

NPTEL MOOC

PROGRAMMING, DATA STRUCTURES AND ALGORITHMS IN PYTHON

Week 4, Lecture 3

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Merge Sort: Shortcomings

- * Merging **A** and **B** creates a new array **C**
 - * No obvious way to efficiently merge in place
- * Extra storage can be costly
- * Inherently recursive
 - * Recursive call and return are expensive

Alternative approach

- * Extra space is required to merge
- * Merging happens because elements in left half must move right and vice versa
- * Can we divide so that everything to the left is smaller than everything to the right?
 - * No need to merge!

Divide and conquer without merging

- * Suppose the median value in A is m
- * Move all values $\leq m$ to left half of A
 - * Right half has values $> m$
 - * This shifting can be done in place, in time $O(n)$
- * Recursively sort left and right halves
- * A is now sorted! No need to merge
 - * $T(n) = 2T(n/2) + n = O(n \log n)$

Divide and conquer without merging

- * How do we find the median?
 - * Sort and pick up middle element
 - * But our aim is to sort!
- * Instead, pick up some value in **A** — **pivot**
 - * Split **A** with respect to this pivot element

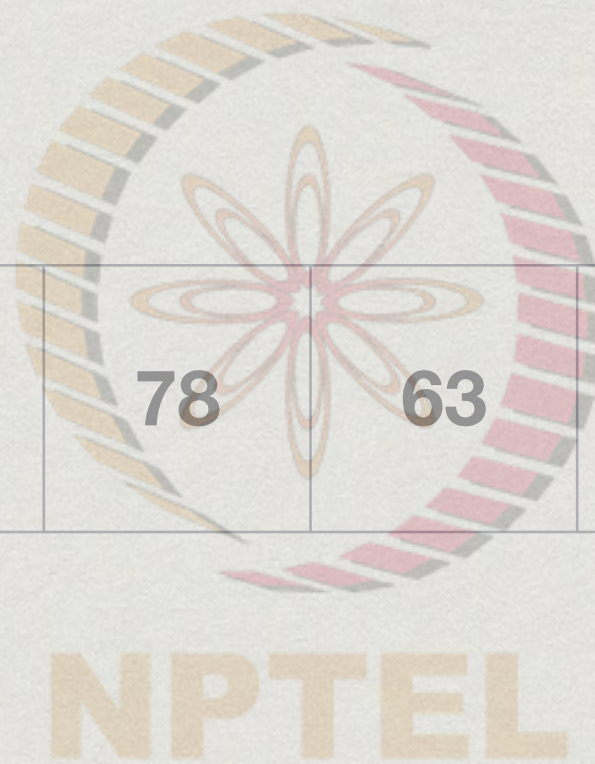
Quicksort

- * Choose a pivot element
 - * Typically the first value in the array
- * Partition **A** into lower and upper parts with respect to pivot
- * Move pivot between lower and upper partition
- * Recursively sort the two partitions

Quicksort

- * High level view

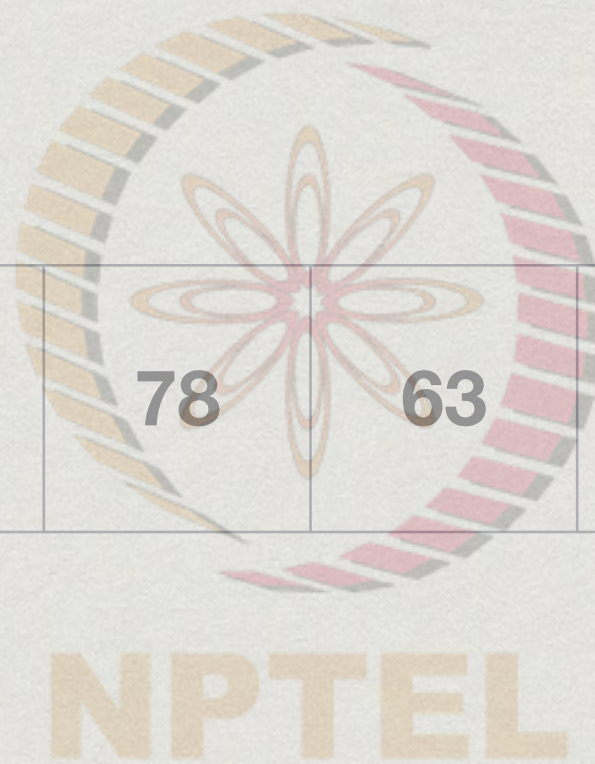
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Quicksort

- * High level view

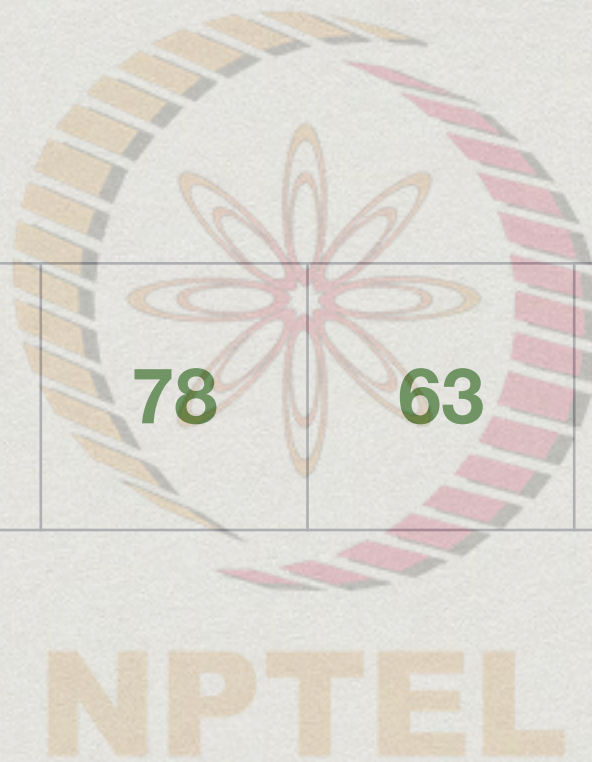
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Quicksort

- * High level view

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Quicksort

- * High level view

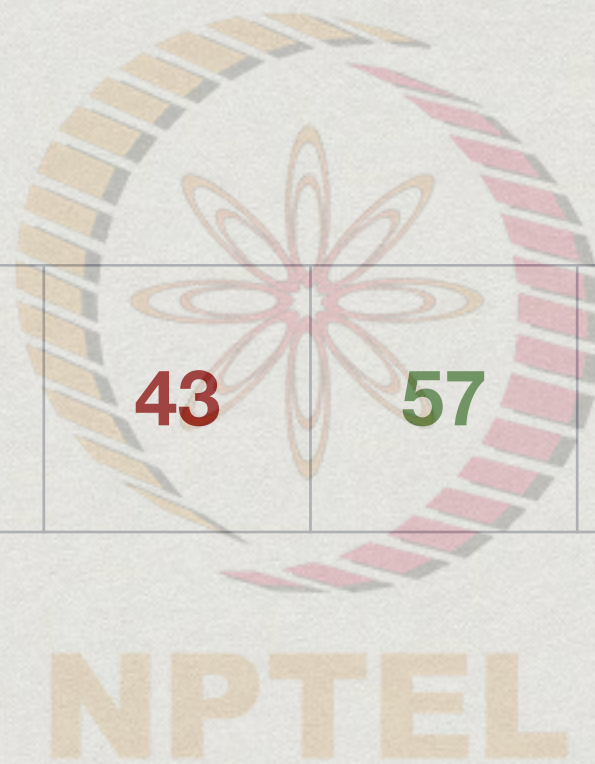
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Quicksort

- * High level view

13	22	32	43	57	63	78	91
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Quicksort: Partitioning

43	32	22	78	63	57	91	13
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Quicksort: Partitioning

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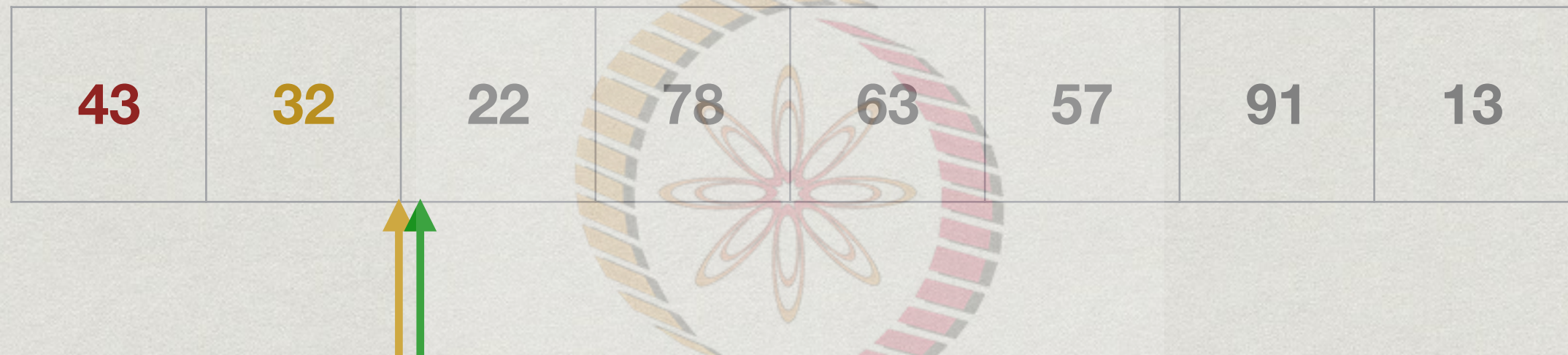


Quicksort: Partitioning

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Quicksort: Partitioning



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Quicksort: Partitioning

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Quicksort: Partitioning

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Quicksort: Partitioning

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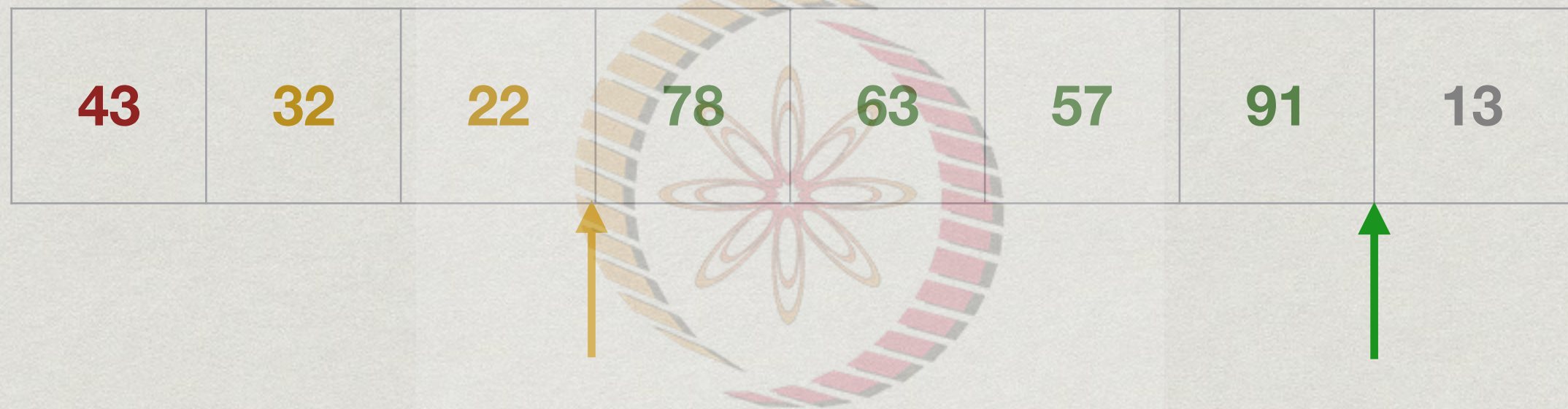
Quicksort: Partitioning

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Quicksort: Partitioning



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Quicksort: Partitioning

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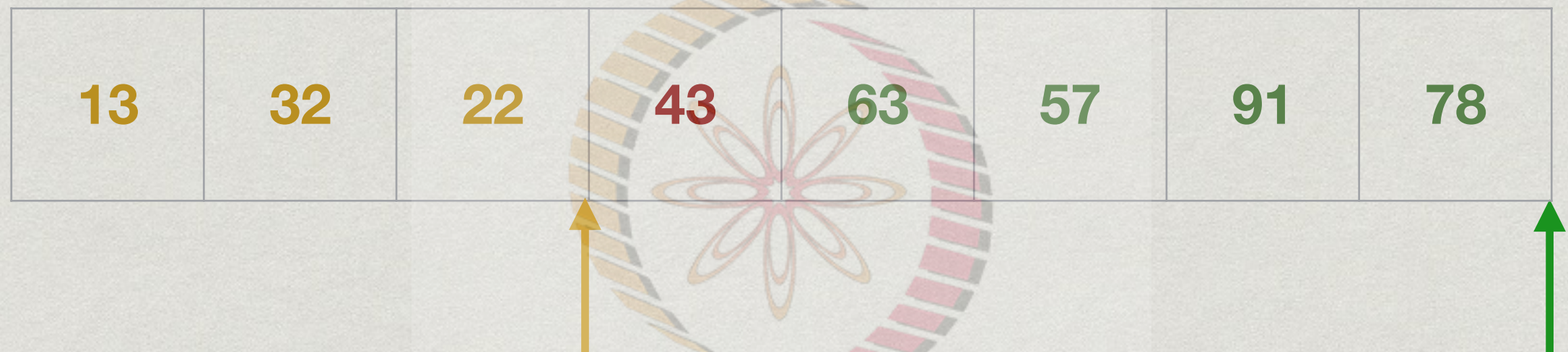
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Quicksort: Partitioning



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Quicksort: Partitioning



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Quicksort in Python

```
def Quicksort(A,l,r): # Sort A[l:r]
    if r - l <= 1: # Base case
        return ()

    # Partition with respect to pivot, a[l]
    yellow = l+1

    for green in range(l+1,r):
        if A[green] <= A[l]:
            (A[yellow],A[green]) = (A[green],A[yellow])
            yellow = yellow + 1

    # Move pivot into place
    (A[l],A[yellow-1]) = (A[yellow-1],A[l])

    Quicksort(A,l,yellow-1) # Recursive calls
    Quicksort(A,yellow,r)
```