

Rewriting Modulo Symmetric Monoidal Structure

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Joint work with

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CNRS

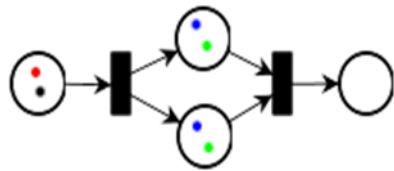
Aleks Kissinger
Radboud U.

Fabio Gadducci
U. Pisa

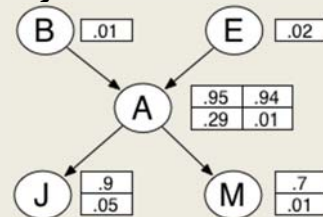
Pawel Sobocinski
U. Southampton

PROPs: algebras of network diagrams

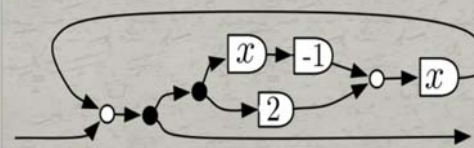
Petri Nets



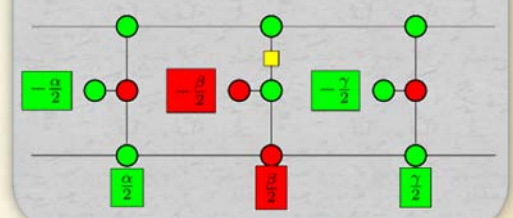
Bayesian Networks



LICS'16
Signal Flow Graphs



LICS'16
Quantum Processes

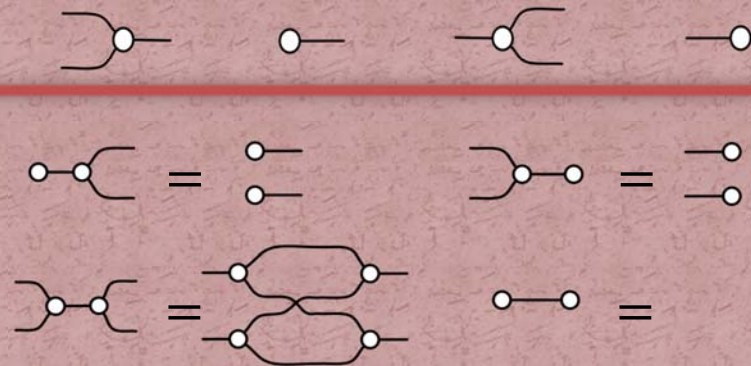


“Linear” Lawvere theory

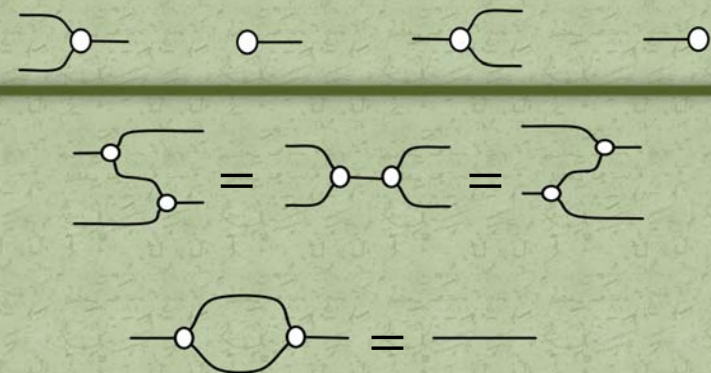
A PROP is (just) a symmetric monoidal category with set of objects \mathbb{N}

PROPs can be freely constructed starting from a signature Σ and equations E

PROP of Bialgebras



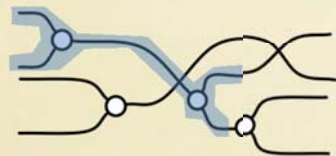
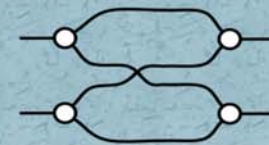
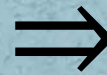
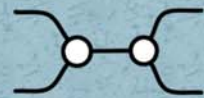
PROP of Separable Frobenius Algebras



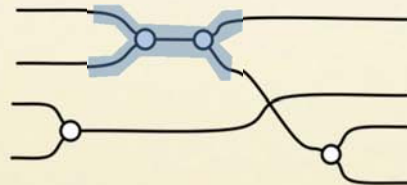
Rewriting in a PROP

Perspective of this work:
see E as a **rewriting system** on network diagrams

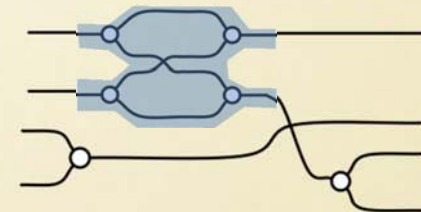
(R)



SMC
 \approx



\Rightarrow_R



Our question

How to implement rewriting modulo symmetric monoidal structure in a simple, yet rigorous way?

Outline

Rewriting modulo symmetric monoidal structure

Rewriting modulo SM +
Frobenius structure

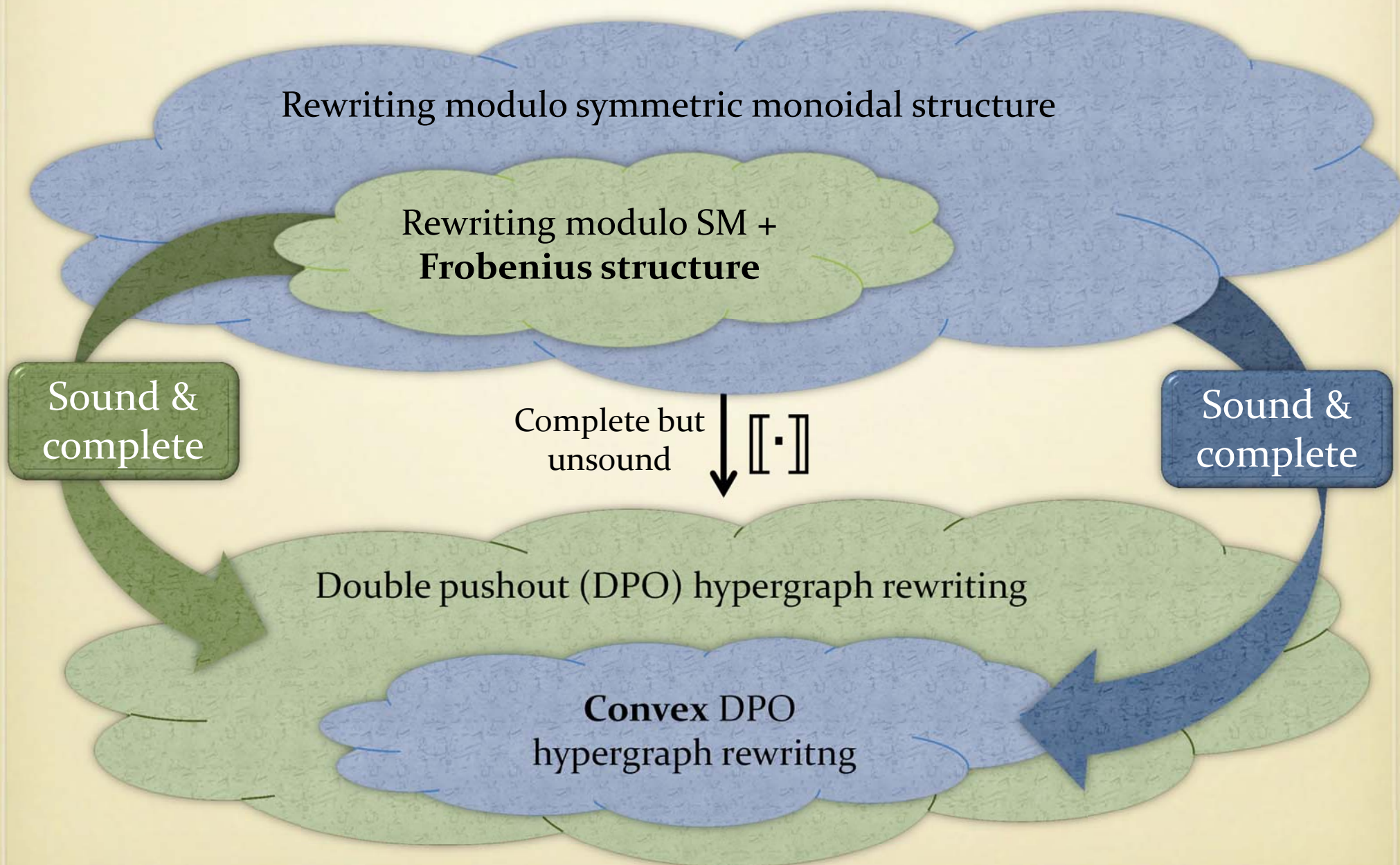
Sound &
complete

Complete but
unsound \downarrow $[\![\cdot]\!]$

Sound &
complete

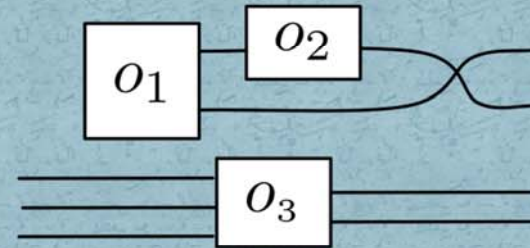
Double pushout (DPO) hypergraph rewriting

Convex DPO
hypergraph rewriting



Hypergraph interpretation

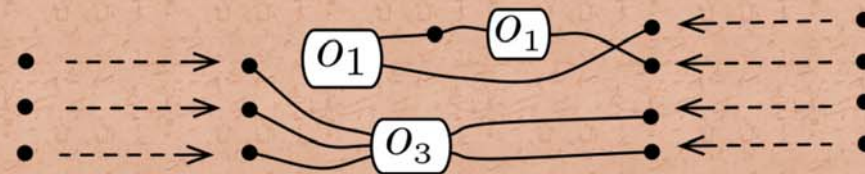
PROP $Syn(\Sigma)$ of syntax
freely generated by
 $\Sigma = \{ \boxed{o_1}, \boxed{o_2}, \boxed{o_3} \}$



Operations from $\Sigma \sim$ hyperedges
Left/right boundary \sim Cospan structure

$\llbracket \cdot \rrbracket$

PROP of (discrete) Cospans
of Σ -labelled hypergraphs
 $Csp(Hyp(\Sigma))$

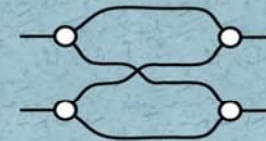
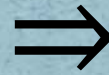
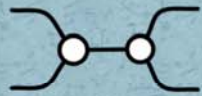


Proposition: $Syn(\Sigma) \xrightarrow{\llbracket \cdot \rrbracket} Csp(Hyp(\Sigma))$ is faithful

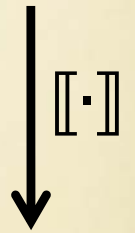
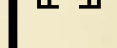
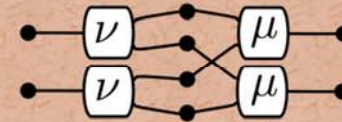
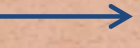
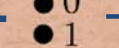
DPO hypergraph rewriting

$Hyp(\Sigma)$ is an *adhesive category* (Lack & Sobocinski)
and thus adapted to double-pushout rewriting.

(R)



$Syn(\Sigma)$

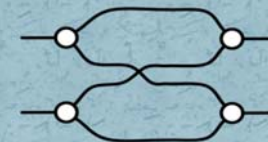
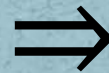
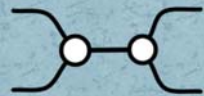


$Hyp(\Sigma)$

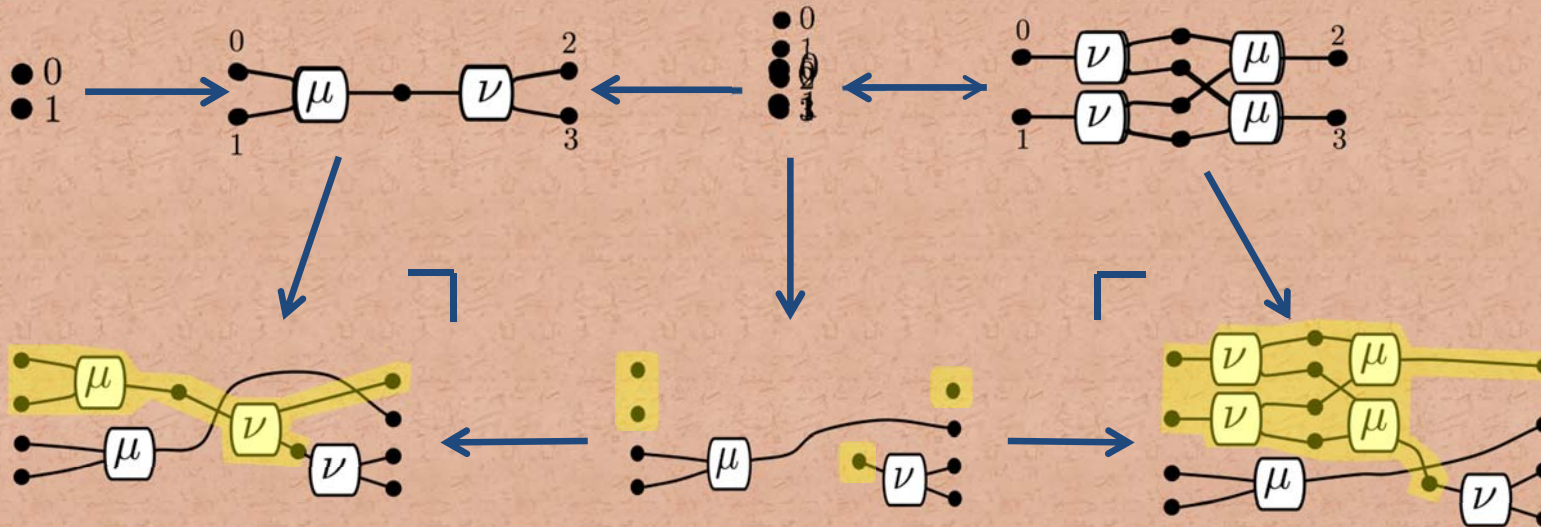
DPO hypergraph rewriting

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(R)



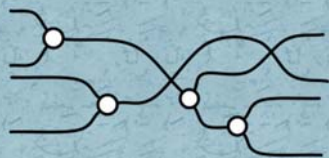
$Syn(\Sigma)$



$\llbracket \cdot \rrbracket$

$Hyp(\Sigma)$

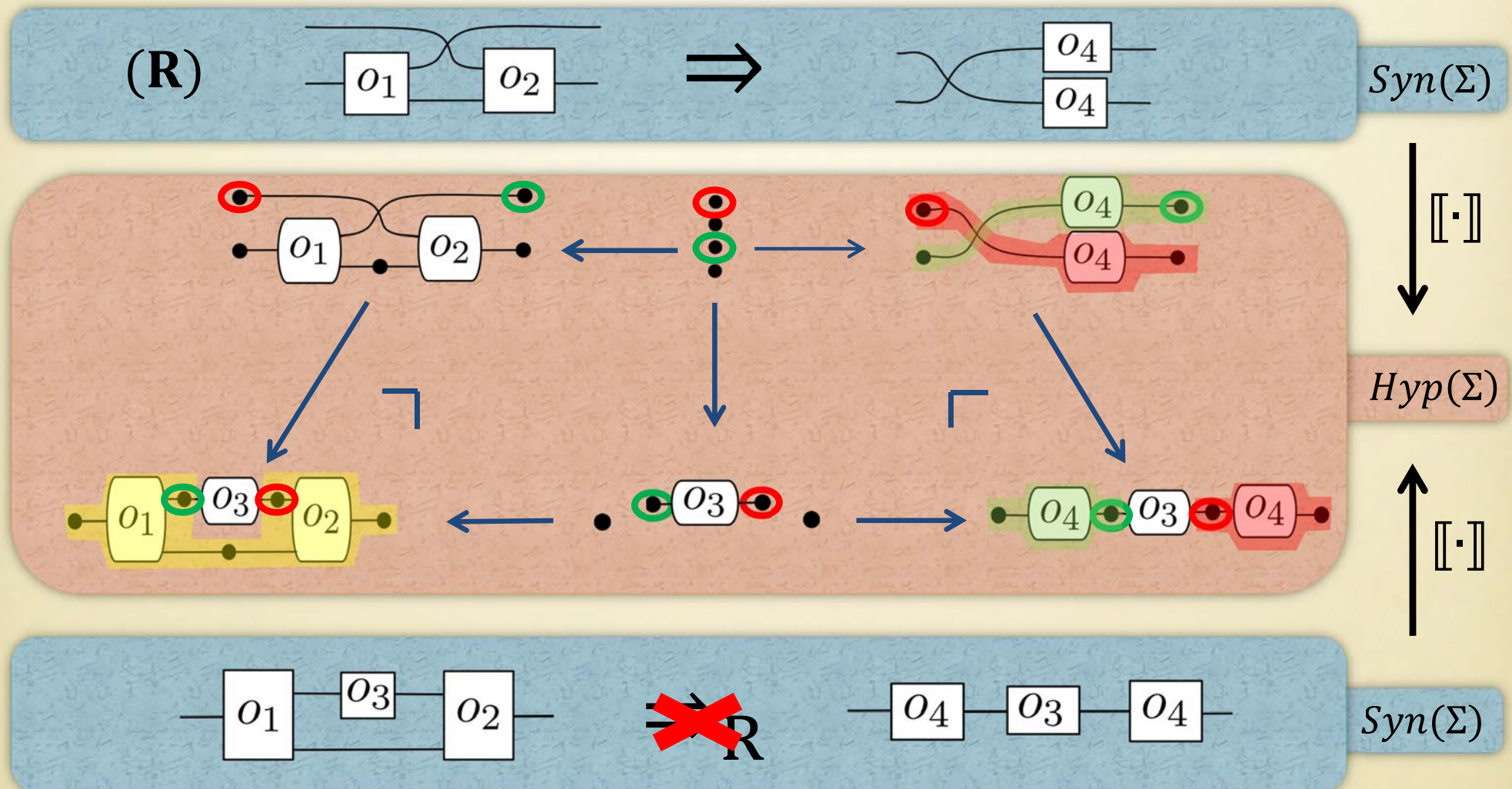
$\llbracket \cdot \rrbracket$



$Syn(\Sigma)$

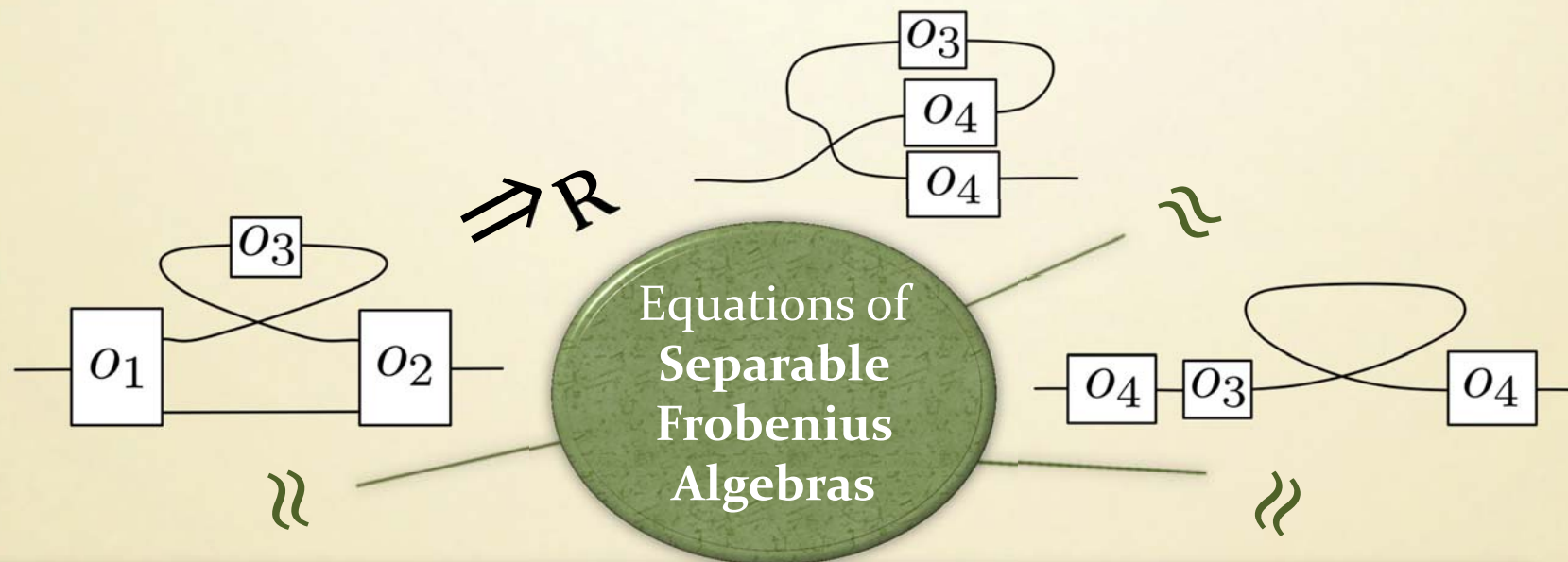
DPO rewriting is unsound

DPO hypergraph rewriting is *complete* but generally not *sound*.



DPO rewriting is unsound

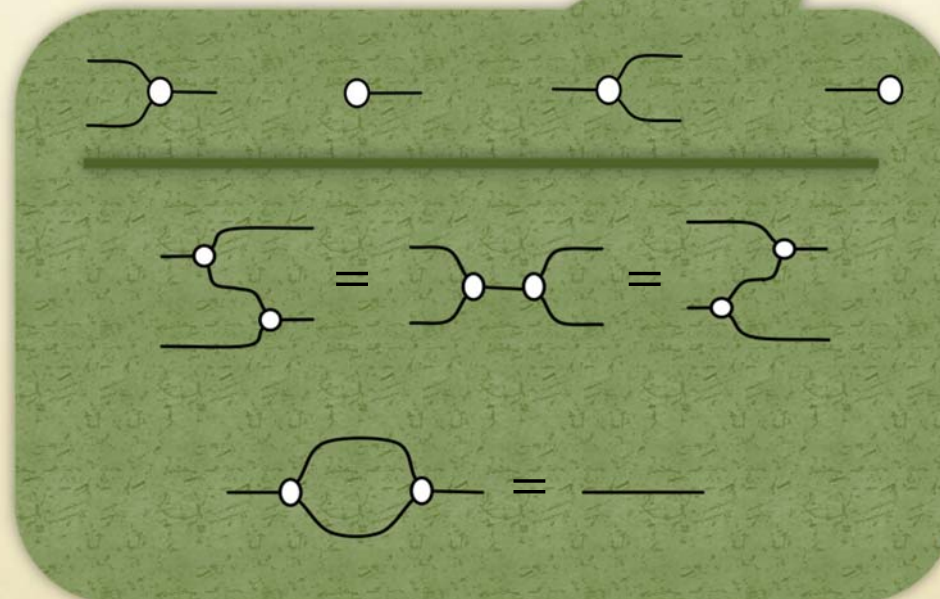
DPO hypergraph rewriting is *complete* but generally not *sound*.



Frobenius makes DPO rewriting sound

Theorem I

DPO hypergraph rewriting is *sound and complete* for symmetric monoidal categories *with a chosen separable Frobenius structure*.



Frobenius makes DPO rewriting sound

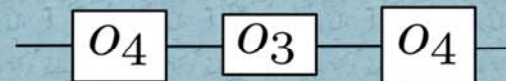
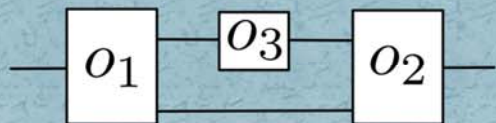
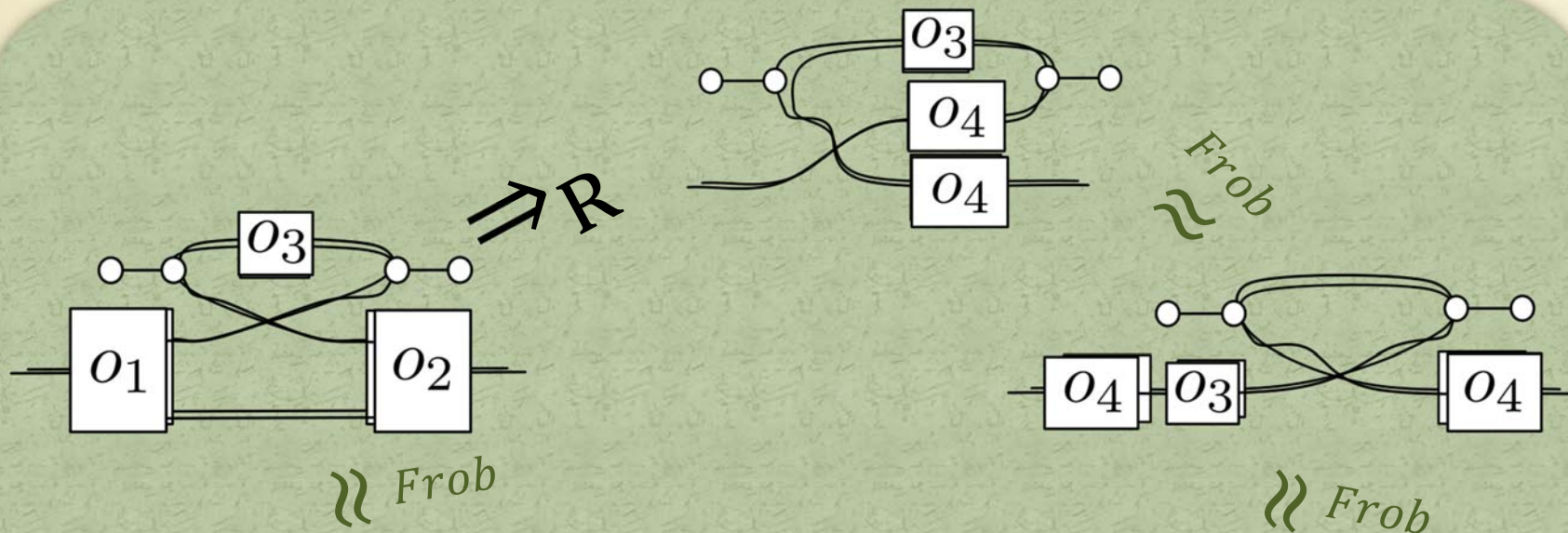
(R)



\Rightarrow

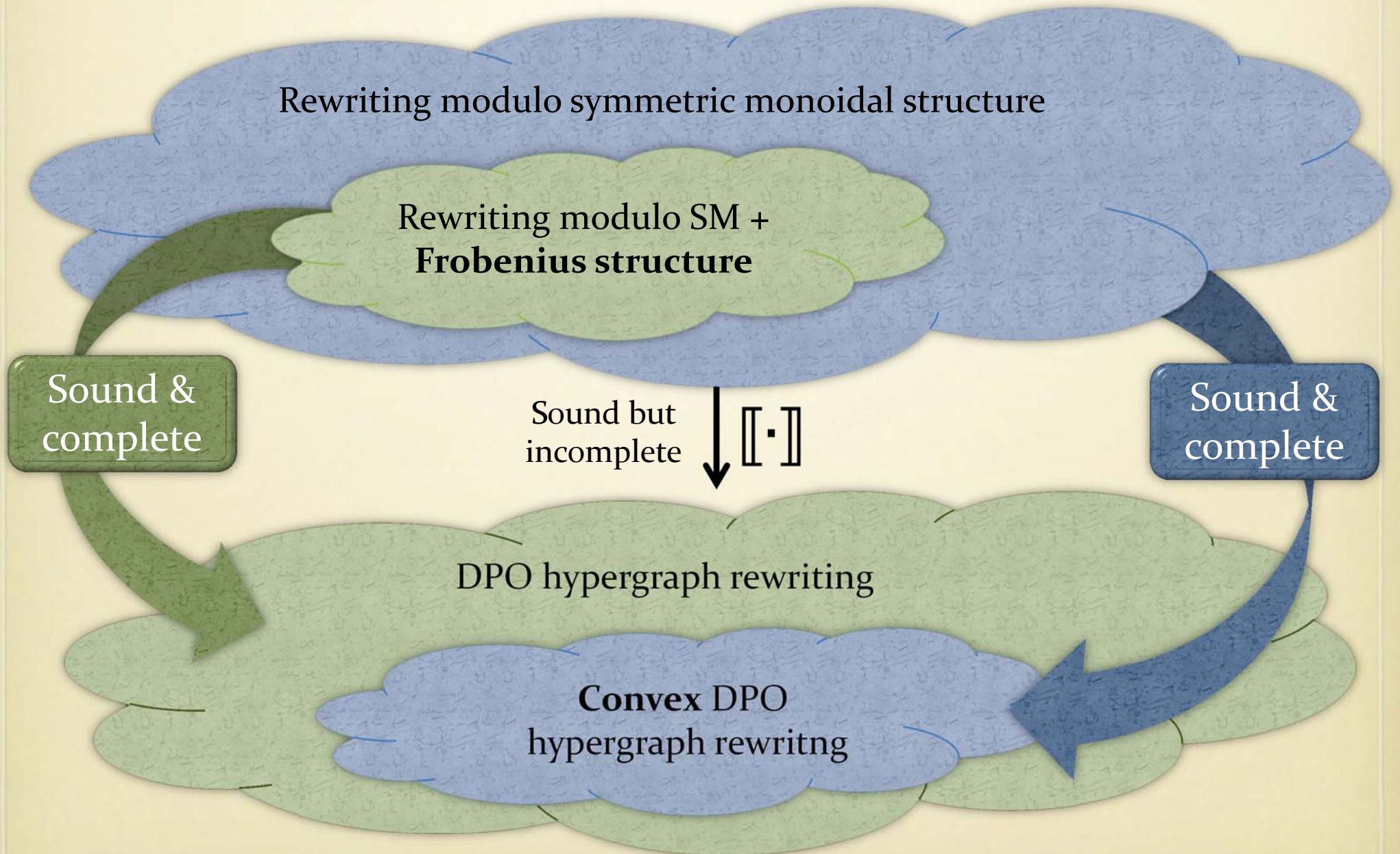


$Syn(\Sigma)$
+
 $Syn(\Sigma)$

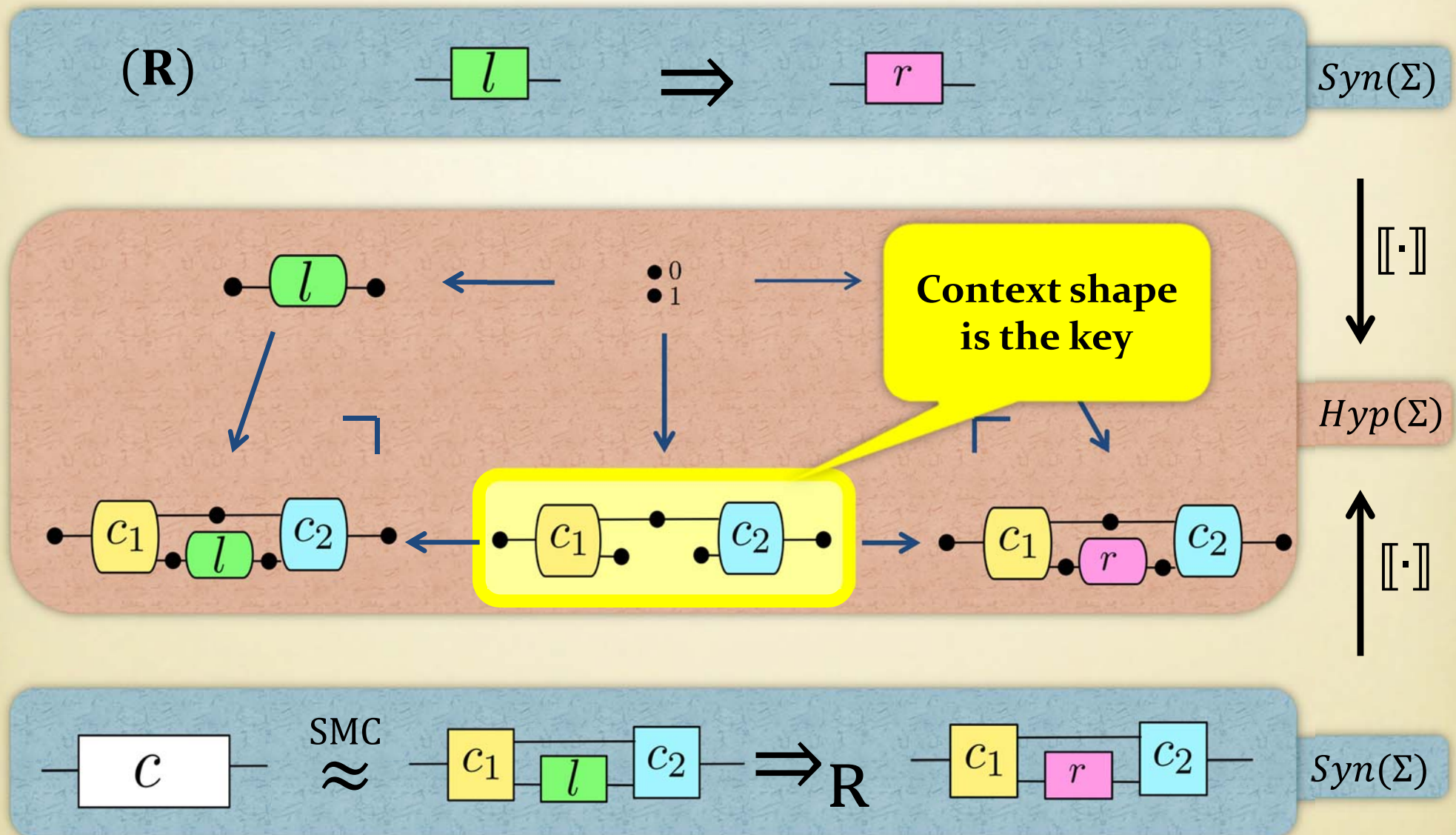


$Syn(\Sigma)$
+
 $Syn(\Sigma)$

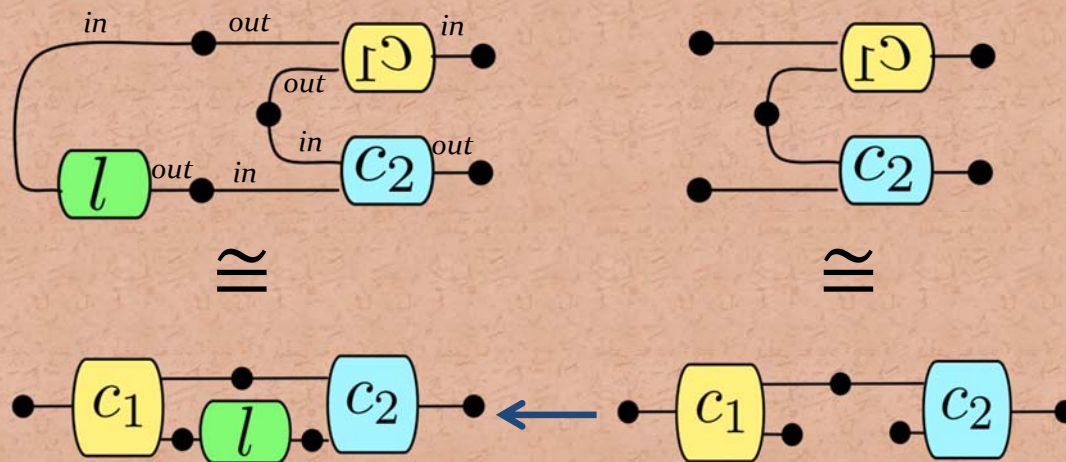
Where we are, so far



How does a sound DPO rewriting look like?



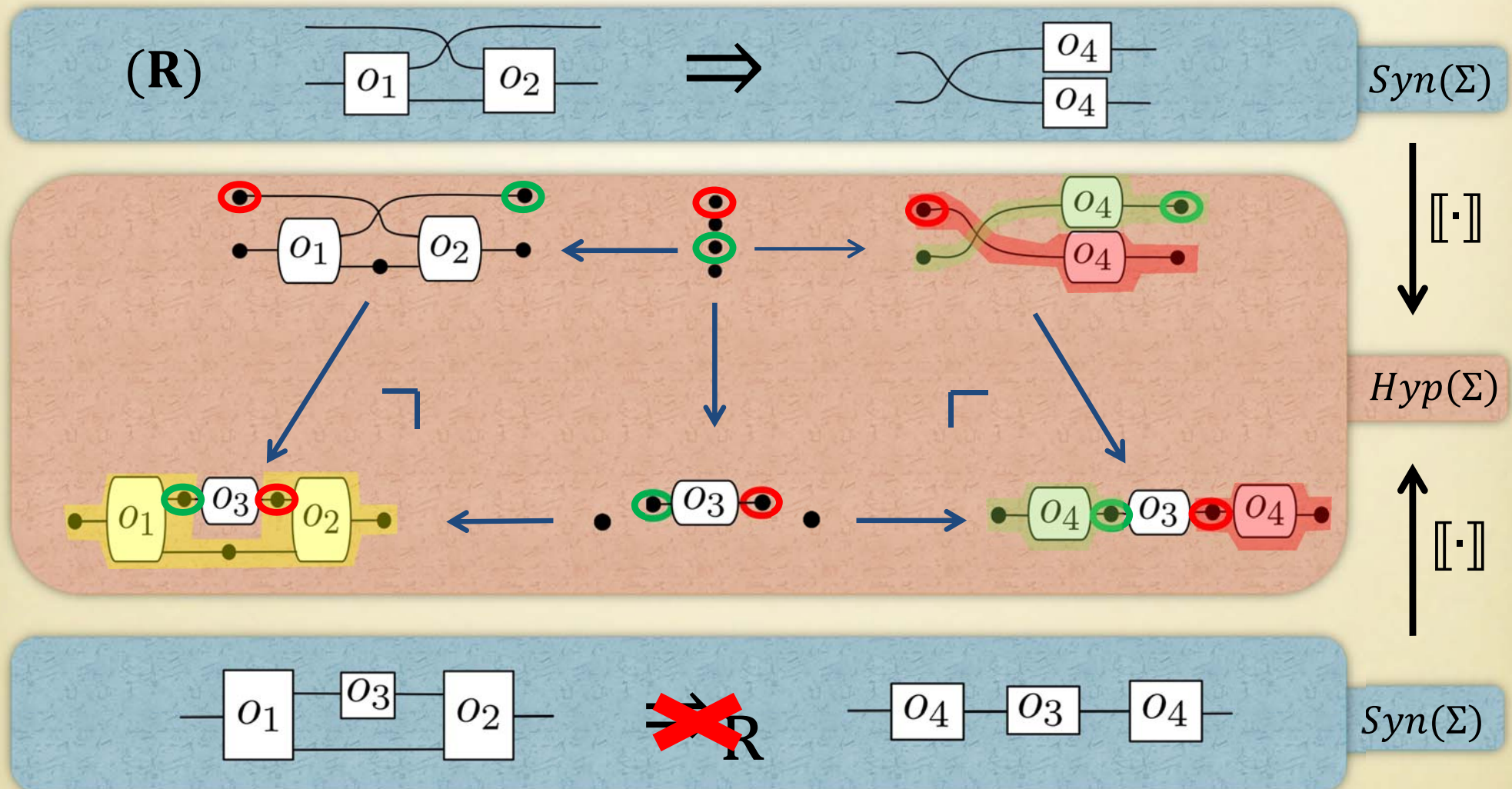
How does a sound DPO rewriting look like?



Leading Intuition
*a rewriting steps
 is sound
 iff
 the rewriting context
 has this shape*

$Hyp(\Sigma)$

Back to the soundness counterexample

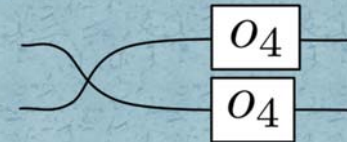


Back to the soundness counterexample

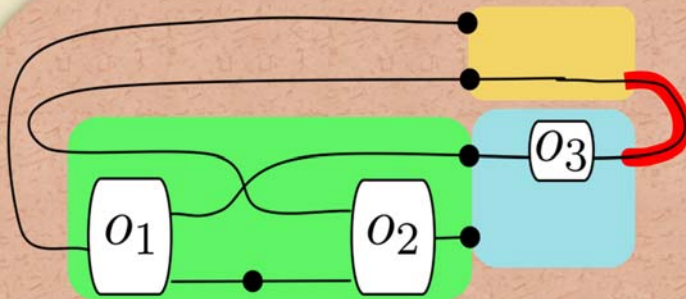
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\Rightarrow

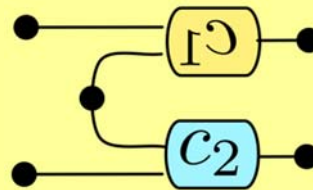
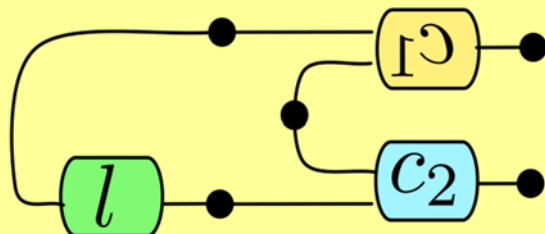
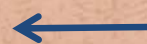
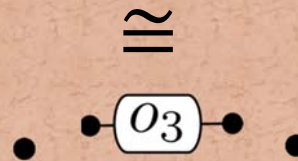
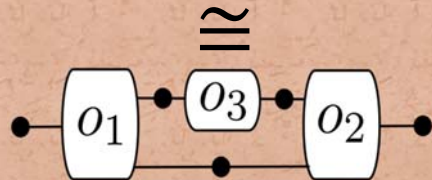


$Syn(\Sigma)$



Unsound context shape

$Hyp(\Sigma)$



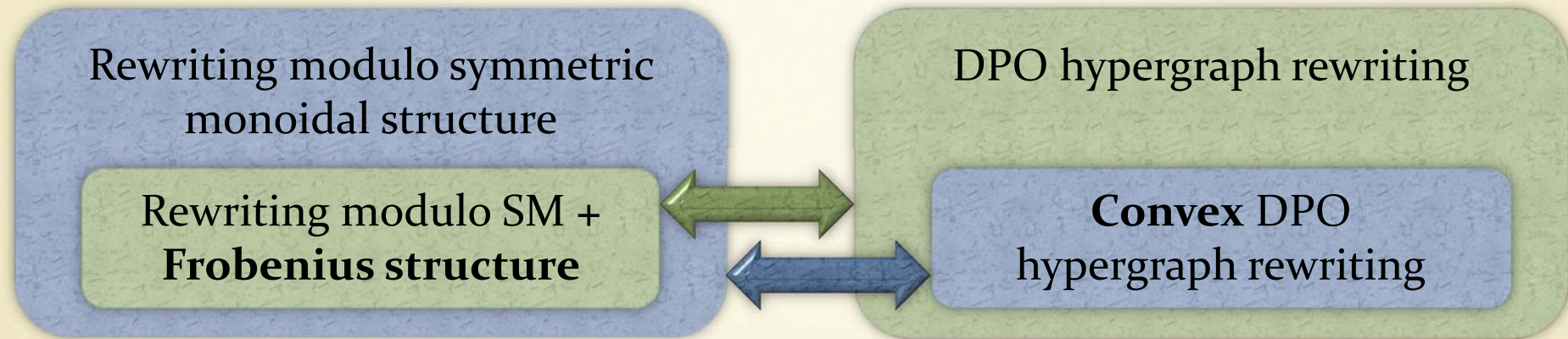
Sound context shape

Convex DPO rewriting is sound

Theorem II

Convex DPO hypergraph rewriting is sound and complete for symmetric monoidal categories.

Discussion



- Ongoing and future work
 - More examples
 - Frobenius structures are commonplace in algebras of circuit diagrams.
 - Study of critical pairs, confluence, termination.
 - Relationship with equational theories generated by distributive laws of PROPs.