

# Do Rust’s compile-time type checks produce equivalent assembly to C while preventing state management errors?

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**Abstract**—Memory safety vulnerabilities—use-after-free, buffer overflows, out-of-bounds write—are at the top of known exploited vulnerabilities[1]. System programming languages like C, which allow for low-level memory control, are needed to write performance critical code. State is the current condition of an object that determines what operations are valid. C cannot enforce valid state transitions, such as reading an open file, at compile-time. Operations need to be checked explicitly at runtime if they are in a valid state before running. Higher-level languages like Java automate detecting potentially unsafe operations, preventing errors but degrading performance. Therefore, developers have to make a tradeoff between safety and speed.

Rust has the ability to move certain behaviors to compile-time execution or analysis. This approach claims to be able to verify state validity during compilation, which could lead to preventing errors that runtime checks would detect while achieving comparable C performance. This work tests whether Rust’s compile-time guarantees produce assembly with equivalent performance to C code that omits all safety checks.

Programs must track state explicitly or implicitly. A file handle must be opened before reading; once closed, reads must fail. Verifying these transitions—the core of a state machine—requires verification logic. This work implements three versions of a file handle state machine, isolating state management from I/O overhead:

- 1) **Defensive C:** Uses an enum to track state, with explicit validation checks before each operation. Safe, but includes runtime conditional branches.
- 2) **Minimal C:** Tracks state with an enum, but omits all validation. Fast but permits invalid operations to compile.
- 3) **Rust:** Encodes each state as a distinct type, making invalid transitions not compile.

The resulting assembly will be compared using total instruction count, conditional branches, and state-tracking overhead.

The deliverables are side-by-side assembly comparisons showing whether Rust eliminates defensive checks present in safe C and an analysis with the above mentioned assembly. The assembly comparison will demonstrate whether encoding state in the type system eliminates defensive branches while matching minimal C’s instruction count.

**Index Terms**—component, formatting, style, styling, insert.

## I. INTRODUCTION

This document is a model and instructions for L<sup>A</sup>T<sub>E</sub>X. Please observe the conference page limits. For more information

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## REFERENCES

- [1] “CWE - 2024 CWE top 10 KEV weaknesses,” CWE - 2024 CWE Top 10 KEV Weaknesses, Accessed: Dec. 10, 2025. [Online]. Available: [https://cwe.mitre.org/top25/archive/2024/2024\\_kev\\_list.html](https://cwe.mitre.org/top25/archive/2024/2024_kev_list.html).