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1 2ECC

```
struct graph {
  int n, t, sz;
  vector<vector<int>> adj;
  vector<int> tin, low, cmp;
  graph(int n): n(n),adj(n),tin(n),low(n),cmp(n){}
  void add_edge(int u, int v){
   adj[u].push_back(v);
    adj[v].push_back(u);
  void dfs(int u, int p){
    tin[u]=low[u]=t++;
    int cnt=0;
    for(int v: adj[u]){
      if(v==p and ++cnt <= 1) continue;</pre>
      if(tin[v]!=-1) low[u] = min(low[u], tin[v]);
        dfs(v,u);
        low[u] = min(low[u], low[v]);
  void dfs2(int u, int p){
   if(p!=-1 \text{ and } tin[p]>=low[u]) cmp[u] = cmp[p];
    else cmp[u] = sz++;
   for(int v: adj[u]){
      if(cmp[v]==-1) dfs2(v,u);
  void process 2ecc(){
   t = 0, sz = 0;
    for (int i = 0; i < n; ++i){
      tin[i] = low[i] = cmp[i] = -1;
    for (int i = 0; i < n; ++i){
      if(tin[i]==-1) dfs(i,-1);
    for (int i = 0; i < n; ++i){
      if(cmp[i]==-1) dfs2(i,-1);
```

2 2SAT

```
//CNF: (a | b) ^{\circ} (c | d) means (!a -> b) ^{\circ}
// (!a or \dot{b}) = (-a, \dot{b}), 1-based indexing
string two sat(int n, vector<array<int, 2>>

    clauses) {
  vector<int> adj[2 * n];
for (auto [a, b]: clauses) {
    if (a > 0) a = 2 * a - 2;
    else a = 2 * -a - 1:
    if (b > 0) b = 2 * b - 2;
    else b = 2 * -b - 1;
    adj[a ^ 1].push_back(b), adj[b ^
 \rightarrow 1].push back(a);
  vector<vector<int>>> sccs = get sccs(2 * n, adj);
  int tot scc = sccs.size();
  vector<int> scc no(2 * n);
  for (int i = 0; i < tot scc; ++i) {
    for (int u: sccs[i]) {
```

```
scc_no[u] = i;
}
string assignment;
for (int u = 0; u < n; u++) {
   if (scc_no[2 * u] == scc_no[2 * u + 1]) {
      return "";
   }
   if (scc_no[2 * u] < scc_no[2 * u + 1]) {
      assignment += '-';
   }
   else {
      assignment += '+';
   }
}
return assignment;
}</pre>
```

3 AHO_CORASICK

struct AC{

```
const int A = 26;
vector<vector<int>>> nxt, idx;
vector<int> lnk, out lnk, ans;
AC(){newNode();}
int newNode(){
  nxt.eb(A, 0), idx.eb(0);
  lnk.eb(0), out lnk.eb(0), ans.eb(0);
  return nxt.size()-1;
void clear(){
  nxt.clear(), idx.clear();
  nxt.clear(), idx.clear();
  lnk.clear(), out lnk.clear(), ans.clear();
  newNode();
// 0(|p|)
void add(string p, int i){
  int v=0;
  v = nxt[v][c-'a'];
  idx[v].eb(i);
// 0(|p1+p2+p3+..|)
void build(){
  queue<int> q; q.push(0);
  while (!q.empty()){
    int u=q.front(); q.pop();
for (int i = 0; i < A; ++i){</pre>
      int v = nxt[u][i];
      if(!v) nxt[u][i] = nxt[lnk[u]][i];
        lnk[v] = u? nxt[lnk[u]][i]: 0;
        out lnk[v] = idx[lnk[v]].empty()?
 out lnk[lnk[v]]: lnk[v];
        q.push(v);
// O(|T|+match)
void trav(string T){
  int v=0:
  for(char c: T){
```

```
if(!nxt[v][c-'a']) v = lnk[v];
    if(nxt[v][c-'a']) v=nxt[v][c-'a'];
    for(auto& i: idx[v]){
        ans[i]++;
    }
    int x = out_lnk[v];
    while(x){
        for(auto& i: idx[x]){
            ans[i]++;
        }
        x = out_lnk[x];
    }
    x = out_lnk[x];
}
//AC ac; ac.add(pi, i); ac.build(); ac.trav(T);
```

4 ARTICULATION_BRIDGE

```
vector<int> adj[N];
int t = 0:
vector<int> tin(N, -1), lo(N);
vector<array<int, 2>> ab;
void dfs (int u, int p) {
 tin[u] = lo[u] = t++;
 for (int v: adj[u]) {
   if (v != p) {
      if (tin[v] != -1)
        lo[u] = min(lo[u], tin[v]);
      else {
        dfs(v, u);
        if (tin[u] < lo[v]) +
          ab.push back({u, v});
        lo[u] = min(lo[u], lo[v]);
dfs(0, -1);
```

5 ARTICULATION_POINT

```
vector<int> adj[N];
int t = 0;
vector<int> tin(N, -1), low(N), ap;
void dfs (int u, int p) {
 tin[u] = low[u] = t++;
 int is ap = 0, child = 0;
 for (int v: adj[u]) {
   if (v != p) {
      if (tin[v] != -1) {
        low[u] = min(low[u], tin[v]);
      else {
        child++;
        dfs(v, u);
        if (tin[u] <= low[v]) {
          is ap = 1;
        low[u] = min(low[u], low[v]);
```

6 BCC

```
struct graph {
 int n, t=0, cno=0;
 vector<vector<int>> q;
 vector<int> tin, lo, bcomp;
 stack <int> st;
 graph(int n):n(n),g(n),lo(n),bcomp(n){}
 void add edge(int u, int v){
   g[u].push back(v);
    g[v].push_back(u);
 void dfs(int v, int p=-1){
   lo[v]=tin[v]=++t;
    st.push(v);
   for(int u:g[v]){
     if(u==p)
                  continue;
      if(!tin[u]){
        dfs(u, v);
        lo[v]=min(lo[v],lo[u]);
     } else{
        lo[v]=min(lo[v],tin[u]);
    if(tin[v]==lo[v]){
      while (!st.empty()){
        int tp=st.top(); st.pop();
        bcomp[tp]=cno;
        if(tp==v)
                     break:
      ćno++:
 vector<int> bcc(){
   tin.assign(n, 0);
   for (int i = 0; i < n; ++i){
     if(!tin[i])
       dfs(i);
    return bcomp;
```

7 BCC_EDGE

```
int ch = 0;
 tin[u] = lo[u] = t++;
  for(auto [v, e] : adj[u]) {
   if (v == p) continue;
if (tin[v] != -1) {
      if (tin[u] > tin[v]) {
        lo[u] = min(lo[u], tin[v]);
        stk.push(e);
    else {
      ch++:
      stk.push(e);
      dfs(v, u);
      if ((p != -1 \text{ or } ch > 1) \text{ and } tin[u] <= lo[v]) {
        is ap[u] = 1;
        pop bcc(e);
      lo[u] = min(lo[u], lo[v]);
void procces bcc(int n) {
 for (int i = 0; i < n; ++i) {
    tin[i] = -1, is ap[i] = 0;
    bcc ed[i].clear();
    bcc[i].clear();
 t = tot = 0;
 for (int u = 0; u < n; ++u) {
    if (tin[u] == -1) {
      dfs(u, -1);
      if (!stk.empty()) {
        while (!stk.empty()) {
          bcc ed[tot].push back(stk.top());
   stk.pop();
        tot++;
 for (int i = 0; i < tot; ++i) {
    for (auto e: bcc ed[i]) {
      auto [u, v] = \overline{e}dges[e];
      bcc[i].push back(u);
      bcc[i].push_back(v);
 for (int i = 0; i < tot; ++i) {
    sort(bcc[i].begin(), bcc[i].end());
    bcc[i].erase(unique(bcc[i].begin(),
   bcc[i].end()), bcc[i].end());
```

8 BIT_TRICKS

3

9 BLOCK_CUT_TREE

```
vector<int> adj[N];
vector\langle int \rangle tin(N, -1), lo(N), is ap(N), bcc[N];
stack<int> stk;
int t = 0, tot = 0;
void pop bcc(int u, int v) {
 bcc[tot].push back(u);
 while (bcc[to\overline{t}].back() != v)
    bcc[tot].push back(stk.top());
    stk.pop();
 tot++:
void dfs (int u, int p) {
 tin[u] = lo[u] = t++;
 stk.push(u);
 int ch = 0;
 for (auto v: adj[u]) {
    if (v != p) {
      if (tin[v] != -1)
        lo[u] = min(lo[u], tin[v]);
      else {
        ch++:
        dfs(v, u);
        if ((p != -1 \text{ or } ch > 1) \text{ and } tin[u] <=
→ lo[v]) -
           // is ap[u] = 1;
           pop b\bar{c}c(u, v);
         lo[u] = min(lo[u], lo[v]);
void process bcc (int n) {
 for (int u = 0; u < n; ++u) {
    tin[u] = -1;
    is_ap[u] = 0;
    bc\overline{c}[u].clear();
 t = tot = 0;
 for (int u = 0; u < n; ++u) {
    if (tin[u] == -1) {
      dfs(u, -1);
      if (!stk.empty()) {
        while (!stk.empty()) {
          bcc[tot].push back(stk.top());
          stk.pop();
        tot++;
      }
```

```
4
```

```
int nn;
vector<int> comp_num(N), bct_adj[N];
void build bct(int n) {
  process bcc(n);
  int nn \equiv tot;
  for (int u = 0; u < n; ++u) {
    if (is ap[u]) {
      comp num[u] = nn++;
  for (int i = 0; i < tot; ++i) {
    for (auto u: bcc[i]) {
      if (is ap[u])
        u = \overline{comp} num[u];
        bct adj[i].push back(u);
        bct adj[u].push back(i);
      else {
        comp num[u] = i;
```

10 CDQ

```
## Problems related to pair
- cdq(l, m)
- cdq(m + 1, r)
- handle influence of (l, m) to (m + 1, r)
## Optimization of 1D DP
- cdq(l, m)
- handle influence of (l, m) to (m + 1, r)
- cdq(m + 1, r)
## Convert dynamic array problems to static array
- problem
```

11 CENTROID_DECOMPOSITION

```
void calc sz(int u, int p) {
  sz[u] = 1;
  for (auto v: adj[u]) {
    if (v != p and !is_cen[v]) {
      calc sz(v, u);
      sz[u] += sz[v];
int get cen(int u, int p, int n) {
  for (auto v: adj[u]) {
    if (v != p \text{ and } !is\_cen[v] \text{ and } 2 * sz[v] > n) {
      return get cen(v, u, n);
  return u;
void decompose(int u=0, int p=-1, int d=0){
  calc sz(u, p);
  int \overline{c} = get cen(u, p, sz[u]);
  is_cen[c] = 1, cpar[c] = p, cdep[c] = d;
  for(int v: adj[c]){
    if(!is cen[v]) {
      decompose(v,c,d+1);
```

```
}
}
decompose();
```

12 CONVOLUTION

```
## FFT
struct cplx {
 ld_a, b;
  cplx(ld a=0, ld b=0):a(a), b(b) {}
  const cplx operator + (const cplx &z) const {

¬ return cplx(a+z.a, b+z.b); }

  const cplx operator - (const cplx &z) const {

→ return cplx(a-z.a, b-z.b); }

 const cplx operator * (const cplx &z) const {
\rightarrow return cplx(a*z.a-b*z.b, a*z.b+b*z.a); }
 const cplx operator / (const ld &k) const {
- return cplx(a/k, b/k); }
const ld PI=acos(-1);
vector<int> rev:
void pre(int sz){
  if(rev.size()==sz) return ;
  rev.resize(sz);
  rev[0]=0;
  int lg n = builtin ctz(sz);
  for (int i = 1; i < sz; ++i) rev[i] = (rev[i>>1]
\rightarrow >> 1) | ((i&1)<<(lq n-1));
|void fft(vector<cplx> &a, bool inv){
  int n = a.size();
  for (int i = 1; i < n-1; ++i) if(i < rev[i])

→ swap(a[i], a[rev[i]]);

  for (int len = 2; len <= n; len <<= 1) {
    ld t = 2*PI/len*(inv? -1: 1);</pre>
    cplx wlen = {cosl(t), sinl(t)};
    int st = 0;
    for (int st = 0; st < n; st += len){</pre>
      cplx w(1);
      for (int i = 0; i < len/2; ++i){
        cplx ev = a[st+i];
        cplx od = a[st+i+len/2]*w;
        a[st+i] = ev+od:
        a[st+i+len/2] = ev-od;
        w = w*wlen;
  if(inv){
    for(cplx &z: a){
      z = z/n;
vector<ll> mul(vector<ll> &a, vector<ll> &b){
  int n = a.size(), m = b.size(), sz = 1;
  while (sz < n+m-1) sz <<= 1;
  vector<cplx> x(sz), y(sz), z(sz);
  for (int i = 0; i < sz; ++i){
    x[i] = cplx(i < n? a[i]: 0, 0);
    y[i] = cplx(i < m? b[i]: 0, 0);
  pre(sz);
```

```
fft(x, 0);
 fft(y, 0);
 for (int i = 0; i < sz; ++i){
   z[i] = x[i] * y[i];
 fft(z, 1);
 vector<ll> c(n+m-1);
 for (int i = 0; i < n+m-1; ++i) {
    c[i] = round(z[i].a);
 return c;
## NTT
const int mod = 998244353;
const int root = 15311432;
const int k = 1 << 23;
int root 1;
vector<int> rev;
ll bigmod(ll a, ll b, ll mod){
 a %= mod;
 ll ret = 1;
 while(b){
    if(b\&1) ret = ret*a%mod;
    a = a*a*mod;
    b >>= 1;
  return ret;
void pre(int sz){
  root 1 = bigmod(root, mod-2, mod);
  if(rev.size()==sz) return ;
  rev.resize(sz);
  rev[0]=0;
 int lq n = builtin ctz(sz);
 for (int i = 1; i < sz; ++i) rev[i] = (rev[i>>1]
   >> 1) | ((i\&1) << (lq n-1));
void fft(vector<int> &a, bool inv){
 int n = a.size();
 for (int i = 1; i < n-1; ++i) if(i<rev[i])

→ swap(a[i], a[rev[i]]);

 for (int len = 2; len <= n; len <<= 1) {
    int wlen = inv ? root 1 : root;
    for (int i = len; i < k; i <<= 1){
      wlen = 1ll*wlen*wlen%mod;
    for (int st = 0; st < n; st += len) {
      int w = 1;
      for (int j = 0; j < len / 2; j++) {
        int ev = a[st+j];
        int od = 1ll*a[st+j+len/2]*w%mod;
        a[st+i] = ev + od < mod ? ev + od : ev + od
   - mod:
        a[st+j+len/2] = ev - od >= 0 ? ev - od : ev
    - od + mod;
       w = 1li * w * wlen % mod;
```

```
if (inv) {
    int n = bigmod(n, mod-2, mod);
    for (int \& x : a)
      x = 111*x*n 1%mod;
vector<int> mul(vector<int> &a, vector<int> &b){
  int n = a.size(), m = b.size(), sz = 1;
  while (sz < n+m-1) sz <<= 1;
  vector<int> x(sz), y(sz), z(sz);
  for (int i = 0; i < sz; ++i){
    x[i] = i < n? a[i]: 0;
    y[i] = i < m? b[i]: 0;
  pre(sz);
  fft(x, 0);
  fft(y, 0);
  for (int i = 0; i < sz; ++i){
    z[i] = 111* x[i] * y[i] % mod;
  fft(z, 1);
  z.resize(n+m-1);
  return z;
## Any mod
const int N = 3e5 + 9, mod = 998244353;
struct base {
  double x, y;
  base() { x = y = 0; }
  base(double x, double y): x(x), y(y) { }
inline base operator + (base a, base b) { return
\rightarrow base(a.x + b.x, a.y + b.y);
inline base operator - (base a, base b) { return
→ base(a.x - b.x, a.y - b.y); }
inline base operator * (base a, base b) { return
   base(a.x * b.x - a.y * b.y, a.x * b.y + a.y *
 → b.x); }
inline base conj(base a) { return base(a.x, -a.y); }
int lim = 1;
vector<br/>vector<br/>vector<br/>se> roots = \{\{0, 0\}, \{1, 0\}\};
vector < int > rev = \{0, 1\}
const double PI = acosl(- 1.0);
void ensure base(int p) {
  if(p <= lim) return;</pre>
  rev.resize(1 << p);
  for(int i = 0; i < (1 << p); i++) rev[i] = (rev[i
\rightarrow >> 1] >> 1) + ((i & 1) << (p - 1));
  roots.resize(1 << p);
  while(lim < p) {</pre>
    double angle = 2 * PI / (1 << (lim + 1));
    for(int i = 1 << (lim - 1); i < (1 << lim);
 roots[i << 1] = roots[i];
      double angle i = angle * (2 * i + 1 - (1 <<
      roots[(i \ll 1) + 1] = base(cos(angle i),
   sin(angle i));
    lim++;
```

```
|void fft(vector<base> \deltaa, int n = -1) {
  if(n == -1) n = a.size():
  assert((n \& (n - 1)) == 0);
  int zeros = builtin ctz(n);
  ensure base(zeros);
  int shift = lim - zeros;
  for(int i = 0; i < n; i++) if(i < (rev[i] >>
shift)) swap(a[i], a[rev[i] >> shift]);
  for(int k = 1; k < n; k <<= 1) {
    for(int i = 0; i < n; i += 2 * k) {
      for(int j = 0; j < k; j++) {
  base z = a[i + j + k] * roots[j + k];</pre>
        a[i + j + k] = a[i + j] - z;
        a[i + j] = a[i + j] + z;
//eq = 0: 4 FFTs in total
//eq = 1: 3 FFTs in total
|vector<int> multiply(vector<int> &a, vector<int>
\rightarrow &b, int eq = 0)
  int need = a.size() + b.size() - 1;
  int p = 0:
  while((1 << p) < need) p++;
  ensure base(p);
  int sz = 1 \ll p;
  vector<base> A, B;
  if(sz > (int)A.size()) A.resize(sz);
  for(int i = 0; i < (int)a.size(); i++) {</pre>
    int x = (a[i] \% mod + mod) \% mod;
    A[i] = base(x \& ((1 << 15) - 1), x >> 15);
  fill(A.begin() + a.size(), A.begin() + sz,
\rightarrow base\{0, 0\});
  fft(A, sz);
  if(sz > (int)B.size()) B.resize(sz);
  if(eq) copy(A.begin(), A.begin() + sz, B.begin());
  else -
    for(int i = 0; i < (int)b.size(); i++) {</pre>
      int x = (b[i] \% mod + mod) \% mod;
      B[i] = base(x \& ((1 << 15) - 1), x >> 15);
    fill(B.begin() + b.size(), B.begin() + sz,
    base{0, 0});
    fft(B, sz);
  double ratio = 0.25 / sz;
  base r^{2}(0, -1), r^{3}(ratio, 0), r^{4}(0, -ratio),
\rightarrow r5(0, 1);
 for(int i = 0; i \le (sz >> 1); i++) {
    int i = (sz - i) \& (sz - 1);
    base a1 = (A[i] + conj(A[j])), a2 = (A[i] -
   conj(A[j])) * r2;
    base b1 = (B[i] + conj(B[j])) * r3, b2 = (B[i])
   - coni(B[i])) * r4;
    if(i != j) {
      base c1 = (A[j] + conj(A[i])), c2 = (A[j] -
   conj(A[i])) * r2;
      base d1 = (B[j] + conj(B[i])) * r3, d2 =
    (B[i] - conj(B[i])) * r4;
      A[i] = c1 * d1 + c2 * d2 * r5;
      B[i] = c1 * d2 + c2 * d1;
```

```
A[i] = a1 * b1 + a2 * b2 * r5;
    B[j] = a1 * b2 + a2 * b1;
 fft(A, sz); fft(B, sz);
 vector<int> res(need);
 for(int i = 0; i < need; i++) {
   long long aa = A[i].x + 0.5;
    long long bb = B[i].x + 0.5;
    long long cc = A[i].y + 0.5;
    res[i] = (aa + ((bb \% mod) << 15) + ((cc \% mod))
   << 30))%mod:
 return res;
vector<int> pow(vector<int>& a, int p) {
 vector<int> res:
 res.emplace back(1);
 while(p) {
   if(p \& 1) res = multiply(res, a);
   a = multiply(a, a, 1);
   p >>= 1;
 return res;
int main() {
 int n, k; cin >> n >> k;
 vector<int> a(10, 0);
 while(k--) {
   int m; cin >> m;
   a[m] = 1;
 vector < int > ans = pow(a, n / 2):
 int res = 0;
 for(auto x: ans) res = (res + 1LL * x * x % mod)

→ % mod:

 cout << res << '\n';
 return 0;
## Online NTT
void solve() {
 f[0]=1; // base case
 for(int i=0; i<=MAX; i++) {
    // Doing the part 1\,
    f[i+1]=(f[i+1]+f[i]*A[0])%mod;
    f[i+2]=(f[i+2]+f[i]*A[1])%mod;
    if(!i) continue;
    // part 2
    int limit=(i&-i);
    for(int p=2; p<=limit; p*=2) {
      convolve(i-p, i-1, p, min(2*p-1, MAX));
void convolve(int l1, int r1, int l2, int r2) {
 int n=max(r1-l1+1,r2-l2+1);
 int t=1:
 while(t<n) t<<=1;
 n=t:
 vector<ll> a(n), b(n);
 for(int i=l1; i<=r1; i++) a[i-l1]=f[i];
 for(int i=l2; i<=r2; i++) b[i-l2]=A[i];
 vector<ll> ret=fft::multiply(a,b);
    for(int i=0; i<ret.size(); i++) {</pre>
```

```
6
```

```
int idx=i+l1+l2+1;
    if(idx>MAX) break;
    // adding to the appropriate entry
    f[idx]+=ret[i];
    f[idx]%=mod;
## FWHT (AND, OR, XOR)

    Time complexity: O(nlogn)

- AND, OR works for any modulo, XOR works for only

→ prime

    size must be power of two

const ll mod = 998244353:
int add (int a, int b) {
  return a + b < mod? a + b: a + b - mod;
int sub (int a, int b) {
  return a - b >= 0? a - b: a - b + mod;
ll poww (ll a, ll p, ll mod){
  a \% = mod;
  ll ret = 1;
  while (p){
   if (p & 1) {
  ret = ret * a % mod;
    a = a * a % mod:
    p >>= 1;
  return ret;
void fwht(vector<int> &a, int inv, int f) {
  int sz = a.size();
  for (int len = 1; 2 * len <= sz; len <<= 1) {
    for (int i = 0; i < sz; i += 2 * len) {
      for (int j = 0; j < len; j++) {
        int x = a[i + j];
        int y = a[i + j + len];
        if (f == 0) {
          if (!inv) a[i + j] = y, a[i + j + len] =
\rightarrow add(x, _y);
          else a[i + j] = sub(y, x), a[i + j +
\rightarrow len] = x;
        else if (f == 1) {
          if (!inv) a[i + j + len] = add(x, y);
          else a[i + j + len] = sub(y, x);
        else {
          a[i + i] = add(x, y);
          a[i + j + len] = sub(x, y);
vector<int> mul(vector<int> a, vector<int> b, int
\rightarrow f) { // 0:AND, 1:OR, 2:XOR
  int sz = a.size();
  fwht(a, 0, f); fwht(b, 0, f);
```

```
vector<int> c(sz);
  for (int i = 0; i < sz; ++i) {
    c[i] = 111 * a[i] * b[i] % mod;
  fwht(c, 1, f);
  if (f) {
    int sz inv = poww(sz, mod - 2, mod);
    for (int i = 0; i < sz; ++i) {
      c[i] = 1ll * c[i] * sz inv % mod;
  return c;
## subset convolution
vector<int> subset conv (vector<int> a, vector<int>
 int n = a.size();
  int lq = log2(n);
  vector<int> cnt(n);
  vector<vector<int>> fa(lg + 1, vector<int> (n)),
    fb(lg + 1, vector<int> (n)), g(lg + 1,
\stackrel{\sim}{\rightarrow} vector<int> (n));
 for (int i = 0; i < n; ++i) {
  cnt[i] = cnt[i >> 1] + (i & 1);
    fa[cnt[i]][i] = a[i] % mod;
fb[cnt[i]][i] = b[i] % mod;
  for (int k = 0; k <= lg; ++k)
    fwht(fa[k], 0, 1); fwht(fb[k], 0, 1);
  for (int k = 0; k \le lg; ++k) {
    for (int j = 0; j \le k; ++j) {
      for (int i = 0; i < n; ++i) {
        g[k][i] = add(g[k][i], 1ll^* fa[j][i] *
\rightarrow fb[k - j][i] % mod);
  for (int k = 0; k \le lg; ++k) {
    fwht(g[k], 1, 1);
  vector<int> c(n);
  for (int i = 0; i < n; ++i) {
    c[i] = q[cnt[i]][i];
  return c;
13 CPP
```

```
maf.max load_factor(0.25);
## gp hash_table:
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
struct chash{
   int operator()(ii p) const {
     return p.first*31 + p.second;
   };
gp_hash_table<ii, int, chash> cnt;
mt19937 rng(chrono::steady_clock::now().time_since_j
   epoch().count());
int x = rng() % 495;
```

14 DETERMINANT

```
const double EPS = 1E-9;
int n;
vector < vector<double> > a (n, vector<double> (n));
double det = 1;
for (int i=0; i<n; ++i) {
 int k = i;
 for (int j=i+1; j<n; ++j)
    if (abs (a[j][i]) > abs (a[k][i]))
  if (abs (a[k][i]) < EPS) {
    det = 0;
    break;
  swap (a[i], a[k]);
  if (i != k)
    det = -det:
  det *= a[i][i];
 for (int j=i+1; j<n; ++j)
a[i][j] /= a[i][i];</pre>
  for (int j=0; j<n; ++j)
    if (j != i \&\& abs (a[j][i]) > EPS)
      for (int k=i+1; k<n; ++k)
        a[j][k] -= a[i][k] * a[j][i];
```

15 DINIC

```
// V^2E, sqrt(E)E, sqrt(V)E(bpm)
// Effective flows are adj[u][3] where adj[u][3] > 0
ll get max flow(vector<array<int, 3>> edges, int n,
\rightarrow int s, int t) {
 vector<array<li, 4>> adj[n];
 for (auto [u, v, c]: edges) {
   adj[u].push back({v, (int)adj[v].size(), c, 0});
   adj[v].push_back({u, (int)adj[u].size() - 1, 0,
   0});
 ll max flow = 0;
 while (true) {
   queue<int> q; q.push(s);
   vector<int> dis(n, -1); dis[s] = 0;
   while (!q.empty()) {
     int u = q.front(); q.pop();
     for (auto [v, idx, c, f]: adj[u]) {
```

```
7
```

```
if (dis[v] == -1 \text{ and } c > f) {
          q.push(v);
          dis[v] = dis[u] + 1;
   if (dis[t] == -1) break;
   vector<int> next(n);
   function<ll(int, ll) > dfs = [\&] (int u, ll)
→ flow) {
     if (u == t) return flow;
     while (next[u] < adj[u].size()) {
  auto &[v, idx, c, f] = adj[u][next[u]++];</pre>
        if (c > f \text{ and } dis[v] == dis[u] + 1) {
          ll bn = dfs(v, min(flow, c - f));
          if (bn > 0) {
            f += bn;
            adj[v][idx][3] -= bn;
            return bn;
       }
     return Oll;
   while (ll flow = dfs(s, LLONG MAX)) {
     max flow += flow;
 return max flow;
```

16 DOMINATOR TREE

```
const int N = 2e5+5;
vector <int> q[N], rg[N], dtree[N], bucket[N];
int sdom[N], par[N], dom[N], dsu[N], lab[N],

    arr[N], rev[N], dpar[N], n, ts, src;

void init(int n, int s) {
 ts = 0, n = n, src = s;
 for (int i = 1; i \le n; ++i) {
    g[i].clear(), rg[i].clear(), dtree[i].clear(),
 → bucket[i].clear();
    sdom[i]=par[i]=dom[i]=dsu[i]=lab[i]=arr[i]=rev[_
   i]=dpar[i]=0;
void dfs(int u) {
   ts++; arr[u] = ts; rev[ts] = u;
   lab[ts] = sdom[ts] = dsu[ts] = ts;
   for(int \&v : g[u]) {
      if(!arr[v]) { dfs(v); par[arr[v]] = arr[u]; }
      rg[arr[v]].push back(arr[u]);
   }
inline int root(int u, int x = 0) {
   if(u == dsu[u]) return x ? -1 : u;
   int v = root(dsu[u], x + 1);
   if(v < 0) return u;
   if(sdom[lab[dsu[u]]] < sdom[lab[u]]) lab[u] =</pre>
   lab[dsu[u]];
   dsu[u] = v; return x ? v : lab[u];
```

```
void build() {
   dfs(src);
   for(int i=n; i; i--) {
      for(int j : rq[i]) sdom[i] =
   min(sdom[i],sdom[root(j)]);
      if(i > 1) bucket[sdom[i]].push back(i);
      for(int w : bucket[i]) {
         int v = root(w);
         if(sdom[v] == sdom[w]) dom[w] = sdom[w];
         else dom[w] = v;
      if(i > 1) dsu[i] = par[i];
   for(int i=2; i<=n; i++) {
      int &dm = dom[i];
      if(dm ^ sdom[i]) dm = dom[dm];
dtree[rev[i]].push_back(rev[dm]);
      dtree[rev[dm]].push back(rev[i]);
      dpar[rev[i]] = rev[dm];
```

17 DP_ON_TREE

```
// Rerooting Technique
vector<array<ll, 2>> down(N), up(N);
void dfs() {
    // calculate down dp
}
void dfs2() {
    ll pref = ?;
    for (auto v: adj[u]) {
        // update up[v] and pref
    }
    reverse(adj[u].begin(), adj[u].end());
    ll suf = ?;
    for (auto v: adj[u]) {
        // update up[v] and suf
    }
    for (auto v: adj[u]) {
        dfs2(v)
    }
}
```

18 DP_OPTIMIZATION

```
## CHT
## Dynamic CHT
const ll IS QUERY = -(1LL << 62);</pre>
struct line {
  ll m. b:
  mutable function <const line*()> succ;
  bool operator < (const line &rhs) const {</pre>
    if (rhs.b != IS QUERY) return m < rhs.m;</pre>
    const line *s = succ();
    if (!s) return 0;
    ll x = rhs.m;
    return b - s -> b < (s -> m - m) * x;
|struct CHT : public multiset <line> {
  bool bad (iterator y) {
    auto z = next(y);
    if (y == begin()) {
      if (z == end()) return 0;
```

```
return y \rightarrow m == z \rightarrow m \&\& y \rightarrow b <= z \rightarrow b;
    auto x = prev(y);
    if (z == end()) return y -> m == x -> m &  y ->
 \rightarrow b <= x -> b;
    return 1.0 \div (x -> b - y -> b) * (z -> m - y ->
    m) >= 1.0 * (y -> b - z -> b) * (y -> m - x ->
   m);
  void add (ll m, ll b) {
    auto y = insert({m, b});
    y \rightarrow succ = [=] \{return \ next(y) == end() ? 0 :
    &*next(y);};
    if (bad(y)) {erase(y); return;}
    while (next(y) != end() \&\& bad(next(y)))
   erase(next(y));
    while (y != begin() && bad(prev(y)))
    erase(prev(y));
  ll eval (ll x) {
    auto l = *lower bound((line) {x, IS_QUERY});
    return l.m * x \mp l.b;
  To find maximum
CHT cht;
cht.add(m, c);
ly max = cht.eval(x);
/7 To find minimum
CHT cht:
cht.add(-m, -c);
y min = -cht.eval(x);
// Divide an array into k parts
// Minimize the sum of squre of each subarray
ll pref[N], dp[N][N];
void compute(int l, int r, int j, int kl, int kr) {
  if (l > r) return ;
  int m = (l + r) / 2;
  array<ll, 2> best = {LLONG MAX, -1};
  for (int k = kl; k \le min(m - 1, kr); ++k) {
best = min(best, \{dp[k][j - 1] + (pref[m] - 1, kr)\}
    pref[k]) * (pref[m] - pref[k]), k});
  dp[m][j] = best[0];
  compute(l, m - 1, j, kl, best[1]);
  compute(m + 1, r, j, best[1], kr);
  Divide an array into n parts.
// Cost of each division is subarray sum
// Minimize the cost
ll dp[n][n], opt[n][n];
for (int i = 0; i < n; ++i) {
  for (int j = 0; j < n; ++j) {
    dp[i][j] = LLONG MAX;
  opt[i][i] = i;
dp[i][i] = 0;
for (int i = n - 2; i >= 0; --i) {
  for (int j = i + 1; j < n; ++j) {
    for (int k = opt[i][j - 1]; k <= min(j - 1ll,</pre>
    opt[i + 1][j]); ++k) {
```

```
if (dp[i][j] >= dp[i][k] + dp[k + 1][j] +
    (pref[j + 1] - pref[i])) {
        dp[i][j] = dp[i][k] + dp[k + 1][j] +
   (pref[j + 1] - pref[i]);
        opt[i][j] = k;
cout << dp[0][n - 1] << "\n";
## Lichao Tree
const int N = int(5e4 + 2);
const ll INF = ll(1e17);
vector<vector<ll> > tree(4*N, {0, INF});
ll f(vector<ll> line, int x){
return line[0] * x + line[1];
void insert(vector<ll> line, int lo = 1, int hi =
\rightarrow N, int i = 1){
 int m = (lo + hi) / 2;
bool left = f(line, lo) < f(tree[i], lo);</pre>
  bool mid = f(line, m) < f(tree[i], m);</pre>
  if(mid) swap(tree[i], line);
  if(hi - lo == 1) return;
  else if(left != mid) insert(line, lo, m, 2*i);
  else insert(line, m, hi, 2*i+1);
ll query(int x, int lo = 1, int hi = N, int i = 1){
  int m = (lo+hi)/2;
  ll curr = f(tree[i], x);
  if(hi-lo==1) return curr;
  if(x<m) return min(curr, query(x, lo, m, 2*i));</pre>
  else return min(curr, query(x, m, hi, 2*i+1));
```

19 DSU ON TREE

```
void dfs(int u, int p) {
  node[tt] = u;
  tin[u] = tt++, sz[u] = 1, hc[u] = -1;
  for (auto v: adj[u]) {
   if (v != p) {
      dfs(v, u);
      sz[u] += sz[v];
      if (hc[u] == -1 \text{ or } sz[hc[u]] < sz[v]) {
        hc[u] = v;
  tout[u] = tt - 1;
void dsu(int u, int p, int keep) {
  for (int v: adj[u]) {
   if (v != p and v != hc[u]) {
      dsu(v, u, 0);
  if (hc[u] != -1) {
    dsu(hc[u], u, 1);
  for (auto v: adj[u]) {
   if (v != p and v != hc[u]) {
      for (int i = tin[v]; i <= tout[v]; ++i) {</pre>
        int w = node[i];
```

```
// get ans in case of ans is related to
- simple path or pair
}
for (int i = tin[v]; i <= tout[v]; ++i) {
    int w = node[i];
    // Add contribution of node w
}
}

// Add contribution of node u
// get ans in case ans is related to subtree
if (!keep) {
    for (int i = tin[u]; i <= tout[u]; ++i) {
        int w = node[i];
        // remove contribution of node w
    }
// Data structure in initial state (empty
- contribution)
}
dfs(0, 0); dsu(0, 0, 0);</pre>
```

20 DS_TRICKS

21 DYNAMIC_CONNECTIVITY

```
const int Q = 1e5+5;
|vector<array<<mark>int</mark>, 2>> t[4 * 0];
vector<int> ans(0);
int q;
struct DSU {
  int n, comps;
  vector<int> par, rnk;
  stack<array<int, 4>> ops;
  DSU(){}
  DSU(int n): n(n), comps(n), par(n), rnk(n) {
    iota(par.begin(), par.end(), 0);
  int find(int u) {
    return (par[u] == u)? u: find(par[u]);
  bool unite(int u, int v)
    u = find(u), v = find(v);
    if (u == v) return false;
    comps - -;
```

```
if (rnk[u] > rnk[v]) swap(u, v);
    ops.push({u, rnk[u], v, rnk[v]});
    par[u] = v;
    if (rnk[u] == rnk[v]) rnk[v]++;
    return true;
 void rollback() 
   if (ops.empty()) return ;
    auto [u, rnku, v, rnkv] = ops.top(); ops.pop();
    par[u] = u, rnk[u] = rnku;
    par[v] = v, rnk[v] = rnkv;
    comps++;
} dsu;
void add(int l, int r, array<int, 2> ed, int u = 1,
\rightarrow int s = 0, int e = q) {
 if (r < s or e < l) return ;</pre>
 if (l <= s and e <= r) {
   t[u].push back(ed);
    return :
 int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
 add(l, r, ed, v, s, m);
 add(l, r, ed, w, m + 1, e);
void qo(int u = 1, int s = 0, int e = q) {
 int rmv = 0;
 for (auto &ed: t[u]) rmv += dsu.unite(ed[0],
\rightarrow ed[1]);
 if (s == e) ans[s] = dsu.comps;
 else {
   int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
    go(v, s, m);
    go(w, m + 1, e);
 while (rmv--) dsu.rollback();
```

8

22 EULER_WALK

```
## Directed graph
vector<int> euler cycle(vector<int> *adj, int s =
→ 0) {
 vector<int> cycle;
 function<void(int)> dfs = [\&] (int u) {
   while (!adj[u].empty()) {
     int v = adi[u].back();
     adj[u].pop back();
     dfs(v);
   cycle.push back(u);
 dfs(s);
 reverse(cycle.begin(), cycle.end());
 return cycle;
## Undirected graph
vector<int> euler cycle(vector<int> *adj,
   vector<int> *des idx, vector<int> *done, int s
\rightarrow = 0) {
 vector<int> cycle;
```

```
function\langle void(int) \rangle dfs = [\&] (int u) {
    while (!adj[u].empty()) {
      int i = adj[u].size() - 1;
      if (done[u][i]) {
        adj[u].pop back();
        continue;
      int v = adj[u][i];
      adj[u].pop back();
      done[u][i] = 1;
      done[v][des idx[u][i]] = 1;
      dfs(v);
    cycle.push back(u);
  dfs(s);
  return cycle;
int n, m; cin >> n >> m;
vector<int> adj[n], des idx[n], done[n];
vector<int> deg(n);
for (int e = 0; e < m; ++e) {
  int u, v; cin >> u >> v; u--, v--;
  des idx[u].push back(adj[v].size());
  des_idx[v].push_back(adj[u].size());
  adj[u].push back(v);
  adj[v].push_back(u);
  done[u] push back(0);
  done[v].push_back(0);
  deg[u]++, de\overline{g}[v]++;
for (int u = 0; u < n; ++u) {
  if (deg[u] & 1)
    cout << "IMPOSSIBLE\n";</pre>
    return ;
vector<int> cycle = euler cycle(adj, des idx, done,
if (cvcle.size() != m + 1) {
  cout << "IMPOSSIBLE\n";</pre>
  return ;
```

23 GEOMETRY

```
#include<bits/stdc++.h>
using namespace std;

const int N = 3e5 + 9;

const double inf = 1e100;
const double eps = 1e-9;
const double PI = acos((double)-1.0);
int sign(double x) { return (x > eps) - (x < -eps);
    }

struct PT {
    double x, y;
    PT() { x = 0, y = 0; }
    PT(double x, double y) : x(x), y(y) {}
    PT(const PT &p) : x(p.x), y(p.y) {}</pre>
```

```
PT operator + (const PT &a) const { return PT(x
    + a.x, y + a.y); }
    PT operator - (const PT &a) const { return PT(x
    - a.x, y - a.y); }
    PT operator * (const double a) const { return
    PT(x * a, y * a); }
    friend PT operator * (const double &a, const PT
    &b) { return PT(a * b.x, a * b.y); }
    PT operator / (const double a) const { return
    PT(x / a, y / a); }
    bool operator == (PT a) const { return sign(a.x
    -x) == 0 \&\& sign(a.y - y) == 0; }
    bool operator != (PT a) const { return !(*this
    == a): }
    bool operator < (PT a) const { return sign(a.x</pre>
    - x) == 0 ? y < a.y : x < a.x; }
bool operator > (PT a) const { return sign(a.x
    -x) == 0 ? y > a.y : x > a.x; }
    double norm() { return sqrt(x * x + y * y); }
    double norm2() \{ return \dot{x} * \dot{x} + \dot{y} * \dot{y}; \}
    PT perp() { return PT(-y, x); } double arg() { return atan2(y, x); }
    PT truncate(double r) { // returns a vector
   with norm r and having same direction
        double k = norm();
        if (!sign(k)) return *this;
        r /= k;
        return PT(x * r, y * r);
inline double dot(PT a, PT b) { return a.x * b.x +
 \rightarrow a.v * b.v; }
inline double dist2(PT a, PT b) { return dot(a - b,
 \rightarrow a - b); }
\rightarrow - b, a - b)); }
inline double cross(PT a, PT b) { return a.x * b.y
\rightarrow - a.y * b.x; }
inline double cross2(PT a, PT b, PT c) { return

    cross(b - a, c - a); }

inline int orientation(PT a, PT b, PT c) { return
\rightarrow sign(cross(b - a, c - a)); }
PT perp(PT a) { return PT(-a',y, a.x); }
PT rotateccw90(PT a) { return PT(-a.y, a.x); }
PT rotatecw90(PT a) { return PT(a.y, -a.x); }
PT rotateccw(PT a, double t) { return PT(a.x *
    cos(t) - a.y * sin(t), a.x * sin(t) + a.y *
   cos(t)); }
PT rotatecw(PT a, double t) { return PT(a.x *
    cos(t) + a.y * sin(t), -a.x * sin(t) + a.y *
    cos(t));
double SQ(double x) { return x * x; }
double rad to deg(double r) { return (r * 180.0 /
double deg to rad(double d) { return (d * PI /
   180.0); }
double get angle(PT a, PT b) {
    double costheta = dot(a, b) / a.norm() /
    return acos(max((double)-1.0, min((double)1.0,
    costheta)));
```

```
bool is point in angle(PT b, PT a, PT c, PT p) { //
does point p lie in angle <bac</p>
    assert(orientation(a, b, c) != 0);
    if (orientation(a, c, b) < 0) swap(b, c);
    return orientation(a, c, p) \geq 0 \& \&
  orientation(a, b, p) <= 0;
bool half(PT p) {
    return p.y > 0.0 \mid \mid (p.y == 0.0 \&\& p.x < 0.0);
void polar sort(vector<PT> &v) { // sort points in

→ counterclockwise

    sort(v.begin(), v.end(), [](PT a,PT b) {
        return make tuple(half(a), 0.0, a.norm2())
  < make tuple(half(b), cross(a, b), b.norm2());</pre>
    });
struct line {
    PT a, b; // goes through points a and b
    PT v; double c; //line form: direction vec
    [cross](x, y) = c
    line() {}
    //direction vector v and offset c
 line(PT v, double c) : v(v), c(c) {
        auto p = get points();
        a = p.first; b = p.second;
 // equation ax + by + c = 0
line(double a, double b, double c) : v({ b,
\rightarrow - a}), c(- c) {
 auto p = get points();
        a = p.first; b = p.second;
 // goes through points p and q
line(PT p, PT q): v(q - p), c(cross(v, p)), a(p),
\rightarrow b(q) {}
pair<PT, PT> get points() { //extract any two
→ points from this line
 PT p, q; double a = -v.y, b = v.x; // ax + by = -c
 if (sign(a) == 0) {
            p = PT(0, -c / b);
            q = PT(1, -c / b);
        else if (sign(b) == 0) {
            p = PT(-c / a, 0);
            q = PT(-c / a, 1);
        else {
            p = PT(0, -c / b);
            q = PT(1, (-c - a) / b);
        return {p, q};
    //ax + by + c = 0
    array<double, 3> get abc() {
        double a = -v.y, b = v.x;
        return {a, b, c};
    // 1 if on the left, -1 if on the right, 0 if
   on the line
    int side(PT p) { return sign(cross(v, p) - c); }
    // line that is perpendicular to this and goes
   through point p
    line perpendicular through(PT p) { return {p, p
    + perp(v)}; }
```

```
// translate the line by vector t i.e. shifting
   it by vector t
    line translate(PT t) { return {v, c + cross(v,
    // compare two points by their orthogonal
 → projection on this line
    // a projection point comes before another if
 → it comes first according to vector v
    bool cmp by projection(PT p, PT q) { return
 \rightarrow dot(v, p) < dot(v, q);
 line shift left(double d)
  PT z = v.\overline{p}erp().truncate(d);
  return line(a + z, b + z);
};
// find a point from a through b with distance d

through b with distance d

through b double d) {
PT point along line(PT a, PT b, double d) {
    return a + (((b - a) / (b - a).norm()) * d);
// projection point c onto line through a and b
→ assuming a != b
PT project from point to line(PT a, PT b, PT c) {
    return a + (b - a) * dot(c - a, b - a) / (b -
   a).norm2();
// reflection point c onto line through a and b

    assuming a != b

PT reflection from point to line(PT a, PT b, PT c) {
    PT p = project from point to line(a,b,c);
    return point along line(c, p, 2.0 * dist(c, p));
// minimum distance from point c to line through a
double dist from point to line(PT a, PT b, PT c) {
    return fabs(cross(b - a, c - a) / (b -
 \rightarrow a).norm()):
// returns true if point p is on line segment ab
bool is point on seg(PT a, PT b, PT p) {
    if (fabs(cross(p - b, a - b)) < eps) {
        if (p.x < min(a.x, b.x) \mid | p.x > max(a.x, b.x))
 → b.x)) return false;
        if (p.y < min(a.y, b.y) \mid | p.y > max(a.y, b.y) \mid |
   b.y)) return false:
        return true;
    return false;
// minimum distance point from point c to segment

    ab that lies on segment ab

PT project from point to seg(PT a, PT b, PT c) {
    double r = dist2(a, b);
    if (fabs(r) < eps) return a;</pre>
    r = dot(c - a, b - a) / r;
    if (r < 0) return a;
    if (r > 1) return b;
    return a + (b - a) * r;
// minimum distance from point c to segment ab
double dist from point to seg(PT a, PT b, PT c) {
    return dist(c, project from point to seq(a, b,
// 0 if not parallel, 1 if parallel, 2 if collinear
bool is_parallel(PT a, PT b, PT c, PT d) {
```

```
double k = fabs(cross(b - a, d - c));
    if (k < eps){
        if (fabs(cross(a - b, a - c)) < eps &&
    fabs(cross(c - d, c - a)) < eps) return 2;
        else return 1;
    else return 0;
 // check if two lines are same
bool are lines same(PT a, PT b, PT c, PT d) {
    if (fabs(cross(a - c, c - d)) < eps &&
    fabs(cross(b - c, c - d)) < eps) return true;
    return false:
 // bisector vector of <abc
PT angle bisector(PT &a, PT &b, PT &c){
    PT p^{-} = a - b, q = c - b;
    return p + q * sqrt(dot(p, p) / dot(q, q));
 // 1 if point is ccw to the line, 2 if point is cw
 → to the line, 3 if point is on the line
int point line relation(PT a, PT b, PT p) {
    int c = sign(cross(p - a, b - a));
    if (c < 0) return 1;
    if (c > 0) return 2;
    return 3;
// intersection point between ab and cd assuming
   unique intersection exists
bool line line intersection(PT a, PT b, PT c, PT d,
 → PT &ans) {
    double a1 = a.y - b.y, b1 = b.x - a.x, c1 =
   cross(a. b):
    double a2 = c.y - d.y, b2 = d.x - c.x, c2 =
    cross(c, d);
    double det = a1 * b2 - a2 * b1;
    if (det == 0) return 0;
    ans = PT((b1*c2 - b2*c1) / det, (c1*a2 -
    a1 * c2) / det);
    return 1;
 // intersection point between segment ab and
   segment cd assuming unique intersection exists
bool seg seg intersection(PT a, PT b, PT c, PT d,
 → PT &ans) {
    double oa = cross2(c, d, a), ob = cross2(c, d,
    double oc = cross2(a, b, c), od = cross2(a, b,
    if (oa * ob < 0 && oc * od < 0){
        ans = (a * ob - b * oa) / (ob - oa);
        return 1:
    else return 0:
// intersection point between segment ab and
    segment cd assuming unique intersection may not
    exists
   se.size()==0 means no intersection
// se.size()==1 means one intersection
 // se.size()==2 means range intersection
|set<PT> seg seg intersection inside(PT a, PT b,
    PT c, \overline{PT} d\overline{I} {
    PT ans;
```

```
if (seg seg intersection(a, b, c, d, ans))
    return {ans};
    set<PT> se:
    if (is point on seg(c, d, a)) se.insert(a);
    if (is point on seg(c, d, b)) se.insert(b);
    if (is point on seg(a, b, c)) se.insert(c);
    if (is point on seq(a, b, d)) se.insert(d);
    return se;
   intersection between segment ab and line cd
// 0 if do not intersect, 1 if proper intersect, 2

→ if segment intersect

int seq line relation(PT a, PT b, PT c, PT d) {
    double p = cross2(c, d, a);
    double q = cross2(c, d, b);
    if (sign(p) == 0 \&\& sign(q) == 0) return 2;
    else if (p * q < 0) return 1;
    else return 0;
// intersection between segament ab and line cd

→ assuming unique intersection exists

bool seg line intersection(PT a, PT b, PT c, PT d,
→ PT &ans) {
    bool k = seg line relation(a, b, c, d);
    assert(k = \overline{2});
    if (k) line line intersection(a, b, c, d, ans);
    return k;
// minimum distance from segment ab to segment cd
double dist from seg to seg(PT a, PT b, PT c, PT d)
    PT dummy;
    if (seg seg intersection(a, b, c, d, dummy))
   return 0.0:
    else return min({dist from point to seg(a, b,
    c), dist from point to seg(a, b, d),
        dist from point to seq(c, d, a),
    dist from point to sea(c, d, b)):
// minimum distance from point c to ray (starting
→ point a and direction vector b)
double dist from point to ray(PT a, PT b, PT c) {
    b = a + b:
    double r = dot(c - a, b - a);
    if (r < 0.0) return dist(c, a);</pre>
    return dist from point to line(a, b, c);
// starting point as and direction vector ad
bool ray ray intersection(PT as, PT ad, PT bs, PT
    double dx = bs.x - as.x, dy = bs.y - as.y;
    double det = bd.x * ad.y - bd.y * ad.x;
    if (fabs(det) < eps) return 0;</pre>
    double u = (dv * bd.x - dx * bd.y) / det;
    double v = (dy * ad.x - dx * ad.y) / det;
    if (sign(u) \ge 0 \&\& sign(v) \ge 0) return 1;
    else return 0;
double ray ray distance(PT as, PT ad, PT bs, PT bd)
    if (ray ray intersection(as, ad, bs, bd))
   return 0.0:
    double ans = dist from point to ray(as, ad, bs);
```

```
ans = min(ans, dist from point to ray(bs, bd,
   as));
return ans;
struct circle {
    PT p; double r;
    circle() {}
    circle(PT p, double r): p( p), r( r) {};
    // center (x, y) and \overline{r} adius \overline{r}
    circle(double x, double y, double r): p(PT(x,
- y)), r(_r) {};
// circumcircle of a triangle
    // the three points must be unique
    circle(PT a, PT b, PT c) {
        b = (a + b) * 0.5;
        c = (a + c) * 0.5;
        line line intersection(b, b + rotatecw90(a
   - b), c, c + rotatecw90(a - c), p);
        r = dist(a, p);
    // inscribed circle of a triangle
    circle(PT a, PT b, PT c, bool t) {
        line u, v;
        double m = atan2(b.y - a.y, b.x - a.x), n =
   atan2(c.y - a.y, c.x - a.x);
        u.b = u.a + (PT(cos((n + m)/2.0), sin((n +
 \rightarrow m)/2.0)));
        v.a = b;
        m = atan2(a.y - b.y, a.x - b.x), n =
 → atan2(c.y - b.y, c.x - b.x);
        v.b = v.a + (PT(cos((n + m)/2.0), sin((n +
        line line intersection(u.a, u.b, v.a, v.b,
        r = dist from point to seg(a, b, p);
    bool operator == (circle v) { return p == v.p
   && sign(r - v.r) == 0;
    double area() { return PI * r * r; }
    double circumference() { return 2.0 * PI * r; }
//O if outside, 1 if on circumference, 2 if inside
→ circle
int circle point relation(PT p, double r, PT b) {
    double d = dist(p, b);
    if (sign(d - r) < 0) return 2;
    if (sign(d - r) == 0) return 1;
    return 0;
// 0 if outside, 1 if on circumference, 2 if inside
int circle line relation(PT p, double r, PT a, PT
    double d = dist from point to line(a, b, p);
    if (sign(d - r) < 0) return 2;
    if (sign(d - r) == 0) return 1;
    return 0;
//compute intersection of line through points a and
//circle centered at c with radius r > 0
vector<PT> circle line intersection(PT c, double r,
→ PT a, PT b) {
    vector<PT> ret;
```

```
b = b - a; a = a - c;
    double A = dot(b, b), B = dot(a, b);
    double C = dot(a, a) - r * r, D = B * B - A * C;
    if (D < -eps) return ret;</pre>
    ret.push back(c + a + b * (-B + sgrt(D + eps))
    if (D > eps) ret.push back(c + a + b * (-B - a))
    sqrt(D)) / A);
    return ret:
//5 - outside and do not intersect
//4 - intersect outside in one point
//3 - intersect in 2 points
//2 - intersect inside in one point
//1 - inside and do not intersect
int circle circle relation(PT a, double r, PT b,
   double R) {
    double d = dist(a, b);
    if (sign(d - r - R) > 0) return 5;
    if (sign(d - r - R) == 0) return 4;
    double l = fabs(r - R);
    if (sign(d - r - R) < 0 \&\& sign(d - l) > 0)
    return 3;
    if (sign(d - l) == 0) return 2;
    if (sign(d - l) < 0) return 1;
    assert(0); return -1;
vector<PT> circle circle intersection(PT a, double
    r, PT b, double R) {
    if (a == b \&\& sign(r - R) == 0) return
    {PT(1e18, 1e18)};
    vector<PT> ret;
    double d = sqrt(dist2(a, b));
    if (d > r + R \mid | d + min(r, R) < max(r, R))
    return ret;
    double x = (d * d - R * R + r * r) / (2 * d);
    double y = sqrt(r * r - x * x);
    PT v = (b - a) / d:
    ret.push back(a + v * x + rotateccw90(v) * y);
    if (y > \overline{0}) ret.push back(a + v * x -
    rotateccw90(v) * v):
    return ret;
// returns two circle c1, c2 through points a, b
   and of radius r
// 0 if there is no such circle, 1 if one circle, 2

→ if two circle

int get circle(PT a, PT b, double r, circle &c1,

    circle &c2) {

    vector<PT> v = circle circle intersection(a, r,
    b, r);
    int t = v.size();
    if (!t) return 0;
    c1.p = v[0], c1.r = r;
    if (t == 2) c2.p = v[1], c2.r = r;
    return t:
// returns two circle c1, c2 which is tangent to
   line u, goes through
// point q and has radius r1; 0 for no circle, 1 if
\rightarrow c1 = c2 , 2 if c1 != c2
int get circle(line u, PT q, double r1, circle &c1,

    circle &c2) {
    double d = dist from point to line(u.a, u.b, q);
```

```
if (sign(d - r1 * 2.0) > 0) return 0;
    if (sign(d) == 0) {
  cout << u.v.x << ' ' << u.v.y << '\n';</pre>
        c1.p = q + rotateccw90(u.v).truncate(r1);
        c2.\dot{p} = \dot{q} + rotatecw90(u.v).truncate(r1);
        c1.r = c2.r = r1;
        return 2;
    line u1 = line(u.a +
    rotateccw90(u.v).truncate(r1), u.b +
    rotateccw90(u.v).truncate(r1));
    line u2 = line(u.a +
    rotatecw90(u.v).truncate(r1), u.b +
    rotatecw90(u.v).truncate(r1));
    circle cc = circle(q, r1);
    PT p1, p2; vector<PT> v;
    v = circle line intersection(q, r1, u1.a, u1.b);
    if (v.siz\overline{e}()) \overline{v} = circle line intersection(q,
    r1, u2.a, u2.b);
    v.push back(v[0]);
    p1 = v[0], p2 = v[1];
    c1 = circle(p1, r1);
    if (p1 == p2) {
        c2 = c1;
        return 1;
    c2 = circle(p2, r1);
    return 2:
// returns area of intersection between two circles
double circle circle area(PT a, double r1, PT b,

→ double r2) {
    double d = (a - b).norm();
    if(r1 + r2 < d + eps) return 0;
    if(r1 + d < r2 + eps) return PI * r1 * r1;
    if(r2 + d < r1 + eps) return PI * r2 * r2;
    double theta 1 = acos((r1 * r1 + d * d - r2 *
   r2) / (2 * r1 * d)).
     theta 2 = acos((r2 * r2 + d * d - r1 * r1)/(2
    * r2 * d));
return r1 * r1 * (theta_1 - sin(2 *
    theta 1)/2.) + r2 * r2 \bar{*} (theta 2 - \sin(2 *)
    theta 2)/2.);
 / tangent lines from point q to the circle
int tangent lines from point(PT p, double r, PT q,
    line &u, line &v) {
    int x = sign(dist2(p, q) - r * r);
    if (x < 0) return 0; // point in cricle</pre>
    if (x == 0) { // point on circle
        u = line(q, q + rotateccw90(q - p));
        \ddot{v} = u;
        return 1:
    double d = dist(p, q);
    double l = r * r / d;
    double h = sqrt(r * r - l * l);
    u = line(q, p + ((q - p).truncate(l) +
    (rotateccw90(q - p).truncate(h))))
    v = line(q, p + ((q - p).truncate(l) +
    (rotatecw90(q - p).truncate(h))));
    return 2;
// returns outer tangents line of two circles
```

```
// if inner == 1 it returns inner tangent lines
int tangents lines from circle(PT c1, double r1, PT
if (inner) r2 = -r2;
    PT d = c2 - c1;
    double dr = r1 - r2, d2 = d.norm(), h2 = d2 - rac{1}{2}
   dr * dr:
    if (d2 == 0 | | h2 < 0)  {
        assert(h2 != 0);
        return 0:
   vector<pair<PT, PT>>out;
   for (int tmp: {- 1, 1}) {
        PT v = (d * dr + rotateccw90(d) * sqrt(h2)
   * tmp) / d2;
        out.push back(\{c1 + v * r1, c2 + v * r2\});
    u = line(out[0].first, out[0].second);
    if (out.size() == 2) v = line(out[1].first,
   out[1].second);
    return 1 + (h2 > 0);
//O(n^2 \log n)
struct CircleUnion {
    int n;
    double x[2020], y[2020], r[2020];
    int covered[2020];
    vector<pair<double, double> > seg, cover;
    double arc, pol;
    inline int sign(double x) {return x < -eps ? -1</pre>
   : x > eps;
    inline int sign(double x, double y) {return
\rightarrow sign(x - y);}
    inline double SQ(const double x) {return x * x;}
    inline double dist(double x1, double y1, double
   x2, double y2) {return sqrt(SQ(x1 - x2) + SQ(y1
    - y2));}
    inline double angle(double A, double B, double
        double val = (SQ(A) + SQ(B) - SQ(C)) / (2 *

→ A * B);

        if (val < -1) val = -1;
        if (val > +1) val = +1;
        return acos(val):
    CircleUnion() {
        n = 0;
        seg.clear(), cover.clear();
        arc = pol = 0;
    void init() {
        n = 0:
        seg.clear(), cover.clear();
        arc = pol = 0;
    void add(double xx, double yy, double rr) {
        x[n] = xx, y[n] = yy, r[n] = rr, covered[n]
   = 0, n++;
    void getarea(int i, double lef, double rig) {
        arc += 0.5 * r[i] * r[i] * (rig - lef -

    sin(rig - lef));

        double x1 = x[i] + r[i] * cos(lef), y1 =
\rightarrow v[i] + r[i] * sin(lef);
```

```
double x2 = x[i] + r[i] * cos(rig), y2 =
  y[i] + r[i] * sin(rig);
       pol += x1 * y2 - x2 * y1;
   double solve() {
       for (int i = 0; i < n; i++) {
           for (int j = 0; j < i; j++) {
               if (!sign(x[i] - x[j]) &&
   !sign(y[i] - y[j]) && !sign(r[i] - r[j])) {
                    r[i] = 0.0;
                    break;
       for (int i = 0; i < n; i++) {
           for (int j = 0; j < n; j++) {
               if (i != j \&\& sign(r[j] - r[i]) >=
   0 && sign(dist(x[i], y[i], x[j], y[j]) - (r[j]
  - r[i])) <= 0) {
                    covered[i] = 1;
                    break;
       for (int i = 0; i < n; i++) {
           if (sign(r[i]) && !covered[i]) {
                seq.clear();
                for (int j = 0; j < n; j++) {
                    if (i != j) {
                        double d = dist(x[i], y[i],

→ x[j], y[j]);

                        if (sign(d - (r[j] + r[i]))
\Rightarrow >= 0 || sign(d - abs(r[j] - r[i])) <= 0) {
                            continue;
                        double alpha = atan2(y[j] -
\rightarrow y[i], x[i] - x[i]);
                        double beta = angle(r[i],
\rightarrow d, r[j]);
                        pair<double, double>
  tmp(alpha - beta, alpha + beta);
                        if (sign(tmp.first) \leq 0 &&
  sign(tmp.second) \ll 0) {
   seg.push back(pair<double, double>(2 * PI +
   tmp.first, 2 * PI + tmp.second));
                        else if (sign(tmp.first) <</pre>
  0) {
   seg.push back(pair<double, double>(2 * PI +
   tmp.first. 2 * PI)):
   seg.push back(pair<double, double>(0,
  tmp.second));
                        else {
                            seg.push back(tmp);
               sort(seq.begin(), seq.end());
               double rig = 0;
```

```
for (vector<pair<double, double>
    >::iterator iter = seq.begin(); iter !=
   seq.end(); iter++)
                     if (sign(rig - iter->first) >=
    0) {
                         rig = max(rig,

    iter->second);

                     élse {
                         getarea(i, rig,

    iter->first);

                         rig = iter->second;
                 if (!sign(rig)) {
                     arc += r[i] * r[i] * PI;
                 else {
                     getarea(i, rig, 2 * PI);
        return pol / 2.0 + arc;
} CU:
double area of triangle(PT a, PT b, PT c) {
    return \overline{fabs}(cross(b - a, c - a) * 0.5);
// -1 if strictly inside, 0 if on the polygon, 1 if

→ strictly outside

int is point_in_triangle(PT a, PT b, PT c, PT p) {
    if^-(sign(cross(b - a, c - a)) < 0) swap(b, c);
    int c1 = sign(cross(b - a, p - a));
    int c2 = sign(cross(c - b, p - b));
    int c3 = sign(cross(a - c, p - c));
    if (c1<0 || c2<0 || c3 < 0) return 1;</pre>
    if (c1 + c2 + c3 != 3) return 0;
    return -1:
double perimeter(vector<PT> &p) {
    double ans=0; int n = p.size();
    for (int i = 0; i < n; i++) ans += dist(p[i],
    p[(i + 1) % n]);
    return ans;
double area(vector<PT> &p) {
    double ans = 0; int n = p.size();
    for (int i = 0; i < n; i++) ans += cross(p[i],
    p[(i + 1) % n]);
return fabs(ans) * 0.5;
   centroid of a (possibly non-convex) polygon,
   assuming that the coordinates are listed in a
    clockwise or
   counterclockwise fashion. Note that the
    centroid is often known as
   the "center of gravity" or "center of mass".
|PT centroid(vector<PT> &p) {
    int n = p.size(); PT c(0, 0);
    double sum = 0;
    for (int i = 0; i < n; i++) sum += cross(p[i],
    p[(i + 1) % n]);
double scale = 3.0 * sum;
    for (int i = 0; i < n; i++) {
        int j = (i + 1) % n;
```

```
c = c + (p[i] + p[j]) * cross(p[i], p[j]);
    return c / scale;
}
// 0 if cw, 1 if ccw
bool get direction(vector<PT> &p)
    double ans = 0; int n = p.size();
    for (int i = 0; i < n; i++) ans += cross(p[i],
   p[(i + 1) % n]);
    if (sign(ans) > 0) return 1;
    return 0:
// it returns a point such that the sum of distances
// from that point to all points in p is minimum
// O(n log^2 MX)
PT geometric median(vector<PT> p) {
 auto tot dist = [\&](PT z) {
     double res = 0;
     for (int i = 0; i < p.size(); i++) res +=
 \rightarrow dist(p[i], z);
     return res;
 };
 auto findY = [\&] (double x)
     double yl = -1e5, yr = 1e5;
     for (int i = 0; i < 60; i++) {
         double ym1 = yl + (yr - yl) / 3;
         double ym2 = yr - (yr - yl) / 3;
         double d1 = tot dist(PT(x, ym1));
         double d2 = tot dist(PT(x, ym2));
         if (d1 < d2) yr = ym2;
         else yl = ym1;
     return pair<double, double> (yl,
    tot dist(PT(x, yl)));
    double xl = -1e5, xr = 1e5;
    for (int i = 0; i < 60; i++) {
        double xm1 = xl + (xr - xl) / 3;
        double xm2 = xr - (xr - xl) / 3;
        double y1, d1, y2, d2;
        auto z = findY(xm1); y1 = z.first; d1 =

    z.second;

        z = findY(xm2); y2 = z.first; d2 = z.second;
        if (d1 < d2) xr = xm2;
        else xl = xm1:
    return {xl, findY(xl).first };
vector<PT> convex hull(vector<PT> &p) {
if (p.size() \ll \overline{1}) return p;
 vector < PT > v = p:
    sort(v.begin(), v.end());
    vector<PT> up, dn;
    for (auto& p : v) {
        while (up.size() > 1 \&\&
    orientation(up[up.size() - 2], up.back(), p) >=
→ 0) {
            up.pop back();
        while (dn.size() > 1 \&\&
   orientation(dn[dn.size() - 2], dn.back(), p) <=
dn.pop back();
        up.push back(p);
```

```
dn.push back(p);
    \dot{v} = dn;
    if (v.size() > 1) v.pop back();
    reverse(up.begin(), up.end());
    up.pop back();
    for (\overline{auto}\& p : up) {
        v.push back(p);
    if (v.size() == 2 \&\& v[0] == v[1]) v.pop back();
    return v;
 //checks if convex or not
bool is convex(vector<PT> &p) {
    bool s[3]; s[0] = s[1] = s[2] = 0;
    int n = p.size();
    for (int i = 0; i < n; i++) {
        int j = (i + 1) \% n;
        int k = (j + 1) \% n;
        s[sign(cross(p[j] - p[i], p[k] - p[i])) +
   1] = 1;
        if (s[0] && s[2]) return 0;
    return 1;
// -1 if strictly inside, 0 if on the polygon, 1 if
   strictly outside
// it must be strictly convex, otherwise make it

→ strictly convex first

int is point in convex(vector<PT> &p, const PT& x)
    \{ // O(\log n) \}
    int n = p.size(); assert(n >= 3);
    int a = orientation(p[0], p[1], x), b =
    orientation(p[0], p[n - 1], x);
    if (a < 0 | | b > 0) return 1;
    int l = 1, r = n - 1;
    while (l + 1 < r) {
        int mid = l + r >> 1:
        if (orientation(p[0], p[mid], x) \geq 0) l =
   mid;
        else r = mid:
    int k = orientation(p[l], p[r], x);
    if (k <= 0) return -k;
    if (l == 1 \&\& a == 0) return 0;
    if (r == n - 1 \&\& b == 0) return 0;
    return -1:
|bool is point on polygon(vector<PT> &p, const PT&
 \rightarrow z) \overline{\{}
    int n = p.size();
    for (int i = 0; i < n; i++) {
     if (is point on seg(p[i], p[(i + 1) % n], z))
   return \overline{1}:
    return 0;
// returns le9 if the point is on the polygon
int winding number(vector<PT> &p, const PT& z) { //
   O(n)
    if (is point on polygon(p, z)) return 1e9;
    int n = p.size(), ans = 0;
    for (int i = 0; i < n; ++i) {
        int j = (i + 1) % n;
```

```
bool below = p[i].y < z.y;</pre>
        if (below != (p[j].y < z.y)) {
             auto orient = orientation(z, p[i],
\rightarrow p[i]);
            if (orient == 0) return 0;
            if (below == (orient > 0)) ans += below

→ ? 1 : -1:

    return ans;
// -1 if strictly inside, 0 if on the polygon, 1 if

→ strictly outside

int is point in polygon(vector<PT> &p, const PT& z)
\rightarrow { // O(n)
    int k = winding number(p, z);
    return k == 1e9^{-}? 0 : k == 0 ? 1 : -1;
// id of the vertex having maximum dot product with
// polygon must need to be convex
// top - upper right vertex
// for minimum dot prouct negate z and return
\rightarrow -dot(z, p[id])
int extreme vertex(vector<PT> &p, const PT &z,
    const int top) \{ // O(\log n) \}
    int n = p.size();
    if (n == 1) return 0;
 double ans = dot(p[0], z); int id = 0;
    if (dot(p[top], z) > ans) ans = dot(p[top], z),
   id = top:
    int l = 1, r = top - 1;
    while (l < r) {
    int mid = l + r >> 1;
        if (dot(p[mid + 1], z) >= dot(p[mid], z)) l
   = mid + 1;
        else r = mid:
    if (dot(p[l], z) > ans) ans = dot(p[l], z), id
\rightarrow = 1;
    l = top + 1, r = n - 1;
    while (l < r) {
        int mid = l + r \gg 1;
        if (dot(p[(mid + 1) % n], z) >= dot(p[mid],
\rightarrow z)) l = mid + 1;
        else r = mid:
    if (dot(p[l], z) > ans) ans = dot(p[l], z), id
   = l:
    return id:
double diameter(vector<PT> &p) {
    int n = (int)p.size();
    if (n == 1) return 0;
    if (n == 2) return dist(p[0], p[1]);
    double ans = 0:
    int i = 0, j = 1;
    while (i < n) {
        while (cross(p[(i + 1) % n] - p[i], p[(j +
\rightarrow 1) % nl - p[i]) >= 0) {
         ans = max(ans, dist2(p[i], p[j]));
         j = (j + 1) \% n;
```

```
ans = max(ans, dist2(p[i], p[j]));
    return sqrt(ans);
double width(vector<PT> &p) {
    int n = (int)p.size();
    if (n <= 2) return 0;
    double ans = inf;
    int i = 0, j = 1;
   while (i < n) {
        while (cross(p[(i + 1) % n] - p[i], p[(j +
\rightarrow 1) % n] - p[i]) >= 0) i = (i + 1) % n;
        ans = min(ans)
   dist from point to line(p[i], p[(i + 1) % n],
   p[j]));
        1++;
    return ans;
// minimum perimeter
double minimum enclosing rectangle(vector<PT> &p) {
int n = p.size();
if (n <= 2) return perimeter(p);</pre>
int mndot = 0; double tmp = dot(p[1] - p[0], p[0]);
 for (int i = 1; i < n; i++) {
 if (dot(p[1] - p[0], p[i]) <= tmp) {</pre>
  tmp = dot(p[1] - p[0], p[i]);
  mndot = i;
 double ans = inf;
int i = 0, j = 1, mxdot = 1;
while (i < n) {
 PT cur = p[(i + 1) \% n] - p[i];
        while (cross(cur, p[(j + 1) % n] - p[j]) >=
\rightarrow 0) j = (j + 1) % n;
        while (dot(p[(mxdot + 1) % n], cur) >=
\rightarrow dot(p[mxdot], cur)) mxdot = (mxdot + 1) % n;
        while (dot(p[(mndot + 1) % n], cur) <=
\rightarrow dot(p[mndot], cur)) mndot = (mndot + 1) % n;
        ans = min(ans, 2.0 * ((dot(p[mxdot], cur)))
    cur.norm() - dot(p[mndot], cur) / cur.norm()) +
   dist from point to line(p[i], p[(i + 1) % n],
   p[j])));
        1++:
    return ans:
// given n points, find the minimum enclosing
   circle of the points
// call convex hull() before this for faster
// expected O(n)
circle minimum enclosing circle(vector<PT> &p) {
    random shuffle(p.begin(), p.end());
    int n = p.size();
    circle c(p[0], 0);
   for (int i = 1; i < n; i++) {
        if (sign(dist(c.p, p[i]) - c.r) > 0) {
            c = circle(p[i], 0);
            for (int j = 0; j < i; j++) {
                if (sign(dist(c.p, p[j]) - c.r) >
→ 0) {
```

```
dist(p[i], p[j]) / 2);
                     for (int k = 0; k < j; k++) {
                          if (sign(dist(c.p, p[k]) -
 \rightarrow c.r) > 0) {
                              c = circle(p[i], p[j],
 \rightarrow p[k]);
    return c;
## Closest Pair of Points
ll min dis(vector<array<int, 2>> &pts, int l, int
 ¬ r) {
  if (l + 1 >= r) return LLONG MAX;
  int m = (l + r) / 2;
  ll my = pts[m-1][1];
  ll d = min(min dis(pts, l, m), min dis(pts, m,
  inplace merge(pts.begin()+l, pts.begin()+m,
 → pts.begin()+r);
  for (int i = l; i < r; ++i) {
  if ((pts[i][1] - my) * (pts[i][1] - my) < d) {</pre>
      pts[j][0]) * (pts[i][0] - pts[j][0]) < d; ++j) {
         [l] dx = pts[i][0] - pts[j][0], dy =
    pts[i][1] - pts[i][1];
         d = min(d, dx * dx + dy * dy);
  return d;
|vector<array<<mark>int</mark>, 2>> pts(n);
sort(pts.begin(), pts.end(), [\&] (array<int, 2> a,
 \rightarrow array<int, 2> b){
  return make pair(a[1], a[0]) < make pair(b[1],</pre>
 \rightarrow b[0]);
});
## Angular Sort
inline bool up (point p) {
 return p.y > 0 or (p.y == 0 \text{ and } p.x >= 0);
|sort(v.begin(), v.end(), [] (point a, point b) {
  return up(a) == up(b) ? a.x * b.y > a.y * b.x :
 \rightarrow up(a) < up(b);
});
## Convex Hull
struct pt {
 int x, y;
ll cross(pt a, pt b, pt c) { //ab*ac
  return 1ll*(b.x-a.x)*(c.y-a.y) -
 \rightarrow 111*(c.x-a.x)*(b.y-a.y);
vector<pt> convexHull(vector<pt>& p) {
  sort(p.begin(), p.end(), [\&] (pt a, pt b) {
    return (a.x==b.x? a.y<b.y: a.x<b.x);
  int n = p.size(), m = 0;
```

c = circle((p[i] + p[j]) / 2,

24 GRAY_CODE

```
int gc(int n){ return n^(n>>1); }
int gc to dec(int g) {
   int d=0;
   while (g) { d ^= g; g >>= 1; }
   return d;
}
```

25 HASHING

```
// Hashing
// Hashvalue(l...r) = hsh[l] - hsh[r + 1] * base ^
   (r - l + 1);
// Must call preprocess
#include<bits/stdc++.h>
using namespace std;
typedef long long ll;
const int MAX = 100009;
ll\ mods[2] = \{10000000007, 10000000009\}
//Some back-up primes: 1072857881, 1066517951,
- 1040160883
ll bases[2] = \{137, 281\};
ll pwbase[3][MAX];
void Preprocess(){
  pwbase[0][0] = pwbase[1][0] = 1;
  for(ll i = 0; i < 2; i++){
    for(ll j = 1; j < MAX; j++)
      pwbase[i][j] = (pwbase[i][j - 1] * bases[i])
    % mods[i];
struct Hashing{
 ll hsh[2][MAX];
  string str;
  Hashing(){}
  Hashing(string _str) {str = _str; memset(hsh, 0,

¬ sizeof(hsh)): build():}

  void Build(){
    for(ll i = str.size() - 1; i >= 0; i--){
      for(int j = 0; j < 2; j++){
  hsh[j][i] = (hsh[j][i + 1] * bases[j] +</pre>

    str[i]) % mods[i];

        hsh[j][i] = (hsh[j][i] + mods[j]) % mods[j];
    }
```

```
pair<ll,ll> GetHash(ll i, ll j){
    assert(i <= j);</pre>
    ll tmp1 = (hsh[0][i] - (hsh[0][i + 1] *
    pwbase[0][j - i + 1]) % mods[0]) % mods[0];
    ll tmp2 = (hsh[1][i] - (hsh[1][i + 1])
   pwbase[1][i - i + 1]) % mods[1]) % mods[1];
    if(tmp1 < 0) tmp1 += mods[0];
    if(tmp2 < 0) tmp2 += mods[1]:
    return make pair(tmp1, tmp2);
};
/***
    * Everything is 0 based
    * Call precal() once in the program
    * Call update(1,0,n-1,i,j,val) to update the

→ value of position

      i to j to val, here n is the length of the
   string
    * Call query(1,0,n-1,L,R) to get a node

→ containing hash

      of the position [L:R]
    * Before any update/query
        - Call init(str) where str is the string to

→ be hashed

        - Call build(1,0,n-1)
namespace strhash {
  int n;
  const int MAX = 100010;
  int ara[MAX]:
  const int MOD[] = {2078526727, 2117566807};
  const int BASE[] = {1572872831, 1971536491};
  int BP[2][MAX], CUM[2][MAX];
  void init(char *str) {
   n = strlen(str);
    for(int i=0;i<n;i++) ara[i] = str[i]-'0'+1; ///</pre>

→ scale str[i] if needed
  void precal()
    BP[0][0] = BP[1][0] = 1;
    for(int i=1;i<MAX;i++)</pre>
      BP[0][i] = (BP[0][i-1] * (long long) BASE[0]
 → ) % MOD[0];
      BP[1][i]' = (BP[1][i-1] * (long long) BASE[1]
     % MOD[1];
  struct node {
    int sz;
    int h[2];
    node() {}
  } tree[4*MAX];
  int lazy[4*MAX];
  inline node Merge(node a, node b) {
    node ret;
    ret.h[0] = ((a.h[0] * (long long) BP[0][b.sz]
 \rightarrow ) + b.h[0] ) % MOD[0];
```

```
ret.h[1] = ( (a.h[1] * (long long) BP[1][b.sz]
    ) + b.h[1] ) % MOD[1];
    ret.sz = a.sz + b.sz:
    return ret;
  inline void build(int n,int st,int ed) {
    if(st==ed) ·
      tree[n].h[0] = tree[n].h[1] = ara[st];
      tree[n].sz = 1;
      return;
    int mid = (st+ed)>>1;
    build(n+n,st,mid);
    build(n+n+1,mid+1,ed);
    tree[n] = Merge(tree[n+n], tree[n+n+1]);
  inline void update(int n,int st,int ed,int id,int
    if(st>id or ed<id) return;</pre>
    if(st==ed and ed==id) {
      tree[n].h[0] = tree[n].h[1] = v;
      return;
    int mid = (st+ed)>>1;
    update(n+n,st,mid,id,v);
    update(n+n+1,mid+1,ed,id,v);
    tree[n] = Merge(tree[n+n], tree[n+n+1]);
  inline node query(int n,int st,int ed,int i,int
    if(st>=i and ed<=j) return tree[n];</pre>
    int mid = (st+ed)/2:
    if(mid<i) return query(n+n+1,mid+1,ed,i,j);</pre>
    else if(mid>=j) return query(n+n,st,mid,i,j);
    else return Merge(query(n+n,st,mid,i,j),query(n)
   +n+1,mid+1,ed,i,i));
26 HLD
```

```
int tt, tin[N], tout[N], sz[N], par[N][LG], hvc[N];
void dfs(int u, int p) {
 tin[u] = tt++, sz[u] = 1, par[u][0] = p;
  for (int j = 1; j < LG; ++j) {
    par[u][j] = par[par[u][j-1]][j-1];
  int mx = 0;
  for (int \&v: adj[u]) {
    if (v != p) {
      dfs(v, u);
      sz[u] += sz[v];
      if (sz[v] > mx) {
        mx = sz[v];
        hvc[u] = v;
  tout[u] = tt-1;
int ch cnt, idx cnt, chno[N], chd[N], idx[N];
|void hld(int u. int p) {
 if(chd[ch cnt] == -1) {
```

```
chd[ch cnt] = u;
  chno[u] = ch cnt, idx[u] = idx cnt++;
 if(hvc[u] != -1) {
    hld(hvc[u], u);
 for (int &v: adj[u]) {
    if (v != p and v != hvc[u]) {
      ch cnt++:
      hld(v, u);
void ?node update(int u, int x) {
  ?update(Idx[u], x);
void ?pupdate up(int u, int anc) {
 if (chno[u] == chno[anc]) {
    return ?rupdate(idx[anc], idx[u]);
  ?rupdate(idx[chd[chno[u]]], idx[u]);
  ?pupdate up(par[chd[chno[u]]][0], anc);
void ?pupdate(int u, int v) {
 int l = lca(u, v);
  ?pupdate up(u, l);
  ?pupdate_up(v, l);
ll ?node query(int u) {
 return ?query(idx[u]);
int ?pquery up(int u, int anc) {
 if (chno[\overline{u}] == chno[anc]) {
    return ?rquery(idx[anc], idx[u]);
 return f(?rquery(idx[chd[chno[u]]], idx[u]),
   ?pquery up(par[chd[chno[u]]][0], anc));
int ?rquery(int u, int v) {
 int l = lca(u, v):
 return f(?pquery_up(u, l), ?pquery_up(v, l));
adj[u].clear(); hvc[u] = -1;
tt = 0; dfs(0, 0);
chd[ch] = -1;
ch cnt = 0, idx cnt = 0; hld(0, 0);
```

27 HOPCROFT_KARP

```
// 1-based
const int N = 1e5+5, INF = 1e8 + 5;
vector <int> g[N];
int n, e, match[N], dist[N];
bool bfs() {
   queue <int> q;
   for (int i = 1; i <= n; ++i) {
      if (!match[i]) dist[i] = 0, q.emplace(i);
      else dist[i] = INF;
   }
   dist[0] = INF;
   while (!q.empty()) {
      int u = q.front(); q.pop();
   }
}</pre>
```

28 HUNGARIAN_ALGORITHM

```
16
```

```
if (!u) continue;
    for (int v : g[u]) {
      if (dist[match[v]] == INF) {
   dist[match[v]] = dist[u] + 1,
        q.emplace(match[v]);
  return dist[0] != INF;
bool dfs (int u) {
  if (!u) return 1;
  for (int v : q[u]) {
    if (dist[match[v]] == dist[u] + 1 and
→ dfs(match[v])) {
      match[u] = v, match[v] = u;
      return 1;
  dist[u] = INF;
  return 0;
int hopcroftKarp() {
  int ret = 0;
  while (bfs()) {
    for (int i = 1; i \le n; ++i) {
      ret += !match[i] and dfs(i);
  return ret;
```

```
= idx[k];
        if (R[j] < 0) goto aug;
    int q = idx[s++], i = R[q];
    for (int k = t; k < m; ++k) {
      i = idx[k];
      \dot{T} h = residue(i,j) - residue(i,q) + w;
      if (h < dist[j]) {
        dist[j] = h; prev[j] = i;
        if (h == w) {
          if (R[j] < 0) goto aug;
          idx[k] = idx[t]; idx[t++] = j;
aug:
  for(int k = 0; k < l; ++k)
   v[idx[k]] += dist[idx[k]] - w;
  int i;
  do {
   R[j] = i = prev[j];
    swap(j, L[i]);
  } while (i != f);
T ret = 0;
for (int i = 0; i < n; ++i) {
  ret += c[i][L[i]]; // (i, L[i]) is a solution
return {ret, L};
```

29 KMP

```
template<typename T>
pair<T, vector<int>> MinAssignment(const

→ vector<vector<T>> &c) {
 int n = c.size(), m = c[0].size();
                                            //
\rightarrow assert(n <= m);
 vector<T> v(m), dist(m);
                                            // v:

→ potential

 vector<int> L(n, -1), R(m, -1);
                                            //

→ matching pairs

 vector<int> idx(m), prev(m);
 iota(idx.begin(), idx.end(), 0);
 auto residue = [&](int i, int j) { return c[i][j]
→ - v[j]; };
 for (int f = 0; f < n; ++f) {
   for (int j = 0; j < m; ++j) {
     dist[j] = residue(f, j); prev[j] = f;
   T w; int j, l;
   for (int s = 0, t = 0;;) {
     if (s == t) {
        l = s; w = dist[idx[t++]];
        for (int k = t; k < m; ++k) {
           = idx[k]; T h = dist[j];
          if (h <= w) {
            if (h < w) { t = s; w = h; }
            idx[k] = idx[t]; idx[t++] = j;
        for (int k = s; k < t; ++k) {
```

```
vector<<mark>int</mark>> get pi(string& s){
int n = s.size();
 vector<int> pi(n);
 for (int k = 0, i = 1; i < n; ++i){
   if(s[i] == s[k]) pi[i] = ++k;
   else if(k == 0) pi[i] = 0;
   else k = pi[k-1], --i;
return pi;
pi.back(): n
// Borders = pi.back(), pi[pi.back() - 1], ...
// Prefix palindrome: s + "#" + rev(s)
// Number of occurrences of each prefix:
|vector<int> pref occur(vector<int> &pi) {
  int n = pi.size();
  vector<int> pref occur(n + 1);
  for (int i = 0; \bar{i} < n; ++i) {
    pref occur[pi[i]]++;
  for (int len = n; len > 0; --len) {
    pref occur[pi[len - 1]] += pref occur[len];
    pref occur[len]++;
  return pref occur;
// Find the length of the longest proper suffix of
   a suffix which also its prefix
// Reverse -> Find prefix function -> Reverse
```

```
// Find minimum length string such that given

→ strings occur as substring
```

30 MANACHER

```
// p[0][i] = half length of longest even palindrome
   around pos i-1, i and starts at i-p[0][i] and
   ends at i+p[0][i]-1
// p[1][i] = longest odd (half rounded down)
   palindrome around pos i and starts at i-p[1][i]
  and ends at i+p[1][i]
vector<vector<int>> manacher(string &s) {
 int n = s.size();
 vector<vector<int>>> p(2, vector<int> (n));
 for (int z = 0; z < 2; ++z) {
    for (int i=0, l=0, r=0; i<n; ++i) {
      int t = r-i+!z;
      if (i<r) {
        p[z][i] = min(t, p[z][l+t]);
      int L = i-p[z][i], R = i+p[z][i]-!z;
      while (L>=1 \text{ and } R+1< n \text{ and } s[L-1]==s[R+1]) {
        p[z][i]++, L--, R++;
      if (R>r) {
        l=L, r=\tilde{R};
 return p;
```

31 MATRIX_EXPO

```
using row = vector<int>;
using matrix = vector<row>;
matrix unit mat(int n) {
 matrix I(n, row(n));
 for (int i = 0; i < n; ++i){
   I[i][i] = 1;
 return I;
matrix mat mul(matrix a, matrix b) {
 int m = a.size(), n = a[0].size();
 int p = b.size(), q = b[0].size();
 // assert(n==p);
 matrix res(m, row(q));
 for (int i = 0; i < m; ++i){
    for (int j = 0; j < q; ++j){
      for (int k = 0; k < n; ++k){
        res[i][j] = (res[i][j] + a[i][k]*b[k][j]) %
   mod;
 return res:
matrix mat exp(matrix a, int p) {
 int m = a.size(), n = a[0].size();
 // assert(m==n);
 matrix res = unit mat(m);
```

```
17
```

```
while (p) {
   if (p&1) res = mat_mul(a, res);
   a = mat_mul(a, a);
   p >>= 1;
}
return res;
}
```

32 MCF

```
struct MCF {
 int n;
 vector<vector<array<ll, 5>>> adj;
                                       // v, pos of

→ u in v, cap, cost, flow
 vector<ll> dis, par, pos;
 MCF(int n): n(n), adj(n), dis(n), par(n), pos(n)
→ {}
 void add edge(int u, int v, int cap, int cost) {
    adj[u] push back({v, adj[v].size(), cap, cost,
   adj[v].push back({u, adj[u].size() - 1, 0,
   -cost, 0});
 ll spfa(int s, int t) {
    dis.assign(n, INF);
   vector<ll> mn_cap(n, INF), inq(n);
   queue<int> q;
   q.push(s), inq[s] = 1, dis[s] = \theta;
   while (!q.empty()) {
     int u = q.front(); q.pop();
      inq[u] = 0;
      for (int i = 0; i < adj[u].size(); ++i)</pre>
        auto [v, idx, cap, cost, flow] = adj[u][i];
        if (cap > flow and dis[v] > dis[u] + cost) { |
          dis[v] = dis[u] + cost;
          par[v] = u;
          pos[v] = i;
          mn cap[v] = min(mn cap[u], cap - flow);
          q.push(v);
         inq[v] = 1;
   return (mn cap[t] == INF? 0: mn cap[t]);
 array<ll, 2> get(int s, int t, ll max flow = INF)
   Il flow = 0, mc = 0;
   while (ll f = min(spfa(s, t), max flow - flow))
     flow += f;
      mc += f * dis[t];
     int u = t;
      while (u != s) {
        int p = par[u];
        adi[p][pos[u]][4] += f;
        adi[u][adi[p][pos[u]][1]][4] -= f;
    return {flow, mc};
```

33 MO ALOGO

```
vector<array<int, 4>> cu(m);
for (int i = 0; i < m; ++i) {
   auto &[b, l, r, idx] = cu[i];
   cin >> l >> r; l--;
   b = r / B;
   idx = i;
}
sort(cu.begin(), cu.end());
int s = 0, e = -1;
for (auto [b, l, r, i]: cu) {
   while (l < s) add(--s);
   while (e < r) add(++e);
   while (s < l) remove(s++);
   while (r < e) remove(e--);
   ans[i] = cur_ans;
}</pre>
```

34 NUMBER THEORY

```
## Floor
|ll floor (ll n, ll k) {
  if (n \ge 0) return n / k;
  return (n - (k - 1)) / k:
ll ceil (ll n, ll k) {
  if (n >= 0) return (n + k - 1) / k;
  return n / k;
## Highly Composite Number
1e6(240), 1e9(1344), 1e12(6720), 1e14(17280)
## Harmonic Lemma (ceill)
lli=1;
|while (i < n) {
  ll cval = (n + i - 1) / i;
  ll j = (n + cval - 2) / (cval - 1);
  // ceil(n/i)...ceil(n/(i - 1)) = cval
  cout << i << " " << j - 1 << ": " << cval <<
 i = j;
ĺl bezout(ll a, ll b, ll &x, ll &y){
  if(b == 0){
    x=1, y=0;
    return a;
  1l q = bezout(b, a%b, y, x);
  y = a/b*x;
  return g;
|ll mod inv(ll a, ll m){
 ll x, y;
```

```
ll g = bezout(a, m, x, y);
  if(q != 1) return -1; //no solution exists
  return (x%m+m)%m;
## Linear-sieve
int lpf[N], pm[N], pcnt = 0;
for (int i = 2; i < N; ++i) {
  if (!lpf[i])  lpf[i] = i, pm[pcnt++] = i;</pre>
  for (int j = 0; j < sz; ++j) {
    int p = pm[j];
    if (lpf[i] = N) break;
    lpf[i * p] = p;
## Miller-Rabin
bool isp(ll n){
  if(n==2 || n == 3) return 1;
  if(n<=1 | n%2==0) return 0;
  for (int k = 0; k < 10; ++k){
    ll a = 2+rand()%(n-2);
    ll s = n-1:
    while(!(s\&1)) s>>=1;
    if(powmod(a, s, n) == 1) continue;
    int iscomp = 1:
    while (s!=n-1)
      if(powmod(a, s, n)==n-1){
        iscomp = 0;
        break:
      s=s<<1:
    if(iscomp) return 0;
  return 1;
## Miller-Rabin Deterministic:
bool check composite(u64 n, u64 a, u64 d, int s) {
  u64 x = \overline{binpower}(a, d, n);
  if (x == 1 | | x == n - 1)
    return false;
  for (int r = 1; r < s; r++) {
    x = (u128)x * x % n;
    if (x == n - 1)
      return false;
  return true;
bool isp(u64 n) {
  if (n' < 2)
    return false;
  int r = 0;
  u64 d = n - 1:
  while ((d \& 1) == 0) {
    d >>= 1;
    r++;
  for (int a : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29,

→ 31, 37}) {
    if (n == a)
      return true;
    if (check composite(n, a, d, r))
```

```
return false;
  return true:
## Prime Factorize of large number(Pollard Rho):
ll f(ll x, ll c, ll n){
  return (mulmod(x,x,n)+c)%n;
if(n == \overline{1}) return 1;
  if(n\%2 == 0) return 2;
  ll x = rand()%(n-2)+2;
  ll y = x;
  ll c = rand()%(n-1)+1;
  ll q = 1;
  while (g == 1){
   x = f(x, c, n);
   y = f(y, c, n);
   y = f(y, c, n);
    g = \underline{gcd(abs(x-y), n)};
  return q;
vector<ll> prime factorize(ll n){
 if(n<=1) return vector<ll>();
  if(isp(n)) return vector<ll> ({n});
  ll d = pollard rho(n);
  vector<ll> v = factorize(d);
  vector<ll> w = factorize(n/d);
  v.insert(v.end(), w.begin(), w.end());
  sort(v.begin(), v.end());
  return v;
// auto pf = prime factorize(n);
## Number of divisors of n O(n^1/3):
int nod(ll n){
 sieve();
  int ret = 1;
  for (int i = 2; 1LL*i*i*i <= n; ++i){
   if(isp[i]){
      int e = 0;
      while(n\%i == 0){
        e++;
        n /= i;
      ret *= e+1;
  ll sq = sqrt(1.0L*n);
  if(isprime(n)) ret *= 2;
  else if(n == sq*sq and isprime(sq)) ret *= 3;
 else if(n!=1) ret *= 4;
  return ret;
## Smallest inverse phi
ll inv phi(ll phi, ll n, int pc) {
  if (\overline{phi} == 1) return n;
  if (pc == -1) return INF;
  ll ret = inv_phi(phi, n, pc - 1);
  if (phi % (p[pc] - 1) == 0) {
    phi /= (p[pc] - 1);
    n = n / (p[pc] - 1) * p[pc];
   while (phi % p[pc] == 0) {
      phi /= p[pc];
```

```
ret = min(ret, inv phi(phi, n, pc - 1));
  return ret;
ll phi; cin >> phi;
if (phi & 1) {
  cout << (phi == 1) << "\n";
  for (int i = 1; i * i <= phi; ++i) {
    if (phi % i == 0) {
       if (isp(i + 1)) \{
         p.push back(i + 1);
       if (i * i != phi and isp(phi / i + 1)){
         p.push back(phi / i + 1);
  sort(p.begin(), p.end());
ll ans = inv_phi(phi, phi, p.size() - 1);
  cout << (ans == INF? 0; ans) << "\n";
## GCD sum function from 1 to N:
|ll phi[N], q[N];
void pcgsm(){ //pre calculate gcd sum fucntion
  pcphi();
  for (int i = 1; i < N; ++i){
    for (int j = i; j < N; j+=i){
  g[j] += i*phi[j/i];</pre>
## All Pair gcd sum:
for (int i = 1; i < N; ++i) {
  for (int j = i; j < N; j += i) {
  gcd_sum[j] += 1ll * phi[i] * (j / i);</pre>
  gcd sum[i] -= i;
  pref_gcd_sum[i] = pref_gcd_sum[i - 1] +

    qcd sum[i];

## LCM sum function of n:
|ll lsm(ll n){
  ll ret=0;
  for(ll d=1; d*d<=n; d++){
    if(n%d==0){
       ret += d*phi(d);
       if(n/d!=d) ret += n/d*phi(n/d);
  return (ret+1)*n/2;
## LCM sum function from 1 to N
|ll phi[N], l[N];
void pclsm(){ //pre calculate lcm sum function
  pcphi();
  for (int i = 1; i < N; ++i){
    for (int j = i; j < N; j+=i){
    l[j] += i*phi[i];</pre>
  for (int i = 1; i < N; ++i){
    l[i] = (l[i]+1)*i/2;
```

```
## All pair lcm sum:
for (int i = 1; i < N; ++i) {
 for (int j = i; j < N; j += i) {
   lcm_sum[j] += i * phi[i];</pre>
  lcm sum[i]++;
  lcm_sum[i] /= 2;
  lcm sum[i] *= i;
  lcm sum[i] -= i;
  pref lcm sum[i] = lcm sum[i];
  pref lcm sum[i] += pref lcm sum[i - 1];
## Number of co-prime pairs of an array:
vector<ll> cnt(A);
for (int xi: x) {
  for (int d = 1; d * d <= xi; ++d) {
    if (xi % d == 0) {
      cnt[d]++;
      if (xi / d != d) {
        cnt[xi / d]++;
ll ans = 0;
for (int i = 1; i < A; ++i) {
 if (!sq_free[i]) continue;
ll ways = cnt[i] * (cnt[i] - 1) / 2;
  if (pf[i].size() \& 1 ^{\circ} 1) ans += ways;
  else ans -= ways;
## All pair gcd sum of an array:
vector<ll> cnt(A);
for (auto ai: a) {
 for (int d = 1; d * d <= ai; ++d) {
    if (ai % d == 0) {
      cnt[d]++;
      if (ai / d != d) {
        cnt[ai / d]++;
ll sum = 0;
vector<ll> left(A);
iota(left.begin(), left.end(), 0);
for (int i = 1; i < A; ++i) {
    ll add = left[i] * cnt[i] * (cnt[i] - 1) / 2;</pre>
  sum += add:
 for (int j = 2 * i; j < A; j += i) {
    left[j] -= left[i];
ll crt(ll r1, ll m1, ll r2, ll m2){
  if(m1<m2) swap(r1, r2), swap(m1, m2);
  ll p, q, g = bezout(m1, m2, p, q);
  if((r2-r1)%q !=0) return -1; //no solution
  ll x = (r2-r1)m2*pm2*m1/g + r1;
  return x<0? x+m1*m2/g: x;
ll crt(vector<ll>& r, vector<ll>& m){
 ll x = r[0], M=m[0];
```

```
for (int i = 1; i < r.size(); ++i){</pre>
   x = crt(x, M, r[i], m[i]);
   ll q = gcd(M, m[i]);
   M = (M/\overline{q})^* (m[i]/q);
 return x;
## Discrete Logarithm
ll discrete log(ll a, ll b, ll m) {
 a %= m, b %= m;
 if(a == 0){
   return (b == 0? 1: -1);
 ll k = 1, add = 0, g;
 while ((q = qcd(a, m)) > 1)  {
   if (b == k) return add;
   if (b % q) return -1;
   b /= g, m /= g, k = (k * a / g) % m, ++add;
 int n = sqrt(m) + 1;
 unordered map<int. int> vals:
 for (ll q = 0, cur = b; q <= n; ++q) {
   vals[cur] = q;
   cur = (cur * a) % m;
 il an = 1;
 for (int i = 0; i < n; ++i) {
   an = (an * a) % m;
 for (ll p = 1, cur = k; p \le n; ++p) {
   cur = (cur * an) % m;
   if (vals.count(cur)) {
     return n * p - vals[cur] + add;
 return -1;
```

```
pt[x].link;
 if(!pt[at].next.count(c)){
    pt[at].next[c] = ++sz;
    pt[sz].len = pt[at].len + 2
    // cnt[pt[at].len+2]++; //for finding number
   of distinct palindrome of lenght k
    pt[sz].link = (pt[sz].len == 1)? 1 :
   pt[x].next[c];
    // pt[sz].no of suf pal = 1 +
   pt[pt[sz].link].no of suf pal; //for finding
   number of palindrome which last position is si
 // cnt[pt[at].len + 2]++; //for finding number

→ of palindrome of lenght k

 at = pt[at].next[c];
int num of pal(int ai){  //distinct palindrome,

→ arrav index

 int ret = pt[at].ans;
 for(auto x : pt[ai].next)
    ret += num of pal(x.second);
  return ret;
int main(){
 scanf("%s", s);
 pt init();
 for (int i = 0; s[i]; ++i){
    pt_extend(i);
 int ans = num of pal(0) + num of pal(1) - 2;
 printf("%d\n", ans);
  return 0;
```

while (s[si - pt[x].len - 1] != s[si]) x =

```
root[0] = new node(); // initialization
root[k] = add(root[k], i, x, 0, sz - 1);
root[ver++] = root[k];
cout << rsum(root[k], l, r, 0, sz - 1) << "\n";
## count numbers > k in a range
root[0] = new node();
for (int i = 0; i < n; ++i) {
root[i + 1] = add(root[i], a[i], 1);
while (q--) {
  int l, r, k; cin >> l >> r >> k; l--, r--;
  int ans = rsum(root[r + 1], k, E - 1) -
\rightarrow rsum(root[l], k, E - 1);
  cout << ans << "\n":
## kth number in a range: O(logn)
int kth(node *ul, node *ur, int k, int s = 0, int e
\rightarrow = E - 1) {
  if (s == e) return s;
  int m = (s + e) / 2;
  int cnt left = ur->left->sum - ul->left->sum;
  if (cnt<sup>-</sup>left >= k) return kth(ul->left,
\rightarrow ur->left, k, s, m);
  else return kth(ul->right, ur->right, k -
   cnt left, m + 1, e);
root[0] = new node();
for (int i = 0; i < n; ++i) {
 root[i + 1] = add(root[i], a[i + ], 1);
while (q--) {
  int l, r, k; cin >> l >> r >> k; l--, r--;
  int x = kth(root[l], root[r + 1], k);
```

35 PALINDROMIC TREE

```
const int N = 1e5+10;
struct vertex
  int len, link, no of suf pal;
  map<char, int> next;
}pt[N];
int sz, at, cnt[N];
char s[N];
void pt init(){
  for (int i = 0; i < N; ++i){
   pt[i].next.clear();
 memset(cnt, 0, sizeof(cnt));
 pt[0].len = -1, pt[0].link = 0,

→ pt[0].no of suf pal = 0;
 pt[1].len = 0, pt[1].link = 0,
\rightarrow pt[1].no of suf pal = 0;
 sz = at = I;
void pt_extend(int si){  //string index
 while (s[si - pt[at].len - 1] != s[si]) at =

    pt[at].link:
 int x = pt[at].link, c = s[si]-'a';
```

36 PERSISTENT_SEGMENT_TREE

```
## Point Addition & Range Sum:
struct node {
  ll sum;
  node *l, *r;
  node(ll s = 0, node *l = NULL, node *r = NULL):
 \rightarrow sum(s), l(l), r(r) {}
|node* add(node *u, int i, int x, int s, int e) {
  if (s == e) return new node(u->sum + x);
  if (|u->l) u->l = new node(), u->r = new node();
  node *nu = new node(u->sum, u->l, u->r);
  int m = (s + e) / 2;
  if (i \le m) nu > l = add(nu - > l, i, x, s, m);
  else nu \rightarrow r = add(nu \rightarrow r, i, x, m + 1, e);
  nu->sum = nu->l->sum + nu->r->sum:
  return nu:
ll rsum(node *u, int l, int r, int s, int e) {
  if (!u) return 0;
  if (s > r \text{ or } e < l) return 0;
  if (l <= s and e <= r) return u->sum;
  int m = (s + e) / 2;
return rsum(u->l, l, r, s, m) + rsum(u->r, l, r,
 \rightarrow m + 1, e);
vector<node*> root(VER);
```

37 PERSISTENT_TRIE

```
struct node
 node *nxt[2];
 node() { fill(nxt, nxt + 2, nullptr); }
node* add(node *prev, int x) {
 node *new root = new node();
 node * cur = new root;;
 for (int idx = I\overline{D}X - 1; idx >= 0; --idx) {
   int f = (x \gg idx) \& 1;
   if (prev and prev->nxt[!f]) cur->nxt[!f] =
→ prev->nxt[!f];
    cur->nxt[f] = new node():
    cur = cur->nxt[f];
    if (prev) prev = prev->nxt[f];
 return new root;
int get max(node *root, int x) {
 if (!root) return 0;
 node *u = root;
 int ret = 0;
 for (int idx = IDX - 1; idx \geq 0; --idx) {
   int f = (x >> idx) \& 1;
   if (u->nxt[!f]) ret += (1 << idx), u =
   u->nxt[!f];
    else u = u - nxt[f];
```

```
}
return ret;
}
```

38 POLYNOMIAL_INTERPOLATION

```
// P(x) = a0 + a1x + a2x^2 + ... + anx^n
// y[i] = P(i)
ll eval (vector<ll> y, ll k) {
          int n = y.size() - 1;
          if (k <= n) {
                    return y[k];
          vector<ll> L(n + 1, 1);
          for (int x = 1; x <= n; ++x) {
                  L[0] = L[0] * (k - x) % mod;
                   L[0] = L[0] * inv(-x) % mod;
          for (int x = 1; x <= n; ++x) {
                  L[x] = L[x - 1] * inv(k - x) % mod * (k - (x - x)) % mod * (k - x) % mod * (x - x) % mod * (
               1)) % mod;
                   L[x] = L[x] * ((x - 1) - n + mod) % mod *

→ inv(x) % mod;

         il yk = 0;
         for (int x = 0; x <= n; ++x) {
                  yk = add(yk, L[x] * y[x] % mod);
          return yk;
```

39 SCC

```
void dfs1(int u, vector<int> *adj, vector<int>
vis[u] = 1;
 for (int \&v: adj[u]) {
   if (!vis[v]) {
     dfs1(v, adj, vis, order);
 order.emplace back(u);
void dfs2(int u, vector<int> *rev adj, vector<int>

→ &vis, vector<int> &scc) {
 scc.emplace back(u);
 vis[u] = 1;
 for (int &v: rev adj[u]) {
   if (!vis[v]) {
     dfs2(v, rev adj, vis, scc);
vector<vector<int>> get sccs(int n, vector<int>
→ *adj) {
 vector<int> vis(n), order;
 for (int u = 0; u < n; ++u) {
   if (!vis[u]) {
     dfs1(u, adj, vis, order);
 vector<int> rev adj[n];
 for (int u = 0; u < n; ++u) {
   for (int v: adi[u]) {
```

```
rev_adj[v].emplace_back(u);
}

vector<vector<int>> sccs;
reverse(order.begin(), order.end());
vis.assign(n, 0);
for (int u: order) {
   if (!vis[u]) {
      sccs.emplace_back(0);
      dfs2(u, rev_adj, vis, sccs.back());
   }
}
return sccs;
}
vector<vector<int>> sccs = get_sccs(n, adj);
int tot scc = sccs.size();
vector<int> scc_no(n);
for (int i = 0; i < tot_scc; ++i) {
   for (int u: sccs[i]) {
      scc_no[u] = i;
   }
}</pre>
```

40 SEGMENT_TREE

```
## Range Addition and Range Assign and Range sum
int n;
ll t[3 * N], p[3 * N], p2[3 * N]; //t for sum, p

→ for assign & p2. for add

void pull(int v) {
 t[v] = t[2 * v] + t[2 * v + 1];
void push(int v, int st, int ed) {
 int lc = 2 * v, rc = 2 * v + 1, md = (st + ed) /
 if (p[v] != -1) {
    t[lc] = p[v] * (md - st + 1);
    t[rc] = p[v] * (ed - md);
    p[lc] = p[rc] = p[v];
    p2[lc] = p2[rc] = 0;
    p[v] = -1;
 if (p2[v]) {
    t[[c] + p2[v] * (md - st + 1);
    t[rc] += p2[v] * (ed - md);
    p2[lc] += p2[v];
    p2[rc] += p2[v];
    p2[v] = 0;
void assign(int l, int r, int x, int v = 1, int st
\rightarrow = 0, int ed = n - 1) {
 if (l > ed or r < st) return;</pre>
  if (l <= st and ed <= r) {
    t[v] = 1LL * (ed - st + 1) * x;
    y(x) = x
    p2[v] = 0;
    return;
  int lc = 2 * v, rc = 2 * v + 1, md = (st + ed) /

→ 2;

 push(v, st, ed);
  assign(l, r, x, lc, st, md);
  assign(l, r, x, rc, md + 1, ed);
  pull(v);
```

```
void add(int l, int r, int x, int v = 1, int st =
\rightarrow 0, int ed = n - 1) {
  if (l > ed or r < st) return;</pre>
 if (l <= st and ed <= r) {
    t[v] += 1LL * (ed - st + 1) * x;
    p2[v] += x;
    return ;
  push(v, st, ed);
  int lc = 2 * v, rc = 2 * v + 1, md = (st + ed) /
  add(l, r, x, lc, st, md);
  add(l, r, x, rc, md + 1, ed);
  pull(v);
il rsum(int l, int r, int v = 1, int st = 0, int ed
\rightarrow = n - 1) +
 if (l > ed or r < st) return 0;</pre>
  if (l <= st and ed <= r) return t[v];</pre>
  push(v, st, ed);
  int lc = 2 * v, rc = 2 * v + 1, md = (st + ed) /
 ll lret = rsum(l, r, lc, st, md);
  ll rret = rsum(l, r, rc, md + 1, ed);
  return lret + rret;
## Make All Elements <= k and Make all elements >=
→ k on range & Point Query:
const int I = 1e9 + 9;
int t[3 * N], pa[3 * N], pr[3 * N], ar[3 * N]; //pa
    for propagate adding, pr for propagate remove,
ar for check last on is adding(1) or remove(0)
void fg(int x, int u) { //function for

→ make greater

 t[u] = max(t[u], x);
  pa[u] = max(pa[u], x);
  pr[u] = max(pr[u], x);
  ar[u] = 1;
void fl(int x, int u) { //function for make less
 t[u] = min(t[u], x);
  pr[u] = min(pr[u], x);
  pa[u] = min(pa[u], x);
  ar[u] = 0;
void push(int u) {
  int v = 2 * u, w = 2 * u + 1;
  if (ar[u] == 0) {
    if (pa[u] != -1) {
      fg(pa[u], v); fg(pa[u], w);
    if (pr[u] != I) -
      fl(pr[u], v); fl(pr[u], w);
  } else {
    if (pr[u] != I)
      fl'(pr[u], v); fl(pr[u], w);
    if (pa[u] != -1) {
      fg(pa[u], v); fg(pa[u], w);
```

```
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```

```
pa[u] = -1; pr[u] = I:
void make greater(int l, int r, int x, int u = 1,
\rightarrow int s = 0, int e = N - 1) {
 if (l > e or r < s) return;</pre>
  if (l <= s and e <= r) {
    fg(x, u);
    return ;
  push(u);
 int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
make_greater(l, r, x, v, s, m);
 make\_greater(l, r, x, w, m + 1, e);
void make less(int l, int r, int x, int u = 1, int
\rightarrow s = 0, int e = N - 1) {
 if (l > e or r < s) return;</pre>
 if (l <= s and e <= r) {
    fl(x, u);
    return;
  push(u);
  int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
 make_less(l, r, x, v, s, m);
 make^{-less(l, r, x, w, m + 1, e)};
int at(int i, int u = 1, int s = 0, int e = N - 1) {
 if (s == e) return t[u]:
  push(u);
  int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
 if (i <= m) return at(i, v, s, m);</pre>
  else return at(i, w, m + 1, e);
```

41 SHORTEST_PATH

```
## Diikstra
priority queue<array<ll, 2>> pq;
vector<ll>dis(n, INF), vis(n);
while (!pq.empty()) {
  auto [d, u] = pq.top(); pq.pop();
  if (vis[u]) continue;
  vis[u] = 1;
  for (auto [v, c]: next[u]) {
   if (dis[v] > d + c) {
      dis[v] = d + c;
      pq.push({dis[v], v});
## Bellman-ford
vector<int> bellman ford(int s){
  vector<int> dis(n. I):
  dis[s]=0;
  while(1){}
    int any=0;
    for (auto& e: ed){
      if(dis[e.u]<I){</pre>
        if(dis[e.u]+e.cost < dis[e.v]){</pre>
          dis[e.v] = dis[e.u]+e.cost;
          any=1;
    if(!any) break;
```

```
return dis;
## Floy-Warshall
|for (int k = 0; k < n; ++k) {
  for (int i = 0; i < n; ++i) {
    for (int j = 0; j < n; ++j) {
      dis[i][j] = min(dis[i][j], dis[i][k] +

    dis[k][j]);

42 SOS DP
## Count over subset
for (int i = 0; i < n; ++i) f[a[i]] = ?;
for (int i = 0; i < i; ++i) {
  for (int mask = 0; mask < (1 << n); ++mask) {
    if (mask\&(1<<i)) {
      f[mask] += f[mask^(1<<i)];
## Count over superset
for (int i = 0; i < n; ++i) f[a[i]] = ?;</pre>
for (int i = 0; i < n; ++i) {
  for (int mask = (1 << n) - 1; mask >= 0; --mask)
    {
if (!(mask&(1<<i))) {
      f[mask] += f[mask^{(1<<i))};
## How many pairs in ara[] such that (ara[i] &
\rightarrow ara[j]) = 0
/// N --> Max number of bits of any array element
const int N = 20;
int inv = (1<<N) - 1;
int F[(1<<N) + 10];
int ara[MAX];
/// ara is 0 based
long long howManyZeroPairs(int n,int ara[]) {
    CLR(F):
    for(int i=0;i<n;i++) F[ara[i]]++;</pre>
    for(int i = 0; i < N; ++i)
        for(int mask = 0; mask < (1 << N); ++mask){
            if(mask & (1<<i))
                F[mask] += F[mask^(1<<i)];
    long long ans = 0;
    for(int i=0;i<n;i++) ans += F[ara[i] ^ inv];</pre>
    return ans;
/// To get
    for(int mask = 0; mask < (1 << N); ++mask)
        for(int i = 0; i < (1 << N); ++i)
            if( (mask & i) == mask ) { /// i is a
 F[mask] += A[i];
/// The code is the following
```

```
for(int i = 0; i < (1 << N); ++i) F[i] = A[i];
    for(int i = 0; i < N; ++i)
        for(int mask = (1 << N) - 1; mask >= 0; --mask){
             \mathbf{if} (!(mask & (1<<i)))
                 F[mask] += F[mask \mid (1 << i)];
## Number of subsequences of ara[0:n-1] such that
## sub[0] \& sub[2] \& ... \& sub[k-1] = 0
const int N = 20;
int inv = (1 << N) - 1;
int F[(1 << \dot{N}) + \dot{1}0];
int ara[MAX];
int p2[MAX]; /// p2[i] = 2^i
///0 based array
int howManyZeroSubSequences(int n,int ara[]) {
    CLR(F);
    for(int i=0;i<n;i++) F[ara[i]]++;</pre>
    for(int i = 0; i < N; ++i)
        for(int mask = (1 << N) - 1; mask >= 0; --mask){
             if (!(mask & (1<<i)))
                 F[mask] += F[mask \mid (1 << i)];
    int ans = 0;
    for(int mask=0; mask<(1<< N); mask++) {
        if( builtin popcount(mask) \& 1) ans =
    sub(ans, p2[F[mask]]);
        else ans = add(ans, p2[F[mask]]);
    return ans;
## Number of subsequences of ara[0:n-1] such that
## sub[0] | sub[2] | ... | sub[k-1] = 0
int F[(1<<20) + 10], ara[MAX];
int p2[MAX]; /// p2[i] = 2^i
/// ara is 0 based
int howManySubsequences(int n, int ara[], int m,
→ int Q) {
    CLR(F);
    for(int i=0;i<n;i++) F[ara[i]]++;</pre>
    if(0 == 0) return sub(p2[F[0]], 1);
    for(int i = 0; i < m; ++i)
        for(int mask = 0; mask < (1<<m); ++mask){</pre>
             if (mask \& (1<<i))
                 F[mask] += F[mask ^ (1<<i)];
    int ans = 0;
    for(int mask=0; mask<(1<< m); mask++) {
        if(mask & Q != mask) continue;
        if( builtin popcount(mask ^{\circ} Q) \& 1) ans =

    sub(ans, p2[F[mask]]);

        else ans = add(ans, p2[F[mask]]);
    return ans:
```

43 SPARSE_TABLE

```
int n, a[N], lg[N], st[N][K];
for (int i = 2; i < N; ++i) {
    lg[i] = lg[i / 2] + 1;
}
void build() {
    for (int k = 0; k < K; ++k) {</pre>
```

```
22
```

44 SPARSE_TABLE_2D

```
int st[N][N][LG][LG];
int a[N][N], lg2[N];
int yo(int x1, int y1, int x2, int y2) {
  x2++;
  y2++;
  int a = \lg 2[x2 - x1], b = \lg 2[y2 - y1];
  return max(
         \max(st[x1][y1][a][b], st[x2 - (1 <<

→ a)][y1][a][b])
         \max(st[x1][y2 - (1 << b)][a][b], st[x2 -
   (1 << a)][y2 - (1 << b)][a][b])
       );
void build(int n, int m) { // 0 indexed
 for (int i = 2; i < N; i++) lg2[i] = lg2[i >> 1]

→ + 1;

  for (int i = 0; i < n; i++) {
    for (int j = 0; j < m; j++) {
      st[i][j][0][0] = a[i][j];
  for (int a = 0; a < LG; a++) {
    for (int b = 0; b < LG; b++) {
      if (a + b == 0) continue:
      for (int i = 0; i + (1 << a) <= n; i++) {
        for (int j = 0; j + (1 << b) <= m; j++) {
          if (!a) {
            st[i][i][a][b] = max(st[i][i][a][b -
\rightarrow 1], st[i][j + (1 << (b - 1))][a][b - 1]);
          } else {
            st[i][j][a][b] = max(st[i][j][a -
\rightarrow 1][b], st[i + (1 << (a - 1))][j][a - 1][b]);
```

45 SQRT_DECOMPOSITION

```
const int SZ2 = 2e5, SZ = sqrt(SZ2+.0)+1, N = SZ*SZ;
int n, a[N], b[SZ];
void build() {
   for (int i = 0; i < SZ; ++i){
      b[i] = INT_MAX;
   }
   for (int i = 0; i < n; ++i){
      b[i/SZ] = min(b[i/SZ], a[i]);
   }</pre>
```

```
int rmq(int l, int r) {
    int lb = l/SZ, rb = r/SZ;
    int ret = INT_MAX;
    if(lb==rb) {
        for (int i = l; i <= r; ++i) {
            ret = min(ret, a[i]);
        }
    } else {
        for (int i = l; i < (lb+1)*SZ; ++i) {
            ret = min(ret, a[i]);
        }
        for (int i = lb+1; i < rb; ++i) {
            ret = min(ret, b[i]);
        }
        for (int i = rb*SZ; i <= r; ++i) {
            ret = min(ret, a[i]);
        }
    }
    return ret;
}</pre>
```

46 STRESS_TESTING

```
set -e
g++ -02 -static -std=gnu++17 gen.cpp -o gen
g++ -02 -static -std=gnu++17 main.cpp -o main
g++ -02 -static -std=gnu++17 brute.cpp -o brute
for((i = 1; ; ++i)); do
    echo $i
        ./gen $i > in
        # ./main < in > out
        # ./brute < in > out2
        # diff -w out out2 || break
        diff -w <(./main < in) <(./brute < in) || break
done</pre>
```

47 SUFFIX_ARRAY

```
|array<vector<<mark>int</mark>>, 2> get sa(string& s, int
    lim=128) { // for integer, just change string
    to vector<int> and minimum value of vector must
   be >= 1
  int n = s.size() + 1, k = 0, a, b;
  vector<int> x(begin(s), end(s)+1), y(n), sa(n),

    lcp(n), ws(max(n, lim)), rank(n);

  x.back() = 0;
  iota(begin(sa), end(sa), 0);
  for (int j = 0, p = 0; p < n; j = max(1, j * 2),
\rightarrow lim = p) {
    p = j, iota(begin(y), end(y), n - j
    for (int i = 0; i < n; ++i) if (sa[i] >= j)
    y[p++] = sa[i] - j;
fill(begin(ws), end(ws), 0);
    for (int i = 0; i < n; ++i) ws[x[i]]++;
    for (int i = 1; i < lim; ++i) ws[i] += ws[i -
   1];
    for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
    swap(x, y), p = 1, x[sa[0]] = 0;
    for (int i = 1; i < n; ++i) a = sa[i - 1], b =
   sa[i], x[b] =
      (y[a] == y[b] \&\& y[a + j] == y[b + j]) ? p -
    1 : p++;
```

```
for (int i = 1; i < n; ++i) rank[sa[i]] = i;
for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)
    for (k && k--, j = sa[rank[i] - 1]; s[i + k] ==
        s[j + k]; k++);
    sa.erase(sa.begin()), lcp.erase(lcp.begin());
    return {sa, lcp};
}</pre>
```

48 SUFFIX_AUTOMATON

```
struct state {
 int len, link, cnt tmp = 0, cnt = 0;
 map<char, int> next;
const int MAXLEN = 100000:
state st[2*MAXLEN];
int dp[2*MAXLEN];
int sz, last;
void sa_init() {
 st[0].len = 0;
 st[0].link = -1;
 SZ++;
 last = 0:
 memset(dp, -1, sizeof(dp));
void sa extend(char c) {
 int cur = sz++;
 st[cur].len = st[last].len + 1;
  int p = last;
 while (p != -1 \&\& !st[p].next.count(c)) {
    st[p].next[c] = cur;
    p = st[p].link;
 if (p == -1) {
    st[cur].link = 0;
  } else {
    int q = st[p].next[c];
    if (st[p].len + 1 == st[q].len) {
      st[cur].link = q;
    } else {
      int clone = sz++;
      st[clone].len = st[p].len + 1;
      st[clone].next = st[q].next;
      st[clone].link = st[q].link;
      while (p != -1 \&\& st[p].next[c] == q) {
        st[p].next[c] = clone;
        p = st[p].link;
      st[q].link = st[cur].link = clone;
  last = cur:
## Count Occurence
int occurence(string p){
 int at = 0:
 for (int i = 0; p[i]; ++i){
    if(st[at].next.count(p[i]) == 0){
      return 0;
```

```
else{
        at = st[at].next[p[i]];
 return st[at].cnt tmp;
vector<int> used;
void dfs(int x){
  used[x]=1;
  for(auto it:st[x].next) {
    if(!used[it.second]) dfs(it.second);
    if(it.first=='#') st[x].cnt tmp++;
    else st[x].cnt tmp+=st[it.second].cnt tmp;
    st[x].cnt+=st[īt.second].cnt;
  st[x].cnt+=st[x].cnt tmp;
## number of distinct substring O(n)
long long disub(int at){
  if(dp[at] != -1)
    return 0:
  dp[at] = 0;
  long long ret = 0;
  for(auto x : st[at].next)
    ret += disub(x.second);
  if(at != 0)
    ret += (st[at].len - st[st[at].link].len);
  return ret;
## longest common substring: O(|T|)
int lcs (string S, string T) {
  sa init();
  for (int i = 0; i < S.size(); i++)
    sa extend(S[i]);
  int v = 0, l = 0, best = 0, bestpos = 0;
  for (int i = 0; i < T.size(); i++) {
    while (v \&\& !st[v].next.count(T[i]))  {
      v = st[v].link;
      l = st[v].len ;
    if (st[v].next.count(T[i])) {
      v = st [v].next[T[i]];
      l++;
    if (l > best) {
      best = l:
      bestpos = i;
  return best;
## Distinct Substring
long long disub(int at){
 long long ret = 1;
  for(auto x : st[at].next){
      ret += disub(x.second);
  return ret-1;
int main(){
 int T, caseno = 0;
```

```
scanf("%d", &T);
  while(T--){
    int a:
              cin >> a:
    sa init();
    string s; cin >> s;
    cout << s << endl:
    s += "#";
    for (int i = 0; s[i]; ++i){
      sa extend(s[i]);
    used.assign(sz,0);
    dfs(0);
    printf("Case %d:\n", ++caseno);
    while (q--){
      string p;
                   cin >> p;
      int ans = occurence(p);
      cout << ans << endl;
  return 0;
1. Finding Pattern
2. Frequency of each stat
3. First Occurrence
4. Last Occurrence
5. All Occurrenc
Longest Repeated substring:
7. Count number of different substring
Total length of different substring
9. k-th smallest distinct substring
10. K-th smallest substring
11. Smallest Cyclic Shift
12. Find borders
13. Find Periods:
14. Longest Common Substring
```

49 TREAP

```
## Typical TEAP
struct node {
  ll val, prior, sz, sum;
  node *ĺ, *r;
  node(int val, int prior, int sz) : val(val),
    prior(prior), sz(sz), sum(0), l(nullptr),

¬ r(nullptr){}
|using pnode = node*;
pnode root;
pnode new node(ll val){
 return new node(val, rand(), 1);
| int get sz(pnode u){
  return u? u->sz: 0;
void update(pnode u){
  if (!u) return ;
  u->sz = get sz(u->l) + 1 + get sz(u->r);
  u->sum = u->val + (u->l? u->l->sum: 0) + (u->r?
 \rightarrow u->r->sum: 0):
void split(pnode u, pnode &l, pnode &r, ll val){
  if(!u) l = r = NULL;
  else if(val > u->val) split(u->r, u->r, r, val),
  else split(u \rightarrow l, l, u \rightarrow l, val), r = u;
```

```
update(u);
void merge(pnode &u, pnode l, pnode r){
  if(!l or !r) u = l? l: r;
  if(l->prior > r->prior) merge(l->r, l->r, r), u
  else merge(r->l, l, r->l), u = r;
  update(u);
void insert(pnode &u, pnode it){
  if(!u) u = it;
  else if(it->prior > u->prior) split(u, it->l,
\rightarrow it->r, it->val), u = it;
  else insert(it->val <= u->val ? u->l: u->r, it);
  update(u);
void erase(pnode &u, ll val){
  if(!u) return ;
  if(val == u->val) merge(u, u->l, u->r);
  else erase(val < u->val ? u->l: u->r, val);
  update(u);
bool present(pnode u, int x){
  if(!u) return false;
  if(u->val == x) return true;
  if(u->val < x) return present(u->r, x);
  return present(u->l, x);
ll kth(pnode u, int k){
  if(get sz(u) < k) return INT MIN;</pre>
  if(get_sz(u->l) == k-1) return u->val;
  if(get sz(u->l) < k-1) return kth(u->r, k -
\rightarrow get \overline{sz}(u->1) - 1);
  return kth(u->l, k);
int cnt less(pnode u, ll x){
  if(!u) return 0;
  if(x <= u->val) return cnt less(u->l, x);
  return get sz(u->l) + 1 + cnt less(u->r, x);
ll sum less(pnode u, ll x) {
  if (Tu) return 0;
  if (x <= u->val) return sum less(u->l, x);
  return u \rightarrow val + (u \rightarrow l? u \rightarrow l \rightarrow sum: 0) +
\rightarrow sum less(u->r, x);
## Implicit TREAP
struct node {
  ll val, sum;
  int prior, sz, rev;
  node *l, *r;
  node(){}
  node(ll val): val(val), sum(val), prior(rand()),
\rightarrow sz(1), rev(0), l(nullptr), r(nullptr) {}
|using pnode = node*;
pnode root;
int get sz(pnode t) {
  return t? t->sz: 0;
11 get sum(pnode t)
  return t? t->sum: 0;
void update(pnode &t) {
```

```
24
```

```
if (!t) return ;
  t->sz = get sz(t->l) + 1 + get sz(t->r);
  t \rightarrow sum = ge\overline{t} sum(t \rightarrow l) + t \rightarrow va\overline{l} + get sum(t \rightarrow r);
void push(pnode t) {
  if (t and t->rev) {
    swap(t->l, t->r);
    t - rev = 0;
    if (t->l) {
      t->l->rev ^= 1;
    if (t->r) {
 t->r->rev ^= 1;
void merge(pnode &t, pnode l, pnode r){
  push(l);
  push(r);
  if(!l or !r) t=l?l:r;
  else if(l->prior > r->prior) merge(l->r, l->r,
   r), t=l;
  else merge(r->l,l,r->l) , t=r;
  update(t);
void split(pnode t, pnode &l, pnode &r, int pos,

    int add=0) {

  push(t);
  if(!t) return void(r=l=NULL);
  int cur pos = get sz(t->l)+add;
  if(pos > cur pos) split(t->r, t->r, r, pos,
   cur pos+1), l = t;
  else = split(t->l, l, t->l, pos, add), r=t;
  update(t);
void insert(pnode &t, pnode it, int i) {
  pnode t1, t2;
  split(t, t1, t2, i);
  merge(t1, t1, it);
  merge(t, t1, t2);
void reverse(pnode &t, int l, int r) {
  pnode lt, mt, rt;
  split(t, t, rt, r + 1);
split(t, lt, mt, l);
  mt -> rev = 1;
  merge(mt, mt, rt);
  merge(t, lt, mt);
ll rsum(pnode& t, int l, int r) {
  pnode lt, mt, rt;
  split(t, t, rt, r + 1);
split(t, lt, mt, l);
  ll ret = mt->sum;
  merge(mt, mt, rt);
  merge(t, lt, mt);
  return ret;
int n, q; cin >> n >> q;
vector<ll> a(n);
for (auto &ai: a) {
  cin >> ai;
for (int i = 0; i < n; ++i) {
  insert(root, new node(a[i]), i);
```

```
|while (q--) {
  int tp, l, r; cin >> tp >> l >> r; l--, r--;
  if (tp == 1) {
    reverse(root, l, r);
  else {
    cout << rsum(root, l, r) << "\n";</pre>
```

$50 ext{ XOR_BASIS}$

```
ll rnk, basis[D];
void insert_vector(ll mask){
 if(!basis[i]){
    basis[i] = mask, rnk++;
    return;
   } else mask ^= basis[i];
```

51 Z_ALGORITHM

```
vector<int> get z(string s){
  int n=s.size(), l=1, r=0;
  vector\langle int \rangle z(n); z[0]=n;
  for (int i = 1; i < n; ++i){
  if(i<=r)  z[i]=min(z[i-l], r-i+1);</pre>
    while(s[i+z[i]]==s[z[i]]) z[i]++;
     if(i+z[i]-1>r) l=i, r=i+z[i]-1;
  return z;
```

52 note

Binomial Coefficent

```
- Factoring in: \binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}
```

Sum over $k: \sum_{k=0}^{n} {n \choose k} = 2^n$

Alternating sum: $\sum_{k=0}^{n} (-1)^k \binom{n}{k} = 0$

Even and odd sum: $\sum_{k=0}^{n} {n \choose 2k} = \sum_{k=0}^{n} {n \choose 2k+1} 2^{n-1}$

The Hockey Stick Identity

(Left to right) Sum over n and k: $\sum_{k=0}^{m} {n+k \choose k} = {n+m-1 \choose m}$

(Right to left) Sum over $n: \sum_{m=0}^{n} {m \choose k} = {n+1 \choose k+1}$

Sum of the squares: $\sum_{k=0}^{n} {\binom{n}{k}}^2 = {\binom{2n}{n}}$

Weighted sum: $\sum_{k=1}^{n} k \binom{n}{k} = n2^{n-1}$

Connection with the fibonacci numbers: $\sum_{k=0}^{\infty} {n-k \choose k} = F_{n+1}$

Fibonacci Number

$$k = A - B$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1}$$
(1)

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n \tag{2}$$

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - (-1)^n \tag{3}$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1}$$
 (4)

$$GCD(F_m, F_n) = F_{GCD(m,n)}$$
 (5)

$$\sum_{0 \le k \le n} \binom{n-k}{k} = Fib_{n+1} \tag{6}$$

$$\gcd(F_n, F_{n+1}) = \gcd(F_n, F_{n+2}) = \gcd(F_{n+1}, F_{n+2}) = 1$$
 (7)

Lucas Theorem

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p}$$

 $\binom{m}{n}$ is divisible by p if and only if at least one of the base-p digits of n is greater than the corresponding base-p digit of m.

- The number of entries in the nth row of Pascal's triangle that are not divisible by $p = \prod_{i=0}^{k} (n_i + 1)$

All entries in the $(p^k-1)th$ row are not divisble by p. $\binom{n}{m} \equiv \lfloor \frac{n}{p} \rfloor \pmod{p}$

2nd Kaplansky's Lemma

The number of ways of selecting k objects, no two consecutive, from *n* labelled objects arrayed in a circle is $\frac{n}{h}\binom{n-k-1}{h-1} =$ $\left(\frac{n}{n-k}\binom{n-k}{k}\right)$

Distinct Objects into Distinct Bins

- *n* distinct objects into *r* distinct bins = r^n

- Among *n* distinct objects, exactly *k* of them into r distincts bins = $\binom{n}{k} r^k$

- *n* distinct objects into *r* distinct bins such that each bin contains at least one object = $\sum_{i=0}^{r} (-1)^{i} {r \choose i} (r-i)^{n}$

Stirling Number 2nd Kind

- Count the number of ways to partition a set of n labelled objects into k nonempty unlabelled subsets.

$$S(n,k) = S(n-1,k-1) + k * S(n-1,k)$$

$$S(0,0) = 1, S(>0,0) = 0, S(0,>0) = 0$$

- Time Complexity: $O(k \log n)$

```
ll get sn2(int n, int k) {
 ll sn2 = 0;
 for (int i = 0; i \le k; ++i) {
   ll now = nCr(k, i) * powmod(k - i, n, mod) % mod;
   if (i&1)
     now = now * (mod - 1) % mod:
   sn2 = (sn2 + now) % mod;
 sn2 = sn2 * ifact[k] % mod;
 return sn2;
```

- Number of ways to color a 1n grid using k colors such that each color is used at least once = k!.sn2(n,k)

Bell Numbers

Counts the number of partitions of a set.

$$B_{n+1} = \sum_{k=0}^{n} \left(\frac{n}{k}\right) \cdot B_k \tag{8}$$

 $B_n = \sum_{k=0}^n S(n,k)$, where S(n,k) is stirling number of second

kind. Partition Number

- Time Complexity: $O(n\sqrt{n})$

```
for (int i = 1; i <= n; ++i) {
  pent[2 * i - 1] = i * (3 * i - 1) / 2;
  pent[2 * i] = i * (3 * i + 1) / 2;</pre>
p[0] = 1;
for (int i = 1; i \le n; ++i) {
   p[i] = 0:
   for (int j = 1, k = 0; pent[j] <= i; ++j) {
     if (k < 2) p[i] = add(p[i], p[i - pent[j]]);</pre>
     else p[i] = sub(p[i], p[i - pent[j]]); ++k, k &=
```

- The number of partitions of a positive integer n into exactly k parts equals the number of partitions of n whose largest part equals k

$$p_{k}(n) = p_{k}(n-k) + p_{k-1}(n-1)$$

- The number of labeled undirected graphs with n vertices

 $2^{n(n-1)}$

- The number of connected labeled undirected graphs $16x^4 + 25x^5 + ... e^x = 1 + x + (x^2)/2! + (x^3)/3! + (x^4)/4! + ...$ with *n* vertices, $C_n = 2^{\binom{n}{2}} - \frac{1}{n} \sum_{k=1}^{n-1} k \binom{n}{k} 2^{\binom{n-k}{2}} C_k = 2^{\binom{n}{2}}$ $\sum_{k=1}^{n-1} {n-1 \choose k-1} 2^{{n-k \choose 2}} C_k$

- The number of k-connected labeled undirected graphs with n

= the number of spanning trees of a complete graph with n la- $\left|\sum_{i=1}^{n} f_k(i) = \frac{1}{k+1} n(n+1)(n+2)...(n+k) = \frac{1}{k+1} \frac{(n+k)!}{(n-1)!}\right|$ beled vertices = n^{n-2}

Deleg vertices = n^{n-2} - Number of ways to color a graph using k color such that no $\sum_{i=0}^{n} nix^{i} = 1 + 2x^{2} + 3x^{3} + 4x^{4} + 5x^{5} + ... + nx^{n} = \frac{(x - (n+1)x^{n+1} + nx^{n+2})}{(x-1)^{2}}$ two adjacent nodes have same color

Complete graph = k(k-1)(k-2)...(k-n+1)

Tree =
$$k(k-1)^{n-1}$$

Cycle = $(k-1)^n + (-1)^n(k-1)$

- Number of trees with n labeled nodes: n^{n-2}

Lucas Number

Number of edge cover of a cycle graph C_n is L_n L(n) = L(n-1) + L(n-2); L(0) = 2, L(1) = 1

Catalan Number

$$C_{n+1} = C_0C_n + C_1C_{n-1} + C_2C_{n-2} + \dots + C_nC_0$$

$$C_n = \binom{2n}{n} - \binom{2n}{n+1}$$

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

Derangement

$$D_n = (n-1)(D_{n-1} + D_{n-2}) = nD_{n-1} + (-1)^n$$
$$D_0 = 1.D_1 = 0$$

1,0,1,2,9,44,265,...

Ballot Theorem

Suppose that in an election, candidate A receives a votes and candidate B receives b votes, where a kb for some positive integer k. Compute the number of ways the ballots can be ordered so that A maintains more than k times as many votes as B throughout the counting of the ballots.

The solution to the ballot problem is ((a kb)/(a+b)) * C(a+b, a)**Classical Problem** F(n,k) = number of ways to color n objects using exactly k colors

Het G(n,k) be the number of ways to color n objects using no more than k colors. Then, F(n,k) = G(n,k) - C(k,1) * G(n,k-1) + C(k,2) * G(n,k-2) -

C(k,3) * G(n,k-3)...

Determining G(n, k):

Suppose, we are given a 1 * n grid. Any two adjacent cells can not have same color. Then, $G(n, k) = k * ((k-1)^{(n-1)})$ If no such condition on adjacent cells. Then, $G(n, k) = k^n$

Generating Function

 $\begin{vmatrix} 1/(1-x) = 1 + x + x^2 + x^3 + \dots & 1/(1-ax) = 1 + ax + (ax)^2 + (ax)^3 + \dots \\ 1/(1-x)^2 = 1 + 2x + 3x^2 + 4x^3 + \dots & 1/(1-x)^3 = C(2,2) + C(3,2)x + \dots \end{vmatrix}$ $(k,k)(ax)^2 + C(3+k,k)(ax)^3 + \dots + x(x+1)(1-x)^3 = 1+x+4x^2+9x^3+$

with
$$n$$
 vertices, $C_n = 2^{\binom{n}{2}} - \frac{1}{n} \sum_{k=1}^{n-1} k \binom{n}{k} 2^{\binom{n-k}{2}} C_k = 2^{\binom{n}{2}} - \frac{1}{n} \sum_{k=1}^{n-1} k \binom{n}{k} 2^{\binom{n-k}{2}} C_k = 2^{\binom{n}{2}} - \frac{1}{n} \sum_{k=1}^{n-1} \binom{n-1}{k-1} 2^{\binom{n-k}{2}} C_k$

The number of k-connected labeled undirected graphs with n vertices, $D[n][k] = \sum_{s=1}^{n} \binom{n-1}{s-1} C_s D[n-s][k-1]$

Cayley's formula: the number of trees on n labeled vertices n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates n the number of spanning trees of a complete graph with n lates

Probability

$$P(A|B) = \frac{P(A \cap B)}{P(B)}$$

$$P(A|B) = \frac{P(B|A)P(A)}{P(B)}$$

Matching Formula Normal Graph

MM + MEC = n (without isolated vertex)

IS + VC = GMaxIS + MVC = G

Bipartite Graph

Solution of $x^2 \equiv a \pmod{p}$:

 $-ca \equiv cb \pmod{m} \iff a \equiv b \pmod{\frac{n}{\gcd(n,c)}}$

 $ax \equiv b \pmod{m}$ has a solution $\iff \gcd(a,m)|b|$

- If $ax \equiv b \pmod{m}$ has a solution, then it has gcd(a,m) solutions and they are separated by $\frac{n}{\gcd(a,m)}$

 $-ax \equiv 1 \pmod{m}$ has a solution or a is invertible $\pmod{m} \iff$ gcd(a, m) = 1

 $x^2 \equiv 1 \pmod{p}$ then $x \equiv \pm 1 \pmod{p}$

- There are $\frac{p-1}{2}$ has no solution.

There are $\frac{p^{-1}}{2}$ has exactty two solutions.

- When $p\%4 = 3, x \equiv \pm a^{\frac{p+1}{4}}$

- When p%8 = 5, $x \equiv a^{\frac{p+3}{8}}$ or $x \equiv 2^{\frac{p-1}{4}}a^{\frac{p+3}{8}}$

Totient

- If p is a prime $(p^k) = p^k - p^{k-1}$

- If a b are relatively prime, $\phi(ab) = \phi(a)\phi(b)$

 $-\phi(n) = n(1-\frac{1}{n_1})(1-\frac{1}{n_2})(1-\frac{1}{n_2})...(1-\frac{1}{n_k})$

- Sum of coprime to $n = n * \frac{\phi(n)}{2}$

- If $n = 2^k$, $\phi(n) = 2^{k-1} = \frac{n}{2}$

- For a b, $\phi(ab) = \phi(a)\phi(b)\frac{d}{\phi(d)}$

 $\phi(ip) = p\phi(i)$ whenever p is a prime and it divides i

The number of $a(1 \le a \le N)$ such that gcd(a, N) = d is $\phi(\frac{n}{d})$

If n > 2, $\phi(n)$ is always even

Sum of gcd, $\sum_{i=1}^{n} gcd(i,n) = \sum_{d|n} d\phi(\frac{n}{d})$

Sum of lcm, $\sum_{i=1}^{n-1} nlcm(i,n) = \frac{n}{2} (\sum_{d|n} (d\phi(d)) + 1)$

 $\phi(1) = 1$ and $\phi(2) = 1$ which two are only odd ϕ

 $\phi(3) = 2$ and $\phi(4) = 2$ and $\phi(6) = 2$ which three are only prime

- Find minimum n such that $\frac{\phi(n)}{n}$ is maximum- Multiple of small primes- 2*3*5*7*11*13*...

Mobius

$$\sum_{i=1}^n \sum_{j=1}^n [gcd(i,j)=1] = \sum_{k=1}^n \mu(k) \lfloor \frac{n}{k} \rfloor^2$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = \sum_{k=1}^{n} k \sum_{l=1}^{\lfloor \frac{n}{k} \rfloor} \mu(l) \lfloor \frac{n}{kl} \rfloor^{2}$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} gcd(i,j) = \sum_{k=1}^{n} \left(\frac{\lfloor \frac{n}{k} \rfloor (1 + \lfloor \frac{n}{k} \rfloor)}{2}\right)^{2} \sum_{d \mid k} \mu(d)kd$$