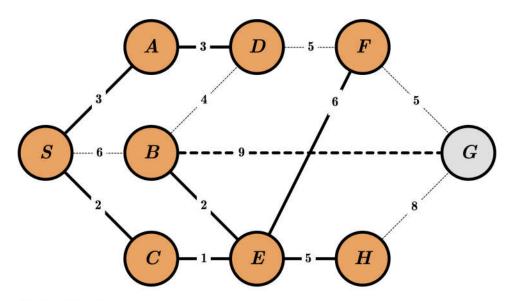


Priority Queue:



Order of Visit:

S C A E B D H

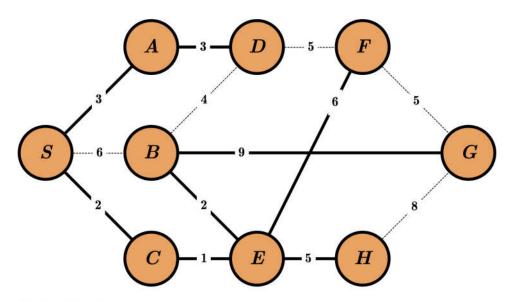


Priority Queue:



Order of Visit:

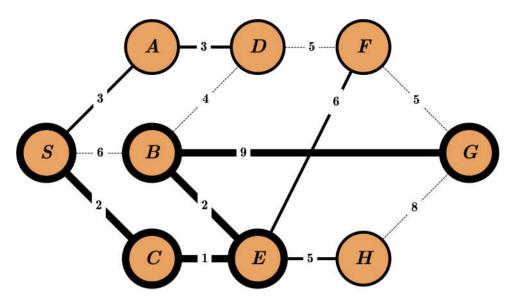
S C A E B D H F



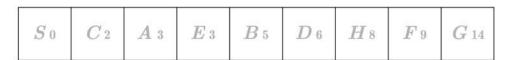
Priority Queue:



Order of Visit:



Priority Queue:

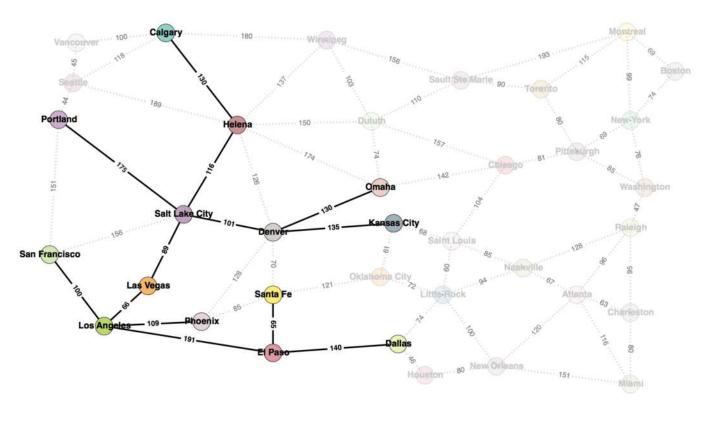


Order of Visit:

Examples using the map

Start: Las Vegas

Goal: Calgary



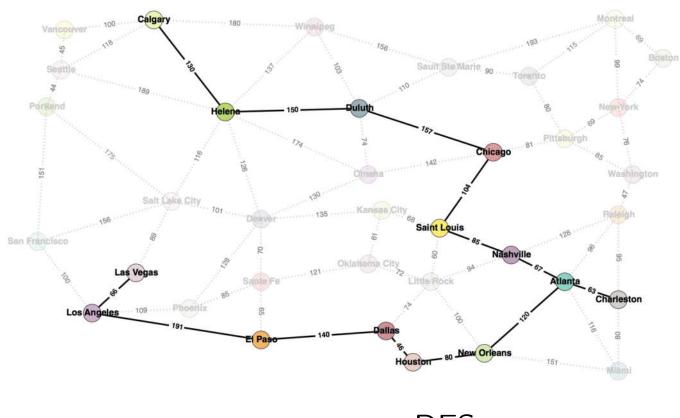
BFS

Order of Visit: Las Vegas, Los Angeles, Salt Lake City, El Paso, Phoenix, San Francisco, Denver, Helena, Portland, Dallas, Santa Fe, Kansas City, Omaha, Calgary.

Examples using the map

Start: Las Vegas

Goal: Calgary



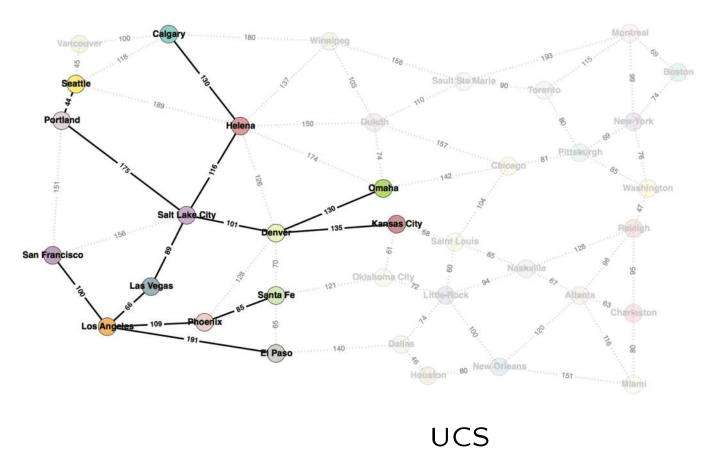
DFS

Order of Visit: Las Vegas, Los Angeles, El Paso, Dallas, Houston, New Orleans, Atlanta, Charleston, Nashville, Saint Louis, Chicago, Duluth, Helena, Calgary.

Examples using the map

Start: Las Vegas

Goal: Calgary



Order of Visit: Las Vegas, Los Angeles, Salt Lake City, San Francisco, Phoenix, Denver, Helena, El Paso, Santa Fe, Portland, Seattle, Omaha, Kansas City, Calgary.

Credit

• Artificial Intelligence, A Modern Approach. Stuart Russell and Peter Norvig. Third Edition. Pearson Education.

http://aima.cs.berkeley.edu/