

# Report First Project IAJ

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## 1 Introduction

The goal of this project was to create different levels of path finding algorithms, and compare their performance. We compared 4 different algorithms: Basic A\*(unordered list for open set, unordered list for closed set), Basic A\* but using tiebreaking (unordered list for open set, unordered list for closed set), NodeArray A\* (NodeArray for open and closed set) and NodeArray A\* with Goal Bounding.

## 2 Basic A\*

### 2.1 Algorithm

The A\* is a search algorithm that uses a heuristic to find the best path between 2 nodes. Even though it's a basic algorithm, it shows relatively good performance when compared to other algorithms like djikstra. Despite being a basic algorithm, it's performance can be improved by using better data structures, or using other optimizations like we're going to show in the next sections.

### 2.2 Data

Table 1: Basic A\* performance (Path 1)

Method	Calls	Execution Time (ms)
A*Pathfinding.Search	1	21432.88
GetBestAndRemove	2890	114.64
AddToOpen	3020	2.54
SearchInOpen	28291	759.53
RemoveFromOpen	0	0
Replace	0	0
AddToClosed	2890	1.76
SearchInClosed	27990	20319.63
RemoveFromClosed	0	0

Table 2: Basic A\* performance (Path 2)

Method	Calls	Execution Time (ms)
A*Pathfinding.Search	1	21432.88
GetBestAndRemove	2890	114.64
AddToOpen	3020	2.54
SearchInOpen	28291	759.53
RemoveFromOpen	0	0
Replace	0	0
AddToClosed	2890	1.76
SearchInClosed	27990	20319.63
RemoveFromClosed	0	0

## 3 Basic A\* with tiebraking

### 3.1 Algorithm

This algorithm is basically the previous one, but we use tiebraking to break ties between nodes with the same f value. This way, it makes for better ordering of the nodes in the open set, and we can get the best node faster. When 2 nodes have the same f value, we use order the nodes by smallest h value. This way we first pick the node that is closer to the goal.

### 3.2 Data

Table 3: Basic A\* with tiebraking performance (Path 1)

Method	Calls	Execution Time (ms)
A*Pathfinding.Search	1	10029.02
GetBestAndRemove	1904	112.83
AddToOpen	1954	1.56
SearchInOpen	18564	260.07
RemoveFromOpen	0	0
Replace	0	0
AddToClosed	1904	1.27
SearchInClosed	18460	9524.91
RemoveFromClosed	0	0

Table 4: Basic A\* with tiebraking performance (Path 2)

Method	Calls	Execution Time (ms)
A*Pathfinding.Search	1	21970.11
GetBestAndRemove	2890	282.13
AddToOpen	3021	1.94
SearchInOpen	28291	768.2
RemoveFromOpen	0	0
Replace	0	0
AddToClosed	2890	1.61
SearchInClosed	27990	20761.34
RemoveFromClosed	0	0

## 4 NodeArray A\*

### 4.1 Algorithm

NodeArray A\* is an A\* implementation that uses a NodeArray to store the nodes. We use this array as our open and closed set, and we change the status property of the nodes when we add them to the open or closed set. This way, we can search for nodes in the open and closed set in constant time.

### 4.2 Data

Table 5: NodeArray A\* performance (Path 1)

Method	Calls	Execution Time (ms)
A*Pathfinding.Search	1	5.48
GetBestAndRemove	200	1.96
AddToOpen	216	1.13
SearchInOpen	235	0
RemoveFromOpen	0	0
Replace	0	0
AddToClosed	100	0.01
SearchInClosed	1009	0.04
RemoveFromClosed	0	0

Table 6: NodeArray A\* performance (Path 2)

Method	Calls	Execution Time (ms)
A*Pathfinding.Search	1	159.11
GetBestAndRemove	2890	38.36
AddToOpen	3019	9.31
SearchInOpen	28247	1.84
RemoveFromOpen	0	0
Replace	0	0
AddToClosed	2885	0.58
SearchInClosed	27932	1.73
RemoveFromClosed	0	0

## 5 NodeArray A\* with Goal Bounding

### 5.1 Algorithm

By using precomputation of the grid, we can make bounding boxes for each node and improve the NodeArray A\* algorithm. We do this by using djikstra to calculate fastest path from each node to all other nodes. This way, we know which direction we should choose when trying to go to a specific node. This optimization causes, sometimes, a heavy increase on the starting time, due to the precomputation, but it improves the runtime of the algorithm by a lot.

### 5.2 Data

Table 7: NodeArray A\* with Goal Bounding performance (Path 1)

Method	Calls	Execution Time (ms)
A*Pathfinding.Search	1	10.62
GetBestAndRemove	200	0.43
AddToOpen	216	0.34
SearchInOpen	235	0
RemoveFromOpen	0	0
Replace	0	0
AddToClosed	100	0.01
SearchInClosed	126	0
RemoveFromClosed	0	0

Table 8: NodeArray A\* with Goal Bounding performance (Path 2)

Method	Calls	Execution Time (ms)
A*Pathfinding.Search	1	17.12
GetBestAndRemove	158	0.40
AddToOpen	165	0.35
SearchInOpen	388	0.01
RemoveFromOpen	0	0
Replace	0	0
AddToClosed	158	0.03
SearchInClosed	282	0.01
RemoveFromClosed	0	0

## 6 Bonus Level

## 7 Conclusions

We can infer that A\* is pretty slow when compared to its optimizations.

Also, we can notice that adding pre-processing to the algorithm can improve its runtime by a lot, even though it takes some time to do it.