A bijection between 3-dimensional Catalan objects

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• (Bidimensional) Catalan numbers:

$$1, 1, 2, 5, 14, 42, 132, 429, \dots$$

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1	2	5	6

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7	8	11	12
3	4	9	10
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Ex. 2 4 3 1 6 7 5

• Shifted concatenation product *

• Shifted concatenation product \star

$$\sigma = 364512$$

$$\tau = 2314$$

$$\begin{bmatrix} 3 & 4 \\ 2 & 1 \end{bmatrix}$$

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Schützenberger involution

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$$\sigma = 364512 \rightarrow 215463$$

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Schützenberger involution

$$\sigma = 364512 \rightarrow 215463 \rightarrow 562314 = S(\sigma)$$

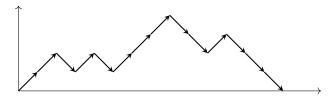
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$$\sigma = 364512 \rightarrow 215463 \rightarrow 562314 = S(\sigma)$$

Rem. $1234 \rightarrow 4321 \rightarrow 1234 \Rightarrow$ preserves the avoidance

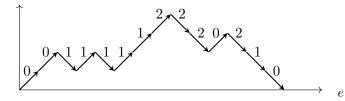
• Dyck path



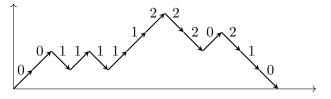
- Dyck path
- Weight wd(e) on each step e



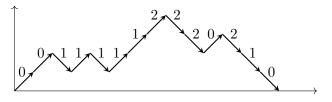
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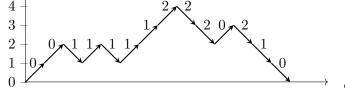
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- Weight wd(e) on each step e such that
 - 1. $wd(e) \leq wd(e+1)$ on upward slopes



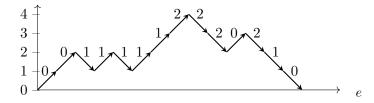
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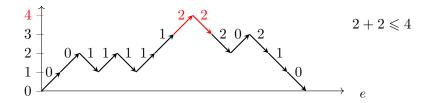
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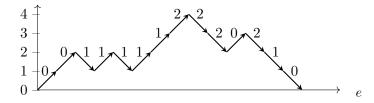
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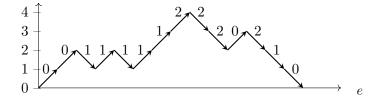
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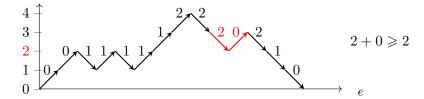
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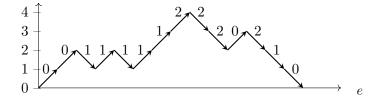
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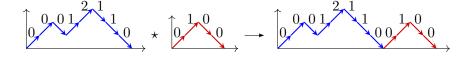


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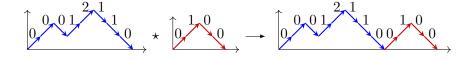


Concatenation product \star

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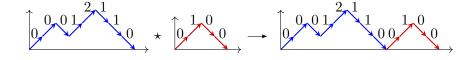


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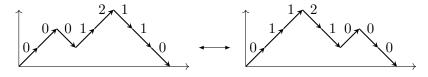


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- compatibility with the concatenation products and Schützenberger involutions
- the positions of up-steps correspond to the elements of the bottom word in the image permutation

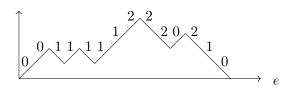
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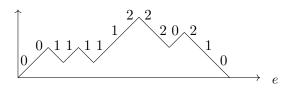
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Based on computational exploration (matching numbers of elements according to these statistics).

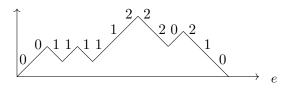


From a weighted Dyck path wd

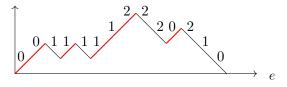
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- First: positions of up-steps, from left to right $\to Bot(\sigma)$
- Second: positions of down-steps, from right to left $\to Top(\sigma)$



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$$\beta \downarrow \qquad \qquad \downarrow \beta$$

$$\beta(wd) \qquad \longleftrightarrow \qquad S(\beta(wd))$$

S(wd)

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Handling of down-steps \rightarrow "same" as up-steps

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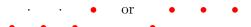
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- Idea: the (non-decreasing) weights could measure this number \rightarrow if the weight has increased, we insert u at this distance from the right-hand end

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 ⇒ the number of elements to the right of the 1st ascent increases
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 → if the weight has increased, we insert u at this distance
 from the right-hand end; if not, at the left-hand end.

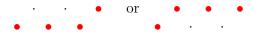
Pattern avoidance (1)

1234 patterns we are able to avoid:



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3 other configurations to avoid...

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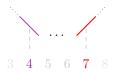


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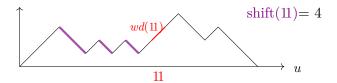




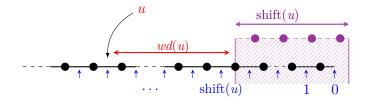
- \rightarrow Every down-step d to the left of a given up-step u should be inserted as a top element to the right of u (bottom element).
- \rightarrow shift(u) = # down-steps to the left of the up-step u

We add shift(u) to the weight wd(u) to obtain the insertion position.

Computation of the insertion position



Bottom word



Insertion position from the right

Recall: Same weight as previous one \Rightarrow no new leftmost ascent \Rightarrow jump

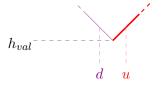
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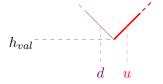
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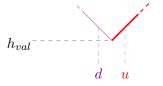
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Valley condition to decide the jump: $wd(d) + wd(u) \ge h_{val}$ If equality, u jumps to the left-hand end.

 \rightarrow General jump rule: compare wd(u) to its minimal possible value to decide the jump

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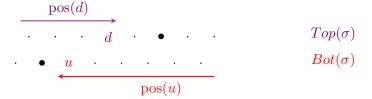


And the peak condition:



Valley condition and insertion position





Peak condition and pattern avoidance

Similarly, the peak condition should prevent the last pattern.

$$h_{peak}$$
 ---- u d

$$wd(d) + \frac{wd(u)}{wd(u)} \leqslant h_{peak}$$

 $u < d$

$$\xrightarrow[d]{\operatorname{pos}(d)}$$

Heuristical solution for every possible pattern situation... Is this over?

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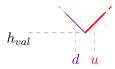
Is this over?

Not yet!

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$$wd(d) + wd(u) \geqslant h_{val}$$

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 \rightarrow Not bijective!

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Adjacent slopes (one upward, one downward)

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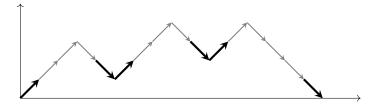
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Until now:



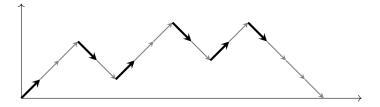
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Another idea... Not compatible with S!



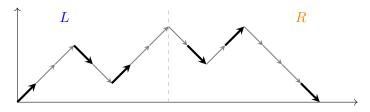
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Finally:



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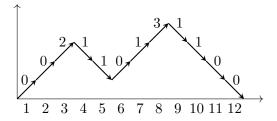
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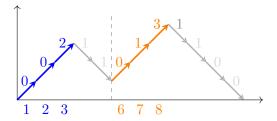
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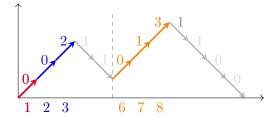
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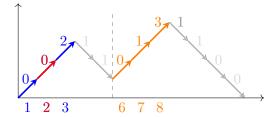
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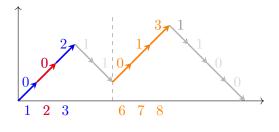
Rem. The image permutation being up-down is checked by a more thorough look at every distance involved.



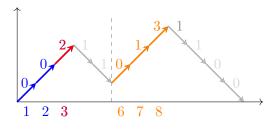




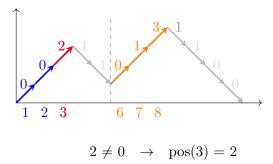


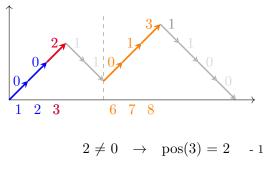


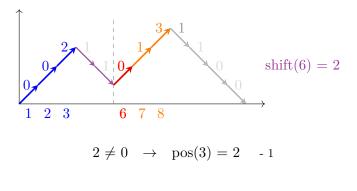
$$0 = 0 \rightarrow 2 \text{ jumps}$$

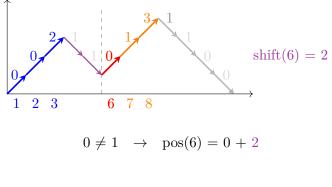


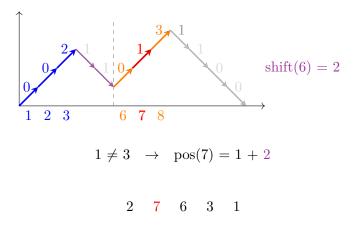
$$0 = 0 \rightarrow 2 \text{ jumps}$$

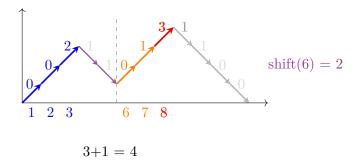




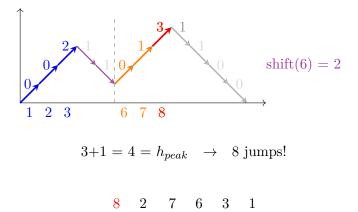


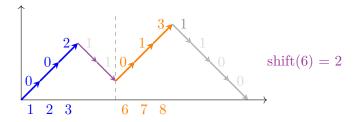




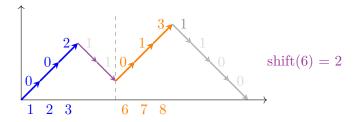


2 7 6 3 1





 $Top(\sigma)$ is obtained by scanning the down-steps from right to left



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Thank you for your attention!