CH-231-A Algorithms and Data Structures ADS

Lecture 22

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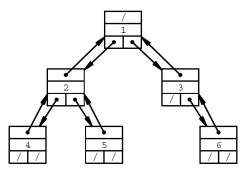
Spring 2022

Representing Rooted Trees

- ► Traversing a rooted tree requires us to know about the hierarchical relationships of their nodes.
- Similar to linked list implementations, such relationships can be stored by using pointers.

Binary Tree

- ▶ Binary trees *T* have an attribute *T.root*.
- ► They consist of nodes x with attributes x.parent (short x.p), x.left, and x.right in addition to x.key.



d-ary Trees

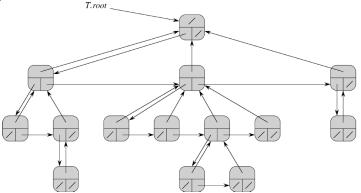
- ightharpoonup d-ary trees are rooted trees with at most d children per node.
- ▶ They can be handled analogously to binary trees.

```
struct node {
    int val;
    node* parent;
    node* child[d];
};

typedef node* tree;
```

Rooted Trees with Arbitrary Branching

Rooted trees T with arbitrary branching consist of nodes x with attributes x.p, x.leftmost-child, and x.right-sibling in addition to x.key.



Discussion

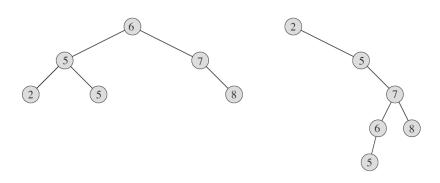
- ▶ Representing trees with pointers allows for a simple and intuitive representation.
- It also allows for a dynamic data management.
- Modifying operations can be implemented efficiently.
- ► However, extra memory requirements exist for storing the pointers.

Binary Search Tree: Definition

- ➤ A binary search tree (BST) is a binary tree with the following property:
 - Let x be a node of the BST.
 - ▶ If y is a node in the left subtree of x, then $y.key \le x.key$.
 - ▶ If y is a node in the right subtree of x, then $x.key \le y.key$.
- ▶ The idea of a BST data structure is to support efficient dynamic set operations, many in O(h), where h is the tree's height.

Binary Search Trees

Binary Search Tree: Examples



Query: In Order Visit

▶ Visit all nodes in order and execute an operation:

```
Function DFS-Inorder-Visit(Node n)

1 if n = N/L then return;
2 DFS-Inorder-Visit(n.left);
3 n.Operation();
4 DFS-Inorder-Visit(n.right);
```

- ▶ The operation could, e.g., be printing the key.
- This tree traversal is also referred to as in-order tree walk.
- Time complexity (n = number of nodes): O(nk) when assuming that the operation is in O(k).

Query: Searching

► Recursive tree search:

```
TREE-SEARCH(x, k)

1 if x == NIL or k == x.key

2 return x

3 if k < x.key

4 return TREE-SEARCH(x.left, k)

5 else return TREE-SEARCH(x.right, k)
```

Iterative tree search:

```
ITERATIVE TREE-SEARCH(x,k)

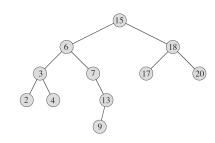
1 while x \neq \text{NIL} and k \neq x.key

2 if k < x.key

3 x = x.left

4 else x = x.right

5 return x
```



Time complexity: O(h)