

Overview LP

1. Examples: politics, flow, shortest paths
2. General form and converting to it
3. Simplex algorithm
 - Iterating over slack form

1/ Politics

How to campaign to win an election? -> estimate vote obtained per dollar spent advertising in support of a particular issue.

Policy	Urban	Suburban	Rural
Build roads [x1]	-2	5	3
Gun Control [x2]	8	2	-5
Farm Subsidies [x3]	0	0	10
Gasoline Tax [x4]	10	0	2

want a win by spending the min. amount of money.

Policy	Urban	Suburban	Rural
population	100,000	200,000	50,000
majority	50,000	100,000	25,000

Algebraic Setup

Let x_1, x_2, x_3, x_4 denote \$ spent per issue.

Minimize $x_1 + x_2 + x_3 + x_4$

subject to

1. $-2x_1 + 8x_2 + 0x_3 + 10x_4 \geq 50,000$
2. $5x_1 + 2x_2 + 0x_3 + 0x_4 \geq 100,000$
3. $3x_1 - 5x_2 + 10x_3 + 2x_4 \geq 25,000$

and $x_1, x_2, x_3, x_4 \geq 0$

(n variables) (n constraints)

optimum:

$$x_1 = 2050\,000/111$$

$$x_2 = 425\,000/111$$

$$x_3 = 0$$

$$x_4 = 625\,000/111$$

$$x_1 + x_2 + x_3 + x_4 = 3,100,000/111$$

x_i are real numbers

2/ standard form for LP

-> minimize or maximize linear objective fn subject to linear inequalities (or equations)

variable : $\langle x \rangle = \langle x_1 \ x_2 \ x_3 \ \dots \ x_n \rangle$

objective fn : $\langle c \rangle . \langle x \rangle = c_1 x_1 + \dots + c_n x_n$

inequalities : $A . \langle x \rangle \leq \langle b \rangle$

max. $\langle c \rangle . \langle x \rangle$ s.t. $A . \langle x \rangle \leq \langle b \rangle$ and $\langle x \rangle \geq 0$

3/ Certificate of Optimality (but why ?)

Is there a short certificate that shows LP solu. is optimal?

=> Consider $25/222 \cdot (1) + 36/222 \cdot (2) + 14/222 \cdot (3)$ [equations]

=> $x_1 + x_2 + 140/222 x_3 + x_4 \geq 3100000/111$

SINCE, $x_1 + x_2 + x_3 + x_4 \geq x_1 + x_2 + 140/222 x_3 + x_4 \geq 3100000/111$

4/ LP Duality

Theorem:

**** Equivalent ****

primal form	dual form
$\max \langle c \rangle . \langle x \rangle$	$\min \langle b \rangle . \langle y \rangle$
s.t. $A . \langle x \rangle \leq \langle b \rangle$	s.t. $A^T . \langle y \rangle \geq \langle c \rangle$
$\langle x \rangle \geq 0$	$\langle y \rangle \geq 0$

5/ Converting to Standard Form

1. Minimize $-2x_1 + 3x_2$, negate to $2x_1 - 3x_2$, maximize
2. Suppose x_j does NOT have a non-negative constraint, x_j replace with $x_j' - x_j''$, $x_j' \geq 0$ $x_j'' \geq 0$
3. equality constraint $x_1 + x_2 = 7$, $x_1 + x_2 \leq 7$, $-x_1 - x_2 \leq -7$
4. \geq constraint translates to \leq by (-1) multiply

6/ Maximum Flow

$$\max \sum_{v \in V} \{f(s, v) = |f|\}$$

1. Skew symmetry: $f(u, v) = -f(v, u)$ [for all u, v in V]
2. Conservation : $\sum_{v \in V} \{f(u, v) = 0\}$ [for all u in $V - \{s, t\}$]
3. Capacity : $f(u, v) \leq c(u, v)$ [for all u, v in V]

Two Commodities (1 & 2)

f_1, c_1, f_2, c_2 (two distinct disjoint opts)

$$f_1, f_2, \text{ single capacity } c, w_1 f_1(u, v) + w_2 f_2(u, v) \leq c(u, v)$$

7/ Shortest Path

From Vertex S \

[triangular inequality] \ $d[v] - d[u] \leq w(u,v)$ [for all (u,v) in E] \

$$d[s] = 0$$

$$w(u_1, v) \ \& \ w(u_2, v)$$

$$u_1 \rightarrow v \leftarrow u_2$$

$$d[v] - d[u_1] \leq w(u_1, v)$$

$$d[v] - d[u_2] \leq w(u_2, v)$$

$$d[v] = \text{minimum} \{ \dots \}$$

objective: max SUM[V]{d[v]}

** change minimize to maximize ** (becoz all constraints serve as & operator)

8/ Simplex Algorithm | worst case ($m+n$ Choose m)

Flow: Represent LP in slack form

convert one slack form into an equivalent slack form

whose objective value has not decreased, and has likely increased. Keep going till the optimal solution become obvious.

Example

Maximize $3x_1 + x_2 + x_3$

s.t.

$$x_1 + x_2 + 3x_3 \leq 30$$

$$2x_1 + 2x_2 + 5x_3 \leq 24$$

$$4x_1 + x_2 + 2x_3 \leq 36$$

$$x_1, x_2, x_3 \geq 0$$

slackform

$$z = 3x_1 + x_2 + x_3 \text{ [non-basic variables]}$$

$$\text{[basic variables][equation I]}$$

$$x_4 = 30 - x_1 - x_2 - 3x_3$$

$$x_5 = 24 - 2x_1 - 2x_2 - 5x_3$$

$$x_6 = 36 - 4x_1 - x_2 - 2x_3 \setminus$$

$$x_1, x_2, x_3, x_4, x_5, x_6$$

Basic Solution

set all non-basic variables to 0

calculate values of basic variables

$$\text{obj fn: } 3(0) + 1(0) + 2(0) = 0$$

$$\langle 0, 0, 0, 30, 24, 36 \rangle \setminus$$

Pivoting

1. Select a non-basic x_e whose coefficient in the obj. fn is positive
2. Increase the value of x_e as much as possible without violating any of the constraints
3. Variable x_e become basic
4. x_e becomes basic and other variable becomes non-basic. (value of other basic variable)

Iteration

select x_1 nonbasic, increase the value of x_1

3rd constraint is tightest one:

$$x_1 = 9 - x_2/4 - x_3/2 - x_6/4$$

Rewrite the other equations with x_6 on the R.H.S.

i.e. replace x_1 with above equation \

$$z = 27 + x_2/4 + x_3/2 - 3x_6/4$$

$$x_1 = 9 - x_2/4 - x_3/2 - x_6/4$$

$$x_4 = 21 - 3x_2/4 - 5x_3/2 + x_6/4$$

$$x_5 = 6 - 3x_2/2 - 4x_3 + x_1/2$$

original basic solution: $\langle 0, 0, 0, 30, 24, 36 \rangle$

satisfies [equation II] and has

$$\text{objective value: } 27 + 1/4 * 0 + 1/2 * 0 - 3/4 * 36 = 0 \setminus$$

basic solution II: set nonbasic values on 0

$$\langle 9, 0, 0, 21, 6, 0 \rangle$$

$$\text{objective value: } 9 * 3 = 27$$