

# Tuple Relational Calculus

Based on first order logic

- Declarative language. (non-procedural)

- Use tuple variables:  $t, u, \dots$

- Use  $t[A]$  to denote value of attribute  $A$  in tuple  $t$ .

Query specified as

$$\{t \mid P(t)\}$$

where  $P$  is

a predicate and  $t$  a tuple variable.

$t$  is called an output tuple.

Consider:

accounts (acct#, bal, br-name)

customer (cust-name, street, cust-city)

deposit (cust-name, acct#)

branch (br-name, assets, br-city)

loans (loan#, amount, br-name)

borrow (cust-name, <sup>loan</sup>~~acct~~#)

Examples:

1 —  $\{ t \mid t \in \text{accounts} \wedge t[\text{balance}] > 1000 \}$   
 - all account tuples with balance > 1000

2 - Retrieve names of customers having an account with balance > 1000.

Use existential quantifiers.

$$\{ t \mid \exists u \in \text{deposit} \exists v \in \text{accounts} \\ ( u[\text{acct-}\cancel{\text{nbr}}] = v[\text{acct\#}] \wedge \\ v[\text{balance}] > 1000 \wedge \\ t[\text{cust-name}] = \cancel{u[\text{cust-name}]} \\ u[\text{cust-name}] \}$$

$t$  is the output tuple.

It has only one attribute cust-name

3. Retrieve names of customers and their cities, for those customers having an account with bal > 1000.

$\{ t \mid \exists c \in \text{Customer}, \exists u \in \text{deposit}, \exists v \in \text{accounts}$   
 $( t[\text{cust-name}] = c[\text{cust-name}] \wedge t[\text{cust-city}] = c[\text{cust-city}] \wedge$

$u[\text{cust-name}] = u[\text{cust-name}] \wedge$

$u[\text{acct\#}] = v[\text{acct\#}] \wedge$

$v[\text{balance}] > 1000 \}.$

$t$  has more than one attribute - cust-name, cust-city

4. Find names of customers having an account or a loan.

$\{ t \mid \exists u \in \text{deposit} ( t[\text{cust-name}] = u[\text{cust-name}]$   
 $\vee \exists v \in \text{borrow} ( t[\text{cust-name}] = v[\text{cust-name}] )$

( we use  $\vee$  operator )

5. Find names of customers having an account but not a loan

$$\{ t \mid \exists u \in \text{deposit} (u[\text{cust-name}] = t[\text{cust-name}]) \\ \wedge \neg (\exists v \in \text{borrow} (v[\text{cust-name}] = t[\text{cust-name}])) \}$$

( use '¬' operator )

6. ( use  $\Rightarrow$  )  
remember  $p \Rightarrow q \equiv \neg p \vee q$

Find names of customers having an account in every branch located in Chicago.

$$\{ t \mid \forall b \in \text{branch} ( b[\text{br-city}] = \text{'Chicago'} \Rightarrow \\ ( \exists u \in \text{deposit} \exists v \in \text{accounts} \\ ( v[\text{br-name}] = b[\text{br-name}] \wedge \\ u[\text{acct\#}] = v[\text{acct\#}] \wedge \\ u[\text{cust-name}] = t[\text{cust-name}] ) ) ) \}$$

Query -  $\{t \mid P(t)\}$

$P(t)$  is a formula built from atoms of the form:

$s \in r$  , (  $s$  is a tuple variable,  $r$  is a relation )

$s[x] \theta u[y]$  (  $s, u$  - tuple variables

$s[x] \theta c$

$x, y$  - attributes,

$\theta$  - comparison )

$c$  - Constant .

Complex formula  $P$  is

- an atom

or is of the form

$\neg P_1$  ,  $P_1 \vee P_2$  ,  $P_1 \wedge P_2$  ,  $P_1 \Rightarrow P_2$

$\exists s \in r (P_1(s))$  ,  $\forall s \in r (P_1(s))$

where  $P_1, P_2$  are formulas .