# Report First Project IAJ

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## 1 Introduction

The goal of this project was to create different levels of path finding algorithms, and compare their performance. We compared 4 differente algorithms: Basic  $A^*$  (unordered list for open set, unordered list for closed set), Basic  $A^*$  but using tiebreaking (unordered list for open set, unordered list for closed set), NodeArray  $A^*$  (NodeArray for open and closed set) and NodeArray  $A^*$  with Goal Bounding.

## 2 Basic A\*

## 2.1 Algorithm

The A\* is a search algorithm that uses a heuristic to find the best path between 2 nodes. Even though it's a basic algorithm, it shows relatively good performance when compared to other algorithms like djikstra. Despite being a basic algorithm, it's performance can be improved by using better data structures, or using other optimizations like we're going to show in the next sections.

## 2.2 Data

Table 2: Basic A\* performance (Path 2) Table 1: Basic A\* performance (Path 1) Method Execution Time (ms) Method Calls Execution Time (ms) Calls A\*Pathfinding.Search 10181.85 A\*Pathfinding.Search 21432.88  ${\bf GetBestAndRemove}$ 1904 44.09GetBestAndRemove 2890 114.64 AddToOpen 1954 1.82 AddToOpen 3020 2.54SearchInOpen 18564 269.64 SearchInOpen 28291 759.53 RemoveFromOpen RemoveFromOpen 0 0 0 0 0 Replace 0 0 Replace 0 AddToClosed 1904 1.45 AddToClosed 2890 1.76

SearchInClosed

RemoveFromClosed

9710.25

0

Table 3: Basic A\* grid information (Path 1)

TotalPNodes | TotalPTime | MaxOpenNodes

1904 | 0.02 | 77

18460

0

Table 4: Basic A\* grid information (Path 2)

TotalPNodes | TotalPTime | MaxOpenNodes

2889 | 0.02 | 133

27990

0

20319.63

0

# 3 Basic A\* with tiebraking

## 3.1 Algorithm

SearchInClosed

RemoveFromClosed

This algorithm is basically the previous one, but we use tiebraking to break ties between nodes with the same f value. When 2 nodes have the same f value, we use order the nodes by smallest h value. This way we first pick the node that is closer to the goal.

### 3.2 Data

Table 5: Basic A\* with tiebraking performance (Path 1)

Table 6: Basic A\* with tiebraking performance (Path 2)

Method	Calls	Execution Time (ms)	Method	Calls	Execution Time (ms)
A*Pathfinding.Search	1	10029.02	A*Pathfinding.Search	1	21970.11
GetBestAndRemove	1904	112.83	GetBestAndRemove	2890	282.13
$\operatorname{AddToOpen}$	1954	1.56	AddToOpen	3021	1.94
SearchInOpen	18564	260.07	SearchInOpen	28291	768.2
RemoveFromOpen	0	0	RemoveFromOpen	0	0
Replace	0	0	Replace	0	0
AddToClosed	1904	1.27	AddToClosed	2890	1.61
SearchInClosed	18460	9524.91	SearchInClosed	27990	20761.34
${\bf Remove From Closed}$	0	0	RemoveFromClosed	0	0

Table 7: Basic  $A^*$  with tiebraking grid information (Path 1)

TotalPNodes	TotalPTime	MaxOpenNodes
1904	0.02	77

Table 8: Basic A\* with tiebraking\* grid information (Path 2)

TotalPNodes	${\bf Total PTime}$	MaxOpenNodes		
2889	0.02	133		

## 3.3 Comparison

Looking at the data from this and the previous algorithm, that there's a little optimization on the data of path 1 even though there's more spent time getting the best node from the open set, due to more comparisons. In path 2 we still see an increase on the time for GetBestAndRemove, but surprisingly there's more time spent this may be due to bad measurement of data, since this contradicts what is expected.

## 4 NodeArray A\*

## 4.1 Algorithm

NodeArray A\* is an A\* implementation that uses a NodeArray to store the nodes. We use this array as our open and closed set, and we change the status property of the nodes when we add them to the open or closed set. This way, we can search for nodes in the open and closed set in constant time.

### 4.2 Data

Table 9: NodeArray A\* performance (Path 1)

${f Method}$	Calls	Execution Time (ms)	${f Method}$	Calls	Execution Time (ms)
A*Pathfinding.Search	1	5.48	A*Pathfinding.Search	1	159.11
GetBestAndRemove	200	1.96	GetBestAndRemove	2890	38.36
$\operatorname{AddToOpen}$	232	1.13	AddToOpen	3019	9.31
SearchInOpen	1044	0.04	SearchInOpen	28247	1.84
RemoveFromOpen	0	0	RemoveFromOpen	0	0
Replace	0	0	Replace	0	0
AddToClosed	100	0.01	AddToClosed	2885	0.58
SearchInClosed	1009	0.04	SearchInClosed	27932	1.73
${\bf Remove From Closed}$	0	0	${\bf Remove From Closed}$	0	0

Table 11: Node Array A\* grid information (Path 1)

${\bf Total PNodes}$	TotalPTime	MaxOpenNodes
1904	0.4738739	77

Table 12: NodeArray A\* grid information (Path 2)

${\bf Total PNodes}$	${\bf Total PTime}$	MaxOpenNodes
2884	0.4738739	135

## 4.3 Comparison

NodeArray A\* is faster than the previous algorithms, due to the fact that we can search for nodes in the open and closed set in constant time, as we can see by the reduction of the SearchInOpen and SearchInClosed time, even though it spends more time on the AddToOpen due to the use of a PriorityHeap.

# 5 NodeArray A\* with Goal Bounding

### 5.1 Algorithm

By using precomputation of the grid, we can make bounding boxes for each node and improve the NodeArray A\* algorithm. We do this by using djikstra to calculate fastest path from each node to all other nodes. This way, we know which direction

we should choose when trying to go to a specific node. This optimization causes, sometimes, a heavy increase on the starting time, due to the precomputation, but it improves the runtime of the algorithm by a lot.

#### 5.2 Data

Table 13: NodeArray A\* with Goal Bounding performance (Path 1)

mance (Path 2)		
Method	Calls	Execution Time (ms)
A*Pathfinding.Search	1	17.12
${\bf GetBestAndRemove}$	158	0.40
$\operatorname{AddToOpen}$	165	0.35
SearchInOpen	388	0.01
RemoveFromOpen	0	0
Replace	0	0
AddToClosed	158	0.03
	Method  A*Pathfinding.Search GetBestAndRemove AddToOpen SearchInOpen RemoveFromOpen Replace	MethodCallsA*Pathfinding.Search1GetBestAndRemove158AddToOpen165SearchInOpen388RemoveFromOpen0Replace0

0.01

Table 14: NodeArray A\* with Goal Bounding perfor-

Method	Calls	Execution Time (ms)	Method	Calls	Execut
A*Pathfinding.Search	1	10.62	A*Pathfinding.Search	1	
${\bf GetBestAndRemove}$	200	0.43	GetBestAndRemove	158	
AddToOpen	216	0.34	AddToOpen	165	
SearchInOpen	235	0	SearchInOpen	388	
RemoveFromOpen	0	0	RemoveFromOpen	0	
Replace	0	0	Replace	0	
AddToClosed	100	0.01	AddToClosed	158	
SearchInClosed	126	0	SearchInClosed	282	
RemoveFromClosed	0	0	RemoveFromClosed	0	
	!	'		1	

Table 15: NodeArray A\* with Goal Bounding grid information (Path 1)

${\bf Total PNodes}$	TotalPTime	MaxOpenNodes
228	0.1092952	9

Table 16: NodeArray A\* with Goal Bounding grid information (Path 2)

${\bf Total PNodes}$	TotalPTime	MaxOpenNodes
157	0.4738739	8

#### 5.3 Comparison

Comparing this data with the previous ones, we can see that this is by far the best optimization in terms of runtime. This is due to the use of bounding boxes, that shorten the amounts of nodes we process, and thus the amount of calls to add, remove and search in the open and closed set.

### Bonus Level 6

#### 7 Conclusions

Analysing all algorithms we can acess that A\* by itself is already a good algorithm, but it's optimizations make it must faster, without compromising finding the best path.

Also, we can notice that adding pre-processing to the algorithm can improve it's runtime by a lot, even though it takes some time to do it. We can see this in the use of precomputation on the Goal Bounding algorithm and on the calculation of the Dead-End heurisitic.