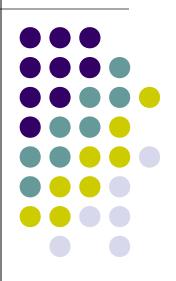
Cipher Techniques



Road Map

- Basic Terminology
- Cryptosystem
- Classical Cryptography
- Algorithm Types and Modes
- Data Encryption Standard
- Other Stream & Block Ciphers

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- plaintext the original message
- ciphertext the coded message
- cipher algorithm for transforming plaintext to ciphertext
- key info used in cipher known only to sender/receiver
- encipher (encrypt) converting plaintext to ciphertext
- decipher (decrypt) recovering ciphertext from plaintext
- cryptography study of encryption principles/methods
- cryptanalysis (codebreaking) the study of principles/ methods of deciphering ciphertext without knowing key
- cryptology the field of both cryptography and cryptanalysis

Cryptosystem



- A *cryptosystem* is a five-tuple $(\mathcal{P}, \mathcal{C}, \mathcal{K}, \mathcal{E}, \mathcal{D})$, where the following are satisfied:
- p is a finite set of possible plaintexts.
- 2. \mathcal{C} is a finite set of possible *ciphertexts*.
- the key space, is a finite set of possible keys
- 4. $\forall K \in \mathcal{K}, \exists E_{\kappa} \in \mathcal{E} \text{ (encryption rule)}, \exists D_{\kappa} \in \mathcal{D}$ (decryption rule). Each $E_{\kappa} \colon \mathcal{P} \rightarrow \mathcal{C}$ and $D_{\kappa} \colon \mathcal{C} \rightarrow \mathcal{P}$ are functions such that $\forall x \in \mathcal{P}, D_{\kappa}(E_{\kappa}(x)) = x$.

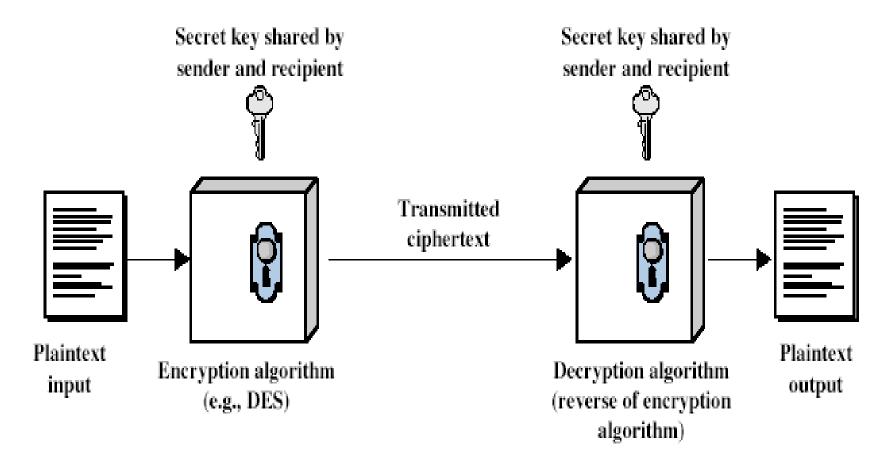




- Cryptography
 - Symmetric / private key / single key
 - Asymmetric / public-key / two key



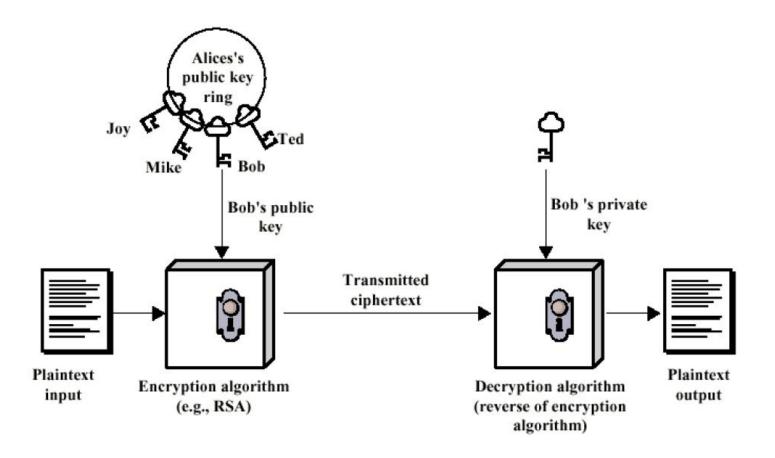




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Requirements



- Two requirements for secure use of symmetric encryption:
 - a strong encryption algorithm
 - a secret key known only to sender / receiver
 Y = E_k(X)

$$X = D_{\kappa}(Y)$$

- assume encryption algorithm is known
- implies a secure channel to distribute key





- Transposition Techniques
- Substitution techniques
 - Caesar Cipher
 - Monoalphabetic Cipher
 - Polyalphabethic Cipher
 - Playfair Cipher

Types of Cryptanalytic Attacks



adversary needs strongest attack

ciphertext only

only know algorithm / ciphertext, statistical, can identify plaintext, or worse: the key

known plaintext

know/suspect plaintext & ciphertext to attack cipher

chosen plaintext

select plaintext and obtain ciphertext to attack cipher

chosen ciphertext

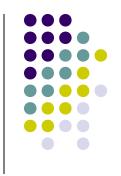
select ciphertext and obtain plaintext to attack cipher

chosen text

• select either plaintext or ciphertext to en/decrypt 18

adversary's attacks can be weaker





- always possible to simply try every key
- most basic attack, proportional to size of key space
- assume either know / recognise plaintext

Key Size (bits)	Number of Alternative Keys	Time required at 1 encryption/µs	Time required at 10 ⁶ encryptions/ <i>µ</i> s
32	$2^{32} = 4.3 \times 10^9$	$2^{31} \mu s = 35.8 \text{ minutes}$	2.15 milliseconds
56	$2^{56} = 7.2 \times 10^{16}$	$2^{55} \mu s = 1142 \text{ years}$	10.01 hours
128	$2^{128} = 3.4 \times 10^{38}$	$2^{127} \mu s = 5.4 \times 10^{24} \text{ years}$	5.4×10^{18} years
168	$2^{168} = 3.7 \times 10^{50}$	$2^{167} \mu s = 5.9 \times 10^{36} \text{ years}$	5.9×10^{30} years
26 characters (permutation)	$26! = 4 \times 10^{26}$	$2 \times 10^{26} \mu \mathrm{s} = 6.4 \times 10^{12} \mathrm{years}$	6.4×10^6 years

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- Consider classical transposition or permutation ciphers
- these hide the message by rearranging the letter order
- without altering the actual letters used
- can recognise these since have the same frequency distribution as the original text

Rail Fence cipher



- write message letters out diagonally over a number of rows
- then read off cipher row by row
- eg. write message out as:

```
m e m a t r h t g p r y e t e f e t e o a a t
```

giving ciphertext

MEMATRHTGPRYETEFETEOAAT





- a more complex scheme
- write letters of message out in rows over a specified number of columns
- then reorder the columns according to some key before reading off the rows

```
      Key:
      4
      3
      1
      2
      5
      6
      7

      Plaintext:
      a
      t
      t
      a
      c
      k
      p

      o
      s
      t
      p
      o
      n
      e

      d
      u
      n
      t
      i
      t
      t

      w
      o
      a
      m
      x
      y
      z
```

Ciphertext: TTNAAPTMTSUOAODWCOIXKNLYPETZ





- where letters of plaintext are replaced by other letters or by numbers or symbols
- or if plaintext is viewed as a sequence of bits, then substitution involves replacing plaintext bit patterns with ciphertext bit patterns





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- earliest known substitution cipher
- by Julius Caesar
- first attested use in military affairs
- replaces each letter by 3rd letter after it
- example:

```
meet me after the toga party PHHW PH DIWHU WKH WRJD SDUWB
```

Caesar Cipher



can define transformation as:

```
abcdefghijklmnopqrstuvwxyz
DEFGHIJKLMNOPORSTUVWXYZABC
```

mathematically give each letter a number

```
abcdefghijklm
0 1 2 3 4 5 6 7 8 9 10 11 12
n opqrstuvwxyZ
13 14 15 16 17 18 19 20 21 22 23 24 25
```

• then have Caesar cipher as:

$$C = E(p) = (p + k) \mod (26)$$

 $p = D(C) = (C - k) \mod (26)$





- rather than just shifting the alphabet
- could shuffle (jumble) the letters arbitrarily
- each plaintext letter maps to a different random ciphertext letter
- hence key is 26 letters long

Plain: abcdefghijklmnopqrstuvwxyz

Cipher: DKVQFIBJWPESCXHTMYAUOLRGZN

Plaintext: ifwewishtoreplaceletters

Ciphertext: WIRFRWAJUHYFTSDVFSFUUFYA

Playfair Cipher



- not even the large number of keys in a monoalphabetic cipher provides security
- one approach to improving security was to encrypt multiple letters
- the Playfair Cipher is an example
- invented by Charles Wheatstone in 1854, but named after his friend Baron Playfair





- a 5X5 matrix of letters based on a keyword
 - (I and J aren't distinguished)
- fill in letters of keyword (sans duplicates)
- fill rest of matrix with other letters
- eg. using the keyword MONARCHY

```
MONAR
```

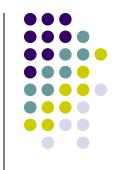
CHYBD

EFGIK

LPQST

UVWXZ

Encrypting and Decrypting



- plaintext encrypted two letters at a time:
 - each letter is replaced by the one in its row in the column of the other letter of the pair, eg. "hs" encrypts to "BP", and "ea" to "IM" or "JM" (as desired). Except when that doesn't work!
 - if a pair is a repeated letter, insert a filler like 'X', eg. "balloon" transformed to "ba lx lo on"
 - if both letters fall in the same row, replace each with letter to right (wrapping back to start from end), eg. "ar" encrypts as "RM"
 - if both letters fall in the same column, replace each with the letter below it (again wrapping to top from bottom), eg. "mu" encrypts to "CM"

Polyalphabetic Ciphers



- another approach to improving security is to use multiple cipher alphabets
- called polyalphabetic substitution ciphers
- makes cryptanalysis harder with more alphabets to guess and flatter frequency distribution
- use a key to select which alphabet is used for each letter of the message
- use each alphabet in turn
- repeat from start after end of key is reached





- simplest polyalphabetic substitution cipher is the Vigenère Cipher
- effectively multiple caesar ciphers
- key is multiple letters long K = k1 k2 ... kd
- ith letter specifies ith alphabet to use
- use each alphabet in turn
- repeat from start after d letters in message
- decryption simply works in reverse

Example



- write the plaintext out
- write the keyword repeated above it
- use each key letter as a caesar cipher key
- encrypt the corresponding plaintext letter
- eg using keyword deceptive

```
key: deceptivedeceptive
plaintext: wearediscoveredsaveyourself
ciphertext:ZICVTWQNGRZGVTWAVZHCQYGLMGJ
```

Autokey Cipher



- ideally want a key as long as the message
- Vigenère proposed the autokey cipher
- with keyword is prefixed to message as key
- knowing keyword can recover the first few letters
- use these in turn on the rest of the message
- but still have frequency characteristics to attack
- eg. given key deceptive

```
key: deceptivewearediscoveredsav plaintext: wearediscoveredsaveyourself ciphertext: ZICVTWQNGKZEIIGASXSTSLVVWLA
```