

Operating Systems

Threads and Concurrency

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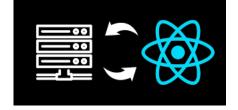
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Fall 2022

Motivation

- Most modern applications are multithreaded
- Multiple tasks with the application can be implemented by threads
 - Update display
 - Fetch data
 - Spell checking



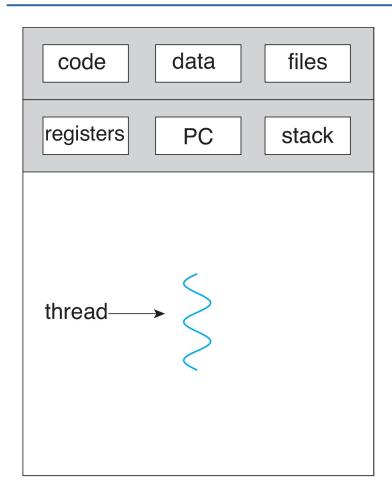




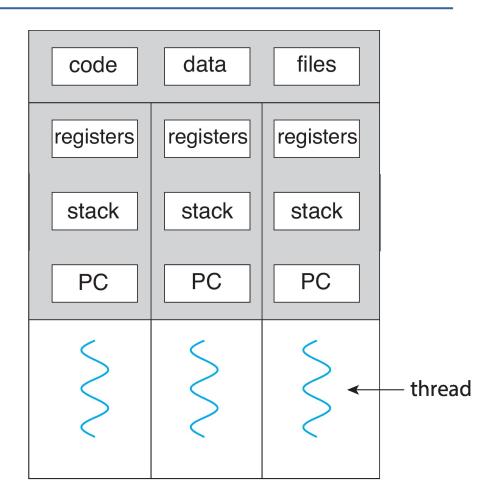
- Process creation is heavy-weight while thread creation is light-weight
- Kernels are generally multithreaded



Single and Multithreaded Processes



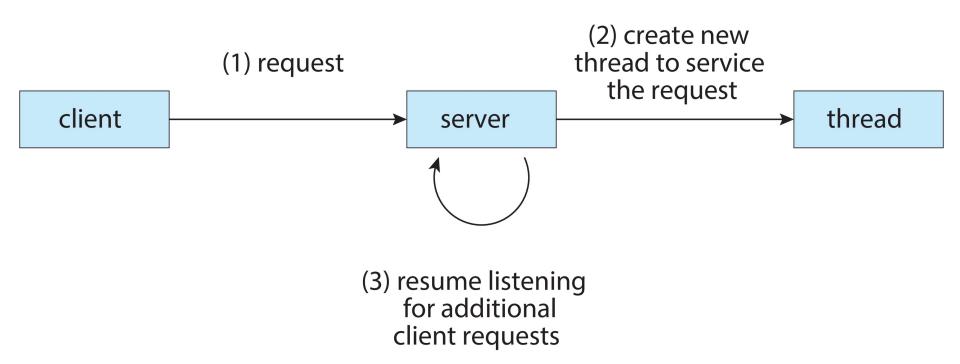
single-threaded process



multithreaded process



Multithreaded Server Architecture





Benefits

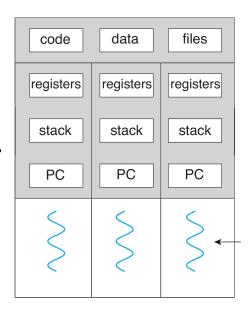
Responsiveness

- Allow continued execution if part of process is blocked
- Especially important for user interfaces



Resource Sharing

- Threads share resources of process
- Easier than shared memory or message passing.





Benefits (cont.)

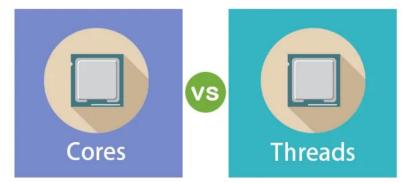
Economy

- Cheaper than process creation
- Thread switching lower overhead than context switching.



Scalability

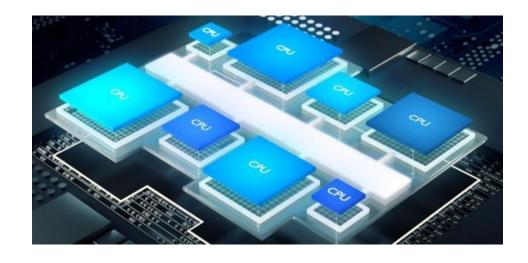
Process can take advantage of multicore architectures.





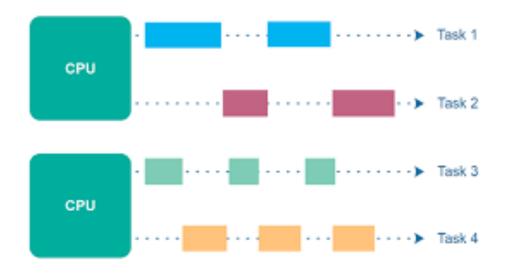
Multicore Programming

- Multicore or multiprocessor systems putting pressure on programmers, challenges include:
 - Dividing activities
 - Balance: ensuring that the tasks perform equal work of equal value.
 - Data splitting
 - Data dependency
 - Testing and debugging



Multicore Programming (cont.)

Parallelism implies performing more than one task simultaneously

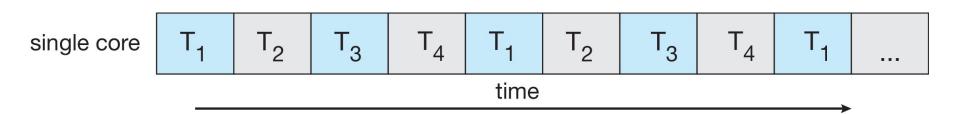


- Concurrency supports more than one task making progress
 - Single processor or core, scheduler providing concurrency

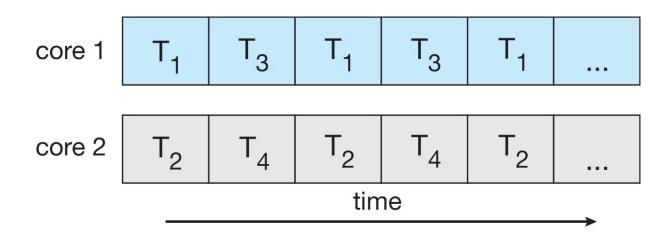


Concurrency vs. Parallelism

Concurrent execution on single-core system:



Parallelism on a multi-core system:



Multicore Programming-Types of Parallelism

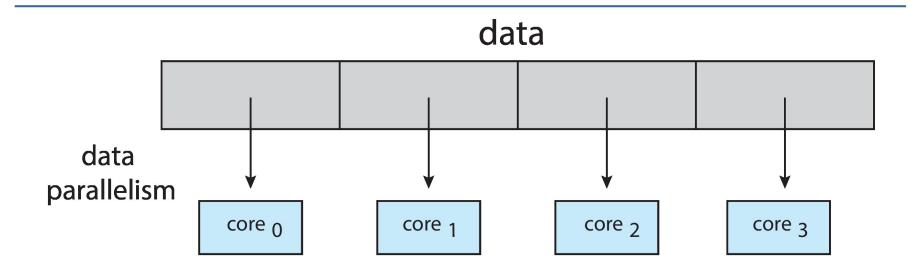
Data parallelism

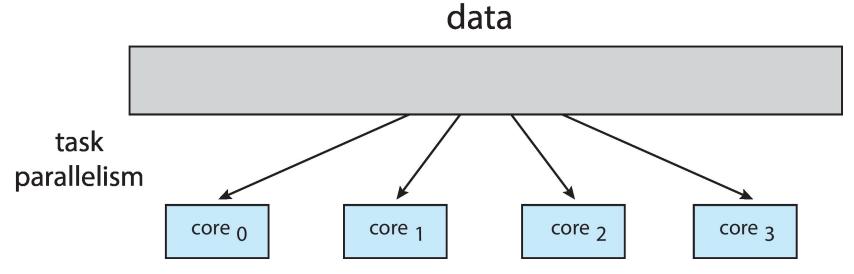
- Distributes subsets of the same data across multiple cores, same operation on each
- Example: summing the contents of an array of size N.

Task parallelism

- Distributing threads across cores, each thread performing unique operation
- Example: Unique statistical operation on the array of elements.

Data and Task Parallelism





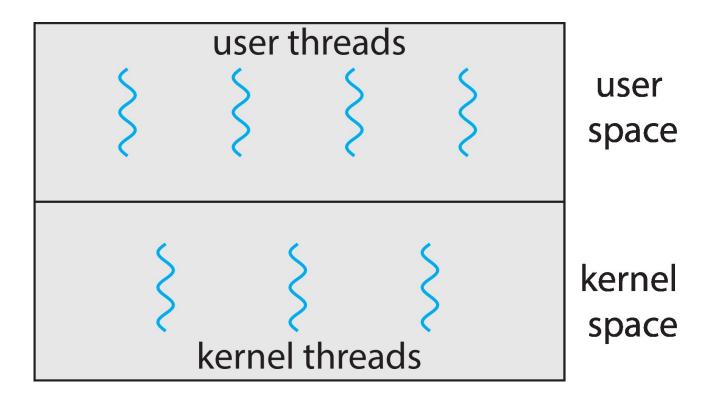


User Threads and Kernel Threads

- User threads: management done by user-level threads library.
- Three primary thread libraries:
 - POSIX Pthreads
 - Windows threads
 - Java threads
- Kernel threads: supported by the Kernel
 - Examples virtually all general -purpose operating systems,
 including: Windows, Linux, Mac OS X, iOS, Android



User and Kernel Threads



Additional review: https://www.geeksforgeeks.org/difference-between-user-level-thread-and-kernel-level-thread/

Multithreading Models

How to map user threads to kernel threads?

- Many-to-One
- One-to-One
- Many-to-Many

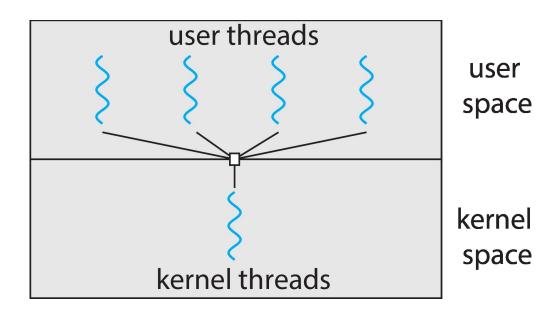
Additional review:

https://stackoverflow.com/questions/14791278/threads-why-must-all-user-threads-be-mapped-to-a-kernel-thread



Many-to-One

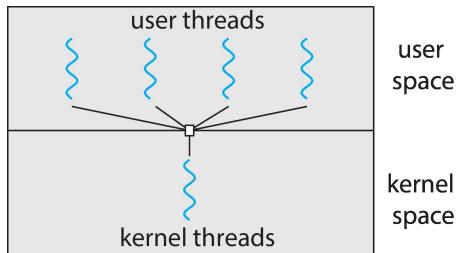
- Many user-level threads mapped to single kernel thread
- Thread management is done by the thread library in user space
 - So it is efficient





Many-to-One (cont.)

- One thread blocking causes all to block
- Multiple threads may not run in parallel on multicore system
 - Because only one may be in kernel at a time
- Few systems currently use this model
- **Examples:**
 - **Solaris Green Threads**
 - **GNU Portable Threads**

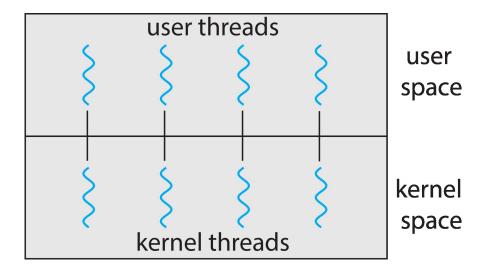


kernel



One-to-One

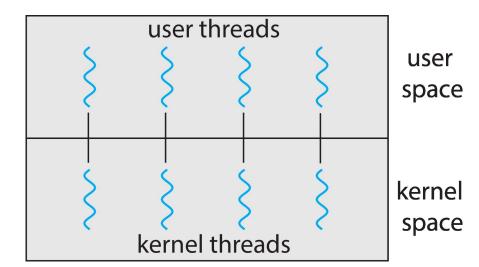
- Each user-level thread maps to kernel thread
- Creating a user-level thread creates a kernel thread





One-to-One (cont.)

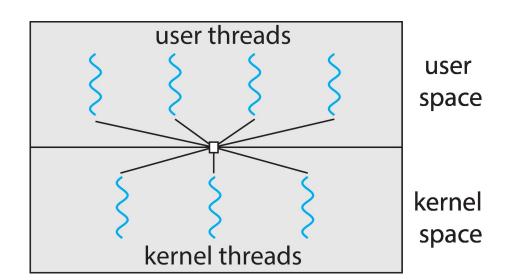
- More concurrency than many-to-one
- Number of threads per process may be restricted due to overhead
- Examples
 - Windows
 - Linux





Many-to-Many Model

- Many user level threads to be mapped to many kernel threads
- Operating system can create a sufficient number of kernel threads



- Examples
 - Windows with the ThreadFiber package
 - Otherwise not very common



Two-level Model

 Similar to M:M, except that it allows a user thread to be bound to kernel thread

