Operating System Concepts (cont'd)





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Lecture Outline



- 1. I/O Structure
- 2. Storage Structure
- 3. How a Modern Computer Works
- 4. Direct Memory Access Structure
- 5. Computer-System Architecture
- 6. A Dual-Core Design
- 7. Clustered Systems
- 8. Operating-System Operations
- 9. Multiprogramming and Multitasking
- 10. Dual-mode and Multimode Operation
- 11. Process Management
- 12. Memory Management
- 13. File-system Management
- 14. Caching

I/O Structure

- After I/O starts, control returns to user program only upon I/O completion
 - Wait instruction idles the CPU until the next interrupt
 - Wait loop (contention for memory access)
 - At most one I/O request is outstanding at a time, no simultaneous I/O processing
- After I/O starts, control returns to user program without waiting for I/O completion
 - □ System call request to the OS to allow user to wait for I/O completion
 - □ Device-status table contains entry for each I/O device indicating its type, address, and state
 - OS indexes into **I/O device table** to determine device status and to modify table entry to include interrupt

I/O Structure

Device-status table

Device ID	Device Type	Address	State
0	Disk Drive	0x1000	Busy
1	Printer	0x1001	Idle
2	Keyboard	0x1002	Ready
3	Network Card	0x1003	Error

Storage Structure

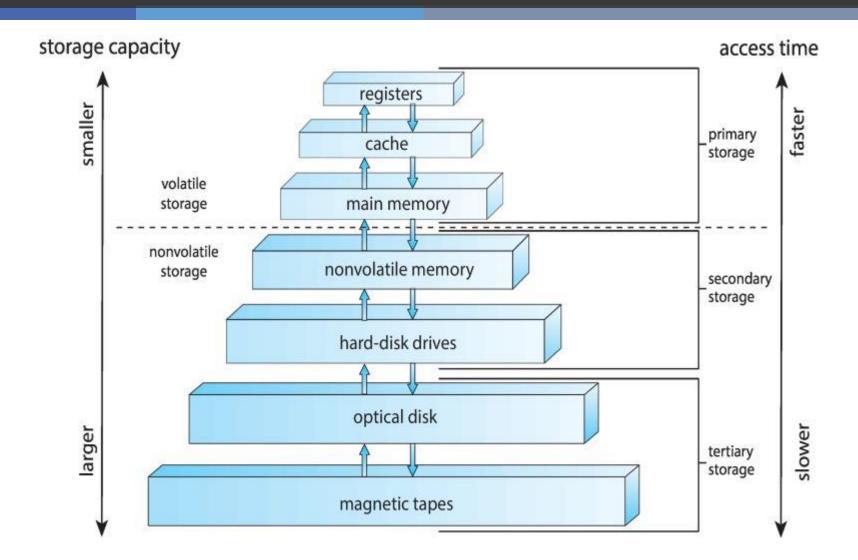
Main memory – only large storage media that the CPU can access directly **Random access** Typically volatile Typically random-access memory (RAM) in the form of Dynamic Random-access Memory (DRAM) **Secondary storage** – extension of main memory that provides large **nonvolatile** storage capacity Hard Disk Drives (HDD) – rigid metal or glass platters covered with magnetic recording material **Disk surface** is logically divided into tracks, which are subdivided into sectors The disk controller determines the logical interaction between the device and the computer Non-volatile memory (NVM) devices—faster than hard disks, Sector Track Various technologies

Becoming more popular as capacity and performance increases, price drops

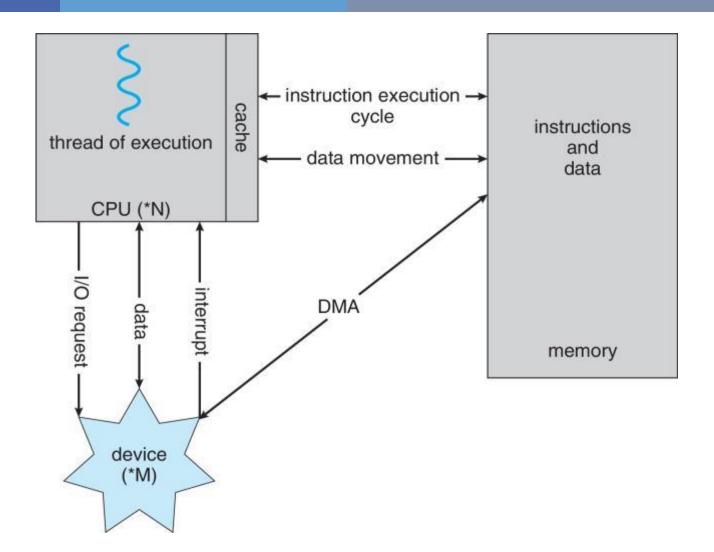
Storage Hierarchy

- □ Storage systems organized in hierarchy
 - Speed
 - □ Cost
 - Volatility
- □ Caching copying information into faster storage system; main memory can be viewed as a cache for secondary storage
- □ Device Driver for each device controller to manage I/O
 - Provides uniform interface between controller and kernel

Storage-Device Hierarchy



How a Modern Computer Works



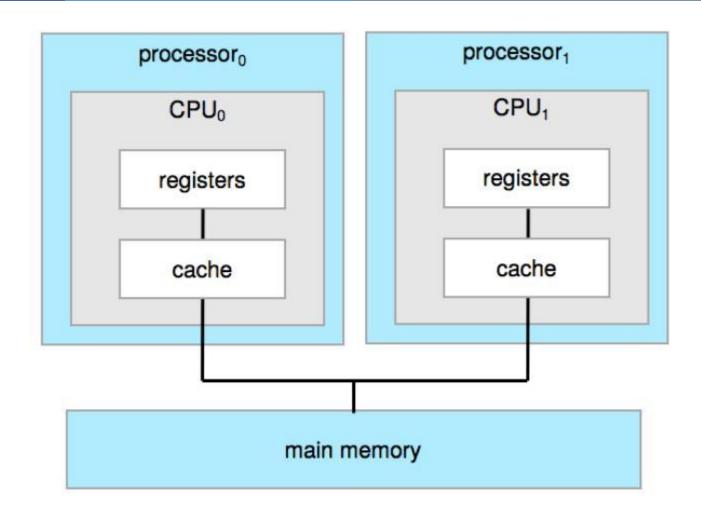
Direct Memory Access (DMA) Structure

- Used for high-speed I/O devices able to transmit information at close to memory speeds
- Device controller **transfers blocks of data from buffer storage directly to main memory** without CPU intervention
- Only one interrupt is generated per block, rather than the one interrupt per byte

Computer-System Architecture

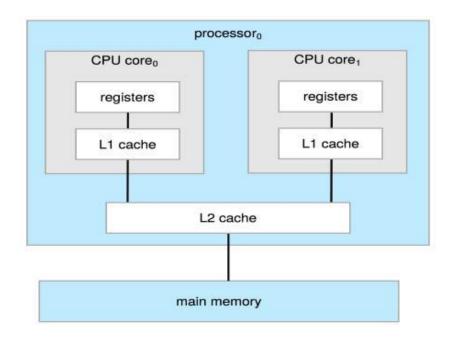
- ☐ Most systems use a single general-purpose processor
 - ☐ Most systems have special-purpose processors as well
- Multiprocessors systems growing in use and importance
 - ☐ Also known as **parallel systems**, **tightly-coupled systems**
 - ☐ Advantages include:
 - 1. Increased throughput
 - 2. Economy of scale
 - 3. **Increased reliability** graceful degradation or fault tolerance
 - ☐ Two types:
 - 1. **Asymmetric Multiprocessing** each processor is assigned a specie task.
 - 2. Symmetric Multiprocessing each processor performs all tasks

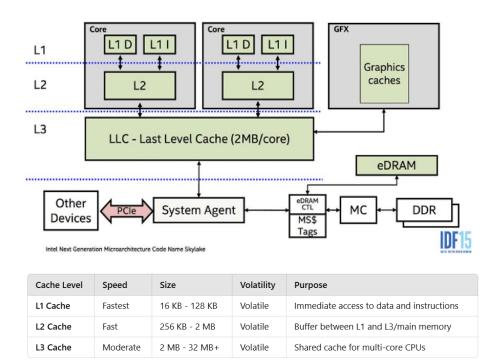
Symmetric Multiprocessing Architecture



A Dual-Core Design

- Multi-chip and multicore
- Systems containing all chips
 - ☐ Chassis containing multiple separate systems

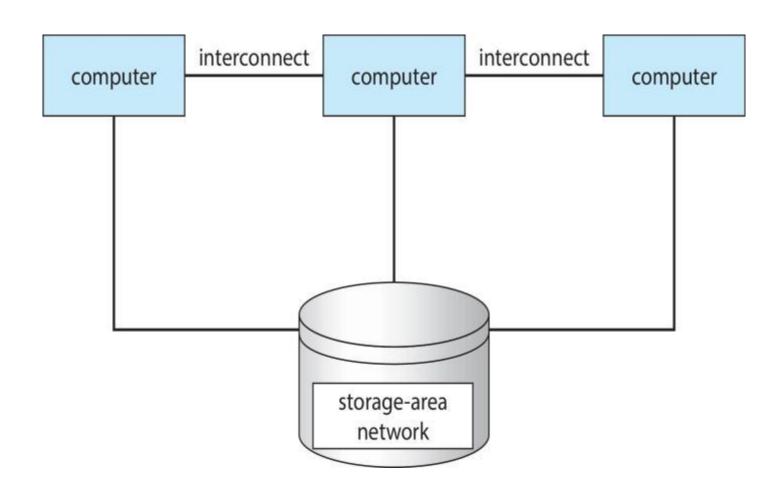




Clustered Systems

- ☐ Like multiprocessor systems, but multiple systems working together
 - ☐ Usually sharing storage via a **storage-area network (SAN)**
 - ☐ Provides a **high-availability** service which survives failures
 - □ **Asymmetric clustering** has one machine in hot-standby mode
 - Symmetric clustering has multiple nodes running applications, monitoring each other
 - □ Some clusters are for **high-performance computing (HPC)**
 - ☐ Applications must be written to use **parallelization**
 - □ Some have distributed lock manager (DLM) to avoid conflicting operations

Clustered Systems



Operating-System Operations

- Bootstrap program simple code to initialize the system, load the kernel
- Kernel loads
- ☐ Starts **system daemons** (services provided outside of the kernel)
- ☐ Kernel **interrupt driven** (hardware and software)
 - ☐ Hardware interrupt by one of the devices
 - □ Software interrupt (exception or trap):
 - □ Software error (e.g., division by zero)
 - Request for operating system service system call
 - Other process problems include infinite loop, processes modifying each other or the operating system

Multiprogramming and Multitasking

Mu	ultiprogramming (Batch system) needed for efficiency		
	Single user cannot keep CPU and I/O devices busy at all times		
	Multiprogramming organizes jobs (code and data) so CPU always has one		
	to execute		
	A subset of total jobs in system is kept in memory		
	One job selected and run via job scheduling		
	When it has to wait (for I/O for example), OS switches to another job		
Tin	nesharing (multitasking) is logical extension in which CPU switches jobs		
so i	frequently that users can interact with each job while it is running, creating		
interactive computing			
	Response time should be < 1 second		
	Each user has at least one program executing in memory process		
	If several jobs ready to run at the same time CPU scheduling		
	If processes doesn't fit in memory, swapping moves them in and out to		
	run		
	Virtual memory allows execution of processes not completely in memory		

Memory Layout for Multiprogrammed System

max operating system process 1 process 2 process 3 process 4

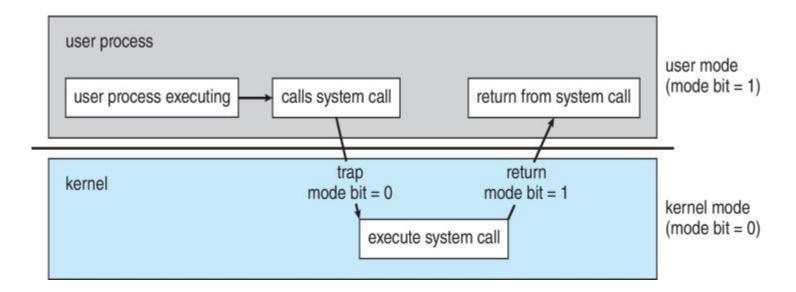
Dual-mode and Multimode Operation

- Dual-mode operation allows OS to protect itself and other system components
 - User mode and kernel mode
 - Mode bit provided by hardware
 - Provides ability to distinguish when system is running user code or kernel code
 - □ Some instructions designated as **privileged**, only executable in kernel mode
 - □ System call changes mode to kernel, return from call resets it to user
- ☐ Increasingly CPUs support multi-mode operations
 - ☐ i.e. virtual machine manager (VMM) mode for guest VMs

Transition from User to Kernel Mode

- ☐ Timer to prevent infinite loop / process hogging resources
 - ☐ Timer is set to interrupt the computer after some time period
 - ☐ Keep a counter that is decremented by the physical clock
 - Operating system set the counter (privileged instruction)
 - ☐ When counter zero generate an interrupt
 - ☐ Set up before scheduling process to regain control or terminate program that exceeds allotted time

Transition from User to Kernel Mode



Process Management

A process is a program in execution. It is a unit of work within the system.
Program is a passive entity, process is an active entity.
Process needs resources to accomplish its task
☐ CPU, memory, I/O, files
☐ Initialization data
Process termination requires reclaim of any reusable resources
Single-threaded process has one program counter specifying location of
next instruction to execute
☐ Process executes instructions sequentially, one at a time, until completion
Multi-threaded process has one program counter per thread
Typically system has many processes, some user, some operating system
running concurrently on one or more CPUs
☐ Concurrency by multiplexing the CPUs among the processes / threads

Process Management Activities

The operating system is responsible for the following activities in connection with process management:

- Creating and deleting both user and system processes
- Suspending and resuming processes
- Providing mechanisms for process synchronization
- Providing mechanisms for process communication
- Providing mechanisms for deadlock handling

Memory Management

- ☐ To execute a program all (or part) of the instructions must be in memory
- All (or part) of the data that is needed by the program must be in memory
- Memory management determines what is in memory and when
 - Optimizing CPU utilization and computer response to users
- Memory management activities
 - ☐ Keeping track of which parts of memory are currently being used and by whom
 - ☐ Deciding which processes (or parts thereof) and data to move into and out of memory
 - ☐ Allocating and deallocating memory space as needed

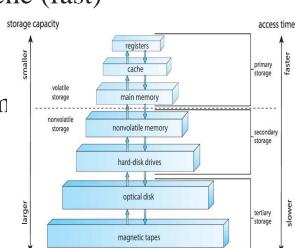
File-system Management

OS provides uniform, logical view of information storage Abstracts physical properties to logical storage unit - file Each medium is controlled by device (i.e., disk drive, tape drive) Varying properties include access speed, capacity, data-transfer rate, access method (sequential or random) File-System management Files usually organized into directories Access control on most systems to determine who can access what OS activities include **Creating and deleting** files and directories Primitives to **manipulate** files and directories ☐ **Mapping files** onto secondary storage

Backup files onto stable (non-volatile) storage media

Caching

- ☐ Important principle, performed at many levels in a computer (in hardware, operating system, software)
- ☐ Information in use copied from slower to faster storage temporarily
- Faster storage (cache) checked first to determine if information is there
 - ☐ If it is, information used directly from the cache (fast)
 - ☐ If not, data copied to cache and used there
- Cache smaller than storage being cached
 - □ Cache management important design problem
 - Cache size and replacement policy



Books



- Operating Systems Concept
 - ☐ Written by Galvin and Silberschatz
 - ☐ Edition: 9th

References

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- Operating Systems Concept
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