

# An Introduction to Z3

Huixing Fang

National Trusted Embedded Software  
Engineering Technology Research Center

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1 SMT

2 Z3

# Outline

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2 Z3

# SMT: Satisfiability Modulo Theory

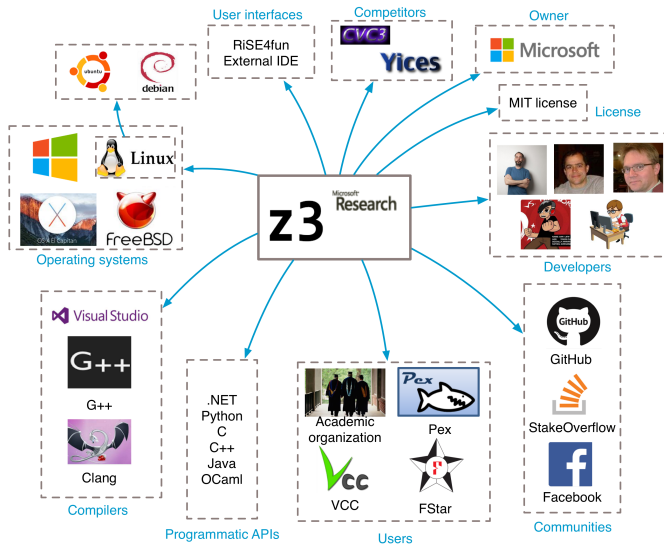
- ① Satisfiability is the problem of determining whether a formula  $\phi$  has a model
  - ① If  $\phi$  is propositional, a model is a truth assignment to Boolean variables
  - ② If  $\phi$  is a first-order formula, a model assigns values to variables and interpretations to the function and predicate symbols
- ② SAT Solvers: check satisfiability of propositional formulas
- ③ SMT Solvers: check satisfiability of formulas in a decidable first-order theory (e.g., linear arithmetic, array theory, bitvectors)

# Outline

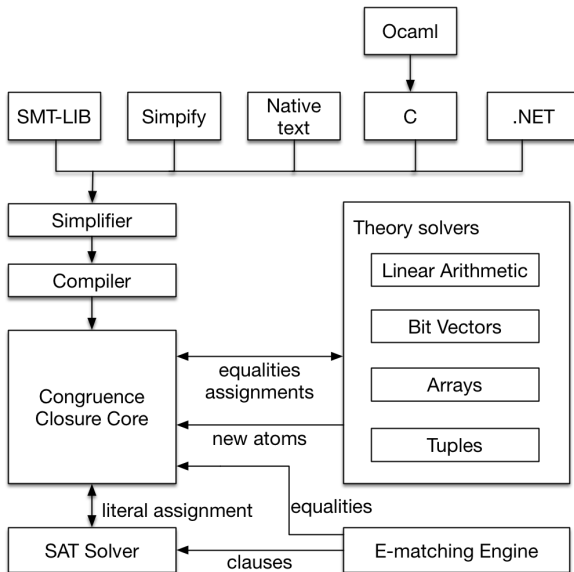
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2 Z3

# Z3 Context



# Z3 Architecture



- 1 Simplifier : Input formulas are first processed using an incomplete, but efficient simplification
  - $p \wedge \text{true} \mapsto p$
  - $x = 4 \wedge q(x) \mapsto x = 4 \wedge q(4)$
- 2 Compiler: The simplified abstract syntax tree representation of the formula is converted into a different data-structure comprising of a set of clauses and congruence-closure nodes
- 3 Congruence Closure Core: Handles equalities and uninterpreted functions. Receives the assignment from the SAT solver, and processed using E-matching



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- ❼ check: `(check-sat)` determines whether the current formulas on the Z3 stack are satisfiable or not
- ❽ model: `(get-model)` is used to retrieve an interpretation that makes all formulas on the Z3 internal stack true



## Basic Building Blocks

- ① The basic building blocks of SMT formulas are constants and functions. Constants are just functions that take no arguments. So everything is really just a function.
- ② An uninterpreted function or function symbol is one that has no other property than its name and n-ary form.

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## Valid and Satisfiable

- 1 A formula  $F$  is valid if  $F$  always evaluates to true for any assignment of appropriate values to its uninterpreted function and constant symbols.
- 2 A formula  $F$  is satisfiable if there is some assignment of appropriate values to its uninterpreted function and constant symbols under which  $F$  evaluates to true.

# Propositional Logic

The pre-defined sort `Bool` is the sort (type) of all Boolean propositional expressions. Z3 supports the usual Boolean operators `and`, `or`, `xor`, `not`, `=>` (implication), `ite` (if-then-else).

## Example

Formulas:

```
(declare-const p Bool)
(declare-const q Bool)
(declare-const r Bool)
(define-fun conjecture () Bool
  (=> (and (=> p q) (=> q r))
    (=> p r)))
(assert (not conjecture))
(check-sat)
```

Result: unsat

# Uninterpreted functions and constants

The basic building blocks of SMT formulas are constants and functions. Constants are just functions that take no arguments. So everything is really just a function.

## Example

Formulas:

```
(declare-fun f (Int) Int)
(declare-fun a () Int)
; a is a constant
(declare-const b Int)
; syntax sugar for (declare-fun b () Int)
(assert (> a 20))
(assert (> b a))
(assert (= (f 10) 1))
(check-sat)
(get-model)
```

Z3 has builtin support for integer and real constants. These two types (sorts) represent the mathematical integers and reals.

```
(declare-const a Int)
(declare-const b Int)
(declare-const d Real)
(declare-const e Real)
(assert (> a (+ b 2)))
(assert (>= d e))
(check-sat)
(get-model)
```

```
sat
(model
  (define-fun e () Real 0.0)
  (define-fun d () Real 0.0)
  (define-fun a () Int 1)
  (define-fun b () Int (- 2))
)
```

# Nonlinear arithmetic

- 1 A formula is nonlinear if it contains expressions of the form  $( * t s )$  where  $t$  and  $s$  are not numbers.
- 2 Nonlinear real arithmetic is very expensive, and Z3 is not complete for this kind of formula.

```
(declare-const a Int)
(assert (> (* a a) 3))
(check-sat)
(get-model)

(declare-const b Real)
(declare-const c Real)
(assert (= (+ (* b b b)
(* b c)) 3.0))
(check-sat)
```

```
sat
(model
  (define-fun a () Int
    (- 8)
  )
)

unknown
```

## Bitvector Literals

Bitvector literals may be defined using binary, decimal and hexadecimal notation.

- 1 #b010 in binary format is a bitvector of size 3
- 2 bitvector literal #x0a0 in hexadecimal format is a bitvector of size 12

The size must be specified for bitvector literals in decimal format.

- (`_ bv10 32`) is a bitvector of size 32 that represents the numeral 10

Z3 supports Bitvectors of arbitrary size. (`_ BitVec n`) is the sort of bitvectors whose length is `n`. Declare 64-bit bitvector constant `x` as follow:

```
(declare-const x (_ BitVec 64))
```

# Basic Bitvector Arithmetic

- 1 `bvadd` : addition
- 2 `bvsub` : subtraction
- 3 `bvneg` : unary minus
- 4 `bvmul` : multiplication
- 5 `bvurem` : unsigned remainder
- 6 `bvsrem` : signed remainder
- 7 `bvsmod` : signed modulo
- 8 `bvshl` : shift left
- 9 `bvlshr` : unsigned (logical) shift right
- 10 `bvashr` : signed (arithmetical) shift right
- 11 `bvor` : bitwise or
- 12 `bvand` : bitwise and
- 13 `bvnot` : bitwise not
- 14 `bvnand` : bitwise nand
- 15 `bvnor` : bitwise nor
- 16 `bvxnor` : bitwise xnor



# Arrays

- 1 The expression `(select a i)` returns the value stored at position `i` of the array `a`;
- 2 and `(store a i v)` returns a new array identical to `a`, but on position `i` it contains the value `v`.

```
(declare-const x Int)
(declare-const y Int)
(declare-const a1 (Array Int Int))
(assert (= (select a1 x) x))
(assert (= (store a1 x y) a1))
(check-sat)
(get-model)
```

```
sat
(model
  (define-fun y () Int 1)
  (define-fun a1 () (Array Int Int) (_ as-array k!0))
  (define-fun x () Int 1)
  (define-fun k!0 ((x!1 Int)) Int (ite (= x!1 1) 1 0)) )
```

# Records

A record is specified as a datatype with a single constructor and as many arguments as record elements.

## Example

Parametric type `Pair`, with constructor `mk-pair` and two arguments that can be accessed using the selector functions `first` and `second`.

```
(declare-datatypes (T1 T2)
  ((Pair (mk-pair (first T1) (second T2)))))
(declare-const p1 (Pair Int Int))
(declare-const p2 (Pair Int Int))
(assert (= p1 p2))
(assert (> (second p1) 20))
(check-sat)
(get-model)
(assert (not (= (first p1) (first p2))))
(check-sat)
```

# Model-based Quantifier Instantiation

The model-based quantifier instantiation (MBQI) is essentially a counter-example based refinement loop, where candidate models are built and checked.

```
(set-option :smt.mbqi true)
(declare-fun f (Int Int) Int)
(declare-const a Int)
(declare-const b Int)
(assert (forall ((x Int))
  (>= (f x x) (+ x a))))
(assert (< (f a b) a))
(assert (> a 0))
(check-sat)
(get-model)

(eval (f (+ a 10) 20))
```

```
sat
(model
  (define-fun b () Int 2)
  (define-fun a () Int 1)
  (define-fun f
    ((x!1 Int)
     (x!2 Int))
    Int
    (ite (and (= x!1 1)
              (= x!2 2))
          0
          (+ 1 x!1)))
)

12
```

# Relations, rules and queries

The default fixed-point engine is a bottom-up Datalog engine. It works with finite relations and uses finite table representations as hash tables as the default way to represent finite relations.

```
(declare-rel a ())  
(declare-rel b ())  
(declare-rel c ())  
(rule (=> b a))  
(rule (=> c b))  
(set-option :fixedpoint.engine datalog)  
(query a)  
(rule c)  
(query a:print-answer true)
```

```
unsat  
sat  
true
```

# Engine for Property Directed Reachability

The PDR engine is targeted at applications from symbolic model checking of software. The systems may be infinite state.

## Example

McCarthy's 91 function illustrates a procedure that calls itself recursively twice

$$mc(x) = \text{if } x > 100 \text{ then } x - 10 \text{ else } mc(mc(x+11))$$

```
(declare-rel mc (Int Int))
(declare-var n Int)
(declare-var m Int)
(declare-var p Int)

(rule (=> (> m 100) (mc m (- m 10))))
(rule (=> (and (<= m 100) (mc (+ m 11) p) (mc p n))
        (mc m n)))
```

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```
(declare-rel mc (Int Int))
(declare-var n Int)
(declare-var m Int)
(declare-var p Int)
(rule (=> (> m 100) (mc m (- m 10))))
(rule (=> (and (<= m 100) (mc (+ m 11) p) (mc p n))
        (mc m n)))

(declare-rel q (Int Int))
(rule (=> (and (mc m n) (< n 92)) (q m n)))
(query q :print-certificate true)
```

# Engine for Property Directed Reachability

## Example

McCarthy's 91 function illustrates a procedure that calls itself recursively twice

$$mc(x) = \text{if } x > 100 \text{ then } x - 10 \text{ else } mc(mc(x+11))$$

```
sat
(and
  (not (<= 92 (:var 5)))
  (mc (:var 4) (:var 5))
  (= (:var 5) (+ (- 10)
    (:var 4)))
  (not (<= (:var 4) 100))
)
```

$$v5 < 92$$
$$mc(v4) = v5$$
$$v5 = v4 - 10$$
$$v4 > 100$$

# End

Thanks for your attention.