

Roll No. 21L-5421Section 3F

National University of Computer and Emerging Sciences, Lahore Campus



Course: Data Structures
 Program: BS(CS)
 Duration: 10 Minutes
 Paper Date: 5 Sep 2021
 Section: F

Course Code: CS 201
 Semester: Fall 2022
 Total Marks: 10
 Exam: Quiz 1

Instruction/Notes: Solve the exam on this question paper.

Question: Consider the following program

i) Give an estimate of $T(N)$. (Show your work and give a $T(N)$ estimate for each line of code.)

⑩
 Good!

	$T(N)$ for each line
void Func(){	
1) \checkmark int count = 0;	1 \checkmark
for(int i = 1; i < n; i=i+2){	$i=1 - 1 \checkmark$
for(int k = i; k > 0; k=k/2)	$i < n - \log n/2 + 1 \checkmark$
count=count++;	$i = i + 2 - n/2 \checkmark$
}	
2) \checkmark 1 \leftarrow int mul3=0;	$\text{int } k = i - n/2 \checkmark$
\checkmark $n+1 \leftarrow$ for(int j=0; j < n; j++)	$k > 0 - (\log_2 n + 1) n/2 \checkmark$
\checkmark $n \leftarrow$ if(j % 3 == 0)	$k = k/2 - (\log_2 n) n/2 \checkmark$
\checkmark $(n+1) \Delta i^2 \leftarrow$ for(int m = 1; m < j*j; m++)	$\log_2 n (n/2) \checkmark$
\checkmark Δi^2 mul3++;	
}	

ii) Find the tight big Oh for the Best-case and Worst-case scenario. Explain in one line how you drive it.ii) Best Case ~~$O(n)$~~ $O(n)$ \checkmark Worst Case $O \Delta i^2 \Rightarrow \frac{n(n+1)(2n+1)}{6} = n^3 \checkmark$

Best Case would still require a loop to check all values till n if they are divisible by 3 however worst Case would be that every number is divisible of 3. \checkmark

i) Best Case $O(\log_2 n/2)$

Rough

i	k	# of lines
1	1	1
3	3	1
5	5, 2.5	2
	1.25	
i	k	# of
1	1, 0.5	
	0.25	
m	$m \leq j \times j$	m++
1	$m \leq 0$	
j 1	m	m++
j 2	m	1, 2, 3, 4
j 3	m	1-9



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i) Give an estimate of $T(N)$. (Show your work and give a $T(N)$ estimate for each line of code.)

<pre>int val = 0; for(int i = 1; i < n; i=i+2){ for(int j = i; j > 0; j=j/3) val++; } int no=0, k=0; while (k < n/2){ if(k % 2 == 0) for(int j = 1; j < k*k; j++) no=no+j; k=k+2; }</pre>	<p>$T(N)$ for each line</p> <p>$i=1 \rightarrow 1, i < n \rightarrow \frac{n}{2}, i=i+2 \rightarrow \frac{n}{2}$ $j=i \rightarrow \frac{n}{2}, j > 0 \rightarrow \frac{n}{2} \log_3 i$</p> <p>Basically $\sum_1^n \log_3 i$ and $i = \text{odd}$ $\log_3 1 + \log_3 3 + \dots \log_3 (n)$ $= n \log_3 n$</p> <p>$\frac{n}{4} k^2$ $\therefore \sum_0^{\frac{n}{2}} i^2$ when $i = \text{even}$ Sum of squares of first $\frac{n}{2}$ even numbers $= n^3$</p>
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ii) Find the tight big Oh for the Best-case and Worst-case scenario. Explain in one line how you drive it.

Best Case : $\frac{n}{4} = O(n)$ Inner loop does not run
 $k \% 2 \neq 0$

Worst Case : n^3 Inner loop does run