## 抽象解释 及其在静态分析中的应用

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### 目录

- 一、抽象解释概述
- 二、抽象解释理论的数学基础
- 三、具体语义下的静态分析
- 四、抽象语义下的静态分析
- 五、基于抽象解释的静态分析工具

# 简单数值程序的语法

#### **Arithmetic expressions:**

```
\begin{array}{lll} \mathtt{exp} & ::= & \mathtt{V} & \mathtt{variable} \ \mathtt{V} \in \mathbb{V} \\ & | & -\mathtt{exp} & \mathtt{negation} \\ & | & \mathtt{exp} \diamond \mathtt{exp} & \mathtt{binary} \ \mathtt{operation:} \ \diamond \in \{+,-,\times,/\} \\ & | & [c,c'] & \mathtt{constant} \ \mathtt{range}, \ c,c' \in \mathbb{I} \cup \{\pm\infty\} \\ & & (c \ \mathsf{is} \ \mathsf{a} \ \mathsf{shorthand} \ \mathsf{for} \ [c,c]) \end{array}
```

#### **Commands:**

$$\begin{array}{lll} \mathsf{com} & ::= & \mathsf{V} & := & \mathsf{exp} & & \mathsf{assignment into} \; \mathsf{V} \in \mathbb{V} \\ & | & \mathsf{exp} \bowtie \mathsf{0} & & \mathsf{test}, \; \bowtie \, \in \{\, =, <, >, \le, \ge, \ne \, \} \end{array}$$

programs: control-flow graphs

$$P \stackrel{\text{def}}{=} (L, e, A)$$

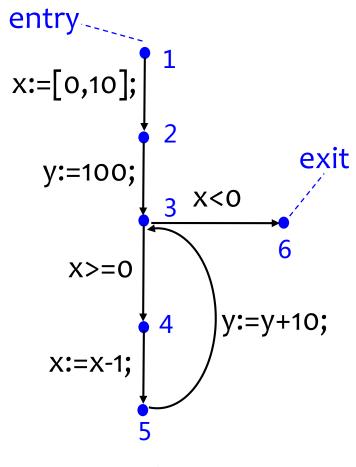
$$P = (L, e, A)$$

$$P$$

# 程序示例

```
1 x:=[0,10];
2 y:=100;
3 while (x>=0) do
4    x:=x-1;
5    y:=y+10;
    done 6
```

结构化的程序



控制流图

• 表达式的具体语义

```
\stackrel{\text{def}}{=} \quad \{ x \in \mathbb{I} \mid c \le x \le c' \}
\mathbb{E}[[c,c']]\rho
                                                \stackrel{\text{def}}{=}
\mathbb{E}[V]\rho
                                                                 \{ \rho(V) \}
                                             \stackrel{\mathrm{def}}{=}
\mathbb{E}\llbracket -e 
bracket 
ho
                                                               \{-v \mid v \in \mathsf{E} \llbracket e \rrbracket \rho \}
                                           \stackrel{\mathrm{def}}{=}
\mathbb{E}[\![e_1 + e_2]\!] \rho
                                                                 \{ v_1 + v_2 \mid v_1 \in \mathsf{E} \llbracket e_1 \rrbracket \rho, v_2 \in \mathsf{E} \llbracket e_2 \rrbracket \rho \}
                                           \stackrel{\mathrm{def}}{=}
\mathbb{E}[\![e_1 - e_2]\!] \rho
                                                                 \{ v_1 - v_2 \mid v_1 \in \mathsf{E} \llbracket e_1 \rrbracket \rho, v_2 \in \mathsf{E} \llbracket e_2 \rrbracket \rho \}
                                           \stackrel{\mathrm{def}}{=}
\mathbb{E}[\![e_1 \times e_2]\!] \rho
                                                                 \{ v_1 \times v_2 \mid v_1 \in \mathsf{E} \llbracket e_1 \rrbracket \rho, v_2 \in \mathsf{E} \llbracket e_2 \rrbracket \rho \}
                                                \stackrel{\text{def}}{=}
\mathbb{E}[\![e_1/e_2]\!]\rho
                                                                  \{ v_1/v_2 \mid v_1 \in \mathsf{E} \llbracket e_1 \rrbracket \rho, v_2 \in \mathsf{E} \llbracket e_2 \rrbracket \rho, v_2 \neq \mathbf{0} \}
```

• 语句的具体语义

$$\mathsf{C}\llbracket c \rrbracket : \mathcal{D} \to \mathcal{D} \text{ where } \mathcal{D} \stackrel{\mathrm{def}}{=} \mathcal{P}(\mathbb{V} \to \mathbb{I})$$

语句c的迁移函数 语句c执行前的状态 语句c执行后的状态

```
C[V := e] \mathcal{X} \stackrel{\text{def}}{=} \{ \rho[V \mapsto V] | \rho \in \mathcal{X}, V \in E[e] \rho \} \\
C[e \bowtie 0] \mathcal{X} \stackrel{\text{def}}{=} \{ \rho | \rho \in \mathcal{X}, \exists V \in E[e] \rho : V \bowtie 0 \}
```

条件测试语句:过滤一些环境

赋值语句:更新变量的值

• 程序的具体语义(聚集语义)

$$P[(L, e, A)]: L \to D \text{ where } D \stackrel{\text{def}}{=} \mathcal{P}(V \to I)$$
程序点 状态的集合

P[(L, e, A)]ℓ 是程序点  $l \in L$  处最精确的不变式 也是如下递归方程系统的最小解

给定的初始状态

• 程序的具体语义(聚集语义)

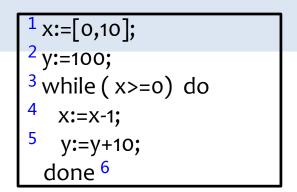
$$\mathsf{P}\llbracket (\mathsf{L}, e, \mathsf{A}) 
rbracketeta : \mathsf{L} o \mathcal{D} \text{ where } \mathcal{D} \stackrel{\mathrm{def}}{=} \mathcal{P}(\mathbb{V} o \mathbb{I})$$

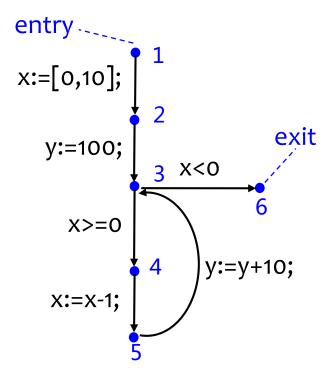
给定的初始状态

- $\triangleright (\mathcal{D}, \subseteq, \cup, \cap, \emptyset, (\mathbb{V} \to \mathbb{I}))$  是完全格
- ho  $M_{\ell}: \mathcal{X}_{\ell} \mapsto \bigcup_{k} \mathbb{C}[\![c]\!] \mathcal{X}_{\ell'}$  是D上的单调函数  $(\ell',c,\ell)\in A$

根据Tarski定理,函数  $\mathsf{M}_\ell$  的最小不动点存在且唯一  $orall \ell$  :  $M_\ell(\mathcal{X}_\ell) = \mathcal{X}_\ell$ 

## 具体语义—示例





$$egin{aligned} \mathcal{X}_1 &= (\{\,\mathtt{X},\mathtt{Y}\,\} 
ightarrow \mathbb{Z}) \ \mathcal{X}_2 &= \mathsf{C}[\![\,\mathtt{X}\, := [0,10]\,]\!]\,\mathcal{X}_1 \ \mathcal{X}_3 &= \mathsf{C}[\![\,\mathtt{Y}\, := 100\,]\!]\,\mathcal{X}_2 \cup \ \mathsf{C}[\![\,\mathtt{Y}\, := \mathtt{Y} + 10\,]\!]\,\mathcal{X}_5 \ \mathcal{X}_4 &= \mathsf{C}[\![\,\mathtt{X}\, \geq 0\,]\!]\,\mathcal{X}_3 \ \mathcal{X}_5 &= \mathsf{C}[\![\,\mathtt{X}\, := \mathtt{X} - 1\,]\!]\,\mathcal{X}_4 \ \mathcal{X}_6 &= \mathsf{C}[\![\,\mathtt{X}\, < 0\,]\!]\,\mathcal{X}_3 \end{aligned}$$

语义方程系统

循环不变式

控制流图

 $\mathcal{X}_3 = \{ \rho \mid \rho(X) \in [-1,10], \ 10\rho(X) + \rho(Y) \in [100,200] \cap 10\mathbb{Z} \}$ 

# 具体语义下求解最小不动点

$$egin{array}{cccc} \mathcal{X}_e \ \mathcal{X}_{\ell 
eq e} &=& igcup_{(\ell',c,\ell) \in A} lacksquare A \ && \ \end{array}$$

• 计算最小不动点: Kleene迭代

$$\left\{ \begin{array}{ll} \mathcal{X}_{e}^{0} & \stackrel{\mathrm{def}}{=} & \mathcal{X}_{e} \\ \mathcal{X}_{\ell \neq e}^{0} & \stackrel{\mathrm{def}}{=} & \emptyset \end{array} \right. \quad \left\{ \begin{array}{ll} \mathcal{X}_{e}^{n+1} & \stackrel{\mathrm{def}}{=} & \mathcal{X}_{e} \\ \mathcal{X}_{\ell \neq e}^{n+1} & \stackrel{\mathrm{def}}{=} & \bigcup_{(\ell',c,\ell) \in \mathcal{A}} \mathbb{C}[\![\,c\,]\!] \mathcal{X}_{\ell'}^{n} \end{array} \right.$$

### 第 0 次迭代

$$\mathcal{X}_1 = \mathbb{Z}^2$$

$$Z^2$$

$$\mathcal{X}_2 = \mathsf{C}[\![\, \mathtt{X} := [0,10] \,]\!]\, \mathcal{X}_1$$

$$\boldsymbol{\Phi}$$

$$\mathcal{X}_3 = \mathsf{C} \llbracket \, \mathsf{Y} := \mathsf{100} \, \rrbracket \, \mathcal{X}_2 \cup \qquad \mathcal{C} \llbracket \, \mathsf{Y} := \mathsf{Y} + \mathsf{10} \, \rrbracket \, \mathcal{X}_5$$

$$\mathcal{X}_4 = C \llbracket \, X \geq 0 \, \rrbracket \, \mathcal{X}_3$$

$$\mathcal{X}_5 = \mathsf{C}[\![\, \mathsf{X} := \mathsf{X} - 1\,]\!]\,\mathcal{X}_4 \qquad \boldsymbol{\phi}$$

$$\mathcal{X}_6 = \mathsf{C} \llbracket \, \mathtt{X} < \mathsf{0} \, \rrbracket \, \mathcal{X}_3$$

### 1 x:=[0,10]; 2 y:=100; 3 while (x>=0) do 4 x:=x-1; 5 y:=y+10; done 6

第1次迭代

$$\mathcal{X}_1 = \mathbb{Z}^2$$

$$Z^2$$

$$\mathcal{X}_2 = \mathsf{C} \llbracket \, \mathtt{X} := [0, 10] \, \rrbracket \, \mathcal{X}_1$$

$$[0,10] \times Z$$

$$\mathcal{X}_3 = \mathsf{C} \llbracket \, \mathsf{Y} := \mathsf{100} \, \rrbracket \, \mathcal{X}_2 \cup \mathsf{C} \llbracket \, \mathsf{Y} := \mathsf{Y} + \mathsf{10} \, \rrbracket \, \mathcal{X}_5$$

$$\mathcal{X}_4 = C \llbracket \, X \geq 0 \, \rrbracket \, \mathcal{X}_3$$

$$\mathcal{X}_5 = \mathsf{C} \llbracket \, \mathtt{X} := \mathtt{X} - 1 \, \rrbracket \, \mathcal{X}_4$$

$$\Phi$$

$$\mathcal{X}_6 = \mathsf{C} \llbracket \, \mathtt{X} < \mathsf{0} \, \rrbracket \, \mathcal{X}_3$$

第2次迭代

$$\mathcal{X}_1 = \mathbb{Z}^2$$

$$Z^2$$

$$\mathcal{X}_2 = \mathsf{C}[\![\, \mathtt{X} := [0,10] \,]\!]\,\mathcal{X}_1$$

$$[0,10] \times Z$$

$$\mathcal{X}_3 = \mathsf{C} \llbracket \, \mathsf{Y} := \mathsf{100} \, \rrbracket \, \mathcal{X}_2 \cup \mathsf{C} \llbracket \, \mathsf{Y} := \mathsf{Y} + \mathsf{10} \, \rrbracket \, \mathcal{X}_5$$

$$\mathcal{X}_4 = C \llbracket \, X \geq 0 \, \rrbracket \, \mathcal{X}_3$$

$$\mathcal{X}_5 = \mathsf{C} \llbracket \, \mathtt{X} := \mathtt{X} - 1 \, \rrbracket \, \mathcal{X}_4$$

$$\Phi$$

$$\mathcal{X}_6 = \mathsf{C} \llbracket \, \mathtt{X} < \mathsf{0} \, \rrbracket \, \mathcal{X}_3$$

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第3次迭代

$$\mathcal{X}_1 = \mathbb{Z}^2$$

$$Z^2$$

$$\mathcal{X}_2 = \texttt{C}[\![\, \texttt{X} := [0,10] \,]\!]\, \mathcal{X}_1$$

$$[0,10] \times Z$$

$$\mathcal{X}_3 = \mathsf{C} \llbracket \, \mathsf{Y} := \mathsf{100} \, \rrbracket \, \mathcal{X}_2 \cup \mathsf{C} \llbracket \, \mathsf{Y} := \mathsf{Y} + \mathsf{10} \, \rrbracket \, \mathcal{X}_5$$

$$\{(0,100),...,(10,100)\}$$

$$\mathcal{X}_4 = C \llbracket \, X \geq 0 \, \rrbracket \, \mathcal{X}_3$$

$$\mathcal{X}_5 = \mathsf{C} \llbracket \, \mathtt{X} := \mathtt{X} - 1 \, \rrbracket \, \mathcal{X}_4$$

$$\mathcal{X}_6 = \mathsf{C} \llbracket \, \mathtt{X} < \mathsf{0} \, \rrbracket \, \mathcal{X}_3$$

第 4 次迭代

$$\mathcal{X}_1 = \mathbb{Z}^2$$

$$\mathcal{X}_2 = \mathsf{C} \llbracket \, \mathtt{X} := [0,10] \, \rrbracket \, \mathcal{X}_1$$

$$\mathcal{X}_3 = \mathsf{C} \llbracket \, \mathtt{Y} := \mathtt{100} \, \rrbracket \, \mathcal{X}_2 \cup \mathsf{C} \llbracket \, \mathtt{Y} := \mathtt{Y} + \mathtt{10} \, \rrbracket \, \mathcal{X}_5$$

$$\mathcal{X}_4 = C \llbracket \, X \geq 0 \, \rrbracket \, \mathcal{X}_3$$

$$\mathcal{X}_5 = \mathsf{C} \llbracket \, \mathtt{X} := \mathtt{X} - 1 \, \rrbracket \, \mathcal{X}_4$$

$$\mathcal{X}_6 = C \llbracket \, \mathtt{X} < 0 \, \rrbracket \, \mathcal{X}_3$$

$$Z^2$$

$$[0,10] \times Z$$

```
{ (0,100),...,(10,100) }
```

```
{ (0,100),...,(10,100) }
```

```
{ (-1,100),...,(9,100) }
```

第 5 次迭代

 $\mathcal{X}_6 = C \mathbb{I} X < 0 \mathbb{I} \mathcal{X}_3$ 

```
1 x:=[0,10];
2 y:=100;
3 while ( x>=0) do
4     x:=x-1;
5     y:=y+10;
    done 6
```

$$\left\{ \begin{array}{ll} \mathcal{X}_{1} = \mathbb{Z}^{2} & \mathcal{Z}^{2} \\ \mathcal{X}_{2} = \mathsf{C} \llbracket \, \mathsf{X} := [0, 10] \, \rrbracket \, \mathcal{X}_{1} & [0, 10] \, \mathsf{x} \, \mathcal{Z} \\ \mathcal{X}_{3} = \mathsf{C} \llbracket \, \mathsf{Y} := 100 \, \rrbracket \, \mathcal{X}_{2} \cup & \{ \, (0, 100), ..., (10, 100), \\ \mathsf{C} \llbracket \, \mathsf{Y} := \, \mathsf{Y} + 10 \, \rrbracket \, \mathcal{X}_{5} & (-1, 110), ..., (9, 110) \, \} \\ \mathcal{X}_{4} = \mathsf{C} \llbracket \, \mathsf{X} \geq 0 \, \rrbracket \, \mathcal{X}_{3} & \{ \, (0, 100), ..., (10, 100), \\ (0, 110), ..., (9, 110) \, \} \\ \mathcal{X}_{5} = \mathsf{C} \llbracket \, \mathsf{X} := \, \mathsf{X} - 1 \, \rrbracket \, \mathcal{X}_{4} & \{ \, (-1, 100), ..., (9, 100) \, \} \\ \mathcal{X}_{6} = \mathsf{C} \llbracket \, \mathsf{X} < 0 \, \rrbracket \, \mathcal{X}_{3} & \{ \, (-1, 110) \, \} \end{array} \right.$$

```
1 x:=[0,10];
2 y:=100;
3 while ( x>=0) do
4    x:=x-1;
5    y:=y+10;
    done 6
```

#### 第7次迭代

$$\left\{ \begin{array}{ll} \mathcal{X}_{1} = \mathbb{Z}^{2} & \mathcal{Z}^{2} \\ \mathcal{X}_{2} = \mathsf{C} \llbracket \, \mathsf{X} := [0, 10] \, \rrbracket \, \mathcal{X}_{1} & [0, 10] \, \mathsf{X} \, \mathcal{Z} \\ \mathcal{X}_{3} = \mathsf{C} \llbracket \, \mathsf{Y} := 100 \, \rrbracket \, \mathcal{X}_{2} \cup & \{ (0, 100), ..., (10, 100), \\ \mathsf{C} \llbracket \, \mathsf{Y} := \, \mathsf{Y} + 10 \, \rrbracket \, \mathcal{X}_{5} & (-1, 110), ..., (9, 110) \} \\ \mathcal{X}_{4} = \mathsf{C} \llbracket \, \mathsf{X} \geq 0 \, \rrbracket \, \mathcal{X}_{3} & \{ (0, 100), ..., (10, 100), \\ (0, 110), ..., (9, 110) \} \\ \mathcal{X}_{5} = \mathsf{C} \llbracket \, \mathsf{X} := \, \mathsf{X} - 1 \, \rrbracket \, \mathcal{X}_{4} & \{ (-1, 100), ..., (9, 100), \\ (-1, 110), ..., (8, 110) \} \\ \mathcal{X}_{6} = \mathsf{C} \llbracket \, \mathsf{X} < 0 \, \rrbracket \, \mathcal{X}_{3} & \{ (-1, 110) \} \end{array} \right.$$

```
1 x:=[0,10];
2 y:=100;
3 while ( x>=0) do
4    x:=x-1;
5    y:=y+10;
    done 6
```

第8次迭代

$$\left\{ \begin{array}{ll} \mathcal{X}_{1} = \mathbb{Z}^{2} & \mathcal{Z}^{2} \\ \mathcal{X}_{2} = \mathsf{C} \llbracket \, \mathsf{X} := [0, 10] \, \rrbracket \, \mathcal{X}_{1} & [0, 10] \, \mathsf{X} \, \mathcal{Z} \\ \mathcal{X}_{3} = \mathsf{C} \llbracket \, \mathsf{Y} := 100 \, \rrbracket \, \mathcal{X}_{2} \cup & \{ (0, 100), ..., (10, 100), \\ \mathsf{C} \llbracket \, \mathsf{Y} := \, \mathsf{Y} + 10 \, \rrbracket \, \mathcal{X}_{5} & (-1, 110), ..., (9, 110), \\ \mathcal{X}_{4} = \mathsf{C} \llbracket \, \mathsf{X} \geq 0 \, \rrbracket \, \mathcal{X}_{3} & \{ (0, 100), ..., (10, 100), \\ (0, 110), ..., (9, 110) \, \} \\ \mathcal{X}_{5} = \mathsf{C} \llbracket \, \mathsf{X} := \, \mathsf{X} - 1 \, \rrbracket \, \mathcal{X}_{4} & \{ (-1, 100), ..., (9, 100), \\ (-1, 110), ..., (8, 110) \, \} \\ \mathcal{X}_{6} = \mathsf{C} \llbracket \, \mathsf{X} < 0 \, \rrbracket \, \mathcal{X}_{3} & \{ (-1, 110) \, \} \end{array} \right.$$

```
1 x:=[0,10];
2 y:=100;
3 while ( x>=0) do
4    x:=x-1;
5    y:=y+10;
    done 6
```

### יני ההענשיוייני

第9次迭代

```
Z^2
\mathcal{X}_1 = \mathbb{Z}^2
                                                                      [0,10] \times Z
\mathcal{X}_2 = \mathsf{C}[\![ \mathtt{X} := [0, 10] \!]\!] \mathcal{X}_1
                                                                     \{(0,100),...,(10,100),
\mathcal{X}_3 = \mathsf{C} \llbracket \, \mathsf{Y} := \mathsf{100} \, \rrbracket \, \mathcal{X}_2 \, \cup \,
              \mathsf{C} \llbracket \, \mathsf{Y} := \mathsf{Y} + \mathsf{10} \, \rrbracket \, \mathcal{X}_{\mathsf{5}}
                                                                      (-1,110),...,(9,110),
                                                                        (-1,120),...,(8,120)
                                                                     { (0,100),...,(10,100),
\mathcal{X}_4 = \mathsf{C}[\![ \mathsf{X} \geq \mathsf{0} ]\!] \mathcal{X}_3
                                                                        (0,110),...,(9,110),
                                                                        (0,120),...,(8,120)
\mathcal{X}_5 = \mathsf{C} \llbracket \, \mathsf{X} := \mathsf{X} - \mathsf{1} \, \rrbracket \, \mathcal{X}_4
                                                                     { (-1,100),...,(9,100),
                                                                        (-1,110),...,(8,110)
\mathcal{X}_6 = C \llbracket X < 0 \rrbracket \mathcal{X}_3
                                                                      { (-1,110), (-1,120) }
```

```
1 x:=[0,10];
2 y:=100;
3 while (x>=0) do
4 x:=x-1;
5 y:=y+10;
done 6
```

第 10 次迭代

```
1 x:=[0,10];
2 y:=100;
3 while (x>=0) do
4  x:=x-1;
5  y:=y+10;
  done 6
```

第 ... 次迭代

```
Z^2
\mathcal{X}_1 = \mathbb{Z}^2
                                                            [0,10] \times Z
\mathcal{X}_2 = \mathsf{C}[\![ \mathtt{X} := [0, 10] \!]\!] \mathcal{X}_1
                                                            \{(0,100),...,(10,100),
\mathcal{X}_3 = \mathsf{C} \llbracket \, \mathsf{Y} := \mathsf{100} \, \rrbracket \, \mathcal{X}_2 \, \cup \,
            C[Y := Y + 10]\mathcal{X}_5
                                                           (-1,110),...,(9,110),
                                                              (-1,120),...,(8,120),...
                                                           \{(0,100),...,(10,100),
\mathcal{X}_4 = \mathsf{C}[\![ \mathsf{X} \geq \mathsf{0} ]\!] \mathcal{X}_3
                                                              (0,110),...,(9,110),
                                                              (0,120),...,(8,120),...
\mathcal{X}_5 = \mathsf{C} \llbracket \mathsf{X} := \mathsf{X} - 1 \rrbracket \mathcal{X}_4
                                                           { (-1,100),...,(9,100),
                                                              (-1,110),...,(8,110),
                                                              (-1,120),...,(7,120),...
\mathcal{X}_6 = C \mathbb{I} X < 0 \mathbb{I} \mathcal{X}_3
                                                            { (-1,110),...,(-1,120),... }
```

```
^{1} x:=[0,10];
<sup>2</sup> y:=100;
^{3} while (x>=0) do
    X:=X-1;
    y:=y+10;
  done <sup>6</sup>
```

# 本讲内容介绍

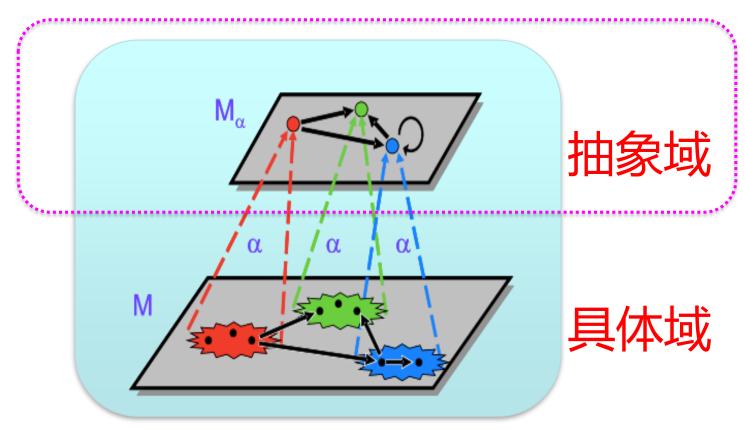
- 一、抽象解释理论的概述
- 二、抽象解释理论的数学基础
- 三、具体语义下的静态分析
- 四、抽象语义下的静态分析
- 五、基于抽象解释的静态分析工具

# 本讲内容介绍

- 一、抽象解释理论的概述
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  - ▶抽象域
  - >基于抽象域的静态分析
- 五、基于抽象解释的静态分析工具

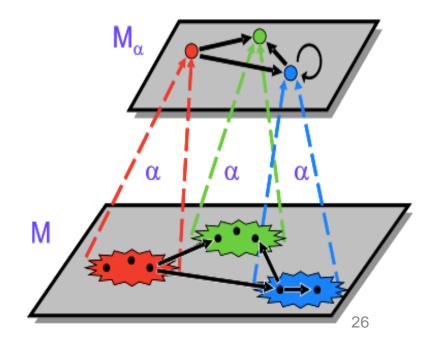
# 抽象域:抽象解释的核心要素

• 静态分析相关的计算都在抽象域上开展



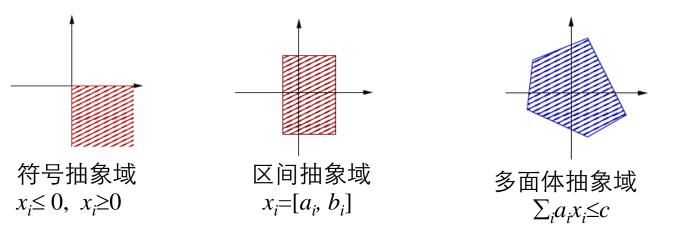
# 抽象域的构成

- 域元素: 对程序状态进行抽象
  - > 表示方法: 约束形式,...
  - ▶ E.g. 区间: *a*<=*x*<= *b*
- 域操作:对程序语义动作进行抽象
  - > 交 (assume语句)
  - > 控制流接合 (if-then-else-endif)
  - 投影(非确定赋值,过程间分析)
  - > 迁移函数
    - 赋值迁移语句 (赋值语句)
    - 测试迁移语句 (if 语句)
  - ▶ 加宽(循环)
  - **>** ...



# 数值抽象域

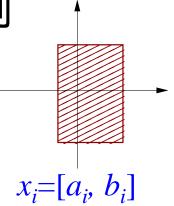
- 刻画程序变量之间的数值关系
- 用途:发现程序中某程序点处的数值不变式,即每次程序执行均满足的数值关系
  - 》除零错、数组越界、整数溢出等运行时错误
  - > 安全方面"缓冲区溢出"问题: 地址(指针)和长度(范围)之间的数值关系



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# 数值抽象域

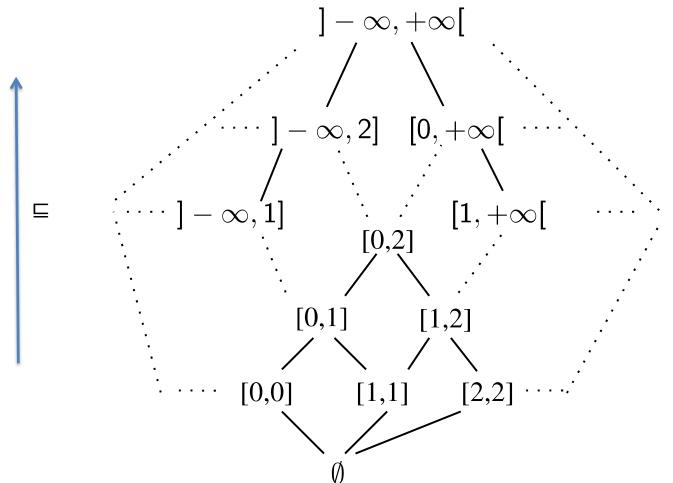
- 以两个数值抽象域为例
  - > 区间抽象域



> 线性等式抽象域

$$\sum_{i} a_{i} x_{i} = c$$

● 区间格 B# = {[a,b]|a∈R∪{-∞}, b∈R∪{+∞}, a≤b} ∪ {⊥#}



Galois连接

$$\wp(R) \xrightarrow{\gamma_b} B^{\#}$$

$$\gamma_{b}([a,b]) \triangleq \{x \in \mathbb{R} \mid a \leq x \leq b\}$$

$$\alpha_{b}(X) \triangleq \begin{cases} \perp^{\#} & \text{if } X = \emptyset \\ [\min X, \max X] & \text{otherwise} \end{cases}$$

### • 格相关操作

$$[a,b] \subseteq^{\sharp} [c,d] \qquad \stackrel{\text{def}}{\Longrightarrow} \qquad a \ge c \text{ and } b \le d$$

$$\top^{\sharp} \qquad \stackrel{\text{def}}{\equiv} \qquad ]-\infty,+\infty[$$

$$[a,b] \cup^{\sharp} [c,d] \qquad \stackrel{\text{def}}{\equiv} \qquad [\min(a,c),\max(b,d)]$$

$$[a,b] \cap^{\sharp} [c,d] \qquad \stackrel{\text{def}}{\equiv} \qquad \left\{ \begin{array}{l} [\max(a,c),\min(b,d)] & \text{if } \max \le \min \\ \bot^{\sharp} & \text{otherwise} \end{array} \right.$$

℘(R)对应的(B♯,⊆♯,∪♯,∩♯,⊥♯,T♯) 是一个完全格

### • 区间算术操作

$$[c, c']^{\sharp} \stackrel{\text{def}}{=} [c, c']$$

$$-^{\sharp} [a, b] \stackrel{\text{def}}{=} [-b, -a]$$

$$[a, b] +^{\sharp} [c, d] \stackrel{\text{def}}{=} [a + c, b + d]$$

$$[a, b] -^{\sharp} [c, d] \stackrel{\text{def}}{=} [a - d, b - c]$$

$$[a, b] \times^{\sharp} [c, d] \stackrel{\text{def}}{=} [\min(ac, ad, bc, bd), \max(ac, ad, bc, bd)]$$

$$[\min(ac, ad, bc, bd), \max(ac, ad, bc, bd)]$$

$$[\min(a/c, a/d, b/c, b/d), \text{ if } 0 \le c$$

$$\max(a/c, a/d, b/c, b/d)$$

$$[-b, -a]/^{\sharp} [-d, -c] \text{ if } d \le 0$$

$$([a, b]/^{\sharp} [c, 0]) \cup^{\sharp} ([a, b]/^{\sharp} [0, d]) \text{ otherwise}$$

where 
$$\begin{vmatrix} \pm \infty \times 0 = 0, & 0/0 = 0, & \forall x : x/\pm \infty = 0 \\ \forall x > 0 : x/0 = +\infty, & \forall x < 0 : x/0 = -\infty \end{vmatrix}$$

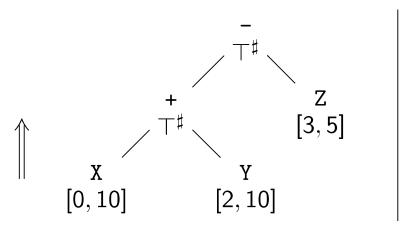
• 赋值操作

$$C^{\sharp} \llbracket \mathbf{V} := e \rrbracket \mathcal{X}^{\sharp} \stackrel{\mathrm{def}}{=} \left\{ \begin{array}{l} \bot^{\sharp} & \text{if } \mathcal{V}^{\sharp} = \bot^{\sharp} \\ \mathcal{X}^{\sharp} [\mathbf{V} \mapsto \mathcal{V}^{\sharp}] & \text{otherwise} \end{array} \right.$$
 where  $\mathcal{V}^{\sharp} = \mathsf{E}^{\sharp} \llbracket e \rrbracket \mathcal{X}^{\sharp}.$ 

> 其中表达式e的值的计算

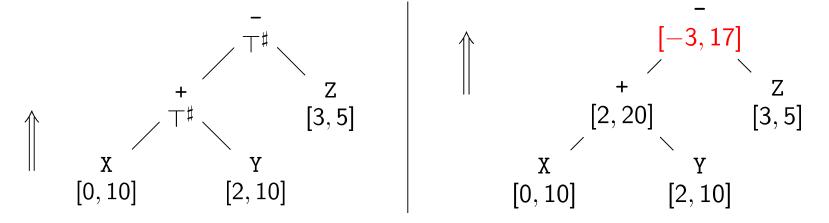
• 赋值操作-示例

$$\mathcal{Y}^{\sharp} \stackrel{\mathrm{def}}{=} C^{\sharp} \llbracket \, \mathbf{X} := \mathbf{X} + \mathbf{Y} - \mathbf{Z} \, \rrbracket \, \mathcal{X}^{\sharp}$$
 with  $\mathcal{X}^{\sharp} = \{ \, \mathbf{X} \mapsto [0, 10], \mathbf{Y} \mapsto [2, 10], \mathbf{Z} \mapsto [3, 5] \, \}$ 



• 赋值操作-示例

$$\mathcal{Y}^{\sharp} \stackrel{\mathrm{def}}{=} C^{\sharp} \llbracket \mathbf{X} := \mathbf{X} + \mathbf{Y} - \mathbf{Z} \rrbracket \mathcal{X}^{\sharp}$$
with  $\mathcal{X}^{\sharp} = \{ \mathbf{X} \mapsto [0, 10], \mathbf{Y} \mapsto [2, 10], \mathbf{Z} \mapsto [3, 5] \}$ 



$$\mathcal{Y}^{\sharp} = \{ \mathbf{X} \mapsto [-3, 17], \mathbf{Y} \mapsto [2, 10], \mathbf{Z} \mapsto [3, 5] \}$$

### • 条件测试

$$\mathcal{X}^{\sharp}(\mathtt{X}) = [a, b]$$
  $\mathcal{X}^{\sharp}(\mathtt{Y}) = [c, d]$ 

$$\bowtie \in \{=,<,>,\leq,\geq,\neq\}$$

• 区间加宽(区间格的高度是无穷的)

把增长的上界变成+oo 把减小的下界变成-oo

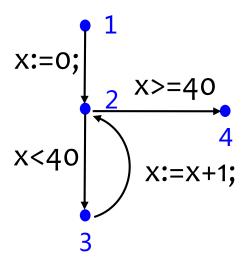
• 区间加宽(区间格的高度是无穷的)

$$\downarrow^{\sharp} \quad \nabla \quad X^{\sharp} \quad \stackrel{\text{def}}{=} \quad X^{\sharp} \\
[a,b] \quad \nabla \quad [c,d] \quad \stackrel{\text{def}}{=} \quad \left[ \left\{ \begin{array}{cc} a & \text{if } a \leq c \\ -\infty & \text{otherwise} \end{array} \right., \left\{ \begin{array}{cc} b & \text{if } b \geq d \\ +\infty & \text{otherwise} \end{array} \right]$$

把增长的上界变成+oo把减小的下界变成-oo

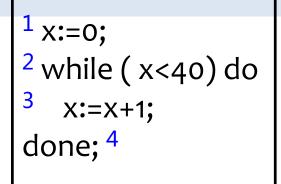
不稳定的要素

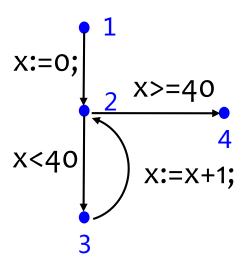
• 区间加宽一示例



```
1 x:=0;
2 while ( x<40) do
3     x:=x+1;
done; 4</pre>
```

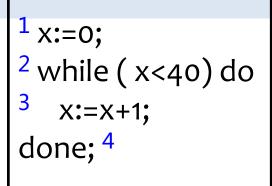
• 区间加宽一示例

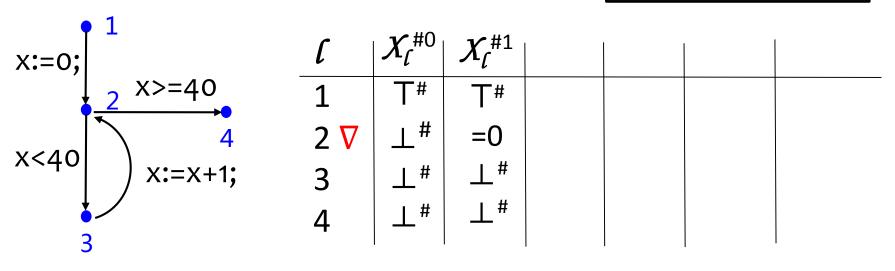




l	$X_{\ell}^{\#0}$			
1	T#			
2 ∇	⊥#			
3	⊥#			
4	⊥#			

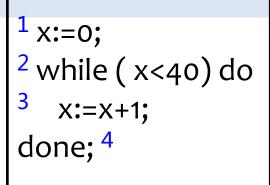
● 区间加宽—示例

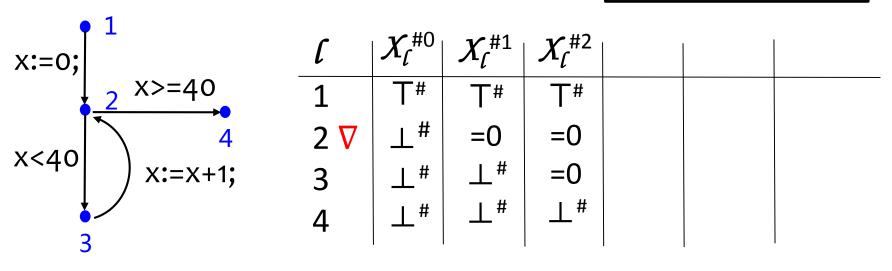




$$X_2^{\#1} = \bot^{\#} \nabla ([0,0] \cup^{\#} \bot^{\#}) = \bot^{\#} \nabla [0,0] = [0,0]$$

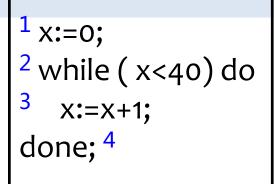
• 区间加宽一示例

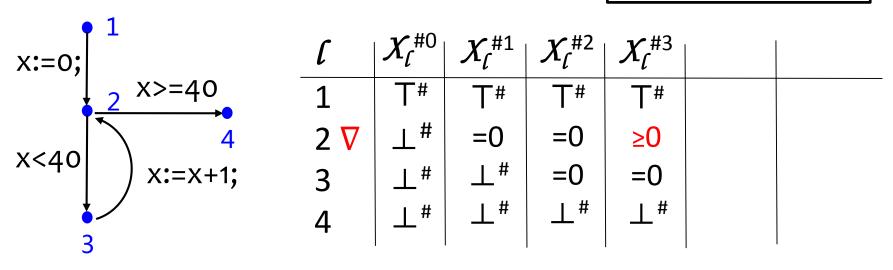




$$X_2^{\#1} = \bot^\# \nabla ([0,0] \cup^\# \bot^\#) = \bot^\# \nabla [0,0] = [0,0]$$
  
 $X_2^{\#2} = [0,0] \nabla ([0,0] \cup^\# \bot^\#) = [0,0] \nabla [0,0] = [0,0]$ 

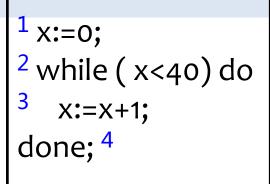
● 区间加宽—示例

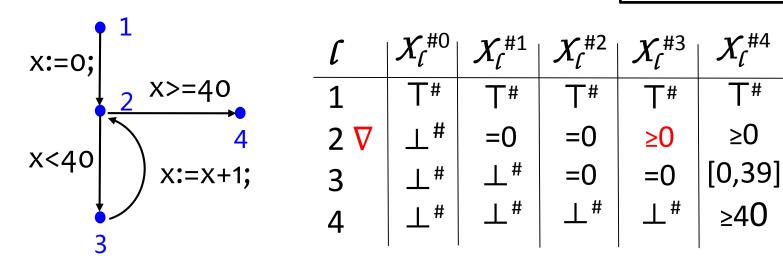




$$X_2^{\#1} = \bot^\# \nabla ([0,0] \cup^\sharp \bot^\#) = \bot^\# \nabla [0,0] = [0,0]$$
  
 $X_2^{\#2} = [0,0] \nabla ([0,0] \cup^\sharp \bot^\#) = [0,0] \nabla [0,0] = [0,0]$   
 $X_2^{\#3} = [0,0] \nabla ([0,0] \cup^\sharp [1,1]) = [0,0] \nabla [0,1] = [0,+\infty[$ 

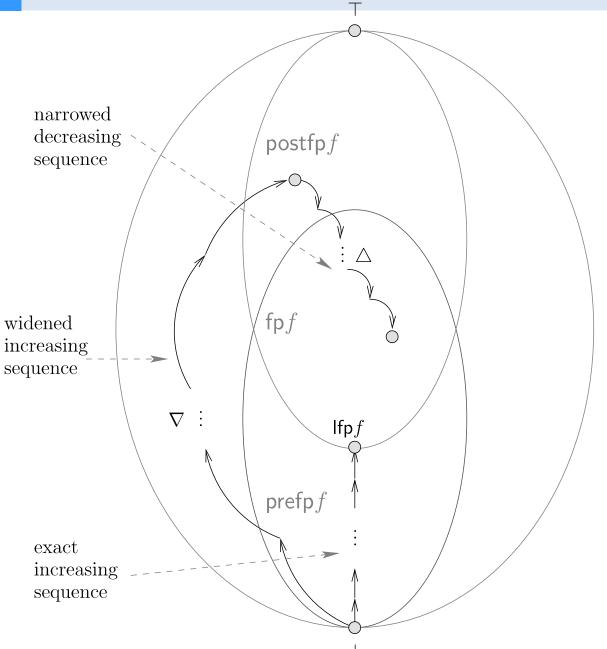
● 区间加宽—示例



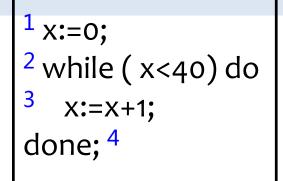


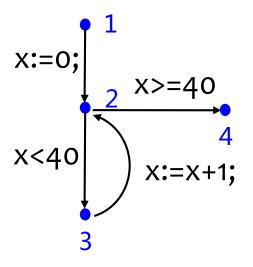
$$X_2^{\#1} = \bot^\# \nabla ([0,0] \cup^\sharp \bot^\#) = \bot^\# \nabla [0,0] = [0,0]$$
  
 $X_2^{\#2} = [0,0] \nabla ([0,0] \cup^\sharp \bot^\#) = [0,0] \nabla [0,0] = [0,0]$   
 $X_2^{\#3} = [0,0] \nabla ([0,0] \cup^\sharp [1,1]) = [0,0] \nabla [0,1] = [0,+\infty[$   
 $X_2^{\#4} = [0,+\infty[ \nabla ([0,0] \cup^\sharp [1,40]) = [0,+\infty[ \nabla [0,40] = [0,+\infty[$ 

# 基于加宽/变窄的不动点迭代



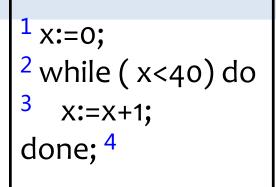
• 区间变窄一示例

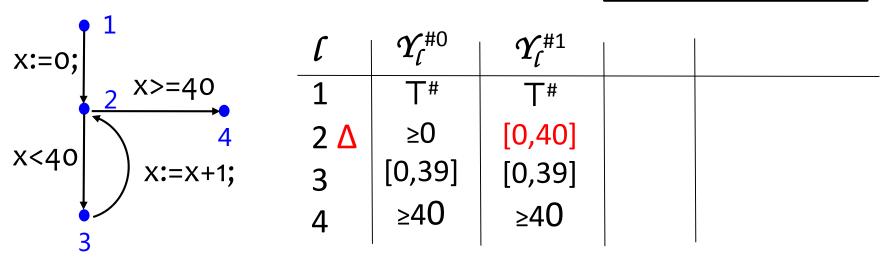




l	$Y_{\!\ell}^{\sharp 0}$		
1	Τ#		
2 △	≥0		
3	≥0 [0,39] ≥40		
4	≥40		

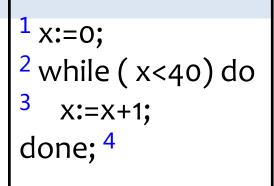
• 区间变窄一示例

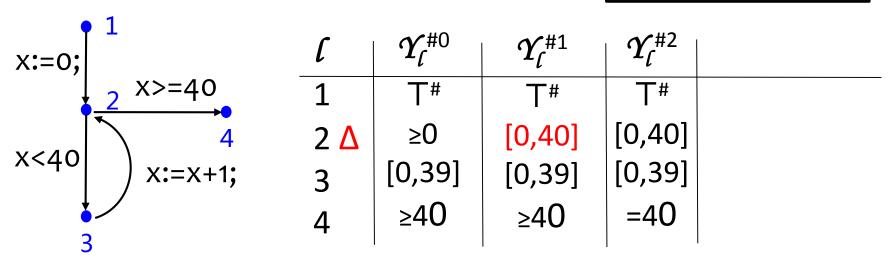




$$\Upsilon_2^{\#1} = [0,+\infty[ \Delta ([0,0] \cup^{\#} [1,40]) = [0,+\infty[ \Delta [0,40] = [0,40]]$$

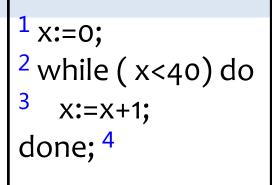
● 区间变窄—示例

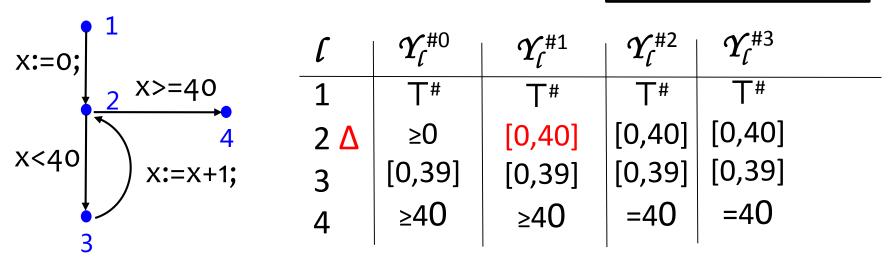




$$\Upsilon_2^{\#1} = [0,+\infty[ \Delta ([0,0] \cup^{\sharp} [1,40]) = [0,+\infty[ \Delta [0,40] = [0,40]$$
  
 $\Upsilon_2^{\#2} = [0,40[ \Delta ([0,0] \cup^{\sharp} [1,40]) = [0,40] \Delta [0,40] = [0,40]$ 

● 区间变窄—示例





$$\Upsilon_2^{\#1} = [0,+\infty[ \Delta ([0,0] \cup^{\sharp} [1,40]) = [0,+\infty[ \Delta [0,40] = [0,40]$$
  
 $\Upsilon_2^{\#2} = [0,40[ \Delta ([0,0] \cup^{\sharp} [1,40]) = [0,40] \Delta [0,40] = [0,40]$ 

从而得到:在2处 x∈[0,40], 在4处x=40

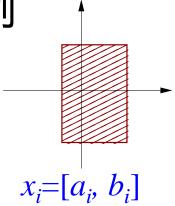
### 区间抽象域的局限性

- 区间抽象域是非关系型抽象域
  - > 只能表达单个变量的取值范围
  - > 不能表达多个变量之间的关系

### 数值抽象域

• 以两个数值抽象域为例

> 区间抽象域



> 线性等式抽象域

$$\sum_{i} a_{i} x_{i} = c$$

- 域表示
  - > 约束表示:线性等式系统 Ax=b
    - -x ∈ R<sup>n</sup> 表示程序变量x1, ..., xn构成的向量
    - A ∈ R<sup>m×n</sup> , b ∈ R<sup>m</sup> 为系数 , 由静态分析自动分析得到

$$\begin{cases} 2X + Y + Z = 19 \\ 2X + Y - Z = 9 \\ 3Z = 15 \end{cases}$$

- 域表示
  - > 约束表示:线性等式系统 Ax=b
  - > 规范型:唯一表示
    - 化简后的行阶梯形矩阵 (Reduced Row Echelon Form)
    - 高斯消元法 ( Gaussian Elimination )
      - 把线性等式系统转换为其规范型

$$\begin{cases} 2X + Y + Z = 19 \\ 2X + Y - Z = 9 \\ 3Z = 15 \end{cases} \longrightarrow \begin{cases} X + 0.5Y = 7 \\ Z = 5 \end{cases}$$

● 域表示: <A, b>

$$\wp(R) \xrightarrow{\gamma_b} LE^\#$$
 $\gamma(\langle A,b \rangle) \triangleq \{x \in R^n \mid Ax=b\}$ 
 $g(X) \triangleq \{\langle A,b \rangle \mid$ 

#### • 域操作

$$\mathcal{X}^{\sharp} \cap^{\sharp} \mathcal{Y}^{\sharp} \triangleq \textit{Gauss} \left( \left\langle \begin{bmatrix} \mathsf{A}_{\mathcal{X}^{\sharp}} \\ \mathsf{A}_{\mathcal{Y}^{\sharp}} \end{bmatrix}, \begin{bmatrix} \mathsf{b}_{\mathcal{X}^{\sharp}} \\ \mathsf{b}_{\mathcal{Y}^{\sharp}} \end{bmatrix} \right\rangle \right)$$

$$\mathcal{X}^{\sharp} = {}^{\sharp} \mathcal{Y}^{\sharp} \iff \mathsf{A}_{\mathcal{X}^{\sharp}} = \mathsf{A}_{\mathcal{Y}^{\sharp}} \text{ and } \mathsf{b}_{\mathcal{X}^{\sharp}} = \mathsf{b}_{\mathcal{Y}^{\sharp}}$$

$$\mathcal{X}^{\sharp} \sqsubseteq^{\sharp} \mathcal{Y}^{\sharp} \iff \mathcal{X}^{\sharp} \cap^{\sharp} \mathcal{Y}^{\sharp} = {}^{\sharp} \mathcal{X}^{\sharp}$$

$$\mathsf{C}[\![ \sum_{j} \alpha_{j} \mathsf{V}_{j} - \beta = 0 ]\!]^{\sharp} (\mathcal{X}^{\sharp}) \triangleq \textit{Gauss} \left( \left\langle \begin{bmatrix} \mathsf{A}_{\mathcal{X}^{\sharp}} \\ \alpha_{1} \cdots \alpha_{n} \end{bmatrix}, \begin{bmatrix} \mathsf{b}_{\mathcal{X}^{\sharp}} \\ \beta \end{bmatrix} \right\rangle \right)$$

$$\mathsf{C}[\![ e \bowtie 0 ]\!]^{\sharp} (\mathcal{X}^{\sharp}) \triangleq \mathcal{X}^{\sharp} \qquad \text{for other tests}$$

$$\sqsubseteq^{\sharp}, =^{\sharp}, \cap^{\sharp}, =^{\sharp}, C[\![\sum_{j} \alpha_{j} V_{j} - \beta = 0]\!]^{\sharp}$$
是精确的,因为
$$\mathcal{X}^{\sharp} \sqsubseteq^{\sharp} \mathcal{Y}^{\sharp} \iff \gamma(\mathcal{X}^{\sharp}) \subseteq \gamma(\mathcal{Y}^{\sharp}), \quad \gamma(\mathcal{X}^{\sharp} \cap^{\sharp} \mathcal{Y}^{\sharp}) = \gamma(\mathcal{X}^{\sharp}) \cap \gamma(\mathcal{Y}^{\sharp})$$

#### • 域操作

$$\begin{split} \mathsf{C}[\![\, \mathsf{V}_{\mathtt{j}} &:= \sum_{i} \alpha_{i} \mathsf{V}_{\mathtt{i}} + \beta \,]\!]^{\sharp} (\mathcal{X}^{\sharp}) & \stackrel{\triangle}{=} \\ & \text{if } \alpha_{j} \neq 0, \mathcal{X}^{\sharp} \text{ where } \mathsf{V}_{\mathtt{j}} \text{ is replaced with} (\mathsf{V}_{\mathtt{j}} - \sum_{i \neq j} \alpha_{i} \mathsf{V}_{\mathtt{i}} - \beta) / \alpha_{j} \\ & \text{if } \alpha_{j} = 0, (\mathsf{C}[\![\, \sum_{i} \alpha_{i} \mathsf{V}_{\mathtt{i}} - \mathsf{V}_{\mathtt{j}} + \beta = 0 \,]\!]^{\sharp} \circ \mathsf{C}[\![\, \mathsf{V}_{\mathtt{j}} := ?(-\infty, +\infty) \,]\!]^{\sharp}) (\mathcal{X}^{\sharp}) \end{split}$$

$$\mathbb{C}[V_{j} := ?(-\infty, +\infty)]^{\sharp}(\mathcal{X}^{\sharp}) \stackrel{\Delta}{=} \text{GuassElimination}(\langle A_{\chi^{\#}}, b_{\chi^{\#}} \rangle, V_{j})$$

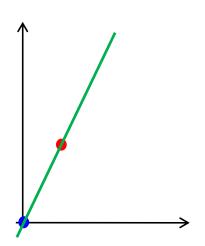
$$\mathsf{C}[\![\,\mathsf{V}_{\mathtt{j}} := e\,]\!]^{\sharp}(\mathcal{X}^{\sharp}) \stackrel{\Delta}{=} \mathsf{C}[\![\,\mathsf{V}_{\mathtt{j}} := ?(-\infty, +\infty)\,]\!]^{\sharp}(\mathcal{X}^{\sharp})$$
 for other assignments

$$C[V_j := \sum_i \alpha_i V_i + \beta]^{\sharp}, C[V_j := ?(-\infty, +\infty)]^{\sharp}$$
 是精确的

域操作: X<sup>♯</sup> ∪<sup>♯</sup> Y<sup>♯</sup>

目标:给定  $\gamma(\mathcal{X}^{\sharp}) = \{x \mid Ax = b\}, \gamma(\mathcal{Y}^{\sharp}) = \{x \mid A'x = b'\}$ 

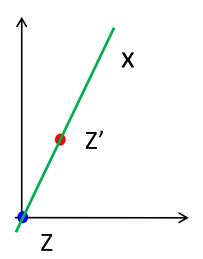
求仿射闭包  $\mathcal{X}_H^{\sharp}$  使得  $\gamma(\mathcal{X}^{\sharp}) \subseteq \gamma(\mathcal{X}_H^{\sharp})$ 且  $\gamma(\mathcal{Y}^{\sharp}) \subseteq \gamma(\mathcal{X}_H^{\sharp})$ 



域操作: X<sup>♯</sup> ∪<sup>♯</sup> y<sup>♯</sup>

目标:给定  $\gamma(\mathcal{X}^{\sharp}) = \{x \mid Ax = b\}, \gamma(\mathcal{Y}^{\sharp}) = \{x \mid A'x = b'\}$  求仿射闭包  $\mathcal{X}_{H}^{\sharp}$  使得  $\gamma(\mathcal{X}^{\sharp}) \subseteq \gamma(\mathcal{X}_{H}^{\sharp})$  且  $\gamma(\mathcal{Y}^{\sharp}) \subseteq \gamma(\mathcal{X}_{H}^{\sharp})$ 

$$\gamma(\mathcal{X}_{H}^{\sharp}) = \left\{ x \middle| \begin{array}{c} x = \sigma_{1}z + \sigma_{2}z' \wedge \sigma_{1} + \sigma_{2} = 1 \wedge \\ Az = b \wedge A'z' = b' \end{array} \right\}$$



域操作: X<sup>♯</sup> ∪<sup>♯</sup> Y<sup>♯</sup>

目标:给定  $\gamma(\mathcal{X}^{\sharp}) = \{x \mid Ax = b\}, \gamma(\mathcal{Y}^{\sharp}) = \{x \mid A'x = b'\}$  求仿射闭包  $\mathcal{X}_{H}^{\sharp}$  使得  $\gamma(\mathcal{X}^{\sharp}) \subseteq \gamma(\mathcal{X}_{H}^{\sharp})$  且  $\gamma(\mathcal{Y}^{\sharp}) \subseteq \gamma(\mathcal{X}_{H}^{\sharp})$ 

$$\gamma(\mathcal{X}_{H}^{\sharp}) = \left\{ x \middle| \begin{array}{c} x = \sigma_{1}z + \sigma_{2}z' \wedge \sigma_{1} + \sigma_{2} = 1 \wedge \\ Az = b \wedge A'z' = b' \end{array} \right\}$$
引入变量 
$$y = \sigma_{1}z$$

$$y' = \sigma_{2}z'$$

$$\gamma(\mathcal{X}_{AH}^{\sharp}) = \left\{ x \middle| \begin{array}{c} x = y + y' & \wedge \sigma_1 + \sigma_2 = 1 & \wedge \\ Ay \leq \sigma_1 b & \wedge A'y' \leq \sigma_2 b' \end{array} \right\}$$

从  $\gamma(\mathcal{X}_{AH}^\sharp)$  中投影掉变量  $\sigma_1,\sigma_2,y,y'$  即可

 $\mathcal{X}^{\sharp} \cup^{\sharp} \mathcal{Y}^{\sharp}$  是最佳的但不是精确的

域操作: X<sup>♯</sup> ∪<sup>♯</sup> Y<sup>♯</sup>

**目标**: 给定 
$$\gamma(\mathcal{X}^{\sharp}) = \{x \mid Ax = b\}, \gamma(\mathcal{Y}^{\sharp}) = \{x \mid A'x = b'\}$$

求仿射闭包  $\mathcal{X}_H^{\sharp}$  使得  $\gamma(\mathcal{X}^{\sharp}) \subseteq \gamma(\mathcal{X}_H^{\sharp})$ 且  $\gamma(\mathcal{Y}^{\sharp}) \subseteq \gamma(\mathcal{X}_H^{\sharp})$ 

$$\gamma(\mathcal{X}_{H}^{\sharp}) = \left\{ x \middle| \begin{array}{c} x = \sigma_{1}z + \sigma_{2}z' \wedge \sigma_{1} + \sigma_{2} = 1 \wedge \\ Az = b \wedge A'z' = b' \end{array} \right\}$$
引入变量 
$$y = \sigma_{1}z$$

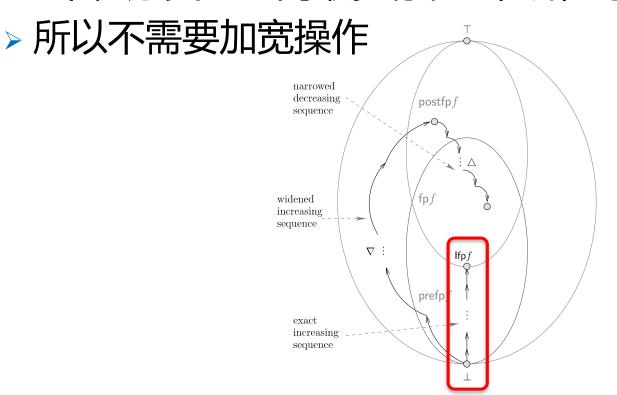
$$y' = \sigma_{2}z'$$

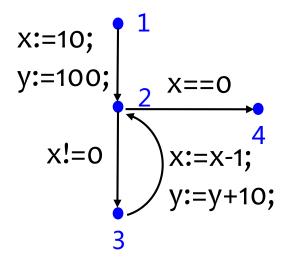
$$\gamma(\mathcal{X}_{AH}^{\sharp}) = \left\{ x \left| \begin{array}{cc} x = y + y' & \wedge \sigma_1 + \sigma_2 = 1 & \wedge \\ Ay \leq \sigma_1 b & \wedge A'y' \leq \sigma_2 b' \end{array} \right. \right\}$$

从  $\gamma(\mathcal{X}_{AH}^{\sharp})$  中

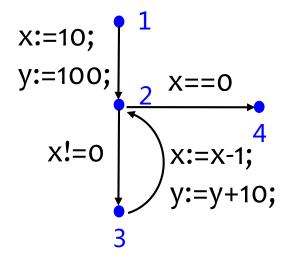
从  $\gamma(\mathcal{X}_{AH}^{\sharp})$  中投影掉变量  $\sigma_1, \sigma_2, y, y'$  即可

- 程序变量间线性等式集合构成的格的高度是有穷的
  - » n个程序变量之间最多存在n个线性等式



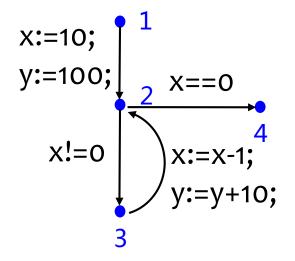


```
1 x:=10; y:=100;
2 while ( x!=0) do
3     x:=x-1; y:=y+10;
    done 4
```



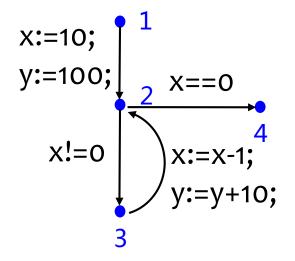
l	$ \mathcal{X}_{\ell}^{\sharp 0} $		
1	T#		
2	⊥#		
3	#		
4	#		

```
1 x:=10; y:=100;
2 while ( x!=0) do
3     x:=x-1; y:=y+10;
    done 4
```



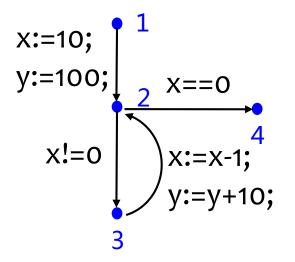
l	$X_{\ell}^{\#0}$	$X_{\ell}^{\sharp 1}$		
1	T#	T#		
2	│	(10,100)		
3	⊥#	⊥#		
4	#			

```
1 x:=10; y:=100;
2 while (x!=0) do
3 x:=x-1; y:=y+10;
done 4
```



l	$X_{\ell}^{\#0}$	$X_{\ell}^{\sharp 1}$	$X_l^{\#2}$	
1	T#	T#	T#	
2	│		(10,100)	
3	#	⊥#	(10,100)	
4	│		│	

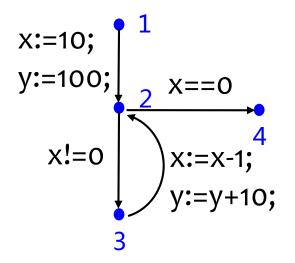
```
1 x:=10; y:=100;
2 while (x!=0) do
3 x:=x-1; y:=y+10;
done 4
```



l	$\mathcal{X}_{l}^{\sharp 0}$	$X_{\ell}^{\sharp 1}$	$X_l^{\#2}$	$X_l^{\#3}$
1	T#	T#	T#	Т#
2	⊥#	(10,100)	(10,100)	10x+y=200
3	上#	上#	(10,100)	(10,100)
4	│	⊥#	⊥#	⊥#

```
1 x:=10; y:=100;
2 while (x!=0) do
3 x:=x-1; y:=y+10;
done 4
```

$$X_2^{#3} = \{(10,100)\} \cup \#\{(9,110)\} = \{(x,y) \mid 10x+y=200\}$$



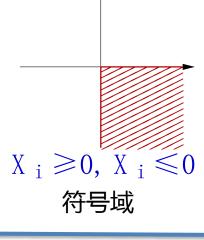
l	$ \mathcal{X}_{\ell}^{ ext{#0}} $	$\mathcal{X}_{l}^{\sharp 1}$	$\mathcal{X}_{l}^{ ext{#2}}$	$X_l^{\#3}$	$\mathcal{X}_{\ell}^{$ #4
1	T#	T#	Τ#	Τ#	T#
2	⊥#		(10,100)	10x+y=200	10x+y=200
3	上#	上#	(10,100)	(10,100)	10x+y=200
4	⊥#	#	⊥#	⊥#	(0,200)

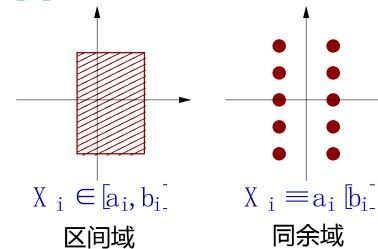
```
    1 x:=10; y:=100;
    2 while (x!=0) do
    3 x:=x-1; y:=y+10;
    done 4
```

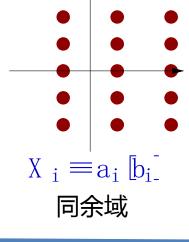
$$X_2^{#3} = \{(10,100)\} \cup \#\{(9,110)\} = \{(x,y) \mid 10x+y=200\}$$

### 数值抽象域—谱系

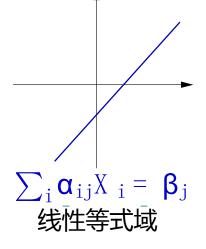
非关系型

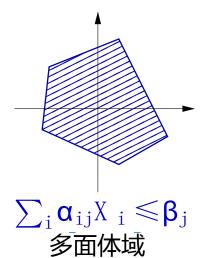


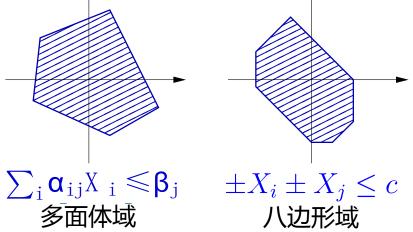




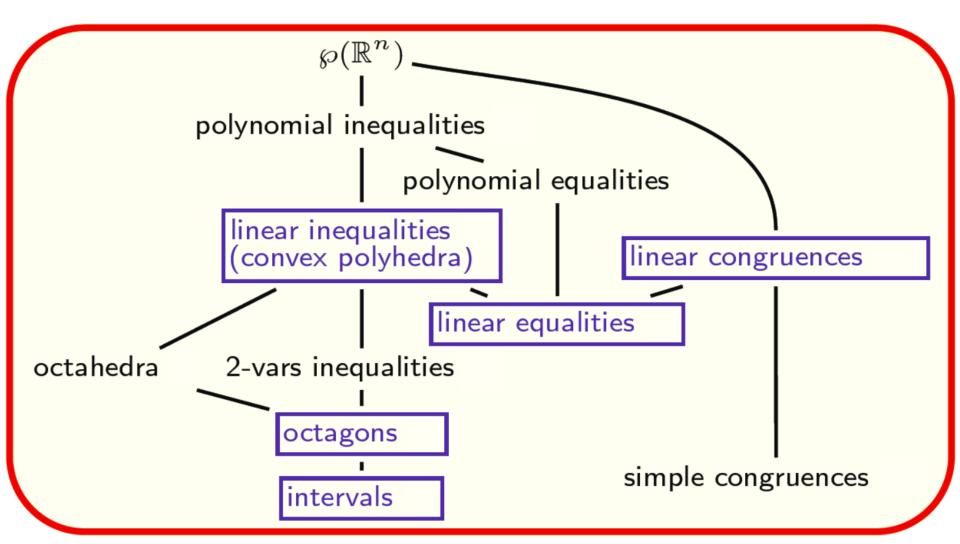
关系型







### 数值抽象域—谱系



## 谢谢!