

# Mathematical Preliminaries

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# Approximation Error

True value:  $a$

Approximate value:  $\tilde{a}$

$$a \approx \tilde{a}$$

Approximation Error  $\Delta a \equiv \tilde{a} - a$

Absolute Error is  $|\Delta a|$

An upper bound is any (known) number  $\Delta_a$  such that  $|\Delta a| \leq \Delta_a$

$$\tilde{a} - \Delta_a \leq a \leq \tilde{a} + \Delta_a \longrightarrow a = \tilde{a} \pm \Delta_a$$

# Acceleration in Sweden

Acceleration  $g$  in Sweden  $9.81666 \leq g \leq 9.82008$

Best guess  $\tilde{g} = \frac{9.82008 + 9.81666}{2} = 9.81837$

Uncertainty  $\Delta_g = \frac{9.82008 - 9.81666}{2} = 0.00171$

# Neutron Mass

NIST reports

$$\begin{aligned}\tilde{m} &= 1.674\,927\,498\,04 \times 10^{-27} \text{ kg} \\ \Delta_m &= 0.000\,000\,000\,95 \times 10^{-27} \text{ kg} \\ \tilde{m} &= 1.674\,927\,498\,04(95) \times 10^{-27} \text{ kg}\end{aligned}$$

(Exact uncertainty)  $|\Delta m| \leq \Delta_m = 0.000\,000\,000\,95 \times 10^{-27} \text{ kg}$

$$\tilde{m} - \Delta_m \leq m \leq \tilde{m} + \Delta_m$$

# Sensor Readings

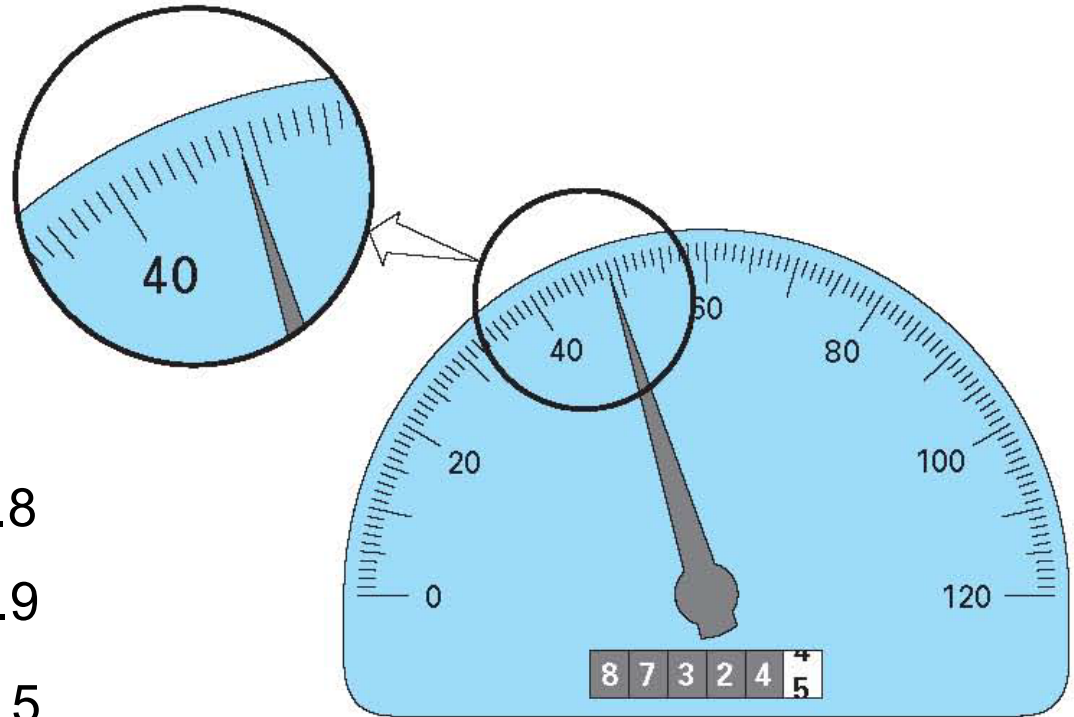
Visual inspection:

$$48 \leq v \leq 49$$

One person insists:  $v = 48.8$

Another insists:  $v = 48.9$

Commonly accepted:  $v = 48.5$



The estimated digit is one-half of the smallest scale division.

# Significant Digits

The significant digits of a number are those that can be used with confidence.

Significant digits are certain digits plus one estimated digit.

$$v = 48.5$$

Zeros are not always significant: 0.2021, 0.02021, 0.002021

Unclear: 202100 may have 4, 5 or 6 significant digits.

Resolve:  $2.02100 \times 10^5$  (6 significant digits).

# Round-off Errors

Computers may retain only limited numbers of digits.

Specific numbers:  $\sqrt{2}$ ,  $\pi$ ,  $e$ , ... have infinitely many significant digits.

$$\pi = 3.141592653589793238462643\dots$$

The omission of the remaining significant figures is called round-off error.

# Two Rounding Rules

Chopping:  $1.650 \approx 1.6$

Nearest:  $1.650 \approx 1.7$

			$x - \tilde{x}$	$x - \tilde{x}$
$x$	Chop	Nearest	Chop	Nearest
1.649	1.6	1.6	0.049	0.049
1.650	1.6	1.7	0.050	-0.050
1.651	1.6	1.7	0.051	-0.049
1.699	1.6	1.7	0.099	-0.001
1.749	1.7	1.7	0.049	0.049
1.750	1.7	1.8	0.050	-0.050



# Relative Error

$$\begin{aligned} l &= 1 \text{ cm} \pm 1 \text{ cm} \\ l &= 100 \text{ cm} \pm 1 \text{ cm} \end{aligned} \quad \forall a \neq 0 \quad \delta \equiv \frac{\Delta a}{a} = \frac{\tilde{a} - a}{a} \longrightarrow \tilde{a} = a(1 + \delta)$$

An upper bound is any (known) number  $\delta_a$  such that  $|\delta| \leq \delta_a$

$$|\delta| = \frac{|\Delta a|}{|a|} \longrightarrow |\Delta a| = |a| \cdot |\delta| \leq |a| \delta_a \longrightarrow \Delta_a = |a| \delta_a$$

$$\Delta_a = |a| \delta_a \approx |\tilde{a}| \delta_a \longrightarrow \tilde{a}(1 - \delta_a) \leq a \leq \tilde{a}(1 + \delta_a) \longrightarrow a = \tilde{a}(1 \pm \delta_a)$$

# Approximation Error of a Sum

$$\tilde{x}_i = x_i + \Delta x_i \quad \left| \Delta x_i \right| \leq \Delta_i$$

$$\tilde{x} = \tilde{x}_1 + \tilde{x}_2 + \dots + \tilde{x}_n = x + \Delta x$$

$$\Delta x = \Delta x_1 + \Delta x_2 + \dots + \Delta x_n$$

$$\left| \Delta x \right| \leq \Delta_1 + \Delta_2 + \dots + \Delta_n = \Delta_x$$

# Relative Error of a Sum

$$\text{All } x_i > 0; \quad |\Delta x_i| \leq \Delta_i; \quad \frac{|\Delta x_i|}{x_i} \leq \delta_i; \quad \delta_{\max} \equiv \max \delta_i$$

$$\delta \equiv \frac{\Delta x_1 + \Delta x_2 + \dots + \Delta x_n}{x_1 + x_2 + \dots + x_n}$$

$$|\delta| \leq \frac{x_1 \delta_1 + x_2 \delta_2 + \dots + x_n \delta_n}{x_1 + x_2 + \dots + x_n} \leq \delta_{\max}$$

# Relative Error of Product

$$\tilde{x} = \tilde{x}_1 \cdot \tilde{x}_2 \cdot \dots \cdot \tilde{x}_n; \quad \text{All } x_i > 0;$$

$$\delta \equiv \left| \frac{\Delta x}{x} \right| \leq \delta_1 + \delta_2 + \dots + \delta_n$$

# Selected Cases

$$\tilde{u} = k \cdot \tilde{x} \longrightarrow \delta_u = \delta_x; \quad \Delta u = k \cdot \Delta x$$

$$\tilde{u} = \tilde{x}_1 / \tilde{x}_2 \longrightarrow \delta_u = \delta_1 + \delta_2$$

$$\tilde{u} = \tilde{x}^m \longrightarrow \delta_u = m\delta_x$$

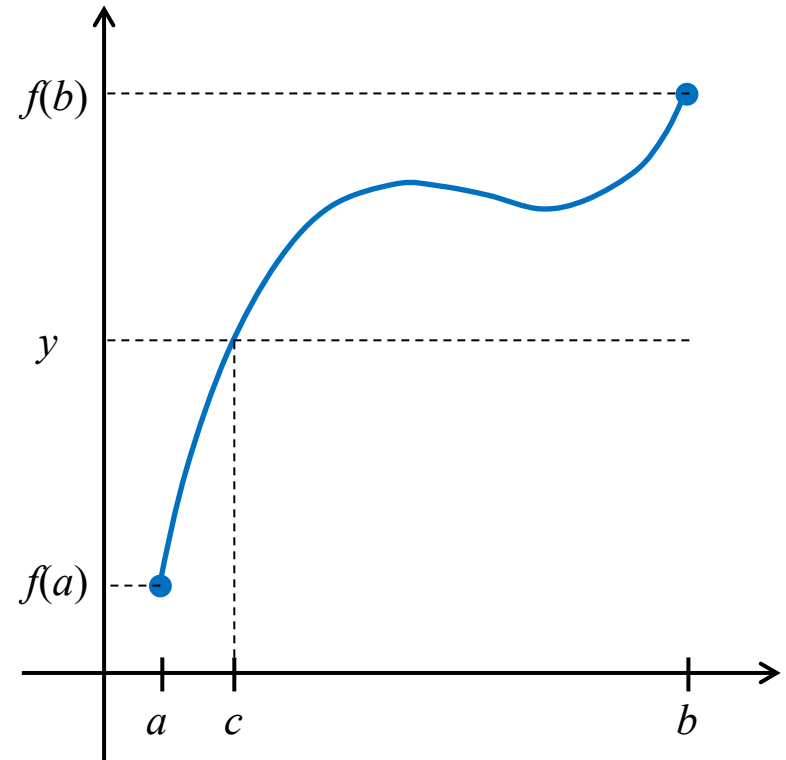
$$\tilde{u} = \sqrt[m]{\tilde{x}} \longrightarrow \delta_u = \frac{1}{m}\delta_x$$

# Intermediate Value Theorem

## Continuous function

$$\lim_{x \rightarrow c} f(x) = f(c)$$

Let  $f(x)$  be a continuous function on  $[a, b]$  then  $f$  realises every value between  $f(a)$  and  $f(b)$ . More precisely, if  $y$  is a number between  $a$  and  $b$ , then there exists a number  $c$ ,  $a \leq c \leq b$ , such that  $y = f(c)$ .



# Continuous Limit Theorem

Let  $f(x)$  be a continuous function in a neighborhood of  $x_0$  and  $\lim_{n \rightarrow \infty} x_n = x_0$  then

$$\lim_{n \rightarrow \infty} f(x_n) = f\left(\lim_{n \rightarrow \infty} x_n\right) = f(x_0)$$

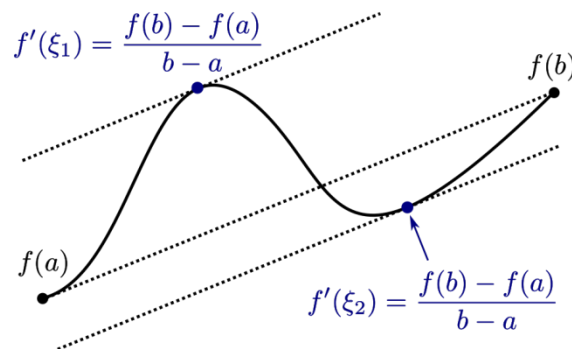
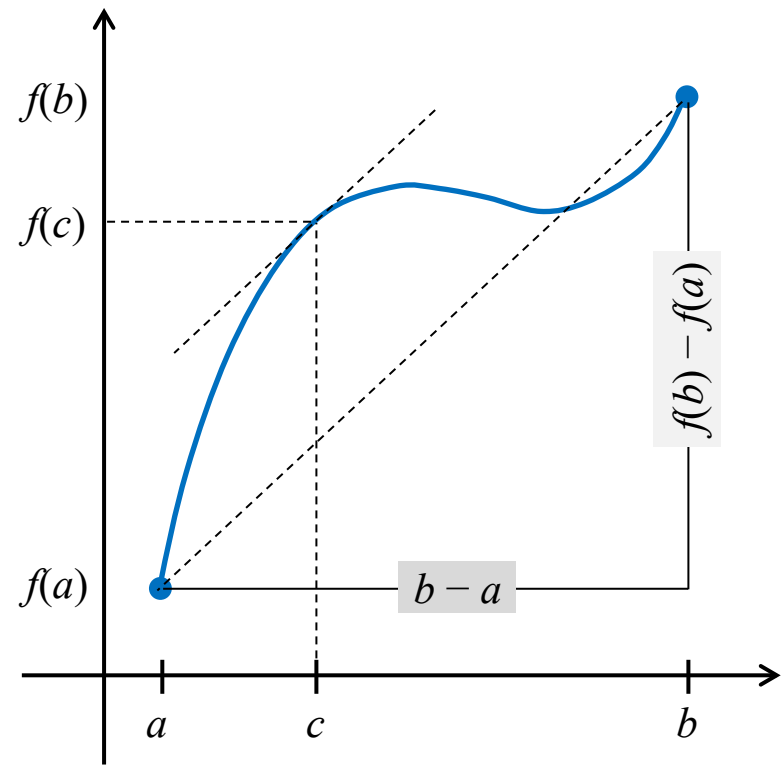
More precisely, limits may be brought inside continuous functions.

# Mean Value Theorem

Let  $f(x)$  be a continuously differentiable function on  $[a, b]$ .

Then there exists a number  $c$  between  $a$  and  $b$  such that

$$\frac{f(b) - f(a)}{b - a} = f'(c)$$

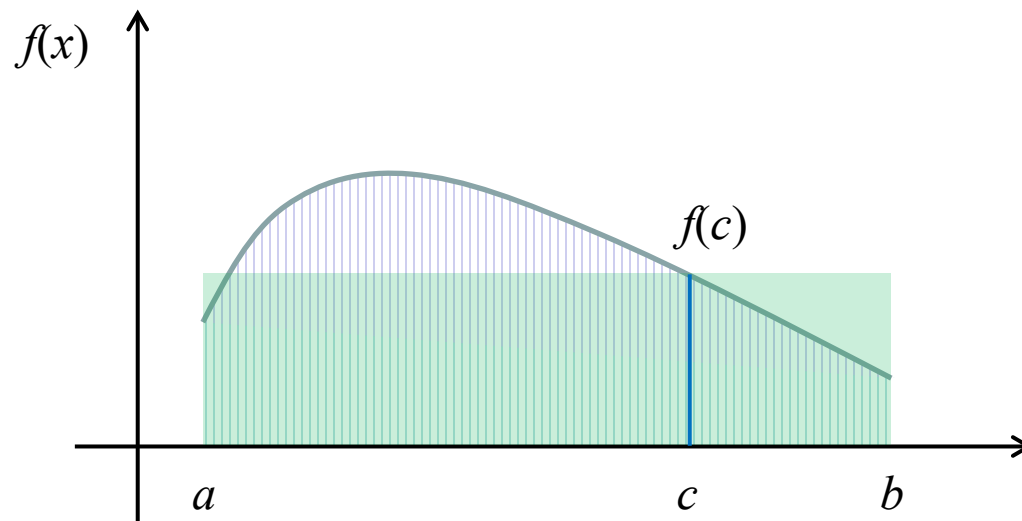


There could be several such numbers.



# Mean Value for Integrals

Let  $f(x)$  be a continuous function on a closed bounded interval  $[a, b]$  then there exists at least one number  $c$  such that 
$$\int_a^b f(x) dx = f(c)(b - a)$$



# Mean-Value Theorem

$$f(x), g(x) \in C[a, b] \quad g(x) \geq 0 \quad \forall x$$

$$\int_a^b f(x)g(x)dx = f(c)\int_a^b g(x)dx$$

$$\int_a^b f(x)dx = f(c)(b-a)$$

# Uncertainty Sources

Input data

$$x = 0.1$$

$$x = 9.81$$

$$x = \text{sqrt}(a)$$

$$x := \tilde{x} = 0.1 + \Delta(0.1)$$

$$x := \tilde{x} = g_{\text{true}} + \Delta(g)$$

$$x := \tilde{x} = \sqrt{a} + \Delta(\sqrt{a})$$

Representation

Experimental

Calculations

$$y = F(x)$$

$$y = F(x)$$

Error propagation

$$y + \Delta y = F(x + \Delta x)$$

How uncertainty in  $y$  is related to uncertainty in  $x$  ?

Assumption: no rounding errors (exact arithmetic)

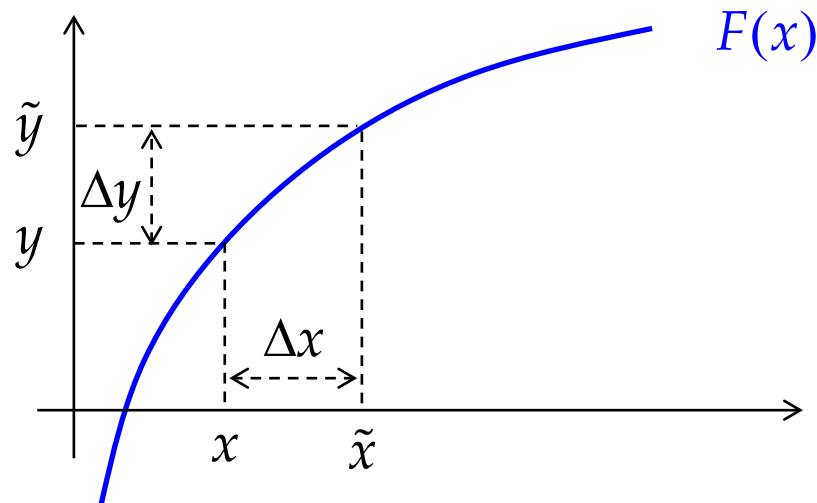
# Error Bounds

$$\tilde{x} = x + \Delta x \quad |\Delta x| \leq \Delta_x \quad \text{Error bound in input}$$

$$\tilde{y} = y + \Delta y \quad |\Delta y| \leq \Delta_y \quad \text{Error bound in output}$$

Notation

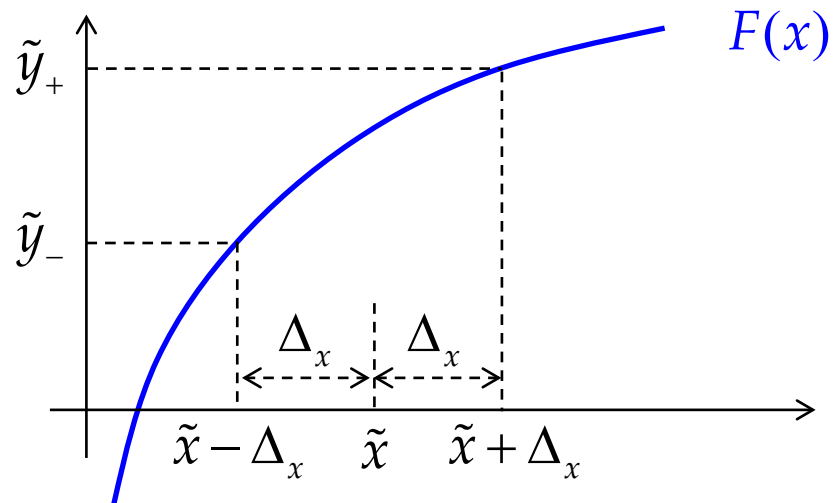
$$x = \tilde{x} \pm \Delta_x \quad y = \tilde{y} \pm \Delta_y$$



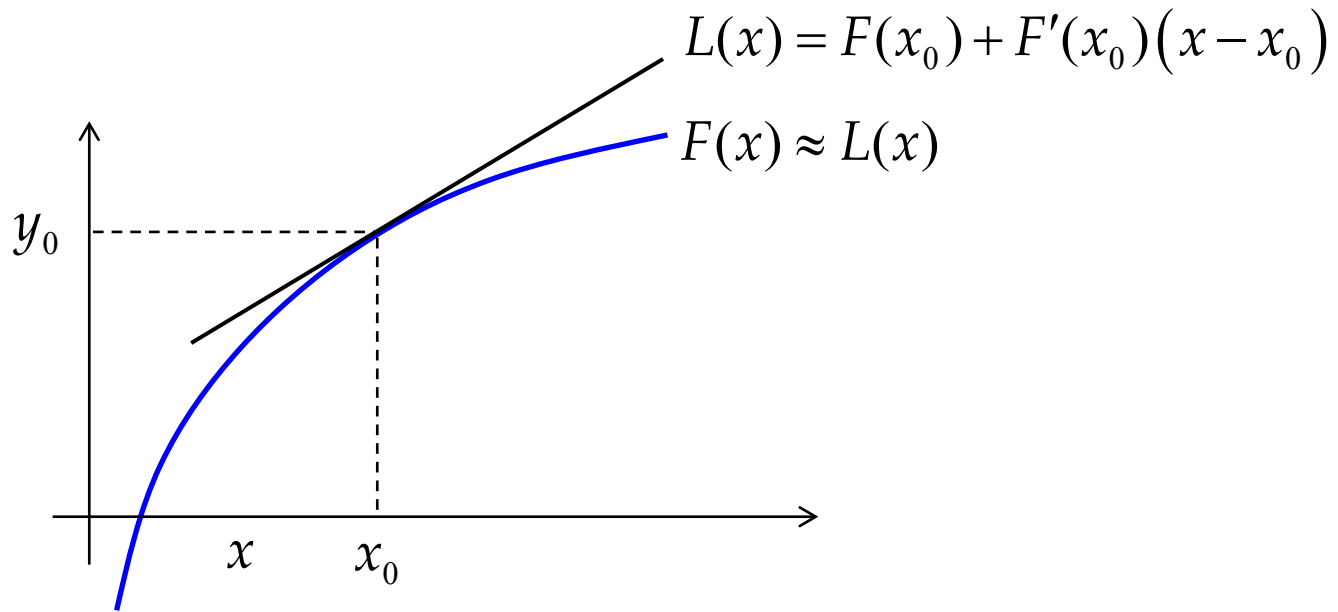
# Exact Error Estimate

$$\tilde{x} = x + \Delta x \quad |\Delta x| \leq \Delta_x$$

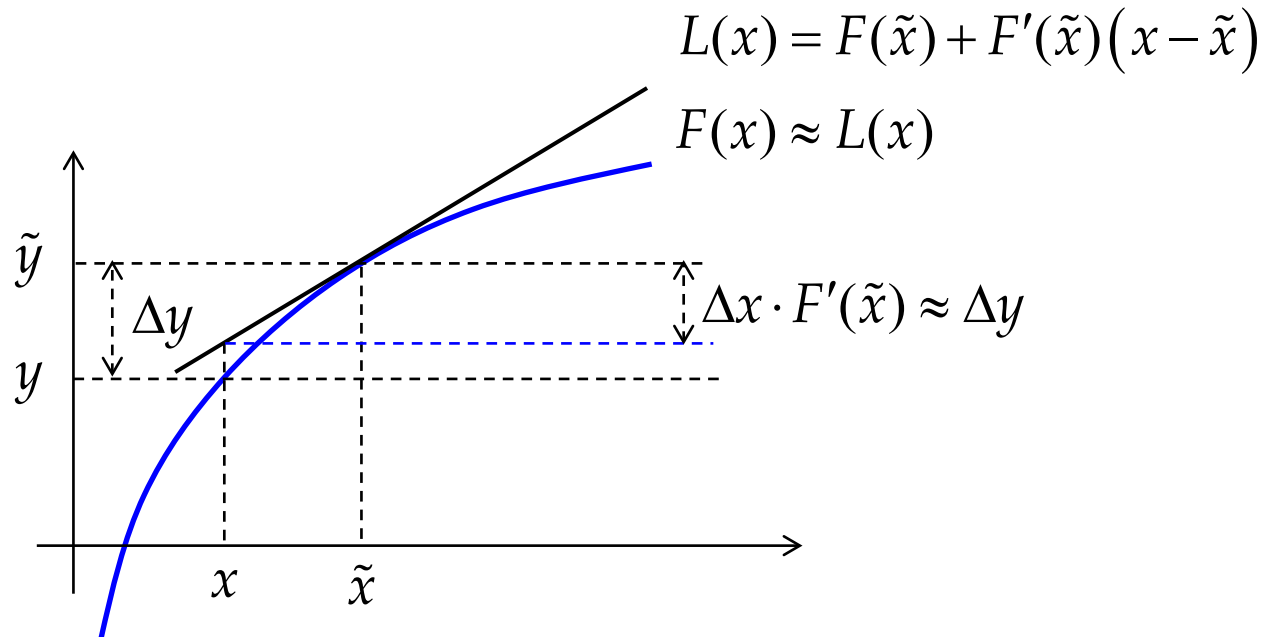
$$\tilde{y} = y + \Delta y \quad |\Delta y| \leq \Delta_y$$



# Linearization

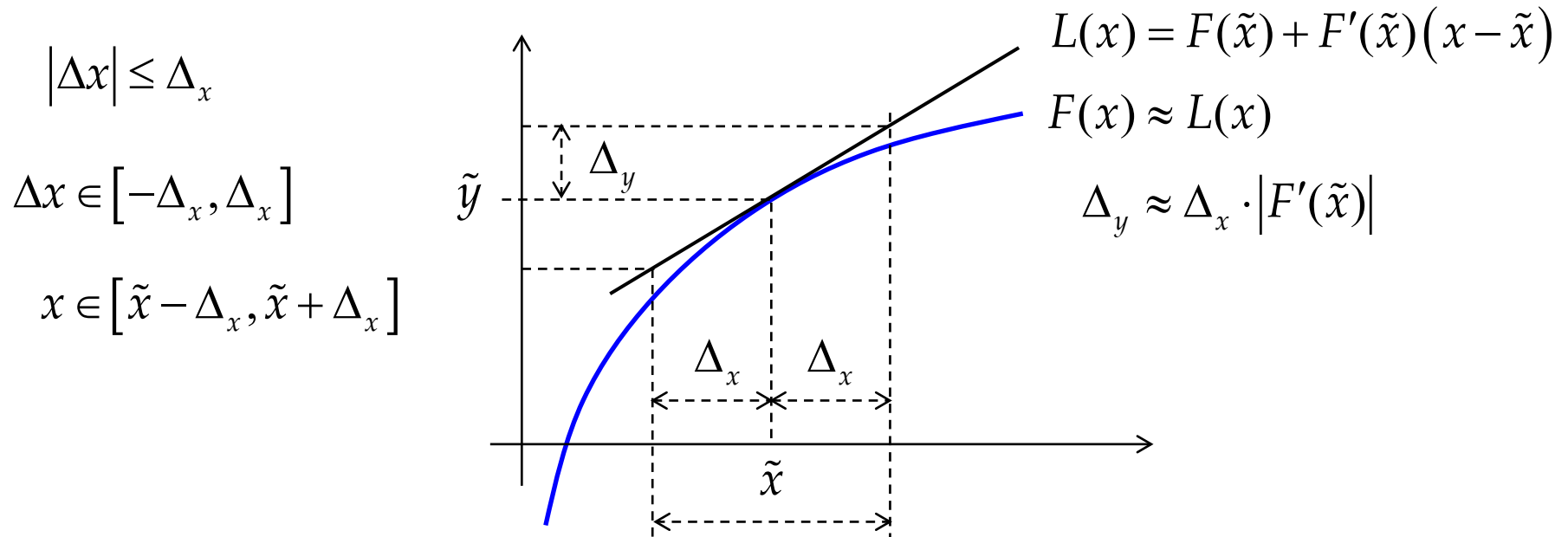


# Approximate Error Estimate



Error propagation  $\Delta y \approx \Delta x \cdot F'(\tilde{x})$

# Approximate Bounds



Error propagation  $\Delta_y \approx \Delta_x \cdot |F'(\tilde{x})|$



# Conditioning

- Well-conditioned: small errors in  $x$  induce small errors in  $y$   
Condition:  $|F'(\tilde{x})|$  is moderate.
- Ill-conditioned: small errors in  $x$  induce large errors in  $y$   
Condition:  $|F'(\tilde{x})|$  is large.

$$\Delta_y \approx \Delta_x \cdot |F'(\tilde{x})|$$

Relative

error in  $x$ :  $\delta_x \equiv \Delta_x / |x| \approx \Delta_x / |\tilde{x}|$

$$\delta_y \approx \frac{\Delta_y}{|\tilde{y}|} \approx \frac{|F'(\tilde{x})| \cdot \Delta_x}{|\tilde{y}|} = |F'(\tilde{x})| \frac{|\tilde{x}|}{|\tilde{y}|} \frac{\Delta_x}{|\tilde{x}|} = \kappa \delta_x$$

Relative

error in  $y$ :  $\delta_y \equiv \Delta_y / |y| \approx \Delta_y / |\tilde{y}|$

$$\kappa \equiv |F'(\tilde{x})| \frac{|\tilde{x}|}{|\tilde{y}|} \longrightarrow \delta_y \approx \kappa \delta_x$$

# Multivariate Functions

Bivariate

$$z = F(x, y) \quad \Delta z \approx \Delta x \cdot F_x(\tilde{x}, \tilde{y}) + \Delta y \cdot F_y(\tilde{x}, \tilde{y})$$

$$\Delta_z \approx \Delta_x \cdot |F_x(\tilde{x}, \tilde{y})| + \Delta_y \cdot |F_y(\tilde{x}, \tilde{y})|$$

$$|\Delta z| \leq \Delta_z$$

Multivariate

$$z = F(x_1, x_2, \dots, x_n)$$

$$\Delta z \approx \Delta x_1 F_{x_1}(\tilde{x}_1, \tilde{x}_2, \dots, \tilde{x}_n) + \Delta x_2 F_{x_2}(\tilde{x}_1, \tilde{x}_2, \dots, \tilde{x}_n) + \dots + \Delta x_n F_{x_n}(\tilde{x}_1, \tilde{x}_2, \dots, \tilde{x}_n)$$

$$\Delta_z \approx \Delta_{x_1} |F_{x_1}(\tilde{x}_1, \tilde{x}_2, \dots, \tilde{x}_n)| + \Delta_{x_2} |F_{x_2}(\tilde{x}_1, \tilde{x}_2, \dots, \tilde{x}_n)| + \dots + \Delta_{x_n} |F_{x_n}(\tilde{x}_1, \tilde{x}_2, \dots, \tilde{x}_n)|$$

$$|\Delta z| \leq \Delta_z$$

# Example

$$z = F(x, y) = \sqrt{1 + x^2 + y}$$

$$x = 1.0 \pm 0.1$$

$$y = 2.0 \pm 0.5$$

$$\tilde{x} = 1.0 \quad \Delta_x = 0.1$$

$$\tilde{y} = 2.0 \quad \Delta_y = 0.5$$

$$F_x(x, y) = \frac{x}{\sqrt{1 + x^2 + y}}$$

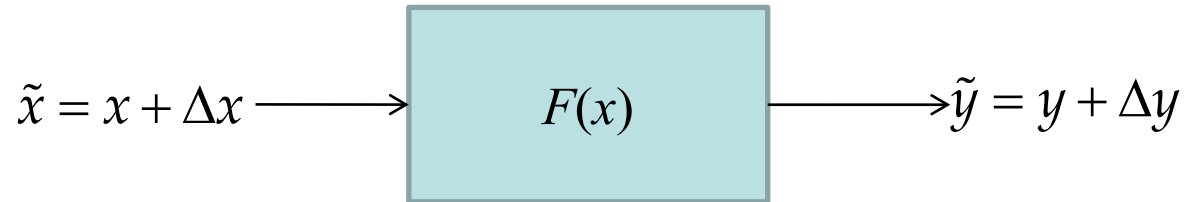
$$F_y(x, y) = \frac{1}{2\sqrt{1 + x^2 + y}}$$

$$F_x(\tilde{x}, \tilde{y}) = \frac{1}{\sqrt{1 + 1^2 + 2}} = 0.5 \quad F_y(\tilde{x}, \tilde{y}) = \frac{1}{2\sqrt{1 + 1^2 + 2}} = 0.25$$

$$\Delta_z \approx \Delta_x \cdot |F_x(\tilde{x}, \tilde{y})| + \Delta_y \cdot |F_y(\tilde{x}, \tilde{y})| = 0.1 \times 0.5 + 0.5 \times 0.25 = 0.175$$

$$\delta_z \approx \Delta_z / F(\tilde{x}, \tilde{y}) = 0.175 / 2 = 0.0875$$

# Perturbation Experiment 1D



- $F(x)$  is often:
  - Complicated, expensive, external code
- $\tilde{x} = x + \Delta x$  is input data:
  - Initial value, physical constant, problem parameter
- $\tilde{y} = y + \Delta y$  is output data:
  - Result, numbers, arrays

$$\tilde{y} = F(\tilde{x}) \quad y_{\text{exp}} = F(\tilde{x} + \Delta_x) \longrightarrow \Delta_y \approx |y_{\text{exp}} - \tilde{y}|$$

# Perturbation Experiment 2D

$$z = F(x, y)$$

1) Best guess  $\tilde{z} = F(\tilde{x}, \tilde{y})$

2) Exp1  $z_{1,\text{exp}} = F(\tilde{x} + \Delta_x, \tilde{y})$

3) Exp2  $z_{2,\text{exp}} = F(\tilde{x}, \tilde{y} + \Delta_y)$

4) Sum up  $E_z \approx |z_{1,\text{exp}} - \tilde{z}| + |z_{2,\text{exp}} - \tilde{z}|$

# Taylor's Theorem

$$f \in C^n[a, b] \text{ \& } f^{(n+1)} \text{ \exists on } (a, b) \quad \forall x, c \in [a, b]$$

$\exists \xi$  |  $\xi$  is between  $x$  and  $c$

$$f(x) = f(c) + \sum_{k=1}^n \frac{1}{k!} f^{(k)}(c) (x-c)^k + E_n(x)$$

$$E_n(x) = \frac{1}{(n+1)!} f^{(n+1)}(\xi) (x-c)^{n+1}$$

# Example

$$f(x) = \ln x; \quad a = 1, \quad b = 2, \quad c = a \quad f^{(k)}(x) = (-1)^{k-1} (k-1)! x^{-k}$$

$$\ln x = (x-1) - \frac{1}{2}(x-1)^2 + \frac{1}{3}(x-1)^3 - \dots + (-1)^{n-1} \frac{1}{n}(x-1)^n + E_n(x)$$

$$E_n(x) = \frac{1}{\xi^n} \frac{(-1)^n}{(n+1)} (x-1)^{n+1}; \quad 1 < \xi < x \longrightarrow |E_n(x)| \leq \frac{1}{n+1} (x-1)^{n+1}$$

# Estimating Accuracy

$$\ln 2 = 1 - \frac{1}{2} + \frac{1}{3} - \frac{1}{4} + \dots + (-1)^{n-1} \frac{1}{n} + E_n(2)$$

$$|E_n(2)| \leq \frac{1}{n+1} \leq 10^{-6} \longrightarrow n+1 \geq 10^6$$



# Another Form

$$f(x+h) = f(x) + \sum_{k=1}^n \frac{1}{k!} f^{(k)}(x) h^k + E_n(h)$$

$$E_n(h) = \frac{1}{(n+1)!} f^{(n+1)}(\xi) h^{n+1} \quad \xi \text{ is between } x \text{ and } x+h$$

$$E_n(h) = \frac{1}{(n+1)!} f^{(n+1)}(x + \theta h) h^{n+1} \quad 0 < \theta < 1$$

# Taylor's Theorem in 2D

$$f : \mathbb{R}^2 \rightarrow \mathbb{R}$$

$$\begin{aligned} f(a+h, b+k) = & f(a, b) + \\ & + h \frac{\partial f(a, b)}{\partial x} + k \frac{\partial f(a, b)}{\partial y} + \\ & + \frac{1}{2} h^2 \frac{\partial^2 f(a, b)}{\partial x^2} + hk \frac{\partial^2 f(a, b)}{\partial x \partial y} + \frac{1}{2} k^2 \frac{\partial^2 f(a, b)}{\partial y^2} + \\ & + E_2(h, k) \end{aligned}$$

# General 2D Form

$$f : \mathbb{R}^2 \rightarrow \mathbb{R}$$

$$f(a+h, b+k) = \sum_{i=0}^n \frac{1}{i!} \left( h \frac{\partial}{\partial x} + k \frac{\partial}{\partial y} \right)^i f(a, b) + E_n(h, k)$$

$$E_n(h, k) = \frac{1}{(n+1)!} \left( h \frac{\partial}{\partial x} + k \frac{\partial}{\partial y} \right)^{n+1} f(a + \theta h, b + \theta k) \quad 0 < \theta < 1$$

# Binominal Theorem

$$(x + y)^n = x^n + nx^{n-1}y + \dots \binom{n}{k} x^{n-k} y^k + \dots + nxy^{n-1} + y^n$$

$$\binom{n}{k} \equiv \frac{n!}{k!(n-k)!}; \quad \binom{n}{k} = \binom{n}{n-k}; \quad \binom{n}{0} = 1; \quad \binom{n}{1} = n.$$

$$(x + y)^n = \sum_{k=0}^n \binom{n}{k} x^{n-k} y^k = \sum_{k=0}^n \binom{n}{k} x^k y^{n-k}$$

# Big O

$$\{x_n\}, \{\alpha_n\}$$

$$x_n = O(\alpha_n) \quad \exists C, N: \quad |x_n| \leq C|\alpha_n| \quad \forall n \geq N$$

$$\left| \frac{x_n}{\alpha_n} \right| \leq C \quad \text{when } n \rightarrow \infty$$

# Little o

$$\{x_n\}, \{\alpha_n\}$$

$$x_n = o(\alpha_n) \quad \lim_{n \rightarrow \infty} \frac{x_n}{\alpha_n} = 0$$

$$\frac{n+1}{n^2} = O\left(\frac{1}{n}\right) \quad \frac{1}{n \ln n} = o\left(\frac{1}{n}\right) \quad e^{-n} = o\left(\frac{1}{n^2}\right)$$

# O Notation

$$\sin x = x - \frac{x^3}{6} + O(x^5) \quad (x \rightarrow 0)$$

$$\left| \sin x - x + \frac{x^3}{6} \right| \leq Cx^5 \quad \text{in an neighbourhood of } 0$$

# O Notation

$$f(x) = O(g(x)) \quad (x \rightarrow 0, \quad x \rightarrow x_0, \quad x \rightarrow \infty)$$

$$|f(x)| \leq C|g(x)|$$

$$f(x) = o(g(x)) \iff \lim_{x \rightarrow x_0} \frac{f(x)}{g(x)} = 0$$



# Orders of Convergence

$$x_n \xrightarrow{n \rightarrow \infty} L$$

At least  
linear

$$|x_{n+1} - L| \leq c |x_n - L| \quad (c < 1, \quad n \geq N)$$

Superlinear

$$|x_{n+1} - L| \leq \varepsilon_n |x_n - L| \quad (\varepsilon_n \rightarrow 0)$$

At least  
quadratic

$$|x_{n+1} - L| \leq c |x_n - L|^2$$

# Polynomials

$$p = p(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0 \quad a_n \neq 0$$

$$\deg p(x) \equiv n$$

$$\Pi_n \equiv \{p \mid \deg p \leq n\}$$

$$\forall p, q \in \Pi_n \longrightarrow p + q \in \Pi_n \quad \& \quad \lambda p \in \Pi_n$$

$$p(x) = d(x) \cdot q(x) + r(x) \quad \deg q < \deg p \quad \deg r < \deg p$$

# Horner's Method

$$p(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0 = \sum_{k=0}^n \left( a_k \prod_{j=1}^k x \right)$$

$$p(x) = \left( \left( \left( \dots \left( \left( x a_n + a_{n-1} \right) x + a_{n-2} \right) x + \dots + a_3 \right) x + a_2 \right) x + a_1 \right) x + a_0$$

```
p = a(n);  
for k = n-1:-1:0  
    p = p*x + a(k);  
end
```

```
p = a[n]  
for k in range(n-1,-1,-1):  
    p = p*x + a[k]
```

# Important

- Absolute/Relative Error
- Error Sources
- Significant Digits
- Propagation of Errors
- Linearization/Taylor's Theorem
- Taylor's Theorem in 2D
- Order of Convergence