

RECAP

OCCUPANCY / BALLS-INTO-BINS

PROBLEM: $m \geq n$ identical balls
thrown randomly and independently
into n bins (boxes)

1) Maximum number of balls into any bin

Let $X_b = \text{"# balls into bin } b\text{"}$, $1 \leq b \leq n$
 $X_b \sim \text{Binomial}(m, \frac{1}{n})$
$$\Pr(X_b = i) = \binom{m}{i} \frac{1}{n^i} \left(1 - \frac{1}{n}\right)^{m-i}$$

$C_b(k) = \text{"bin } b \text{ has } \geq k \text{ balls"}$

$$\Pr(C_b(k)) = \sum_{i \geq k} \Pr(X_b = i) \leq 2 \left(\frac{e m}{k n} \right)^k$$

$m = n$

 $\Pr(C_b(k)) \leq \frac{1}{n^2} \text{ for } k = \Theta(\ln n / \ln \ln n)$

$$\Rightarrow \Pr(\exists b : C_b(k) \text{ holds}) \leq \frac{1}{n}$$

(union bound)

\Rightarrow Insertion/search in hash table with
chaining ($m = n$ slots/keys): $O(\log n / \log \log n)$
w.h.p.

Case 2: $m = n \ln n$

$$\text{Set } k^* = e^2 \ln n \left(> 2 \frac{e m}{n} \right)$$

$$\Pr(C_b(k^*)) \leq 2 \left(\frac{e m}{k^* n} \right)^{k^*}$$

$$< 2 \left(\frac{e \cancel{n} \cancel{\ln n}}{e^2 \cancel{n} \cancel{\ln n}} \right)^{e^2 \ln n} \leq \frac{2}{n^{e^2}} < \frac{1}{n^2}$$

Thus

$$\Pr(\exists \text{ bin with } \geq k^* = e^2 \ln n \text{ balls}) = \Pr\left(\bigcup_{b=1}^n C_b(k^*)\right) \leq \sum_{b=1}^n \Pr(C_b(k^*)) < \frac{n}{n^2} = \frac{1}{n}$$

Therefore, the maximum number of balls into a bin, when $m = n \ln n$, is $O(\log n) = O\left(\frac{m}{n}\right)$ w.h.p.

This means that balls are roughly balanced among the bins!

This phenomenon also holds for $m = \Omega(n \log n)$! (exercise)

FURTHER RESULTS

$m = n$ \exists bin b with $\Theta(\log n / \log \log n)$
balls w.h.p. (no proof)

$m = n \ln n$:

$\Pr(\exists \text{ bin } b \text{ with } \Theta(\log n) \text{ balls}) = 1$
(PIGEON HOLE)

APPLICATION: Scheduling $m \gg n$
($\Omega(n \log n)$) work units among
 n workers. A random allo-
cation guarantees weak
balancing: $\Theta(\frac{m}{n})$ units/worker whp!

2. Number of empty bins ($m=n$)

For a fixed bin b :

$$\Pr(X_b=0) = \left(1 - \frac{1}{n}\right)^n$$

Consider an indicator variable

$$Y_b = 1 \iff X_b = 0 \quad \Pr(Y_b=1) = \left(1 - \frac{1}{n}\right)^n$$

Then $Y = \sum_{b=1}^m Y_b = \# \text{ empty bins}$

The Y_b 's are not independent, but

$$E[Y] = \sum_{b=1}^m E[Y_b] = n \left(1 - \frac{1}{n}\right)^n$$

$$\text{For } n \geq 2: \frac{1}{4} \leq \left(1 - \frac{1}{n}\right)^n < e^{-1}$$

$$\left(\text{for large } n: \left(1 - \frac{1}{n}\right)^n \approx e^{-1}\right)$$

On average there are

$$\text{thru } \frac{n}{4} \leq E[Y] \leq \frac{n}{e}$$

empty bins!

MORALE: When $m=n$, randomly distributing work units results into a constant fraction of idle workers! \Rightarrow MANY CONFLICTS

Case $m \ll n$: BIRTHDAY PARADOX

- Class of $m = 30$ students
- What is the probability that two students have the same birthday?

(assume fully random birthdays)

e.g. no leap years
no twins
etc...

Experiment: throwing $m = 30$ balls (students) into $n = 365$ bins (birthdays):

evaluate $P = \Pr(\text{no bin with } > 1 \text{ balls})$

There are $\overset{n}{m} \binom{365}{30}$ choices of picking distinct birthdays, and $30!$ ways

to distribute the distinct birth days to the students.

The number of all possible birth day configurations is 365^{30}

Therefore

$$P = \frac{{}^n (365)_{30}}{365^{30}} = \frac{365 \cdot 364 \cdot \dots \cdot (365 - 29)}{365^{30}}$$

$$= 1 \cdot \left(1 - \frac{1}{365}\right) \cdot \left(1 - \frac{2}{365}\right) \cdot \dots \cdot \left(1 - \frac{29}{365}\right)$$

$$< 0.3!$$

In $> 70\%$ of the cases, there will be two students with the same birthday!

GENERALIZING:

$$\begin{aligned} P_{m,n} &= \prod_{j=1}^{m-1} (1 - j/n) \approx \prod_{j=1}^{m-1} e^{-j/n} = e^{-\sum_{j=1}^{m-1} j/n} \\ &= e^{-\frac{(m-1)m}{2n}} \approx e^{-\frac{m^2}{2n}} \end{aligned}$$

The probability stays $\geq \frac{1}{2}$ for m up to $\approx \sqrt{2 \ln 2 n} = \Theta(\sqrt{n})$.

It becomes vanishing $\left(\xrightarrow[n \rightarrow \infty]{} 0 \right)$ for $m = \Omega(\sqrt{n \log n})$