

Set Theory

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Chapter 1

Axioms and Operations

1.1 Basic Axioms

For sets x and y , we write $x \in y$ to say that x is an element of y , and we write $x = y$ to say that x and y are equal. Furthermore, we define

$$\begin{aligned}x \notin y &\Leftrightarrow \neg(x \in y) \\x \neq y &\Leftrightarrow \neg(x = y).\end{aligned}$$

Axiom I (Extensionality). Two sets are equal if they have exactly the same elements. Formally,

$$\forall x \forall y (\forall z (z \in x \leftrightarrow z \in y) \rightarrow x = y).$$

Definition 1.1. Let x and y be sets. We say that x is a **subset** of y , denoted by $x \subseteq y$, if every element of x belongs to y . Formally,

$$x \subseteq y \Leftrightarrow \forall z (z \in x \rightarrow z \in y).$$

Furthermore, x is a **proper subset** of y , denoted by $x \subsetneq y$, if both $x \subseteq y$ and $x \neq y$ hold.

Definition 1.2. The **empty set**, denoted by \emptyset , is the set that has no elements.

Axiom II (Pairing). For any two sets x and y , there is a set that consists of exactly x and y . Formally,

$$\forall x \forall y \exists z \forall w (w \in z \leftrightarrow (w = x \vee w = y)).$$

Definition 1.3. The **pair set** of two sets x and y , denoted by $\{x, y\}$, is the set that consists of exactly x and y .

Axiom III (Power Set). For any set x , there is a set whose members are exactly the subsets of x . Formally,

$$\forall x \exists y \forall z (z \subseteq x \leftrightarrow z \in y).$$

Definition 1.4. The **power set** of a set x , denoted by $\mathcal{P}(x)$, is the set that consists of exactly the subsets of x .

Axiom IV (Separation Scheme). Let $\phi(z)$ be a formula. For any set x , there exists a set y such that for any set z , we have $z \in y$ if and only if both $z \in x$ and $\phi(z)$ hold. Formally,

$$\forall x \exists y \forall z (z \in y \leftrightarrow z \in x \wedge \phi(z)).$$

Definition 1.5. Let x, y be sets and let $\phi(z)$ be a formula. If for any set z , we have $z \in y$ if and only if $z \in x$ and $\phi(z)$, then we write

$$y = \{z \in x : \phi(z)\}.$$

Theorem 1.6. There is no set to which every set belongs. Formally,

$$\forall x \exists y (y \notin x).$$

Proof. Let x be a set and let $y = \{z \in x : z \notin z\}$. Then

$$y \in y \quad \Leftrightarrow \quad y \in x \wedge y \notin y.$$

If $y \in x$, then

$$y \in y \quad \Leftrightarrow \quad y \notin y,$$

contradiction. Thus $y \notin x$, which completes the proof. \square

Axiom V (Union). For any set x , there exists a set whose elements are exactly the elements of the elements of x . Formally,

$$\forall x \exists y \forall z (z \in y \leftrightarrow \exists w (w \in x \wedge z \in w)).$$

Definition 1.7. Let x be a set.

- We define the **union** of x , denoted by $\bigcup x$, to be the set that consists of the sets that belongs to at least one element of x . Formally, for any set z we have

$$z \in \bigcup x \quad \Leftrightarrow \quad \exists w (w \in x \wedge z \in w).$$

- If x is nonempty, we define the **intersection** of x , denoted by $\bigcap x$, to be the set that consists of the sets that belongs to all elements of x . Formally, for any set z we have

$$z \in \bigcap x \quad \Leftrightarrow \quad \forall w (w \in x \rightarrow z \in w).$$

Definition 1.8. For any sets x and y , we define

$$x \cup y = \bigcup \{x, y\}, \quad x \cap y = \bigcap \{x, y\}.$$

Chapter 2

Relations and Functions

2.1 Ordered Pairs

Definition 2.1. For sets x and y , we define

$$\langle x, y \rangle = \{\{x\}, \{x, y\}\}.$$

Lemma 2.2. Let x, y, y' be sets. If $\{x, y\} = \{x, y'\}$, then $y = y'$.

Proof. Suppose that $y \neq y'$. Since $y \in \{x, y\} = \{x, y'\}$ and $y \neq y'$, we have $y = x$. Then we have $y' \in \{x, y'\} = \{x, y\} = \{x\}$, implying $y' = x = y$, contradiction. Thus, $y = y'$. \square

Theorem 2.3. For sets x, x', y, y' , we have

$$\langle x, y \rangle = \langle x', y' \rangle$$

if and only if $x = x'$ and $y = y'$.

Proof. (\Leftarrow) Straightforward. (\Rightarrow) Suppose that $x \neq x'$. Since

$$\{\{x\}, \{x, y\}\} = \{\{x'\}, \{x', y'\}\},$$

either $\{x\} = \{x', y'\}$ or $\{x\} = \{x'\}$ holds. For both cases we all have $x' \in \{x\}$, implying $x' = x$, contradiction. Hence we have $x = x'$, and it follows that $\{x\} = \{x'\}$, implying $\{x, y\} = \{x', y'\}$, and thus $y = y'$. \square

Lemma 2.4. If $x, y \in C$, then $\langle x, y \rangle \in \mathcal{P}(\mathcal{P}(C))$.

Proof. Since $\{x\}$ and $\{y\}$ are subsets of C , we have $\{x\}, \{x, y\} \in \mathcal{P}(C)$. It follows that $\{\{x\}, \{x, y\}\}$ is a subset of $\mathcal{P}(C)$, implying

$$\langle x, y \rangle = \{\{x\}, \{x, y\}\} \in \mathcal{P}(\mathcal{P}(C)).$$

Theorem 2.5. For any sets A and B , there is a set whose members are exactly the pairs $\langle x, y \rangle$ with $x \in A$ and $y \in B$.

Proof. Since $x, y \in A \cup B$, the set of pairs $\langle x, y \rangle$ with $x \in A$ and $y \in B$ can be given by

$$\{z \in \mathcal{P}(\mathcal{P}(A \cup B)) : z = \langle x, y \rangle \text{ for some } x \in A \text{ and } y \in B\}.$$

Definition 2.6. For any sets A and B , the **Cartesian product** of A and B , denoted by $A \times B$, is the set whose members are exactly the pairs $\langle x, y \rangle$ with $x \in A$ and $y \in B$.

2.2 Relations

Definition 2.7. A **relation** is a set of ordered pairs. For any relation R , the **domain** and the **range** of R , denoted by $\text{dom}(R)$ and $\text{ran}(R)$, respectively, are defined as follows.

- $\text{dom}(R)$ is the collection of sets x with $\langle x, y \rangle \in R$ for some y .
- $\text{ran}(R)$ is the collection of sets y with $\langle x, y \rangle \in R$ for some x .

Definition 2.8. Let R and S be relations and let X be a set.

- The **inverse** of R , denoted by R^{-1} , is the set of all pairs $\langle y, x \rangle$ with $\langle x, y \rangle \in R$.
- The **restriction** of R to X , denoted by $R \upharpoonright X$, is the set of all pairs $\langle x, y \rangle \in R$ with $x \in X$.
- The **composition** of R and S , denoted by $R \circ S$, is the set of all pairs $\langle x, z \rangle$ with $\langle x, y \rangle \in R$ and $\langle y, z \rangle \in S$.

Definition 2.9. A **function** is a relation f such that for any set $x \in \text{dom}(f)$, there exists a unique set y such that $\langle x, y \rangle \in f$. The unique set y with respect to x is called the **value** of f at x and is denoted $f(x)$.

- We say that f is a function from A to B , denoted by $f : A \rightarrow B$, if $\text{dom}(f) = A$ and $\text{ran}(f) \subseteq B$.
- f is **one-to-one** if for any $y \in \text{ran}(f)$, there exists a unique set $x \in \text{dom}(f)$ with $f(x) = y$.

Definition 2.10. For any sets A and B , the set of functions from A to B is denoted by B^A .

2.3 Equivalence Relations and Ordering Relations

Definition 2.11. Let A be a set. An **equivalence relation** on A is a relation $R \subseteq A \times A$ that satisfies the following three conditions.

- Reflexivity: $\langle x, x \rangle \in R$ for any $x \in A$.
- Symmetry: $\langle x, y \rangle \in R$ implies $\langle y, x \rangle \in R$ for any $x, y \in A$.
- Transitivity: $\langle x, y \rangle \in R$ and $\langle y, z \rangle \in R$ implies $\langle x, z \rangle \in R$ for any $x, y, z \in A$.

Chapter 3

Natural Numbers

3.1 Inductive Sets

Definition 3.1. The **successor** of a set x , denoted x^+ , is defined by

$$x^+ = x \cup \{x\}.$$

We say that a set A is **inductive** if $\emptyset \in A$ and for any $x \in A$, we have $x^+ \in A$.

Axiom VI (Infinity). There exists an inductive set.

Definition 3.2. A **natural number** is a set belonging to all inductive sets. The set of natural numbers is denoted by ω .

Theorem 3.3. ω is inductive.

Proof. First, $\emptyset \in \omega$ since \emptyset belongs to all inductive sets by definition. For any set $x \in \omega$, x belongs to all inductive sets, implying that x^+ belongs to all inductive sets, and thus $x^+ \in \omega$. Thus, ω is inductive. \square

3.2 Recursion

3.3 Arithmetic

3.4 Ordering