Gathering Atmospheric Data

Using an Unmanned Air Vehicle

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Abstract

This report looks at three dimensional energy based path planning for unmanned air vehicles in a predetermined area, with particular consideration to quality of data produced.

Contents

1 Introduction

The following section outline the process in which a minimum cost route through a sample space can be obtained that provides the best data collection quality

2 Plane Properties

Name	Wingspan	height
Plane 1	20m	5m
Plane 3	40m	6m

Table 1: Table of Plane Properties

To considere th minimum cost of circumnavigating a particular route the specifications of a plane must be considered in table 1 the properties of differnt planes is shown

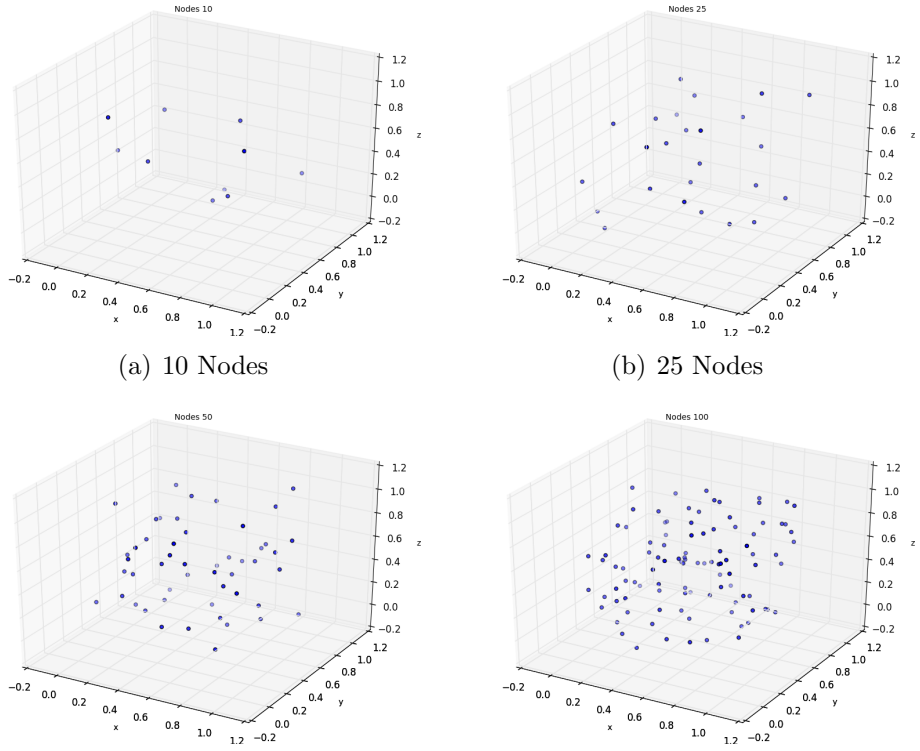
3 Energy Model

$$E = \alpha D + \beta H \tag{1}$$

From these plane properties the following energy model has been definened in equation 1 where α and β are coeficients that are determined by the plane. For the current plane shown in table 1 α and β take values of 10 and 6 respectively.

4 Latin Hypercubes

Latin hypercubes are sampling pland that provide the best space fillingness while limiting the total number of sampling points required. This is generally applied to testing of computer similations where the collection of each point is expensive. In this situation however the travel bertween the points the expensive component.



(a) 10 Nodes (b) 25 Nodes

(c) 50 Nodes (d) 100 Nodes

Figure 1: Latin Hypercubes with Varying Numbers of Nodes

Figure 1 shows a number of latin hypercubes with differnt numbers of nodes. All the Latin Hypercubes are within a unit cube. For collection of data in a required area these cubes can be stretched to fill the desired space. This does not provide an even spacing in each direction however means that eeach vertex of data collection is equally considered.

For this project the idea is to follow this logic to utilise Latin Hypercubes:

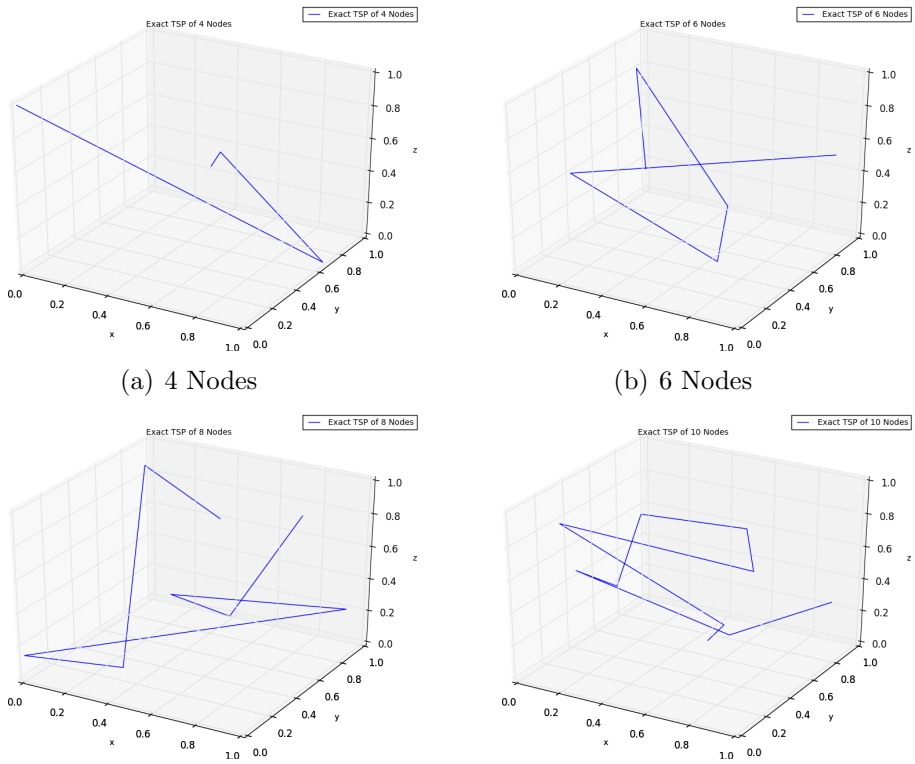
1. Specifiy area of interest to researcher
2. Estimate number of nodes able to be circumnavigated given the UAV total energy and the area of sample space
3. Fit Latin Hypercube of given nodes to sample area
4. Calcualte least energy route through the sample space
5. After first flight asses areas of encertainty to plan route through for next flight

5 Route Planning

Given a set of points within a sample space the next stage of the proceedings is to compute the least cost route through these points. This problem presents itself

in the forn of the travelling salesman problem. The travelling salesman problem is the problem of finding the least cost route through a set of points. There is lots of work done on the elicidean travelling salesman problem and introducing heroustics to improe the time taken to compute. This is due to the problem being an NP hard problem (the computing time required increases exponentially with the number of points in the route)

5.1 Exact Travelling Salesman



(a) 4 Nodes (b) 6 Nodes

(c) 8 Nodes (d) 10 Nodes

Figure 2: Exact routes calculated by travelling salesman

Figure 4 shows the optimal routes for differnt numbers of nodes. These optimal routes are found by computing the exact cost of each and every route option. Although this yield the shortest routes this approach is not efficient in terms of the computation time required.

Number of points	4	6	8	10
Number of possible routes	24	720	40320	3628800
Computation time (ms)	0.0	5.0	307.0	32927.0
Best route cost (J)	110.38	106.71	106.72	108.76

Table 2: Comparison of route calculation

Table 3 shows the number of possible routes and the resulting computation time given different numbers of nodes. It can easily be seen that the number of route options is

equivilent to $n!$ where n is the number of nodes. The number of routes directly realtes to the computation time.

The computation time of the exact TSP approach can be reduced in a number of ways. Primarily the start node of the calculation can be defined however this only reduces the complexity to the equivilent of removoing a node form the computation. Other approaches involve producing a best guess and improving upon that. The approach taken in this project is taken from considering the best routes in figure 4 generally comprise of a climbing component and a decending component.

5.2 Improving Travelling Salesman

The progressive travelling salesman is an approach to computing a best guess route for the least energy route for a number of points. The computation process is as follows

1. Order the nodes by there vertical location from lowest to highest
2. Considering the lowest N nodes compute all posible combinations to produce two routes from the given nodes
3. Compare all permutations of routes to return the least cost combination
4. Add first nodes of both routes to final route
5. Consider lowest N nodes that are not in the final route and reitterate
6. Work through all nodes until no nodes are left without a route

Nodes in each route	2	3	4	5
Computation time (ms)	2.0	3.0	52.0	4127.0
Best route cost (J)	120.49	115.48	115.48	114.79

Table 3: Comparison of route calculation

Table 3 shows the varying computation time for routes through a 10 node latin hypercube that have a differnet numbers of nodes in each route. This refers to the number of nodes that are consideered in each progressive itteration of the code.

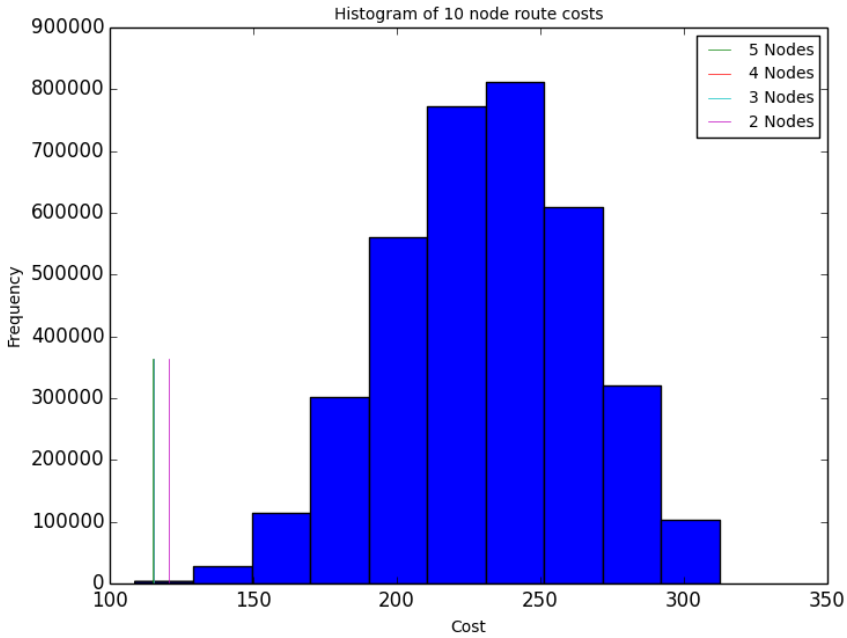
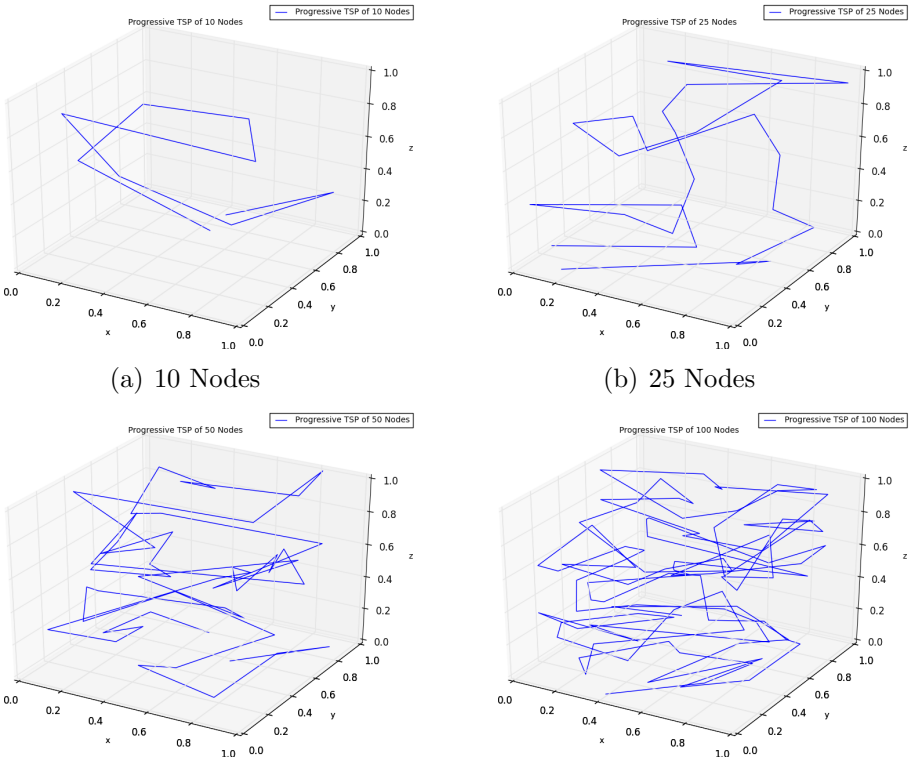


Figure 3: Histogram of 10 node route costs

Figure 3 shows a histogram of differnt route costs for a 10 node latin hypercube. The lines on this histogram plot represent the best cost routes with different numbers of nodes in the route

5.3 Progressive Travelling Salesman

In section 5.2 a best guess approach to the travelling salesman was presented and compared with the exact results. Given the results were close to the optimum even when the progressive algorithem only considered 2 nodes in the route it can be brought forward to be considered in reference to larger route problems. In these routes 4 nodes are considered for both the up route and the down route are used at each stage of the progressive algorithm.



(a) 10 Nodes

(b) 25 Nodes

(c) 50 Nodes

(d) 100 Nodes

Figure 4: Exact routes calculated by travelling salesman

Figure 4 shows a number of optimal routes for varing numbers of nodes whos order is defined be the progressive travelling salesman algorithm. Though the 100 node route is difficult to see the exact routing the other routes show a logical approach to the routing problem

Number of points	10	25	50	100
Computation time (ms)	52.0	235.0	593.0	1332.0
Best route cost (J)	115.48	143.62	202.65	333.73

Table 4: Comparison of progressive route calculation

Table 4 shows the computation time and route cost for the differnt routing situations. The computation time is far bellow that of the exact travelling salesman algorithm due to the complexity of this problem being broken down into 4 node chunks