

15.09.2021

SHEET #01

wednesday.

Introduction to Computer Science.

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30005

Problem 1.1

$$\rightarrow E = \{ \}$$

$$A = \{ \{a\}, \{b\}, \{c\}, \{d\}, \{e\}, \{f\} \}$$

$$C = 0$$

$$\rightarrow E = \{ \{e, f\} \}$$

$$A = \{ \{a\}, \{b\}, \{c\}, \{d\}, \{e, f\} \}$$

$$C = 1$$

$$\rightarrow E = \{ \{e, f\}, \{d, f\} \}$$

$$A = \{ \{a\}, \{b\}, \{c\}, \{d, e, f\} \}$$

$$C = 1 + 2 = 3$$

$$\rightarrow E = \{ \{e, f\}, \{d, f\}, \{c, f\} \}$$

$$A = \{ \{a\}, \{b\}, \{c, d, e, f\} \}$$

$$C = 1 + 2 + 4 = 7$$

$$\rightarrow E = \{ \{e, f\}, \{d, f\}, \{c, f\}, \{b, f\} \}$$

$$A = \{ \{a\}, \{b, c, d, e, f\} \}$$

$$C = 1 + 2 + 4 + 6 = 13$$

Minimal Spanning
Tree.

$$\rightarrow E = \{ \{e, f\}, \{d, f\}, \{c, f\}, \{a, f\} \}$$

$$A = \{ a, b, c, d, e, f \}$$

$$C = 1 + 2 + 4 + 6 + 8 = \underline{21} \text{ Total Minimum Cost}$$

t = FFLFLFRFRFFLFRF p = FFLFR

```
F F L F L F R F R F F L F R F
F F L F R
  F F l f r
    F f l f r
      F F l f r
        F f l f r
          F F l f r
            F f l f r
              F F l f r
                F f l f r
                  F F L F R
```

alignments used: 10 comparisons made: 22

b.

t = FFLFLFRFRFFLFRF p = FFLFR

F	F	L	F	L	F	R	F	R	F	F	L	F	R	F	skip
f	f	l	f	R											1
		F	F	L	F	R									0
			f	f	l	f	R								0
						f	f	l	f	R					2
							f	f	L	F	R				1
									F	F	L	F	R		

alignments used: 6 comparisons made: 16

c. lookup table

	F	F	L	F	R
	0	1	2	3	4

F	-	-	0	-	0
L	0	1	-	0	1
R	0	1	2	3	-
P	0	1	2	3	4

Problem 1.3

a) $\rightarrow f_8(n) = \log \log n$

$\rightarrow f_6(n) = 2 \log n$

$\rightarrow f_1(n) = \frac{1}{2} n \log n$

$\rightarrow f_3(n) = \sqrt{n^3}$

$\rightarrow f_2(n) = n^2$

$\rightarrow f_5(n) = 100n^2 + 10n^3$

$\rightarrow f_7(n) = (n^2)^2$

$\rightarrow f_4(n) = n^n$

b) $f \text{ is } o(g), g \text{ is } O(h) \rightarrow f \text{ is } o(h)$

PF:

- $f(n) \leq c g(n) \mid c > 0 \ n \geq 0 \ \forall n \geq n_0$

- $g(n) \leq c' h(n) \mid c' > 0 \ n \geq 0 \ \forall n \geq n_1$

if it follows

$$f(n) \leq c g(n) \leq c c' h(n)$$

$$f(n) \leq c c' h(n) \quad \forall n \geq \max(n_0, n_1)$$