### **Iraya Use Case Design Document**

This document contains a proposed design for the use case in Neil Ongkingco's Iraya Backend Engineer Interview Test.

### **General Overview**

The system is designed to be a cluster of 5 nodes (1 parent VM and 4 worker VMs) using the Kubernetes and Docker cluster and container frameworks. Each node will host several software containers that will provide the functionality required to fulfill the responsibilities of the node.

In addition to the 5 nodes, the system uses a keyed **FileStore** to store the raw document data and a **Transaction Database** (SQL) for storing file info, computational results and task status for each document.

The **parent VM** has the following responsibilities:

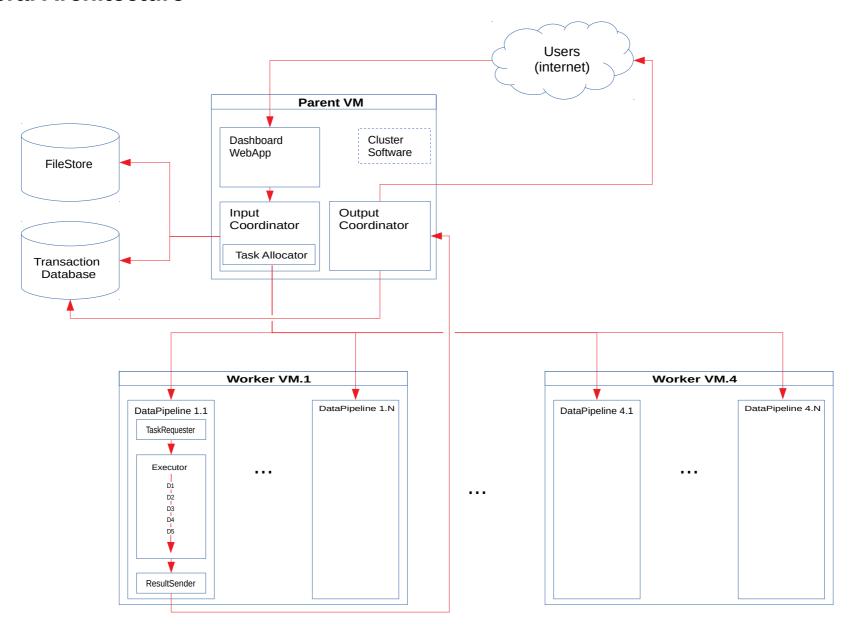
- host the container that provides the web interface to the users to upload their files
- store the uploaded documents to a FileStore. At the same time, it will create a
   *Task* record for the file and record relevant file info and task state in the
   Transaction Database
- package and dispatch document processing tasks to the DataPipeline containers in the worker nodes
- receive and store the results computed by the workers into the transaction database
- provide notifications to the users on the progress of the document processing tasks (total progress, duplicate document and processing failure emails)

In addition the parent VM will also act as the master of the cluster and host the Kubernetes control plane.

The **worker nodes** will host multiple **DataPipeline** containers that will actually execute the steps of the data pipeline on the documents. The goal here is to parallelize document processing to maximize throughput. These containers have the following responsibilities:

- Request and process document processing tasks from the master node
- Send each document through the 5 steps (D1 to D5) of the data processing pipeline
- Package the results upon completion of the 5 step pipeline and send them to the parent VM for processing and storage
- Handle any failures in data processing by retrying or upon too many failures notify the master node of the failure of the task

### **General Architecture**



The figure above shows the general architecture of the proposed design, and details the major components of the system as well as the high-level data-flows between the components of the system and the outside world.

The system consists of the following major components:

- Parent Virtual Machine (Parent VM): The virtual machine that hosts the containers
  that coordinate the functions of the components in the worker nodes, as well as the
  WebApp container that runs the Web Dashboard that users will use to uploading
  documents. The Parent VM will also be the master node that hosts the Kubernetes
  control plane that manages the cluster. The parent node will host the following
  containers:
  - Dashboard WebApp: runs the dashboard web application that the users will access to upload files
  - *Input Coordinator*: receives the uploaded file from the Dashboard WebApp and performs the following:
    - create a unique FileID for the document (by computing the SHA256 hash on the file contents)
    - detect if the document is a duplicate and send the necessary email notifications to the users while aborting storage of the document data
    - store the file contents in the FileStore keyed with FileID
    - create a *task entry* in the Transaction database that gives assigns a document a *taskID* and contains basic info on the task (*FileID*, name, size, type, date uploaded, task state).

The *Input Coordinator* will also have a *Task Allocator* sub-component that packages the document into *tasks* (*taskID* + *FileID* + file contents) and orders them into a queue. It then listens to requests from DataPipeline components and serves tasks if any are available in the queue.

- Output Coordinator: receives the computed results for the document from the DataPipeline component on the worker VMs, stores the results in the Transaction Database using the corresponding transaction ID, as well as handling any errors caused by processing failures in the worker nodes. This includes notifying users of any documents that have failed to process after the appropriate number of retries.
- Worker Virtual Machines 1 to 4 (Worker VM.1 to VM.4): These virtual machines
  will contain multiple instances of *DataPipeline* containers which will actually perform
  the data processing tasks on a document. The goal is to have as many *DataPipeline* containers running on a worker VM to parallelize the processing of the
  documents and overall throughput.

- The **DataPipeline** containers have the following sub-components:
  - TaskRequester: Constantly requests tasks from the Input Coordinator as long as the Executor is not currently busy on a task. Upon receiving a task, it extracts the document data and passes it to the Executor for processing.
  - **Executor**: Runs the steps D1 to D5 in succession on the document data and collects the results. The executor is also responsible for retrying a step (up to 3 times) in case of failure. The results (or failure) are then passed to the *ResultSender*.
  - **ResultSender**: packages the results (or failure details) with the corresponding taskID and sends them to the *Output Coordinator* on the *Parent VM* for final processing.

### **Technology Review**

The two major cloud technologies that the proposed system will use are **Docker** for software containers and **Kubernetes** for cluster management.

**Docker** is a platform that allows the packaging of software services (i.e. web servers, ML compute tasks) into a *container* that has all the dependencies needed by the service and be able to run it in a lightweight, standalone execution environment. Multiple containers can be hosted in a single virtual machine, therefore decoupling software services from specific hosts and allowing for the deployment of multiple instances of the same service in a single machine for improved reliability, increased throughput and higher resource utilization. For our purposes, each Docker container can be treated as a separate host, and can communicate (with some configuration and some help from Kubernetes) to other containers using standard web messaging interfaces (HTTP requests).

Using Docker will allow for easier development of the features required by the system as they can be coded and tested in isolation on developer machines with a reasonable guarantee that with the correct deployment configuration they will also work correctly on the cloud. The ability to create multiple instances of the same container will also be essential to allow for the creation of as many *DataPipeline* components as the worker VMs can handle to parallelize the processing of the documents through the 5 step data pipeline.

<u>Kubernetes</u> allows the setup and management of clusters: a set of nodes (i.e. virtual machines) that host multiple containers. Kubernetes has several features that will significantly affect the ease of implementation and deployment of the proposed system.

The first feature of Kubernetes that is relevant to the design system is its ability to deploy multiple instances of a component over all worker nodes with ease (you only need to specify the component image to use and how many replicas you want to deploy and it can

be done in a single configuration file without having to reference the explicit nodes were the components will be deployed). This relieves a lot the coordination required from the master node when initializing the *DataPipeline* components, and will also allow for changing the number and resource allocations of the deployed containers on the fly to maximize document throughput.

The second relevant feature is Kubernetes' ability to allow components to directly connect to each other without going through the container's node first. Kubernetes' network is designed so a container is identifiable by its assigned IP address *across all nodes*. This virtually allows the system to treat each component as a separate host, using standard web messaging techniques (HTTP requests) to transfer data. This will significantly ease the development of the logic required by the various containers required by the system.

A tertiary feature that will be relevant when the system is deployed in a real environment is Kubnertes' ability to recover from system failures (both node and container crashes). With careful design of the logic in the coordinator and *DataPipeline* containers (by enforcing idempotence and atomicity of all write operations to the *FileStore* and *Transaction Database*). Kubnertes allows for the restarting of failed nodes and containers and these can just look at the state of the tasks in the Transaction Database and continue execution of any unfinished tasks.

The technologies to be used for the systems main data storage components (*FileStore* and *Transaction Database*) are also worth discussing.

The *FileStore* only requires that the documents' contents be stored in a way that they can be retrieved using a key that uniquely identifies that document. Using a cryptographic hash (such as SHA256) over the document's contents should be sufficient to produce this key. As for the storage technology to use, a cloud BLOB storage that has atomic reads and writes should be sufficient for the purpose, again using the files SHA256 hash as the key.

The use of NoSQL databases for storing file contents has been explored but from preliminary review it seems that the most popular NoSQL solutions for storing large files have limits that could be exceeded by the expected size of the documents, or have write performance that degrades as document sizes become larger (MongoDB). Hence the decision to use a simple keyed BLOB store.

The *Transaction Database* will use a traditional RDBMS (SQL). This satisfies the requirements of the system for task management and result storage (atomicity and the ability to quickly perform join and filter queries).

**Designer's note**: These reviews are based on a relatively light overview of the documentation and beginner presentations for Docker and Kubernetes. Due to tight time constraints and limited experience, I have not been able to do proof of concept implementations to properly try out their features (particularly for Kubernetes). There could be hidden caveats in these technologies that would require changes to the design.

### **Data Stores and Schemas**

First we start with a description of the data stores used by the system and the structure of the data stored within. This will provide a foundation for all the processes and dataflows that will occur in the system. The system uses two major data stores: the **FileStore** for storing the raw file content, and the **Transaction Database** for storing file, task and results data.

### **FileStore**

The *FileStore* is a simple BLOB (binary large object) store which has the following properties:

- Allows storage and retrieval of a document's contents based on a FileID unique to that document
- · Atomicity (all or nothing) of read and write operations

A simple cloud blob storage solution (like Azure blob storage) would satisfy these requirements with the appropriate choice of *FileID*.

The *FileID* itself could simply be the SHA256 hash of the contents of the file. Given the properties of the SHA256 cryptographic function (which turns any size file into a 256-bit binary string with negligible chances ( $\sim 10^{-60}$  for identical hash values for 1 billion random documents) it should be a sufficient choice for use as the *FileID*.

#### **Transaction Database**

The transaction database is intended to store the following information:

- A record of a document's basic information (*FileID*, name, type, size, owner/uploader among others)
- A record of the task that will process the document (this maps a task to the *FileID* of the document that it is processing as well has having the current status of the task)
- The results of running a document through the pipeline mapped to the document's FileID
- A log of the task's status changes and the times when they happened, which will be useful for reporting, diagnostics and error recovery

The rationale for this structure is the creation of a logical *Task* entity that will contain all the state required by the system to process and keep track of a document's progress as it passes through the system. As a high-level object, a task can be thought of as

```
(TaskID, FileID, FileInfo, TaskStatus, <results>)
```

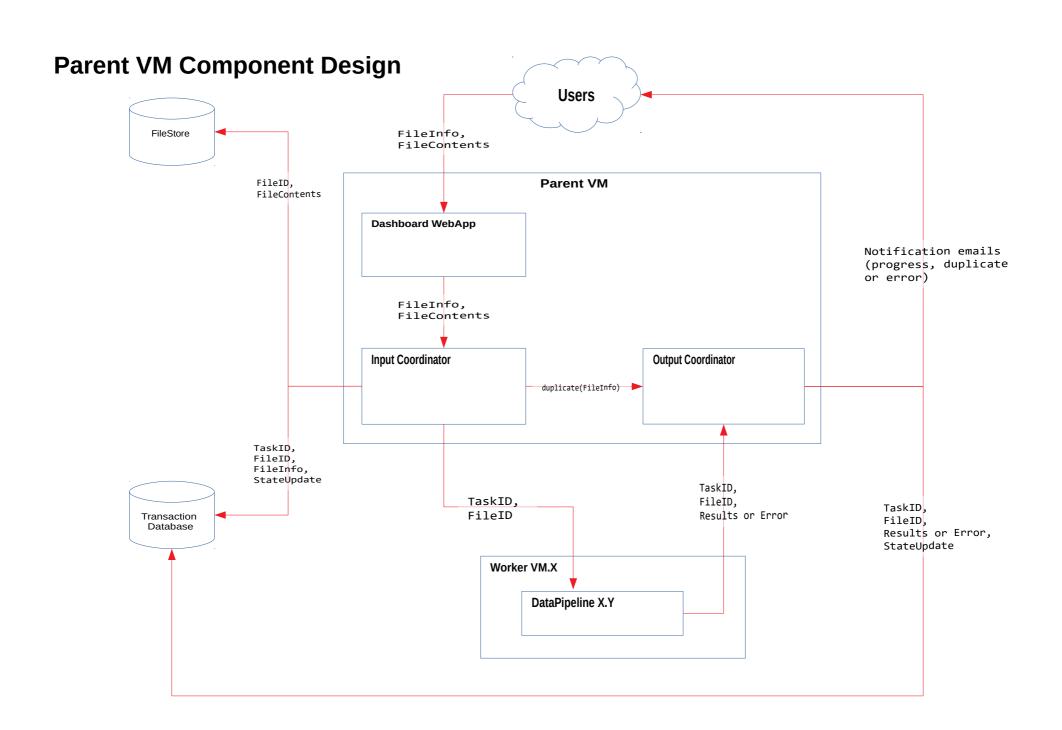
The schema of the database is an implementation of this high level object as relational database tables. A more detailed description of the Transaction Database's schema (the structure of the tables in the database) is shown below in pseudo-SQL format. Comments are in light green.

```
table Task(
     TaskID: int identity
     TaskStatus: string (one of UPLOADED|PROCESSING|DONE|ERROR)
)
table FileInfo(
     TaskID: int (foreign key to Task)
     FileID: string (identical to FileID used in FileStore)
     Name: string
     Size: numeric
     Type: string
     Owner: string (the user that uploaded the file)
     (other fields as needed)
)
table Results(
     TaskID: int (foreign key to Task)
     FileID: string
     <results>: varies
)
table TaskLog(
     TaskID: int (foreign key to Task)
     StatusChange: string (same values as TaskStatus)
     ChangeTime: timestamp (time when the status change occurred)
     Message: string (any info/diagnostic message attached to the
status change)
)
```

Note that the Results table is just a generalized abstraction for how the results would be stored in the database. It is likely that different types of results will actually be stored in

different tables on the database (e.g. table MapData(TaskID, FileID, latitude, longitude)).

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The diagram above shows a more detailed view of the dataflows to and from the components of the Parent VM containers.

#### **Definition of Data Labels**

Before going into a detailed description of how the components work, we define the different kinds of data used by the system.

- **FileInfo**: a collective label for all the basic information about a file (name, size, type, etc)
- FileContents: raw binary contents of a file
- FileID: a key that uniquely identifies a file
- TaskID: a key that uniquely identifies a task. Is generated as a new task is written to the database
- **StateUpdate**: a collective label for the changes required to update the TaskStatus and TaskLog entries of a Task to reflect a new state
- Results: a collective label for the computed results after the FileContents have been processed by the 5 step data pipeline
- **Error**: the information of the details of an error that occurred when trying to process the data at any point in the system

### **Component Descriptions**

The Parent VM has 3 major components, the **Dashboard WebApp**, the **Input Coordinator** and the **Output Coordinator**. The following sections will describe the functionality of each using pseudocode followed by a short section on the rationale for the choices made in their design.

### **Dashboard WebApp**

Pseudocode:

```
while true:
    wait for upload from user
    send FileInfo, FileContents to InputCoordinator
```

From the standpoint of the system, the dashboard app is a simple pass-through that just accept uploads from a user and passes it on to the Input Coordinator.

### **Input Coordinator**

Pseudocode:

```
(UploadHandler: handling upload receipt from WebApp)
while true:
     wait for FileInfo, FileContents from WebApp
     compute FileID from FileContents
     if FileID already in FileStore:
           send dup_file(FileInfo) event to OutputCoordinator
     else:
           store FileID, FileContents in FileStore
           add Task in DB from FileID, FileContents with state UPLOADED
(TaskAllocator: handling task request from DataPipeline)
while true:
     wait for TaskRequest from DataPilelines
     if exists a Task with state == UPLOADED:
           find oldest Task with state == UPLOADED
           reply with TaskID, FileID of the Task
           apply a StateUpdate to change the Task to the PROCESSING state
     else:
           reply with No Task available
```

The input coordinator has 2 subcomponents running at the same time: the **UploadHandler** which accepts uploads from the WebApp, and the **TaskAllocator** which handles TaskRequests from the system's DataPipeline containers. In implementation these will just be HTTP request handlers listening on 2 different ports.

The UploadHandler is pretty straightforward, it just receives the upload, checks if it is a duplicate and if it is not stores the file and updates the Transaction Database with the appropriate information.

The TaskAllocator deals with the other side: it listens for TaskRequests from the system's DataPipelines and allocates any Tasks in the UPLOADED state to these requests in a first-in-first-out order.

A pull architecture (DataPipelines sending requests to the Input Coordinator instead of the coordinator pushing tasks out to the pipelines) was chosen to minimize the amount of state that the coordinator has to store, simplifying its implementation while also providing better error recovery in case either the Input Coordinator container or Parent VM crashes. This is also in line with the general principles of the design of Kubernetes, which uses a similar architecture, which should create less framework friction (i.e. trying to do what the framework doesn't want you to do) during implementation.

### **Output Coordinator**

Pseudocode:

```
while true:
    listen to events from Input Coordinator and DataPipelines
    if event == dup_file(FileInfo) from Input Coordinator:
        send duplicate document email to Users
    elseif event == (TaskID, FileID, results) from DataPipeline:
        store (TaskID, FileID, results) in DB
        perform StateUpdate on Task to set state to DONE
        if (number of DONE Tasks % 1000 == 0):
            send progress email to Users
    elseif event == (TaskID, FileID, error) from DataPipeline:
        perform StateUpdate on Task to set state to ERROR
        send document error email to Users
```

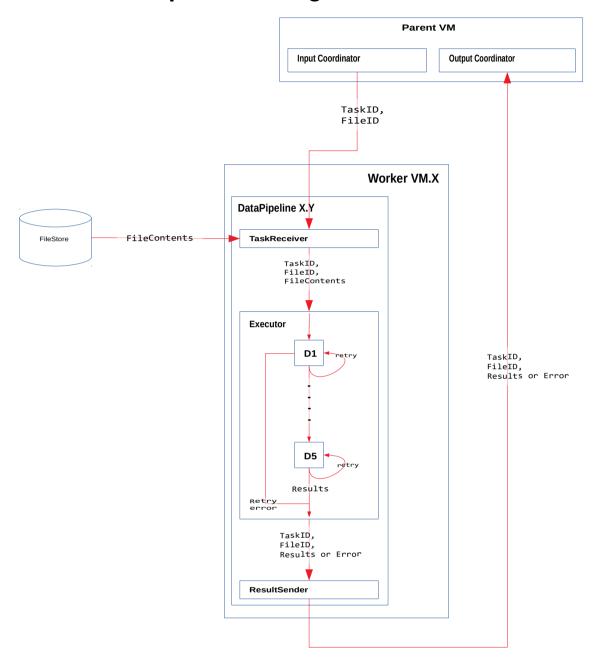
The Output Coordinator is a simple event handler that listens for events from DataPipelines and the Input Coordinator and performs the corresponding action. In implementation this is just an HTTP request handler listening at a given port.

#### **Additional Considerations**

Since the Parent VM hosts both the coordination infrastructure and the user-facing dashboard webapp, it is best to overprovision its containers to ensure that there are no resource shortages that could affect either the user experience or bring down coordination tasks that will have system-wide effects. Given the small number of users and an expected low number of messages per second to process, CPU utilization is expected to be pretty low, while memory and intermediate storage use could potentially be high if file sizes get too large.

However, if actual usage shows that there is enough unused resources on the Parent VM, Kubernetes should make it possible to run a few DataPipeline containers on the parent VM as well to increase utilization and throughput, as long as the appropriate quotas are set to make sure the coordination and user facing containers have enough resources for smooth operation.

## **Worker VM Component Design**



The diagram above is a more detailed view of the dataflows of the worker VM components.

### **DataPipeline**

Pseudocode:

```
(TaskRequester)
while true:
     wait(POLL INTERVAL)
     if Executor is READY:
           send TaskRequest to Input Coordinator
           if Task(TaskID, FileID) received
                 fetch FileContents from FileStore using FileID
                 send (TaskID, FileID, FileContents) to Executor
(Executor)
while true:
     wait for (TaskID, FileID, FileContents) from TaskRequester
     set state to BUSY, clear Errors
     stepInput = fileContents
     for dataStep in [D1...D5]:
           retries = 0
           while retries <= 3:
                 stepResult = dataStep(stepInput)
                 if no error:
                       stepInput = stepResult
                       continue to next dataStep
           set RetryError (maximum retries reached)
           break (get out for loop, maximum retries)
     if no error:
           Results = stepInput
           send (TaskID, FileID, Results) to ResultSender
     else:
           send (TaskID, FileID, RetryError) to ResultSender
     set state to READY
(ResultSender)
while true:
     wait for (TaskID, FileID, Results or Error) from Executor
     if results:
           send (TaskID, FileID, Results) to Output Coordinator
     else if error:
           send (TaskID, FileID, Error) to Output Coordinator
```

The DataPipeline has 3 subcomponents: the **TaskRequester** that periodically requests tasks from the Input Coordinator if the Executor is not busy, the **Executor** that runs the data steps on the FileContents, handling retries when needed, and the **ResultSender** that handles the results or errors generated by the **Executor** and sends them to the Output Coordinator.

In terms of implementation, though all these components can be implemented in a single process (with the TaskRequester being the outer loop that calls Executor and then ResultSender), there is an argument to be made for implementing them as separate processes as the Executor is expected to consume a lot of resources and can cause process-ending faults like Out of Memory errors. If they are run as separate processes, the TaskRequester and ResultSender would still be running and can monitor the failure of the Executor and take the necessary action (in this case the ResultSender would send a document error message to the Output Coordinator and then restart the Executor).

# **Additional System Design**

**Detailed Logging** 

**System Failure Recovery**