Relational Algebra

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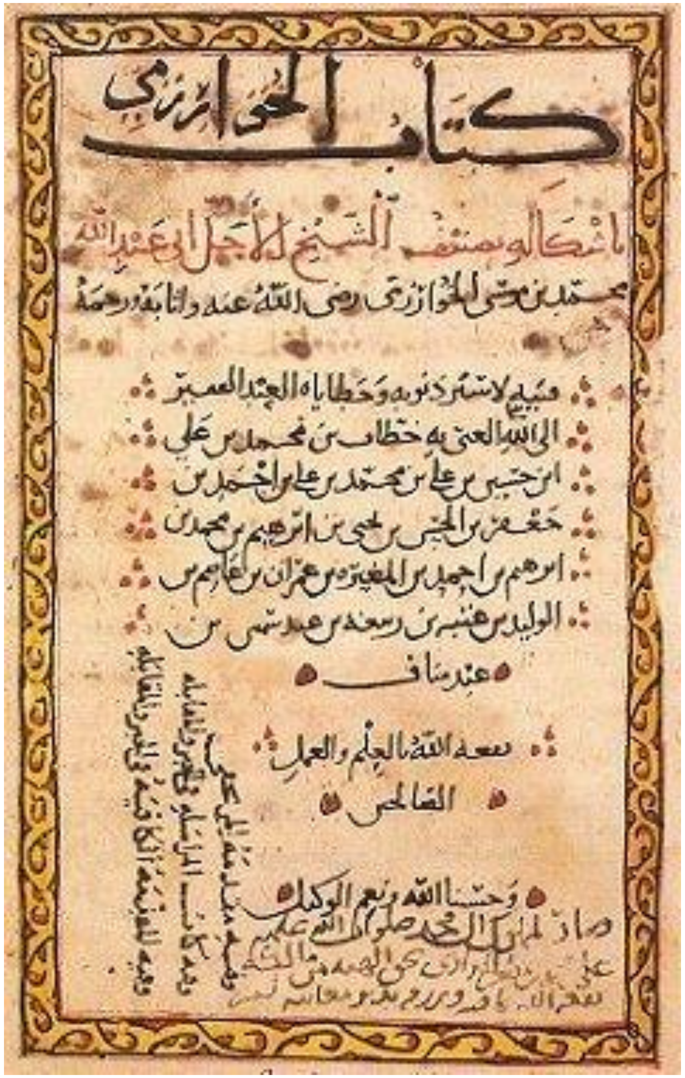
Database Management: Complete Book,

Chapters 2 & 5

Algebra

- 2 "fathers of algebra“:
  - where algebra ≡ theory of equations
    → Greek *Diophantus*
  - where algebra ≡ rules for manipulating & solving equations
    → Persian *al-Khwarizmi*
- Source: Wikipedia

Xorazm,
Usbekistan



What is "Algebra"?

- Mathematical system consisting of:
  - Operands - variables or values from which new values can be constructed
  - Operators - symbols denoting procedures that construct new values from given values
  - Ex: ((x + 7)/(2 - 3)) + x
- Algebra A = (C,OP)
  -- "simplest" mathematical structure:
  - C nonempty carrier set (=value set)
  - OP nonempty operation set
  - C closed under OP expressions

Selection

- R1 := σ_C (R2)
  - *C* : condition on attributes of R2.
  - R1 is all those tuples of R2 that satisfy *C*.

```
sid   name  login       gpa
-----------------------------
53666 Jones jones@cs    3.4
53688 Smith smith@eecs  3.2
53650 Smith smith@math  3.8
```

$\sigma_{gpa<3.8}$(Students):

```
sid   name  login       gpa
-----------------------------
53666 Jones jones@cs    3.4
53688 Smith smith@eecs  3.2
```

Selection: Observations

- unary operation: 1 table
- conditions apply to each tuple individually
  - condition cannot span tuples (how to do that?)
- degree of $\sigma_C(R)$ = degree of R
  - Cardinality?
- Select is commutative: $\sigma_{C1}(\sigma_{C2}(R)) = \sigma_{C2}(\sigma_{C1}(R))$

Projection

- R1 := π_{attr}(R2)
  - *attr* : list of attributes from R2 schema
- For each tuple of R2:
  - extract attributes from list attr in order specified (!) → R1 tuple
- Eliminate duplicate tuples

sid	name	login	gpa
53666	Jones	jones@cs	3.4
53688	Smith	smith@eecs	3.2
53650	Smith	smith@math	3.8

$\pi_{name,login}$(Students) =

name	login
Jones	jones@cs
Smith	smith@eecs

Projection: Observations

- Unary operation: 1 table
- removes duplicates in result
  - Cardinality?
  - Degree?
- Project is not commutative
- Sample algebraic law: $\pi_{L1} (\pi_{L2}(R)) = \pi_{L1}(R)$ if $L1 \subseteq L2$
  - else incorrect expression, syntax error

Exercises

- $\pi_{Name,login}(\sigma_{gpa=3.8}(Students))$ = ?

```
sid   name  login      gpa
----------------------------
53666 Jones jones@cs   3.4
53688 Smith smith@eecs 3.2
53650 Smith smith@math 3.8
```

- "name and rating for sailors with rating > 8"
  - Note explicit operation sequence!

Cartesian Product

- project, select operators operate on single relation
- Cartesian product combines two:
  R3 = R1 x R2
  - Pair each tuple t1 $\in$ R1 with each tuple t2 $\in$ R2
  - Concatenation t1,t2 is a tuple of R3
  - Schema of R3 = attributes of R1 and then R2, in order
  - beware attribute *A* of the same name in R1 and R2: use R1.*A* and R2.*A*

Cross Product ("Cartesian Product")

- Example U := R x S

A	B
1	2
3	4

(a) Relation R

B	C	D
2	5	6
4	7	8
9	10	11

(b) Relation S

A	$R.B$	$S.B$	C	D
1	2	2	5	6
1	2	4	7	8
1	2	9	10	11
3	4	2	5	6
3	4	4	7	8
3	4	9	10	11

(c) Result $R \times S$

Natural Join

- T = R ⋈ S
  - Ex: Reserves $\bowtie_{bid}$ Sailors
- connect two relations:
  - Equate attributes of same name, project out redundant attribute(s)

A	B
1	2
3	4

(a) Relation R

B	C	D
2	5	6
4	7	8
9	10	11

(b) Relation S

A	R.B	S.B	C	D
1	2	2	5	6
1	2	4	7	8
1	2	9	10	11
3	4	2	5	6
3	4	4	7	8
3	4	9	10	11

(c) Result $R \times S$

A	B	C	D
1	2	5	6
3	4	7	8

R ⋈ S

Generalizing Join

- $T = R \bowtie_C S$
  - First build R x S, then apply σ_C
- Generalization of equi-join: A θ B where θ one of =, <, ...
  - Today, more general: σ_C can be any predicate
- Common join types:
  - Left join, right join, natural join, self join, ...

Relational Algebra: Summary

- Mathematical definition of relations + operators
  - Query = Algebraic expression
- Relational algebra $A = (R, OP)$ with relation $R = A_1 \times \ldots \times A_n$, $OP = \{\pi, \sigma, \times\}$
  - Projection: $\pi_{attr}(R) = \{ r.attr \mid r \in R \}$
  - Selection: $\sigma_p(R) = \{ r \mid r \in R, p(r) \}$
  - Cross product: $R_1 \times R_2 = \{(r_{11}, r_{12}, \ldots, r_{21}, r_{22}, \ldots) \mid (r_{11}, r_{12}, \ldots) \in R_1, (r_{21}, r_{22}, \ldots) \in R_2 \}$
  - Further: set operations, join, ...

Relational Calculus

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- Tuple variable = variable over some relation schema
- Query Q = $\{ T \mid T \in R, p(T) \}$
  - R relation schema, p(T) predicate over T
- Example 1: "sailors with rating above 8"
  - Sailors = sid:int $\times$ sname:string $\times$ rating:int $\times$ age:float
  - = $\{ S \mid S \in \text{Sailors} \wedge S.\text{rating} > 8 \}$
- Example 2: "names of sailors who have reserved boat #103":
  - Reserves = sid:int $\times$ bid:int $\times$ day:date
  - = $\{ S.\text{sname} \mid \exists S \in \text{Sailors}\ \exists R \in \text{Reserves}: R.\text{sid}=S.\text{sid} \wedge R.\text{bid}=103 \}$

Comparison of Relational Math

- Relational algebra
  - set-based formalization of selection, projection, cross product (no aggregation!)
  - Operation oriented = procedural = imperative; therefore basis of optimization
- Relational calculus
  - Same, but in predicate logic
  - Describing result = declarative; therefore basis of SQL semantics
- Equally powerful
  - proven by Codd in 1970