

Game theory

A course for the
MSc in ICT for Internet and multimedia

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Static games of complete information

The simplest form of multi-player games

Game with multiple players

- How do multiple players interact? *una per giocatore*
 - ▣ We assume they have a payoff (utility) function
- Remember that rational players move to maximize of their own payoffs
- What is the simplest interaction like this?
 - ▣ certainly not sports or dices where most moves are random – we will not see them in the course
 - ▣ not even board games - they are closer, still they require some extensions

Static games of complete information

- **Static** (all players move together; they do not necessarily play simultaneously, but without knowledge of everybody else's move)
- **Complete information** (meaning anybody's payoff function is known)
 - most games within this class are actually “artificial” games (theoretical models)
 - examples of actual games: Odds & Evens, Matching pennies, Rock/paper/scissors

Static games of complete information

- Each player i in the game **simultaneously** and **independently** chooses an action from its own set of available actions A_i
- The combination of actions chosen by the n players determines the **outcome** of the game
- Outcome (a_1, a_2, \dots, a_n) determines a payoff for each player through an individual utility function of player i : $u_i = u_i(a_1, a_2, \dots, a_n)$
- 3 ingredients = actions + outcome + utility

Action versus strategy

- As will be seen later, it is useful to think of strategies instead of actions
- A strategy is a plan of action *a seconda di condizioni*
 - ▣ e.g.: if these conditions are met, then my action is a, otherwise is either a' or a''
 - ▣ this plan can even be random (we will see why)
- Right now, we just need certain plans
- These are called pure strategies, i.e.: a pure strategy is a deterministic plan of action

Normal-form representation

- Each player i simultaneously chooses a strategy from a set of pure strategies S_i
- This results in a given action chosen by each of the n players that ultimately determines a payoff for each player
- If any player i plays strategy $s_i \in S_i$, the combination of moves is $(s_1, s_2, \dots, s_i, \dots, s_n)$
- Player i gets payoff $u_i(s_1, s_2, \dots, s_i, \dots, s_n) \in \mathbb{R}$
- The **normal form** of the game is specified by $G = \{S_1, \dots, S_n; u_1, \dots, u_n\}$
 ⇒ di norma, si specificano i giocatori e si copisce già da n

Simultaneous and independent

- Simultaneous moves do not really need to happen at the same time
 - ▣ it is just that strategies are chosen without knowledge of everybody else's actions
- These two versions are both “simultaneous”
 - ▣ version A: two players are writing their strategy on opposite sides of a board at the same time
 - ▣ version B: player 1 is asked to write first, while he writes, player 2 is blindfolded; then the board is turned and player 2 writes

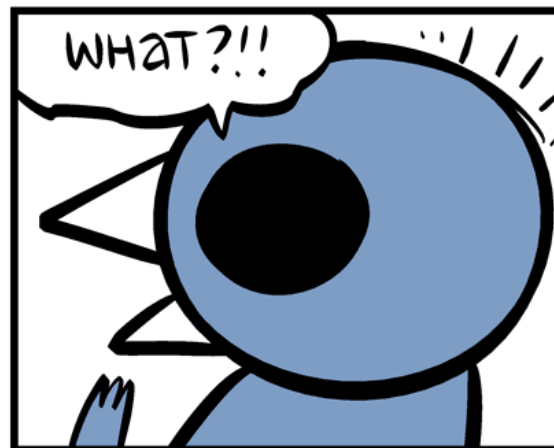
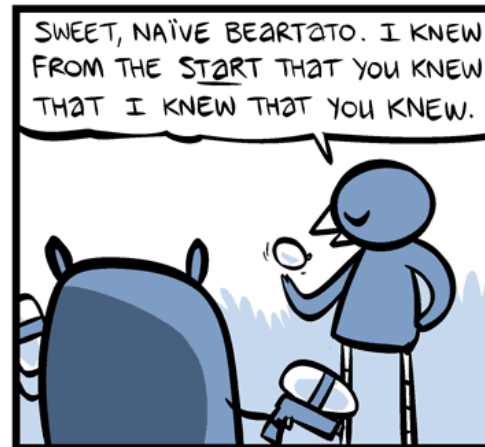
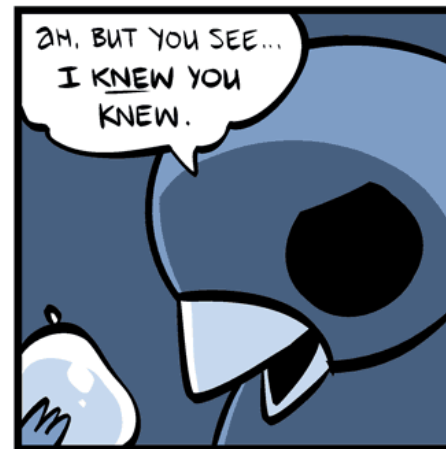
Common knowledge

- We say that E is common knowledge if:
 - everybody knows E
 - everybody knows that everybody knows E
 - and so on, ad infinitum

- This is a powerful but not obvious assumption
 - it requires not only full knowledge on information pertinent to myself, but also on what should be everybody else's concern

Common knowledge

- “Complete information” means that
 - ▣ all possible actions of all players
 - ▣ all possible outcomes resulting from these actions
 - ▣ the individual preferences of all players about these outcomes (i.e., their utilities about them)are common knowledge among the players
- Player rationality is common knowledge
 - ▣ which means that everybody is maximizing their own payoff and everybody **knows** that everybody is maximizing their payoff!



Matrix representation

- n -player games can be represented as functions in $S_1 \times S_2 \times \dots \times S_n$
- If S, s are discrete sets, a n -dimensional matrix can be used, where each cell contains again a n -dimensional value in \mathbb{R}^n
- Usually we have $n=2$, so the representation is simply an ordinary matrix where we put a pair of real numbers in each cell (therefore called a **bi-matrix**)

Example of bi-matrix

Player 2 strategies

Player 1 strategies

	$s_2^{(1)}$	$s_2^{(2)}$	$s_2^{(j)}$
$s_1^{(1)}$	$u_1(s_1^{(1)}, s_2^{(1)}),$ $u_2(s_1^{(1)}, s_2^{(1)})$	$u_1(s_1^{(1)}, s_2^{(2)}),$ $u_2(s_1^{(1)}, s_2^{(2)})$
$s_1^{(2)}$	$u_1(s_1^{(2)}, s_2^{(1)}),$ $u_2(s_1^{(2)}, s_2^{(1)})$
$s_1^{(i)}$	<div> $u_1(s_1^{(i)}, s_2^{(j)}),$ $u_2(s_1^{(i)}, s_2^{(j)})$ </div>

payoff of user 1

payoff of user 2

Example 1

- Player A has three strategies: {U, M, D}
- Player B has two strategies: {L, R}

		player B	
		L	R
player A	U	8, 0	0, 5
	M	1, 0	4, 3
	D	0, 7	2, 0

Example 2

- Player A has three strategies: {U, M, D}
- Player B has three strategies: {L, C, R}

		player B		
		L	C	R
player A	U	0, 5	4, 0	7, 3
	M	4, 0	0, 5	7, 3
	D	3, 7	3, 7	9, 9

Example 3: Odds & Evens

- Players Odd and Even bet 4 euros
- Player Odd has two strategies: $\{0, 1\}$
- Player Even has two strategies: $\{0, 1\}$

		Even	
		0	1
Odd	0	-4, 4	4, -4
	1	4, -4	-4, 4

- “Similar” to: matching pennies (head/tails), penalty kick (left/right), poker (bluff/bet)

Example 4: rock/paper/scissors

- Players A and B bet 4 euros
- Player A has two strategies: {R,P,S}
- Player B has two strategies: {R,P,S}

		player B		
		R	P	S
player A	R	0, 0	-4, 4	4, -4
	P	4, -4	0, 0	-4, 4
	S	-4, 4	4, -4	0, 0

Example 5: Battle of Sexes

- Ann and Brian are partners who agreed to meet at a movie theater not knowing that 2 movies are available: romance (R) or sci-fi (S)
- Main goal of both is to see the other, but Ann prefers movie R and Brian prefers S

		Brian	
		R	S
Ann	R	2, 1	0, 0
	S	0, 0	1, 2

Example 6: Prisoners' Dilemma

- Simple version: both Al and Bob can choose between (M) “lose 1\$” or (F) “let the other pay 20\$”

		Bob	
		M	F
Al	M	-1, -1	-21, 0
	F	0, -21	-20, -20

Example 6: Prisoners' Dilemma

- Original version: involves a theft Al and Bob committed together. Caught by police, they can choose between (M)um and (F)ink
- Their (negative) payoff is the number of months they will spend in jail

		Bob	
		M	F
Al	M	-1, -1	-9, 0
	F	0, -9	-6, -6

Pareto efficiency

- A joint strategy s is **Pareto dominated** by another joint strategy s' if

$$\begin{array}{ll} u_i(s') \geq u_i(s) & \text{for every player } i \\ u_i(s') > u_i(s) & \text{for some player } i \end{array}$$

no altra strategia
strettamente
migliore sotto
ogni aspetto
↓

- A joint strategy s not Pareto dominated by any joint strategy s' , is said to be **Pareto efficient**
- There may be more than one Pareto efficient strategy, none of which dominates the others

Strict dominance

a comparison of strategies

Strictly dominated strategy

- Consider game $G = \{S_1, \dots, S_n; u_1, \dots, u_n\}$
- If $s_i, s_i' \in S_i$, we say that s_i is **strictly dominated** by s_i' if i 's payoff when playing s_i' is greater than when playing s_i for any other move of the other players, i.e.

$$u_i(s_1, s_2, \dots, s_i', \dots, s_n) > u_i(s_1, s_2, \dots, s_i, \dots, s_n)$$

$$\forall (s_1, \dots, s_{i-1}, s_{i+1}, \dots, s_n) \in S_1 \times \dots \times S_{i-1} \times S_{i+1} \times \dots \times S_n$$

- Rational players do not play such strategies

back to Example 1

- Strategy D is strictly dominated by M

		player B	
		L	R
player A	U	6, 0	0, 5
	M	1, 0 ✓	4, 3 ✓
	D	0, 7	2, 0

*M sempre supera
dato altra mossa*

back to Example 1

- Other strictly dominated strategies are found:
first L for player B (dominated by R)
then U for player A (dominated by M)

		player B	
		L	R
player A	U	6, 0	0, 5
	M	1, 0	4, 3
	D	0, 7	2, 0

Annotations in the image: Arrows pointing from R to L for Player B indicate that R is strictly dominant. Arrows pointing from M to U for Player A indicate that M is strictly dominant.

rational players play (M, R), with result (4,3)

back to Example 6

- F dominates M for both Al and Bob.
- The only playable strategies give outcome = F,F
- This justifies the “Dilemma” name

		Bob	
		M	F
Al	M	-1, -1	-21, 0
	F	0, -21	-20, -20

- ▶ The result does not seem that efficient

razionalità rispettata, ma risultato rimorso non è Pareto efficiente

Solving problems via IESDS

- This procedure is called “iterated elimination of strictly dominated strategies” (IESDS)
- Sometimes can find the outcome of a game, and is useful to obtain a “smaller” game by relying on common knowledge
- However, in several cases, it does not provide any solution

back to Example 2

- Here, no dominated strategy can be found.
- However, (D,R) seems to be a good choice.

		player B		
		L	C	R
player A	U	0, 5	4, 0	7, 3
	M	4, 0	0, 5	7, 3
	D	3, 7	3, 7	9, 9

back to Example 3

- Neither 0 or 1 strictly dominates the other
- There seems not to be any “better” strategy

		Even	
		0	1
Odd	0	-4, 4	4, -4
	1	4, -4	-4, 4

back to Example 5

- Here, there are **two** strategies that seem to be “good” for rational players, (R,R) and (S,S)
- But again, no dominated strategy to eliminate

		Brian	
		R	S
Ann	R	2, 1	0, 0
	S	0, 0	1, 2