

Stateless Model Checking for TSO and PSO

Parosh Aziz Abdulla Stavros Aronis Mohamed Faouzi Atig
Bengt Jonsson **Carl Leonardsson** Konstantinos Sagonas

Uppsala University, Sweden

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Goals

Stateless Model Checking

- Find **safety** errors...
 - in given test case (fixed input program)
 - for all interleavings
 - for all reorderings (TSO/PSO)
- Works on real code in C/pthreads

Valid Test Case (in this presentation)

- 1 Terminates in bounded time
- 2 Nondeterminism: Interleavings, Reordering

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Partial Order Reduction (SC)

volatile int x = 0, y = 0;

p	q
x = 1;	y = 1;
int a = y;	int b = x;

Executions

$p:wx1$	$p:wx1$	$p:wx1$
$p:ry0$	$q:wy1$	$q:wy1$
$q:wy1$	$p:ry1$	$q:rx1$
$q:rx1$	$q:rx1$	$p:ry1$
$q:wy1$	$q:wy1$	$q:wy1$
$q:rx0$	$p:wx1$	$p:wx1$
$p:wx1$	$q:rx1$	$p:ry1$
$p:ry1$	$p:ry1$	$q:rx1$

Partial Order Reduction (SC)

```
volatile int x = 0, y = 0;
```

<i>p</i>	<i>q</i>
x = 1;	y = 1;
int a = y;	int b = x;

Executions

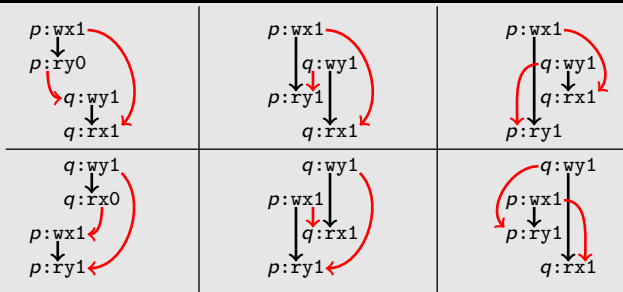
<i>p:wx1</i>	<i>p:wx1</i>	<i>p:wx1</i>
<i>p:ry0</i>	<i>q:wy1</i>	<i>q:wy1</i>
<i>q:wy1</i>	<i>p:ry1</i>	<i>q:rx1</i>
<i>q:rx1</i>	<i>q:rx1</i>	<i>p:ry1</i>
<i>q:wy1</i>	<i>q:wy1</i>	<i>q:wy1</i>
<i>q:rx0</i>	<i>p:wx1</i>	<i>p:wx1</i>
<i>p:wx1</i>	<i>q:rx1</i>	<i>p:ry1</i>
<i>p:ry1</i>	<i>p:ry1</i>	<i>q:rx1</i>

Partial Order Reduction (SC)

```
volatile int x = 0, y = 0;
```

<i>p</i>	<i>q</i>
<code>x = 1;</code>	<code>y = 1;</code>
<code>int a = y;</code>	<code>int b = x;</code>

Executions (Happens-Before)

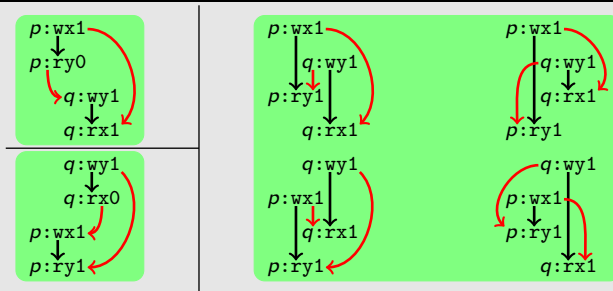


Partial Order Reduction (SC)

```
volatile int x = 0, y = 0;
```

p	q
$x = 1;$	$y = 1;$
$\text{int } a = y;$	$\text{int } b = x;$

Mazurkiewicz Traces ~ Equivalence Classes over Executions



Stateless Model Checking with DPOR

[Flanagan, Godefroid 2005]

Idea

- Explore one execution per Mazurkiewicz trace.
 - Cover all observable behaviours.
- Keep only one execution in memory.
- Examine happens-before relation to find the next trace.


```
volatile int x = 0, y = 0;
```

p	q
$x = 1;$	$y = 1;$
$\text{int } a = y;$	$\text{int } b = x;$

```
volatile int x = 0, y = 0;
```

p	q
x = 1;	y = 1;
int a = y;	int b = x;

$p:wx1$

$p:ry0$

$q:wy1$

$q:rx1$

```
volatile int x = 0, y = 0;
```

p	q
$x = 1;$	$y = 1;$
$\text{int } a = y;$	$\text{int } b = x;$

$p:wx1$



$p:ry0$

$q:wy1$



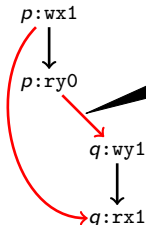
$q:rx1$

HB: Necessary Edges

- Program order (under SC)
- Thread spawning to child's first event
- ...

```
volatile int x = 0, y = 0;
```

p	q
$x = 1;$	$y = 1;$
$\text{int } a = y;$	$\text{int } b = x;$

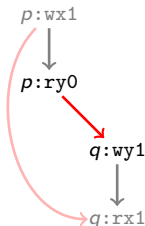


HB: Races

- Conflicting stores and loads.
- ...

```
volatile int x = 0, y = 0;
```

p	q
$x = 1;$	$y = 1;$
$\text{int } a = y;$	$\text{int } b = x;$

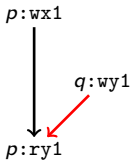


Reverse Races

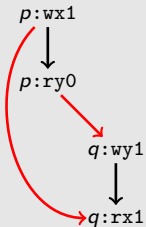
- Start from the end of the execution.

```
volatile int x = 0, y = 0;
```

p	q
$x = 1;$	$y = 1;$
$\text{int } a = y;$	$\text{int } b = x;$

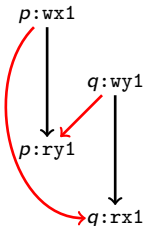


Previous Trace

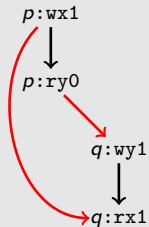


```
volatile int x = 0, y = 0;
```

p	q
$x = 1;$	$y = 1;$
$\text{int } a = y;$	$\text{int } b = x;$

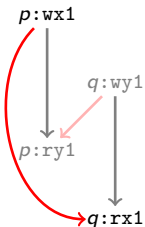


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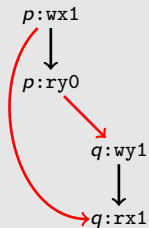


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volatile int x = 0, y = 0;
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p	q
$x = 1;$	$y = 1;$
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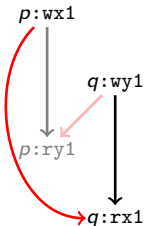


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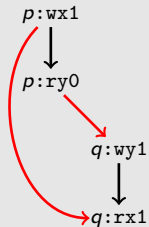



```
volatile int x = 0, y = 0;
```

p	q
$x = 1;$	$y = 1;$
$\text{int } a = y;$	$\text{int } b = x;$

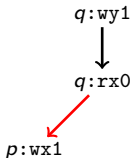


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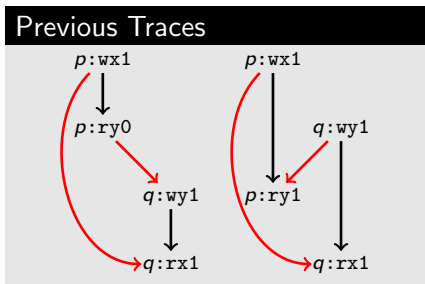


```
volatile int x = 0, y = 0;
```

p	q
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$\text{int } a = y;$	$\text{int } b = x;$

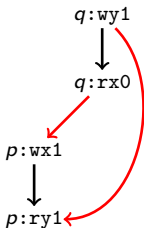


Previous Traces

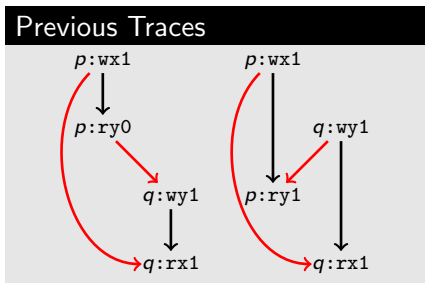


```
volatile int x = 0, y = 0;
```

p	q
$x = 1;$	$y = 1;$
$\text{int } a = y;$	$\text{int } b = x;$



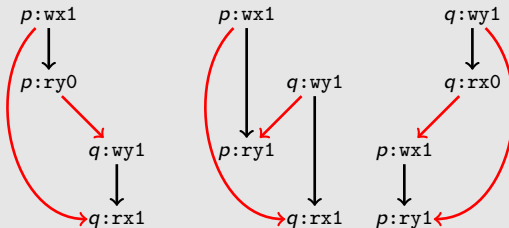
Previous Traces



```
volatile int x = 0, y = 0;
```

p	q
$x = 1;$	$y = 1;$
$\text{int } a = y;$	$\text{int } b = x;$

All Traces



TSO

- Relaxes $W \rightarrow R$
- Store forwarding to own reads (ROWE)
- Operational semantics: store buffer per thread

Defining Traces for TSO

Extend Mazurkiewicz Traces to TSO

- Suitable equivalence classes
- Compatible with DPOR

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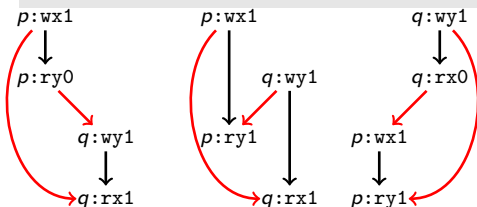
Shasha-Snir Traces

Capture observable order.

Defining Traces for TSO

Extend Mazurkiewicz Traces to TSO

- Suitable equivalence classes
- Compatible with DPOR



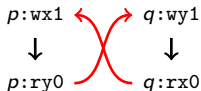
Shasha-Snir Traces

Same as
Mazurkiewicz traces
under SC!

Defining Traces for TSO

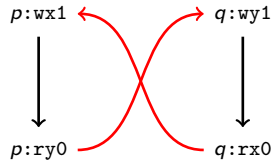
Extend Mazurkiewicz Traces to TSO

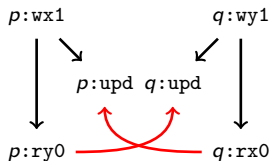
- Suitable equivalence classes
- **Compatible with DPOR**



Shasha-Snir Traces

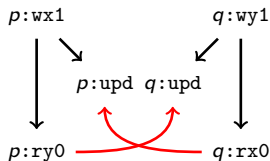
Not partial order.





Getting Rid of Cycles

- Operational semantics
- Ignores *which* events are reordered
- Canonical?



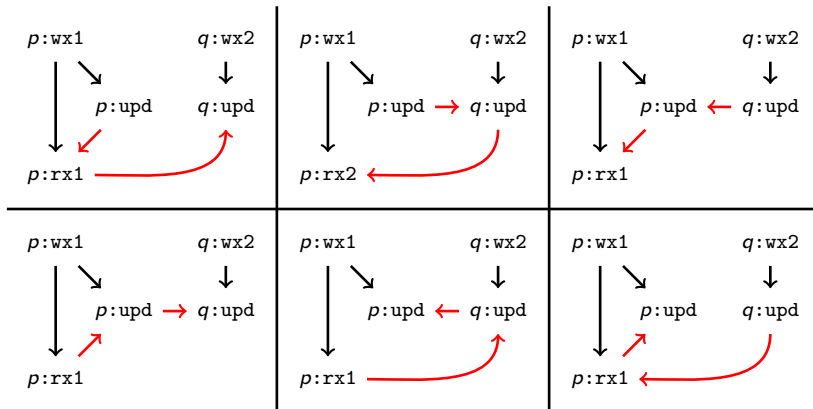
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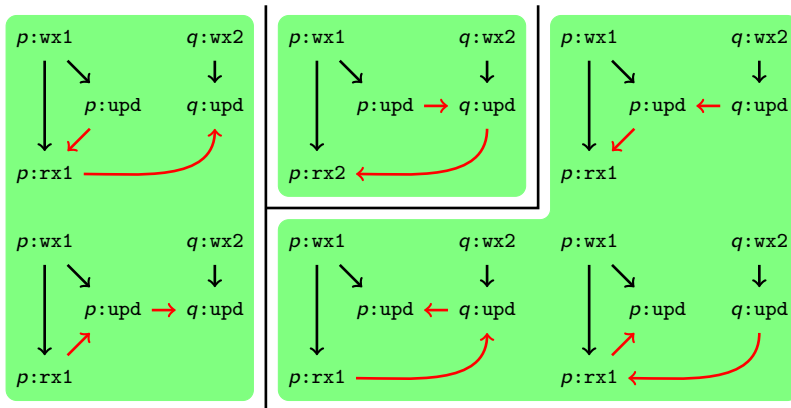
```
volatile int x = 0;
```

p	q
$x = 1;$ $\text{int } a = x;$	$x = 2;$

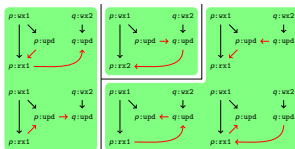
Problem



Problem



Solution: Chronological Traces

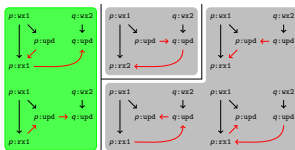


Two Rules

No order between load and update...

- 1 of the same thread.
- 2 when the update is “hidden” from the load.

Solution: Chronological Traces

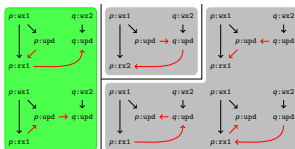
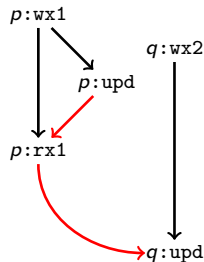
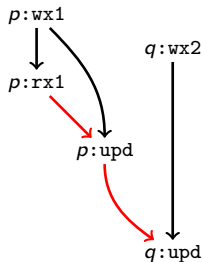


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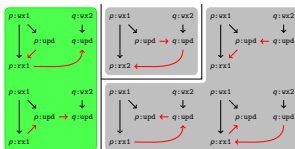
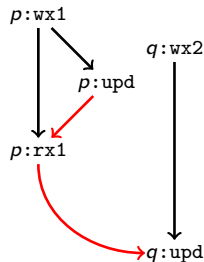
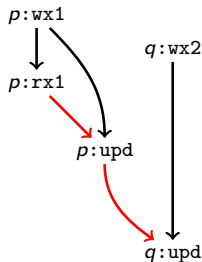


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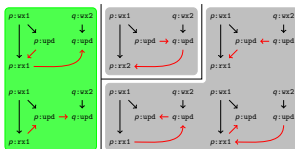
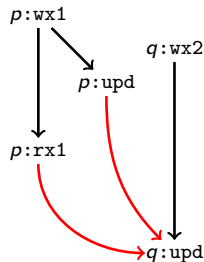
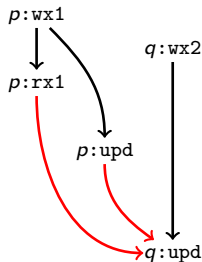


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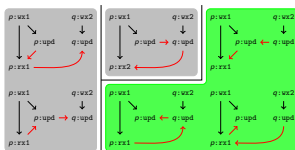
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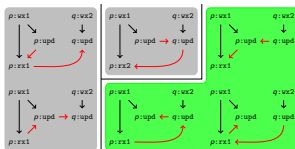
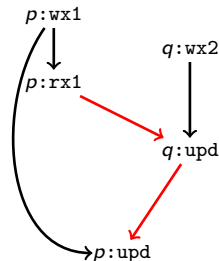
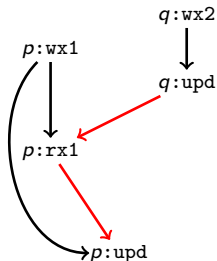
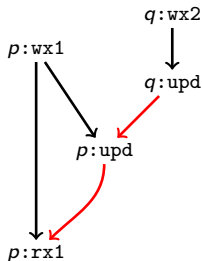
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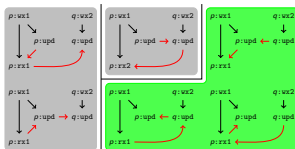
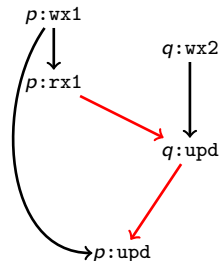
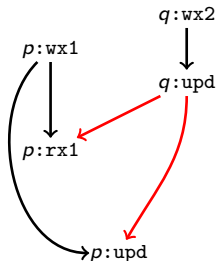
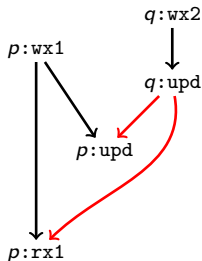


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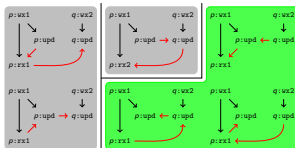
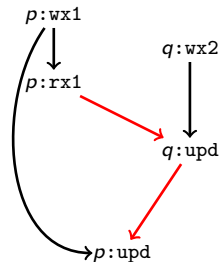
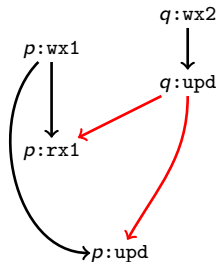
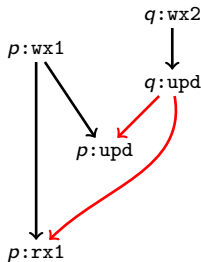


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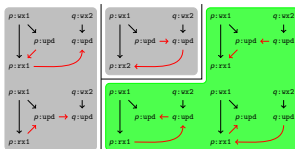
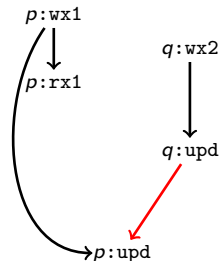
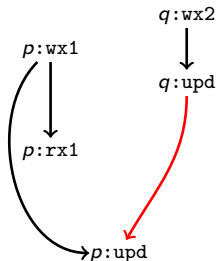
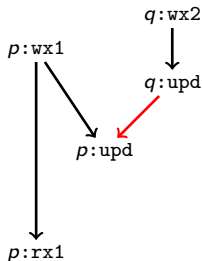


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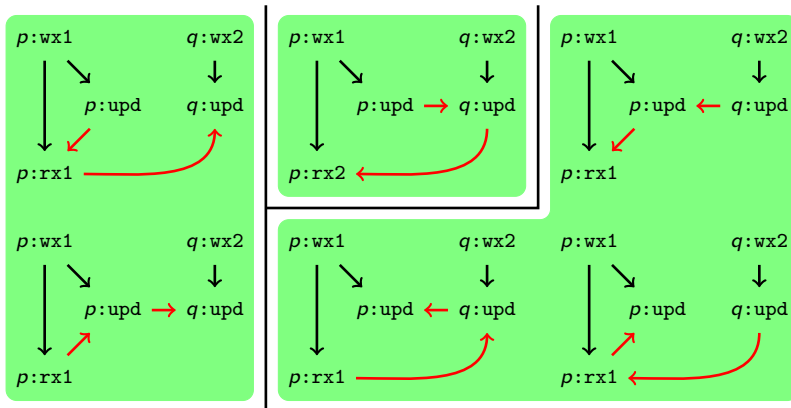


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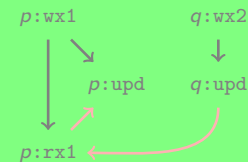
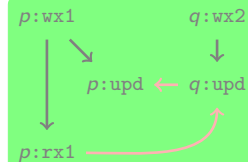
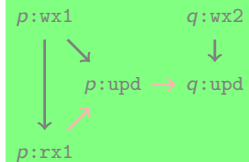
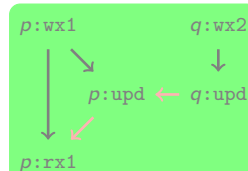
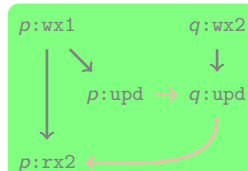
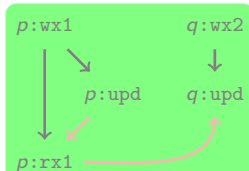
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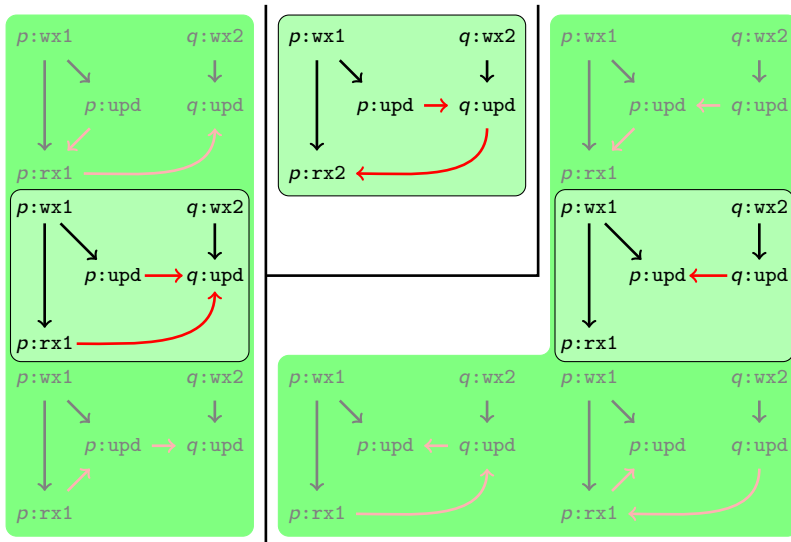
Solution: Chronological Traces



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Our Contribution: Chronological Traces

- 1-to-1 with Shasha-Snir traces
- Compatible with DPOR
 - Executions under operational TSO/PSO semantics
 - Efficiently computable happens-before relation

Implementation: Nidhugg

Nidhugg

- Source-DPOR [Abdulla et al. 2014]
 - Near optimal.
 - Straight-forward.
- Works on C/Pthreads via LLVM IR.
- Runs on *one compilation* of test case.

Open source: <https://github.com/nidhugg/nidhugg>

Results: Time Consumption											
	fence	LB	CBMC			goto-instrument			Nidhugg		
			SC	TSO	PSO	SC	TSO	PSO	SC	TSO	PSO
apr_1.c	-	5	t/o	t/o	t/o	t/o	!	!	5.88	6.06	16.98
apr_2.c	-	5	t/o	t/o	t/o	!	!	!	2.60	2.20	5.39
dcl_singleton.c	-	7	5.95	31.47	*18.01	5.33	5.36	*0.18	0.08	0.08	*0.08
dcl_singleton.c	ps0	7	5.88	30.98	29.45	5.20	5.18	5.17	0.08	0.08	0.08
dekker.c	-	10	2.42	*3.17	*2.84	1.68	*4.00	*220.11	0.10	*0.11	*0.09
dekker.c	tso	10	2.39	5.65	*3.51	1.62	297.62	t/o	0.11	0.12	*0.08
dekker.c	ps0	10	2.55	5.31	4.83	1.72	428.86	t/o	0.11	0.12	0.12
fib_false.c	-	-	*1.63	*3.38	*3.00	*1.60	*1.58	*1.56	*2.39	*5.57	*6.20
fib_false_join.c	-	-	*0.98	*1.10	*1.91	*1.31	*0.88	*0.80	*0.32	*0.62	*0.71
fib_true.c	-	-	6.28	9.39	7.72	6.32	7.63	7.62	25.83	75.06	86.32
fib_true_join.c	-	-	6.61	8.37	10.81	7.09	5.94	5.92	1.20	2.88	3.19
indexer.c	-	5	193.01	210.42	214.03	191.88	70.42	69.38	0.10	0.09	0.09
lamport.c	-	8	7.78	*11.63	*10.53	6.89	t/o	t/o	0.08	*0.08	*0.08
lamport.c	tso	8	7.60	26.31	*15.85	6.80	513.67	t/o	0.09	0.08	*0.07
lamport.c	ps0	8	7.72	30.92	27.51	7.43	t/o	t/o	0.08	0.08	0.08
parker.c	-	10	12.34	*91.99	*86.10	11.63	9.70	9.65	1.50	*0.09	*0.08
parker.c	ps0	10	12.72	141.24	166.75	11.76	10.66	10.64	1.50	1.92	2.94
peterson.c	-	-	0.35	*0.38	*0.35	0.18	*0.20	*0.21	0.07	*0.07	*0.07
peterson.c	tso	-	0.35	0.39	*0.35	0.19	0.18	0.56	0.07	0.07	*0.07
peterson.c	ps0	-	0.35	0.41	0.40	0.18	0.18	0.19	0.07	0.07	0.08
pgsql.c	-	8	19.80	60.66	*4.63	21.03	46.57	*296.77	0.08	0.07	*0.08
pgsql.c	ps0	8	23.93	71.15	121.51	19.04	t/o	t/o	0.07	0.07	0.08
pgsql_bnd.c	ps0	(4)	3.57	9.55	12.68	3.59	t/o	t/o	89.44	106.04	112.60
stack_safe.c	-	-	44.53	516.01	496.36	45.11	42.39	42.50	0.34	0.36	0.43
stack_unsafe.c	-	-	*1.40	*1.87	*2.08	*1.00	*0.81	*0.79	*0.08	*0.08	*0.09
szymanski.c	-	-	0.40	*0.44	*0.43	0.23	*0.89	*1.16	0.07	*0.13	*0.07
szymanski.c	tso	-	0.40	0.50	*0.43	0.23	0.23	2.48	0.08	0.08	*0.07
szymanski.c	ps0	-	0.39	0.50	0.49	0.23	0.24	0.24	0.08	0.08	0.08

Results: Time Consumption											
	fence	LB	CBMC			goto-instrument			Nidhugg		
			SC	TSO	PSO	SC	TSO	PSO	SC	TSO	PSO
apr_1.c	-	5	t/o	t/o	t/o	t/o	!	!	5.88	6.06	16.98
apr_2.c	-	5	t/o	t/o	t/o	!	!	!	2.60	2.20	5.39
dcl_singleton.c	-	7	5.95	31.47	*18.01	5.33	5.36	*0.18	0.08	0.08	*0.08
dcl_singleton.c	pso	7	5.88	30.98	29.45	5.20	5.18	5.17	0.08	0.08	0.08
dekker.c	-	10	2.42	*3.17	*2.84	1.68	*4.00	*220.11	0.10	*0.11	*0.09
dekker.c	tso	10	2.39	5.65	*3.51	1.62	297.62	t/o	0.11	0.12	*0.08
dekker.c	pso	10	2.55	5.31	4.83	1.72	428.86	t/o	0.11	0.12	0.12
fib_false.c			*1.62	*2.28	*2.00	*1.60	*1.58	*1.56	*2.20	*5.57	*6.20
fib_false_join.c										0.62	*0.71
fib_true.c										0.06	86.32
fib_true_join.c										0.88	3.19
indexer.c										0.09	0.09
lambport.c										0.08	*0.08
lambport.c										0.08	*0.07
lambport.c										0.08	0.08
parker.c										0.09	*0.08
parker.c	pso	10	12.72	141.24	166.75	11.76	10.66	10.64	1.50	1.92	2.94
peterson.c	-	-	0.35	*0.38	*0.35	0.18	*0.20	*0.21	0.07	*0.07	*0.07
peterson.c	tso	-	0.35	0.39	*0.35	0.19	0.18	0.56	0.07	0.07	*0.07
peterson.c	pso	-	0.35	0.41	0.40	0.18	0.18	0.19	0.07	0.07	0.08
pgsql.c	-	8	19.80	60.66	*4.63	21.03	46.57	*296.77	0.08	0.07	*0.08
pgsql.c	pso	8	23.93	71.15	121.51	19.04	t/o	t/o	0.07	0.07	0.08
pgsql_bnd.c	pso	(4)	3.57	9.55	12.68	3.59	t/o	t/o	89.44	106.04	112.60
stack_safe.c	-	-	44.53	516.01	496.36	45.11	42.39	42.50	0.34	0.36	0.43
stack_unsafe.c	-	-	*1.40	*1.87	*2.08	*1.00	*0.81	*0.79	*0.08	*0.08	*0.09
szymanski.c	-	-	0.40	*0.44	*0.43	0.23	*0.89	*1.16	0.07	*0.13	*0.07
szymanski.c	tso	-	0.40	0.50	*0.43	0.23	0.23	2.48	0.08	0.08	*0.07
szymanski.c	pso	-	0.39	0.50	0.49	0.23	0.24	0.24	0.08	0.08	0.08

Robust

Nidhugg

SC TSO PSO

stack_safe.c 0.34 0.36 0.43

Future Work

Ongoing Work: POWER

- More relaxed model
- Order enforced by complex event interaction
- #traces probably similar to under SC
- Techniques carry over to ARM, Alpha, etc.