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NetVM: High Performance and Flexible Networking using Virtualization on Commodity Platforms

Jinho Hwang[†], K. K. Ramakrishnan*, Timothy Wood[†]

[†]The George Washington University

*Rutgers University

High Performance Networking

- Line Rate High Performance Machines

Special HW + SW
Very Expensive



PacketShader
[Han:Sigcomm:2010]
GPU



NetMap
[Rizzo:ATC:2012]
Commodity
Servers

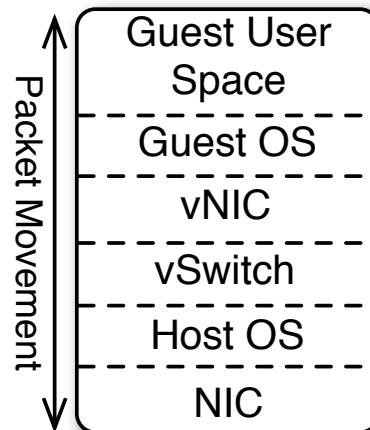


Intel DPDK
Commodity NICs



Data center is
virtualized

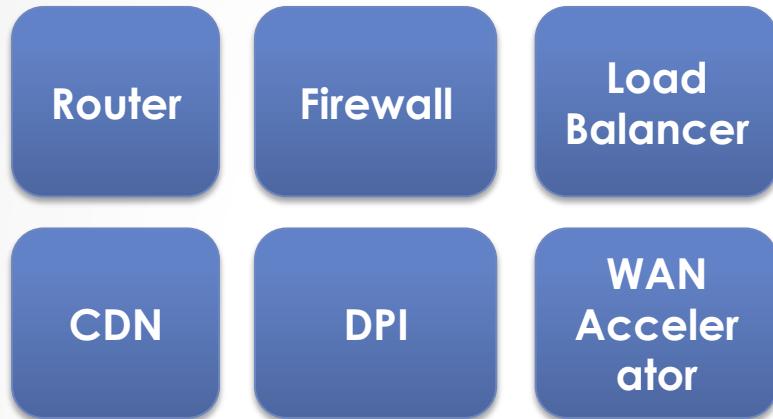
Performance in
Virtualized Platforms?



**Need to overcome
virtualization overheads**

Network Functions Virtualization (NFV)

Existing Network Functions



6 Machines x Power x \$HW x \$SW

NFV-enabled Network Functions

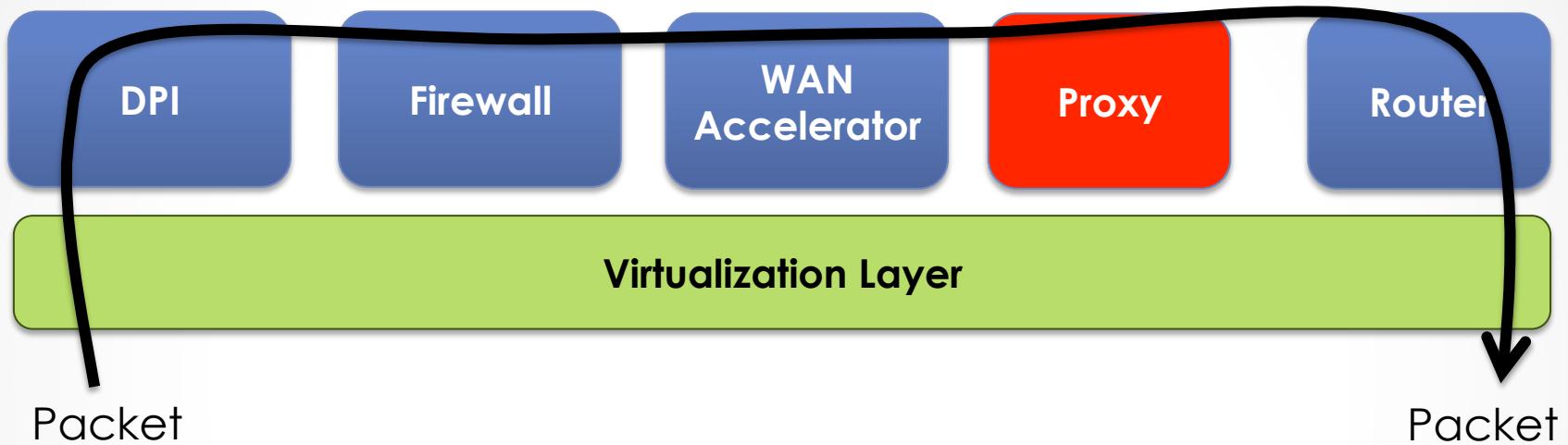


1 Machine x Power x \$SW
+ Flexible deployment

**Need to provide high speed
and low latency**

Chained Functionality

- Functions are often sequential



**Need High Speed
Inter-Function (VM) Communication**

So, What Has NetVM Done?

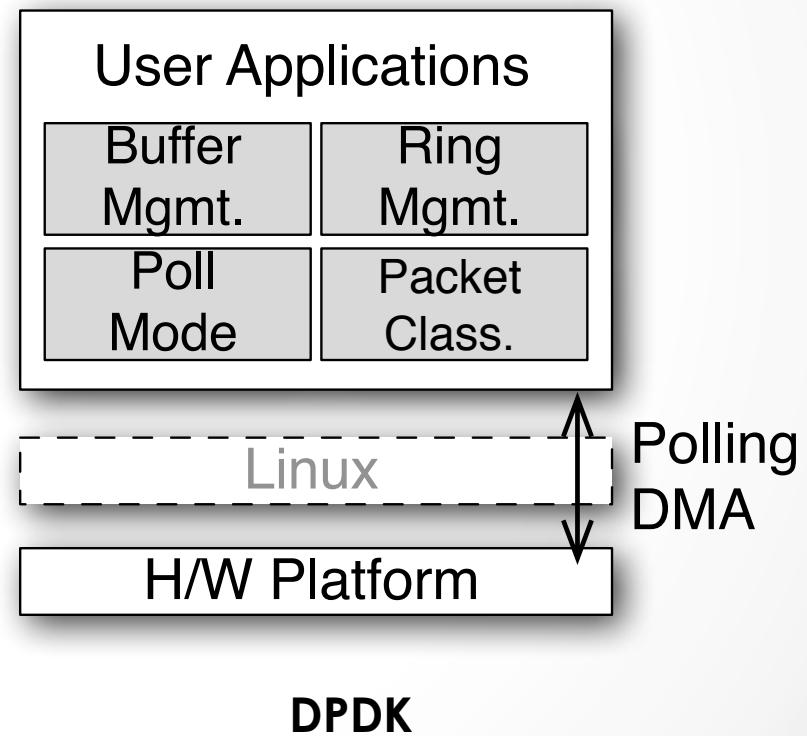
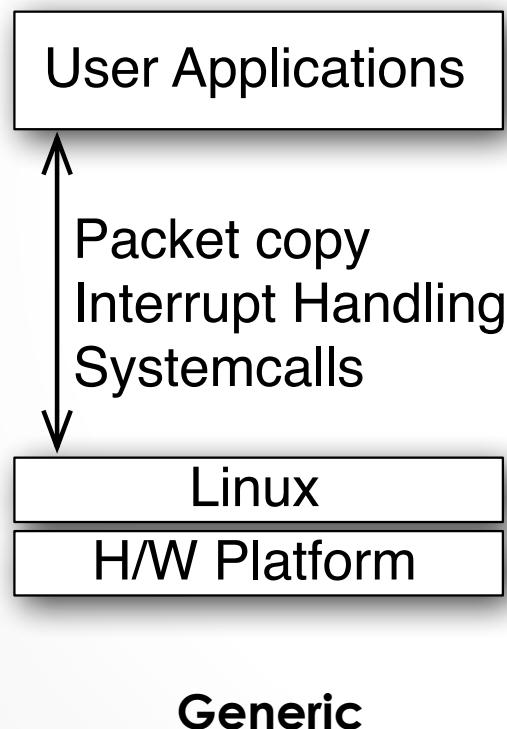
- NetVM is a platform for running complex network functionality at line-speed (10Gbps~) using commodity hardware
 - ✓ A virtualization-based high-speed packet delivery
 - ✓ Memory sharing framework for network data
 - ✓ A hypervisor-based switching
 - ✓ High speed inter-VM communication
 - ✓ Security domains

Outline

- Motivations & Contribution
- Background
- System Design
- Evaluation
- Conclusion & Future Work

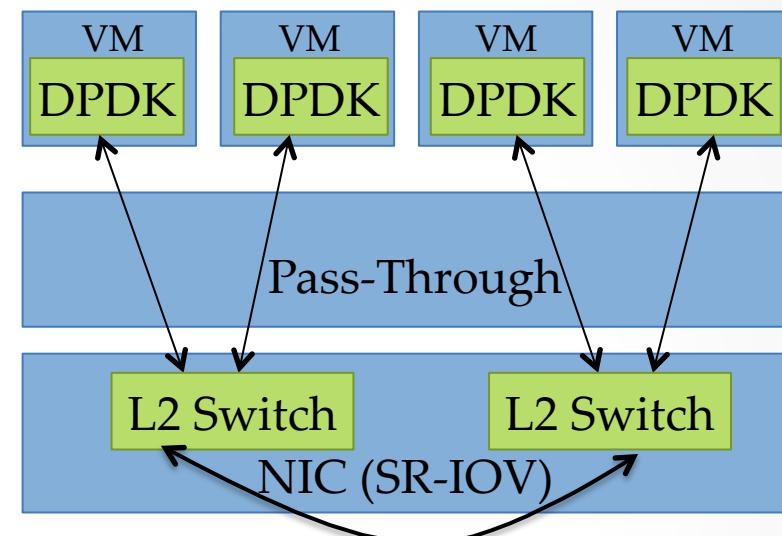
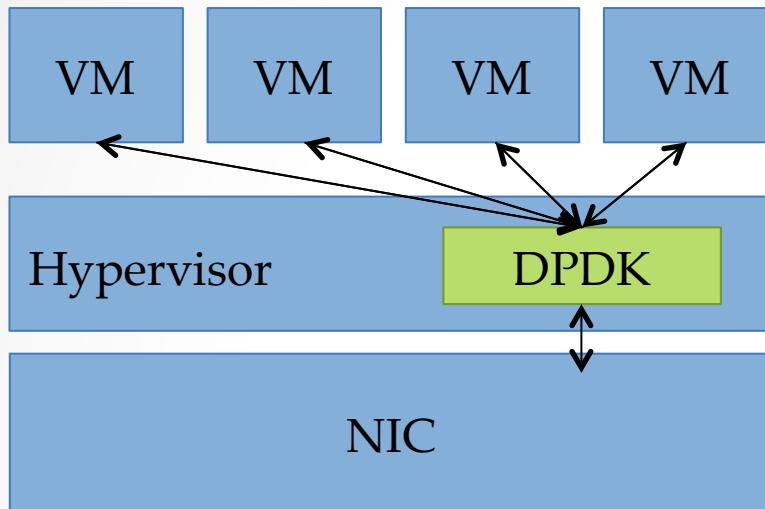
COTS and DPDK (Data Plane Development Kit)

- COTS = commercial, off-the-shelf servers
- DPDK is a set of optimized “**user space**” software libraries and drivers that can be used to accelerate packet processing on Intel architecture



DPDK with Virtualization

- Two Architectural Variations



SR-IOV = Single Root I/O Virtualization

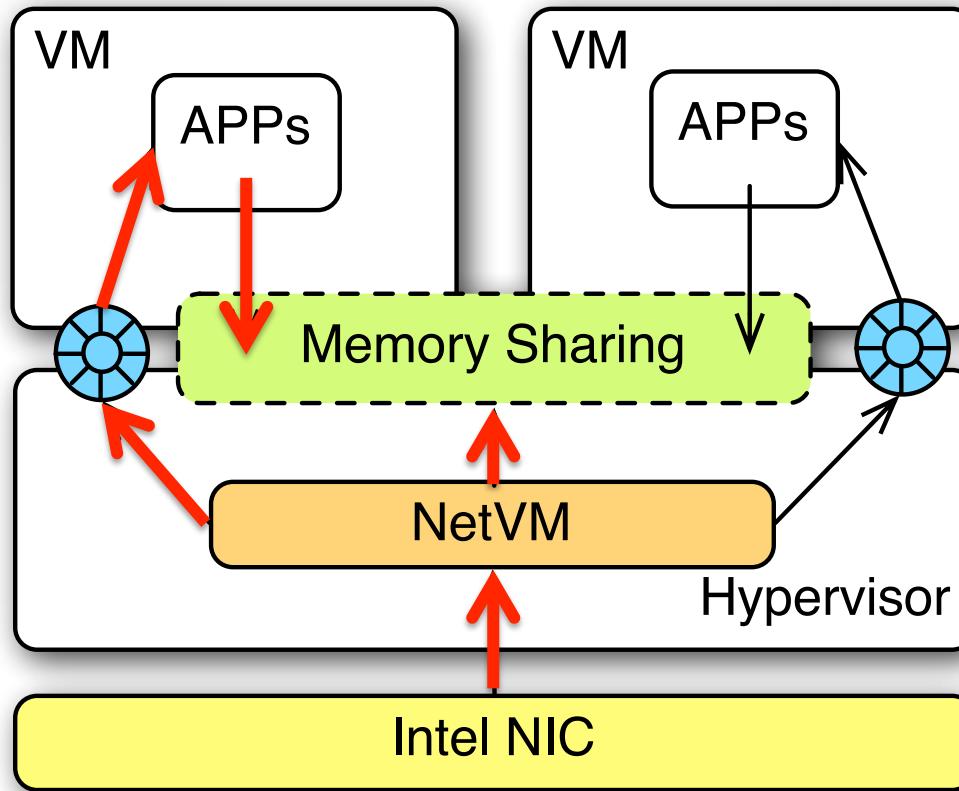
- Flexible(dynamic) Configuration
 - Control over packet switching
 - Control over load-balancing
 - Has more overhead
- Max 63 virtual functions (tx/rx)
 - Static configuration
 - Inter-VM switch is limited per port
 - No control over packet switching
 - No control over load-balancing

Neither can achieve full line-rate network speed in VMs

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NetVM System Overview



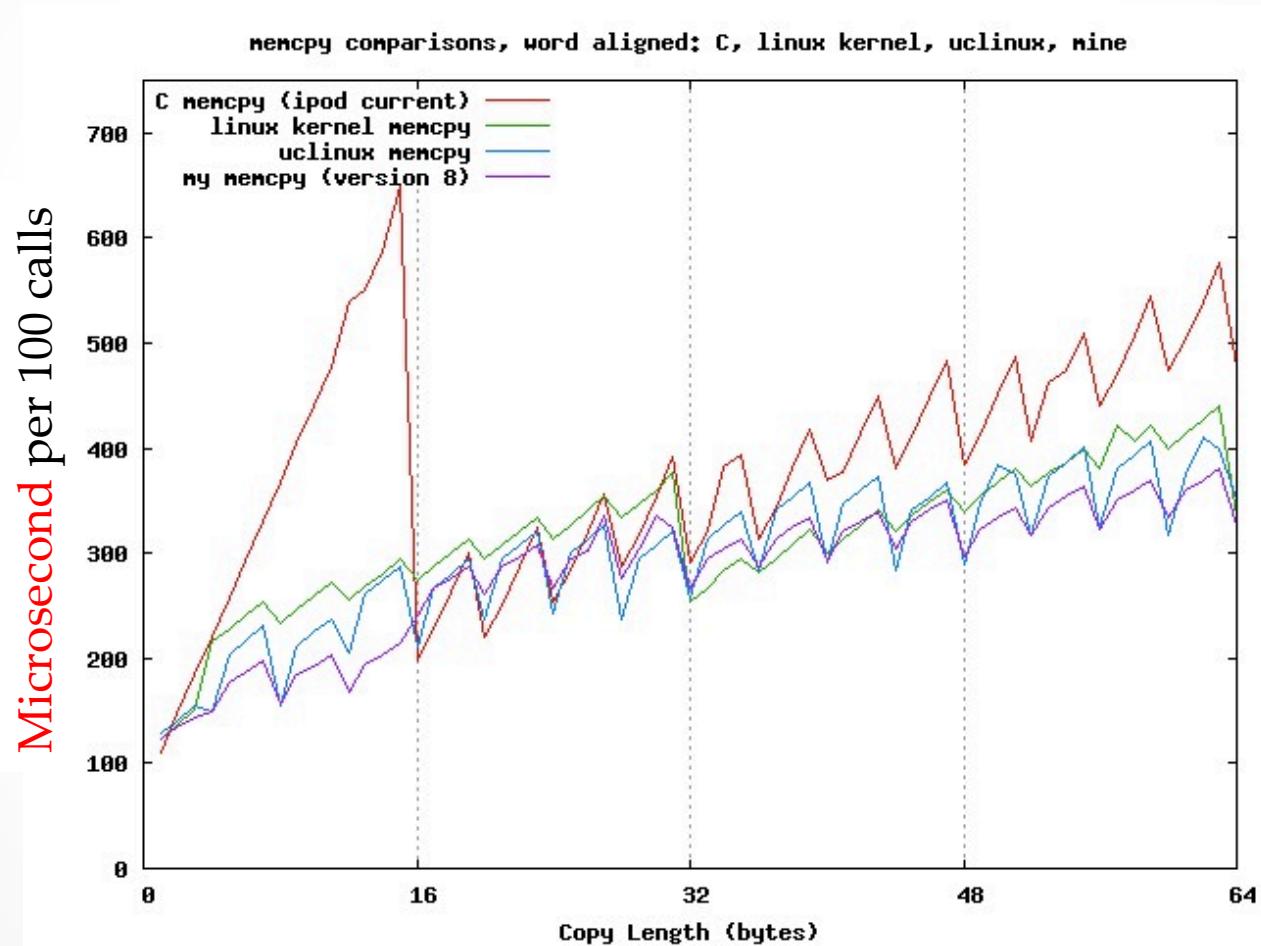
- NetVM (with DPDK) runs in hypervisor User Space
- For zero packet copy, memory is shared for network data
- Each VM has its own ring to receive/transmit a packet descriptor

System Design Challenges

1. Zero-copy
2. A huge-page sharing
3. Lockless & non-uniform memory architecture (NUMA) aware design
4. Security domains

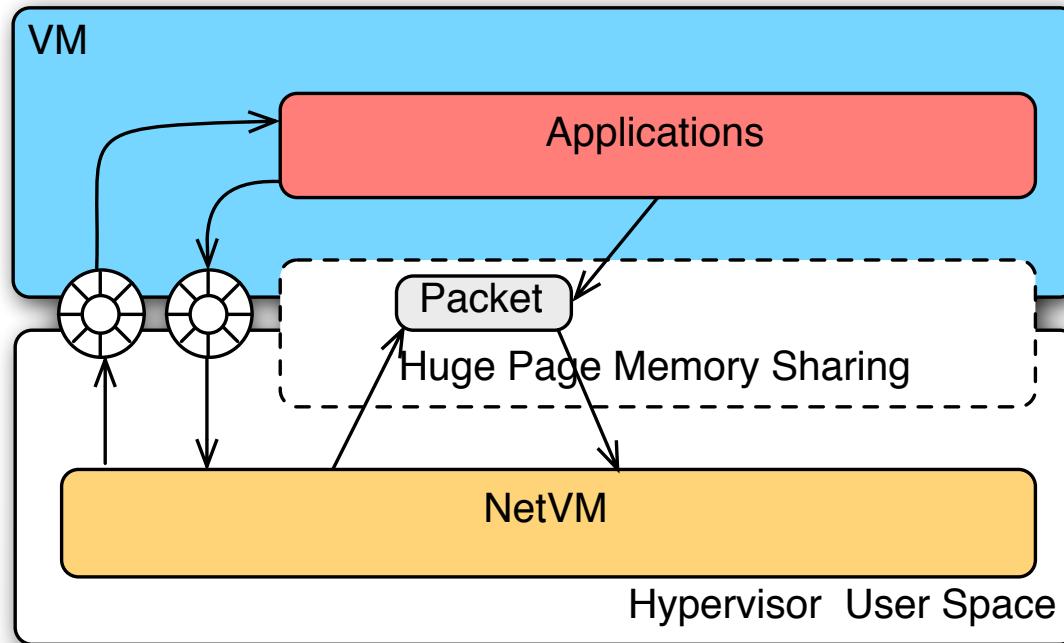
Memory Copy Hazard

- 10Gbps = 14,880,952 packets/s (64 byte)
- Each packet should be processed in 67 ns



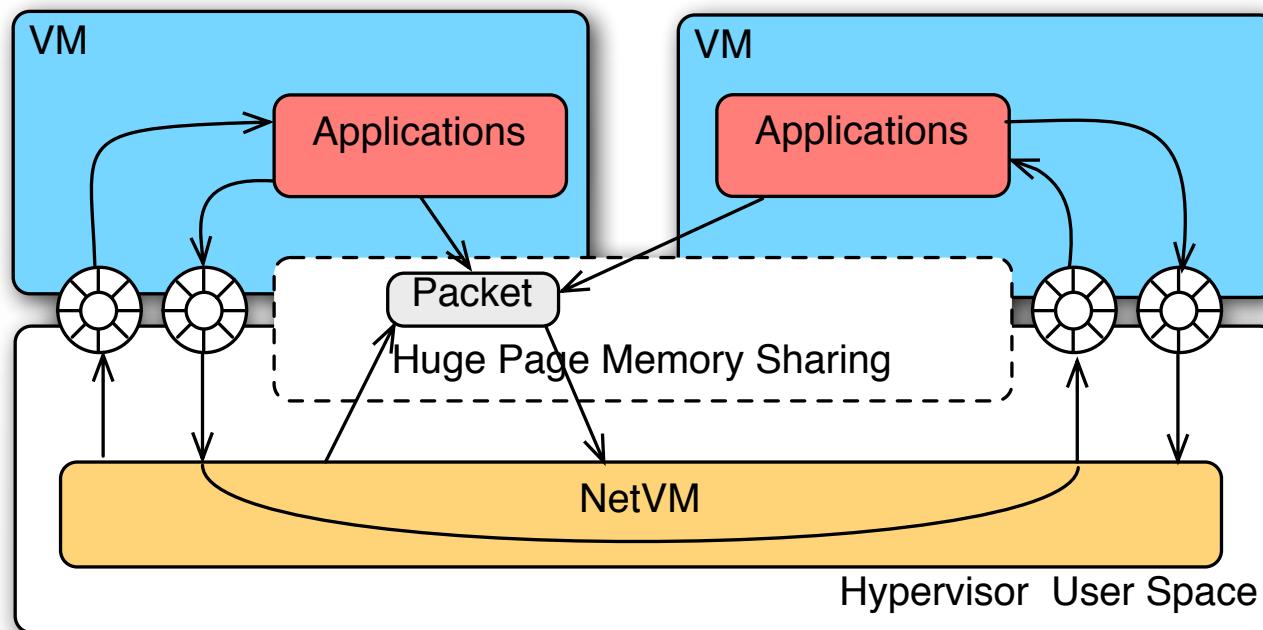
Zero-Copy Packet Delivery

- Packet directly DMAs into huge page memory
- Applications in VM receive references (location) via shared ring buffers
- Packet contents can be manipulated
- Applications decide an action: chain to another VM, send out, discard



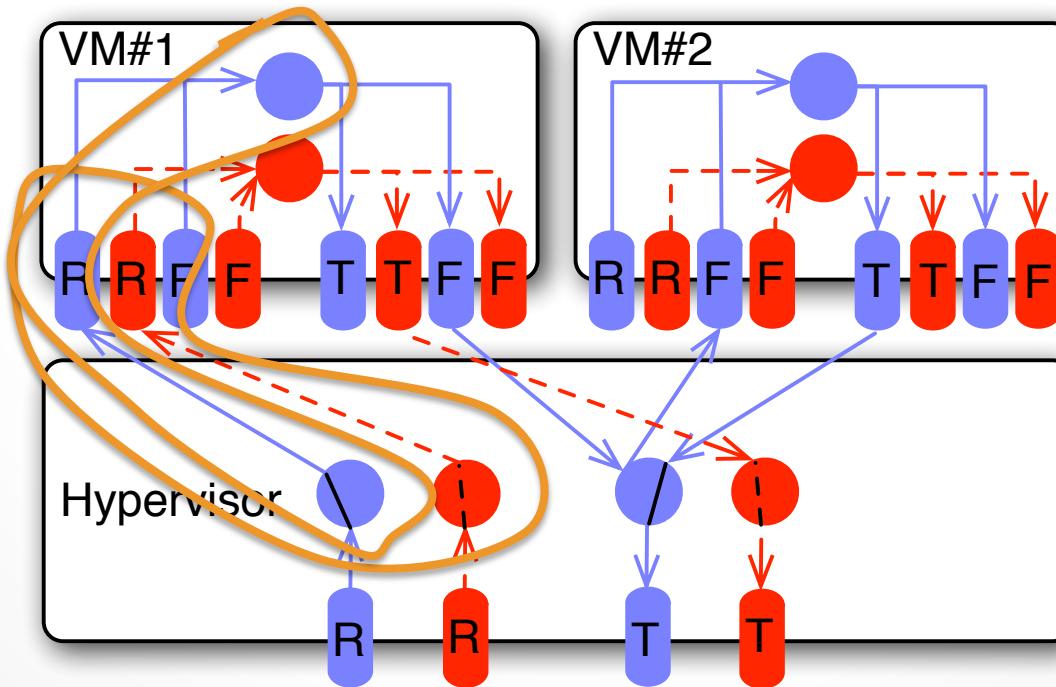
Chained Packet Delivery

- Packets are processed in a sequence/parallel
- Applications in VMs pass packet references to other VM
- Only one application can process a packet at a time (read/write)



Lockless & NUMA-Aware Design

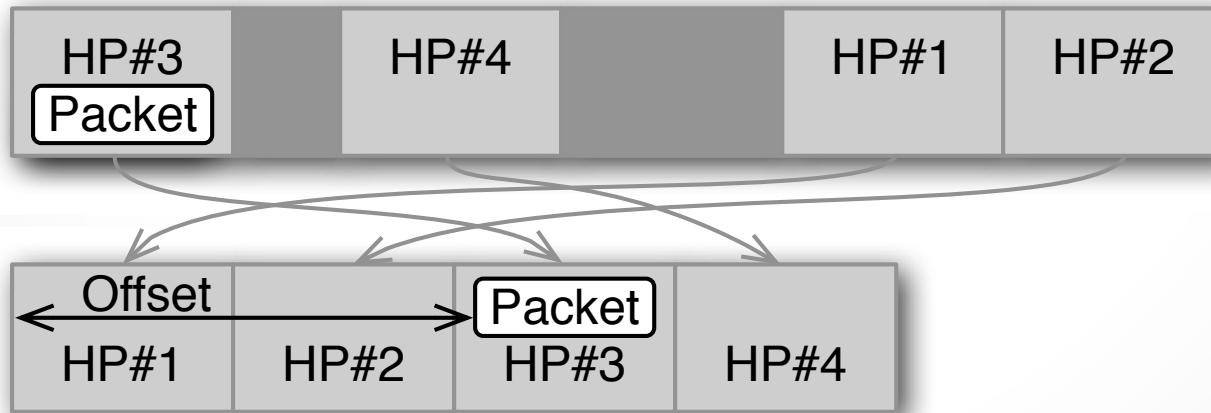
- **Locks** significantly degrade packet transmission performance
 - Core-queue matching (between horizontal cores) and data structure separation (between vertical cores)
- With **NUMA**, reading local and remote memory alternatively invalidates local cache → cache misses
 - This is fraught with Intel Direct I/O (Sandy-bridge)
 - Process in aligned path in both hypervisor and VM



Huge Page VA Mapping

- Hypervisor is often unable to allocate consecutive physical huge pages and map to consecutive virtual addresses
- But, VM can see a mapped huge memory, which makes huge pages consecutive
- Offset Calculation (no looping)
 - ✓ Pre-calculate locations of each huge page
 - ✓ When a packet is received, we can calculate an offset in one programming line without looping (only bit operation and array index)

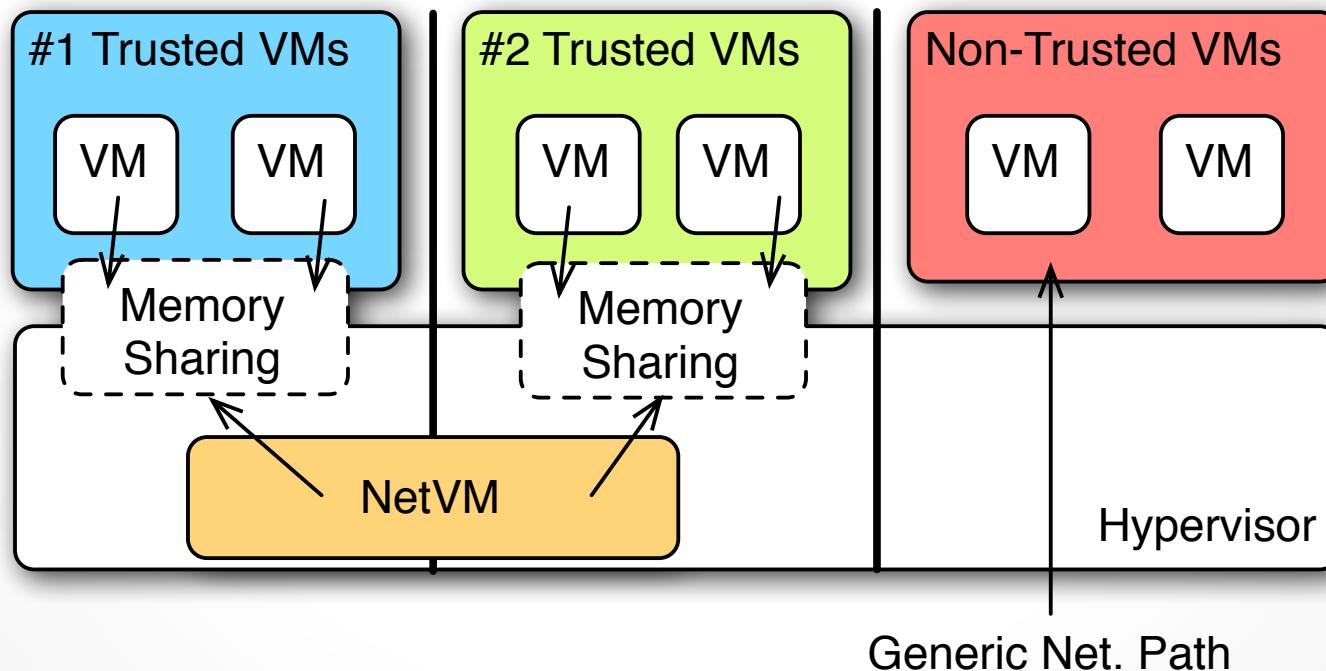
Host Huge Page VA Mapping



VM Huge Page PCI Mapping

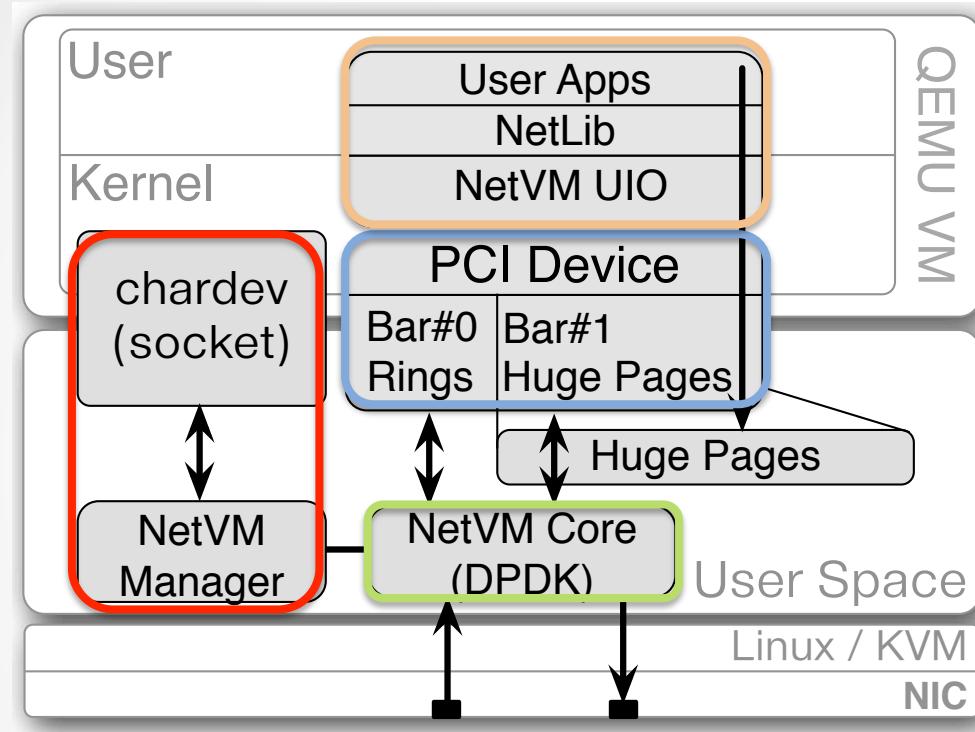
Trusted and Non-trusted Domains

- Virtualization should provide security guarantees among VMs
- Provide a security boundary between trusted VMs and non-trusted VMs
- Non-trusted VMs can not see packets from NetVM
- Groups of trusted VMs via huge page separation



NetVM Implementation

- KVM + DPDK



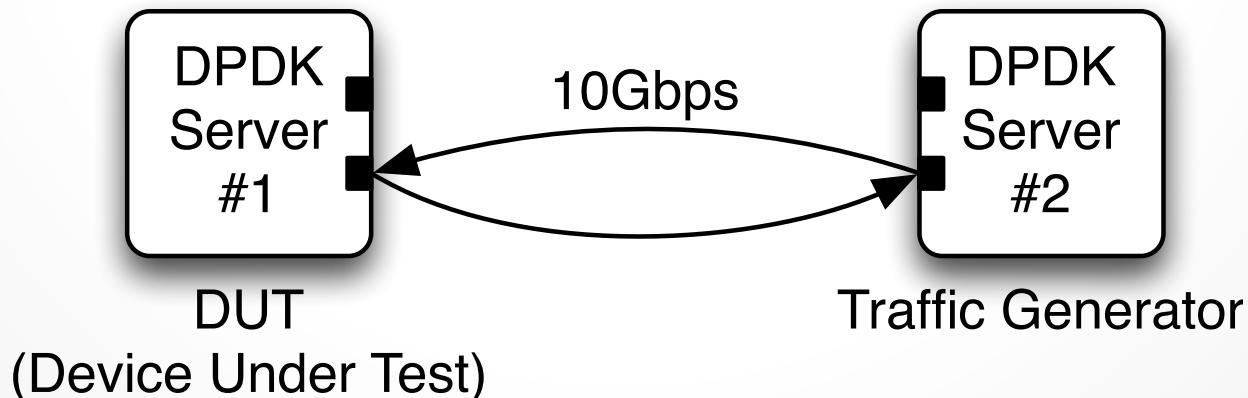
- **NetVM Manager**
 - Hypervisor user space application + QEMU chardev interface
- **NetVM Core**
 - Hypervisor user space application + DPDK
- **Emulated PCI (QEMU)**
 - KVM hardware emulation
- **NetLib**
 - VM user space library + Kernel module

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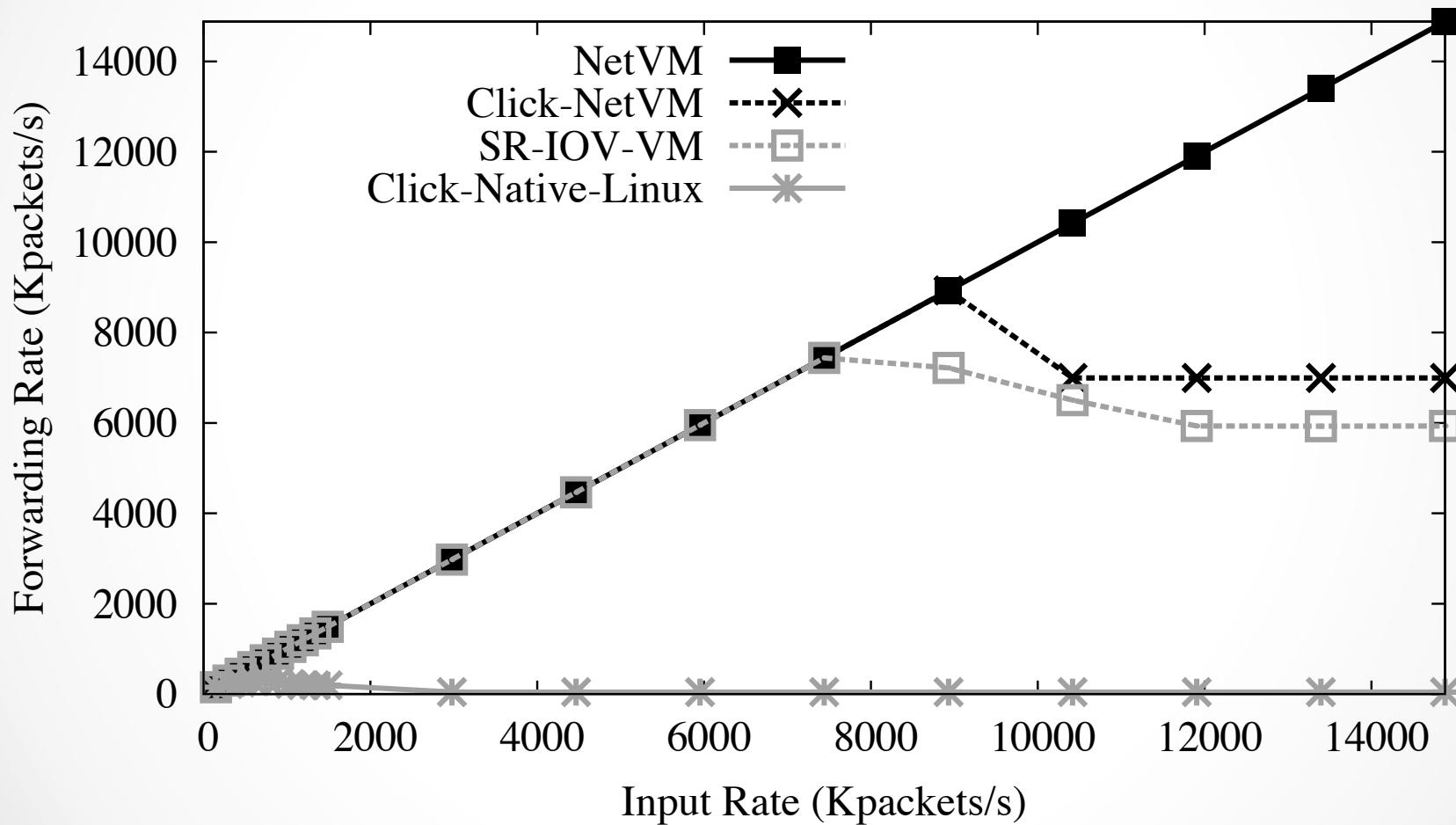
Evaluation Settings

- 2 x Nehalem-architecture (old architecture) machines
- 82599 Intel NIC
- 2 processors * 6 Cores (no hyperthreading)
 - ✓ 2 cores for receiving packets
 - ✓ 4 cores for transmitting/forwarding packets
 - ✓ Rest of cores for VM
- Total 8GB Hugepages (4GB per each processor)
- Only 1 port(out of 2 ports) is used
- Apps: L2/L3 Forwarder, Click Userspace Router, Firewall
- Packet Generator: DPDK-Pktgen (Wind River Systems, Inc.)

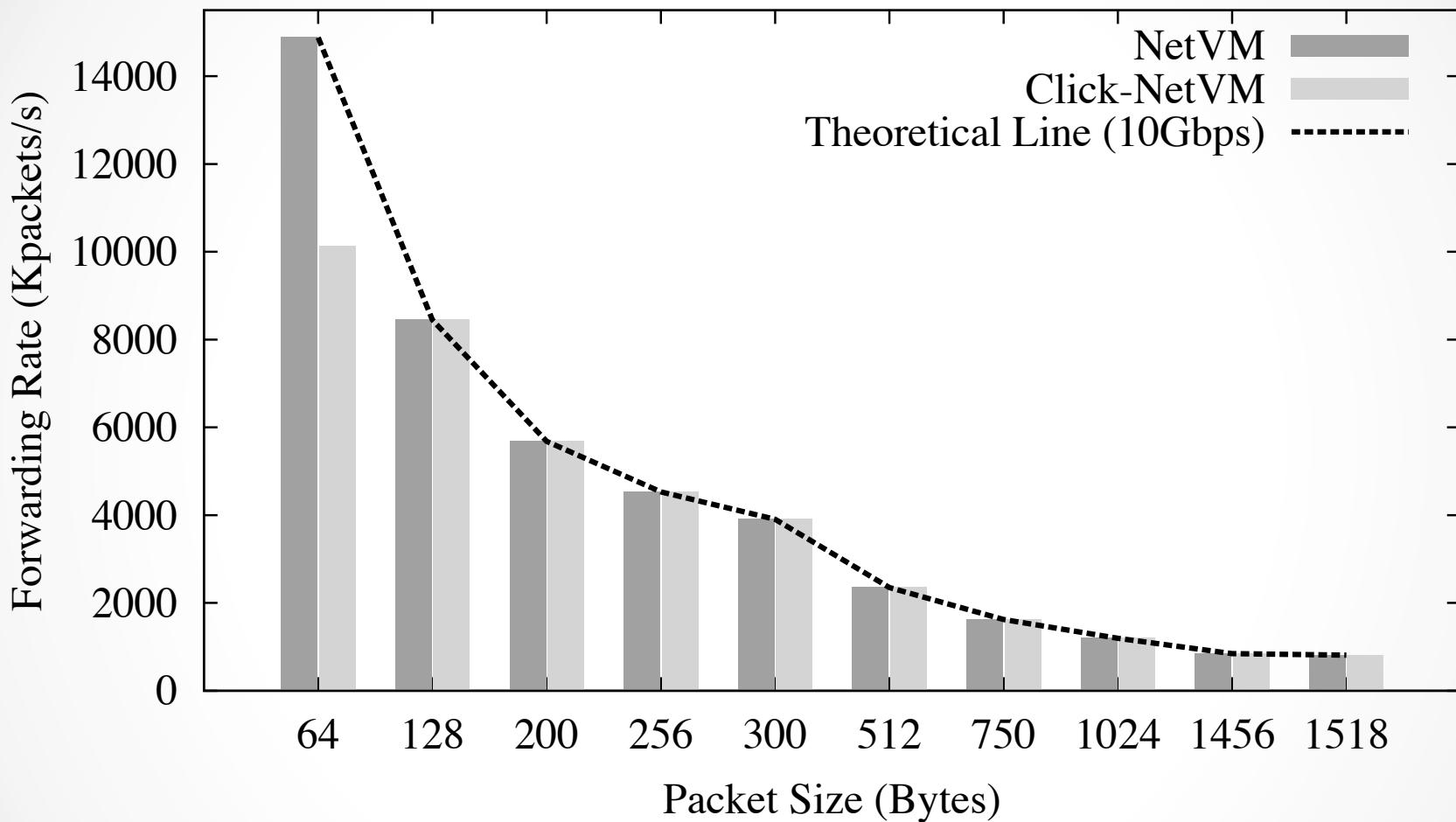


High Speed Packet Delivery

- 64-byte packets, 10Gbps = 14,880,952 packets/s

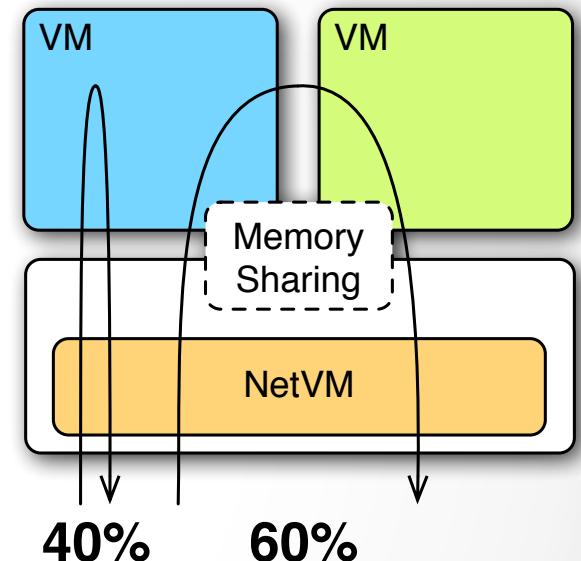
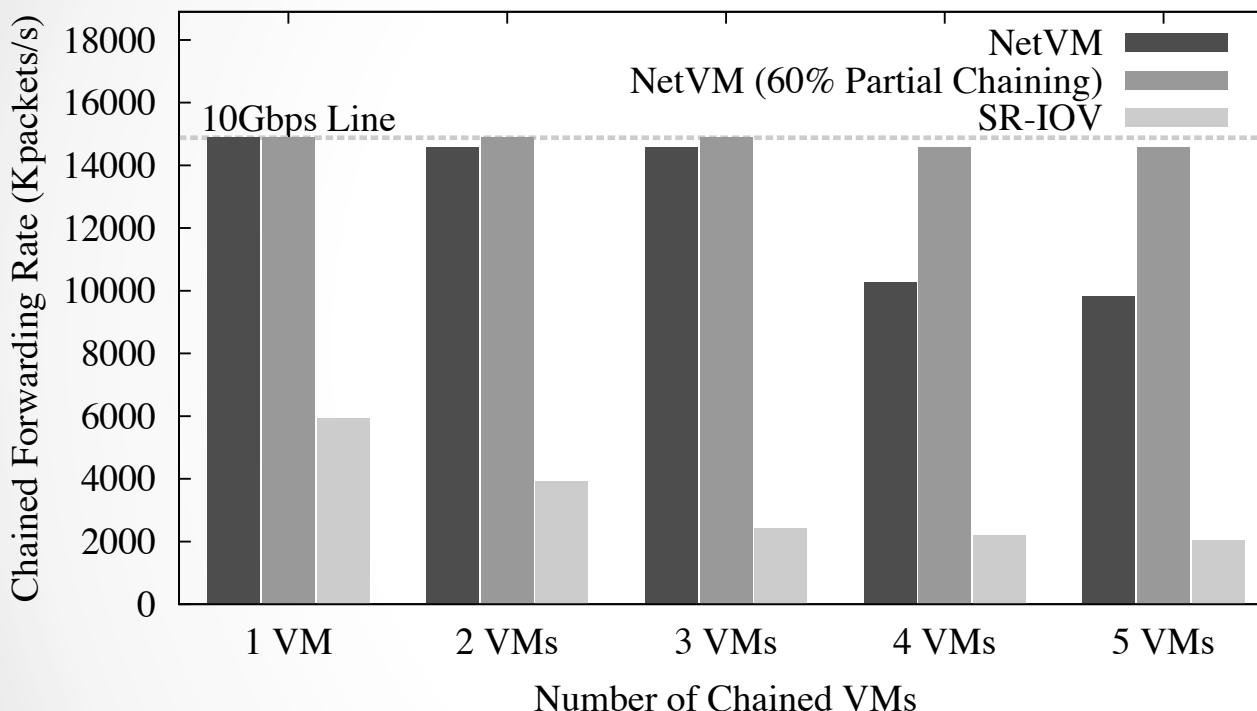


Packet Size vs. Forwarding Rate



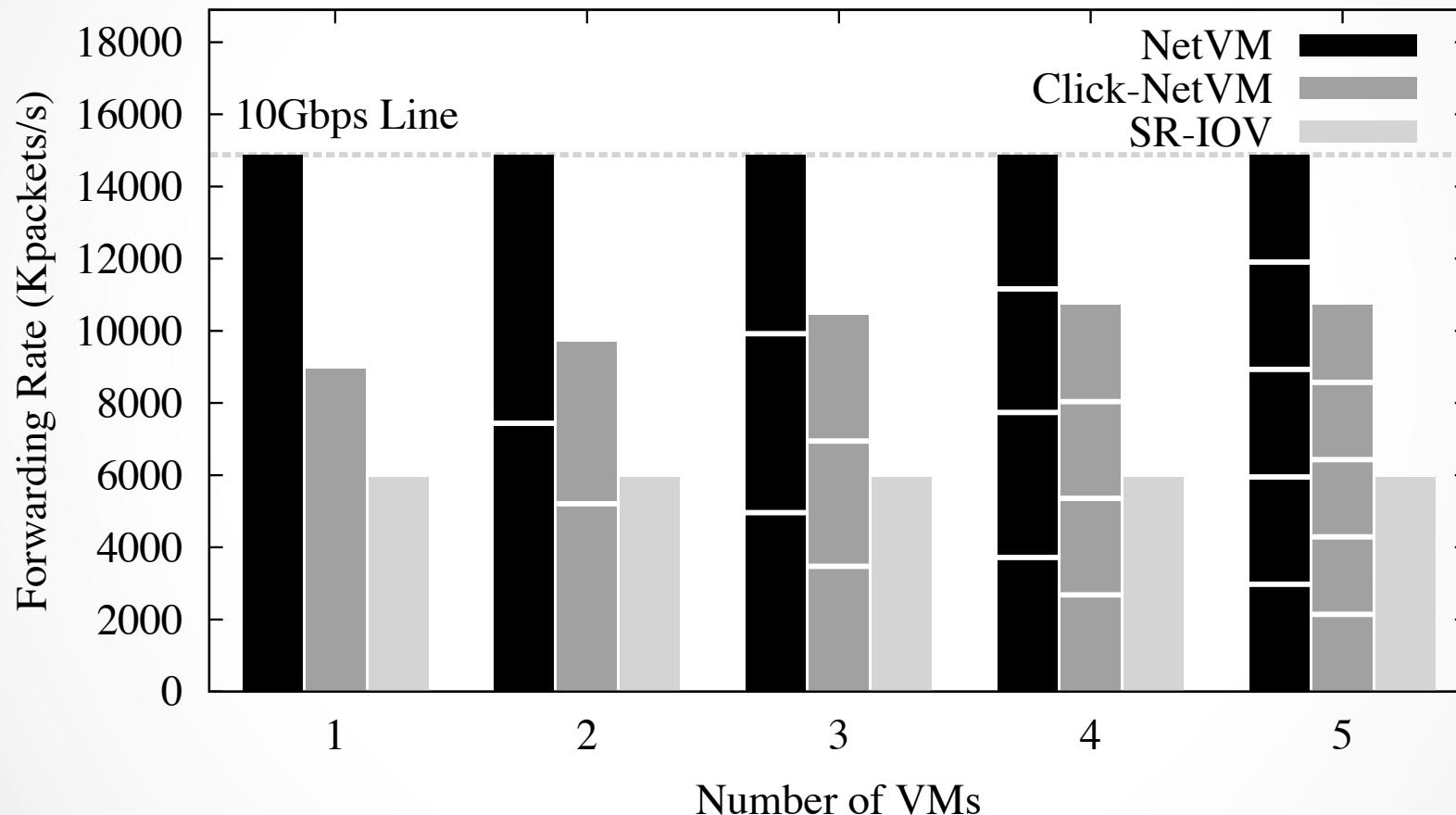
Inter-VM Forwarding

- A mix of L2/L3 forwarding and custom firewall (address filtering)



Switching Flexibility

- State-based load balancing (queue load)



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Conclusion

- NetVM, a high performance network framework in virtualized commodity platforms
- NetVM provides:
 - a virtualization-based high-speed packet delivery
 - memory sharing framework for network data
 - a hypervisor-based switching
 - high speed inter-VM communication
 - security domains

Future Work

- **Scale to next generation machine**
 - Sandy-bridge and Ivy-bridge processors have significant additional hardware capabilities – faster cache, data direct I/O technology, faster AES instructions, PCIe* Gen 3, IO MMU and Enhanced QPI
- **Building edge routers with NetVM**
 - We recognize that the capabilities of NetVM to act as a network element, such as an edge router in an ISP context
- **SDN integration**
 - We aim to integrate Open Flow into our NetVM Manager
- **Server type applications**
 - NetVM focuses on middle- box type applications, but end-host applications can also run in the NetVM platform

