

# Set Theory

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# Chapter 1

## Axioms and Operations

### 1.1 Basic Axioms

**Axiom I (Extensionality).** For any sets  $x$  and  $y$ , if for any set  $z$ , we have  $z \in x$  if and only if  $z \in y$ , then we say that  $x$  and  $y$  are **equal**, denoted  $x = y$ .

**Axiom II (Empty Set).** There is a set  $x$  such that  $y \notin x$  for each set  $y$ . The set  $x$  is called the **empty set** and is denoted by  $\emptyset$ .

**Axiom III (Pairing).** For any sets  $x$  and  $y$ , there is a set  $w$  such that for each set  $z \in w$ , either  $z = x$  or  $z = y$  holds. The set  $w$  is called the **pair set** of  $x$  and  $y$  and is denoted by  $\{x, y\}$ . If  $x = y$ , then we write  $\{x\}$  for short.

**Example.** By axiom of pairing,  $\{\emptyset\}$  and  $\{\emptyset, \{\emptyset\}\}$  are sets.

**Axiom IV (Power Set).** For any set  $x$ , there exists a set  $y$  such that for any set  $z$ ,  $z \in y$  if and only if  $z \subseteq x$ . The set  $y$  is called the **power set** of  $x$  and is denoted by  $\mathcal{P}(x)$ .

**Axiom V (Subset).** Let  $\phi(z)$  be a first-order formula such that  $z$  is the only free variable in  $\phi$ . For any set  $x$ , there exists a set  $y$  such that for any set  $z$ ,  $z \in y$  if and only if both  $z \in x$  and  $\phi(z)$  holds. The set  $y$  will be denoted by

$$y = \{z \in x : \phi(z)\}.$$

**Theorem 1.1.** There is no set to which every set belongs. That is, for any set  $x$ , there exists a set  $y$  such that  $y \notin x$ .

*Proof.* Let  $y = \{z \in x : z \notin z\}$ . Then we have  $y \in y$  if and only if  $y \in x$  and  $y \notin y$ . If  $y \in x$ , then  $y \in y$  if and only if  $y \notin y$ , contradiction. Thus,  $y \notin x$ .  $\square$

## 1.2 Arbitrary Unions and Intersections

**Axiom VI (Union).** For any set  $x$ , there exists a set  $y$  whose elements are exactly the members of the members of  $x$ . That is,

$$\forall x \exists y \forall z (z \in y \leftrightarrow \exists w (w \in x \wedge z \in w)).$$

The set  $y$  is called the **union** of  $x$ , denoted by  $\bigcup x$ .

**Theorem 1.2.** For any nonempty set  $x$ , there exists a unique set  $y$  such that for any set  $z$ ,  $z \in y$  if and only if  $z$  belongs to every member of  $x$ .

*Proof.* Since  $x$  is nonempty, there is a member  $w_0$  of  $x$ . Then by a subset axiom there exists a set  $y$  such that

$$y = \{z \in w_0 : \forall w (w \in x \rightarrow z \in w)\},$$

and uniqueness of  $y$  follows from extensionality. □

**Definition 1.3.** For any nonempty set  $x$ , we define the **intersection** of  $x$  as the set  $y$  such that for any set  $z$ ,  $z \in y$  if and only if  $z$  belongs to every member of  $x$ . Let  $\bigcap x$  denote the intersection of  $x$ .

**Definition 1.4.** For any sets  $x$  and  $y$ , we define

$$\begin{aligned} x \cup y &= \bigcup \{x, y\} \\ x \cap y &= \bigcap \{x, y\}. \end{aligned}$$

# Chapter 2

## Relations and Functions

### 2.1 Ordered Pairs

**Definition 2.1.** For sets  $x$  and  $y$ , we define

$$\langle x, y \rangle = \{\{x\}, \{x, y\}\}.$$

**Lemma 2.2.** Let  $x, y, y'$  be sets. If  $\{x, y\} = \{x, y'\}$ , then  $y = y'$ .

*Proof.* Suppose that  $y \neq y'$ . Since  $y \in \{x, y\} = \{x, y'\}$  and  $y \neq y'$ , we have  $y = x$ . Then we have  $y' \in \{x, y'\} = \{x, y\} = \{x\}$ , implying  $y' = x = y$ , contradiction. Thus,  $y = y'$ .  $\square$

**Theorem 2.3.** For sets  $x, x', y, y'$ , we have

$$\langle x, y \rangle = \langle x', y' \rangle$$

if and only if  $x = x'$  and  $y = y'$ .

*Proof.*  $(\Leftarrow)$  Straightforward.  $(\Rightarrow)$  Suppose that  $x \neq x'$ . Since

$$\{\{x\}, \{x, y\}\} = \{\{x'\}, \{x', y'\}\},$$

either  $\{x\} = \{x', y'\}$  or  $\{x\} = \{x'\}$  holds. For both cases we all have  $x' \in \{x\}$ , implying  $x' = x$ , contradiction. Hence we have  $x = x'$ , and it follows that  $\{x\} = \{x'\}$ , implying  $\{x, y\} = \{x', y'\}$ , and thus  $y = y'$ .  $\square$

**Lemma 2.4.** If  $x, y \in C$ , then  $\langle x, y \rangle \in \mathcal{P}(\mathcal{P}(C))$ .

*Proof.* Since  $\{x\}$  and  $\{y\}$  are subsets of  $C$ , we have  $\{x\}, \{x, y\} \in \mathcal{P}(C)$ . It follows that  $\{\{x\}, \{x, y\}\}$  is a subset of  $\mathcal{P}(C)$ , implying

$$\langle x, y \rangle = \{\{x\}, \{x, y\}\} \in \mathcal{P}(\mathcal{P}(C)).$$

$\square$

**Theorem 2.5.** For any sets  $A$  and  $B$ , there is a set whose members are exactly the pairs  $\langle x, y \rangle$  with  $x \in A$  and  $y \in B$ .

*Proof.* Since  $x, y \in A \cup B$ , the set of pairs  $\langle x, y \rangle$  with  $x \in A$  and  $y \in B$  can be given by

$$\{z \in \mathcal{P}(\mathcal{P}(A \cup B)) : z = \langle x, y \rangle \text{ for some } x \in A \text{ and } y \in B\}.$$

$\square$

**Definition 2.6.** For any sets  $A$  and  $B$ , the **Cartesian product** of  $A$  and  $B$ , denoted by  $A \times B$ , is the set whose members are exactly the pairs  $\langle x, y \rangle$  with  $x \in A$  and  $y \in B$ .

## 2.2 Relations

**Definition 2.7.** A **relation** is a set of ordered pairs. For any relation  $R$ , the **domain** and the **range** of  $R$ , denoted by  $\text{dom}(R)$  and  $\text{ran}(R)$ , respectively, are defined as follows.

- $\text{dom}(R)$  is the collection of sets  $x$  with  $\langle x, y \rangle \in R$  for some  $y$ .
- $\text{ran}(R)$  is the collection of sets  $y$  with  $\langle x, y \rangle \in R$  for some  $x$ .

**Definition 2.8.** Let  $R$  and  $S$  be relations and let  $X$  be a set.

- The **inverse** of  $R$ , denoted by  $R^{-1}$ , is the set of all pairs  $\langle y, x \rangle$  with  $\langle x, y \rangle \in R$ .
- The **restriction** of  $R$  to  $X$ , denoted by  $R \upharpoonright X$ , is the set of all pairs  $\langle x, y \rangle \in R$  with  $x \in X$ .
- The **composition** of  $R$  and  $S$ , denoted by  $R \circ S$ , is the set of all pairs  $\langle x, z \rangle$  with  $\langle x, y \rangle \in R$  and  $\langle y, z \rangle \in S$ .

**Definition 2.9.** A **function** is a relation  $f$  such that for any set  $x \in \text{dom}(f)$ , there exists a unique set  $y$  such that  $\langle x, y \rangle \in f$ . The unique set  $y$  with respect to  $x$  is called the **value** of  $f$  at  $x$  and is denoted  $f(x)$ .

- We say that  $f$  is a function from  $A$  to  $B$ , denoted by  $f : A \rightarrow B$ , if  $\text{dom}(f) = A$  and  $\text{ran}(f) \subseteq B$ .
- $f$  is **one-to-one** if for any  $y \in \text{ran}(f)$ , there exists a unique set  $x \in \text{dom}(f)$  with  $f(x) = y$ .

**Definition 2.10.** For any sets  $A$  and  $B$ , the set of functions from  $A$  to  $B$  is denoted by  $B^A$ .

## 2.3 Equivalence Relations and Ordering Relations

**Definition 2.11.** Let  $A$  be a set. An **equivalence relation** on  $A$  is a relation  $R \subseteq A \times A$  that satisfies the following three conditions.

- Reflexivity:  $\langle x, x \rangle \in R$  for any  $x \in A$ .
- Symmetry:  $\langle x, y \rangle \in R$  implies  $\langle y, x \rangle \in R$  for any  $x, y \in A$ .
- Transitivity:  $\langle x, y \rangle \in R$  and  $\langle y, z \rangle \in R$  implies  $\langle x, z \rangle \in R$  for any  $x, y, z \in A$ .

# Chapter 3

## Natural Numbers

### 3.1 Inductive Sets

**Definition 3.1.** The **successor** of a set  $x$ , denoted  $x^+$ , is defined by

$$x^+ = x \cup \{x\}.$$

We say that a set  $A$  is **inductive** if  $\emptyset \in A$  and for any  $x \in A$ , we have  $x^+ \in A$ .

**Axiom VII (Infinity).** There exists an inductive set.

**Definition 3.2.** A **natural number** is a set belonging to all inductive sets. The set of natural numbers is denoted by  $\omega$ .

**Theorem 3.3.**  $\omega$  is inductive.

*Proof.* First,  $\emptyset \in \omega$  since  $\emptyset$  belongs to all inductive sets by definition. For any set  $x \in \omega$ ,  $x$  belongs to all inductive sets, implying that  $x^+$  belongs to all inductive sets, and thus  $x^+ \in \omega$ . Thus,  $\omega$  is inductive.  $\square$

**3.2 Recursion**

**3.3 Arithmetic**

**3.4 Ordering**