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1 2ECC

```
struct graph {
  int n, t, sz;
  vector<vector<int>> adj;
  vector<int> tin, low, cmp;
  graph(int n): n(n),adj(n),tin(n),low(n),cmp(n){}
  void add_edge(int u, int v){
   adj[u].push back(v);
    adj[v].push_back(u);
  void dfs(int u, int p){
    tin[u]=low[u]=t++;
    int cnt=0;
    for(int v: adj[u]){
      if(v==p and ++cnt <= 1) continue;</pre>
      if(tin[v]!=-1) low[u] = min(low[u], tin[v]);
        dfs(v,u);
        low[u] = min(low[u], low[v]);
  void dfs2(int u, int p){
   if(p!=-1 \text{ and } tin[p]>=low[u]) cmp[u] = cmp[p];
    else cmp[u] = sz++;
   for(int v: adj[u]){
      if(cmp[v]==-1) dfs2(v,u);
  void process 2ecc(){
   t = 0, sz = 0;
    for (int i = 0; i < n; ++i){
      tin[i] = low[i] = cmp[i] = -1;
    for (int i = 0; i < n; ++i){
      if(tin[i]==-1) dfs(i,-1);
    for (int i = 0; i < n; ++i){
      if(cmp[i]==-1) dfs2(i,-1);
```

2 2SAT

```
//CNF: (a | b) ^ (c | d) means (!a -> b) ^ (!b -> a)
// (!a or b) = (-a, b), 1-based indexing
string two sat(int n, vector<array<int, 2>>

    clauses) {
  vector<int> adj[2 * n];
  for (auto [a, b]: clauses) {
    if (a > 0) a = 2 * a - 2;
    else a = 2 * -a - 1:
    if (b > 0) b = 2 * b - 2;
    else b = 2 * -b - 1;
    adj[a ^ 1].push_back(b), adj[b ^

→ 1].push back(a);

  vector<vector<int>>> sccs = get sccs(2 * n, adj);
  int tot scc = sccs.size();
  vector<int> scc no(2 * n);
  for (int i = 0; i < tot scc; ++i) {
   for (int u: sccs[i]) {
```

```
scc_no[u] = i;
}
string assignment;
for (int u = 0; u < n; u++) {
   if (scc_no[2 * u] == scc_no[2 * u + 1]) {
      return "";
   }
   if (scc_no[2 * u] < scc_no[2 * u + 1]) {
      assignment += '-';
   }
   else {
      assignment += '+';
   }
}
return assignment;</pre>
```

3 AHO_CORASICK

```
struct AC {
 const int A = 26;
 vector<vector<int>> nxt, idx;
  vector<int> lnk, out lnk, ans;
  AC () { newNode(); }
 int newNode() {
  nxt.eb(A, 0), idx.eb(0);
    lnk.eb(0), out lnk.eb(0), ans.eb(0);
    return nxt.size() - 1;
 void clear () {
  nxt.clear(), idx.clear();
    lnk.clear(), out lnk.clear(), ans.clear();
    newNode();
 int add (string p, int i) {
    int u = 0;
    for (auto c: p) {
     int id = c - 'a';
      if (!nxt[u][id]) nxt[u][id] = newNode();
      u = nxt[u][id];
    idx[u].eb(i);
    return u;
 void build () {
    queue<int> q; q.push(0);
    while (!q.empty()) {
      int u = q.front(); q.pop();
      for (int i = 0; i < A; ++i) {
        int v = nxt[u][i];
        if (!v) nxt[u][i] = nxt[lnk[u]][i];
          lnk[v] = u? nxt[lnk[u]][i]: 0;
          out lnk[v] = idx[lnk[v]].empty()?
  out lnk[lnk[v]]: lnk[v];
          q.push(v);
          // dp[v] = dp[v] + dp[lnk[v]]
 void trav (string T) {
    int u = 0;
    for (auto c: T) {
```

```
int id = c - 'a';
  while (u and !nxt[u][id])  u = lnk[u];
  u = nxt[u][id];
  int x = u;
  while (x) {
    for (auto i: idx[x]) {
        ans[i]++;
    }
    x = out_lnk[x];
  }
}
//AC ac; ac.add(pi, i); ac.build(); ac.trav(T);
```

4 ARTICULATION_POINT

```
vector<int> adj[N];
int t = 0:
vector<int> tin(N, -1), low(N), ap;
void dfs (int u, int p) {
 tin[u] = low[u] = t++;
  int is ap = 0, child = 0;
 for (int v: adj[u]) {
    if (v != p) {
      if (tin[v] != -1) |
        low[u] = min(low[u], tin[v]);
      else {
   child++;
        dfs(v, u);
        if (tin[u] <= low[v]) {
          is ap = 1;
         low[u] = min(low[u], low[v]);
 if ((p != -1 \text{ or child} > 1) \text{ and is ap})

→ ap.push back(u);

dfs(0, -1);
```

5 BCC

```
struct graph {
 int n, t=0, cno=0;
 vector<vector<int>> g;
 vector<int> tin, lo, bcomp;
 stack <int> st;
 graph(int n):n(n),g(n),lo(n),bcomp(n){}
 void add edge(int u, int v){
   g[u].push back(v);
   g[v].push_back(u);
 void dfs(int v, int p=-1){
   lo[v]=tin[v]=++t;
   st.push(v);
   for(int u:g[v]){
     if(u==p)
                 continue;
     if(!tin[u]){
        dfs(u, v);
        lo[v]=min(lo[v],lo[u]);
     } else{
```

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```

```
lo[v]=min(lo[v],tin[u]);
 if(tin[v]==lo[v]){
   while (!st.empty()){
     int tp=st.top(); st.pop();
      bcomp[tp]=cno;
     if(tp==v)
                  break;
   ćno++:
vector<int> bcc(){
 tin.assign(n, 0);
 for (int^i = 0; i < n; ++i){
   if(!tin[i])
     dfs(i);
  return bcomp;
```

```
6 BCC EDGE
vector<array<int, 2>> edges, adj[N];
vector \langle int \rangle tin(N), lo(N), is ap(N), bcc[N],
→ bcc ed[N]:
int t = 0, tot = 0;
stack<int> stk;
void pop bcc(int e) {
  do {
    bcc ed[tot].push back(stk.top()); stk.pop();
  } while (bcc ed[to\overline{t}].back() != e):
void dfs(int u, int p = -1) {
  int ch = 0;
  tin[u] = lo[u] = t++;
  for(auto [v, e] : adj[u]) {
    if (v == p) continue;
if (tin[v] != -1) {
      if (tin[u] > tin[v]) {
        lo[u] = min(lo[u], tin[v]);
         stk.push(e);
    else {
      ch++;
      stk.push(e);
      dfs(v, u);
      if ((p != -1 \text{ or } ch > 1) \text{ and } tin[u] <= lo[v]) { | void pop_bcc(int u, int v) } {}
        is ap[u] = 1;
        pop bcc(e);
       lo[u] = min(lo[u], lo[v]);
void procces bcc(int n) {
  for (int i = 0; i < n; ++i) {
    tin[i] = -1, is ap[i] = 0;
    bcc_ed[i].clear();
    bccTil.clear();
```

```
t = tot = 0;
 for (int u = 0; u < n; ++u) {
   if (tin[u] == -1) {
     dfs(u, -1);
     if (!stk.empty()) {
        while (!stk.empty()) {
          bcc ed[tot].push back(stk.top());
  stk.pop();
        tot++;
     }
 for (int i = 0; i < tot; ++i) {
  for (auto e: bcc_ed[i]) {</pre>
     auto [u, v] = \overline{e}dges[e];
     bcc[i].push back(u);
     bcc[i].push_back(v);
 for (int i = 0; i < tot; ++i) {
   sort(bcc[i].begin(), bcc[i].end());
   bcc[i].erase(unique(bcc[i].begin(),
  bcc[i].end()), bcc[i].end());
}
```

7 BIT_TRICKS

```
## Next Combination Mask
int next combs mask(int mask) {
 int lsb = -mask & mask;
  return (((mask + lsb) ^ mask) / (lsb << 2))
  (mask + lsb);
## Iterate over submask in decreasing order
for (int submask=mask; submask > 0; submask =
   (submask-1)\&mask) {
```

8 BLOCK_CUT_TREE

```
vector<int> adj[N];
vector<int> tin(N, -1), lo(N), is ap(N), bcc[N];
stack<int> stk;
int t = 0, tot = 0;
  bcc[tot].push back(u);
  while (bcc[tot].back() != v)
    bcc[tot].push back(stk.top());
    stk.pop();
  tot++;
void dfs (int u, int p) {
  tin[u] = lo[u] = t++;
  stk.push(u);
  int ch = 0;
  for (auto v: adj[u]) {
    if (v != p) {
```

```
if (tin[v] != -1) {
        lo[u] = min(lo[u], tin[v]);
      else {
        ch++;
        dfs(v, u);
        if ((p != -1 \text{ or ch} > 1) \text{ and tin}[u] <=
→ lo[v]) {
          // is ap[u] = 1;
           pop bcc(u, v);
         lo[u] = min(lo[u], lo[v]);
void process bcc (int n) {
 for (int u = 0; u < n; ++u) {
    tin[u] = -1;
    is ap[u] = 0;
    bcc[u].clear();
 t = tot = 0;
 for (int u = 0; u < n; ++u) {
    if (tin[u] == -1) {
      dfs(u, -1);
      if (!stk.empty()) {
        while (!stk.empty()) {
  bcc[tot].push_back(stk.top());
          stk.pop();
        tot++;
vector<int> comp_num(N), bct_adj[N];
void build bct(int n) {
 process \overline{b}cc(n);
  int nn = tot;
 for (int u = 0; u < n; ++u) {
    if (is ap[u]) {
      comp^-num[u] = nn++;
  for (int i = 0; i < tot; ++i) {
    for (auto u: bcc[i]) {
      if (is ap[u])
        u = \overline{comp} num[u];
        bct adj[i].push_back(u);
        bct adj[u].push back(i);
      else {
        comp num[u] = i;
```

9 CDQ

```
## Problems related to pair

    cdq(l, m)
```

```
- cdq(m + 1, r)
- handle influence of (l, m) to (m + 1, r)

## Optimization of 1D DP
- cdq(l, m)
- handle influence of (l, m) to (m + 1, r)
- cdq(m + 1, r)

## Convert dynamic array offline problem to static
- array offline problem
```

10 CENTROID_DECOMPOSITION

```
void calc sz(int u, int p) {
  sz[u] = 1;
  for (auto v: adj[u]) {
    if (v != p and !is cen[v]) {
      calc sz(v, u);
      sz[u] += sz[v];
int get cen(int u, int p, int n) {
 for (auto v: adj[u]) {
    if (v != p \text{ and } ! \text{ is cen}[v] \text{ and } 2 * \text{sz}[v] > n)
      return get cen(v, u, n);
  return u;
void decompose(int u=0, int p=-1, int d=0){
  calc_sz(u, p); int c = get_cen(u, p, sz[u]);
  is cen[c] = 1, cpar[c] = p, cdep[c] = d;
  for(int v: adi[c]){
    if(!is cen[v]) decompose(v,c,d+1);
decompose();
```

11 CONVOLUTION

```
## FFT
struct cplx {
  ld a, b;
  cplx(ld a=0, ld b=0):a(a), b(b) {}
  const cplx operator + (const cplx &z) const {

    return cplx(a+z.a, b+z.b); }

  const cplx operator - (const cplx &z) const {
→ return cplx(a-z.a, b-z.b); }
  const cplx operator * (const cplx &z) const {
\rightarrow return cplx(a*z.a-b*z.b, a*z.b+b*z.a); }
  const cplx operator / (const ld &k) const {

→ return cplx(a/k, b/k); }

const ld PI=acos(-1);
vector<int> rev;
void pre(int sz){
  if(rev.size()==sz) return ;
  rev.resize(sz);
  rev[0]=0;
  int lg n = builtin ctz(sz);
  for (int i = 1; i < sz; ++i) rev[i] = (rev[i>>1]
\rightarrow >> 1) | ((i&1)<<(lq n-1));
void fft(vector<cplx> &a. bool inv){
  int n = a.size();
```

```
for (int i = 1; i < n-1; ++i) if(i<rev[i])
   swap(a[i], a[rev[i]]);
  for (int len = 2; len <= n; len <<= 1) {
    ld t = 2*PI/len*(inv? -1: 1);</pre>
    cplx wlen = {cosl(t), sinl(t)};
    for (int st = 0; st < n; st += len){</pre>
      cplx w(1);
      for (int i = 0; i < len/2; ++i){
        cplx ev = a[st+i];
        cplx od = a[st+i+len/2]*w;
        a[st+i] = ev+od;
        a[st+i+len/2] = ev-od;
        w = w*wlen;
  if(inv){
    for(cplx &z: a){
      z = z/n;
vector<ll> mul(vector<ll> &a, vector<ll> &b){
  int n = a.size(), m = b.size(), sz = 1;
  while (sz < n+m-1) sz <<= 1;
  vector<cplx> x(sz), y(sz), z(sz);
  for (int i = 0; i < sz; ++i){
    x[i] = cplx(i < n? a[i]: 0, 0);
    y[i] = cplx(i < m? b[i]: 0, 0);
  pre(sz);
  fft(x, 0);
  fft(y, 0);
  for (int i = 0; i < sz; ++i){}
    z[i] = x[i] * y[i];
  fft(z, 1);
  vector<ll> c(n+m-1);
  for (int i = 0; i < n+m-1; ++i){
   c[i] = round(z[i].a);
  return c;
## NTT
const int mod = 998244353;
|const int root = 15311432;
const int k = 1 << 23;
lint root 1:
vector<int> rev;
ll bigmod(ll a. ll b. ll mod){
  a %= mod;
  ll ret = 1;
  while(b){
    if(b\&1) ret = ret*a%mod;
    a = a*a*mod;
    b >>= 1:
  return ret;
void pre(int sz){
  root 1 = bigmod(root, mod-2, mod);
```

```
if(rev.size()==sz) return ;
  rev.resize(sz);
  rev[0]=0;
  int lg n = __builtin_ctz(sz);
 for (int i = 1; i < sz; ++i) rev[i] = (rev[i>>1]
  >> 1) | ((i\&1) << (lq n-1));
void fft(vector<int> &a, bool inv){
 int n = a.size():
 for (int i = 1; i < n-1; ++i) if(i<rev[i])

    swap(a[i], a[rev[i]]);

  for (int len = 2; len <= n; len <<= 1) {
    int wlen = inv ? root 1 : root;
    for (int i = len; i < k; i <<= 1){
      wlen = 1ll*wlen*wlen%mod;
    for (int st = 0; st < n; st += len) {
      int w = 1;
      for (int j = 0; j < len / 2; j++) {
        int ev = a[st+j];
        int od = 1ll*a[st+j+len/2]*w%mod;
        a[st+j] = ev + od < mod ? ev + od : ev + od
        a[st+j+len/2] = ev - od >= 0 ? ev - od : ev
    - od + mod;
        w = 1li * w * wlen % mod;
 if (inv) {
    int n 1 = bigmod(n, mod-2, mod);
    for (int & x : a)
     x = 111*x*n 1%mod;
vector<int> mul(vector<int> &a, vector<int> &b){
 int n = a.size(), m = b.size(), sz = 1;
 while (sz < n+m-1) sz <<= 1;
 vector<int> x(sz), y(sz), z(sz);
 for (int i = 0; i < sz; ++i){
  x[i] = i < n? a[i]: 0;</pre>
    y[i] = i < m? b[i]: 0;
  pre(sz);
  fft(x, 0);
 fft(y, 0);
 for (int i = 0; i < sz; ++i){
    z[i] = 1ll* x[i] * y[i] % mod;
 fft(z, 1);
 z.resize(n+m-1);
 return z;
## Any mod
const int N = 3e5 + 9, mod = 998244353;
struct base {
 double x, y;
 base() { x = y = 0; }
 base(double x, double y): x(x), y(y) { }
```

```
inline base operator + (base a, base b) { return
\rightarrow base(a.x + b.x, a.y + b.y); }
inline base operator - (base a, base b) { return
→ base(a.x - b.x, a.y - b.y); }
inline base operator * (base a, base b) { return
    base(a.x * b.x - a.y * b.y, a.x * b.y + a.y *
 \stackrel{\sim}{\rightarrow} b.x): }
inline base conj(base a) { return base(a.x, -a.y); }
int lim = 1:
vector<br/>vector<br/>vector<br/>se> roots = \{\{0, 0\}, \{1, 0\}\};
vector < int > rev = \{0, 1\};
const double PI = acosl(- 1.0);
void ensure base(int p) {
  if(p <= lim) return;</pre>
  rev.resize(1 << p);
  for(int i = 0; i < (1 << p); i++) rev[i] = (rev[i
\rightarrow >> 1] >> 1) + ((i & 1) << (p - 1));
  roots.resize(1 << p);</pre>
  while(lim < p) {</pre>
    double angle = 2 * PI / (1 << (lim + 1));
    for(int i = 1 << (lim - 1); i < (1 << lim);</pre>
      roots[i << 1] = roots[i];
      double angle i = angle * (2 * i + 1 - (1 << 
      roots[(i \ll 1) + 1] = base(cos(angle i),
   sin(angle i));
    lim++;
void fft(vector<base> \deltaa, int n = -1) {
 if(n == -1) n = a.size();
  assert((n \& (n - 1)) == 0);
  int zeros = builtin ctz(n);
  ensure base(zeros);
  int shift = lim - zeros;
  for(int i = 0; i < n; i++) if(i < (rev[i] >>
shift)) swap(a[i], a[rev[i] >> shift]);
  for(int k = 1; k < n; k <<= 1) {
    for(int i = 0; i < n; i += 2 * k) {
      for(int j = 0; j < k; j++) {
        base z = a[i + j + k] * roots[j + k];
        a[i + j + k] = a[i + j] - z;
        a[i + j] = a[i + j] + z;
//eq = 0: 4 FFTs in total
//eq = 1: 3 FFTs in total
vector<int> multiply(vector<int> &a, vector<int>
\rightarrow &b, int eq = 0)
  int need = a.size() + b.size() - 1;
  int p = 0;
  while((1 << p) < need) p++;
  ensure base(p);
  int sz^{-} = 1 << p;
  vector<base> A, B;
  if(sz > (int)A.size()) A.resize(sz);
  for(int i = 0; i < (int)a.size(); i++) {</pre>
    int x = (a[i] \% mod + mod) \% mod;
    A[i] = base(x \& ((1 << 15) - 1), x >> 15);
```

```
fill(A.begin() + a.size(), A.begin() + sz,
   base{0, 0});
 fft(A, sz);
 if(sz > (int)B.size()) B.resize(sz);
 if(eq) copy(A.begin(), A.begin() + sz, B.begin());
    for(int i = 0; i < (int)b.size(); i++) {</pre>
      int x = (b[i] \% mod + mod) \% mod;
      B[i] = base(x \& ((1 << 15) - 1), x >> 15);
    fill(B.begin() + b.size(), B.begin() + sz,
   base{0, 0});
    fft(B, sz);
  double ratio = 0.25 / sz;
 base r2(0, -1), r3(ratio, 0), r4(0, -ratio),
\sim r5(0, 1);
 for(int i = 0; i \le (sz >> 1); i++) {
    int j = (sz - i) \& (sz - 1);
    base a1 = (A[i] + conj(A[j])), a2 = (A[i] -
   conj(A[j])) * r2;
    base b1 = (B[i] + conj(B[j])) * r3, b2 = (B[i])
    - conj(B[j])) * r4;
    if(i != j) {
      base c1 = (A[j] + conj(A[i])), c2 = (A[j] -
   coni(A[i])) * r2;
      base d1 = (B[i] + coni(B[i])) * r3, d2 =
   (B[j] - conj(B[i])) * r4;
      A[i] = c1 * d1 + c2 * d2 * r5;
      B[i] = c1 * d2 + c2 * d1;
    A[i] = a1 * b1 + a2 * b2 * r5;
   B[j] = a1 * b2 + a2 * b1:
 fft(A, sz); fft(B, sz);
 vector<int> res(need);
 for(int i = 0; i < need; i++) {
    long long aa = A[i].x + 0.5;
    long long bb = B[i].x + 0.5;
    long long cc = A[i].y + 0.5;
    res[i] = (aa + ((bb \% mod) << 15) + ((cc \% mod))
   << 30))%mod;
 return res;
vector<int> pow(vector<int>& a, int p) {
 vector<int> res;
 res.emplace back(1);
 while(p) {
    if(p \& 1) res = multiply(res, a);
    a = multiply(a, a, 1);
    p >>= 1;
  return res;
int main() {
 int n, k; cin >> n >> k;
 vector<int> a(10, 0);
 while(k--) {
    int m; cin >> m;
    a[m] = 1;
 vector<int> ans = pow(a, n / 2);
```

```
int res = 0;
 for(auto x: ans) res = (res + 1LL * x * x * mod)

→ % mod:

 cout << res << '\n':
 return 0;
## Online NTT
void solve() {
 f[0]=1; // base case
 for(int i=0; i<=MAX; i++) {
    // Doing the part 1\,
    f[i+1]=(f[i+1]+f[i]*A[0])%mod;
    f[i+2]=(f[i+2]+f[i]*A[1])%mod;
    if(!i) continue;
    // part 2
    int limit=(i&-i);
    for(int p=2; p<=limit; p*=2) +</pre>
      convolve(i-p,i-1,p,min(2*p-1,MAX));
void convolve(int l1, int r1, int l2, int r2) {
 int n=max(r1-l1+1,r2-l2+1);
 int t=1:
 while(t<n) t<<=1;
 vector<ll> a(n), b(n);
 for(int i=l1; i<=r1; i++) a[i-l1]=f[i];</pre>
 for(int i=l2: i<=r2: i++) b[i-l2]=A[i]:
 vector<ll> ret=fft::multiply(a,b);
    for(int i=0; i<ret.size(); i++) {</pre>
    int idx=i+l1+l2+1;
    if(idx>MAX) break;
    // adding to the appropriate entry
    f[idx]+=ret[i];
    f[idx]%=mod;
## FWHT (AND, OR, XOR)
- Time complexity: O(nloan)
- AND, OR works for any modulo, XOR works for only
- size must be power of two
const ll mod = 998244353;
int add (int a, int b) {
 return a + b < mod? a + b: a + b - mod;
int sub (int a, int b) {
 return a - b >= 0? a - b: a - b + mod;
ll poww (ll a, ll p, ll mod){
 a %= mod;
 ll ret = 1:
 while (p){
   if (p & 1) {
      ret = ret * a % mod;
   a = a * a % mod;
   p >>= 1;
 return ret:
```

```
void fwht(vector<int> &a, int inv, int f) {
  int sz = a.size();
  for (int len = 1; 2 * len <= sz; len <<= 1) {</pre>
    for (int i = 0; i < sz; i += 2 * len) {
      for (int j = 0; j < len; j++) {
        int x = a[i + j];
        int y = a[i + j + len];
        if (f == 0) {
          if (!inv) a[i + j] = y, a[i + j + len] =
\rightarrow add(x, y);
          else a[i + j] = sub(y, x), a[i + j +
   len] = x;
        else if (f == 1) {
          if (!inv) a[i + j + len] = add(x, y);
          else a[i + j + len] = sub(y, x);
        else {
          a[i + j] = add(x, y);
          a[i + j + len] = sub(x, y);
vector<int> mul(vector<int> a, vector<int> b, int
\rightarrow f) { // 0:AND, 1:0R, 2:X0R
  int sz = a.size();
  fwht(a, 0, f); fwht(b, 0, f);
  vector<int> c(sz);
  for (int i = 0; i < sz; ++i) {
    c[i] = 111 * a[i] * b[i] % mod;
  fwht(c, 1, f);
  if (f) {
    int sz inv = poww(sz, mod - 2, mod);
    for (int i = 0; i < sz; ++i) {
      c[i] = 1ll * c[i] * sz inv % mod;
  return c;
## subset convolution
vector<int> subset conv (vector<int> a, vector<int>
→ b) {
  int n = a.size();
  int lq = log2(n);
  vector<int> cnt(n);
  vector<vector<int>> fa(lg + 1, vector<int> (n)),
    fb(lq + 1, vector < int > (n)), q(lq + 1,
   vector<int> (n));
  for (int i = 0; i < n; ++i) {
    cnt[i] = cnt[i >> 1] + (i & 1);
    fa[cnt[i]][i] = a[i] % mod;
fb[cnt[i]][i] = b[i] % mod;
  for (int k = 0; k \le lq; ++k)
    fwht(fa[k], 0, 1); fwht(fb[k], 0, 1);
  for (int k = 0; k \le lq; ++k) {
    for (int j = 0; j \le k; ++j) {
```

```
for (int i = 0; i < n; ++i) {
    g[k][i] = add(g[k][i], lll * fa[j][i] *
    fb[k - j][i] % mod);
    }
}
for (int k = 0; k <= lg; ++k) {
    fwht(g[k], l, l);
}
vector<int> c(n);
for (int i = 0; i < n; ++i) {
    c[i] = g[cnt[i]][i];
}
return c;
}</pre>
```

12 CPP

```
## Ordered Set
#include <ext/pb ds/assoc container.hpp>
using namespace __gnu_pbds;
typedef tree<int, null type, less<int>, rb tree tag,
tree order statistics node update> oset;
## unordered map
struct chash₹
  size t operator()(const pair<int,int>&x)const{
    return hash<long long>()(((long
   long()x.first()^(((long long()x.second()<<32));</pre>
|unordered map<pair<int, int>, int, chash> maf;
maf.reserve(max len);
|maf.max| load factor(0.25);
## gp hash table:
#include <ext/pb ds/assoc container.hpp>
using namespace gnu pbds;
struct chash{
  int operator()(ii p) const {
    return p.first*31 + p.second;
gp_hash table<ii, int, chash> cnt;
## Random Number
→ epoch().count()):
|int x = rng() % 495;
## Running time
|clock t st = clock();
double t = (clock() - st) / (1.0 * CLOCKS PER SEC);
string line; getline(cin, line);
istringstream iss;
string word;
|while (iss >> word) {
 cout << word << "\n";
#pragma GCC target("popcnt") ulimit -s 65532
```

13 DETERMINANT

```
const double EPS = 1E-9;
int n;
vector < vector<double> > a (n, vector<double> (n));
double det = 1;
```

```
for (int i=0; i<n; ++i) {
 int k = i;
 for (int j=i+1; j<n; ++j)
   if (abs (a[j][i]) > abs (a[k][i]))
     k = j;
 if (abs (a[k][i]) < EPS) {
   det = 0;
   break;
 swap (a[i], a[k]);
 if (i != k)
   det = -det;
 det *= a[i][i];
 for (int j=i+1; j<n; ++j)
   a[i][j] /= a[i][i];
 for (int j=0; j<n; ++j)
   if (j != i \&\& abs (a[j][i]) > EPS)
     for (int k=i+1; k<n; ++k)</pre>
       a[j][k] -= a[i][k] * a[j][i];
```

14 DINIC

```
// V^2E, sqrt(E)E, sqrt(V)E(bpm)
// Effective flows are adj[u][3] where adj[u][3] > 0
ll get max flow(vector<array<int, 3>> edges, int n,
int s, int t) {
vector<array<ll, 4>> adj[n];
 for (auto [u, v, c]: edges) {
   adj[u].push back({v, (int)adj[v].size(), c, 0});
    adj[v].push back({u, (int)adj[u].size() - 1, 0,
   0});
 ll max flow = 0;
 while (true) {
   queue<int> q; q.push(s);
    vector < int > dis(n, -1); dis[s] = 0;
    while (!q.empty()) {
      int u = q.front(); q.pop();
      for (auto [v, idx, c, f]: adj[u]) {
        if (dis[v] == -1 \text{ and } c > f) {
          q.push(v);
          dis[v] = dis[u] + 1;
    if (dis[t] == -1) break;
    vector<int> next(n);
    function<ll(int, ll) > dfs = [\&] (int u, ll)
   flow) {
      if (u == t) return flow;
     while (next[u] < adj[u].size()) {</pre>
        auto \&[v, idx, c, f] = adj[u][next[u]++];
        if (c > f \text{ and } dis[v] == dis[u] + 1) {
          ll bn = dfs(v, min(flow, c - f));
          if (bn > 0) {
            f += bn;
            adj[v][idx][3] -= bn;
            return bn;
```

```
return Oll;
};
while (ll flow = dfs(s, LLONG_MAX)) {
   max_flow += flow;
}
return max_flow;
}
```

15 DOMINATOR_TREE

```
const int N = 2e5+5:
vector <int> g[N], rg[N], dtree[N], bucket[N];
int sdom[N], par[N], dom[N], dsu[N], lab[N],

    arr[N], rev[N], dpar[N], n, ts, src;

void init(int n, int s) {
  ts = 0, n = n, src = s;
  for (int i = 1; i <= n; ++i) {
    g[i].clear(), rg[i].clear(), dtree[i].clear(),
   bucket[i].clear();
    sdom[i]=par[i]=dom[i]=dsu[i]=lab[i]=arr[i]=rev[_
   il=dpar[i]=0;
void dfs(int u) {
   ts++; arr[u] = ts; rev[ts] = u;
   lab[ts] = sdom[ts] = dsu[ts] = ts;
   for(int &v : g[u]) {
      if(!arr[v]) { dfs(v); par[arr[v]] = arr[u]; }
      rg[arr[v]].push back(arr[u]);
inline int root(int u, int x = 0) {
   if(u == dsu[u]) return x ? -1 : u;
   int v = root(dsu[u], x + 1);
   if(v < 0) return u;
   if(sdom[lab[dsu[u]]] < sdom[lab[u]]) lab[u] =</pre>
 → lab[dsu[u]]:
   dsu[u] = v; return x ? v : lab[u];
void build() {
   dfs(src);
   for(int i=n; i; i--) {
      for(int j : rq[i]) sdom[i] =
   min(sdom[i],sdom[root(j)]);
      if(i > 1) bucket[sdom[i]].push back(i);
      for(int w : bucket[i]) {
         int v = root(w);
         if(sdom[v] == sdom[w]) dom[w] = sdom[w];
         else dom[w] = v;
      \} if(i > 1) dsu[i] = par[i];
   for(int i=2; i<=n; i++) {
      int \&dm = dom[i];
      if(dm ^ sdom[i]) dm = dom[dm];
      dtree[rev[i]].push back(rev[dm]);
      dtree[rev[dm]].push back(rev[i]);
      dpar[rev[i]] = rev[\overline{dm}];
   }
```

16 DP_ON_TREE

```
// Rerooting Technique
vector<array<ll, 2>> down(N), up(N);
void dfs() {
    // calculate down dp
}
void dfs2() {
    ll pref = ?;
    for (auto v: adj[u]) {
        // update up[v] and pref
}
    reverse(adj[u].begin(), adj[u].end());
    ll suf = ?;
    for (auto v: adj[u]) {
        // update up[v] and suf
}
for (auto v: adj[u]) {
        dfs2(v)
    }
}
```

17 DP_OPTIMIZATION

```
## CHT
|vector<int> m, c;
 // Insert
int mi, ci; cin >> mi >> ci;
while (sz >= 2) {
  if (111*(ci-c[sz-2])*(m[sz-2]-m[sz-1]) <=
 \rightarrow 1ll*(c[sz-1]-c[sz-2])*(m[sz-2]-mi)) {
     m.pop back(); c.pop back();
  } else break;
m.push back(mi); c.push back(ci); sz++;
 // Query
int x; cin >> x;
int lo = 0, hi = sz-1;
while (lo < hi) {
  int mid = lo+hi>>1;
  if (1ll*m[mid]*x+c[mid] >
 \rightarrow 1ll*m[mid+1]*x+c[mid+1]) lo = mid+1;
  else hi = mid:
cout << 1ll*m[lo]*x+c[lo] << '\n';
## Dynamic CHT
const ll IS QUERY = -(1LL << 62);
struct line {
  ll m, b;
  mutable function <const line*()> succ;
  bool operator < (const line &rhs) const {</pre>
     if (rhs.b != IS QUERY) return m < rhs.m;</pre>
     const line *s = succ();
     if (!s) return 0;
     ll x = rhs.m:
     return b - s -> b < (s -> m - m) * x;
struct CHT : public multiset <line> {
  bool bad (iterator y) {
     auto z = next(y);
     if (y == begin()) {
       if (z == end()) return 0;
       return y \rightarrow m == z \rightarrow m \&\& y \rightarrow b <= z \rightarrow b;
```

```
auto x = prev(y);
    if (z == end()) return y \rightarrow m == x \rightarrow m \& \& y \rightarrow
    b \le x -> b;
    return 1.0 \dot{*} (x -> b - y -> b) * (z -> m - y ->
    m) >= 1.0 * (y -> b - z -> b) * (y -> m - x ->
  void add (ll m, ll b) {
  auto y = insert({m, b});
    y \rightarrow succ = [=] \{return \ next(y) == end() ? 0 :
   &*next(y);};
    if (bad(y)) {erase(y); return;}
    while (next(y) != end() \&\& bad(next(y)))
   erase(next(y));
    while (y != begin() \&\& bad(prev(y)))
    erase(prev(v)):
  ll eval (ll x) {
    auto l = *lower bound((line) {x, IS QUERY});
    return l.m * x + l.b;
};
// To find maximum
CHT cht;
cht.add(m, c);
y max = cht.eval(x);
/7 To find minimum
CHT cht:
cht add(-m, -c);
y min = -cht.eval(x);
// Divide an array into k parts
// Minimize the sum of squre of each subarray
ll pref[N], dp[N][N];
void compute(int l, int r, int j, int kl, int kr) {
  if (l > r) return ;
  int m = (l + r) / 2;
  array<ll, 2> best = {LLONG MAX, -1};
  for (int k = kl; k \leftarrow min(m - 1, kr); ++k) {
best = min(best, \{dp[k][j - 1] + (pref[m] - min(m)\}
 → pref[k]) * (pref[m] - pref[k]), k});
  dp[m][j] = best[0];
  compute(l, m - 1, j, kl, best[1]);
  compute(m + 1, r, j, best[1], kr);
// Divide an array into n parts.
// Cost of each division is subarray sum
// Minimize the cost
(1 dp[n][n], opt[n][n];
for (int i = 0; i < n; ++i)
  for (int j = 0; j < n; ++j) {
    dp[i][j] = LLONG MAX;
  opt[i][i] = i;
  dp[i][i] = 0;
for (int i = n - 2; i >= 0; --i) {
  for (int j = i + 1; j < n; ++j) {
    for (int k = opt[i][j - 1]; k <= min(j - 1ll,
\rightarrow opt[i + 1][j]); ++k) {
```

```
if (dp[i][j] >= dp[i][k] + dp[k + 1][j] +
    (pref[j + 1] - pref[i])) {
        dp[i][j] = dp[i][k]' + dp[k + 1][j] +
   (pref[j + 1] - pref[i]);
        opt[i][j] = k;
cout << dp[0][n - 1] << "\n";
## Lichao Tree
const int N = int(5e4 + 2);
const ll INF = ll(1e17);
vector<vector<ll> > tree(4*N, {0, INF});
ll f(vector<ll> line, int x){
return line[0] * x + line[1];
void insert(vector<ll> line, int lo = 1, int hi =
\rightarrow N, int i = 1){
 int m = (lo + hi) / 2;
bool left = f(line, lo) < f(tree[i], lo);</pre>
  bool mid = f(line, m) < f(tree[i], m);</pre>
  if(mid) swap(tree[i], line);
  if(hi - lo == 1) return;
  else if(left != mid) insert(line, lo, m, 2*i);
  else insert(line, m, hi, 2*i+1);
ll query(int x, int lo = 1, int hi = N, int i = 1){
  int m = (lo+hi)/2;
  \overline{ll} curr \stackrel{\cdot}{=} f(tree[i], x);
  if(hi-lo==1) return curr;
  if(x<m) return min(curr, query(x, lo, m, 2*i));</pre>
  else return min(curr, query(x, m, hi, 2*i+1));
```

18 DSU_ON_TREE

```
void dfs(int u, int p) {
  node[tt] = u;
  tin[u] = tt++, sz[u] = 1, hc[u] = -1;
  for (auto v: adj[u]) {
   if (v != p) {
      dfs(v, u);
      sz[u] += sz[v];
      if (hc[u] == -1 \text{ or } sz[hc[u]] < sz[v]) {
        hc[u] = v;
  tout[u] = tt - 1;
void dsu(int u, int p, int keep) {
  for (int v: adj[u]) {
   if (v != p and v != hc[u]) {
      dsu(v, u, 0);
  if (hc[u] != -1) {
    dsu(hc[u], u, 1);
  for (auto v: adj[u]) {
   if (v != p and v != hc[u]) {
      for (int i = tin[v]; i <= tout[v]; ++i) {</pre>
        int w = node[i];
```

```
// get ans in case of ans is related to
simple path or pair
}
for (int i = tin[v]; i <= tout[v]; ++i) {
    int w = node[i];
    // Add contribution of node w
}
}
// Add contribution of node u
// get ans in case ans is related to subtree
if (!keep) {
    for (int i = tin[u]; i <= tout[u]; ++i) {
        int w = node[i];
        // remove contribution of node w
    }
// Data structure in initial state (empty contribution)
}
dfs(0, 0); dsu(0, 0, 0);</pre>
```

19 DS_TRICKS

```
## Max prefix query with insertion only
a1 < a2 < a3 < ... < and b1 < b2 < b3 < ... < bn
// query
auto it = dp.lower_bound(a);
if (it != dp.begin()) {
   mx = max(now, prev(it)->second);
}
// insert
it = dp.upper_bound(a);
if (it != dp.begin() and prev(it)->second >= b) {
   continue;
}
it = dp.insert(it, {a, b});
it->second = b;
while (next(it) != dp.end() and next(it)->second <=
   b) {
   dp.erase(next(it));
}</pre>
```

20 DYNAMIC_CONNECTIVITY

```
const int Q = 1e5+5;
vector<array<int, 2>> t[4 * 0];
vector<int> ans(0);
int q;
struct DSU {
 int n, comps;
  vector<int> par, rnk;
  stack<array<int, 4>> ops;
 DSU(){}
  DSU(int n): n(n), comps(n), par(n), rnk(n) {
    iota(par.begin(), par.end(), 0);
  int find(int u) {
    return (par[u] == u)? u: find(par[u]);
  bool unite(int u, int v)
    u = find(u), v = find(v);
    if (u == v) return false;
    comps - -;
```

```
if (rnk[u] > rnk[v]) swap(u, v);
    ops.push({u, rnk[u], v, rnk[v]});
    par[u] = v;
    if (rnk[u] == rnk[v]) rnk[v]++;
    return true;
 void rollback() 
   if (ops.empty()) return ;
    auto [u, rnku, v, rnkv] = ops.top(); ops.pop();
    par[u] = u, rnk[u] = rnku;
    par[v] = v, rnk[v] = rnkv;
    comps++;
} dsu;
void add(int l, int r, array<int, 2> ed, int u = 1,
\rightarrow int s = 0, int e = q) {
 if (r < s or e < l) return ;</pre>
 if (l <= s and e <= r) {
   t[u].push back(ed);
    return :
 int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
 add(l, r, ed, v, s, m);
 add(l, r, ed, w, m + 1, e);
void qo(int u = 1, int s = 0, int e = q) {
 int rmv = 0;
 for (auto &ed: t[u]) rmv += dsu.unite(ed[0],
\rightarrow ed[1]);
 if (s == e) ans[s] = dsu.comps;
 else {
    int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
    go(v, s, m);
    go(w, m + 1, e);
 while (rmv--) dsu.rollback();
```

21 EULER_WALK

```
## Directed graph
vector<int> euler cycle(vector<int> *adj, int s =
→ 0) {
 vector<int> cycle;
 function<void(int)> dfs = [\&] (int u) {
   while (!adj[u].empty()) {
     int v = adi[u].back();
     adj[u].pop back();
     dfs(v);
   cycle.push back(u);
 dfs(s);
 reverse(cycle.begin(), cycle.end());
 return cycle;
## Undirected graph
vector<int> euler cycle(vector<int> *adj,
   vector<int> *des idx, vector<int> *done, int s
\rightarrow = 0) {
 vector<int> cycle;
```

```
function<void(int)> dfs = [\&] (int u) {
    while (!adj[u].empty()) {
      int i = adj[u].size() - 1;
      if (done[u][i]) {
        adj[u].pop back();
        continue;
      int v = adj[u][i];
      adj[u].pop back();
      done[u][i] = 1;
      done[v][des idx[u][i]] = 1;
      dfs(v);
    cycle.push back(u);
  dfs(s);
  return cycle;
int n, m; cin >> n >> m;
vector<int> adj[n], des idx[n], done[n];
vector<int> deg(n);
for (int e = 0; e < m; ++e) {
  int u, v; cin >> u >> v; u--, v--;
  des idx[u].push back(adj[v].size());
  des_idx[v].push_back(adj[u].size());
  adj[u].push back(v);
  adj[v].push_back(u);
  done[u] push back(0);
  done[v].push_back(0);
  deg[u]++, de\overline{g}[v]++;
for (int u = 0; u < n; ++u) {
  if (deg[u] & 1)
    cout << "IMPOSSIBLE\n";</pre>
    return ;
vector<int> cycle = euler cycle(adj, des idx, done,
if (cvcle.size() != m + 1) {
  cout << "IMPOSSIBLE\n";</pre>
  return ;
```

22 GEOMETRY

```
const int N = 3e5 + 9;
const double inf = 1e100;
const double eps = 1e-9;
const double PI = acos((double)-1.0);
int sign(double x) { return (x > eps) - (x < -eps);
    }
struct PT {
    double x, y;
    PT() { x = 0, y = 0; }
    PT(double x, double y) : x(x), y(y) {}
    PT(const PT &p) : x(p.x), y(p.y) {}
    PT operator + (const PT &a) const { return PT(x - + a.x, y + a.y); }</pre>
```

```
PT operator - (const PT &a) const { return PT(x
    - a.x, y - a.y); }
    PT operator * (const double a) const { return
    PT(x * a, y * a); }
    friend PT operator * (const double &a, const PT
    &b) { return PT(a * b.x, a * b.y); }
    PT operator / (const double a) const { return
    PT(x / a, y / a); }
    bool operator == (PT a) const { return sign(a.x
    -x) == 0 \&\& sign(a.y - y) == 0; }
    bool operator != (PT a) const { return !(*this
    == a); }
    bool operator < (PT a) const { return sign(a.x</pre>
    -x) == 0 ? y < a.y : x < a.x; }
    bool operator > (PT a) const { return sign(a.x
   - x) == 0 ? y > a.y : x > a.x; }
double norm() { return sqrt(x * x + y * y); }
    double norm2() { return x * x + y * y; }
    PT perp() { return PT(-y, x); }
double arg() { return atan2(y, x); }
    PT truncate(double r) { // returns a vector
    with norm r and having same direction
        double k = norm();
        if (!sign(k)) return *this;
        r /= k:
        return PT(x * r, y * r);
→ a.y * b.y; }
\rightarrow a - b): }
inline double dist(PT a, PT b) { return sqrt(dot(a
 \rightarrow - b, a - b)); }
\rightarrow - a.y * b.x; }
inline double cross2(PT a, PT b, PT c) { return
 inline int orientation(PT a, PT b, PT c) { return
 \rightarrow sign(cross(b - a, c - a)); }
PT perp(PT a) { return PT(-a.y, a.x); }
PT rotateccw90(PT a) { return PT(-a.y, a.x); }
PT rotatecw90(PT a) { return PT(a.y, -a.x); }
PT rotateccw(PT a, double t) { return PT(a.x *
    cos(t) - a.v * sin(t), a.x * sin(t) + a.v *
   cos(t)); }
PT rotatecw(PT a, double t) { return PT(a.x *
    cos(t) + a.y * sin(t), -a.x * sin(t) + a.y *
    cos(t)):
double SQ(double x) { return x * x; }
double rad to deg(double r) { return (r * 180.0 /
 → PI); }
double deg to rad(double d) { return (d * PI /
 -180.0); }
|<mark>double</mark> get angle(PT a, PT b) {
    double costheta = dot(a, b) / a.norm() /
    return acos(max((double)-1.0, min((double)1.0,
    costheta))):
|bool is point in angle(PT b, PT a, PT c, PT p) { //
 → does point p lie in angle <bac
    assert(orientation(a, b, c) != 0);
```

```
if (orientation(a, c, b) < 0) swap(b, c);
    return orientation(a, c, p) >= 0 \& \&
   orientation(a, b, p) <= 0:
bool half(PT p) {
    return p.y > 0.0 \mid \mid (p.y == 0.0 \&\& p.x < 0.0);
void polar sort(vector<PT> &v) { // sort points in

→ counterclockwise

    sort(v.begin(), v.end(), [](PT a,PT b) {
   return make_tuple(half(a), 0.0, a.norm2())
   < make tuple(half(b), cross(a, b), b.norm2());
    });
struct line {
    PT a, b; // goes through points a and b
    PT v; double c; //line form: direction vec
    [cross](x, v) = c
    line() {}
    //direction vector v and offset c
line(PT v, double c) : v(v), c(c) {
        auto p = get points();
        a = p.first; b = p.second;
 // equation ax + by + c = 0
line(double a, double b, double c) : v({ b,
   - a}), c(- c) {
  auto p = get points();
        a = p.first; b = p.second;
 // goes through points p and g
line(PT p, PT q): v(q - p), c(cross(v, p)), a(p),
pair<PT, PT> get points() { //extract any two
\rightarrow points from this line
  PT p, q; double a = -v.y, b = v.x; // ax + by = -c
  if (sign(a) == 0) {
            p = PT(0, -c / b);
            q = PT(1, -c / b);
        else if (sign(b) == 0) {
            p = PT(-c / a, 0);
            q = PT(-c / a, 1);
        else {
            p = PT(0, -c / b);
            q = PT(1, (-c - a) / b);
        return {p, q};
    //ax + by + c = 0
    array<double, 3> get abc() {
        double a = -v.y, b = v.x;
        return {a, b, ĉ};
    // 1 if on the left, -1 if on the right, 0 if
→ on the line
    int side(PT p) { return sign(cross(v, p) - c); }
    // line that is perpendicular to this and goes
    through point p
    line perpendicular through(PT p) { return {p, p
    + perp(v)}; }
    // translate the line by vector t i.e. shifting
→ it by vector t
```

```
line translate(PT t) { return {v, c + cross(v,

    t)}; }

        // compare two points by their orthogonal
 → projection on this line
        // a projection point comes before another if
 → it comes first according to vector v
        bool cmp by projection(PT p, PT q) { return
 \rightarrow dot(v, p) < dot(v, q); }
  line shift left(double d)
    PT z = v.\overline{p}erp().truncate(d);
    return line(a + z, b + z);
// find a point from a through b with distance d
PT point along line(PT a, PT b, double d) {
        return a + (((b - a) / (b - a).norm()) * d);
// projection point c onto line through a and b
 → assuming a != b
PT project from point to line(PT a, PT b, PT c) {
        return a + (b - a) * dot(c - a, b - a) / (b -
 \rightarrow a).norm2():
// reflection point c onto line through a and b

→ assuming a != b

PT reflection from point to line(PT a, PT b, PT c) {
        PT p = project from point to line(a,b,c);
        return point along line(c, p, 2.0 * dist(c, p));
// minimum distance from point c to line through a
double dist from point to line(PT a, PT b, PT c) {
        return fabs(cross(b - a, c - a) / (b -
 \rightarrow a).norm());
}
// returns true if point p is on line segment ab
bool is point on seq(PT a, PT b, PT p) {
        if (fabs(cross(p - b, a - b)) < eps) {
                if (p.x < min(a.x, b.x) | | p.x > max(a.x, b.x) | 
 if (p.y < min(a.y, b.y) \mid \mid p.y > max(a.y, b.y)
 → b.y)) return false;
                return true;
        return false;
// minimum distance point from point c to segment
→ ab that lies on segment ab
PT project from point to seg(PT a, PT b, PT c) {
        double r = dist2(a, b);
        if (fabs(r) < eps) return a;</pre>
        r = dot(c - a, b - a) / r;
        if (r < 0) return a;</pre>
        if (r > 1) return b:
        return a + (b - a) * r;
// minimum distance from point c to segment ab
double dist from point to seg(PT a, PT b, PT c) {
        return dist(c, project from point to seg(a, b,

→ c)):
// 0 if not parallel, 1 if parallel, 2 if collinear
bool is parallel(PT a, PT b, PT c, PT d) {
        double k = fabs(cross(b - a, d - c));
        if (k < eps){
```

```
if (fabs(cross(a - b, a - c)) < eps &&
    fabs(cross(c - d, c - a)) < eps) return 2;
         else return 1:
     else return 0;
 // check if two lines are same
|bool are lines same(PT a, PT b, PT c, PT d) {
     if (fabs(cross(a - c, c - d)) < eps &&
    fabs(cross(b - c, c - d)) < eps) return true;
     return false:
 // bisector vector of <abc</pre>
PT angle bisector(PT &a, PT &b, PT &c){
     PT p = a - b, q = c - b;
     return p + q * sqrt(dot(p, p) / dot(q, q));
 // 1 if point is ccw to the line, 2 if point is cw
   to the line, 3 if point is on the line
int point line relation(PT a, PT b, PT p) {
     int c = sign(cross(p - a, b - a));
     if (c < 0) return 1;
     if (c > 0) return 2:
     return 3;
// intersection point between ab and cd assuming
   unique intersection exists
|bool line line intersection(PT a, PT b, PT c, PT d,
    PT &ans) {
     double a1 = a.y - b.y, b1 = b.x - a.x, c1 = a.y
    cross(a, b);
     double a2 = c.y - d.y, b2 = d.x - c.x, c2 = d.x
    cross(c, d);
     double det = a1 * b2 - a2 * b1;
     if (det == 0) return 0;
     ans = PT((b1 * c2 - b2 * c1) / det, (c1 * a2 -
    a1 * c2) / det);
     return 1;
 // intersection point between segment ab and
    segment cd assuming unique intersection exists
 bool seg seg intersection(PT a, PT b, PT c, PT d,
 → PT &ans) {
     double oa = cross2(c, d, a), ob = cross2(c, d,
     double oc = cross2(a, b, c), od = cross2(a, b,
    d);
     if (oa * ob < 0 && oc * od < 0){
         ans = (a * ob - b * oa) / (ob - oa);
         return 1;
     else return 0;
 // intersection point between segment ab and
    segment cd assuming unique intersection may not
   exists
 // se.size()==0 means no intersection
 // se.size()==1 means one intersection
// se.size()==2 means range intersection
 set<PT> seg seg intersection inside(PT a, PT b,
    PT c, \overline{PT d}
     PT ans;
     if (seg seg intersection(a, b, c, d, ans))
    return {ans}:
```

```
set<PT> se;
    if (is point on seg(c, d, a)) se.insert(a);
    if (is point on seg(c, d, b)) se.insert(b);
    if (is point on seg(a, b, c)) se.insert(c);
    if (is point on seg(a, b, d)) se.insert(d);
    return se;
// intersection between segment ab and line cd
// 0 if do not intersect, 1 if proper intersect, 2

→ if segment intersect

int seq line relation(PT a, PT b, PT c, PT d) {
    double p = cross2(c, d, a);
    double q = cross2(c, d, b);
    if (sign(p) == 0 \&\& sign(q) == 0) return 2;
    else if (p * q < 0) return 1;
    else return 0:
// intersection between segament ab and line cd

→ assuming unique intersection exists

bool seg line intersection(PT a, PT b, PT c, PT d,
 → PT &ans) {
    bool k = seg_line_relation(a, b, c, d);
assert(k != 2);
    if (k) line line intersection(a, b, c, d, ans);
    return k;
// minimum distance from segment ab to segment cd
double dist from seg to seg(PT a, PT b, PT c, PT d)
    PT dummy;
    if (seg seg intersection(a, b, c, d, dummy))
   return 0.0:
    else return min({dist from point to seg(a, b,
   c), dist from point to seq(a, b, d),
        dist from point to seg(c, d, a),
    dist from point to seq(c, d, b)});
// minimum distance from point c to ray (starting
   point a and direction vector b)
double dist from point to ray(PT a, PT b, PT c) {
    b = a + b;
    double r = dot(c - a, b - a);
    if (r < 0.0) return dist(c, a);
    return dist from point to line(a, b, c);
// starting point as and direction vector ad
bool ray ray intersection(PT as, PT ad, PT bs, PT
    double dx = bs.x - as.x, dy = bs.y - as.y;
    double det = bd.x * ad.y - bd.y * ad.x;
    if (fabs(det) < eps) return 0;</pre>
    double u = (dy * bd.x - dx * bd.y) / det;
    double v = (dy * ad.x - dx * ad.y) / det;
    if (sign(u) >= 0 \&\& sign(v) >= 0) return 1;
    else return 0:
double ray ray distance(PT as, PT ad, PT bs, PT bd)
    if (ray_ray_intersection(as, ad, bs, bd))
 \rightarrow return \overline{0}.0;
    double ans = dist from point to ray(as, ad, bs);
    ans = min(ans, dist from point to ray(bs, bd,
    as));
```

```
return ans;
struct circle {
    PT p; double r;
   circle() {}
circle(PT _p, double _r): p(_p), r(_r) {};
// center (x, y) and radius r
    circle(double x, double y, double r): p(PT(x,
- y)), r(_r) {};
// circumcircle of a triangle
    // the three points must be unique
    circle(PT a, PT b, PT c) {
        b = (a + b) * 0.5;
        c = (a + c) * 0.5;
        line line intersection(b, b + rotatecw90(a
   - b), c, c + rotatecw90(a - c), p);
        r = dist(a, p);
    // inscribed circle of a triangle
    circle(PT a, PT b, PT c, bool t) {
        line u, v;
        double m = atan2(b.y - a.y, b.x - a.x), n =
 → atan2(c.y - a.y, c.x - a.x);
        u.b = u.a + (PT(cos((n + m)/2.0), sin((n +
 \rightarrow m)/2.0)));
        v.a = b;
        m = atan2(a.y - b.y, a.x - b.x), n =
 → atan2(c.y - b.y, c.x - b.x);
        v.b = v.a + (PT(cos((n + m)/2.0), sin((n +
 \rightarrow m)/2.0)));
        line line intersection(u.a, u.b, v.a, v.b,
        r = dist from point to seg(a, b, p);
    bool operator == (circle v) { return p == v.p
\rightarrow && sign(r - v.r) == 0:
    double area() { return PI * r * r; }
    double circumference() { return 2.0 * PI * r; }
//O if outside, 1 if on circumference, 2 if inside
int circle point relation(PT p, double r, PT b) {
    double d = dist(p, b);
    if (sign(d - r) < 0) return 2;
    if (sign(d - r) == 0) return 1;
    return 0;
// 0 if outside, 1 if on circumference, 2 if inside
int circle line relation(PT p, double r, PT a, PT
→ b) {
    double d = dist from point to line(a, b, p);
    if (sign(d - r) < 0) return 2;
    if (sign(d - r) == 0) return 1;
    return 0;
//compute intersection of line through points a and
//circle centered at c with radius r > 0
vector<PT> circle line intersection(PT c, double r,
→ PT a. PT b) {
    vector<PT> ret;
    b = b - a; a = a - c;
    double A = dot(b, b), B = dot(a, b);
```

```
double C = dot(a, a) - r * r, D = B * B - A * C;
    if (D < -eps) return ret;</pre>
    ret.push back(c + a + b * (-B + sqrt(D + eps))
    / A):
    if (D > eps) ret.push back(c + a + b * (-B - a))
    sqrt(D)) / A);
    return ret;
//5 - outside and do not intersect
//4 - intersect outside in one point
//3 - intersect in 2 points
//2 - intersect inside in one point
//1 - inside and do not intersect
int circle circle relation(PT a, double r, PT b,
   double R) {
    double d = dist(a, b);
    if (sign(d - r - R) > 0) return 5;
    if (sign(d - r - R) == 0) return 4;
    double l = fabs(r - R);
    if (sign(d - r - R) < 0) \&\& sign(d - l) > 0)
    return 3;
    if (sian(d - l) == 0) return 2:
    if (sign(d - l) < 0) return 1;
    assert(0); return -1;
vector<PT> circle circle intersection(PT a, double
    r, PT b, double R) {
    if (a == b \&\& sign(r - R) == 0) return
    {PT(1e18, 1e18)};
    vector<PT> ret;
    double d = sqrt(dist2(a, b));
    if (d > r + R \mid | d + min(r, R) < max(r, R))
    return ret;
    double x = (d * d - R * R + r * r) / (2 * d);
    double y = sart(r * r - x * x);
    PT v = (b - a) / d;
    ret.push back(a + v * x + rotateccw90(v) * y);
    if (y > \overline{0}) ret.push back(a + v * x -
    rotateccw90(v) * v);
    return ret;
// returns two circle c1, c2 through points a, b
   and of radius r
// 0 if there is no such circle, 1 if one circle, 2
→ if two circle
int get circle(PT a, PT b, double r, circle &c1,

    circle &c2) {

    vector<PT> v = circle circle intersection(a, r,
    b, r);
    int t = v.size();
    if (!t) return 0;
    c1.p = v[0], c1.r = r;
    if (t == 2) c2.p = v[1], c2.r = r;
    return t;
// returns two circle c1, c2 which is tangent to
→ line u, goes through
// point q<sup>'</sup>and has radiŭs r1; 0 for no circle, 1 if
- c1 = c2 , 2 if <math>c1 != c2
int get circle(line u, PT g, double r1, circle &c1,

    circle &c2) {
    if (sign(d - r1^{-*} 2.\overline{0}) > 0) return 0;
    if (sign(d) == 0) {
```

```
cout << u.v.x << ' ' << u.v.y << '\n';
                                                             c1.p = q + rotateccw90(u.v).truncate(r1);
                                                             c2.p = q + rotatecw90(u.v).truncate(r1);
                                                             c1.r = c2.r = r1;
                                                             return 2:
                                                         line u1 = line(u.a +
                                                        rotateccw90(u.v).truncate(r1), u.b +
                                                        rotateccw90(u.v).truncate(r1));
                                                        line u2 = line(u.a +
                                                        rotatecw90(u.v).truncate(r1), u.b +
                                                        rotatecw90(u.v).truncate(r1));
                                                        circle cc = circle(q, r1);
                                                        PT p1, p2; vector<PT> v;
                                                        v = circle line intersection(q, r1, u1.a, u1.b);
                                                        if (|v.size()\rangle) \overline{v} = circle line intersection(q,
                                                        r1, u2.a, u2.b);
                                                        v.push back(v[0]);
                                                        p1 = v[0], p2 = v[1];
                                                        c1 = circle(p1, r1);
                                                        if (p1 == p2) {
                                                             c2 = c1:
                                                             return 1:
                                                        c2 = circle(p2, r1);
                                                        return 2;
                                                     // returns area of intersection between two circles
                                                    double circle circle area(PT a, double r1, PT b,
                                                     → double r2) {
                                                        double d = (a - b).norm();
                                                        if(r1 + r2 < d + eps) return 0;
                                                        if(r1 + d < r2 + eps) return PI * r1 * r1;
                                                        if(r2 + d < r1 + eps) return PI * r2 * r2;
                                                        double theta 1 = acos((r1 * r1 + d * d - r2 *
                                                       r2) / (2 * r1 * d)),
                                                         theta 2 = acos((r2 * r2 + d * d - r1 * r1)/(2
                                                       * r2 * d));
return r1 * r1 * (theta_1 - sin(2 *
                                                        theta 1)/2.) + r2 * r2 \overline{*} (theta 2 - \sin(2 *)
                                                        theta 2)/2.);
                                                    // tangent lines from point q to the circle
                                                    int tangent lines from point(PT p, double r, PT q,

→ line &u, line &v) -

                                                        int x = sign(dist2(p, q) - r * r);
                                                        if (x < 0) return 0; // point in cricle</pre>
                                                        if (x == 0) { // point on circle
                                                             u = line(q, q + rotateccw90(q - p));
                                                             \ddot{v} = \ddot{u};
                                                             return 1:
                                                        double d = dist(p, q);
                                                        double l = r * r / d;
                                                        double h = sqrt(r * r - l * l);
                                                        u = line(q, p + ((q - p).truncate(l) +
                                                        (rotateccw90(q - p).truncate(h))));
                                                        \dot{v} = line(q, \dot{p} + (\dot{q} - p).truncate(\dot{l}) +
                                                        (rotatecw90(g - p).truncate(h))));
                                                        return 2;
                                                     / returns outer tangents line of two circles
double d = dist from_point_to_line(u.a, u.b, q); // if inner == 1 it returns inner tangent lines
```

```
int tangents lines from circle(PT c1, double r1, PT
if (inner) r2 = -r2;
    PT d = c2 - c1;
    double dr = r1 - r2, d2 = d.norm(), h2 = d2 - rac{1}{2}
→ dr * dr:
    if (d2 == 0 | | h2 < 0) {
        assert(h2'!=0);
        return 0:
    vector<pair<PT, PT>>out;
    for (int tmp: {- 1, 1}) {
        PT v = (d * dr + rotateccw90(d) * sqrt(h2)
   * tmp) / d2;
        out.push back(\{c1 + v * r1, c2 + v * r2\});
    u = line(out[0].first, out[0].second);
    if (out.size() == 2) v = line(out[1].first,
  out[1].second);
    return 1 + (h2 > 0);
1/0(n^2 \log n)
struct CircleUnion {
    int n;
    double x[2020], y[2020], r[2020];
    int covered[2020];
    vector<pair<double, double> > seq, cover;
    double arc, pol;
    inline int sign(double x) {return x < -eps ? -1</pre>
   : x > eps;}
    inline int sign(double x, double y) {return
   sign(x - y);
    inline double SQ(const double x) {return x * x;}
    inline double dist(double x1, double y1, double
   x2, double y2) {return sqrt(SQ(x1 - x2) + SQ(y1
   - y2));}
    inline double angle(double A, double B, double
        double val = (SQ(A) + SQ(B) - SQ(C)) / (2 *
        if (val < -1) val = -1;
        if (val > +1) val = +1;
        return acos(val);
    CircleUnion() {
        seq.clear(), cover.clear();
        arc = pol = 0;
   void init() {
        n = 0;
        seg.clear(), cover.clear();
        arc = pol = 0;
   void add(double xx, double yy, double rr) {
        x[n] = xx, y[n] = yy, r[n] = rr, covered[n]
\rightarrow = 0, n++;
    void getarea(int i, double lef, double rig) {
        arc += 0.5 * r[i] * r[i] * (rig - lef -
   sin(rig - lef));
        double x1 = x[i] + r[i] * cos(lef), y1 =
   v[i] + r[i] * sin(lef);
        double x2 = x[i] + r[i] * cos(rig), y2 =
\rightarrow v[i] + r[i] * sin(rig);
```

```
pol += x1 * y2 - x2 * y1;
   double solve() {
       for (int i = 0; i < n; i++) {
            for (int j = 0; j < i; j++) {
                if (!sign(x[i] - x[j]) &&
  !sign(y[i] - y[j]) && !sign(r[i] - r[j])) {
                    r[i] = 0.0;
                    break;
       for (int i = 0; i < n; i++) {
            for (int j = 0; j < n; j++) {
                if (i!=j \&\& sign(r[j]-r[i]) >=
   0 && sign(dist(x[i], y[i], x[j], y[j]) - (r[j]
  - r[i])) <= 0) {
                    covered[i] = 1:
                    break;
       for (int i = 0; i < n; i++) {
           if (sign(r[i]) && !covered[i]) {
                seg.clear();
                for (int j = 0; j < n; j++) {
                    if (i != j) {
                        double d = dist(x[i], y[i],
\rightarrow x[j], y[j]);
                        if (sign(d - (r[j] + r[i]))
\Rightarrow >= 0 || sign(d - abs(r[j] - r[i])) <= 0) {
                             continue:
                        double alpha = atan2(y[j] -
\rightarrow y[i], x[j] - x[i]);
                        double beta = angle(r[i],
\rightarrow d, r[j]);
                        pair<double, double>
  tmp(alpha - beta, alpha + beta);
                        if (sign(tmp.first) \leq 0 &&
  sign(tmp.second) \ll 0) {
   seq.push back(pair<double, double>(2 * PI +
   tmp.first, 2 * PI + tmp.second));
                        else if (sign(tmp.first) <</pre>
  0) {
   seq.push back(pair<double, double>(2 * PI +
   tmp.firs\overline{t}, 2 * PI));
   seg.push back(pair<double, double>(0,
  tmp.second));
                        else {
                             seq.push back(tmp);
                sort(seq.begin(), seq.end());
                double rig = 0;
                for (vector<pair<double, double>
   >::iterator iter = seq.begin(); iter !=

    seq.end(); iter++) {
```

```
if (sign(rig - iter->first) >=
→ 0) {
                        rig = max(rig,

    iter->second);

                    élse {
                        getarea(i, rig,

    iter->first);

                        rig = iter->second;
                if (!sign(rig)) {
                    arc += r[i] * r[i] * PI;
                else {
                    getarea(i, rig, 2 * PI);
        return pol / 2.0 + arc;
} CU:
double area of triangle(PT a, PT b, PT c) {
    return fabs(cross(b - a, c - a) * 0.5);
// -1 if strictly inside, 0 if on the polygon, 1 if

→ strictly outside

int is point in triangle(PT a, PT b, PT c, PT p) {
    if (sign(cross(b - a,c - a)) < 0) swap(b, c);
    int c1 = sign(cross(b - a,p - a));
    int c2 = sign(cross(c - b, p - b));
    int c3 = sign(cross(a - c, p - c));
    if (c1<0 || c2<0 || c3 < 0) return 1;
    if (c1 + c2 + c3 = 3) return 0:
    return -1;
double perimeter(vector<PT> &p) {
    double ans=0; int n = p.size();
    for (int i = 0; i < n; i++) ans += dist(p[i],
    p[(i + 1) % n]);
    return ans;
double area(vector<PT> &p) {
    double ans = 0; int n = p.size();
    for (int i = 0; i < n; i++) ans += cross(p[i],
    p[(i + 1) % n]);
    return fabs(ans) * 0.5;
// centroid of a (possibly non-convex) polygon,
// assuming that the coordinates are listed in a
   clockwise or
// counterclockwise fashion. Note that the
    centroid is often known as
  the "center of gravity" or "center of mass".
PT centroid(vector<PT> &p) {
    int n = p.size(); PT c(0, 0);
    double sum = 0;
    for (int i = 0; i < n; i++) sum += cross(p[i],
    p[(i + 1) % n]);
    double scale = 3.0 * sum;
    for (int i = 0; i < n; i++) {
        int j = (i + 1) \% n;
        c = c + (p[i] + p[j]) * cross(p[i], p[j]);
```

```
return c / scale;
// 0 if cw, 1 if ccw
bool get direction(vector<PT> &p) {
    double ans = 0; int n = p.size();
    for (int i = 0; i < n; i++) ans += cross(p[i],
   p[(i + 1) % n]);
    if (sign(ans) > 0) return 1;
    return 0;
// it returns a point such that the sum of distances
// from that point to all points in p is minimum
// O(n log^2 MX)
PT geometric_median(vector<PT> p) {
auto tot dist = [\&](PT z) {
     double res = 0;
     for (int i = 0; i < p.size(); i++) res +=
\rightarrow dist(p[i], z);
     return res;
};
auto findY = [\&] (double x)
     double yl = -1e5, yr = 1e5;
     for (int i = 0; i < 60; i++) {
         double ym1 = yl + (yr - yl) / 3;
         double ym2 = yr - (yr - yl) / 3;
         double d1 = tot dist(PT(x, ym1));
         double d2 = tot dist(PT(x, ym2));
         if (d1 < d2) yr = ym2;
         else yl = ym1;
     return pair<double, double> (yl,
   tot dist(PT(x, yl)));
    double xl = -1e5, xr = 1e5;
    for (int i = 0; i < 60; i++) {
        double xm1 = xl + (xr - xl) / 3;
        double xm2 = xr - (xr - xl) / 3;
        double y1, d1, y2, d2;
        auto z = findY(xm1); y1 = z.first; d1 =
        z = findY(xm2); y2 = z.first; d2 = z.second;
        if (d1 < d2) xr = xm2;
        else xl = xm1;
    return {xl, findY(xl).first };
vector<PT> convex hull(vector<PT> &p) {
if (p.size() \le \overline{1}) return p:
vector < PT > v = p;
    sort(v.begin(), v.end());
    vector<PT> up, dn;
    for (auto\& p : v) {
        while (up.size() > 1 \&\&
   orientation(up[up.size() - 2], up.back(), p) >=
up.pop back();
        while (dn.size() > 1 \&\&
    orientation(dn[dn.size() - 2], dn.back(), p) <=
            dn.pop back();
        up.push back(p);
        dn.push back(p);
```

```
if (v.size() > 1) v.pop back();
    reverse(up.begin(), up.end());
    up.pop back();
    for (\overline{auto}\& p : up) {
        v.push back(p);
    if (v.size() == 2 \&\& v[0] == v[1]) v.pop back();
    return v;
//checks if convex or not
bool is convex(vector<PT> &p) -
    bool s[3]; s[0] = s[1] = s[2] = 0;
    int n = p.size();
    for (int i = 0; i < n; i++) {
        int j = (i + 1) % n;
        int k = (j + 1) \% n;
        s[sign(cross(p[j] - p[i], p[k] - p[i])) +
   1] = 1;
        if (s[0] && s[2]) return 0;
    return 1:
// -1 if strictly inside, 0 if on the polygon, 1 if
   strictly outside
// it must be strictly convex, otherwise make it

→ strictly convex first

int is point in convex(vector<PT> &p, const PT& x)
 \rightarrow { \overline{7}/\ 0(\log n)
    int n = p.size(); assert(n >= 3);
    int a = orientation(p[0], p[1], x), b =
   orientation(p[0], p[n - 1], x);
    if (a < 0 | | b > 0) return 1;
    int l = 1, r = n - 1;
    while (l + 1 < r) {
        int mid = l + r \gg 1;
        if (orientation(p[0], p[mid], x) >= 0) l =
   mid;
        else r = mid;
    int k = orientation(p[l], p[r], x);
    if (k <= 0) return -k:
    if (l == 1 \&\& a == 0) return 0;
    if (r == n - 1 \&\& b == 0) return 0:
    return -1;
bool is point on polygon(vector<PT> &p, const PT&
int n = p.size();
    for (int i = 0; i < n; i++) {
     if (is point on seg(p[i], p[(i + 1) % n], z))
   return 1;
    return 0;
// returns 1e9 if the point is on the polygon
int winding number(vector<PT> &p, const PT& z) { //
    if (is point on polygon(p, z)) return 1e9;
    int n = p.size(), ans = 0;
    for (int i = 0; i < n; ++i) {
        int j = (i + 1) % n;
        bool below = p[i].y < z.y;
```

```
if (below != (p[j].y < z.y)) {
             auto orient = orientation(z, p[j],
\rightarrow p[i]);
            if (orient == 0) return 0;
            if (below == (orient > 0)) ans += below

→ ? 1 : -1:

    return ans;
// -1 if strictly inside, 0 if on the polygon, 1 if

→ strictly outside

int is point in polygon(vector<PT> &p, const PT& z)
\rightarrow { // 0(n)
    int k = winding_number(p, z);
    return k == 1e9^{-}? 0 : k == 0 ? 1 : -1;
// id of the vertex having maximum dot product with
// polygon must need to be convex
// top - upper right vertex
// for minimum dot prouct negate z and return
\rightarrow -dot(z, p[id])
int extreme vertex(vector<PT> &p, const PT &z,
    const int top) \{ // O(\log n) \}
    int n = p.size();
    if (n == 1) return 0;
 double ans = dot(p[0], z); int id = 0;
    if (dot(p[top], z) > ans) ans = dot(p[top], z),
\rightarrow id = top;
    int l = 1, r = top - 1;
    while (l < r) {
        int mid = l + r >> 1:
        if (dot(p[mid + 1], z) >= dot(p[mid], z))
   = mid + 1:
        else r = mid:
    if (dot(p[l], z) > ans) ans = dot(p[l], z), id
l = top + 1, r = n - 1;
    while (l < r) {
        int mid = l + r >> 1;
        if (dot(p[(mid + 1) % n], z) >= dot(p[mid],
\rightarrow z)) l = mid + 1;
        else r = mid;
    ĺ %= n;
    if (dot(p[l], z) > ans) ans = dot(p[l], z), id
    return id;
double diameter(vector<PT> &p) {
    int n = (int)p.size();
    if (n == 1) return 0;
    if (n == 2) return dist(p[0], p[1]);
    double ans = 0;
    int i = 0, j = 1;
    while (i < n) {
        while (cross(p[(i + 1) % n] - p[i], p[(j +
\rightarrow 1) % n] - p[j]) >= 0) {
         ans = max(ans, dist2(p[i], p[j]));
        j = (j + 1) \% n;
```

```
ans = max(ans, dist2(p[i], p[i]));
    return sqrt(ans);
double width(vector<PT> &p) {
    int n = (int)p.size();
    if (n <= 2) return 0;
    double ans = inf;
    int i = 0, j = 1;
    while (i < n) {
        while (cross(p[(i + 1) % n] - p[i], p[(j +
\rightarrow 1) % n] - p[i]) >= 0) i = (i + 1) % n;
        ans = min(ans,
    dist from point to line(p[i], p[(i + 1) % n],
   p[j]));
        1++;
    return ans;
}
// minimum perimeter
double minimum enclosing rectangle(vector<PT> &p) {
int n = p.siz\overline{e}();
if (n <= 2) return perimeter(p);</pre>
int mndot = 0; double tmp = dot(p[1] - p[0], p[0]);
 for (int i = 1; i < n; i++) {
 if (dot(p[1] - p[0], p[i]) <= tmp) {</pre>
  tmp = dot(p[1] - p[0], p[i]);
  mndot = i;
double ans = inf;
int i = 0, j = 1, mxdot = 1;
while (i < n) {
 PT cur = p[(i + 1) \% n] - p[i];
        while (cross(cur, p[(j + 1) % n] - p[j]) >=
\rightarrow 0) j = (j + 1) % n;
        while (dot(p[(mxdot + 1) % n], cur) >=
\rightarrow dot(p[mxdot], cur)) mxdot = (mxdot + 1) % n;
        while (dot(p[(mndot + 1) % n], cur) <=
\rightarrow dot(p[mndot], cur)) mndot = (mndot + 1) % n;
        ans = min(ans, 2.0 * ((dot(p[mxdot], cur)))
    cur.norm() - dot(p[mndot], cur) / cur.norm()) +
   dist from point to line(p[i], p[(i + 1) % n],
   p[j])));
        i++;
    return ans;
// given n points, find the minimum enclosing
 // call convex hull() before this for faster
 // expected O(n)
circle minimum enclosing circle(vector<PT> &p) {
    random shuffle(p.begin(), p.end());
    int n = p.size();
    circle c(p[0], 0);
    for (int i = 1; i < n; i++) {
        if (sign(dist(c.p, p[i]) - c.r) > 0) {
    c = circle(p[i], 0);
            for (int j = 0; j < i; j++) {
                if (sign(dist(c.p, p[j]) - c.r) >
→ 0) {
```

```
dist(p[i], p[j]) / 2);
                      for (int k = 0; k < j; k++) {
                          if (sign(dist(c.p, p[k]) -
 \rightarrow c.r) > 0) {
                              c = circle(p[i], p[j],
 \rightarrow p[k]);
    return c;
## Closest Pair of Points
ll min dis(vector<array<int, 2>> &pts, int l, int
 ¬ r) {
  if (l + 1 >= r) return LLONG MAX;
  int m = (l + r) / 2;
  ll my = pts[m-1][1];
  ll d = min(min dis(pts, l, m), min dis(pts, m,
  inplace merge(pts.begin()+l, pts.begin()+m,

    pts.begin()+r);

  for (int i = l; i < r; ++i) {
  if ((pts[i][1] - my) * (pts[i][1] - my) < d) {</pre>
      pts[j][0]) * (pts[i][0] - pts[j][0]) < d; ++j) {
         [l] dx = pts[i][0] - pts[j][0], dy =
    pts[i][1] - pts[i][1];
         d = min(d, dx * dx + dy * dy);
  return d;
|vector<array<<mark>int</mark>, 2>> pts(n);
sort(pts.begin(), pts.end(), [&] (array<int, 2> a,
 \rightarrow array<int, 2> b){
  return make pair(a[1], a[0]) < make pair(b[1],</pre>
 \rightarrow b[0]);
});
## Angular Sort
inline bool up (point p) {
 return p.y > 0 or (p.y == 0 \text{ and } p.x >= 0);
sort(v.begin(), v.end(), [] (point a, point b) {
  return up(a) == up(b) ? a.x * b.y > a.y * b.x :
 \rightarrow up(a) < up(b);
});
## Convex Hull
struct pt {
  int x, y;
|ll cross(pt a, pt b, pt c) { //ab*ac
  return 111*(b.x-a.x)*(c.y-a.y) -
 \rightarrow 111*(c.x-a.x)*(b.y-a.y);
vector<pt> convexHull(vector<pt>& p) {
  sort(p.begin(), p.end(), [\&] (pt a, pt b) {
     return (a.x==b.x? a.y<b.y: a.x<b.x);
  int n = p.size(), m = 0;
```

c = circle((p[i] + p[j]) / 2,

23 GRAY_CODE

```
int gc(int n){ return n^(n>>1); }
int gc to dec(int g) {
  int d=0;
  while (g) { d ^= g; g >>= 1; }
  return d;
}
```

24 HASHING

```
// Hashing
// Hashvalue(l...r) = hsh[l] - hsh[r + 1] * base ^
\rightarrow (r - l + 1):
// Must call preprocess
#include<bits/stdc++.h>
using namespace std;
typedef long long li;
const int MAX = 100009;
ll mods[2] = {10000000007, 1000000009};
//Some back-up primes: 1072857881, 1066517951,
- 1040160883
ll\ bases[2] = \{137, 281\};
ll pwbase[3][MAX];
void Preprocess(){
  pwbase[0][0] = pwbase[1][0] = 1;
  for(ll i = 0; i < 2; i++){
    for(ll j = 1; j < MAX; j++)
       pwbase[i][j] = (pwbase[i][j - 1] * bases[i])
    % mods[i];
struct Hashing{
  ll hsh[2][MAX];
  string str;
  Hashing(){}
  Hashing(string _str) {str = _str; memset(hsh, 0,

    sizeof(hsh)); build();}

  void Build(){
    for(ll i = str.size() - 1; i >= 0; i--){
      for(int j = 0; j < 2; j++){
  hsh[j][i] = (hsh[j][i + 1] * bases[j] +</pre>

    str[i]) % mods[j];

         hsh[j][i] = (hsh[j][i] + mods[j]) % mods[j];
```

```
pair<ll,ll> GetHash(ll i, ll j){
    assert(i <= j);</pre>
    ll tmpl = (hsh[0][i] - (hsh[0][j + 1] *
    pwbase[0][j - i + 1]) % mods[0]) % mods[0];
    ll tmp2 = (hsh[1][i] - (hsh[1][j + 1] *
    pwbase[1][j - i + 1]) % mods[1]) % mods[1];
    if(tmp1 < 0) tmp1 += mods[0];
    if(tmp2 < 0) tmp2 += mods[1];
    return make pair(tmp1, tmp2);
};
/***
    * Everything is 0 based
* Call precal() once in the program
* Call update(1,0,n-1,i,j,val) to update the

→ value of position

      i to j to val, here n is the length of the
    strina
    * Call query(1,0,n-1,L,R) to get a node

→ containing hash

      of the position [L:R]
    * Before any update/query
         - Call init(str) where str is the string to

→ be hashed

         - Call build(1,0,n-1)
***/
namespace strhash {
  int n;
  const int MAX = 100010;
  int ara[MAX];
  const int MOD[] = {2078526727, 2117566807};
  const int BASE[] = {1572872831, 1971536491};
  int BP[2][MAX], CUM[2][MAX];
  void init(char *str) {
    n = strlen(str);
    for(int i=0;i<n;i++) ara[i] = str[i]-'0'+1; ///</pre>

→ scale str[i] if needed
  void precal() {
    BP[0][0] = BP[1][0] = 1;
    for(int i=1; i<MAX; i++)</pre>
      BP[0][i] = (BP[0][i-1] * (long long) BASE[0]
 → ) % MOD[0];
      BP[1][i]' = (BP[1][i-1] * (long long) BASE[1]
      % MOD[1]:
  struct node {
    int sz;
    int h[2];
    node() {}
  } tree[4*MAX];
  int lazy[4*MAX];
  inline node Merge(node a, node b) {
    node ret;
    ret.h[0] = ( (a.h[0] * (long long) BP[0][b.sz]
 \rightarrow ) + b.h[0] ) % MOD[0];
    ret.h[1] = ((a.h[1])*(long long) BP[1][b.sz]
\rightarrow ) + b.h[1] ) % MOD[1];
```

```
ret.sz = a.sz + b.sz;
    return ret;
  inline void build(int n,int st,int ed) {
    if(st==ed) {
      tree[n].h[0] = tree[n].h[1] = ara[st];
      tree[n].sz = 1;
      return;
    int mid = (st+ed)>>1;
    build(n+n,st,mid);
    build(n+n+1,mid+1,ed);
    tree[n] = Merge(tree[n+n], tree[n+n+1]);
  inline void update(int n,int st,int ed,int id,int
    if(st>id or ed<id) return;</pre>
    if(st==ed and ed==id) {
      tree[n].h[0] = tree[n].h[1] = v;
      return;
    int mid = (st+ed)>>1;
    update(n+n,st,mid,id,v);
    update(n+n+1,mid+1,ed,id,v);
    tree[n] = Merge(tree[n+n], tree[n+n+1]);
  inline node query(int n,int st,int ed,int i,int
    if(st>=i and ed<=j) return tree[n];</pre>
    int mid = (st+ed)/2;
    if(mid<i) return query(n+n+1,mid+1,ed,i,j);</pre>
    else if(mid>=j) return query(n+n,st,mid,i,j);
    else return Merge(query(n+n,st,mid,i,j),query(n)
    +n+1,mid+1,ed,i,j));
25 HLD
```

```
int tt, tin[N], tout[N], sz[N], par[N][LG], hvc[N];
void dfs(int u, int p) {
 tin[u] = tt++, sz[u] = 1, par[u][0] = p;
 for (int j = 1; j < LG; ++j)
    par[u][j] = par[par[u][j-1]][j-1];
  for (int &v: adj[u]) {
    if (v != p) {
      dfs(v, u);
      sz[u] += sz[v];
      if (sz[v] > mx) {
        mx = sz[v];
        hvc[u] = v;
 tout[u] = tt-1;
int ch cnt, idx cnt, chno[N], chd[N], idx[N];
void hld(int u, int p) {
 if(chd[ch cnt] == -1) {
    chd[ch cnt] = u;
```

```
chno[u] = ch cnt, idx[u] = idx cnt++;
 if(hvc[u] != -1) {
    hld(hvc[u], u);
 for (int &v: adj[u]) {
    if (v != p and v != hvc[u]) {
      ch cnt++:
      hld(v, u);
void ?node update(int u, int x) {
  ?update(idx[u], x);
void ?pupdate up(int u, int anc) {
 if (chno[u] == chno[anc]) {
    return ?rupdate(idx[anc], idx[u]);
  ?rupdate(idx[chd[chno[u]]], idx[u]);
  ?pupdate up(par[chd[chno[u]]][0], anc);
void ?pupdate(int u, int v) {
  int l = lca(u, v);
  ?pupdate up(u, l);
  ?pupdate up(v, l);
ĺl ?node query(int u) {
 return ?query(idx[u]);
int ?pquery up(int u, int anc) {
 if (chno[\overline{u}] == chno[anc]) {
    return ?rquery(idx[anc], idx[u]);
  return f(?rquery(idx[chd[chno[u]]], idx[u]),
   ?pquery up(par[chd[chno[u]]][0], anc));
int ?rquery(int u, int v) {
 int l = lca(u, v);
  return f(?pquery up(u, l), ?pquery up(v, l));
adj[u].clear(); hvc[u] = -1;
tt = 0; dfs(0, 0);
chd[ch] = -1;
ch cnt = 0, idx cnt = 0; hld(0, 0);
```

26 HOPCROFT_KARP

```
// 1-based
const int N = 1e5+5, INF = 1e8 + 5;
vector <int> g[N];
int n, e, match[N], dist[N];
bool bfs() {
   queue <int> q;
   for (int i = 1; i <= n; ++i) {
      if (!match[i]) dist[i] = 0, q.emplace(i);
      else dist[i] = INF;
   }
   dist[0] = INF;
   while (!q.empty()) {
      int u = q.front(); q.pop();
      if (!u) continue;</pre>
```

```
for (int v : q[u]) -
      if (dist[match[v]] == INF) {
        dist[match[v]] = dist[u] + 1,
        q.emplace(match[v]);
  return dist[0] != INF;
bool dfs (int u) {
 if (!u) return 1;
  for (int v : g[u]) {
    if (dist[match[v]] == dist[u] + 1 and

    dfs(match[v])) {

      match[u] = v, match[v] = u;
      return 1;
  dist[u] = INF;
 return 0;
int hopcroftKarp() {
  int ret = 0;
  while (bfs()) {
   for (int i = 1; i \le n; ++i) {
      ret += !match[i] and dfs(i);
  return ret;
```

27 HUNGARIAN_ALGORITHM

```
template<typename T>
pair<T, vector<int>>> MinAssignment(const

→ vector<vector<T>> &c) {
 int n = c.size(), m = c[0].size();
                                            //
\rightarrow assert(n <= m);
                                            // v:
 vector<T> v(m), dist(m);

→ potential

 vector<int> L(n, -1), R(m, -1);
                                            //

→ matching pairs

 vector<int> idx(m), prev(m);
 iota(idx.begin(), idx.end(), 0);
 auto residue = [&](int i, int j) { return c[i][j]
→ - v[j]; };
 for (int f = 0; f < n; ++f) {
   for (int j = 0; j < m; ++j) {
      dist[j] = residue(f, j); prev[j] = f;
   T w; int j, l;
   for (int s = 0, t = 0;;) {
     if (s == t) {
        l = s; w = dist[idx[t++]];
        for (int k = t; k < m; ++k) {
            = idx[k]; T h = dist[j];
          if (h <= w) {
            if (h < w)  { t = s; w = h; }
            idx[k] = idx[t]; idx[t++] = j;
        for (int k = s; k < t; ++k) {
         i = idx[k];
```

```
if (R[j] < 0) goto aug;
    int q = idx[s++], i = R[q];
    for (int k = t; k < m; ++k) {
        = idx[k];
      \tilde{T} h = residue(i,j) - residue(i,q) + w;
      if (h < dist[j]) {
         dist[j] = h; prev[j] = i;
        if (h == w) {
          if (R[j] < 0) goto aug;
idx[k] = idx[t]; idx[t++] = j;
aug:
  for(int k = 0; k < l; ++k)
    v[idx[k]] += dist[idx[k]] - w;
  int i;
    R[j] = i = prev[j];
    swap(j, L[i]);
  } while (i != f);
\dot{\mathsf{T}} ret = 0;
for (int i = 0; i < n; ++i) {
  ret += c[i][L[i]]; // (i, L[i]) is a solution
return {ret, L};
```

28 KMP

```
vector<int> get pi(string& s){
 int n = s.size();
 vector<int> pi(n);
 for (int k = 0, i = 1; i < n; ++i){
   if(s[i] == s[k]) pi[i] = ++k;
   else if(k == 0) pi[i] = 0;
   else k = pi[k-1], --i;
return pi;
pi.back(): n
// Borders = pi.back(), pi[pi.back() - 1], ...
// Prefix palindrome: s + "#" + rev(s)
// Number of occurrences of each prefix:
|vector<int> pref occur(vector<int> &pi) {
  int n = pi.size();
  vector<int> pref occur(n + 1);
  for (int i = 0; \overline{i} < n; ++i) {
    pref occur[pi[i]]++;
  for (int len = n; len > 0; --len) {
    pref occur[pi[len - 1]] += pref occur[len];
    pref occur[len]++;
  return pref occur;
// Find the length of the longest proper suffix of
   a suffix which also its prefix
// Reverse -> Find prefix function -> Reverse
// Find minimum length string such that given

→ strings occur as substring
```

29 MANACHER

```
// p[0][i] = half length of longest even palindrome
   around pos i-1, i and starts at i-p[0][i] and
  ends at i+p[0][i]-1
// p[1][i] = longest odd (half rounded down)
   palindrome around pos i and starts at i-p[1][i]
   and ends at i+p[1][i]
vector<vector<int>> manacher(string &s) {
 int n = s.size();
 vector<vector<int>> p(2, vector<int> (n));
 for (int z = 0; z < 2; ++z) {
  for (int i=0, l=0, r=0; i<n; ++i) {</pre>
      int t = r-i+!z;
      if (i<r) {
        p[z][i] = min(t, p[z][l+t]);
      int L = i - p[z][i], R = i + p[z][i] - !z;
      while (L>=1 \text{ and } R+1< n \text{ and } s[L-1]==s[R+1]) {
        p[z][i]++, L--, R++;
      if (R>r) {
        l=L, r=R;
 return p;
```

30 MATRIX EXPO

```
using row = vector<int>;
using matrix = vector<row>;
matrix unit mat(int n) {
 matrix I(n, row(n));
 for (int i = 0; i < n; ++i){
    I[i][i] = 1;
 return I;
matrix mat mul(matrix a, matrix b) {
 int m = a.size(), n = a[0].size();
 int p = b.size(), q = b[0].size();
 // assert(n==p);
 matrix res(m, row(q));
 for (int i = 0; i < m; ++i){
    for (int j = 0; j < q; ++j){
     for (int k = 0; k < n; ++k){
        res[i][j] = (res[i][j] + a[i][k]*b[k][j]) %

→ mod;

 return res;
matrix mat exp(matrix a, int p) {
 int m = a.size(), n = a[0].size();
 // assert(m==n);
 matrix res = unit mat(m);
 while (p) {
```

DU_NotReadyYet, University of Dhaka

```
17
```

```
if (p&1) res = mat_mul(a, res);
    a = mat_mul(a, a);
    p >>= 1;
}
return res;
}
```

31 MCF

```
struct MCF {
 int n;
 vector<vector<array<ll, 5>>> adj;
                                        // v, pos of

    u in v, cap, cost, flow

 vector<ll> dis, par, pos;
 MCF(int n): n(n), adj(n), dis(n), par(n), pos(n)
→ {}
 void add edge(int u, int v, int cap, int cost) {
    adj[u] push back({v, adj[v].size(), cap, cost,
   adj[v].push back({u, adj[u].size() - 1, 0,
   -cost, 0});
 ll spfa(int s, int t) {
   dis.assign(n, INF);
   vector<ll> mn cap(n, INF), ing(n);
   queue<int> q;
   q.push(s), inq[s] = 1, dis[s] = 0;
   while (!q.empty()) {
     int u = q.front(); q.pop();
      inq[u] = 0;
      for (int i = 0; i < adj[u].size(); ++i)</pre>
        auto [v, idx, cap, cost, flow] = adj[u][i];
        if (cap > flow and dis[v] > dis[u] + cost) { |}
         dis[v] = dis[u] + cost;
par[v] = u;
          pos[v] = i;
          mn cap[v] = min(mn cap[u], cap - flow);
          q.push(v);
          inq[v] = 1;
   return (mn cap[t] == INF? 0: mn cap[t]);
 array<ll, 2> get(int s, int t, ll max flow = INF)
   Il flow = 0, mc = 0;
   while (ll f = min(spfa(s, t), max flow - flow))
      flow += f;
     mc += f * dis[t];
      int u = t;
      while (u != s) {
        int p = par[u];
        adj[p][pos[u]][4] += f;
        adj[u][adj[p][pos[u]][1]][4] -= f;
    return {flow, mc};
```

```
};
MCF mcf(n);
for (int e = 0; e < m; ++e) {
   int u, v, r, c; cin >> u >> v >> r >> c; u--,
   w--;
   mcf.add_edge(u, v, r, c);
}
auto [f, mc] = mcf.get(0, n - 1, k);
```

32 MO_ALOGO

```
vector<array<int, 4>> cu(m);
for (int i = 0; i < m; ++i) {
   auto &[b, l, r, idx] = cu[i];
   cin >> l >> r; l--;
   b = r / B;
   idx = i;
}
sort(cu.begin(), cu.end());
int s = 0, e = -1;
for (auto [b, l, r, i]: cu) {
   while (l < s) add(--s);
   while (e < r) add(++e);
   while (s < l) remove(s++);
   while (r < e) remove(e--);
   ans[i] = cur_ans;
}</pre>
```

33 NUMBER_THEORY

```
## Floor
|ll floor (ll n, ll k) {
  if (n \ge 0) return n / k;
  return (n - (k - 1)) / k;
 ## Ceil
ll ceil (ll n, ll k) {
  if (n \ge 0) return (n + k - 1) / k;
  return n / k;
## Modular Inverse
[inv[1] = 1;
|for(int i = 2; i < N; ++i) -
  inv[i] = -(mod / i) * inv[mod % i] % mod;
  inv[i] += mod;
## Highly Composite Number
1e6(240), 1e9(1344), 1e12(6720), 1e14(17280)
## Harmonic Lemma (ceill)
|ll i = 1;
|while (i < n) {
  ll cval = (n + i - 1) / i;
  ll j = (n + cval - 2) / (cval - 1);
  // ceil(n/i)...ceil(n/(j - 1)) = cval
cout << i << " " << j - 1 << ": " << cval <<
 - "\n";
  i = j;
ĺl bezout(ll a, ll b, ll &x, ll &y){
  if(b == 0){
    x=1. v=0:
    return a;
  ĺl g = bezout(b, a%b, y, x);
```

```
y -= a/b*x;
 return q;
fl mod inv(ll a, ll m){
 ll x, y;
 ll g = bezout(a, m, x, y);
 if(q != 1) return -1; //no solution exists
 return (x%m+m)%m;
## Linear-sieve
int lpf[N], pm[N], pcnt = 0;
for (int i = 2; i < N; ++i) {
 if (!lpf[i]) lpf[i] = i, pm[pcnt++] = i;
 for (int j = 0; j < sz; ++j) {
   int p = pm[j];
   if ([pf[i]] = N) break;
    lpf[i * p] = p;
## Miller-Rabin
bool isp(ll n){
 if(n==2 || n == 3) return 1;
 if(n<=1 | n%2==0) return 0;
 for (int k = 0; k < 10; ++k){
   ll a = 2+rand()%(n-2);
   ll s = n-1;
    while(!(s\&1)) s>>=1;
    if(powmod(a, s, n) == 1) continue;
    int iscomp = 1;
    while (s!=n-1)
      if(powmod(a, s, n)==n-1){
        iscomp = 0;
        break;
      s=s<<1;
    if(iscomp) return 0;
 return 1;
## Miller-Rabin Deterministic:
bool check composite(u64 n, u64 a, u64 d, int s) {
 u64 x = \overline{binpower}(a, d, n);
 if (x == 1 | | x == n - 1)
    return false;
 for (int r = 1; r < s; r++) {
   x = (u128)x * x % n;
   if (\dot{x} == \dot{n} - 1)
      return false;
 return true;
bool isp(u64 n) {
 if (n < 2)
   return false;
 int r = 0;
 u64 d = n - 1;
 while ((d \& 1) == 0) {
   d >>= 1:
    r++;
```

```
for (int a: {2, 3, 5, 7, 11, 13, 17, 19, 23, 29,

→ 31, 37}) {
    if (n == a)
      return true;
    if (check composite(n, a, d, r))
      return false:
  return true;
## Prime Factorize of large number(Pollard Rho):
ll f(ll x, ll c, ll n){
  return (mulmod(x,x,n)+c)%n;
il pollard rho(ll n){
  if(n == 1) return 1;
  if(n%2 == 0) return 2;
  ll x = rand()%(n-2)+2;
  ll y = x;
  ll c = rand()%(n-1)+1;
  ll g = 1;
  while (g == 1){
    x = f(x, c, n);
    y = f(y, c, n);
    y = f(y, c, n);
    g = \underline{gcd(abs(x-y), n)};
  return q;
vector<ll> prime_factorize(ll n){
  if(n<=1) return vector<ll>();
  if(isp(n)) return vector<ll> ({n});
 il d = pollard_rho(n);
vector<ll> v = factorize(d);
vector<ll> w = factorize(n/d);
  v.insert(v.end(), w.begin(), w.end());
  sort(v.begin(), v.end());
  return v;
// auto pf = prime factorize(n);
## Number of divisors of n O(n^1/3):
int nod(ll n){
  sieve();
  int ret = 1;
  for (int i = 2; 1LL*i*i*i <= n; ++i) {
    if(isp[i]){
      int e = 0;
      while (n\%i == 0) {
        e++;
        n /= i;
      ret *= e+1;
  ll sq = sqrt(1.0L*n);
  if(isprime(n)) ret *= 2;
  else if(n == sq*sq and isprime(sq)) ret *= 3;
  else if(n!=1) ret *= 4;
  return ret;
## Smallest inverse phi
ll_inv_phi(ll phi, ll n, int pc) {
  if (phi == 1) return n;
  if (pc == -1) return INF;
  ll ret = inv_phi(phi, n, pc - 1);
```

```
if (phi % (p[pc] - 1) == 0) {
    phi /= (p[pc] - 1);
    n = n / (p[pc] - 1) * p[pc];
    while (phi % p[pc] == 0) {
      phi /= p[pc];
    ret = min(ret, inv phi(phi, n, pc - 1));
  return ret;
ll phi; cin >> phi;
|if (phi & 1) {
  cout << (phi == 1) << "\n";
  for (int i = 1; i * i <= phi; ++i) {
    if (phi % i == 0) {
      if (isp(i + 1)) {
        p.push back(i + 1);
      if (i * i != phi and isp(phi / i + 1)){
        p.push_back(phi / i + 1);
  sort(p.begin(), p.end());
  ll ans = inv_phi(phi, phi, p.size() - 1);
  cout << (ans == INF? 0: ans) << "\n";
## GCD sum function from 1 to N:
ll phi[N], g[N];
void pcgsm(){ //pre calculate gcd sum fucntion
  pcphi();
  for (int i = 1; i < N; ++i){
    for (int j = i; j < N; j+=i){
      g[j] += i*phi[j/i];
## All Pair gcd sum:
for (int i = 1; i < N; ++i) {
  for (int j = i; j < N; j += i) {
    qcd sum[j] += 1ll * phi[i] * (j / i);
  gcd sum[i] -= i;
  pref gcd sum[i] = pref gcd sum[i - 1] +

    qcd sum[i];

## LCM sum function of n:
ll lsm(ll n){
  ll ret=0;
  for(ll d=1; d*d<=n; d++){
    if(n%d==0){
      ret += d*phi(d);
      if(n/d!=d) ret += n/d*phi(n/d);
  return (ret+1)*n/2;
## LCM sum function from 1 to N
ll phi[N], l[N];
void pclsm(){ //pre calculate lcm sum function
  pcphi();
  for (int i = 1; i < N; ++i){
    for (int j = i; j < N; j+=i){
      l[i] += i*phi[i];
```

```
for (int i = 1; i < N; ++i){
    l[i] = (l[i]+1)*i/2;
## All pair lcm sum:
for (int i = 1; i < N; ++i) {
  for (int j = i; j < N; j += i) {
    lcm sum[i] += i * phi[i];
  lcm sum[i]++;
  lcm_sum[i] /= 2;
  lcm_sum[i] *= i;
  lcm sum[i] -= i;
  pref lcm sum[i] = lcm sum[i];
  pref lcm sum[i] += pref lcm sum[i - 1];
## Number of co-prime pairs of an array:
vector<ll> cnt(A);
for (int xi: x) {
 for (int d = 1; d * d <= xi; ++d) {
    if (xi % d == 0) {
      cnt[d]++;
      if (xi / d != d) {
        cnt[xi / d]++;
ll ans = 0;
for (int i = 1; i < A; ++i) {
 if (!sq_free[i]) continue;
ll ways = cnt[i] * (cnt[i] - 1) / 2;
  if (pf[i].size() \& 1 ^{\circ} 1) ans += ways;
  else ans -= ways;
## All pair gcd sum of an array:
vector<ll> cnt(A);
for (auto ai: a) {
 for (int d = 1; d * d <= ai; ++d) {</pre>
    if (ai % d == 0) {
      cnt[d]++;
      if (ai / d != d) {
        cnt[ai / d]++;
vector<ll> left(A);
iota(left.begin(), left.end(), 0);
for (int i = 1; i < A; ++i) {
 ll add = left[i] * cnt[i] * (cnt[i] - 1) / 2;
  sum += add;
  for (int j = 2 * i; j < A; j += i) {
   left[i] -= left[i];
ll crt(ll r1, ll m1, ll r2, ll m2){
  if(m1<m2) swap(r1, r2), swap(m1, m2);
 ll'p, q, g = bezout(m1, m2, p, q);
  if((r2-r1)%g!=0) return -1; //no solution
```

```
ll x = (r2-r1)\%m2*p\%m2*m1/q + r1;
  return x<0? x+m1*m2/q: x;
ĺl crt(vector<ll>& r, vector<ll>& m){
  ll x = r[0], M=m[0];
 for (int i = 1; i < r.size(); ++i){
  x = crt(x, M, r[i], m[i]);</pre>
    ll g = gcd(M, m[i]);
    M = (M/\overline{g})^*(m[i]/g);
  return x;
## Discrete Logarithm
ll discrete log(ll a, ll b, ll m) {
  a \% = m, b \% = m;
  if(a == 0){
    return (b == 0? 1: -1);
  ll k = 1, add = 0, g;
 while ((g = __gcd(a, m)) > 1) {
   if (b == k)    return add;
    if (b % q) return -1;
    b /= q, m /= q, k = (k * a / q) % m, ++add;
  int n = sqrt(m) + 1;
  unordered map<int, int> vals;
  for (ll q = 0, cur = b; q \le n; ++q) {
    vals[cur] = q;
    cur = (cur * a) % m;
  llan = 1;
  for (int i = 0; i < n; ++i) {
    an = (an * a) % m;
  for (ll p = 1, cur = k; p \le n; ++p) {
    cur = (cur * an) % m;
    if (vals.count(cur)) {
      return n * p - vals[cur] + add;
  return -1;
```

34 ONLINE_BRIDGE

```
vector<int> par, dsu 2ecc, dsu cc, dsu cc size;
int bridges;
int lca iteration;
vector<int> last visit;
void init(int n) {
    par.resize(n);
    dsu 2ecc.resize(n);
    dsu cc.resize(n);
    dsu cc size.resize(n);
    lca_{iteration} = 0:
    lasT visit.assign(n, 0);
    for (int i=0; i<n; ++i) {
    dsu_2ecc[i] = i;</pre>
         dsu^{-}cc[i] = i;
        dsu\_cc\_size[i] = 1:
         parTil = -1;
    bridges = 0;
```

```
int find 2ecc(int v) {
    if (\overline{v} == -1)
         return -1;
    return dsu 2ecc[v] == v ? v : dsu 2ecc[v] =
    find 2ecc(dsu 2ecc[v]);
|int find cc(int v) {
    v = find 2ecc(v);
    return dsu cc[v] == v ? v : dsu cc[v] =
   find cc(dsu cc[v]);
void make root(int v) {
    v = find 2ecc(v);
    int root = v;
    int child = -1;
    while (v != -1) {
         int p = find 2ecc(par[v]);
         par[v] = chiTd;
         dsu^{c}c[v] = root;
         child = v:
         v = p;
    dsu cc size[root] = dsu cc size[child];
void merge path (int a, int b) {
    ++lca Iteration;
    vector<int> path a, path b;
    int lca = -1;
    while (lca == -1) {
         if (a != -1) {
    a = find_2ecc(a);
             path a.push back(a);
             if (\lambda ast visit[a] == lca iteration){
                  lca = a;
                  break;
             last visit[a] = lca iteration;
             a = \overline{par[a]};
         if (b != -1) {
             \tilde{b} = find 2ecc(b);
             path b.push back(b);
             if (\lambda ast visit[b] == lca iteration){
                  lca \equiv b:
                  break;
             last visit[b] = lca iteration;
             b = \overline{par[b]};
    for (int v : path a) {
         dsu 2ecc[v] = lca;
         if (v == lca)
             break;
         --bridges;
    for (int v : path b) {
         dsu 2ecc[v] = lca;
         if (v == lca)
             break;
```

```
--bridges;
void add edge(int a, int b) {
    a = find 2ecc(a);
    b = find^2 ecc(b);
    if (a == b)
        return:
    int ca = find cc(a);
    int cb = find cc(b);
    if (ca != cb) {
        ++bridges;
        if (dsu cc size[ca] > dsu cc size[cb]) {
            swap(a, b);
            swap(ca, cb);
        make root(a);
        par[\overline{a}] = dsu cc[a] = b;
        dsu cc size[cb] += dsu cc size[a];
        merge path(a, b);
```

35 PALINDROMIC_TREE

```
const int N = 1e5+10;
struct vertex
 int len, link, no of suf pal;
 map<char, int> next;
}pt[N];
int sz, at, cnt[N];
char s[N];
void pt init(){
 for (int i = 0; i < N; ++i)
   pt[i].next.clear();
 memset(cnt, 0, sizeof(cnt));
 pt[0].len = -1, pt[0].link = 0,
\rightarrow pt[0].no of suf pal = 0;
 pt[1].len = 0, pt[1].link = 0,
\rightarrow pt[1].no of suf pal = 0;
 sz = at = T;
void pt extend(int si){ //string index
 while (s[si - pt[at].len - 1] != s[si]) at =
→ pt[at].link;
 int x = pt[at].link, c = s[si]-'a';
 while (s[si - pt[x].len - 1] != s[si]) x =
\rightarrow pt[x].link;
 if(!pt[at].next.count(c)){
    pt[at].next[c] = ++sz;
    pt[sz].len = pt[at].len + 2;
    // cnt[pt[at].len+2]++; //for finding number

→ of distinct palindrome of lenght k

    pt[sz].link = (pt[sz].len == 1)? 1 :
   pt[x].next[c];
    // pt[sz].no_of_suf pal = 1 +
   pt[pt[sz].link].no of suf pal; //for finding
  number of palindrome which last position is si
```

```
// cnt[pt[at].len + 2]++; //for finding number

→ of palindrome of lenght k

 at = pt[at].next[c];
int num of pal(int ai){  //distinct palindrome,

→ array index

 int ret = pt[at].ans;
  for(auto x : pt[ai].next)
    ret += num of pal(x.second);
  return ret;
int main(){
 scanf("%s", s);
  pt init();
  for (int i = 0; s[i]; ++i){
   pt extend(i);
  int ans = num of pal(0) + num of pal(1) - 2;
  printf("%d\n", ans);
  return 0;
```

36 PERSISTENT_SEGMENT_TREE

```
## Point Addition & Range Sum:
struct node {
  ll sum;
  node *ĺ, *r;
  node(ll \ s = 0, \ node \ *l = NULL, \ node \ *r = NULL):
 \rightarrow sum(s), l(l), r(r) {}
node* add(node *u, int i, int x, int s, int e) {
  if (s == e) return new node(u->sum + x);
  if (!u->l) u->l = new node(), u->r = new node();
  node *nu = new node(u->sum, u->l, u->r);
  int m = (s + e) / 2;
  if (i \le m) nu \rightarrow l = add(nu \rightarrow l, i, x, s, m);
  else nu \rightarrow r = add(nu \rightarrow r, i, x, m + 1, e);
  nu->sum = nu->l->sum + nu->r->sum:
  return nu;
ll rsum(node *u, int l, int r, int s, int e) {
  if (!u) return 0;
  if (s > r \text{ or } e < l) return 0;
  if (l <= s and e <= r) return u->sum;
  int m = (s + e) / 2;
  return rsum(u->1, l, r, s, m) + rsum(u->r, l, r,
 \rightarrow m + 1. e):
vector<node*> root(VER);
root[0] = new node(); // initialization
root[k] = add(root[k], i, x, 0, sz - 1);
root[ver++] = root[k];
cout << rsum(root[k], l, r, 0, sz - 1) << "\n"; ## count numbers > k in a range
root[0] = new node();
for (int i = 0; i < n; ++i) {
  root[i + 1] = add(root[i], a[i], 1);
while (q--) {
  int l, r, k; cin >> l >> r >> k; l--, r--;
  int ans = rsum(root[r + 1], k, E - 1) -
 \rightarrow rsum(root[l], k, E - 1);
```

37 PERSISTENT_TRIE

```
struct node -
 node *nxt[2];
node() { fill(nxt, nxt + 2, nullptr); }
node *new root = new node();
  node * cu\overline{r} = new root;;
  for (int idx = I\overline{D}X - 1; idx >= 0; --idx) {
    int f = (x \gg idx) \& 1;
    if (prev and prev->nxt[!f]) cur->nxt[!f] =
   prev->nxt[!f];
    cur->nxt[f] = new node();
    cur = cur->nxt[f];
    if (prev) prev = prev->nxt[f];
  return new root;
|int get max(node *root, int x) {
 if (!root) return 0;
  node *u = root;
  int ret = 0;
  for (int idx = IDX - 1; idx \Rightarrow 0; --idx) {
    int f = (x >> idx) \& 1;
    if (u->nxt[!f]) ret += (1 << idx), u =
   u->nxt[!f];
    else u = u - nxt[f];
  return ret;
```

38 POLYNOMIAL_INTERPOLATION

```
// P(x) = a0 + a1x + a2x^2 + ... + anx^n
// y[i] = P(i)
ll eval (vector<ll> y, ll k) {
   int n = y.size() - 1;
   if (k <= n) {
      return y[k];
   }
   vector<ll> L(n + 1, 1);
   for (int x = 1; x <= n; ++x) {</pre>
```

```
L[0] = L[0] * (k - x) % mod;
L[0] = L[0] * inv(-x) % mod;
}
for (int x = 1; x <= n; ++x) {
    L[x] = L[x - 1] * inv(k - x) % mod * (k - (x - 1)) % mod;
    L[x] = L[x] * ((x - 1) - n + mod) % mod *
    - inv(x) % mod;
}
ll yk = 0;
for (int x = 0; x <= n; ++x) {
    yk = add(yk, L[x] * y[x] % mod);
}
return yk;
}
```

39 SCC

```
void dfs1(int u, vector<int> *adj, vector<int>
vis[u] = 1;
 for (int \&v: adj[u]) {
   if (!vis[v]) {
     dfs1(v, adj, vis, order);
 order.emplace back(u);
void dfs2(int u, vector<int> *rev adj, vector<int>
scc.emplace back(u);
 vis[u] = 1;
 for (int &v: rev adj[u]) {
   if (!vis[v]) {
     dfs2(v, rev adj, vis, scc);
vector<vector<int>> get sccs(int n, vector<int>
→ *adj) {
 vector<int> vis(n), order;
 for (int u = 0; u < n; ++u) {
   if (!vis[u]) {
     dfs1(u, adj, vis, order);
 vector<int> rev adj[n];
 for (int u = 0; u < n; ++u) {
   for (int v: adi[u]) {
     rev adj[v].emplace back(u);
 vector<vector<int>> sccs;
 reverse(order.begin(), order.end());
 vis.assign(n, 0);
 for (int u: order) {
   if (!vis[u]) {
     sccs.emplace back(0);
     dfs2(u, rev adj, vis, sccs.back());
 return sccs;
```

```
}
vector<vector<int>> sccs = get_sccs(n, adj);
int tot_scc = sccs.size();
vector<int>> scc_no(n);
for (int i = 0; i < tot_scc; ++i) {
    for (int u: sccs[i]) {
        scc_no[u] = i;
    }
}</pre>
```

40 SEGMENT_TREE

```
## Range Addition and Range Assign and Range sum
int n:
ll t[3 * N], p[3 * N], p2[3 * N]; //t for sum, p

→ for assign & p2. for add

void pull(int v) {
 t[v] = \dot{t}[2 * \dot{v}] + t[2 * v + 1];
void push(int v, int st, int ed) {
 int lc = 2 * v, rc = 2 * v + 1, md = (st + ed) /
  if (p[v] != -1) {
    t[lc] = p[v] * (md - st + 1);
   t[rc] = p[v] * (ed - md);
    p[lc] = p[rc] = p[v];
    p2[lc] = p2[rc] = 0;
    p[v] = -1;
 if (p2[v]) {
    t[lc] + p2[v] * (md - st + 1);
    t[rc] += p2[v] * (ed - md);
    p2[lc] += p2[v];
    p2[rc] += p2[v];
    p2[v] = 0;
void assign(int l, int r, int x, int v = 1, int st
\rightarrow = 0, int ed = n - 1) {
  if (l > ed or r < st) return;</pre>
  if (l <= st and ed <= r) {
    t[v] = 1LL * (ed - st + 1) * x;
    p[v] = x;
    p2[v] = 0;
    return;
  int lc = 2 * v, rc = 2 * v + 1, md = (st + ed) /
  push(v, st, ed);
  assign(l, r, x, lc, st, md);
  assign(l, r, x, rc, md + 1, ed);
  pull(v);
void add(int l, int r, int x, int v = 1, int st =
\rightarrow 0, int ed = n - 1) {
  if (l > ed or r < st) return;</pre>
  if (l <= st and ed <= r) {
    t[v] += 1LL * (ed - st + 1) * x;
    p2[v] += x;
    return :
  push(v, st, ed);
  int lc = 2 * v, rc = 2 * v + 1, md = (st + ed)
  add(l, r, x, lc, st, md);
```

```
add(l, r, x, rc, md + 1, ed);
  pull(v);
ll rsum(int l, int r, int v = 1, int st = 0, int ed
\rightarrow = n - 1) {
 if (l > ed or r < st) return 0;</pre>
  if (l <= st and ed <= r) return t[v];</pre>
  push(v, st, ed);
  int lc = 2 * v, rc = 2 * v + 1, md = (st + ed) /
 ll lret = rsum(l, r, lc, st, md);
  ll rret = rsum(l, r, rc, md + 1, ed);
  return lret + rret;
## Make All Elements <= k and Make all elements >=
   k on range & Point Query:
const int I = 1e9 + 9;
int t[3 * N], pa[3 * N], pr[3 * N], ar[3 * N]; //pa
    for propagate adding, pr for propagate remove,
- ar for check last on is adding(1) or remove(0)
void fg(int x, int u) { //function for

→ make greater

 t[u] = \overline{max(t[u], x)};
  pa[u] = max(pa[u], x);
  pr[u] = max(pr[u], x);
  ar[u] = 1:
void fl(int x, int u) { //function for make less
 t[u] = min(t[u], x);
  pr[u] = min(pr[u], x);
  pa[u] = min(pa[u], x);
  ar[u] = 0:
|void push(int u) {
 int v = 2 * u, w = 2 * u + 1;
  if (ar[u] == 0) {
    if (pa[u] != -1) {
      fg(pa[u], v); fg(pa[u], w);
    if (pr[u] != I) +
      fl(pr[u], v); fl(pr[u], w);
  } else {
    if (pr[u] != I) -
      fl(pr[u], v); fl(pr[u], w);
    if (pa[u] != -1) {
      fg(pa[u], v); fg(pa[u], w);
 pa[u] = -1; pr[u] = I;
void make greater(int l, int r, int x, int u = 1,
   int s = 0, int e = N - 1) {
 if (l > e or r < s) return;</pre>
  if (l <= s and e <= r) {
    fq(x, u);
    return ;
  push(u);
  int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
  make greater(l, r, x, v, s, m);
 make_{greater}(l, r, x, w, m + 1, e);
```

```
void make less(int l, int r, int x, int u = 1, int
\rightarrow s = 0, int e = N - 1) {
 if (l > e or r < s) return;</pre>
 if (l <= s and e <= r) {
    fl(x, u):
    return;
 push(u);
 int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
 make less(l, r, x, v, s, m);
 make^{-less(l, r, x, w, m + 1, e)};
int at(int i, int u = 1, int s = 0, int e = N - 1) {
 if (s == e) return t[u];
 push(u);
 int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
 if (i <= m) return at(i, v, s, m);</pre>
 else return at(i, w, m + 1, e);
```

41 SHORTEST_PATH

```
## Dijkstra
priority queue<array<ll, 2>> pq;
vector<ll> dis(n, INF), vis(n);
|while (!pq.empty()) {
  auto [d, u] = pq.top(); pq.pop();
  if (vis[u]) continue;
  vis[u] = 1;
  for (auto [v, c]: next[u]) {
    if (dis[v] > d + c) {
      dis[v] = d + c;
      pq.push({dis[v], v});
## Bellman-ford
vector<int> bellman ford(int s){
 vector<int> dis(n. I):
  dis[s]=0;
  while(1){
    int anv=0:
    for (auto& e: ed){
      if(dis[e.u]<I){</pre>
        if(dis[e.u]+e.cost < dis[e.v]){</pre>
          dis[e.v] = dis[e.u]+e.cost;
          any=1;
    if(!any) break;
  return dis:
## Floy-Warshall
for (int k = 0; k < n; ++k) {
 for (int i = 0; i < n; ++i) {
    for (int j = 0; j < n; ++j) {
      dis[i][j] = min(dis[i][j], dis[i][k] +
   dis[k][j]);
```

42 SOS_DP

```
## Count over subset
for (int i = 0; i < n; ++i) f[a[i]] = ?;
for (int i = 0; i < n; ++i) {
  for (int mask = 0; mask < (1 << n); ++mask) {
    if (\max k (1 << i)) {
      f[mask] += f[mask^(1<<i)];
## Count over superset
for (int i = 0; i < n; ++i) f[a[i]] = ?;
for (int i = 0; i < n; ++i) {
  for (int mask = (1 << n) - 1; mask >= 0; --mask)
    if (!(mask&(1<<i))) {
      f[mask] += f[mask^(1<<i)];
## How many pairs in ara[] such that (ara[i] &
\rightarrow ara[j]) = 0
/// N --> Max number of bits of any array element
const int N = 20:
int inv = (1 << N) - 1:
int F[(1 << N) + 10];
int ara[MAX];
/// ara is 0 based
long long howManyZeroPairs(int n,int ara[]) {
    CLR(F);
    for(int i=0;i<n;i++) F[ara[i]]++;</pre>
    for(int i = 0; i < N; ++i)
        for(int mask = 0; mask < (1<<N); ++mask){</pre>
            if(mask \& (1<<i))
                 F[mask] += F[mask^{(1<<i)}];
    long long ans = 0;
    for(int i=0;i<n;i++) ans += F[ara[i] ^ inv];</pre>
    return ans;
/// To get
    for(int mask = 0; mask < (1<<N); ++mask)
        for(int i = 0; i < (1 << N); ++i)
            if( (mask \& i) == mask ) { /// i is a
 F[mask] += A[i];
/// The code is the following
    for(int i = 0; i < (1 << N); ++i) F[i] = A[i];
    for(int i = 0; i < N; ++i)
        for(int mask = (1 << N) - 1; mask >= 0; --mask){
            if (!(mask & (1<<i)))
                F[mask] += F[mask \mid (1 << i)];
## Number of subsequences of ara[0:n-1] such that
## sub[0] \& sub[2] \& ... \& sub[k-1] = 0
const int N = 20;
int inv = (1 << N)^{'} - 1;
int F[(1 << N) + 10];
int ara[MAX]:
int p2[MAX]; /// p2[i] = 2^i
```

```
///0 based array
int howManyZeroSubSequences(int n,int ara[]) {
    CLR(F);
    for(int i=0;i<n;i++) F[ara[i]]++;</pre>
    for(int i = 0; i < N; ++i)
         for(int mask = (1 << N) - 1; mask >= 0; --mask){
             if (!(mask & (1<<i)))
                 F[mask] += F[mask | (1 << i)];
    int ans = 0;
    for(int mask=0; mask<(1<<N); mask++) {
         if( builtin popcount(mask) \& 1) ans =
    sub(ans, p2[F[mask]]);
         else ans = add(ans, p2[F[mask]]);
    return ans;
## Number of subsequences of ara[0:n-1] such that ## sub[0] | sub[2] | ... | sub[k-1] = Q
int F[(1 << 20) + 10], ara[MAX];
int p2[MAX]; /// p2[i] = 2^i
/// ara is 0 based
int howManySubsequences(int n, int ara[], int m,
   int Q) {
    CLR(F);
    for(int i=0;i<n;i++) F[ara[i]]++;
    if(Q == 0) return sub(p2[F[0]], 1);
    for(int i = 0; i < m; ++i)
         for(int mask = 0; mask < (1<<m); ++mask){</pre>
             if (mask & (1 << i))
                 F[mask] += F[mask ^ (1<<i)];
    int ans = 0;
    for(int mask=0; mask<(1<<m); mask++) {</pre>
         if(mask & Q != mask) continue;
         if( builtin popcount(mask \hat{Q}) & 1) ans =
    sub(ans, p2[F[mask]]);
         else ans = add(ans, p2[F[mask]]);
    return ans;
```

43 SPARSE_TABLE

```
int n, a[N], lg[N], st[N][K];
for (int i = 2; i < N; ++i) {
    lg[i] = lg[i / 2] + 1;
}
void build() {
    for (int k = 0; k < K; ++k) {
        for (int i = 0; i + (1 << k) <= n; ++i) {
            if (k == 0) st[i][k] = a[i];
            else st[i][k] = min(st[i][k - 1], st[i + (1 - 1]);
        }
}
int rmq (int l, int r) {
    int k = lg[r - l + 1];
    return min(st[l][k], st[r - (1 << k) + 1][k]);
}</pre>
```

44 SPARSE TABLE 2D

```
int st[N][N][LG][LG];
int a[N][N], lg2[N];
int yo(int x1, int y1, int x2, int y2) {
 x2++;
 y2++;
 int a = \lg 2[x2 - x1], b = \lg 2[y2 - y1];
         \max(st[x1][y1][a][b], st[x2 - (1 <<
   a)][y1][a][b])
        \max(st[x1][y2 - (1 << b)][a][b], st[x2 -
   (1 \ll a)[y2 - (1 \ll b)][a][b])
       );
void build(int n, int m) { // 0 indexed
 for (int i = 2; i < N; i++) lg2[i] = lg2[i >> 1]
 for (int i = 0; i < n; i++) {
    for (int j = 0; j < m; j++) {
      st[i][j][0][0] = a[i][j];
 for (int a = 0; a < LG; a++) {
   for (int b = 0; b < LG; b++) {
     if (a + b == 0) continue:
      for (int i = 0; i + (1 << a) <= n; i++) {
        for (int j = 0; j + (1 << b) <= m; j++) {
          if (!a) {
            st[i][j][a][b] = max(st[i][j][a][b -
\rightarrow 1], st[i][j + (1 << (b - 1))][a][b - 1]);
          } else {
            st[i][j][a][b] = max(st[i][j][a -
-1][b], st[i + (1 << (a - 1))][j][a - 1][b]);
```

45 SQRT DECOMPOSITION

```
const int SZ2 = 2e5, SZ = sqrt(SZ2+.0)+1, N = SZ*SZ;
int n, a[N], b[SZ];
void build() {
 for (int i = 0; i < SZ; ++i){
   b[i] = INT MAX;
 for (int i = 0: i < n: ++i){
    b[i/SZ] = min(b[i/SZ], a[i]);
int rmq(int l, int r) {
 int lb = l/SZ, rb = r/SZ;
 int ret = INT MAX:
 if(lb==rb) -
    for (int i = l; i <= r; ++i){
      ret = min(ret, a[i]);
 } else {
    for (int i = l; i < (lb+1)*SZ; ++i){</pre>
      ret = min(ret, a[i]);
    for (int i = lb+1; i < rb; ++i){
```

```
ret = min(ret, b[i]);
}
for (int i = rb*SZ; i <= r; ++i){
    ret = min(ret, a[i]);
}
return ret;
}</pre>
```

46 STRESS_TESTING

```
set -e
g++ -std=c++17 gen.cpp -o gen
g++ -std=c++17 main.cpp -o main
g++ -std=c++17 brute.cpp -o brute
for((i = 1; ; ++i)); do
    echo $i
    ./gen $i > in
    ./main < in > out
    ./brute < in > out2
    diff -w out out2 || break
done
```

47 SUFFIX_ARRAY

```
array<vector<int>, 2> get sa(string& s, int
    lim=128) { // for integer, just change string
   to vector<int> and minimum value of vector must
 \rightarrow be >= 1
 int n = s.size() + 1, k = 0, a, b;
  vector<int> x(begin(s), end(s)+1), y(n), sa(n),

    lcp(n), ws(max(n, lim)), rank(n);

  x.back() = 0;
  iota(begin(sa), end(sa), 0);
  for (int j = 0, p = 0; p < n; j = max(1, j * 2),
\rightarrow lim = p) {
    p = j, iota(begin(y), end(y), n - j);
    for (int i = 0; i < n; ++i) if (sa[i] >= j)
   y[p++] = sa[i] - j;
    fill(begin(ws), end(ws), 0);
    for (int i = 0; i < n; ++i) ws[x[i]]++;
    for (int i = 1; i < \lim; ++i) ws[i] += ws[i -
    for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
    swap(x, y), p = 1, x[sa[0]] = 0;
    for (int i = 1; i < n; ++i) a = sa[i - 1], b =
      (y[a] == y[b] \&\& y[a + j] == y[b + j]) ? p -
\rightarrow 1 : p++;
  for (int i = 1; i < n; ++i) rank[sa[i]] = i;</pre>
  for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)
    for (k \&\& k--, j = sa[rank[i] - 1]; s[i + k] ==
\rightarrow s[j + k]; k++);
  sa.erase(sa.begin()), lcp.erase(lcp.begin());
  return {sa, lcp};
```

48 SUFFIX_AUTOMATON

```
int len[N], lnk[N], pos[N], sz, last;
map<char, int> nxt[N];
void init () {
```

```
len[0] = 0, lnk[0] = -1, nxt[0].clear(), last =
 \rightarrow 0, sz = 1;
void add (char c) {
  int cur = sz++;
  len[cur] = len[last] + 1, nxt[cur].clear();
  int u = last;
  while (u != -1 \text{ and } !nxt[u].count(c)) {
    nxt[\dot{u}][c] = cur;
    u = lnk[u];
  if (u == -1) {
    lnk[cur] = 0;
  else {
    int v = nxt[u][c];
    if (len[u] + 1 == len[v]) {
      lnk[cur] = v;
    else {
      int w = sz++;
      len[w] = len[u] + 1, lnk[w] = lnk[v], nxt[w]
   = nxt[v];
      while (u != -1 \text{ and } nxt[u][c] == v) {
        nxt[u][c] = w, u = lnk[u];
       lnk[cur] = lnk[v] = w;
  last = cur;
49 TREAP
## Typical TEAP
struct node {
  ll val, prior, sz, sum;
node *l, *r;
  node(int val, int prior, int sz) : val(val),
    prior(prior), sz(sz), sum(0), l(nullptr),
   r(nullptr){}
using pnode = node*;
|pnode root;
pnode new node(ll val){
  return new node(val, rand(), 1);
int get sz(pnode u){
  return u? u->sz: 0:
void update(pnode u){
  if (!u) return ;
```

u->sz = get sz(u->l) + 1 + get sz(u->r);

else split(u->l, l, u->l, val), r = u;

|void merge(pnode &u, pnode l, pnode r){

 \rightarrow u->r->sum: 0):

update(u);

if(!u) l = r = NULL;

if(!l or !r) u = l? l: r;

u->sum = u->val + (u->l? u->l->sum: 0) + (u->r?

else if(val > u->val) split(u->r, u->r, r, val),

void split(pnode u, pnode &l, pnode &r, ll val){

```
if(l->prior > r->prior) merge(l->r, l->r, r),
 else merge(r->l, l, r->l), u = r;
  update(u);
void insert(pnode &u, pnode it){
  if(!u) u = it;
  else if(it->prior > u->prior) split(u, it->l,
 \rightarrow it->r, it->val), u = it;
  else insert(it->val <= u->val ? u->l: u->r, it);
  update(u);
void erase(pnode &u, ll val){
  if(!u) return ;
  if(val == u \rightarrow val) merge(u, u \rightarrow l, u \rightarrow r);
  else erase(val < u->val ? u->l: u->r, val);
  update(u);
bool present(pnode u, int x){
  if(!u) return false;
  if(u->val == x) return true;
  if(u->val < x) return present(u->r, x);
  return present(u->l, x);
ll kth(pnode u, int k){
  if(get sz(u) < k) return INT MIN;</pre>
  if(qet sz(u->l) == k-1) return u->val;
  if(qet sz(u->1) < k-1) return kth(u->r, k-1)
 \rightarrow get sz(u->l) - 1);
  return kth(u->l, k);
int cnt less(pnode u, ll x){
  if(!u) return 0;
  if(x <= u->val) return cnt less(u->l, x);
  return get sz(u->1) + 1 + c\overline{n}t less(u->r, x);
ll sum less(pnode u, ll x) {
  if (\overline{!}u) return 0;
  if (x <= u->val) return sum less(u->l, x);
  return u \rightarrow val + (u \rightarrow l? u \rightarrow l \rightarrow sum: 0) +
 \rightarrow sum less(u->r, x);
## Implicit TREAP
struct node {
  ll val, sum;
  int prior, sz, rev;
  node *l, *r;
  node(){}
  node(ll val): val(val), sum(val), prior(rand()),
 \rightarrow sz(1), rev(0), l(nullptr), r(nullptr) {}
using pnode = node*;
pnode root:
int get sz(pnode t) {
  return t? t->sz: 0;
ll get sum(pnode t)
  return t? t->sum: 0;
void update(pnode &t) {
  if (!t) return ;
  t \rightarrow sz = get sz(t \rightarrow l) + 1 + get sz(t \rightarrow r);
  t - sum = qe\overline{t} sum(t - sl) + t - sum + qet sum(t - sr);
```

```
void push(pnode t) {
  if (t and t->rev) {
    swap(t->l, t->r);
t->rev = 0;
    if (t->l) {
      t->l->rev ^= 1;
    if (t->r) {
      t->r->rev ^= 1;
void merge(pnode &t, pnode l, pnode r){
  push(l):
  push(r);
  if(!l or !r) t=l?l:r;
  else if(l->prior > r->prior) merge(l->r, l->r,
  else merge(r->l,l,r->l) , t=r;
  update(t);
void split(pnode t, pnode &l, pnode &r, int pos,

    int add=0) {
  push(t);
  if(!t) return void(r=l=NULL);
  int cur pos = get sz(t->l)+add;
  if(pos > cur_pos) split(t->r, t->r, r, pos,
\rightarrow cur pos+1), l = t;
  update(t);
void insert(pnode &t, pnode it, int i) {
 pnode t1, t2;
split(t, t1, t2, i);
merge(t1, t1, it);
  merge(t, t1, t2);
void reverse(pnode &t, int l, int r) {
  pnode lt, mt, rt;
  split(t, t, rt, r + 1);
  split(t, lt, mt, l);
  \mathsf{mt}->rev = 1;
  merge(mt, mt, rt);
  merge(t, lt, mt);
pnode lt, mt, rt;
  split(t, t, rt, r + 1);
  split(t, lt, mt, l);
  ll ret = mt->sum;
  merge(mt, mt, rt);
  merge(t, lt, mt);
  return ret;
int n, q; cin >> n >> q;
vector<ll> a(n);
for (auto &ai: a) {
  cin >> ai;
for (int i = 0; i < n; ++i) {
  insert(root, new node(a[i]), i);</pre>
while (q--) {
  int tp, l, r; cin >> tp >> l >> r; l--, r--;
```

```
if (tp == 1) {
  reverse(root, l, r);
   cout << rsum(root, l, r) << "\n";</pre>
```

50 XOR_BASIS

```
ll rnk, basis[D];
void insert vector(ll mask){
 for (int \bar{1} = D-1; i >= 0; --i){
    if((mask & (111 \ll i)) == 0) continue;
    if(!basis[i]){
      basis[i] = mask, rnk++;
      return;
     else mask ^= basis[i];
```

51 Z_ALGORITHM

```
vector<int> get z(string s){
  int n=s.size(), l=1, r=0;
  vector<int> z(n); z[0]=n;
 for (int i = 1; i < n; ++i){
  if(i<=r)  z[i]=min(z[i-l], r-i+1);</pre>
    while(s[i+z[i]]==s[z[i]]) z[i]++;
    if(i+z[i]-1>r) l=i, r=i+z[i]-1;
  return z;
```

52 note

Binomial Coefficent

```
Factoring in: \binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}
Sum over k: \sum_{k=0}^{n} {n \choose k} = 2^{n}
```

Alternating sum: $\sum_{k=0}^{n} (-1)^k \binom{n}{k} = 0$

Even and odd sum: $\sum_{k=0}^{n} {n \choose 2k} = \sum_{k=0}^{n} {n \choose 2k+1} 2^{n-1}$

The Hockey Stick Identity

(Left to right) Sum over n and k: $\sum_{k=0}^{m} {n+k \choose k} = {n+m-1 \choose m}$

• (Right to left) Sum over n: $\sum_{m=0}^{n} {m \choose k} = {n+1 \choose k+1}$

Sum of the squares: $\sum_{k=0}^{n} {n \choose k}^2 = {2n \choose k}$

Weighted sum: $\sum_{k=1}^{n} k \binom{n}{k} = n2^{n-1}$

Connection with the fibonacci numbers: $\sum_{k=0}^{\infty} {n-k \choose k} = F_{n+1}$

Fibonacci Number

$$k = A - B$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1}$$
(1)

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n$$

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - (-1)^n$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1} \tag{4}$$

$$GCD(F_m, F_n) = F_{GCD(m,n)}$$
 (5)

$$\sum_{0 \le k \le n} \binom{n-k}{k} = Fib_{n+1} \tag{6}$$

$$\gcd(F_n, F_{n+1}) = \gcd(F_n, F_{n+2}) = \gcd(F_{n+1}, F_{n+2}) = 1$$
 (7)

Lucas Theorem

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p}$$

- $\binom{m}{n}$ is divisible by p if and only if at least one of the base-p digits of n is greater than the corresponding base-p digit of m.
- The number of entries in the *n*th row of Pascal's triangle that are not divisible by $p = \prod_{i=0}^{k} (n_i + 1)$
- All entries in the $(p^k-1)th$ row are not divisble by p. $\binom{n}{m} \equiv \lfloor \frac{n}{n} \rfloor \pmod{p}$

2nd Kaplansky's Lemma

The number of ways of selecting k objects, no two consecutive, from n labelled objects arrayed in a circle is $\frac{n}{b}\binom{n-k-1}{b-1} =$

Distinct Objects into Distinct Bins

- n distinct objects into r distinct bins = r^n
- Among *n* distinct objects, exactly *k* of them into r distincts bins = $\binom{n}{k} r^k$
- *n* distinct objects into *r* distinct bins such that each bin contains at least one object = $\sum_{i=0}^{r} (-1)^{i} {r \choose i} (r-i)^{n}$

Stirling Number 2nd Kind

- Count the number of ways to partition a set of n labelled objects into k nonempty unlabelled subsets.

$$S(n,k) = S(n-1,k-1) + k * S(n-1,k)$$

$$S(0,0) = 1, S(>0,0) = 0, S(0,>0) = 0$$

- Time Complexity: $O(k \log n)$

```
ll get sn2(int n, int k) {
     ll sn2 = 0;
     for (int i = 0; i \le k; ++i) {
       ll now = nCr(k, i) * powmod(k - i, n, mod) % mod;
       if (i&1) {
         now = now * (mod - 1) % mod:
       sn2 = (sn2 + now) % mod;
     sn2 = sn2 * ifact[k] % mod;
     return sn2;
(3)
```

- Number of ways to color a 1n grid using k colors such that each color is used at least once = k!.sn2(n,k)

Bell Numbers

Counts the number of partitions of a set.

$$B_{n+1} = \sum_{k=0}^{n} \left(\frac{n}{k}\right) \cdot B_k \tag{8}$$

 $B_n = \sum_{k=0}^n S(n,k)$, where S(n,k) is stirling number of second kind.

Partition Number

- Time Complexity: $O(n\sqrt{n})$ for (int i = 1; i <= n; ++i) {
 pent[2 * i - 1] = i * (3 * i - 1) / 2;
 pent[2 * i] = i * (3 * i + 1) / 2;</pre> p[0] = 1;for (int i = 1; $i \le n$; ++i) { p[i] = 0;for (int j = 1, k = 0; pent[j] <= i; ++j) { if (k < 2) p[i] = add(p[i], p[i - pent[j]]); else p[i] = sub(p[i], p[i - pent[j]]); ++k, k &=

- The number of partitions of a positive integer n into exactly k parts equals the number of partitions of n whose largest par equals k

$$p_k(n) = p_k(n-k) + p_{k-1}(n-1)$$

- The number of labeled undirected graphs with n vertices

- The number of connected labeled undirected graphs $16x^4 + 25x^5 + ... e^x = 1 + x + (x^2)/2! + (x^3)/3! + (x^4)/4! + ...$ with *n* vertices, $C_n = 2^{\binom{n}{2}} - \frac{1}{n} \sum_{k=1}^{n-1} k \binom{n}{k} 2^{\binom{n-k}{2}} C_k = 2^{\binom{n}{2}}$ $\sum_{k=1}^{n-1} {n-1 \choose k-1} 2^{{n-k \choose 2}} C_k$

- The number of k-connected labeled undirected graphs with nvertices, $D[n][k] = \sum_{s=1}^{n} {n-1 \choose s-1} C_s D[n-s][k-1]$ - Cayley's formula: the number of trees on n labeled vertices

= the number of spanning trees of a complete graph with n la- $\sum_{i=1}^{n} f_k(i) = \frac{1}{k+1} n(n+1)(n+2)...(n+k) = \frac{1}{k+1} \frac{(n+k)!}{(n-1)!}$

Defect vertices = n- Number of ways to color a graph using k color such that no $\sum_{i=0}^{n} nix^i = 1 + 2x^2 + 3x^3 + 4x^4 + 5x^5 + ... + nx^n = \frac{(x - (n+1)x^{n+1} + nx^{n+2})}{(x-1)^2}$ two adjacent nodes have same color

Complete graph =
$$k(k-1)(k-2)...(k-n+1)$$

Tree = $k(k-1)^{n-1}$

Cycle = $(k-1)^n + (-1)^n(k-1)$

- Number of trees with n labeled nodes: n^{n-2}

Lucas Number Number of edge cover of a cycle graph C_n is L_n

L(n) = L(n-1) + L(n-2); L(0) = 2, L(1) = 1

Catalan Number

$$C_{n+1} = C_0C_n + C_1C_{n-1} + C_2C_{n-2} + ... + C_nC_0$$

$$C_n = \binom{2n}{n} - \binom{2n}{n+1}$$

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

Derangement

$$D_n = (n-1)(D_{n-1} + D_{n-2}) = nD_{n-1} + (-1)^n$$
$$D_0 = 1.D_1 = 0$$

1,0,1,2,9,44,265,...

Ballot Theorem

Suppose that in an election, candidate A receives a votes and candidate B receives b votes, where a kb for some positive integer k. Compute the number of ways the ballots can be ordered so that A maintains more than k times as many votes as B throughout the counting of the ballots.

The solution to the ballot problem is ((a kb)/(a+b)) * C(a+b, a)**Classical Problem** F(n,k) = number of ways to color n objects using exactly k colors

Let G(n,k) be the number of ways to color n objects using no more than k colors. Then, F(n,k) = G(n,k) - C(k,1) * G(n,k-1) + C(k,2) * G(n,k-2) - G(n,k-1) + G(n,k-1) +|C(k,3)*G(n,k-3)...

Determining G(n, k):

Suppose, we are given a 1 * n grid. Any two adjacent cells can not have same color. Then, $G(n, k) = k * ((k-1)^{(n-1)})$ If no such condition on adjacent cells. Then, $G(n, k) = k^n$

Generating Function

 $1/(1-x) = 1 + x + x^2 + x^3 + \dots + 1/(1-ax) = 1 + ax + (ax)^2 + (ax)^3 + \dots$ $1/(1-x)^2 = 1 + 2x + 3x^2 + 4x^3 + \dots$ $1/(1-x)^3 = C(2.2) + C(3.2)x + \dots$ $(k,k)(ax)^2 + C(3+k,k)(ax)^3 + \dots + x(x+1)(1-x)^3 = 1+x+4x^2+9x^3 + \dots$

 $S(n,p) = \frac{1}{n+1} [(n+1)^{p+1} - 1 - \sum_{i=0}^{p-1} {p+1 \choose i} S(n,i)]$

 $1.2 + 2.3 + 3.4 + \dots = \frac{1}{3}n(n+1)(n+2)$

Probability

$$P(A|B) = \frac{P(A \cap B)}{P(B)}$$

$$P(A|B) = \frac{P(B|A)P(A)}{P(B)}$$

Matching Formula Normal Graph

|MM + MEC = n (without isolated vertex)

 $\frac{\text{IS + VC = G}}{\text{MaxIS + MVC = G}}$

Bipartite Graph

 $\begin{array}{l} MaxIS = n - MBM \\ MVC = MBM \\ MEC = n - MBM \end{array}$

Solution of $x^2 \equiv a \pmod{p}$:

 $ca \equiv cb \pmod{m} \iff a \equiv b \pmod{\frac{n}{\gcd(n,c)}}$

 $ax \equiv b \pmod{m}$ has a solution $\iff \gcd(a, m)|b|$

- If $ax \equiv b \pmod{m}$ has a solution, then it has gcd(a,m) solutions and they are separated by $\frac{n}{\gcd(n,m)}$

 $-ax \equiv 1 \pmod{m}$ has a solution or a is invertible $\pmod{m} \iff$ gcd(a, m) = 1

 $-x^2 \equiv 1 \pmod{p}$ then $x \equiv \pm 1 \pmod{p}$

There are $\frac{p-1}{2}$ has no solution.

- There are $\frac{p-1}{2}$ has exactty two solutions.

When p%4=3, $x\equiv\pm a^{\frac{p+1}{4}}$

- When p%8 = 5, $x \equiv a^{\frac{p+3}{8}}$ or $x \equiv 2^{\frac{p-1}{4}}a^{\frac{p+3}{8}}$

Totient

- If p is a prime $(p^k) = p^k - p^{k-1}$

- If a b are relatively prime, $\phi(ab) = \phi(a)\phi(b)$

 $-\phi(n) = n(1 - \frac{1}{p_1})(1 - \frac{1}{p_2})(1 - \frac{1}{p_3})...(1 - \frac{1}{p_b})$

Sum of coprime to $n = n * \frac{\phi(n)}{2}$

- If $n = 2^k$, $\phi(n) = 2^{k-1} = \frac{n}{2}$

For a b, $\phi(ab) = \phi(a)\phi(b)\frac{d}{\phi(d)}$

 $\phi(ip) = p\phi(i)$ whenever p is a prime and it divides i

- The number of $a(1 \le a \le N)$ such that gcd(a, N) = d is $\phi(\frac{n}{d})$

- If n > 2. $\phi(n)$ is always even

- Sum of gcd, $\sum_{i=1}^{n} gcd(i,n) = \sum_{d|n} d\phi(\frac{n}{d})$

- Sum of lcm, $\sum_{i=1}^{l-1} nlcm(i,n) = \frac{n}{2} (\sum_{d|n} (d\phi(d)) + 1)$

 $\phi(1) = 1$ and $\phi(2) = 1$ which two are only odd ϕ

 $-\phi(3)=2$ and $\phi(4)=2$ and $\phi(6)=2$ which three are only prime

- Find minimum n such that $\frac{\phi(n)}{n}$ is maximum- Multiple of small primes- 2*3*5*7*11*13*...

Mobius

$$\sum_{i=1}^{n} \sum_{j=1}^{n} [gcd(i,j) = 1] = \sum_{k=1}^{n} \mu(k) \lfloor \frac{n}{k} \rfloor^{2}$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = \sum_{k=1}^{n} k \sum_{l=1}^{\lfloor \frac{n}{k} \rfloor} \mu(l) \lfloor \frac{n}{kl} \rfloor^{2}$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = \sum_{k=1}^{n} \left(\frac{\lfloor \frac{n}{k} \rfloor (1 + \lfloor \frac{n}{k} \rfloor)}{2}\right)^{2} \sum_{d \mid k} \mu(d)kd$$

Tree Hashing

 $f(u) = sz[u] * \sum_{i=0}^{\infty} f(v) * p^{i}$; f(v) are sorted f(child) = 1