COMP 8700 Fall 2020 Take-home Exam/Assignment 3

Due: Jan 6., 11:59pm

**7.4** Which of the following are correct?

**a**. False |= True.

**b**. True |= False.

**c**. (A ∧ B) |= (A ⇔ B).

**d**. A ⇔ B |= A ∨ B.

**e**. A ⇔ B |= ￢A ∨ B.

**f**. (A ∧ B) ⇒ C |= (A ⇒ C) ∨ (B ⇒ C).

**g**. (C ∨ (￢A ∧ ￢B)) ≡ ((A ⇒ C) ∧ (B ⇒ C)).

**h**. (A ∨ B) ∧ (￢C ∨￢D ∨ E) |= (A ∨ B).

**i**. (A ∨ B) ∧ (￢C ∨￢D ∨ E) |= (A ∨ B) ∧ (￢D ∨ E).

**j**. (A ∨ B)∧ ￢(A ⇒ B) is satisfiable.

**k**. (A ⇔ B) ∧ (￢A ∨ B) is satisfiable.

**l**. (A ⇔ B) ⇔ C has the same number of models as (A ⇔ B) for any fixed set of

proposition symbols that includes A, B, C.

Correct answers are:

a, e, g, h, j, k, i

**7.5** Prove each of the following assertions:

**a**. α is valid if and only if True |= α.

Answer:

According to the definition, True |= α means in all models where True is true, then also α is true. Since True is true in all models, so α is true in all models too, so α is valid as per the definition of validity.

If α is valid, it is true in all worlds, so anything entails α, thus True entails α

**b**. For any α, False |= α.

Answer:

False |= α means, in all the worlds where False is true, α is also true, since there is no model in which False is true, then it is uncontroversial that α is true in an empty set of models

**c**. α |= β if and only if the sentence (α ⇒ β) is valid.

Answer:

α |= β means in all worlds where αis true, then also β is true, so α |= β holds as α and β are both true. True => True

if sentence (α ⇒ β) is true, then α is False, or αis True and β is True as below:

|  |  |  |
| --- | --- | --- |
| α | β | α |= β |
| False | False | True |
| True | True | True |
| False | True | True |

In all the three cases, α |= β is true according to the definition of entailment explained above.

**d**. α ≡ β if and only if the sentence (α ⇔ β) is valid.

Answer:

α ≡ β means α |= βand β |= α, from preceding question c, we have proved α |= β iff sentence (α ⇒ β) is valid, so α ≡ β means (α ⇒ β) and (β ⇒ α), i.e. (α ⇔ β)

**e**. α |= β if and only if the sentence (α ∧ ￢β) is unsatisfiable.

Answer:

(α ∧ ￢β)

≡ ￢(￢α ∨ β) # De Morgan

≡ ￢ (α ⇒ β) # implication elimination

α |= β means in all worlds where α is true, then also βis true, whereas sentence ￢(α ⇒ β) is ￢True thus False, so unsatisfiable given α |= β

on the other hand, ￢ (α ⇒ β) is unsatisfiable means (α ⇒ β) is always valid due to the definition of negation, thus α |= β

**7.6** Prove, or find a counterexample to, each of the following assertions:

**a**. If α |= γ or β |= γ (or both) then (α ∧ β) |= γ

Answer:

α |= γ or β |= γ

≡ (α => γ) ∨ (β => γ)

≡ (￢α ∨ γ) ∨ (￢β ∨ γ)

≡ ￢α ∨ γ ∨ ￢β

≡ ￢α ∨ ￢β ∨ γ

≡ ￢(α ∧ β) ∨ γ

≡ (α ∧ β) => γ

So, it proves (α ∧ β) |= γ

**b**. If α |= (β ∧ γ) then α |= β and α |= γ.

Answer:

α |= (β ∧ γ)

≡ α => (β ∧ γ)

≡ ￢α ∨ (β ∧ γ) #implication elimination

≡ (￢α ∨ β) ∧ (￢α∨ γ) #distributivity of ∨ over ∧

≡ (α => β) ∧ (α=> γ)

So it proves α |= β and α |= γ

**c**. If α |= (β ∨ γ) then α |= β or α |= γ (or both).

Answer:

α |= (β ∨ γ)

≡ α => (β ∨ γ)

≡ ￢α ∨ (β ∨ γ)

≡ (￢α ∨ β) ∨ γ ≡ (α => β) ∨ γ # equation 1

≡ (￢α ∨γ) ∨ β ≡ (α => γ) ∨ β # equation 2

Then we can infer:

|  |  |  |  |
| --- | --- | --- | --- |
| β | γ | α => β | α => γ |
| True | True | Valid as βis true | Valid as γ is true |
| True | False | Valid as βis true |  |
| False | True |  | Valid as γ is true |
| False | False | Valid, as γ is false, according to equation 1:  (α => β) ∨ γ | Valid, as β is false, according to equation 2:  (α => γ) ∨ β |

So from the table, it proves α |= β or α |= γ (or both).

**7.7** Consider a vocabulary with only four propositions, A, B, C, and D. How many models

are there for the following sentences?

**a**. B ∨ C.

Answer: 12 models

**b**. ￢A∨￢B ∨￢C ∨ ￢D.

Answer: 15 models

**c**. (A ⇒ B) ∧ A∧ ￢B ∧ C ∧ D.

Answer: 0 models

(A ⇒ B) ∧ A∧ ￢B ∧ C ∧ D.

≡ (￢A ∨ B) ∧ A∧ ￢B ∧ C ∧ D

≡ (￢A ∧ A∧ ￢B ∧ C ∧ D) ∨ (B ∧ A∧ ￢B ∧ C ∧ D)

So, it’s unsatisfiable, as in the first disjunct, ￢A ∧ A is there, and in the second B ∧ ￢B, both are conjunctions of complementary variables.

**7.10** Decide whether each of the following sentences is valid, unsatisfiable, or neither. Verify

your decisions using truth tables or the equivalence rules of Figure 7.11 (page 249).

**a**. Smoke ⇒ Smoke

Answer: valid, as by implication elimination it becomes following sentence, which is always true

￢Smoke ∨ Smoke

As there must be one true between a pair of complementary literals.

**b**. Smoke ⇒ Fire

Answer: not valid, but satisfiable

Implication elimination: ≡ ￢Smoke ∨ Fire

Answer: valid

|  |  |  |  |
| --- | --- | --- | --- |
| Smoke | Fire | ￢Smoke | ￢Smoke ∨ Fire |
| T | T | F | T |
| F | T | T | T |
| F | F | T | T |
| T | F | F | F |

**c**. (Smoke ⇒ Fire) ⇒ (￢Smoke ⇒ ￢Fire)

Answer: not valid, but satisfiable

≡ (￢Smoke ∨ Fire) => (￢Smoke ⇒ ￢Fire) implication elimination

≡ (￢Smoke ∨ Fire) => (Smoke ∨ ￢Fire) contraposition

≡ ￢ (￢Smoke ∨ Fire) ∨(Smoke ∨ ￢Fire) implication elimination

≡ (Smoke ∧ ￢Fire) ∨(Smoke ∨ ￢Fire)

≡ (Smoke ∨Smoke ∨ ￢Fire) ∧ (￢Fire ∨ Smoke ∨ ￢Fire) distributivity of ∨ over ∧

≡ (Smoke ∨ ￢Fire) ∧ (￢Fire ∨ Smoke)

≡ (Smoke ∨ ￢Fire)

|  |  |  |  |
| --- | --- | --- | --- |
| Smoke | Fire | ￢Fire | Smoke ∨￢Fire |
| T | T | F | T |
| F | T | F | F |
| F | F | T | T |
| T | F | T | T |

**d**. Smoke ∨ Fire ∨ ￢Fire

Answer: valid

≡ Smoke ∨ (Fire ∨ ￢Fire) associativity of ∨

≡ Smoke ∨ True

Which is always true.

**e**. ((Smoke ∧ Heat ) ⇒ Fire) ⇔ ((Smoke ⇒ Fire) ∨ (Heat ⇒ Fire))

Answer: valid

≡(￢ (Smoke ∧ Heat ) ∨ Fire) ⇔ ((￢ Smoke ∨ Fire) ∨ (￢ Heat ∨ Fire)) implication elimination

≡( ￢ Smoke ∨ ￢ Heat ∨ Fire) ⇔ ((￢ Smoke ∨ Fire) ∨ (￢ Heat ∨ Fire)) De Morgan

≡( ￢ Smoke ∨ ￢ Heat ∨ Fire) ⇔ (￢ Smoke ∨ ￢ Heat ∨ Fire ∨ Fire) Associativity of ∨

≡( ￢ Smoke ∨ ￢ Heat ∨ Fire) ⇔ (￢ Smoke ∨ ￢ Heat ∨ Fire)

**f**. (Smoke ⇒ Fire) ⇒ ((Smoke ∧ Heat ) ⇒ Fire)

Answer: valid

≡( ￢ Smoke ∨ Fire) ⇒ (￢ (Smoke ∧ Heat ) ∨ Fire) implication elimination

≡( ￢ Smoke ∨ Fire) ⇒ ((￢Smoke ∨ ￢Heat ) ∨ Fire) De Morgan

≡( ￢ Smoke ∨ Fire) ⇒ (￢Smoke ∨ Fire ∨ ￢Heat ) associativity of∨

≡ ￢ ( ￢ Smoke ∨ Fire) ∨ (￢Smoke ∨ Fire ∨ ￢Heat ) implication elimination

≡ ( Smoke ∧ ￢Fire) ∨ (￢Smoke ∨ Fire ∨ ￢Heat ) De Morgan

≡ ( Smoke ∨ ￢Smoke ∨ Fire ∨ ￢Heat )∧(￢Fire ∨ ￢Smoke ∨ Fire ∨ ￢Heat ) distributivity

≡ ( Smoke ∨ ￢Smoke ∨ Fire ∨ ￢Heat )∧(￢Fire ∨ Fire ∨ ￢Smoke ∨ ￢Heat )

≡ true ∧ true

≡ true

**g**. Big ∨ Dumb ∨ (Big ⇒ Dumb)

Answer: valid

≡Big ∨ Dumb ∨ (￢Big ∨ Dumb)

≡Big ∨￢Big ∨ Dumb ∨ Dumb associativity of ∨

≡Big ∨￢Big ∨ Dumb

≡true

**7.12** Use resolution to prove the sentence ￢A∧￢B from the clauses in Exercise 7.20.

Answer:

From 7.20 we know

S1: A ⇔ (B ∨ E).

S4: E ⇒ B

According to Modus Ponens

A ⇔ (B ∨ E).

≡A ⇔ (B ∨ B) because E ⇒ B

≡A ⇔ B

≡(A ⇒ B) ∧ (B ⇒ A)

≡ (￢A ∨ B) ∧(￢B∨A)

≡ ((￢A ∨ B) ∧(￢B))∨ ((￢A ∨ B) ∧ A))

≡ (￢A ∧￢B) ∨ (B∧ A)

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| A | B | (￢A ∧￢B) | (B∧ A) | (￢A ∧￢B) ∨ (B∧ A) |
| T | T | F | T | T |
| T | F | F | F | F |
| F | T | F | F | F |
| F | F | T | F | T |

From 3, 5,6 in 7.20, we know

S3: C ∧ F ⇒ ￢B and S5: B ⇒ F and S6: B ⇒ C

≡(￢C ∨ ￢F∨￢B) ∧(￢B ∨ F) ∧(￢B ∨ C)

≡((￢C ∨ ￢F∨￢B) ∧(￢B) ∨ (￢C ∨ ￢F∨￢B) ∧F)) ∧(￢B ∨ C)

≡((￢C∧(￢B) ∨ ￢F∧(￢B)) ∨ (￢C∧F ∨ ￢B∧F)) ∧(￢B ∨ C)

≡((￢C∧(￢B) ∨ ￢F∧(￢B)) ∨ (￢C∧F ∨ ￢B∧F)) ∧(￢B) ∨

((￢C∧(￢B) ∨ ￢F∧(￢B)) ∨ (￢C∧F ∨ ￢B∧F)) ∧C

≡((￢C∧(￢B) ∨ ￢F∧(￢B)) ∨ (￢C∧F∧(￢B) ∨ ￢B∧F)) ∨

((￢F∧(￢B) ∧C) ∨ (￢B∧F∧C))

≡(￢B) ∧ (￢C ∨ ￢F ∨ (￢C∧F ∨ F) ∨(￢F∧C) ∨ (F∧C))

According to And-elimination (￢B) is True, then B is False

So, (￢A ∧￢B) ∨ (B∧ A)

≡ (￢A ∧￢B) ∨ False

≡ (￢A ∧￢B)

Proved.

**7.20** Convert the following set of sentences to clausal form.

Answer:

≡ (A ⇒ (B ∨ E)) ∧ ((B ∨ E) ⇒ A)

≡(￢A ∨ B ∨ E) ∧ ((￢B ∧ ￢E) ∨ A) implication elimination

≡(￢A ∨ B ∨ E) ∧ ((￢B∨ A) ∧ (￢E∨ A) ) distributivity of ∨ over ∧

S2: E ⇒ D.

Answer:

≡￢E ∨ D

S3: C ∧ F ⇒ ￢B.

Answer:

≡ ￢(C ∧ F)∨ ￢B

≡ ￢C ∨ ￢F∨￢B De Morgan

S4: E ⇒ B.

≡￢E ∨ B

S5: B ⇒ F.

≡￢B ∨ F

S6: B ⇒ C

≡￢B ∨ C

Give a trace of the execution of DPLL on the conjunction of these clauses.

(￢A ∨ B ∨ E) ∧ (￢B ∨ A) ∧ (￢E∨ A)

￢E ∨ D

￢C ∨ ￢F∨￢B

￢E ∨ B

￢B ∨ F

￢B ∨ C

Set E to false, then simplified as:

(￢A ∨ B) ∧ (￢B ∨ A)

￢C ∨ ￢F∨￢B

￢B ∨ F

￢B ∨ C

Set B to false:

￢A

Set A to false:

true

**8.2** Consider a knowledge base containing just two sentences: P(a) and P(b). Does this

knowledge base entail ∀x P(x)? Explain your answer in terms of models.

Answer:

Suppose a model contains 3 or more elements in the domain, the relation referred to by P holds solely for a and b, but not for the third element. So, it does not entail ∀x P(x).

**8.6** Which of the following are valid (necessarily true) sentences?

**a**. (∃x x=x) ⇒ (∀ y ∃z y =z).

Answer:

valid, if we do not impose unique-names assumption.

**b**. ∀x P(x) ∨￢P(x).

Answer: valid.

**c**. ∀ x Smart(x) ∨ (x=x).

Answer: valid, because (x=x) is always true.

**8.9** This exercise uses the function MapColor and predicates In(x, y), Borders(x, y), and

Country(x), whose arguments are geographical regions, along with constant symbols for

various regions. In each of the following we give an English sentence and a number of candidate

logical expressions. For each of the logical expressions, state whether it (1) correctly

expresses the English sentence; (2) is syntactically invalid and therefore meaningless; or (3)

is syntactically valid but does not express the meaning of the English sentence.

**a**. Paris and Marseilles are both in France.

(i) In(Paris ∧ Marseilles, France ).

(ii) In(Paris, France ) ∧ In(Marseilles, France ).

(iii) In(Paris, France ) ∨ In(Marseilles, France ).

**b**. There is a country that borders both Iraq and Pakistan.

(i) ∃ c Country(c) ∧ Border (c, Iraq) ∧ Border (c, Pakistan).

(ii) ∃ c Country(c) ⇒ [Border (c, Iraq) ∧ Border (c, Pakistan)].

(iii) [∃ c Country(c)] ⇒ [Border (c, Iraq) ∧ Border (c, Pakistan)].

(iv) ∃ c Border (Country(c), Iraq ∧ Pakistan).

**c**. All countries that border Ecuador are in South America.

(i) ∀c Country(c) ∧ Border (c,Ecuador ) ⇒ In(c, SouthAmerica).

(ii) ∀ c Country(c) ⇒ [Border (c,Ecuador ) ⇒ In(c, SouthAmerica)].

(iii) ∀ c [Country(c) ⇒ Border (c,Ecuador )] ⇒ In(c, SouthAmerica).

(iv) ∀c Country(c) ∧ Border (c,Ecuador ) ∧ In(c, SouthAmerica).

**d**. No region in South America borders any region in Europe.

(i) ￢[∃ c, d In(c, SouthAmerica) ∧ In(d, Europe) ∧ Borders(c, d)].

(ii) ∀ c, d [In(c, SouthAmerica) ∧ In(d, Europe)] ⇒ ￢Borders(c, d)].

(iii) ￢∀ c In(c, SouthAmerica) ⇒ ∃d In(d, Europe)∧ ￢Borders(c, d).

(iv) ∀ c In(c, SouthAmerica) ⇒ ∀d In(d, Europe) ⇒ ￢Borders(c, d).

**e**. No two adjacent countries have the same map color.

(i) ∀ x, y ￢Country(x) ∨ ￢Country(y)∨ ￢Borders(x, y) ∨

￢(MapColor (x) = MapColor (y)).

(ii) ∀ x, y (Country(x) ∧ Country(y) ∧ Borders(x, y) ∧ ￢(x = y)) ⇒

￢(MapColor (x) = MapColor (y)).

(iii) ∀ x, y Country(x) ∧ Country(y) ∧ Borders(x, y) ∧

￢(MapColor (x) = MapColor (y)).

(iv) ∀ x, y (Country(x) ∧ Country(y) ∧ Borders(x, y)) ⇒ MapColor (x \_= y).

**8.28** Consider a first-order logical knowledge base that describes worlds containing people,

songs, albums (e.g., “Meet the Beatles”) and disks (i.e., particular physical instances of CDs).

The vocabulary contains the following symbols:

CopyOf (d, a): Predicate. Disk d is a copy of album a.

Owns(p, d): Predicate. Person p owns disk d.

Sings(p, s, a): Albuma includes a recording of song s sung by person p.

Wrote(p, s): Person p wrote song s.

McCartney, Gershwin, BHoliday, Joe, EleanorRigby, TheManILove, Revolver :

Constants with the obvious meanings.

**13.8** Given the full joint distribution shown in Figure 13.3, calculate the following:

**a**. **P**(toothache) .

**b**. **P**(Cavity) .

**c**. **P**(Toothache | cavity) .

**d**. **P**(Cavity | toothache ∨ catch) .

**13.17** Show that the statement of conditional independence

**P**(X, Y | Z) = **P**(X | Z)**P**(Y | Z)

is equivalent to each of the statements

**P**(X | Y,Z) = **P**(X |Z) and **P**(B | X,Z) = **P**(Y | Z) .

Exercise 14.14 (a) (b) (c)iuj

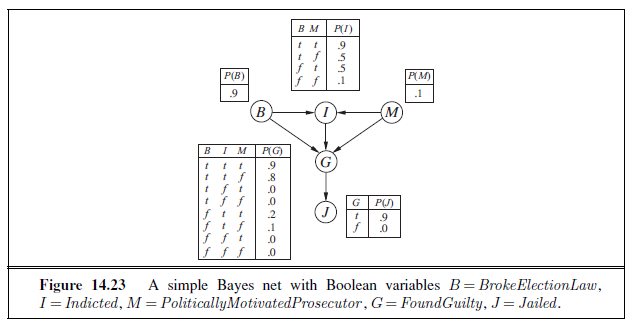
**14.14** Consider the Bayes net shown in Figure 14.23.

**a**. Which of the following are asserted by the network *structure*?

(i) **P**(B, I,M) = **P**(B)**P**(I)**P**(M).

(ii) **P**(J |G) = **P**(J | G, I).

(iii) **P**(M | G,B, I) = **P**(M | G,B, I, J).



**b**. Calculate the value of P(b, i,￢m, g, j).

**c**. Calculate the probability that someone goes to jail given that they broke the law, have

been indicted, and face a politically motivated prosecutor.