

STRONGLY CONNECTED COMPONENT



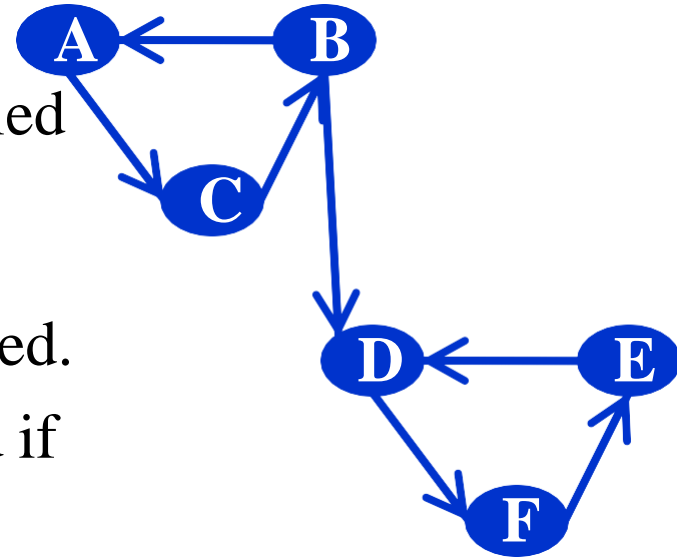
CONNECTIVITY

- Connected Graph

- An **undirected graph** $G(V, E)$ is called connected, if G contains a **path between every pair of vertices**

Otherwise, they are called disconnected.

- A **directed graph** is called connected if every pair of distinct vertices in the graph is connected.



- Connected Components

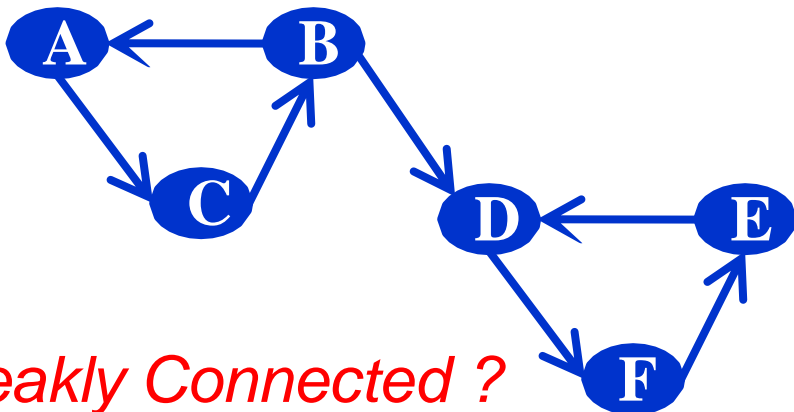
- A connected component is a **maximal connected subgraph** of G . Each vertex belongs to exactly one connected component, as does each edge.

Connected ?

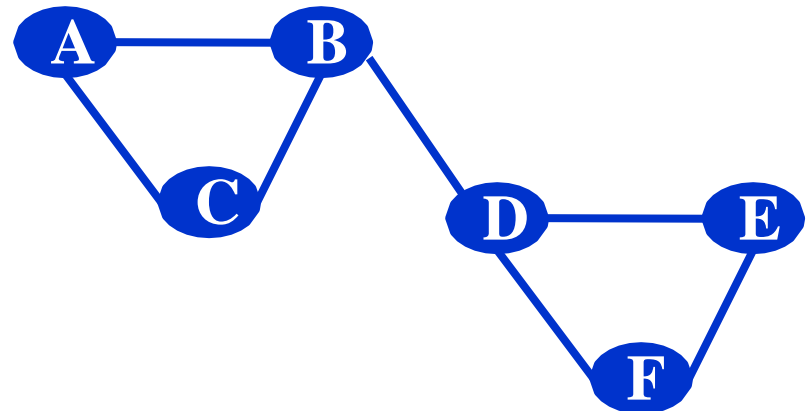


CONNECTIVITY (CONT.)

- Weakly Connected Graph
 - A directed graph is called **weakly connected** if **replacing** all of its directed edges with **undirected edges** produces a connected (undirected) graph.



Weakly Connected ?

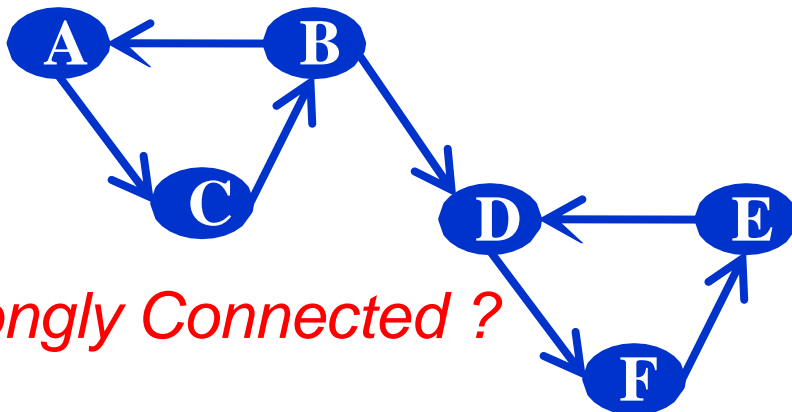


CONNECTIVITY (CONT.)

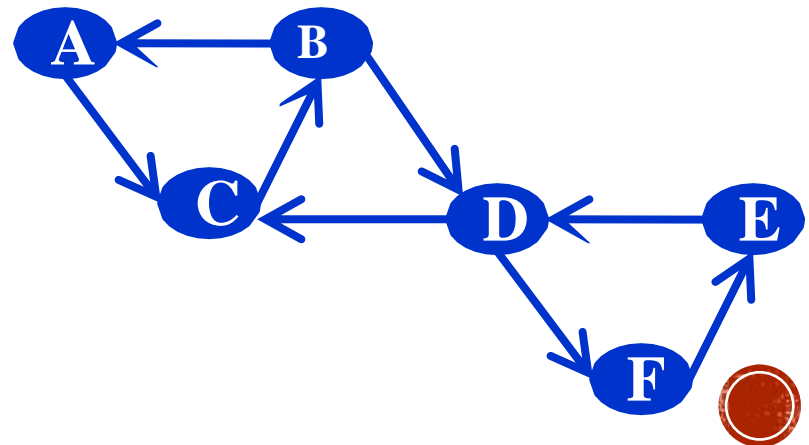
- Strongly Connected Graph

- It is strongly connected or strong if it contains a directed path from u to v for every pair of vertices u, v . The strong components are the maximal strongly connected subgraphs

Strongly Connected ?



Strongly Connected ?



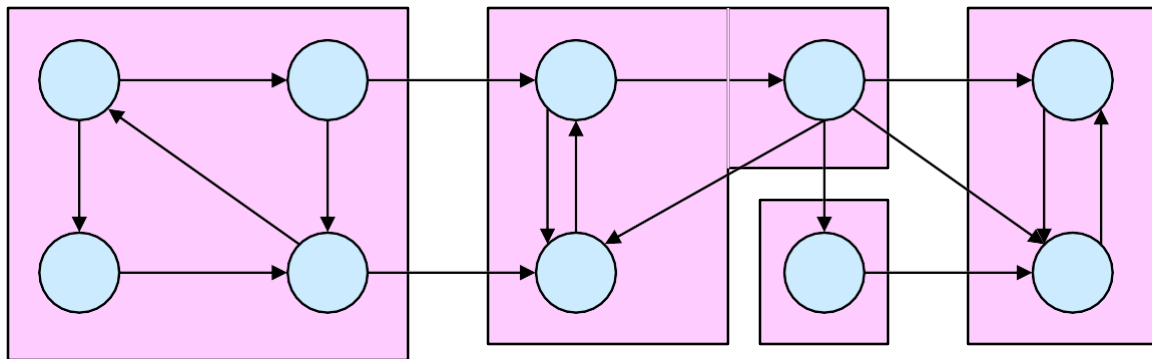
CONNECTED COMPONENTS

- Strongly Connected Components (SCC)
 - The **strongly connected components (SCC)** of a directed graph are its maximal strongly connected subgraphs.
- Here, we work with
 - Directed unweighted graph

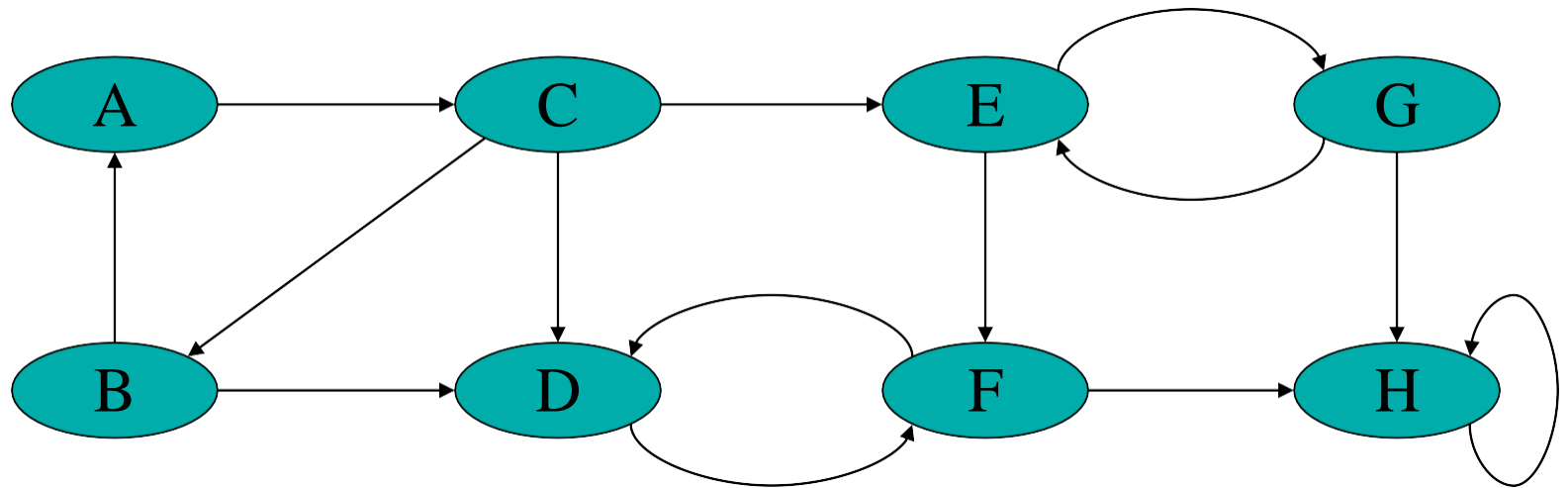


STRONGLY CONNECTED COMPONENTS

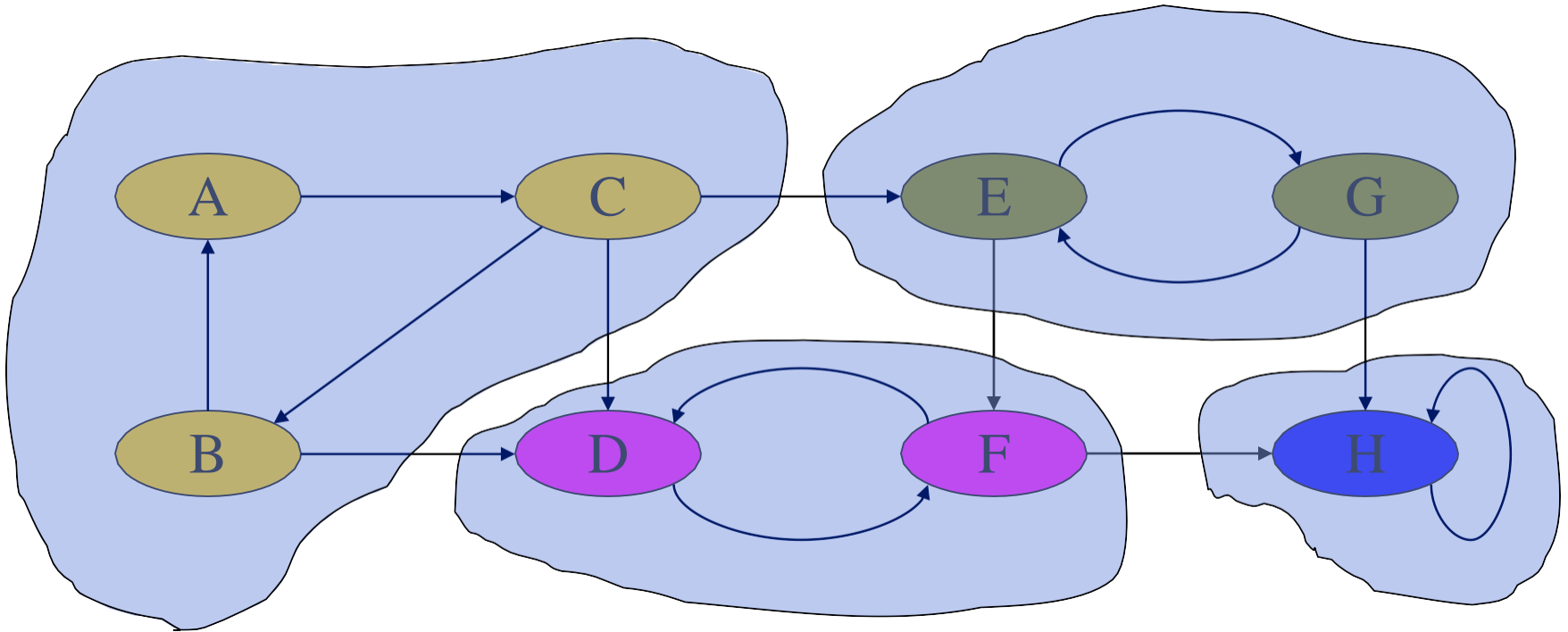
- G is strongly connected if every pair (u, v) of vertices in G is reachable from one another.
- A **strongly connected component** (**SCC**) of G is a maximal set of vertices $C \subseteq V$ such that for all $u, v \in C$, both $u \rightsquigarrow v$ and $v \rightsquigarrow u$ exist.



DFS - STRONGLY CONNECTED COMPONENTS

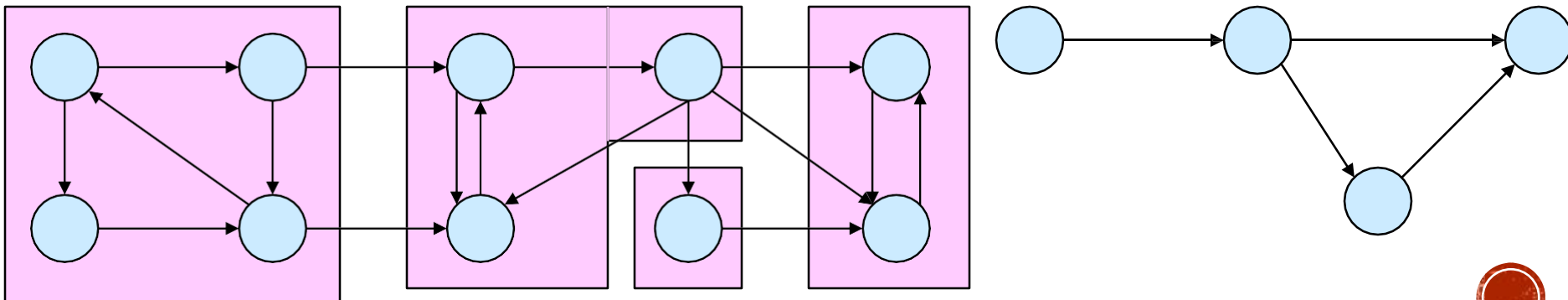


DFS - STRONGLY CONNECTED COMPONENTS



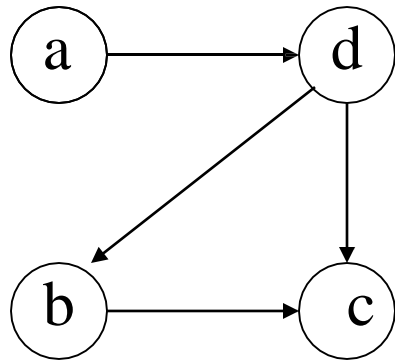
COMPONENT GRAPH

- $G^{\text{SCC}} = (V^{\text{SCC}}, E^{\text{SCC}})$.
- V^{SCC} has one vertex for each SCC in G .
- E^{SCC} has an edge if there's an edge between the corresponding SCC's in G .
- G^{SCC} for the example considered:

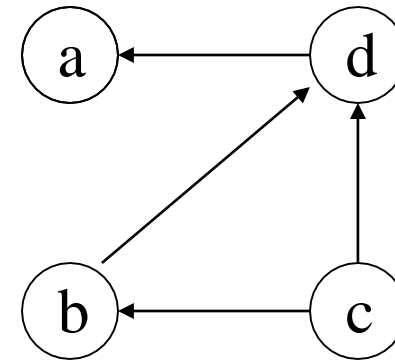


STRONGLY CONNECTED COMPONENTS

The **transpose** M^T of an $N \times N$ matrix M is the matrix obtained when the rows become columns and the column become rows:



G



G^T

Edges
have reverse
direction!

M

	a	b	c	d
a				1
b			1	
c				
d		1	1	

M^T

	a	b	c	d
a				
b				1
c		1		1
d	1			



TRANSPOSE OF A DIRECTED GRAPH

- $G^T =$ **transpose** of directed G .
 - $G^T = (V, E^T)$, $E^T = \{(u, v) : (v, u) \in E\}$.
 - G^T is G with all edges reversed.
- Can create G^T in $\Theta(V + E)$ time if using adjacency lists.
- G and G^T have the *same* SCC's. (u and v are reachable from each other in G if and only if reachable from each other in G^T .)



ALGORITHM TO DETERMINE SCCS

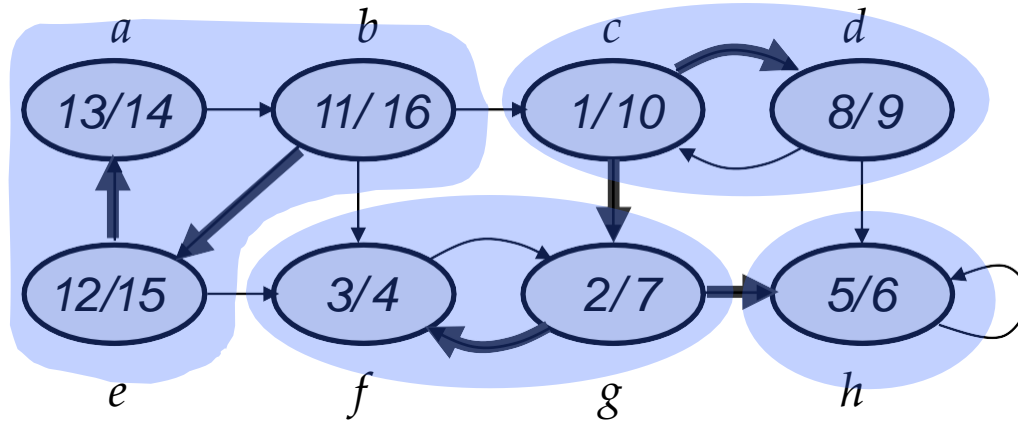
SCC(G)

1. call DFS(G) to compute finishing times $f[u]$ for all u
2. compute G^T
3. call DFS(G^T), but in the main loop, consider vertices in order of decreasing $f[u]$ (as computed in first DFS)
4. output the vertices in each tree of the depth-first forest formed in second DFS as a separate SCC

Time: $\Theta(V + E)$.

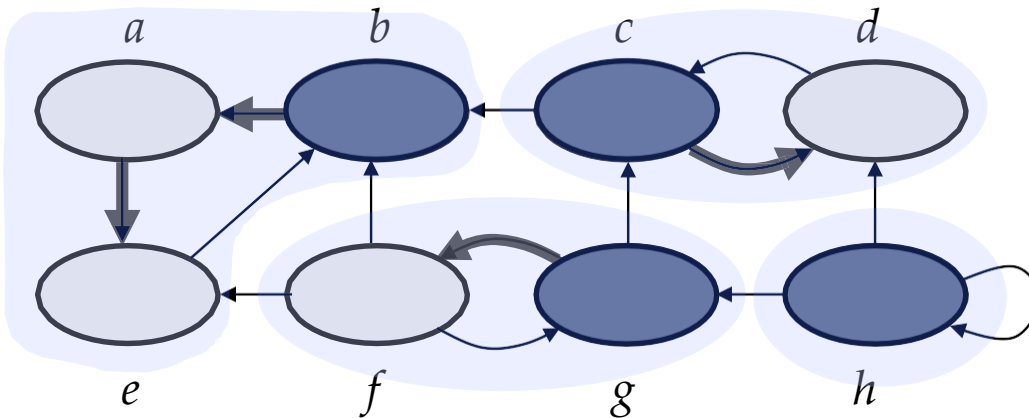


EXAMPLE



DFS on the initial graph G

b	e	a	c	d	g	h	f
16	15	14	10	9	7	6	4



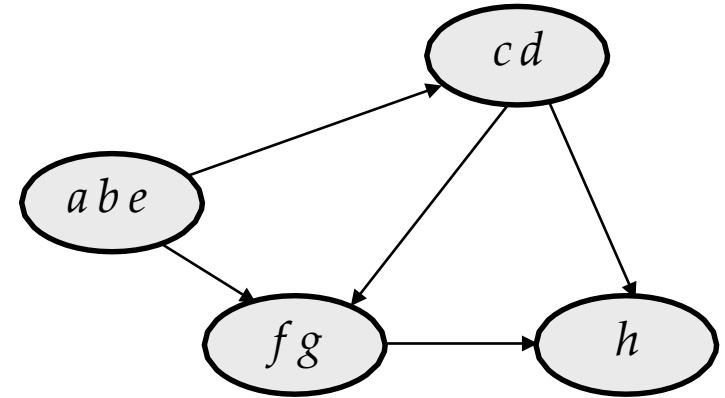
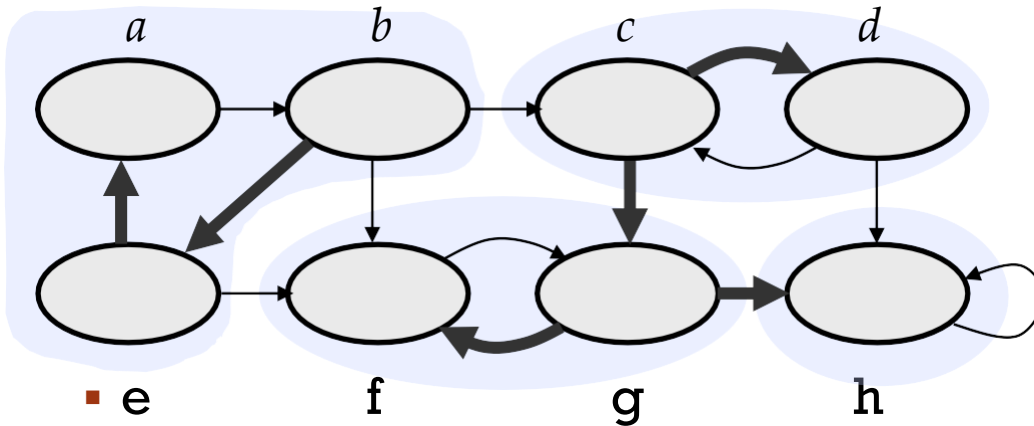
DFS on G^T :

- start at b : visit a, e
- start at c : visit d
- start at g : visit f
- start at h

Strongly connected components: $C_1 = \{a, b, e\}$, $C_2 = \{c, d\}$, $C_3 = \{f, g\}$, $C_4 = \{h\}$



COMPONENT GRAPH



⌘ The **component graph** $G^{\text{SCC}} = (V^{\text{SCC}}, E^{\text{SCC}})$:

- $V^{\text{SCC}} = \{v_1, v_2, \dots, v_k\}$, where v_i corresponds to each
 - strongly connected component \mathcal{C}_i
- There is an edge $(v_i, v_j) \in E^{\text{SCC}}$ if G contains a directed edge (x, y) for some $x \in \mathcal{C}_i$ and $y \in \mathcal{C}_j$

⌘ The component graph is a DAG

