```
#include
#include
const struct device *get_gpio_device(const char *label) {
return device_get_binding(label);
}
```

Zephyr Device Model

Zephyr uses a device model where hardware peripherals (like GPIOs, UARTs, etc.) are represented as struct device objects. These are registered during system initialization.

device_get_binding(label)

This function looks up a device by its **label**, which is defined in the **Device Tree** or board configuration files. For example, "GPIO_O" or "GPIOA".

- It returns a pointer to the struct device representing the GPIO controller.
- If the label is incorrect or the device isn't initialized, it returns NULL.

GPIO Driver Interaction

gpio_pin_set(dev, pin, value);

Once you have the device pointer, you can use it with GPIO driver APIs like:

```
gpio_pin_configure(dev, pin, GPIO_OUTPUT);
```

These functions use the device pointer to interact with the actual hardware registers via the GPIO driver.

```
gpio_setup() → get_gpio_device(label) → device_get_binding(label)

→ DEVICE_DT_DEFINE(...) → gpio_emul_driver (via DEVICE_API)
```

- get_gpio_device() is a helper to fetch the GPIO device.
- It enables your code to be hardware-agnostic, relying on labels from the Device Tree.
- Once you have the device, you can configure pins, read/write values, and handle interrupts using Zephyr's GPIO API.

1. Device Tree Source (DTS) Path for QEMU x86

vi ../../boards/qemu/x86/qemu_x86.dts

```
gpio0: gpio_sim {
       compatible = "zephyr,gpio-emul";
    gpio-controller;
   #gpio-cells = <2>;
    ngpios = <32>;
    label = "GPIO_0";
};
static void gpio_setup(void *fixture){
gpio_dev = get_gpio_device(GPIO_LABEL);
zassert_not_null(gpio_dev, "Failed to get GPIO device");
}
```

- This function is part of a test setup.
- It calls **get_gpio_device()** with a label like "GPIO_0".

• This wraps device_get_binding(), which looks up a device by its label.

device_get_binding(label)

- This searches the global device list for a device with a matching label (from the Device Tree).
- If found, it returns a pointer to the struct device.

Device Tree + Driver Binding

- The label (e.g., "GPIO_O") is defined in the Device Tree (.dts or .overlay).
- The corresponding node is associated with a driver using DEVICE_DT_DEFINE() or DEVICE_DEFINE().

```
DEVICE_DT_INST_DEFINE(_num, gpio_emul_init, \
PM_DEVICE_DT_INST_GET(_num), \
&gpio_emul_data_##_num, \
&gpio_emul_config_##_num, POST_KERNEL, \
CONFIG_GPIO_INIT_PRIORITY, \
&gpio_emul_driver);
```

"../../drivers/gpio/gpio_emul.c

```
static DEVICE_API(gpio, gpio_emul_driver) = {
    .pin_configure = gpio_emul_pin_configure,
#ifdef CONFIG_GPIO_GET_CONFIG
   .pin_get_config = gpio_emul_pin_get_config,
#endif
   .port_get_raw = gpio_emul_port_get_raw,
    .port_set_masked_raw = gpio_emul_port_set_masked_raw,
   .port_set_bits_raw = gpio_emul_port_set_bits_raw,
    .port_clear_bits_raw = gpio_emul_port_clear_bits_raw,
    .port_toggle_bits = gpio_emul_port_toggle_bits,
    .pin_interrupt_configure = gpio_emul_pin_interrupt_configure,
   .manage_callback = gpio_emul_manage_callback,
    .get_pending_int = gpio_emul_get_pending_int,
#ifdef CONFIG_GPIO_GET_DIRECTION
    .port_get_direction = gpio_emul_port_get_direction,
#endif /* CONFIG_GPIO_GET_DIRECTION */
```

- This struct defines the function pointers for the GPIO emulator's API.
- It is passed to DEVICE_DT_DEFINE() to register the device.

Final Mapping

- When device_get_binding("GPIO_0") is called, it returns a struct device whose .api field points to gpio_emul_driver.
- So when you later call gpio_pin_configure(gpio_dev, ...), it internally calls gpio_emul_pin_configure().

How gpio_pin_set() Maps to the Emulator Driver

gpio_pin_set(gpio_dev, TEST_PIN, 1);

This is a public Zephyr API defined in include/zephyr/drivers/gpio.h.

Device API Dispatch

Internally, this calls the function pointer from the device's driver API: gpio_dev->api->port_set_bits_raw(...)

This pointer is set to the emulator's implementation in gpio_emul_driver:

```
static const struct gpio_driver_api gpio_emul_driver = {
...
.port_set_bits_raw = gpio_emul_port_set_bits_raw,
.port_clear_bits_raw = gpio_emul_port_clear_bits_raw,
...
};
```

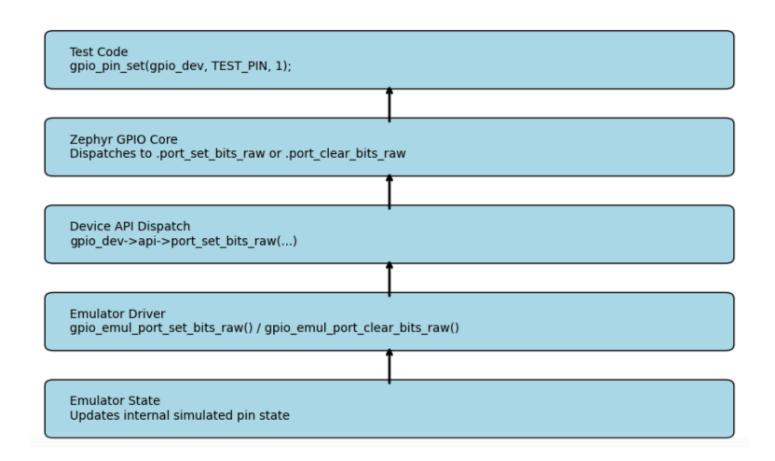


Diagram Layers Explained

1. Test Code

Calls gpio_pin_set(gpio_dev, TEST_PIN, 1); in your test suite.

2. Zephyr GPIO Core

Routes the call to the appropriate function pointer in the GPIO driver API.

3. Device API Dispatch

Internally calls gpio_dev->api->port_set_bits_raw(...) or port_clear_bits_raw(...).

4. Emulator Driver

Executes gpio_emul_port_set_bits_raw() or gpio_emul_port_clear_bits_raw() from gpio_emul.c.

5. Emulator State

Updates the internal simulated pin state (e.g., a bitfield or array) to reflect the pin value.

/zephyrproject/zephyr/samples/Zaphyr_app/my_zephry_app_threadPrio/my_zephry_app_threadPrio

Threads Overview

Thread 1: Sleeping & Suspended

- Function: thread1_fn
- Behavior:
 - Sleeps for 2 seconds repeatedly using k_sleep(), which puts it in the sleeping state.
 - Can be **suspended** by Thread 3 using k_thread_suspend(), which halts its execution regardless of its sleep state.
 - Later **resumed** by Thread 3 using k_thread_resume().

▼ Thread 2: Pending

- Function: thread2_fn
- Behavior:
 - Waits indefinitely on a semaphore using k_sem_take(&sync_sem, K_FOREVER).
 - This puts it in the pending state until another thread gives the semaphore.
 - Once the semaphore is given by Thread 3, it prints a message and loops back to wait again.

✓ Thread 3: Running & Controller

- Function: thread3 fn
- Behavior:
 - Runs continuously and controls the other two threads.
 - Suspends Thread 1, waits 3 seconds, then resumes it.
 - Gives the semaphore to Thread 2 to unblock it.
 - Sleeps for 5 seconds before repeating.
- This thread is always running when active and demonstrates how a thread can manage others.

Thread Priorities

- Thread 1: Priority 1 (highest)
- Thread 3: Priority 4
- Thread 2: Priority 5 (lowest)

In Zephyr, lower numbers mean higher priority. So Thread 1 will preempt others unless it's sleeping or suspended.

Execution Flow Summary

1. Startup:

All three threads are created and start running.

2. Thread 1:

- Sleeps for 2 seconds → enters sleeping state.
- May be **suspended** by Thread 3 during this time.

3. Thread 2:

- Waits on a semaphore → enters **pending state**.
- Remains blocked until Thread 3 gives the semaphore.

4. Thread 3:

- Suspends Thread 1 → Thread 1 enters suspended state.
- Sleeps for 3 seconds.
- Resumes Thread 1 → Thread 1 can now run again.
- Gives semaphore to Thread 2 → Thread 2 wakes up.
- Sleeps for 5 seconds and repeats the cycle.

Thread States Demonstrated

Thread	State(s) Demonstrated	How?
Thread 1	Sleeping, Suspended	k_sleep(), k_thread_suspend()
Thread 2	Pending	k_sem_take()
Thread 3	Running	Actively controls other threads

Zephyr is a **lightweight, scalable, real-time operating system (RTOS)** designed for embedded systems. It supports multiple architectures and is highly modular, allowing developers to include only the components they need.

1. Kernel Architecture

- Small-footprint kernel: Designed for resource-constrained devices.
- Supports multiple CPU architectures: ARM, x86, RISC-V, ARC, MIPS, and more

•

2. Threading Model

Zephyr supports:

- Cooperative threads: Run until they yield or block.
- Preemptive threads: Can be interrupted by higher-priority threads.
- Round-robin scheduling: Among threads of equal priority.
- POSIX pthreads API: Optional compatibility layer.

3. Scheduling Algorithms

- Cooperative and Preemptive Scheduling
- Earliest Deadline First (EDF)
- Meta IRQ scheduling: For deferred interrupt handling.
- Timeslicing: Among equal-priority preemptible threads.

4. Memory Management

- Heap and slab allocators
- Thread stacks: Configurable per thread.
- Custom memory regions: For advanced use cases.

5. Inter-thread Communication

- Semaphores: Binary and counting.
- Mutexes: For mutual exclusion.
- Message Queues and Pipes: For data passing.

6. Power Management

- System Power Management: Application-defined policies.
- Device Power Management: Driver-level control.

7. Modularity

• Zephyr is **highly configurable** via Kconfig and device tree.

• Developers can enable/disable features to optimize footprint.

System Threads in Zephyr

System threads are **automatically spawned by the kernel** during system initialization. They are essential for basic OS operation.

Types of System Threads

1. Main Thread

- Purpose: Executes kernel initialization and then calls the user-defined main() function.
- Priority:
 - Highest preemptible priority (0) if preemption is enabled.
 - Lowest cooperative priority (-1) if not.
- Behavior:
 - If main() is defined and returns normally, the thread terminates without error.
 - If main() is missing or the thread aborts, a fatal error is raised.

2. Idle Thread

- Purpose: Runs when no other thread is ready.
- Behavior:
 - May invoke power-saving routines.
 - Otherwise, runs a "do-nothing" loop.
- Priority: Always the lowest.
- Essential: Cannot terminate; if it aborts, a fatal error occurs.

3. Workqueue Threads

- Spawned if the **system workqueue** is enabled.
- Handle deferred work items submitted by the application.

Writing a main() Function

Example: