

Artificial Intelligence

ENCS 3340

Constraint Satisfaction Problems (Local Search)

Constraint Satisfaction

- Specifies structural properties of the problem
 - may depend on the representation of the problem
- The problem is defined through a set of variables and a set of domains
 - variables can have possible values specified by the problem
 - constraints describe allowable combinations of values for a subset of the variables
- **state** in a CSP
 - defined by **an assignment** of **values** to some or all **variables**
- **solution** to a CSP
 - must assign values to ALL variables
 - must satisfy ALL constraints
 - solutions may be ranked according to an **objective function**

Example1: 3-SAT

Variables:

x_1, x_2, x_3, x_4, x_5

Domains:

$\{\text{True}, \text{False}\}$

Constraints: \vee and \wedge

$(x_1 \vee x_2 \vee x_4),$

$(x_2 \vee x_4 \vee \neg x_5),$

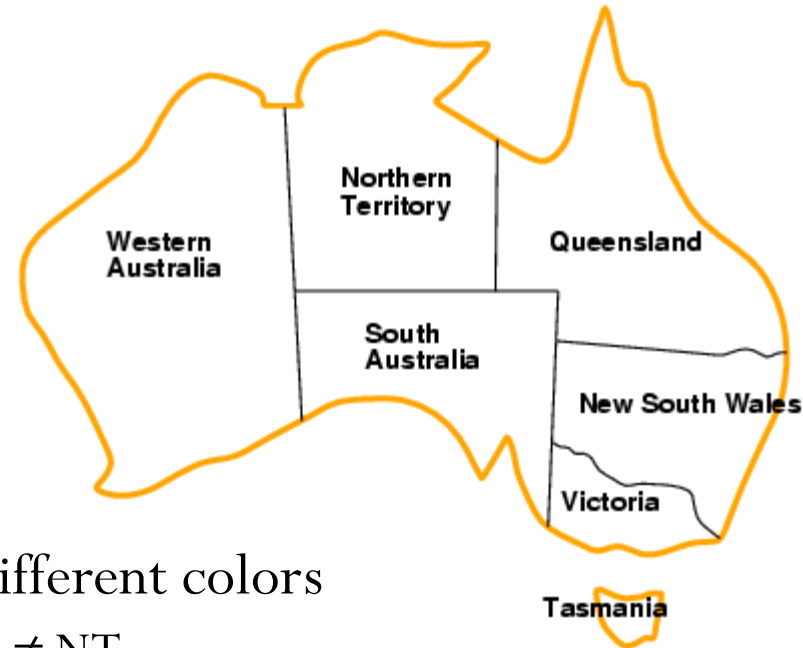
$(x_3 \vee \neg x_4 \vee \neg x_5)$

Suggest a solution!

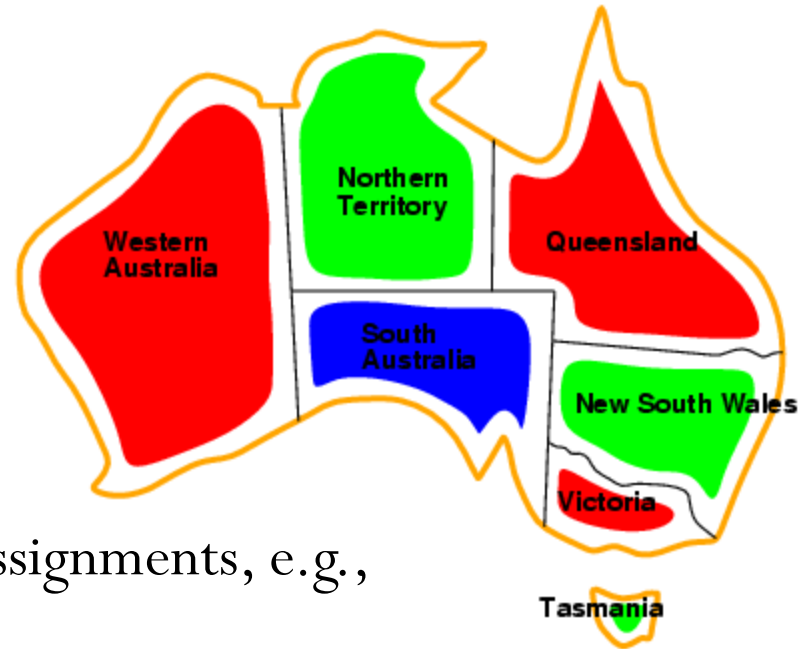
$$(x_1 \vee x_2 \vee x_4) \wedge \\ (x_2 \vee x_4 \vee \neg x_5) \wedge \\ (x_3 \vee \neg x_4 \vee \neg x_5)$$

Example2: Map-Coloring Problem

- **Variables** WA, NT, Q, NSW, V, SA, T
- **Domain** $D_i = \{\text{red, green, blue}\}$
- **Constraints**: adjacent regions must have different colors
 - e.g., $\text{Color}(WA) \neq \text{Color}(NT)$ or in short $WA \neq NT$
 - $(WA, NT) \in \{(\text{red}, \text{green}), (\text{red}, \text{blue}), (\text{green}, \text{red}), (\text{green}, \text{blue}), (\text{blue}, \text{red}), (\text{blue}, \text{green})\}$ OR
 - $(WA, NT) \notin \{(\text{red}, \text{red}), (\text{blue}, \text{blue}), (\text{green}, \text{green})\}$
- Graph Coloring Problem (more general)!



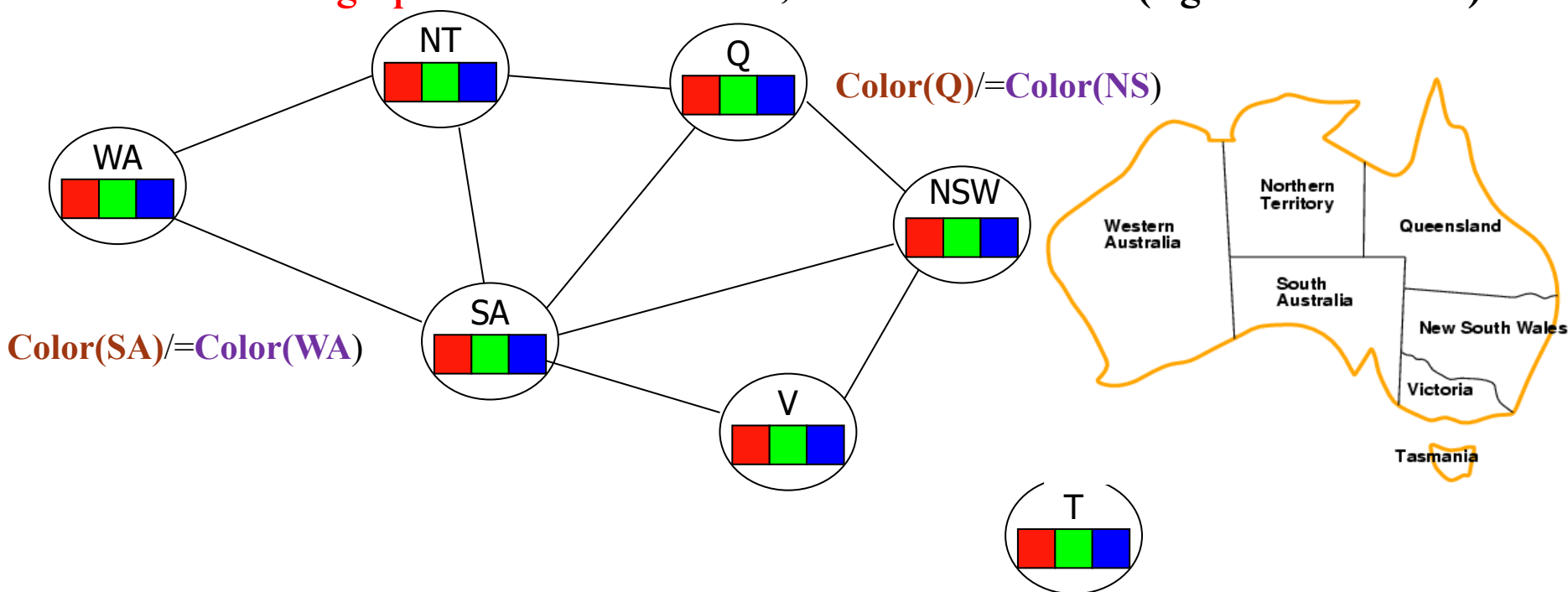
Example: Map-Coloring



- Solutions are **complete** and **consistent** assignments, e.g.,
WA = red, NT = green, Q = red, NSW = green, V = red, SA = blue, T = green
- **Complete**: all are assigned, **consistent**: obeys the constraints.
- A **state** may be incomplete e.g., just WA=red

Constraint graph

- It is helpful to visualize a CSP as a **constraint graph**
 - **Binary CSP:** each constraint relates two variables [here states]
 - **Constraint graph:** nodes are variables, arcs are constraints (e.g. color different)



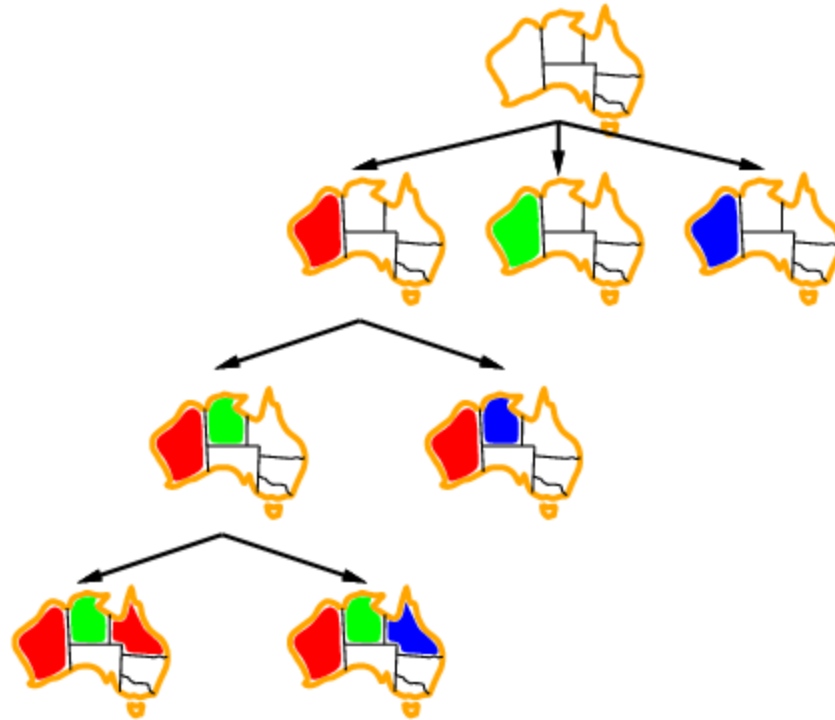
Varieties of CSPs

- Discrete variables
 - finite domains:
 - n variables, domain size d , $O(dn)$ complete assignments
 - e.g., Boolean CSPs, incl. \sim Boolean satisfiability (NP-complete)
 - infinite domains:
 - integers, strings, etc.
 - e.g., job scheduling, variables are start/end days for each job
 - need a constraint language, e.g., $\text{StartJob1} + 5 \leq \text{StartJob3}$
- Continuous variables
 - e.g., Time: start/end times for Hubble Space Telescope observations
 - linear constraints solvable in polynomial time by linear programming

CSP as Incremental Search Problem

- initial state
 - all (or at least some) variables UNassigned
- successor function
 - assign a value to an UNassigned variable
 - must not conflict with previously assigned variables
- goal test
 - all variables have values assigned
 - no conflicts exist (in the assignments)
- path cost
 - e.g. constant for each step [some colors may be expensive]
 - may be problem-specific

Example



CSPs and Search

In principle, any search algorithm can be used to solve a CSP, but:

- awful branching factor
 - $n \cdot d$ for n variables with d values at the top level, $(n-1) \cdot d$ at the next level, etc.
- not very efficient, since they neglect some CSP properties
 - commutativity: the order in which values are assigned to variables is irrelevant, since the outcome is the same

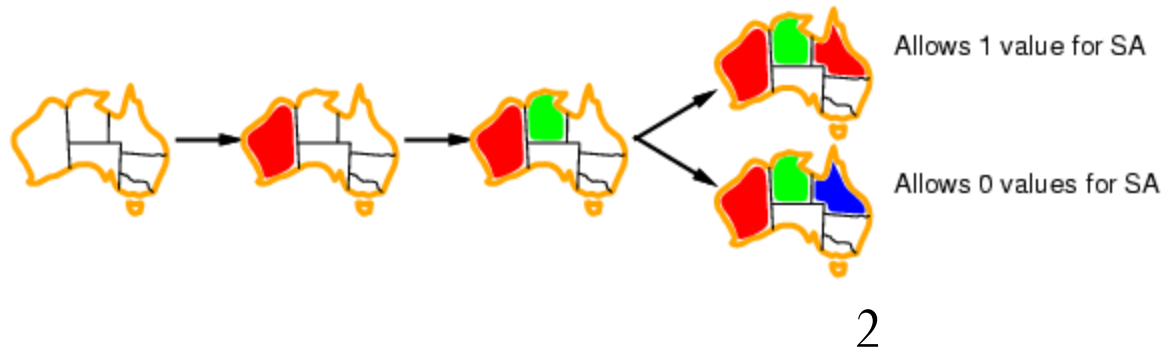
Backtracking Search for CSPs

A variation of depth-first search that is often used for CSPs

- values are chosen for one variable at a time
- if no legal values are left, the algorithm **backs up** and changes a **previous assignment**
- very easy to implement
 - initial state, successor function, goal test are standardized
- not very efficient
 - can be improved by trying to select more **suitable unassigned** variables **first**

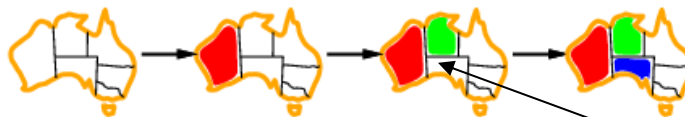
Improving backtracking efficiency

- **General-purpose** methods can give huge gains in speed:
 1. Which variable should be assigned next? $\{WA, NT, Q, NSW, V, SA, T\}$
 2. In what order should its values be tried? $[R, B, G], [R, G, B], \dots$
 3. Can we detect inevitable failure early? Case 2 below



Heuristics for CSP

1. most-constrained variable (Minimum Remaining Values: **MRV**, “fail-first”)



- variable with the **fewest** possible values is selected
- tends to minimize the branching factor

2. most-constraining variable **MCV**

- variable with the **largest** number of constraints on other unassigned variables



3. least-constraining value **LCV**

- for a selected variable, choose the **value** that leaves more freedom for future choices



Most constrained variable

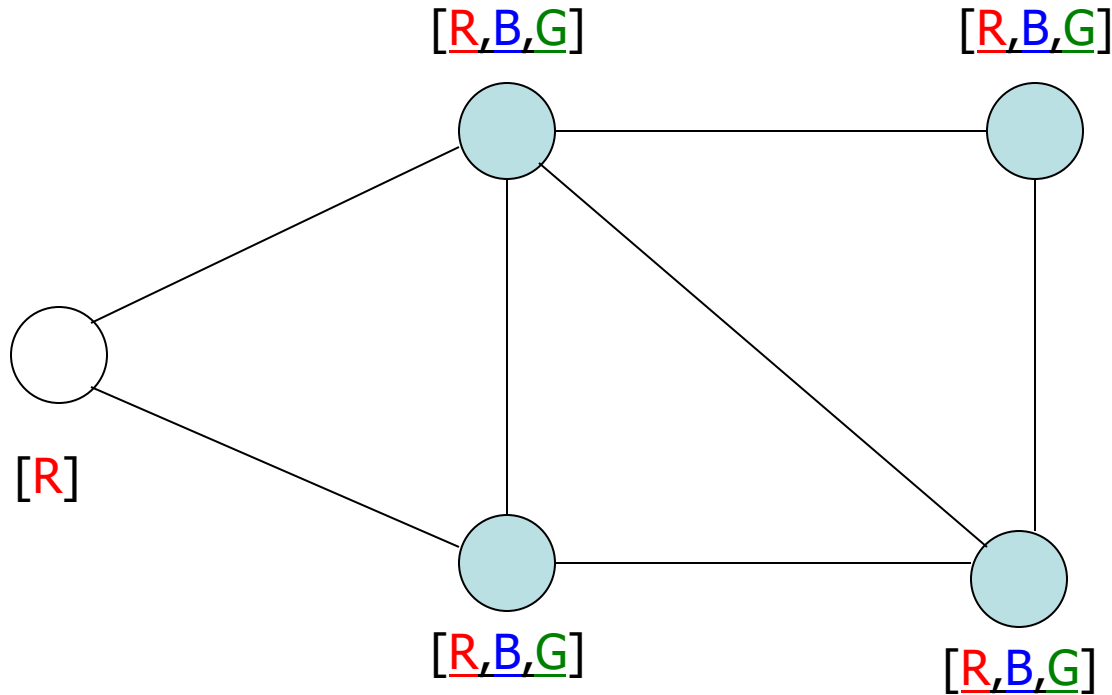
Minimum Remaining Values (MRV)

- Most constrained variable:
choose the variable with the **fewest legal values**

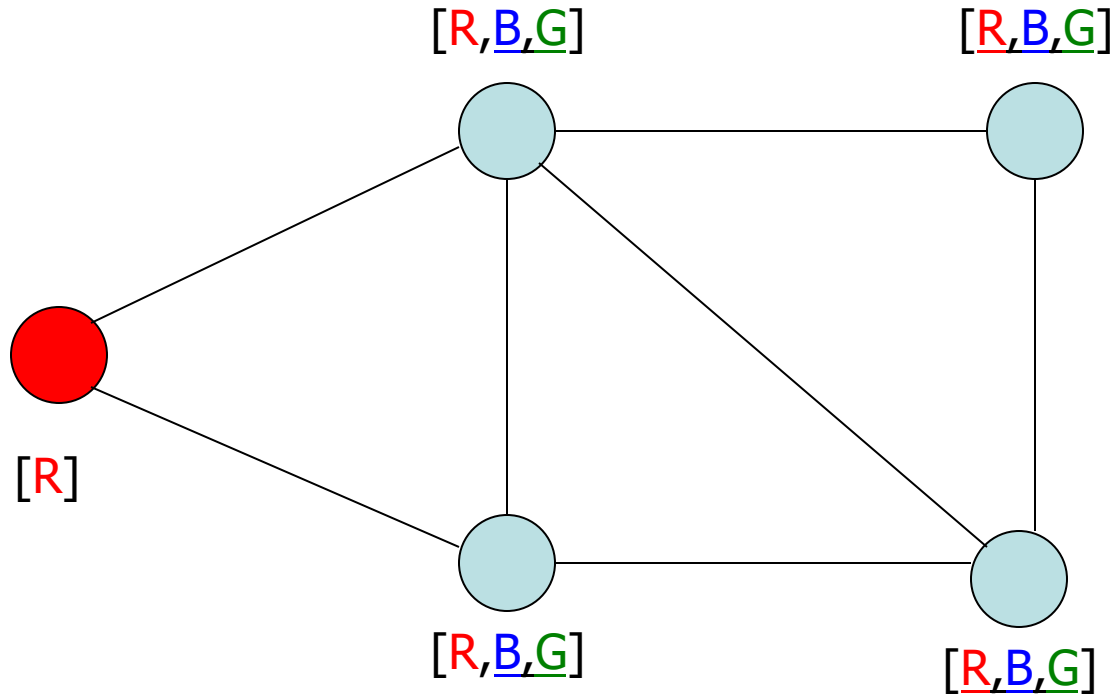


- Called **minimum remaining values (MRV)** heuristic
- “fail-first” heuristic: Picks a variable which will cause failure as soon as possible, allowing the tree to be pruned.

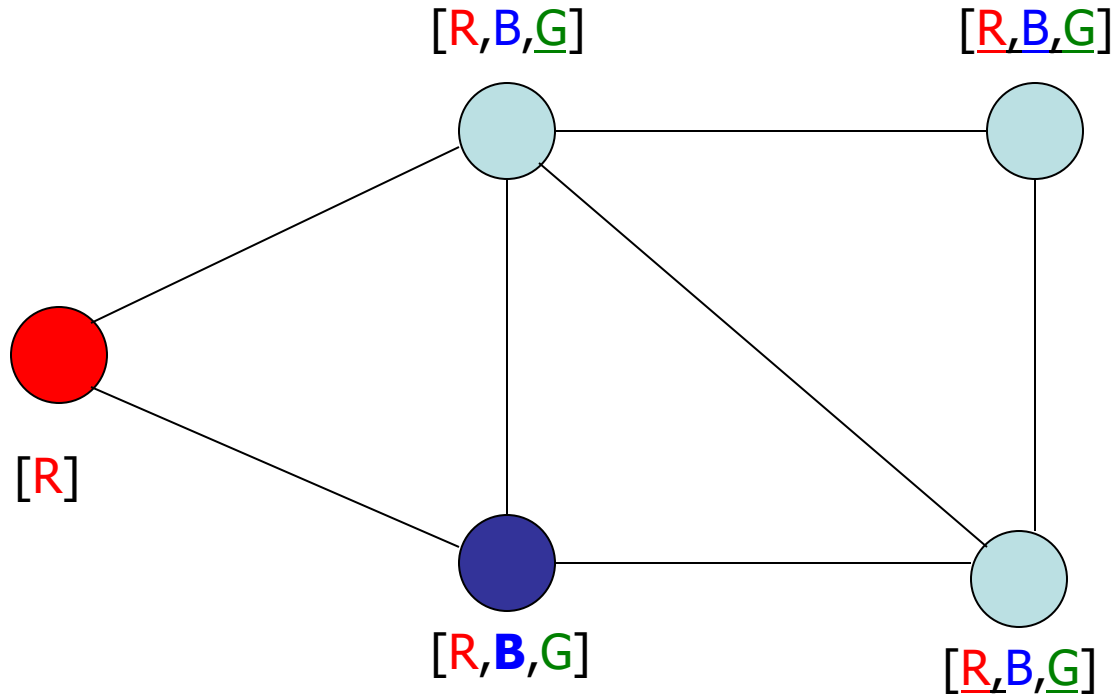
Backpropagation - MRV



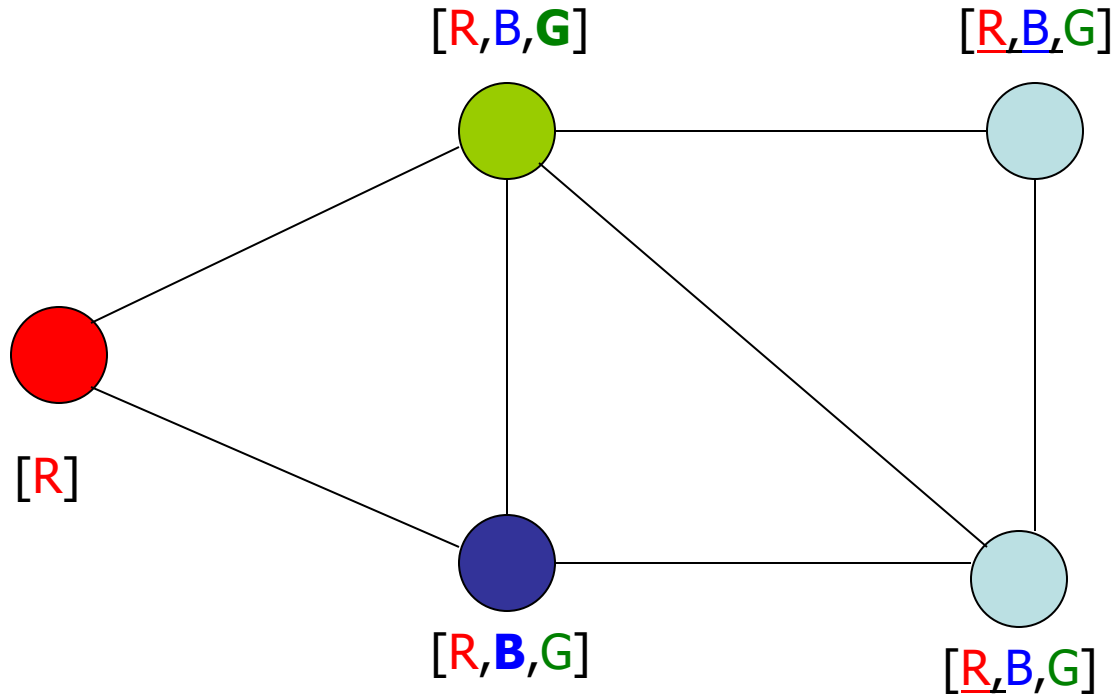
Backpropagation – MRV minimum remaining values



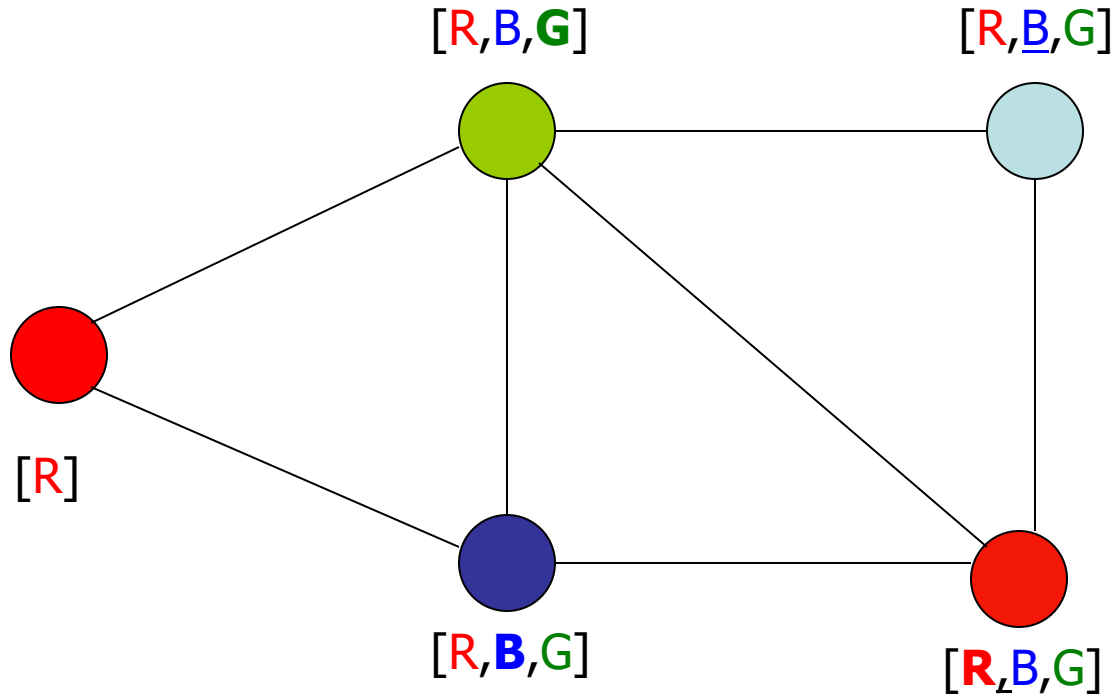
Backpropagation - MRV minimum remaining values



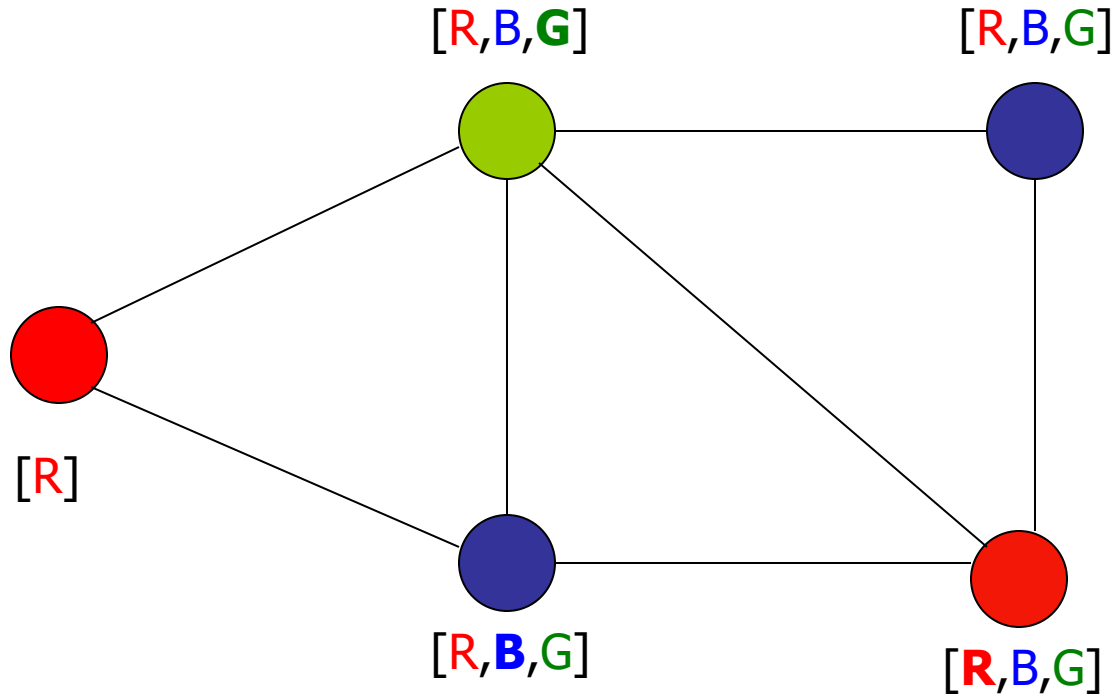
Backpropagation - MRV



Backpropagation - MRV



Backpropagation - MRV minimum remaining values



Solution !!!

Most constraining variable - MCV

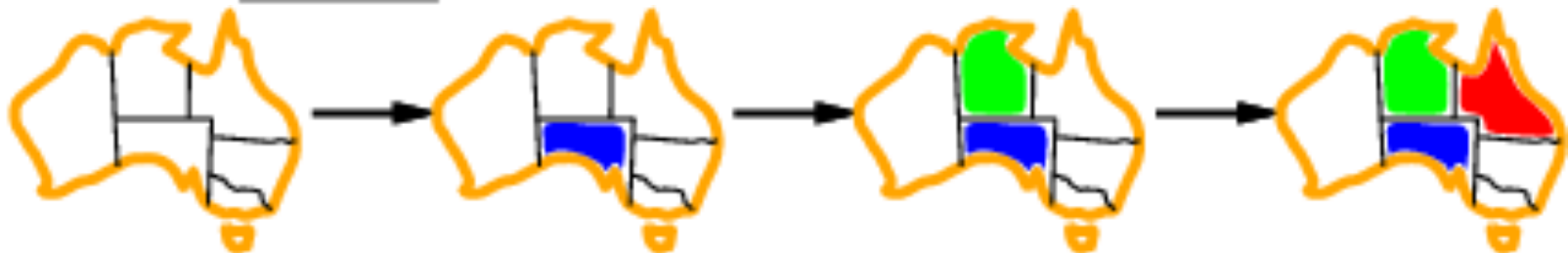
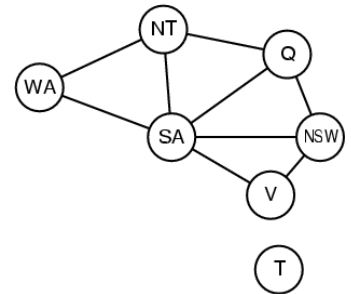
- **Tie-breaker** among most constrained variables (**MRV**)

- Most constraining variable:

- choose the variable **with the most constraints on remaining variables** (select variable that is involved in the largest number of constraints - edges in graph on other unassigned variables: **SA:5**, WA:2, NT:3, Q:3, NSW:3, V:2 then:

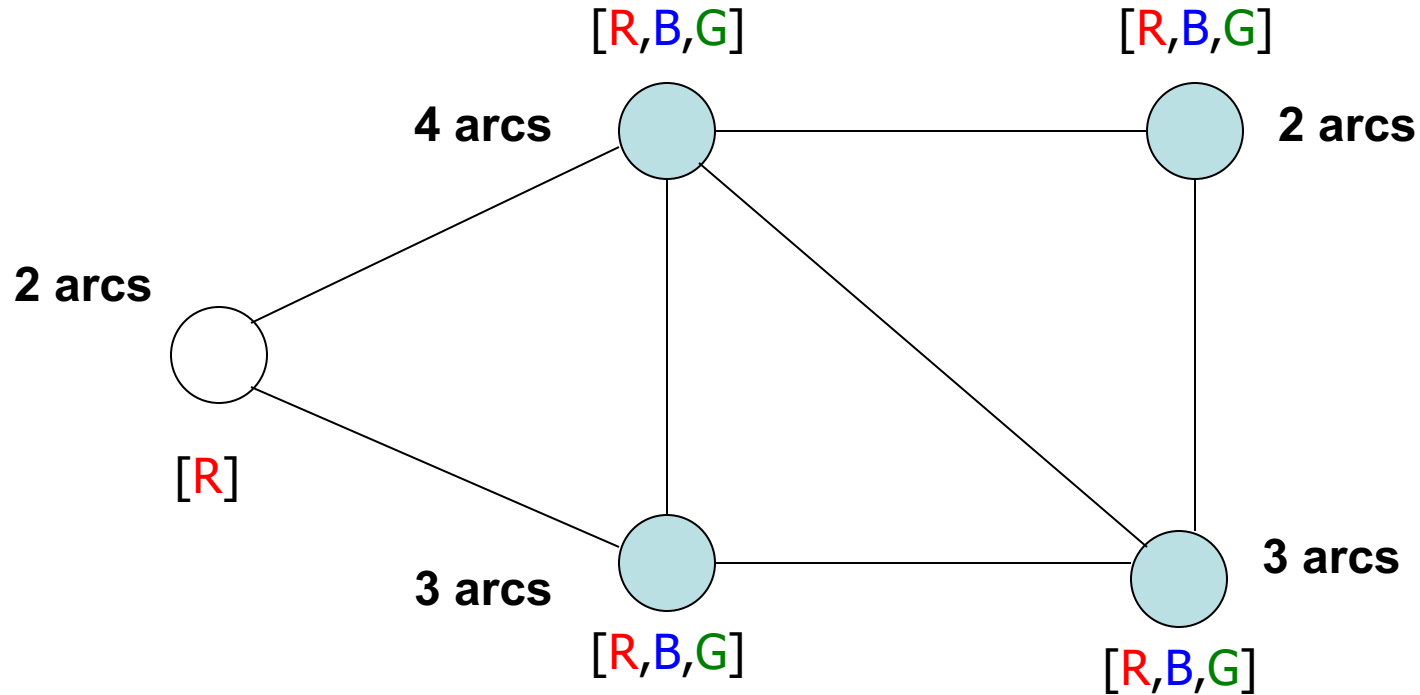
- WA:1, **NT:2**, Q:2, NSW:2, V:1 then

- **Q:1**, NSW:2, V:1 then WA:0, **NT:1**, **V:1** ?? Which?



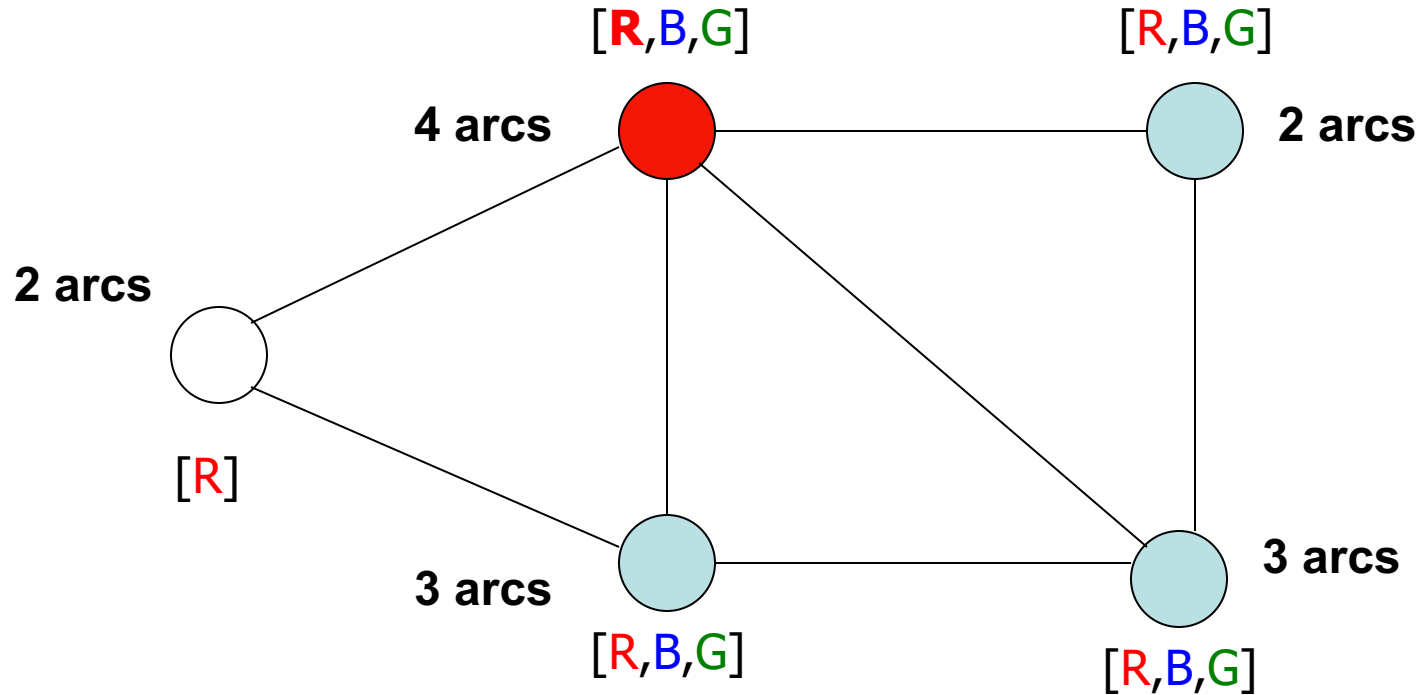
Backpropagation - MCV

Most constraining variable



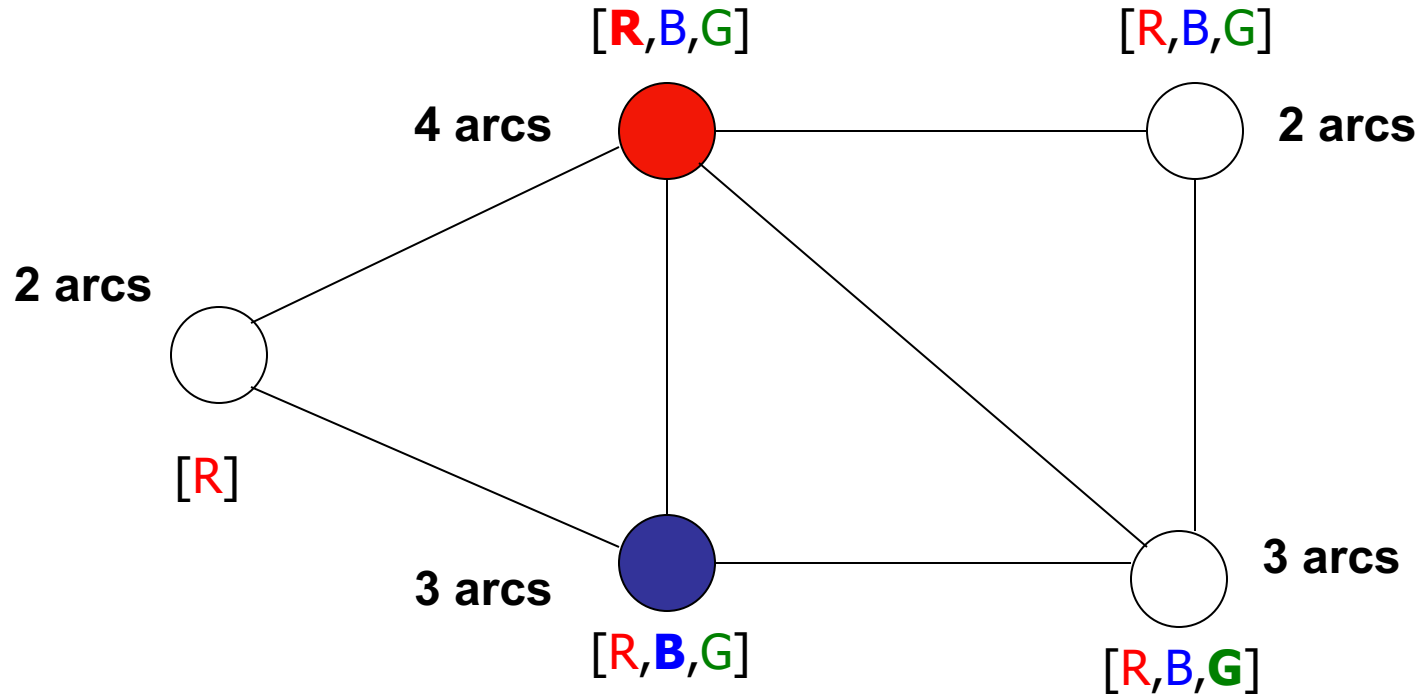
Backpropagation – MCV

Most constraining variable



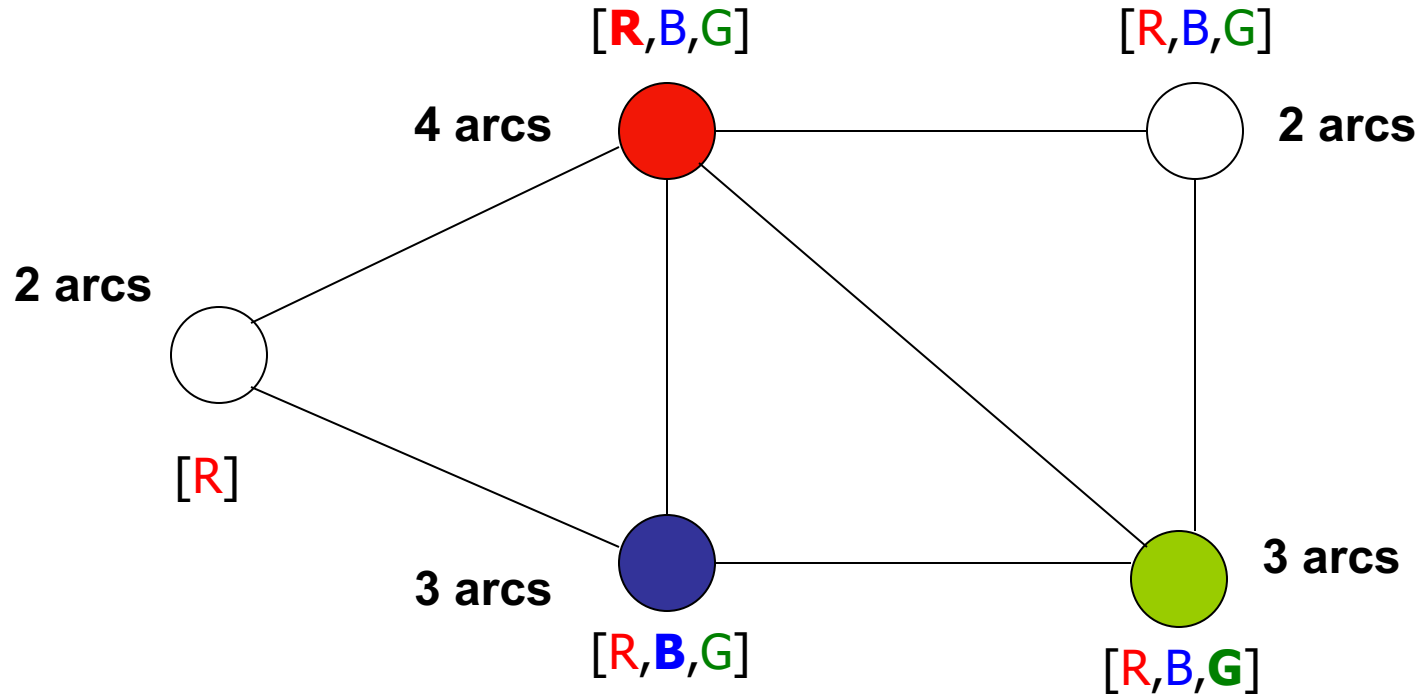
Backpropagation – MCV

Most constraining variable



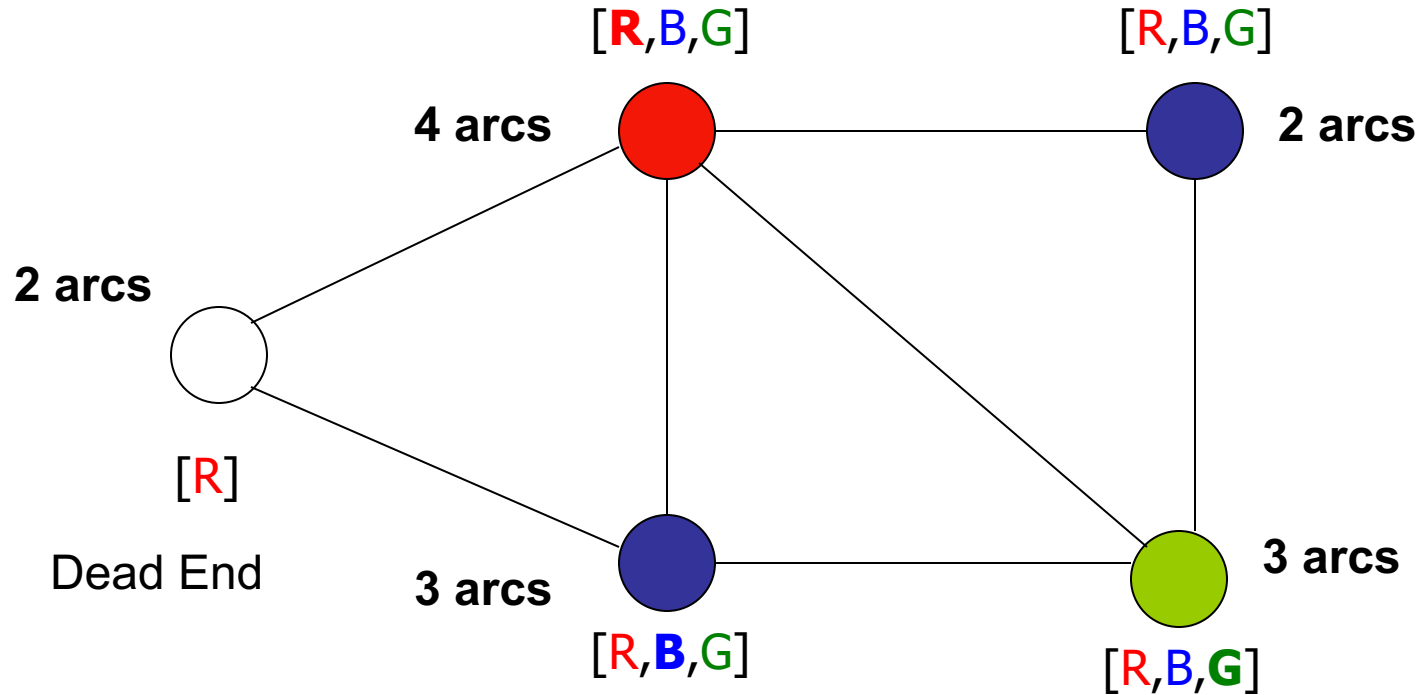
Backpropagation - MCV

Most constraining variable



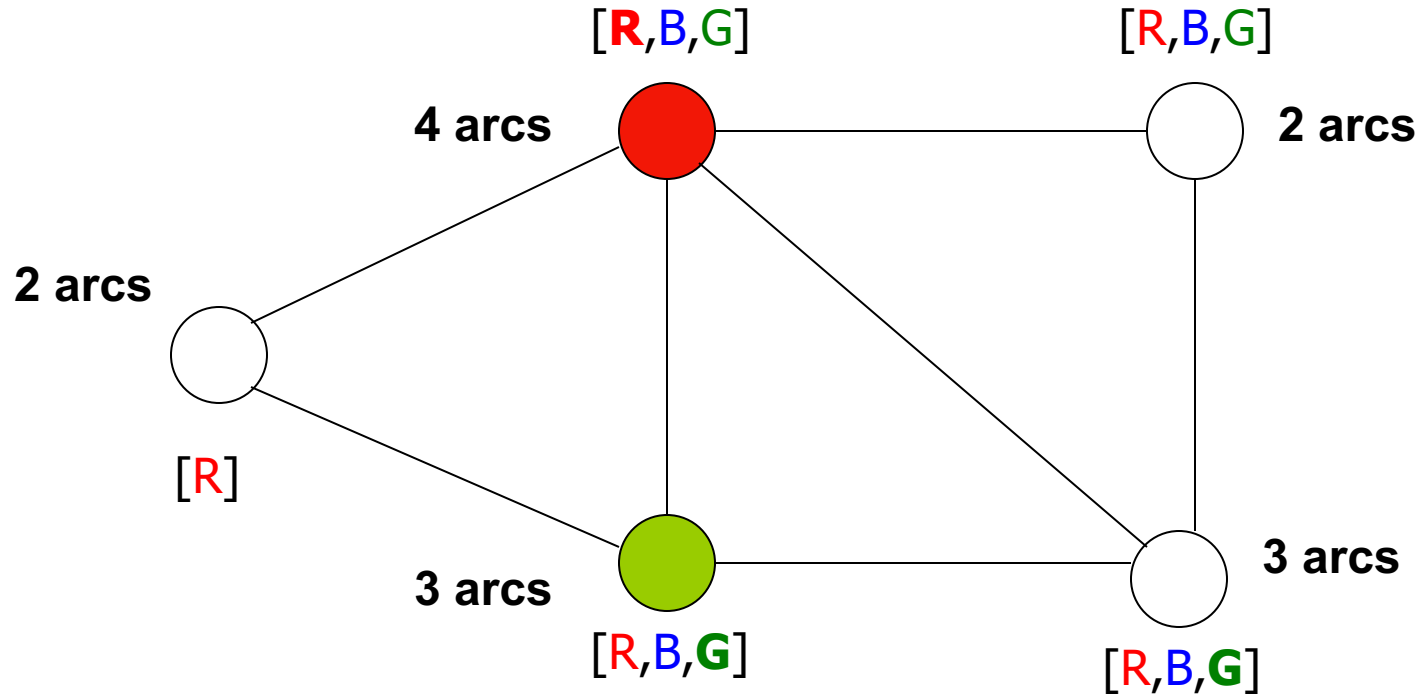
Backpropagation - MCV

Most constraining variable



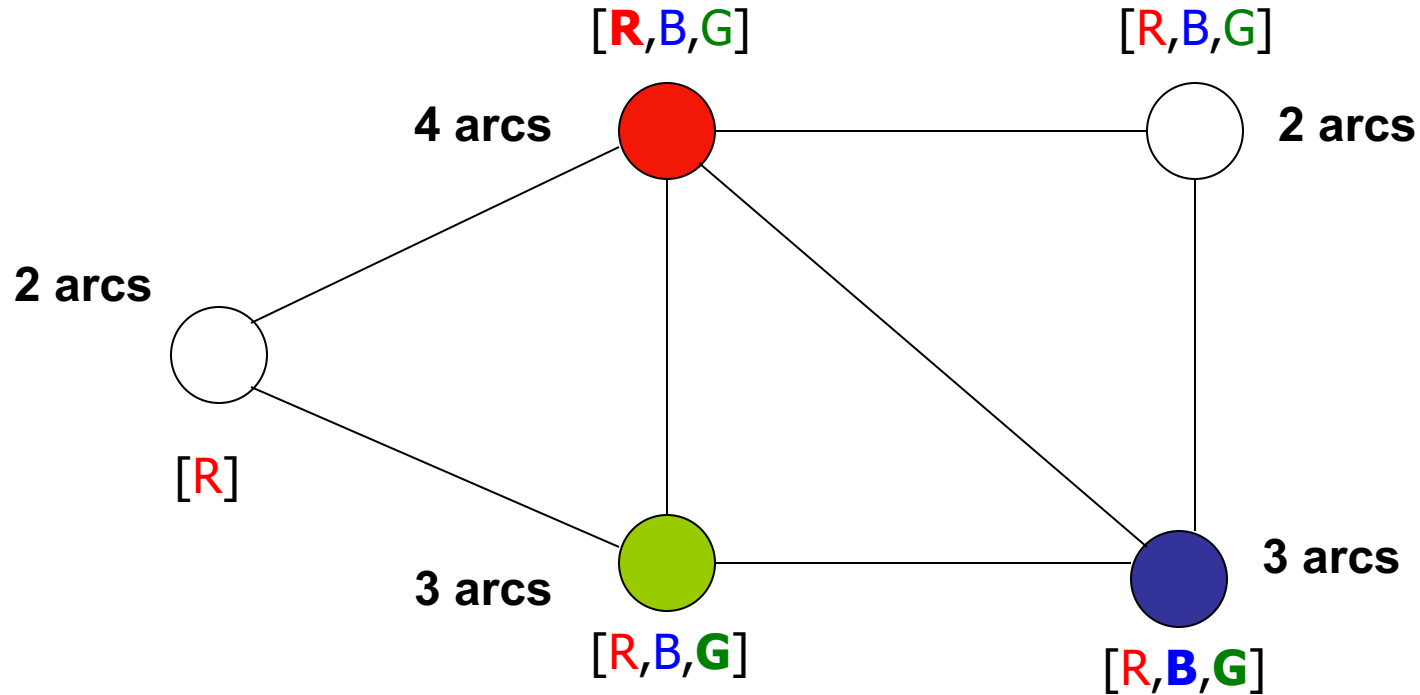
Backpropagation - MCV

Most constraining variable



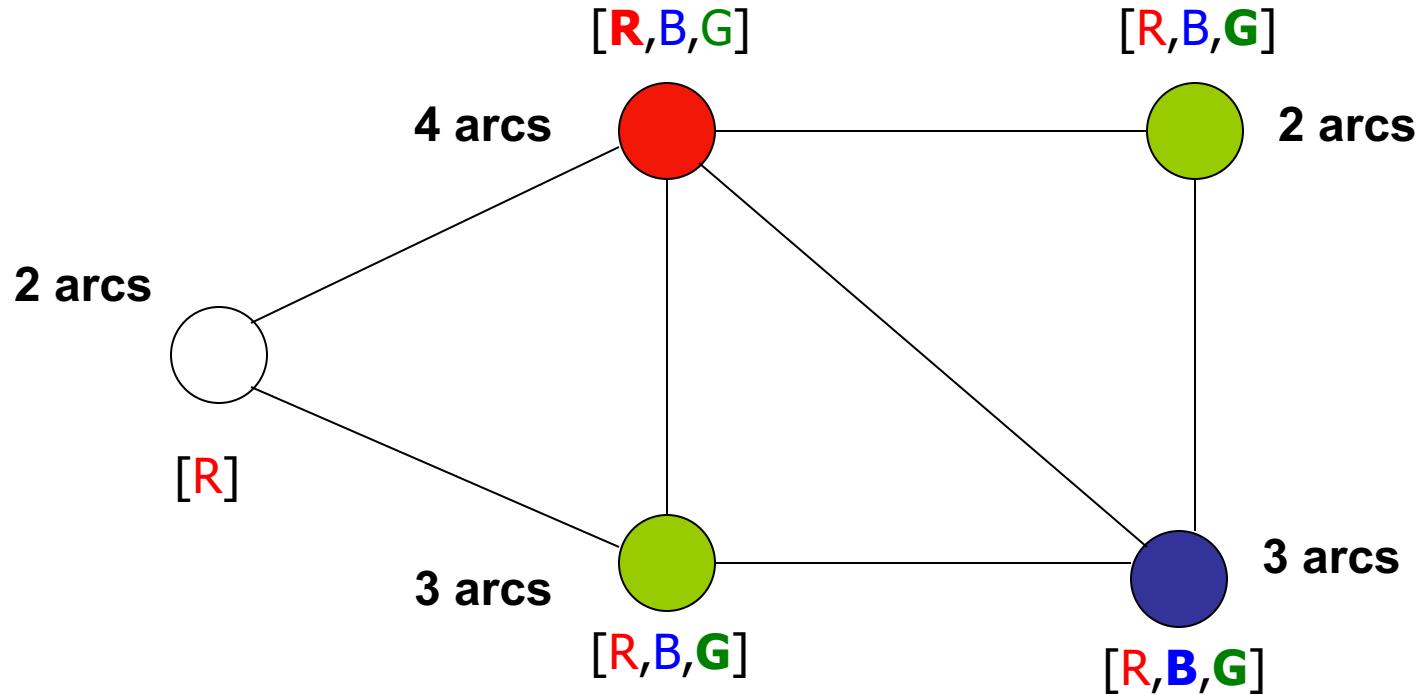
Backpropagation - MCV

Most constraining variable



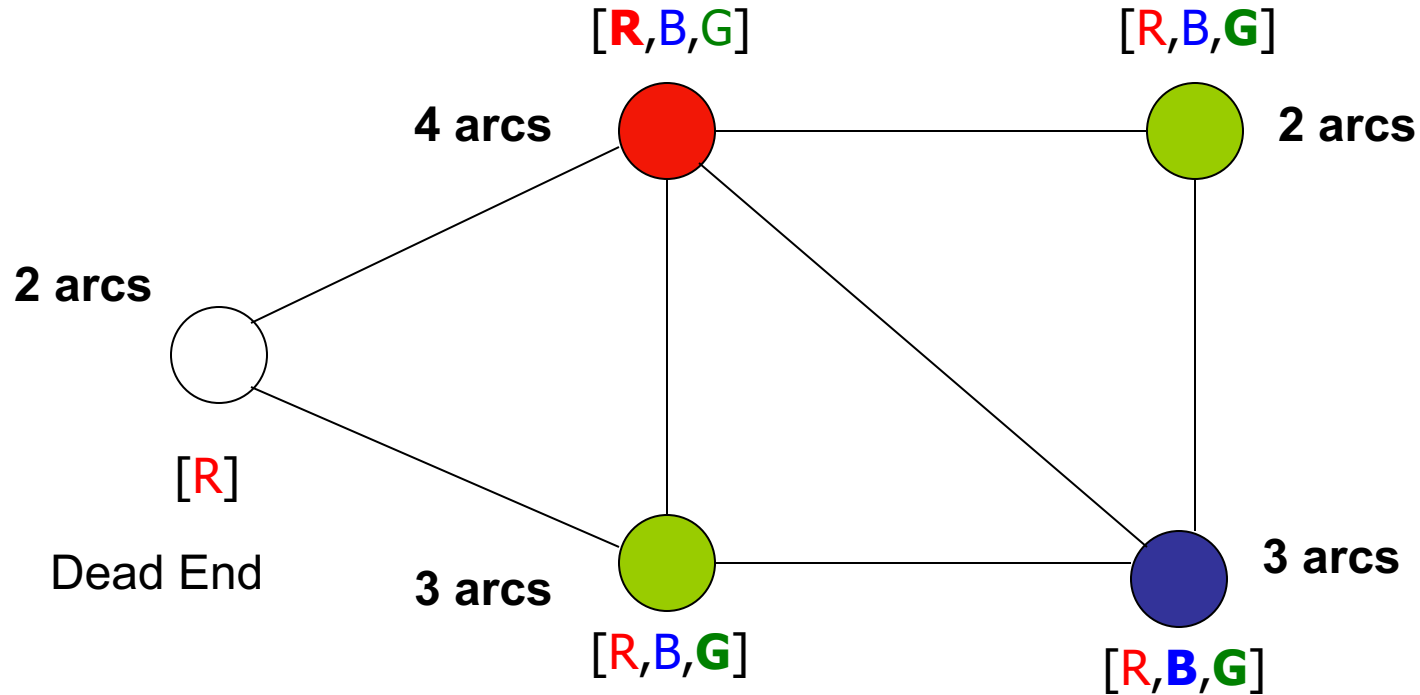
Backpropagation – MCV

Most constraining variable



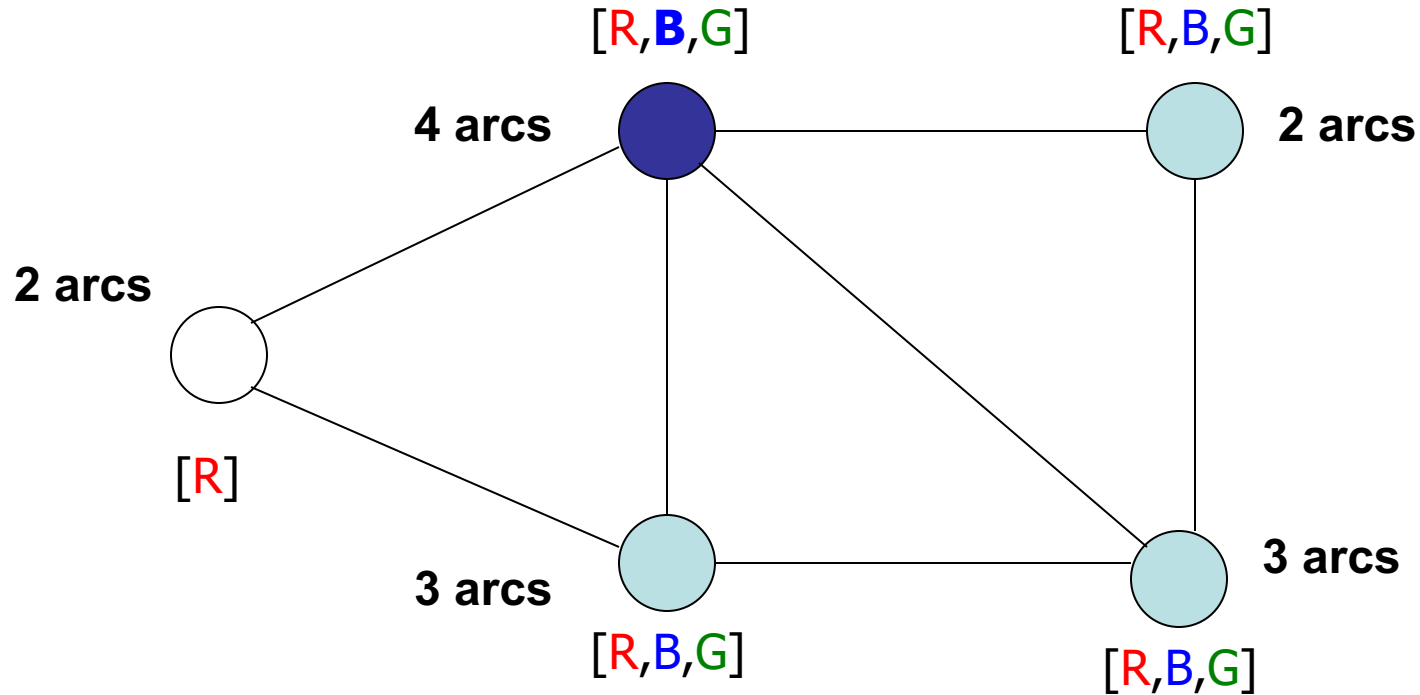
Backpropagation - MCV

Most constraining variable



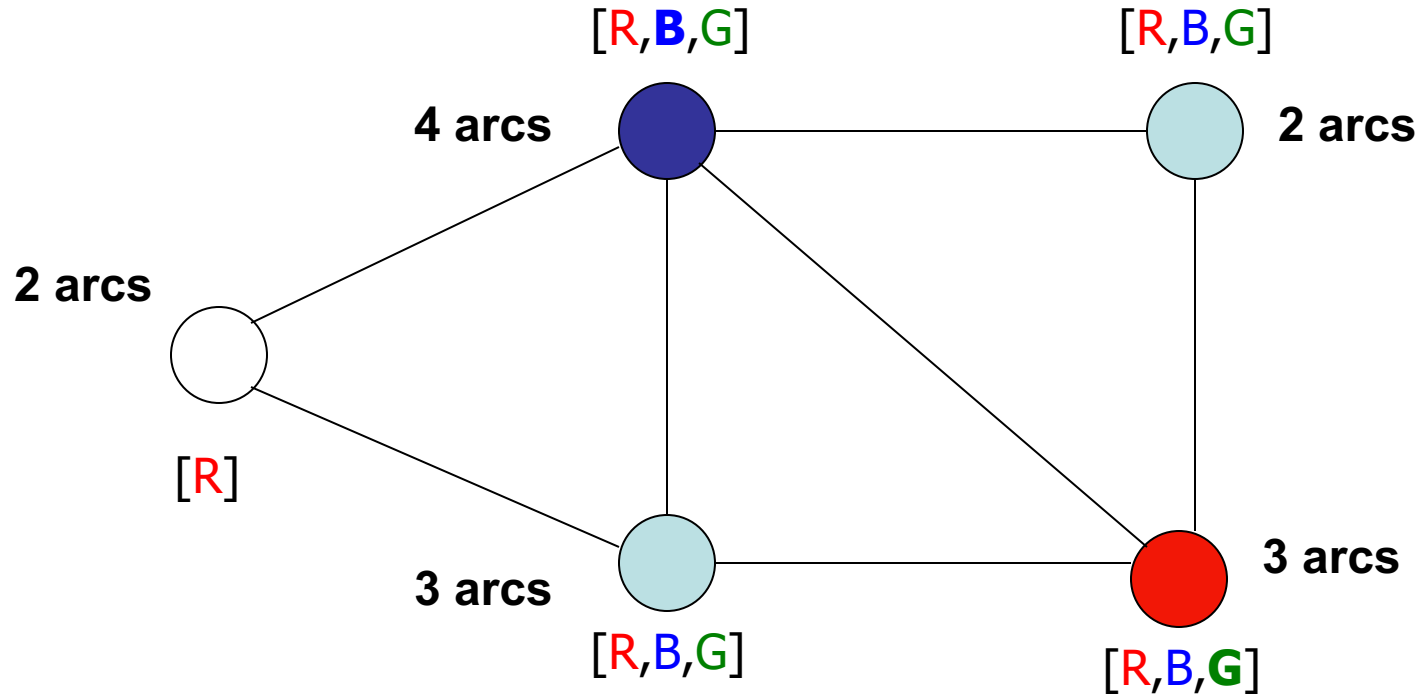
Backpropagation - MCV

Most constraining variable



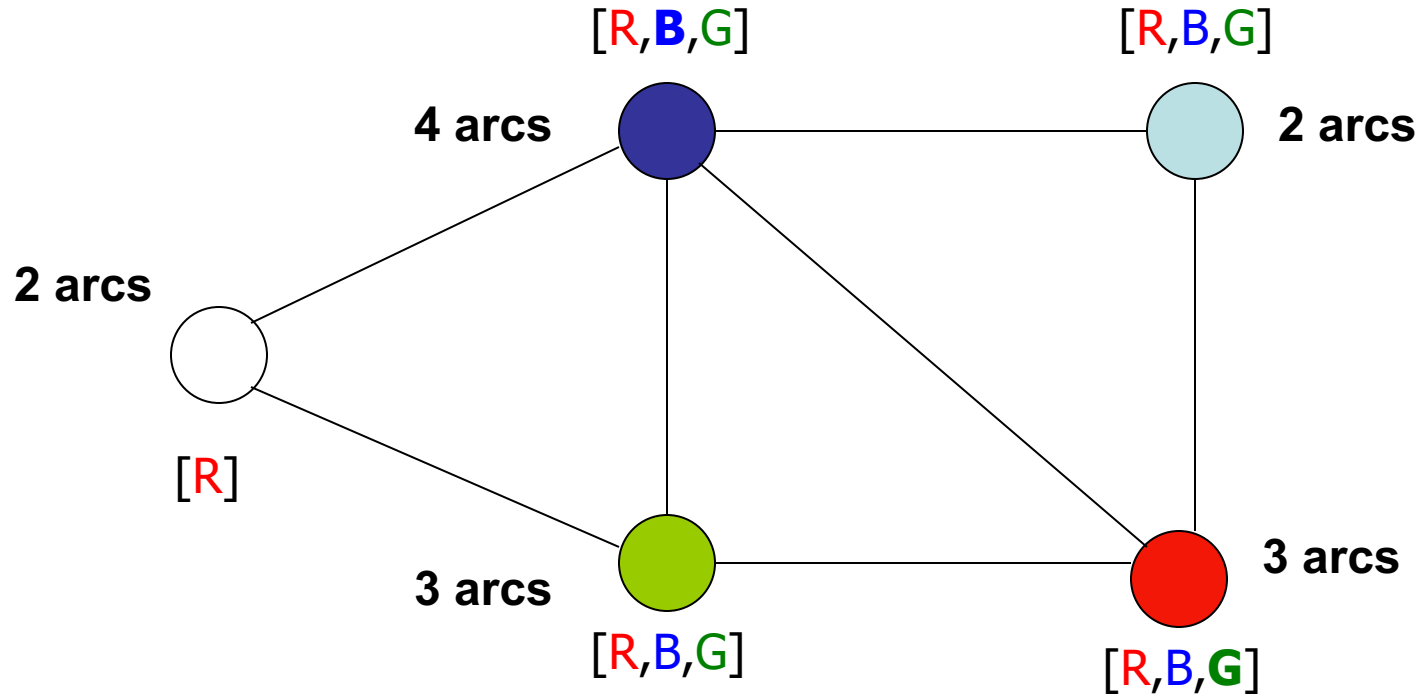
Backpropagation - MCV

Most constraining variable



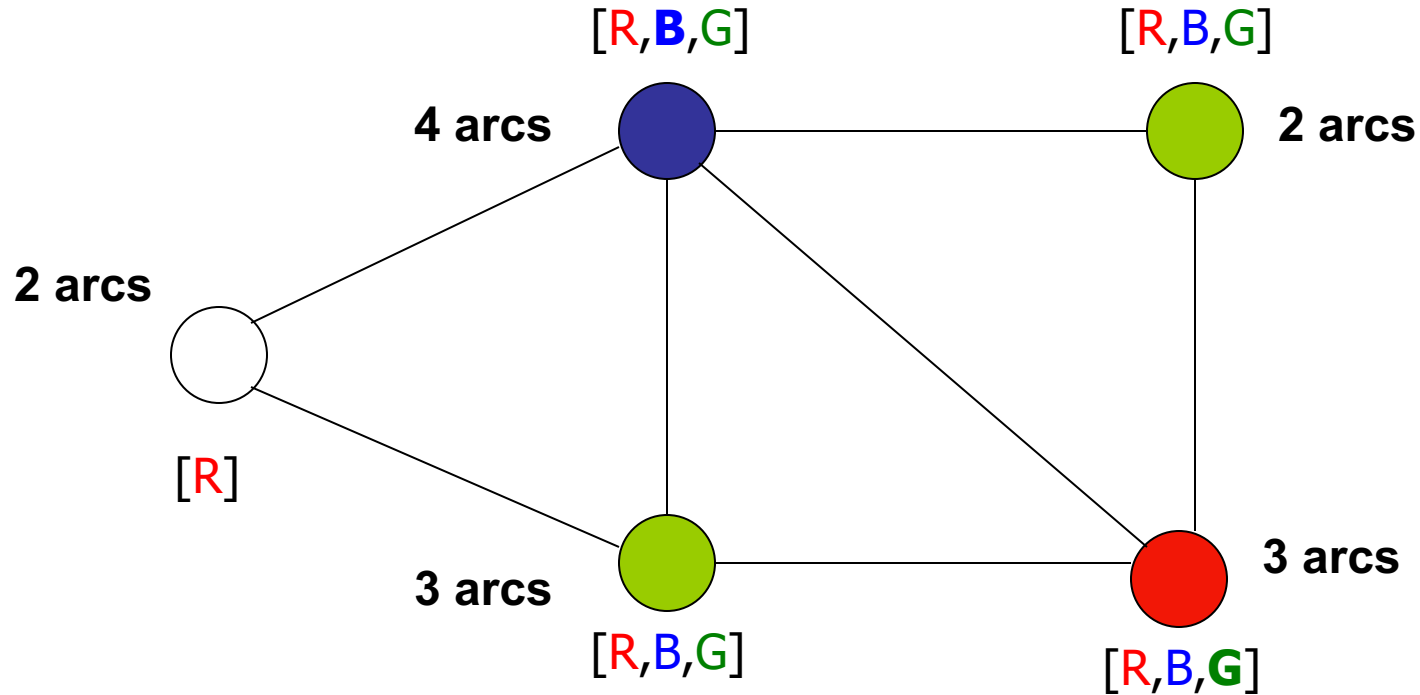
Backpropagation - MCV

Most constraining variable



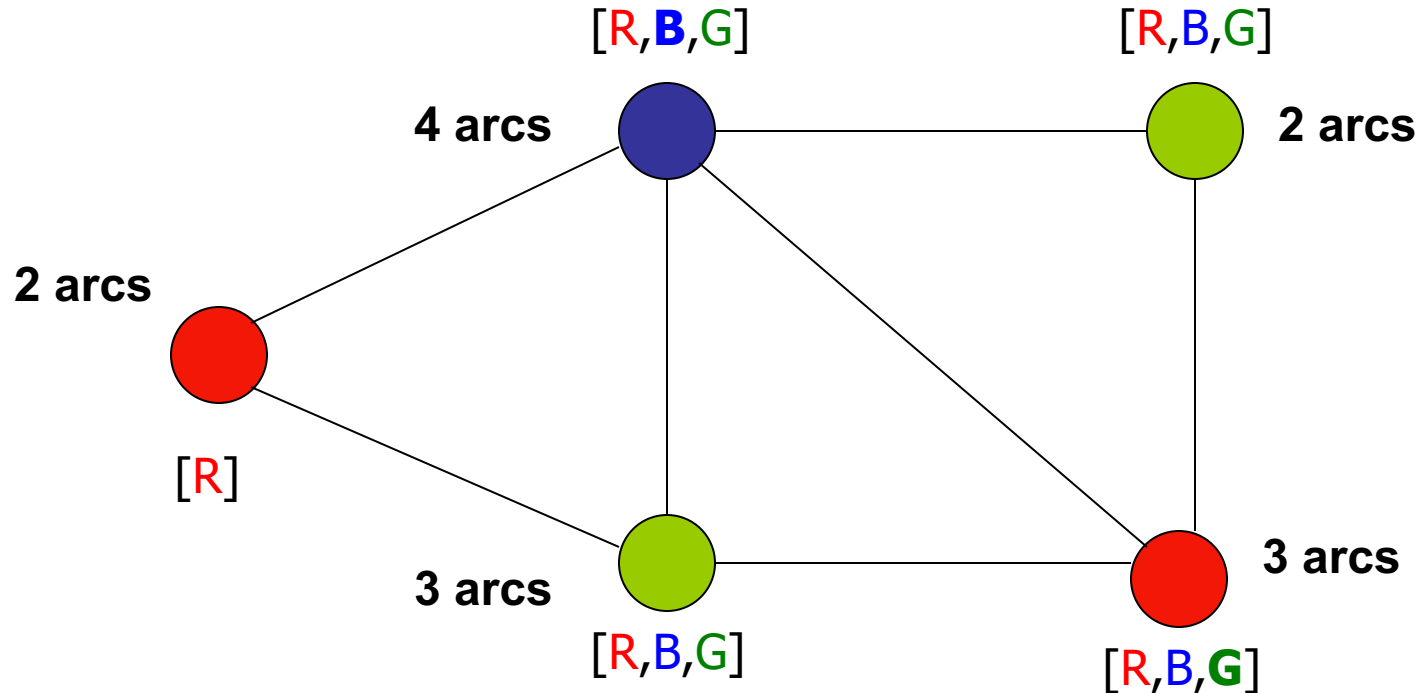
Backpropagation - MCV

Most constraining variable



Backpropagation - MCV

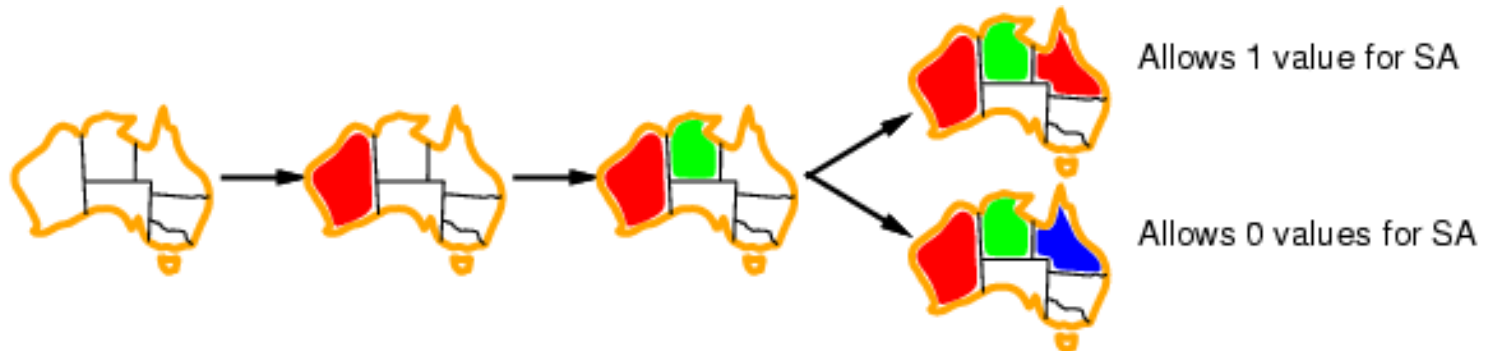
Most constraining variable



Solution !!!

Least constraining value - LCV

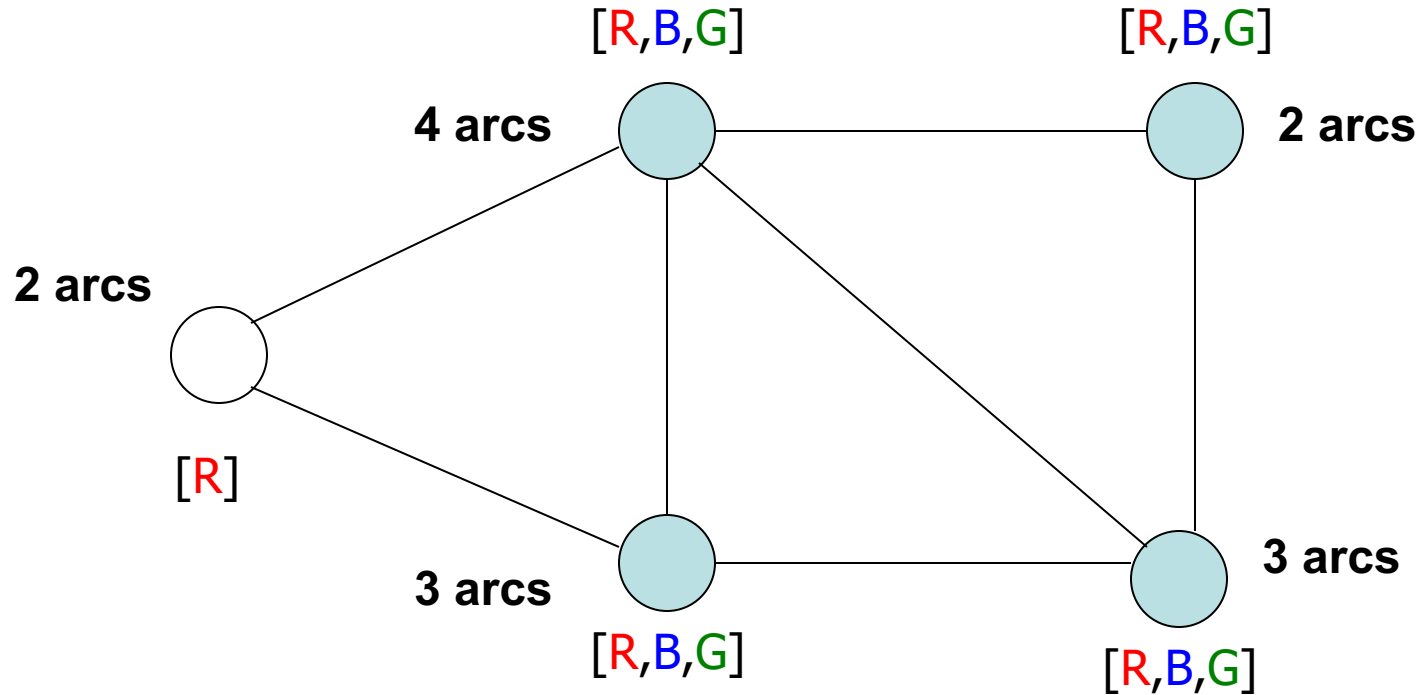
- Given a variable, choose the least constraining value:
 - the one that rules out/eliminates the fewest values in the remaining variables (keeps the most)



- Combining these heuristics makes 1000 queens feasible

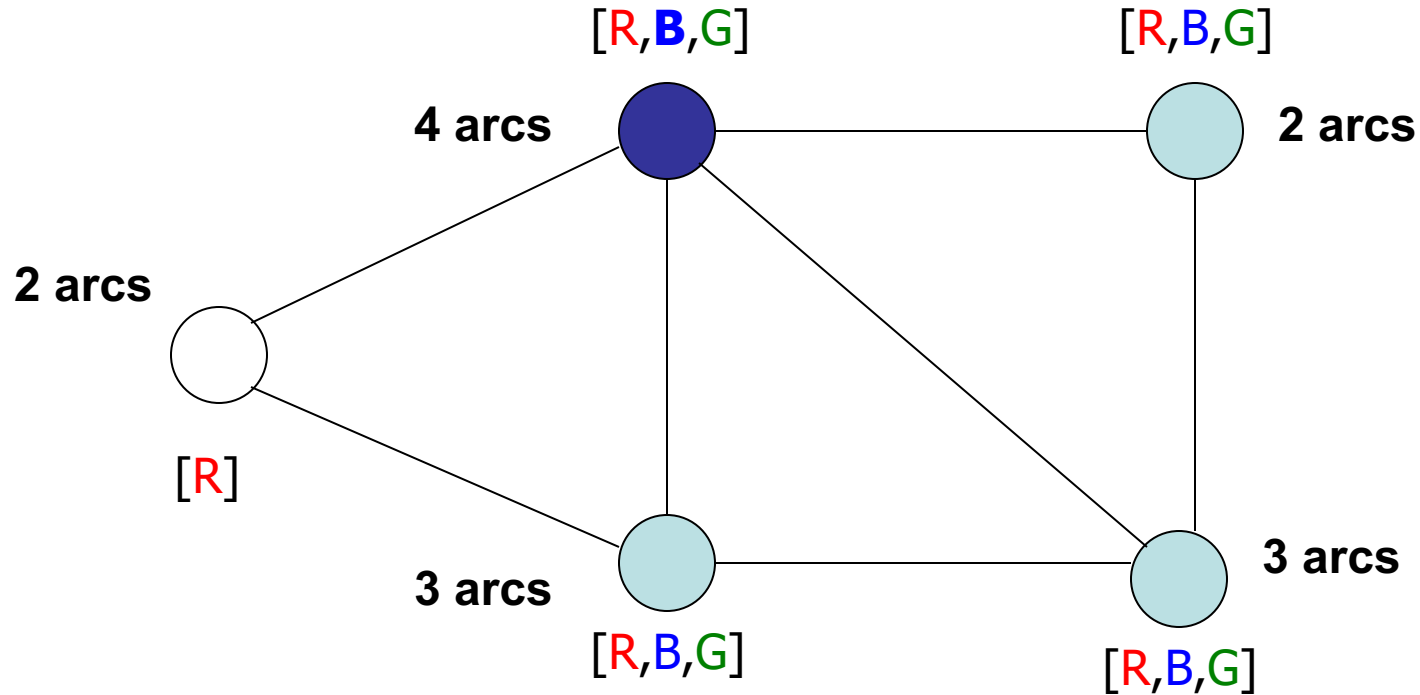
Backpropagation – LCV

Least constraining value



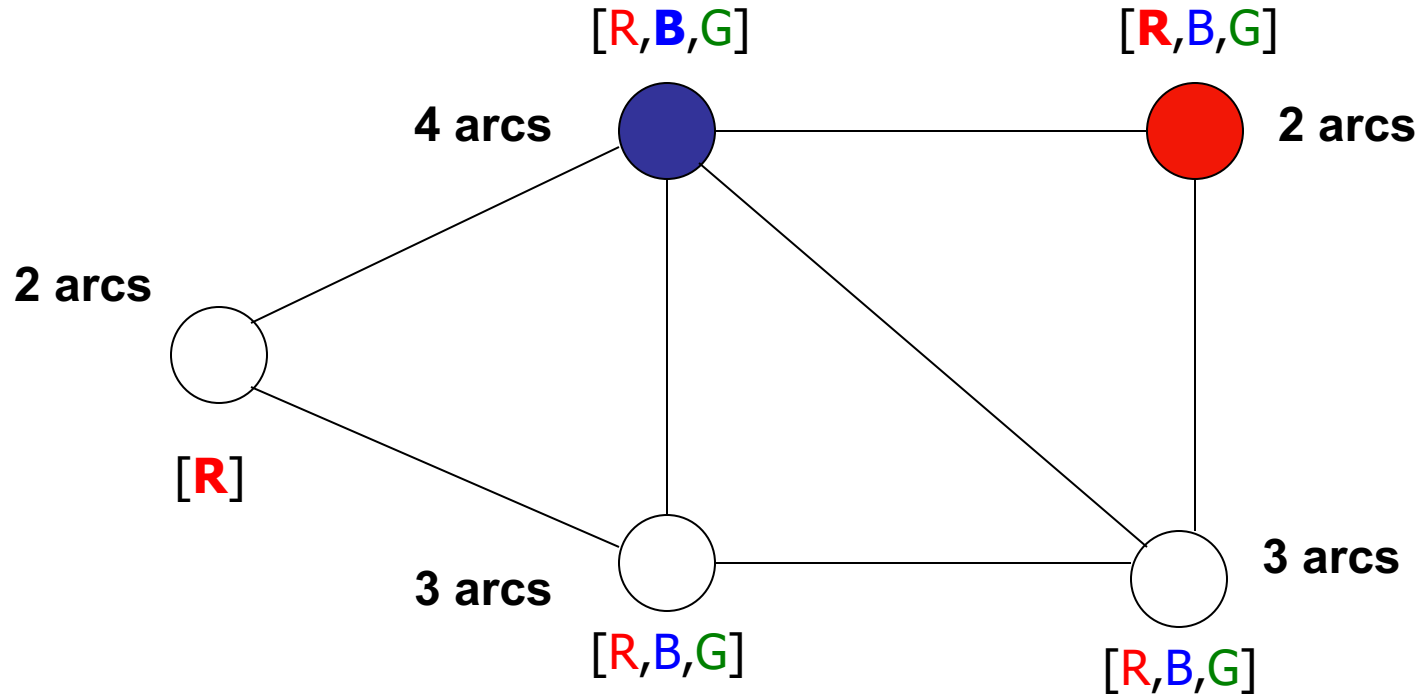
Backpropagation – LCV

Least constraining value



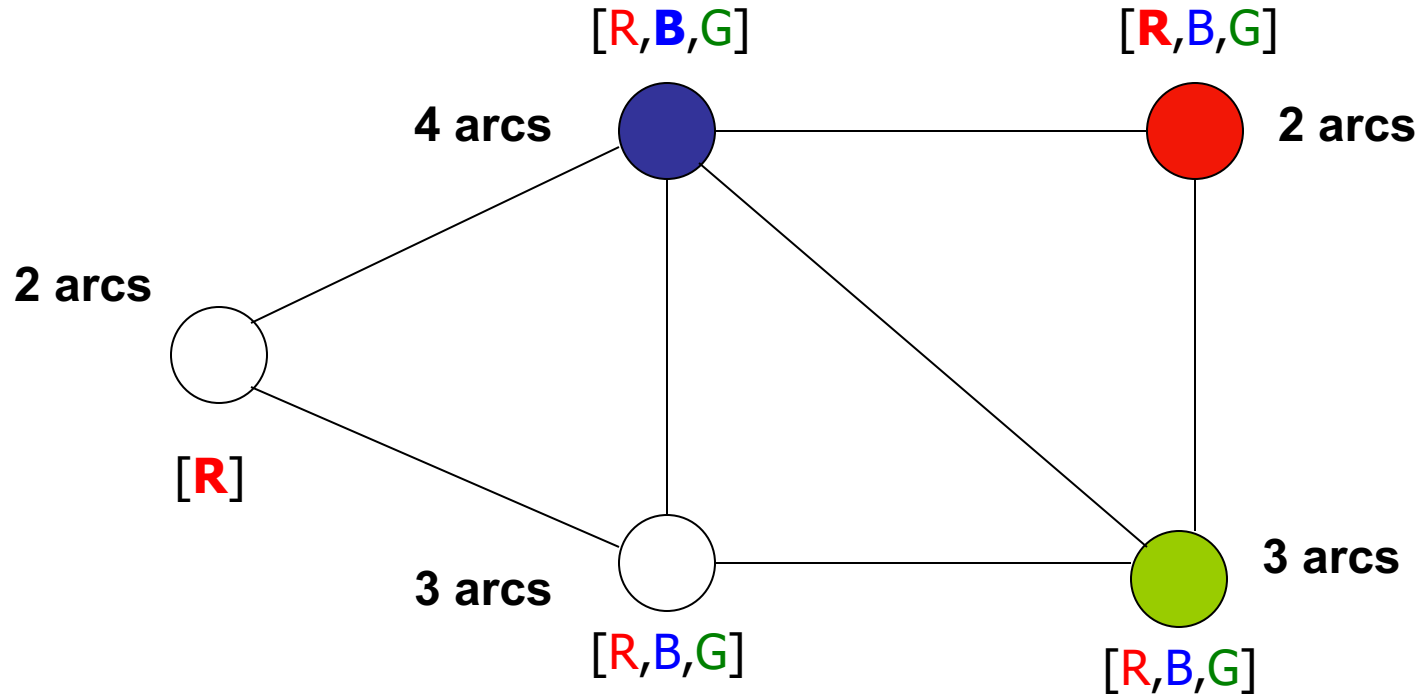
Backpropagation - LCV

Least constraining value



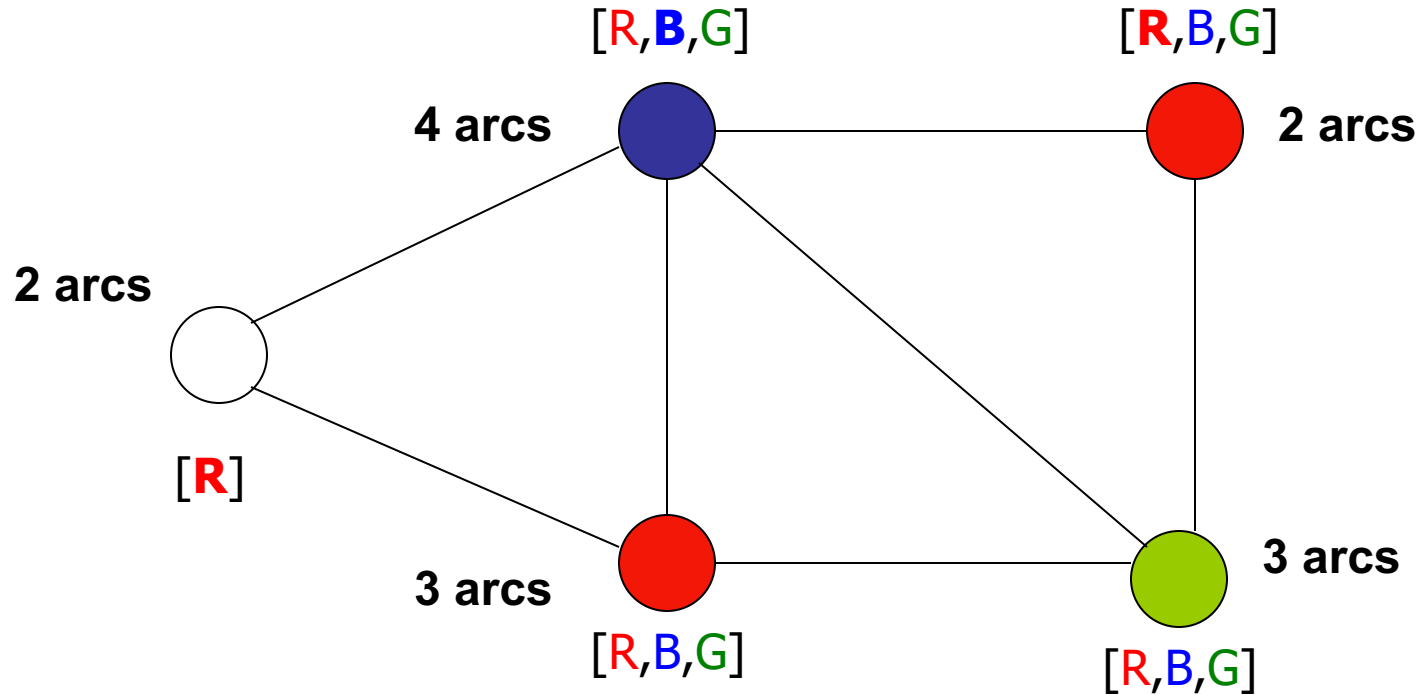
Backpropagation - LCV

Least constraining value



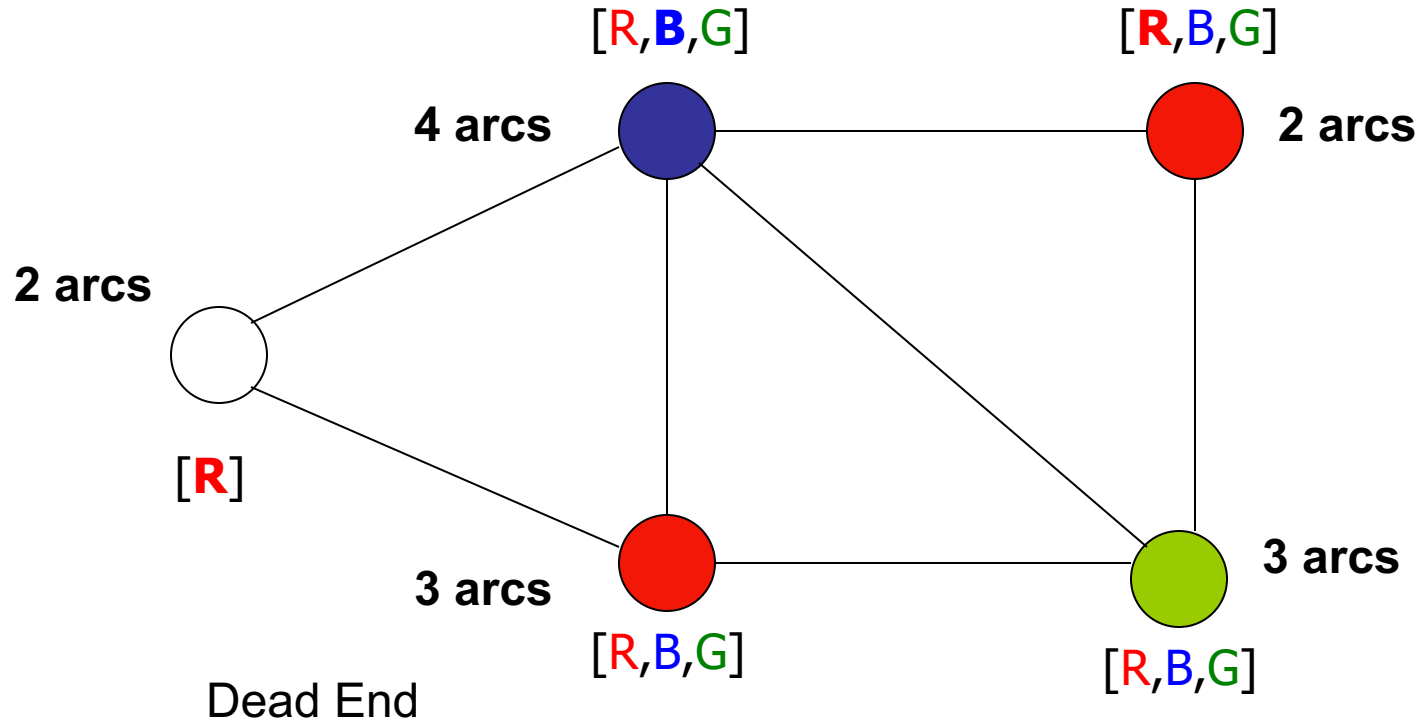
Backpropagation - LCV

Least constraining value



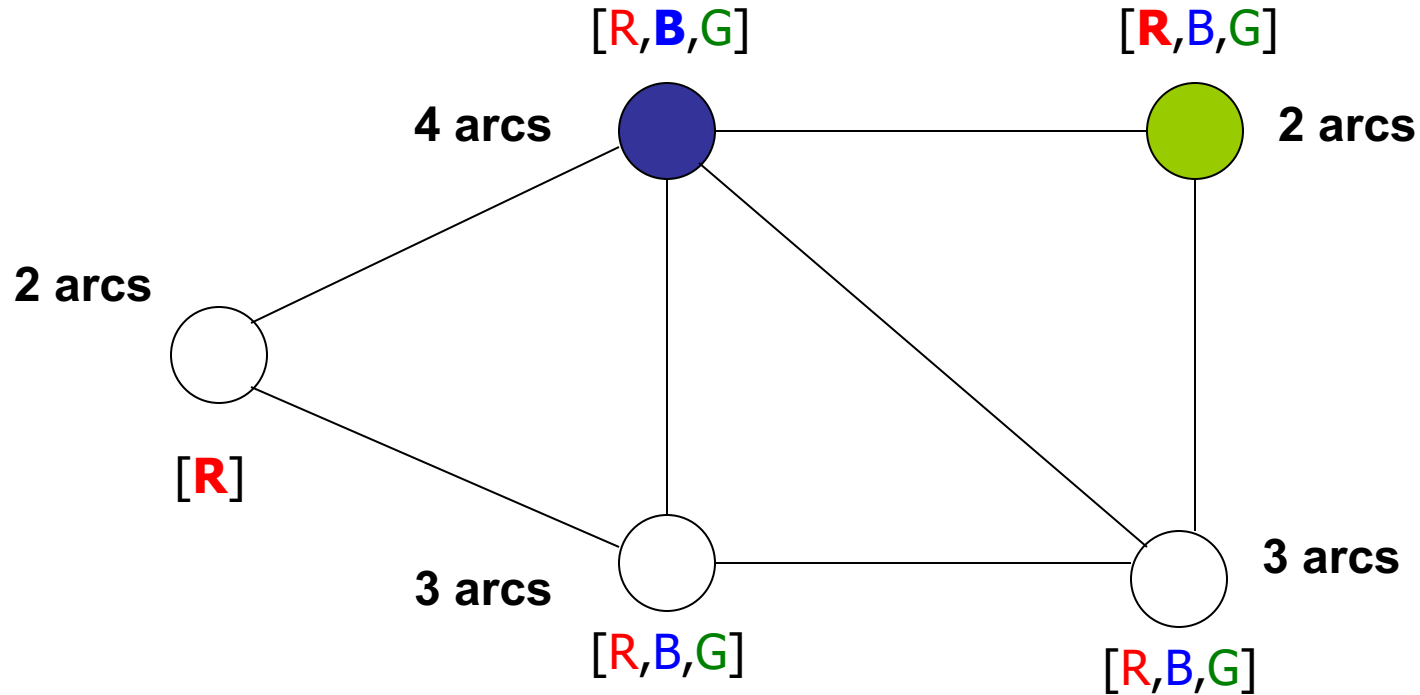
Backpropagation - LCV

Least constraining value



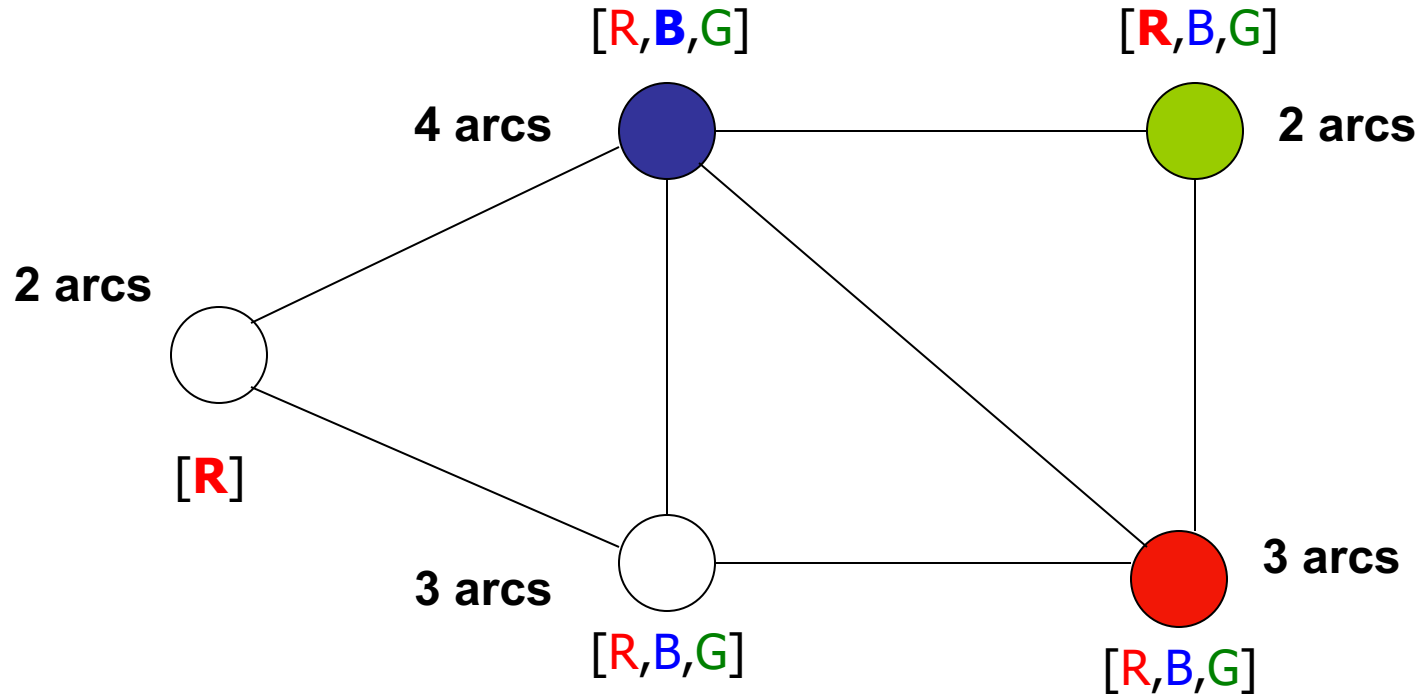
Backpropagation - LCV

Least constraining value



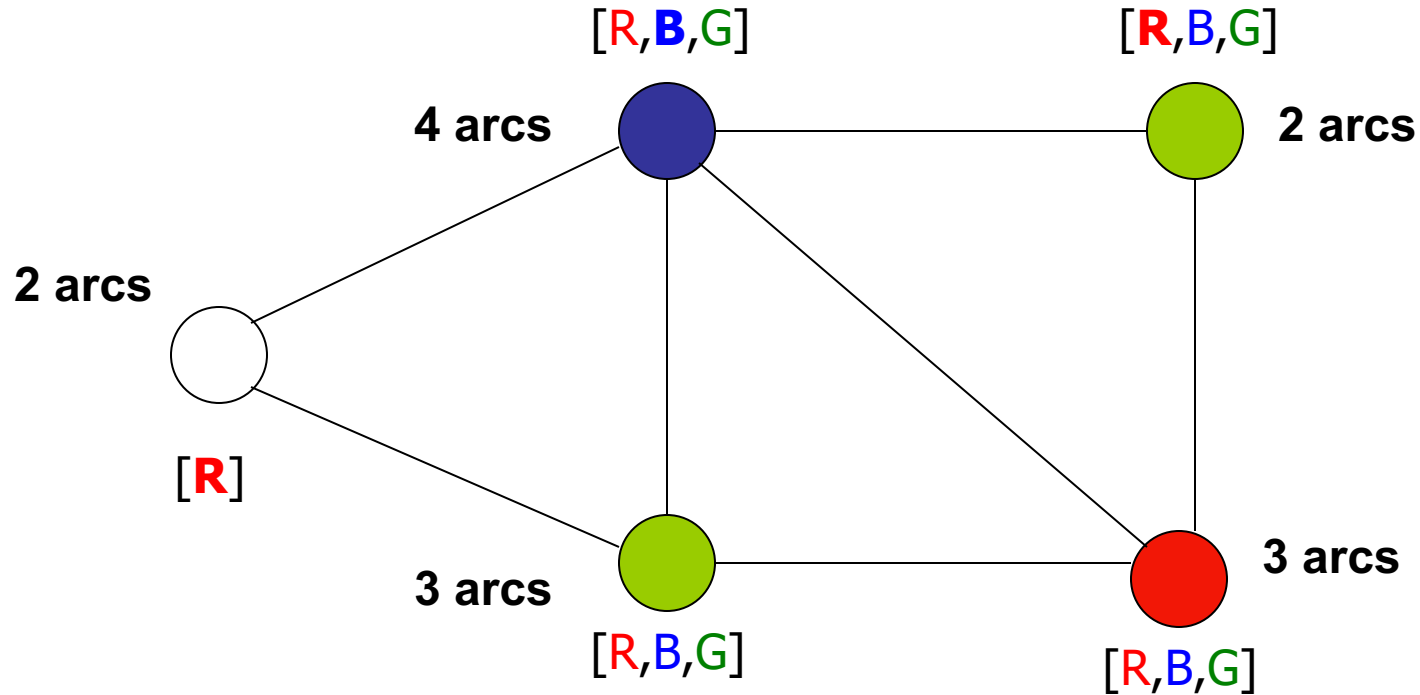
Backpropagation – LCV

Least constraining value



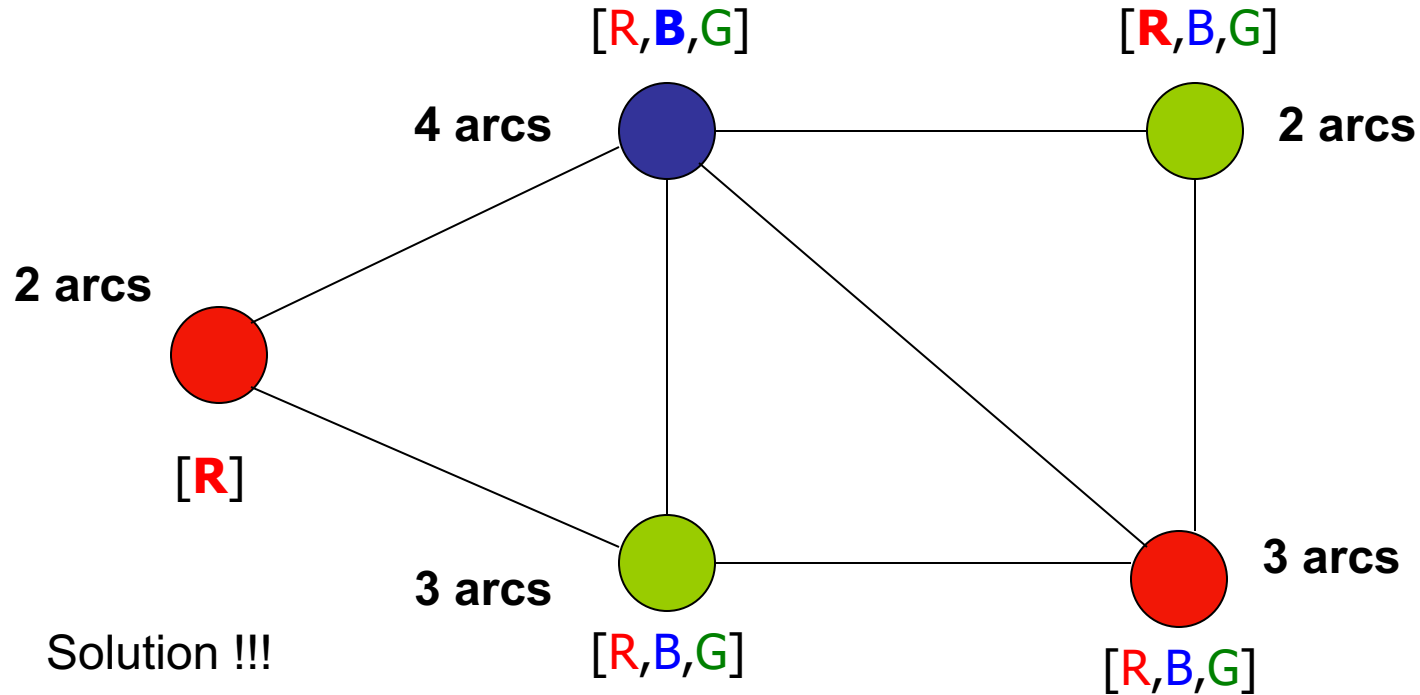
Backpropagation – LCV

Least constraining value



Backpropagation - LCV

Least constraining value

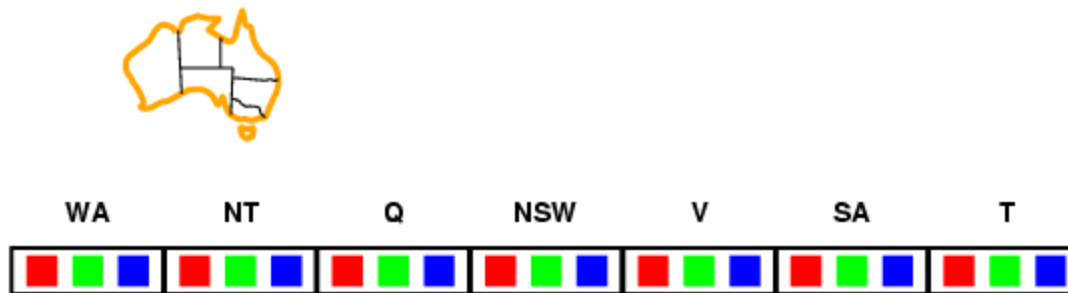


Analyzing Constraints

- forward checking
 - when a value X is assigned to a variable, inconsistent values are eliminated for all variables connected to X **[remove conflicting values]**
 - identifies “dead” branches of the tree before they are visited
- constraint propagation
 - analyses interdependencies between variable assignments via **arc consistency**
 - an arc between X and Y is consistent if for every possible value x of X , there is some value y of Y that is consistent with x
 - more powerful than forward checking, but still reasonably efficient
 - but does not reveal every possible inconsistency

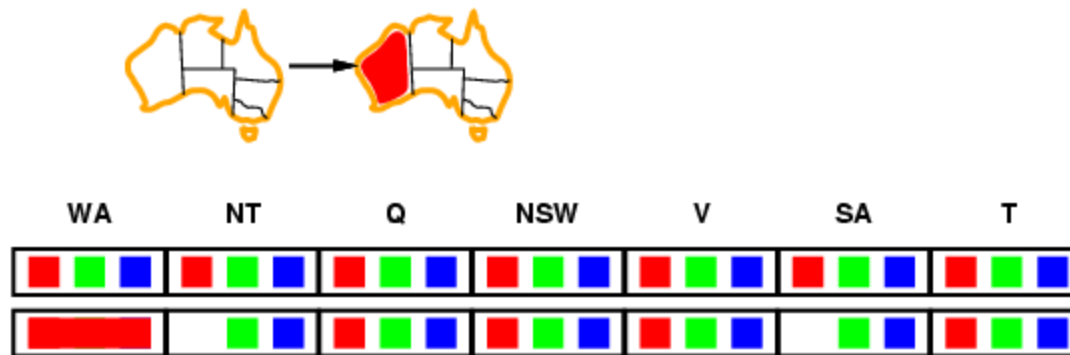
Forward checking

- Idea:
 - Keep track of remaining legal values for unassigned variables
 - Terminate search when any variable has no legal values



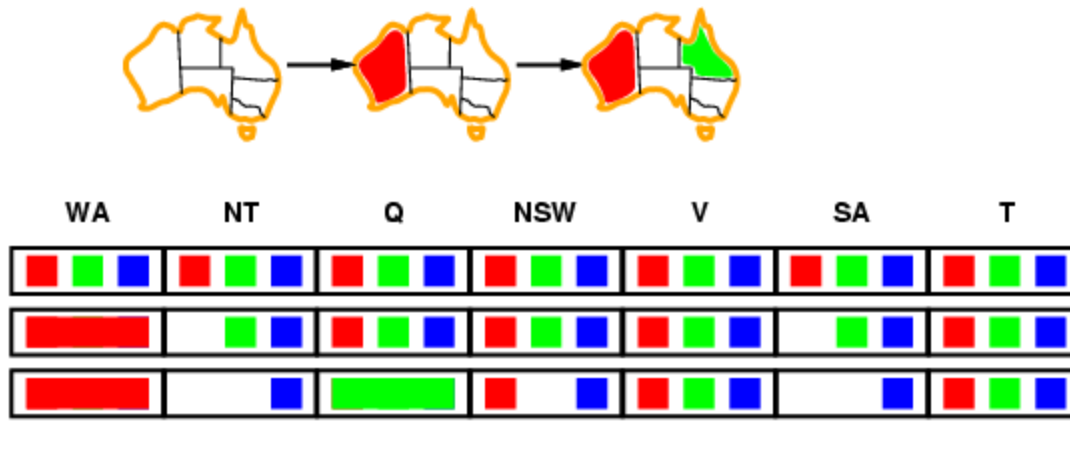
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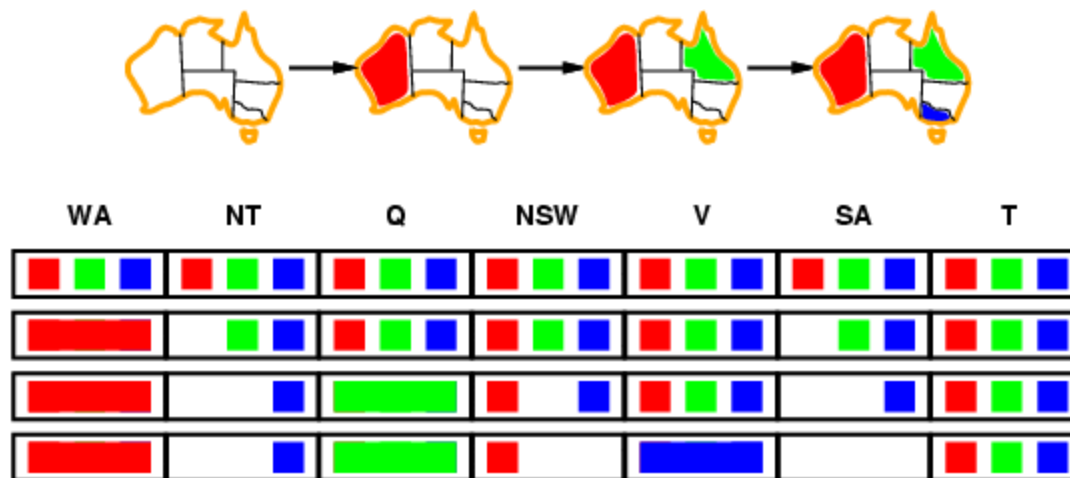
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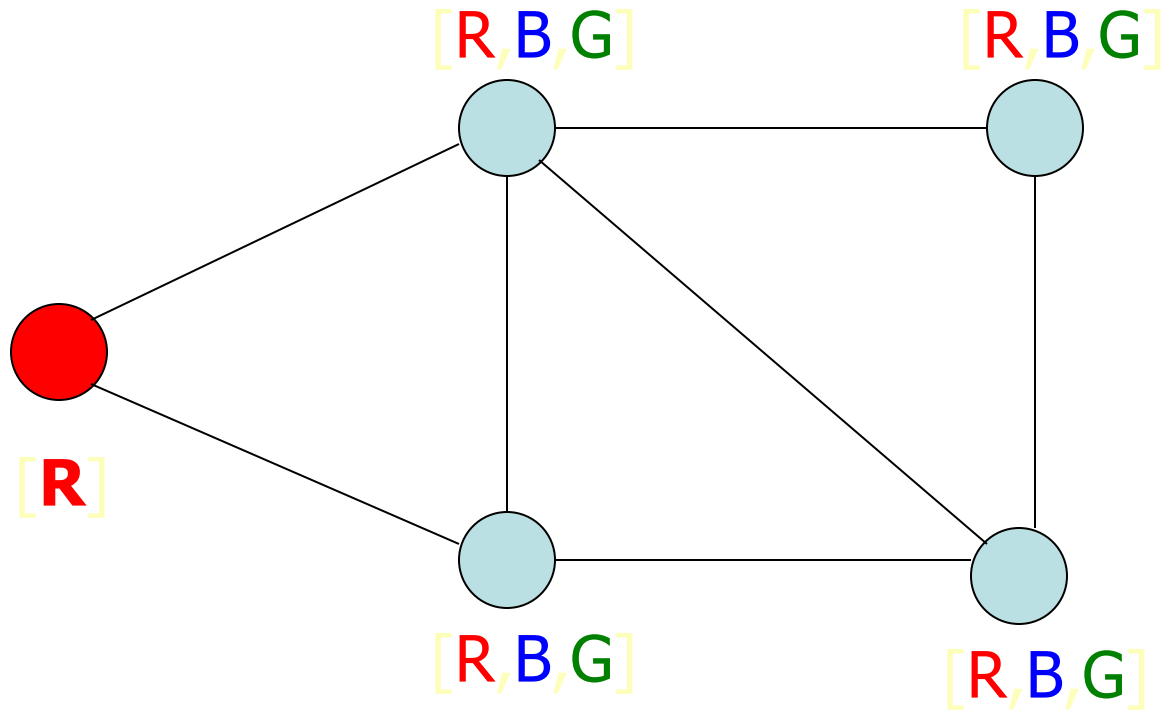


Forward checking

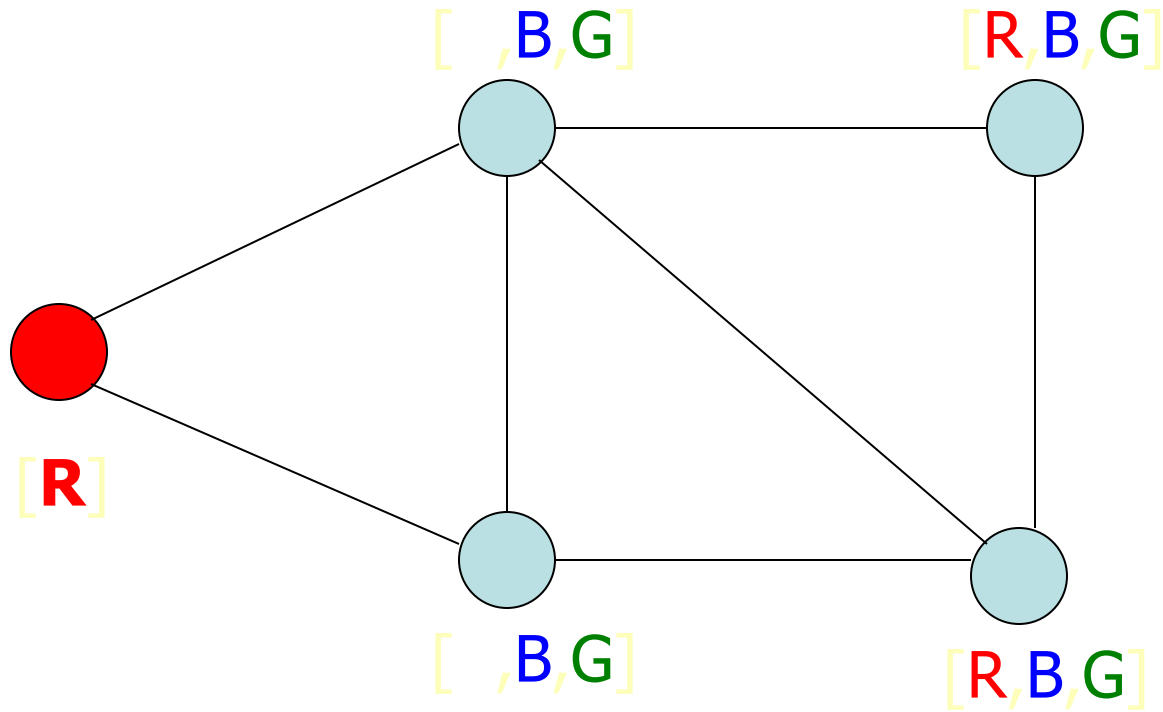
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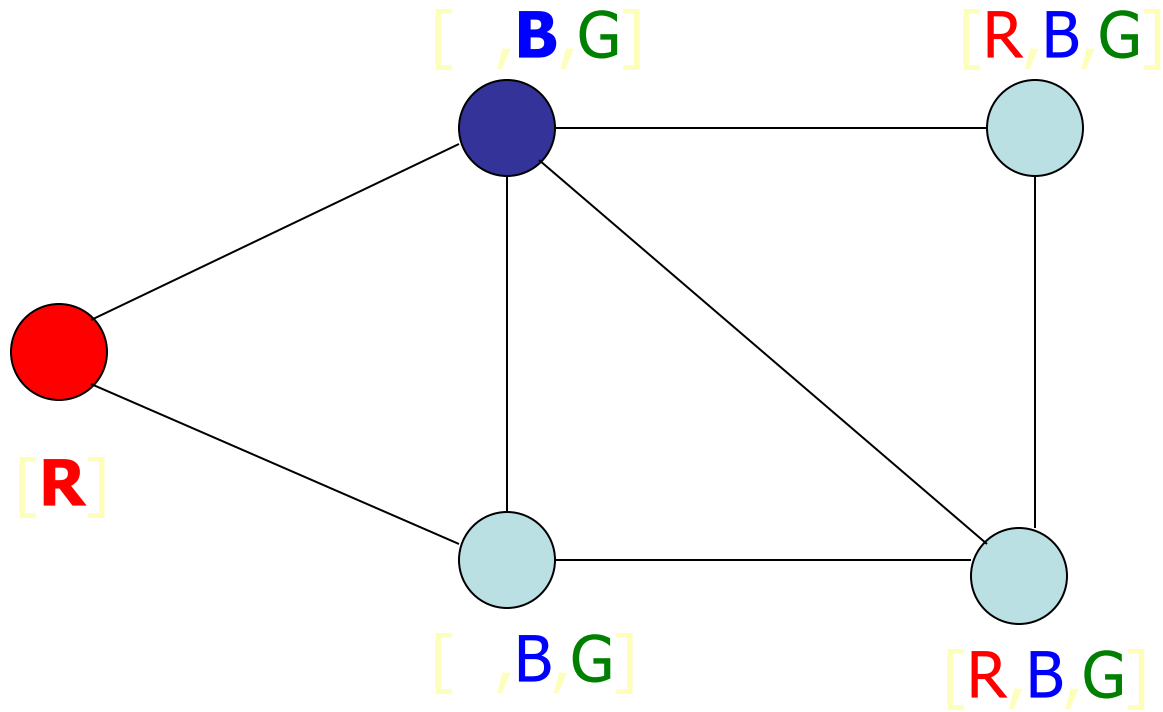
Forward Checking



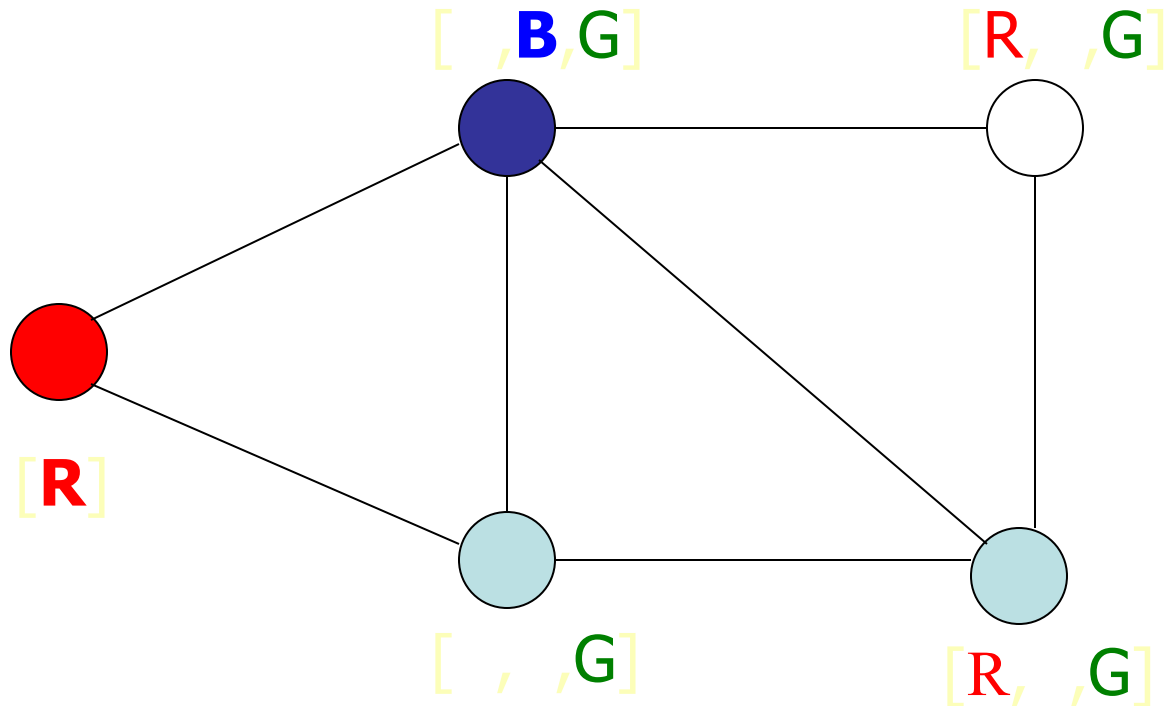
Forward Checking



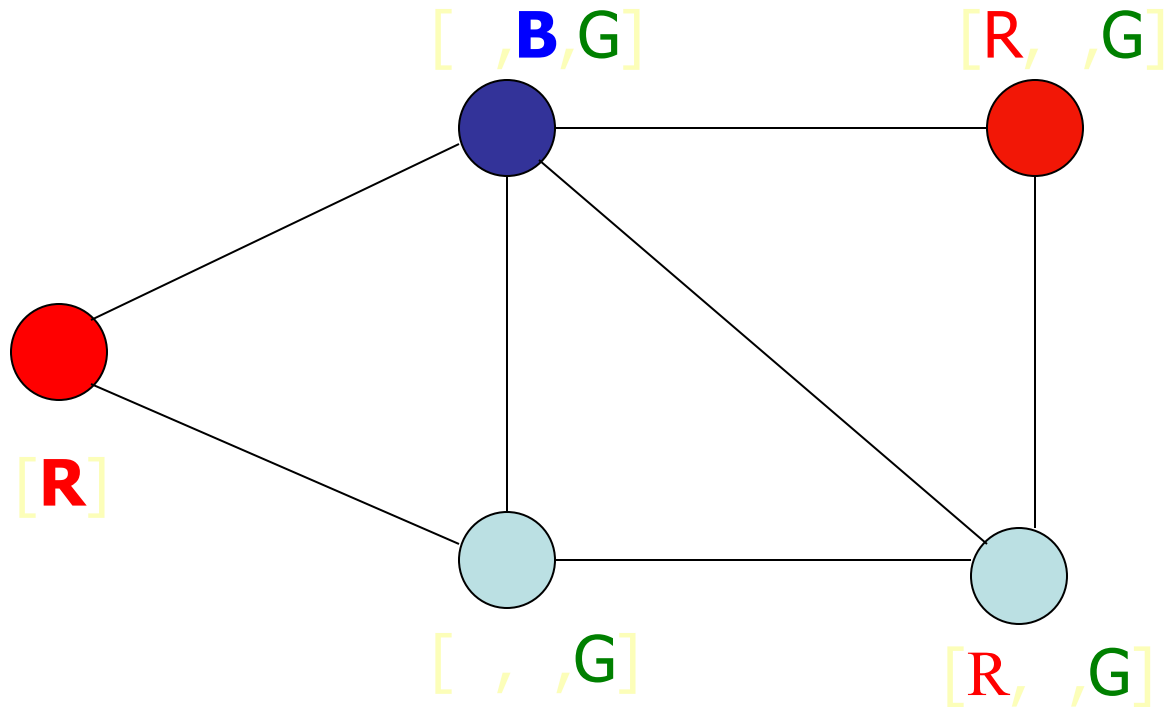
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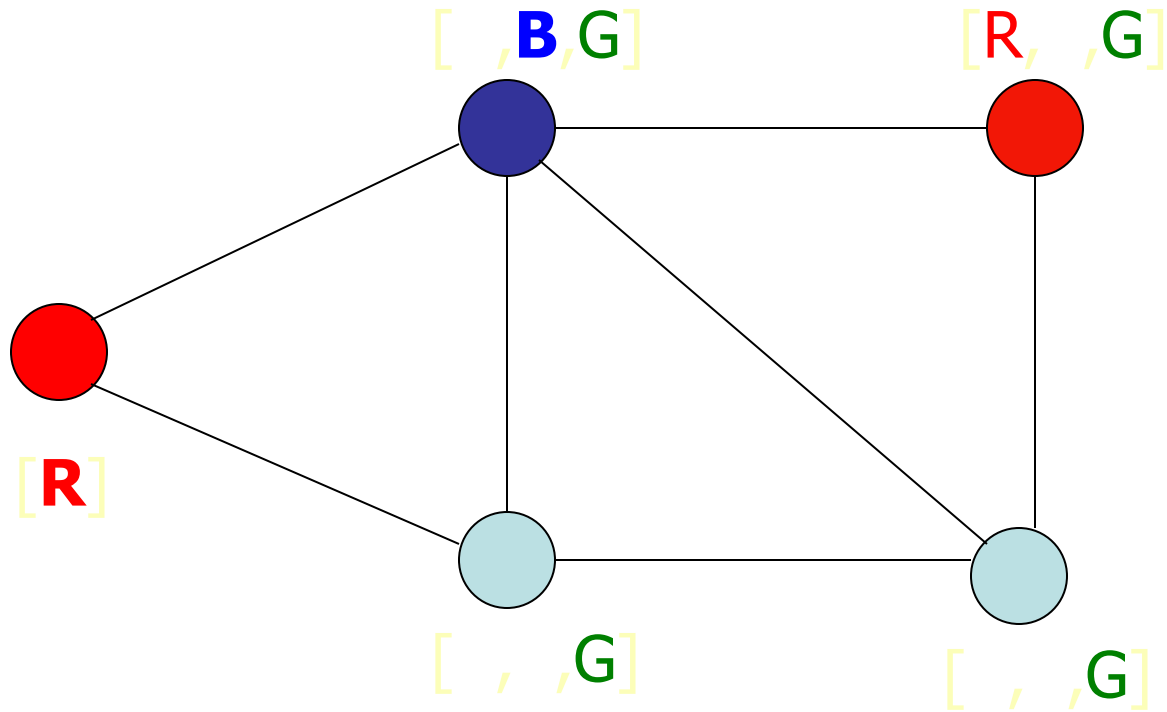
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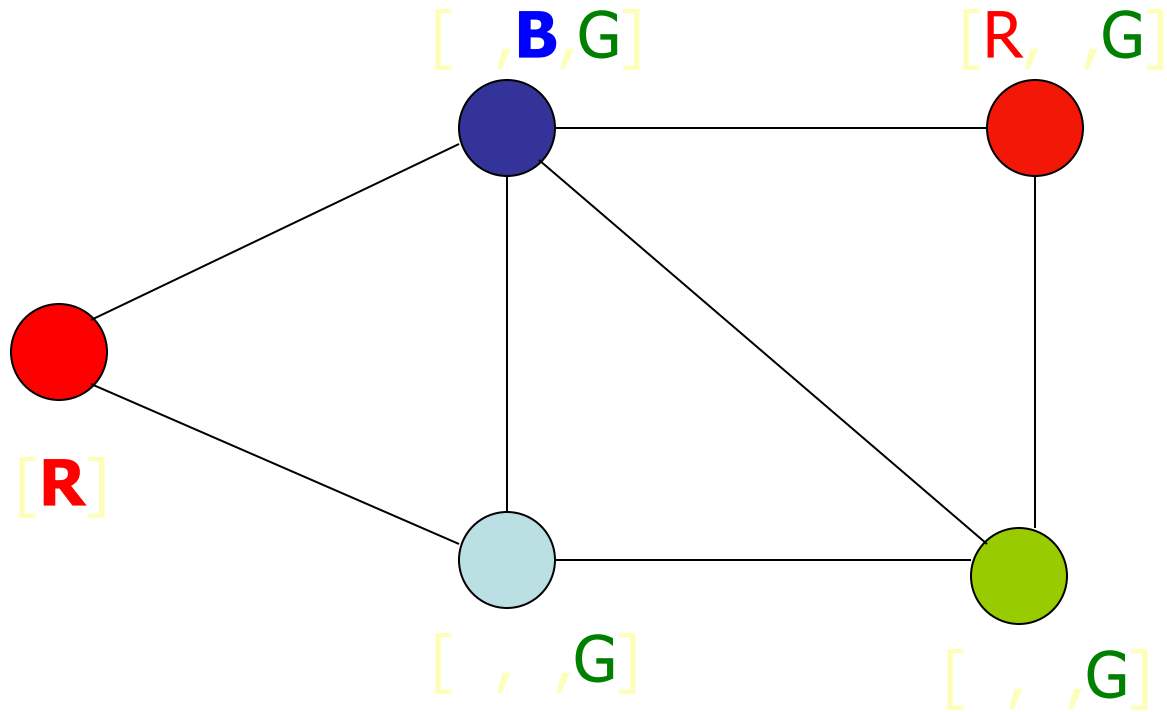
Forward Checking



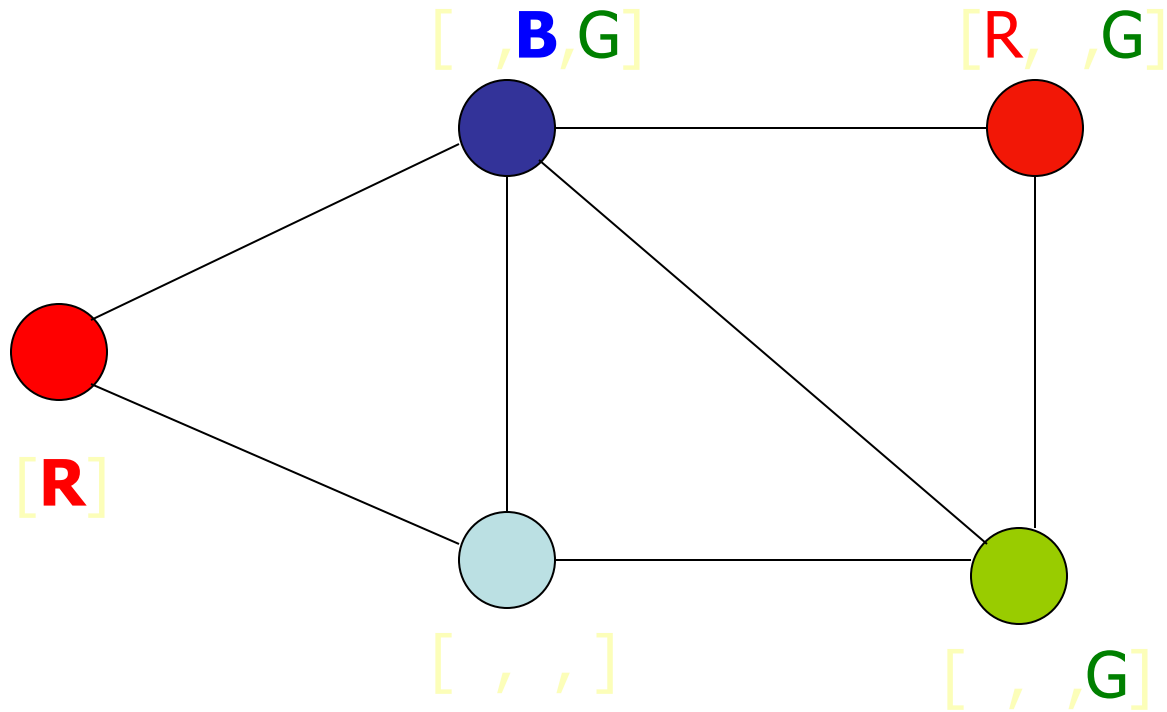
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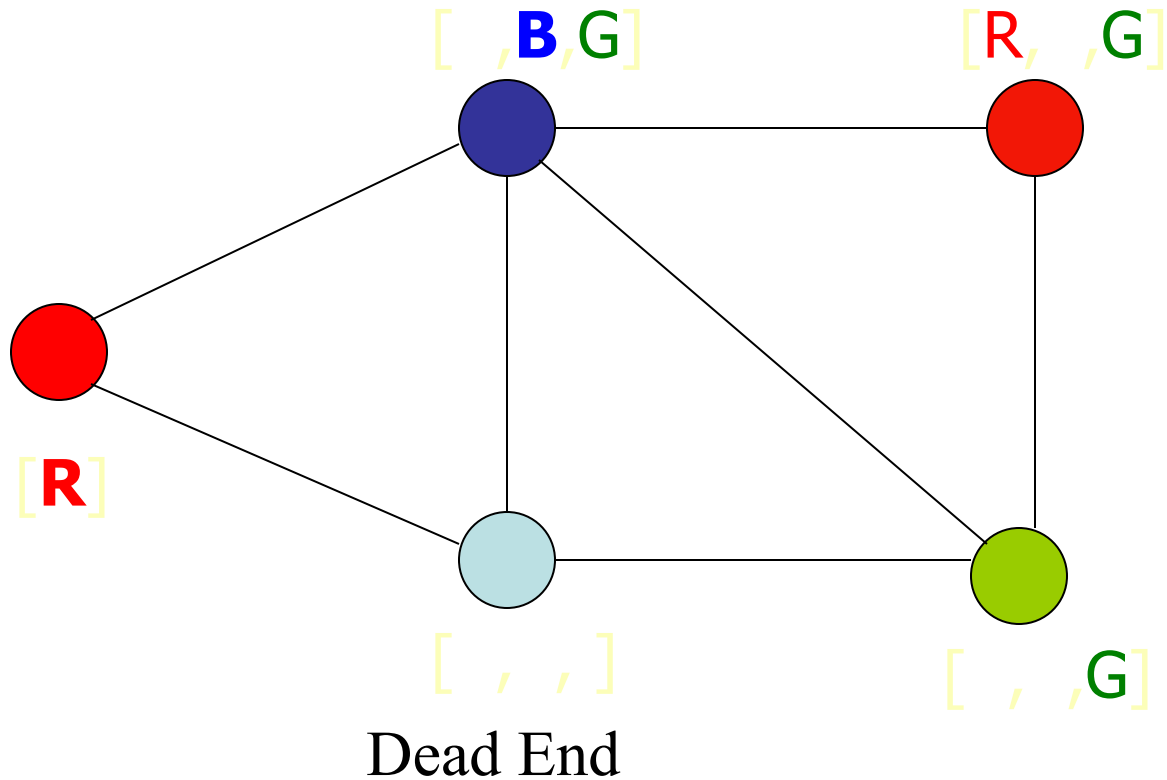
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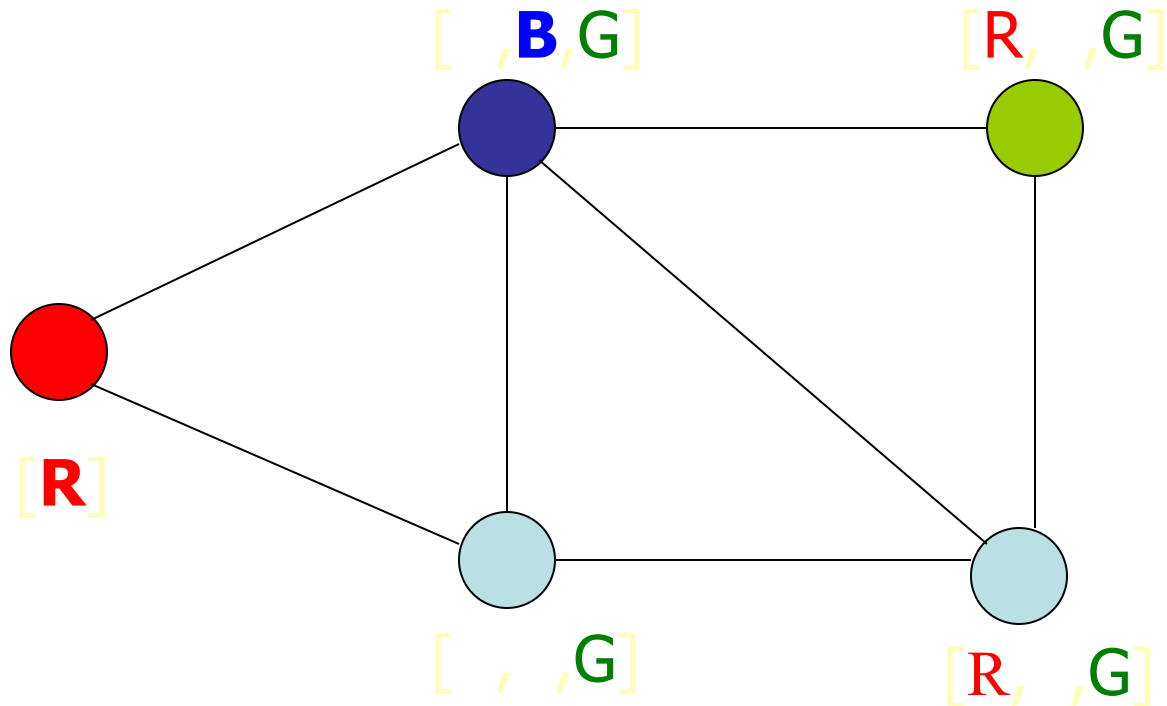
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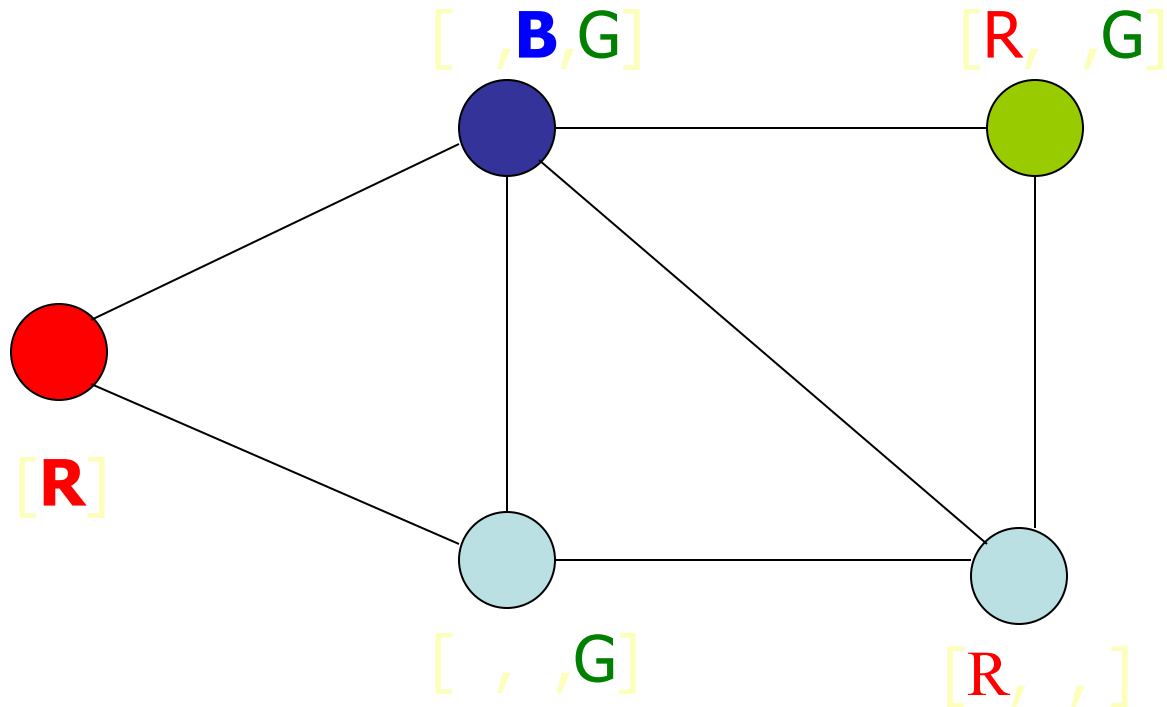
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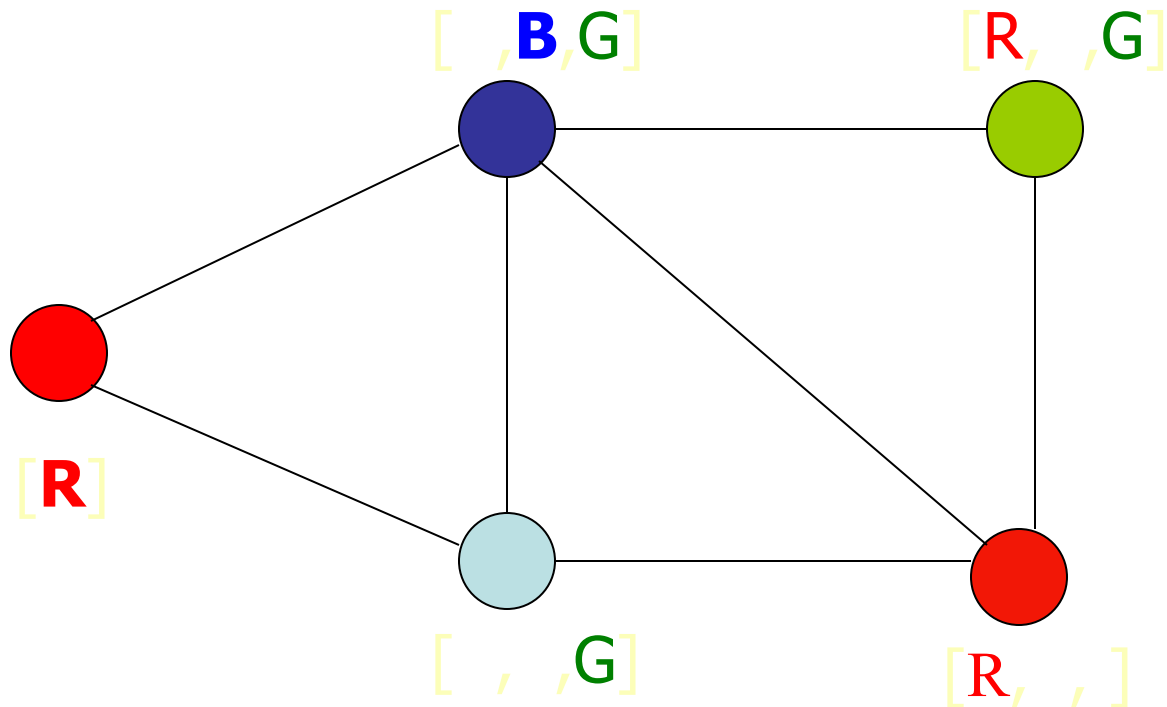
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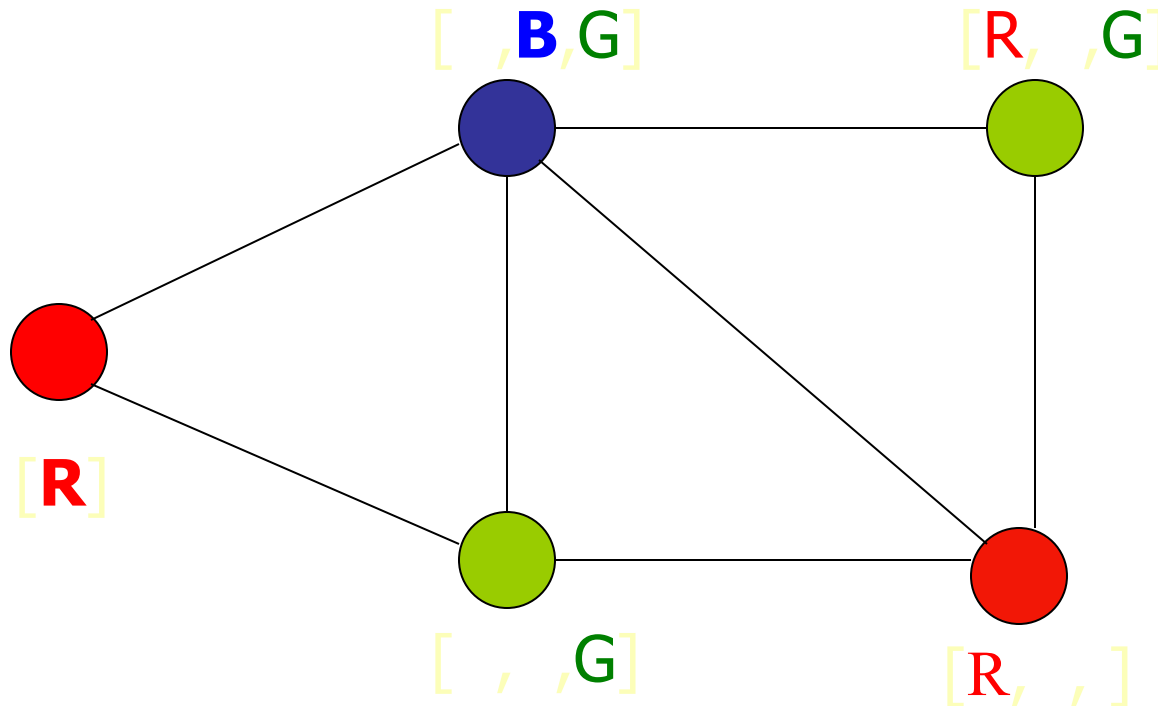
Forward Checking



Forward Checking



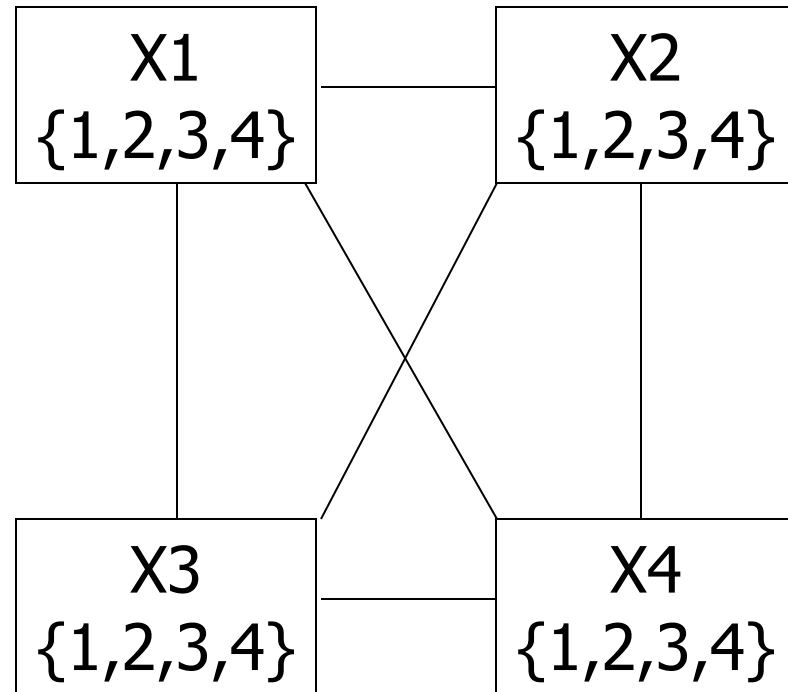
Forward Checking










Solution !!!

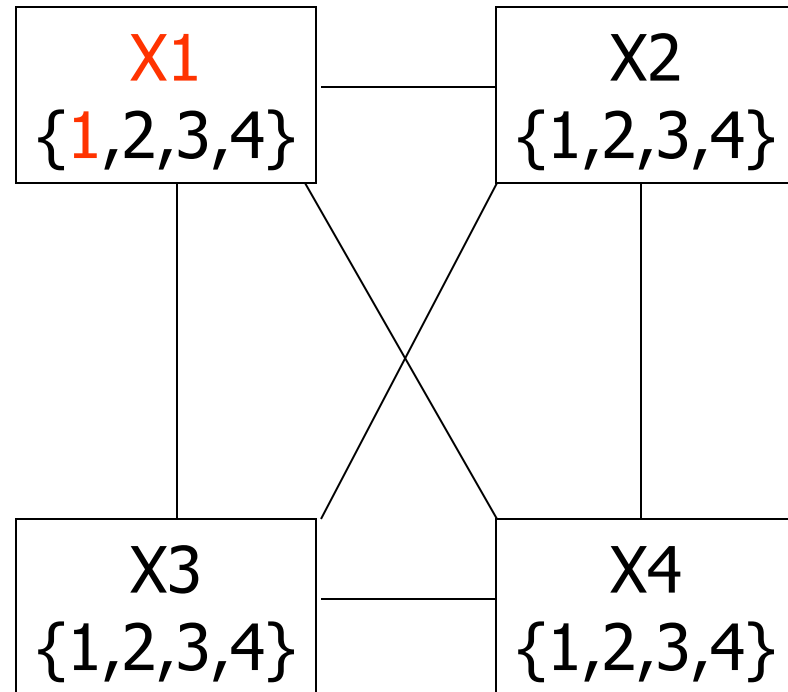
Example: 4-Queens Problem

	1	2	3	4
1				
2				
3				
4				










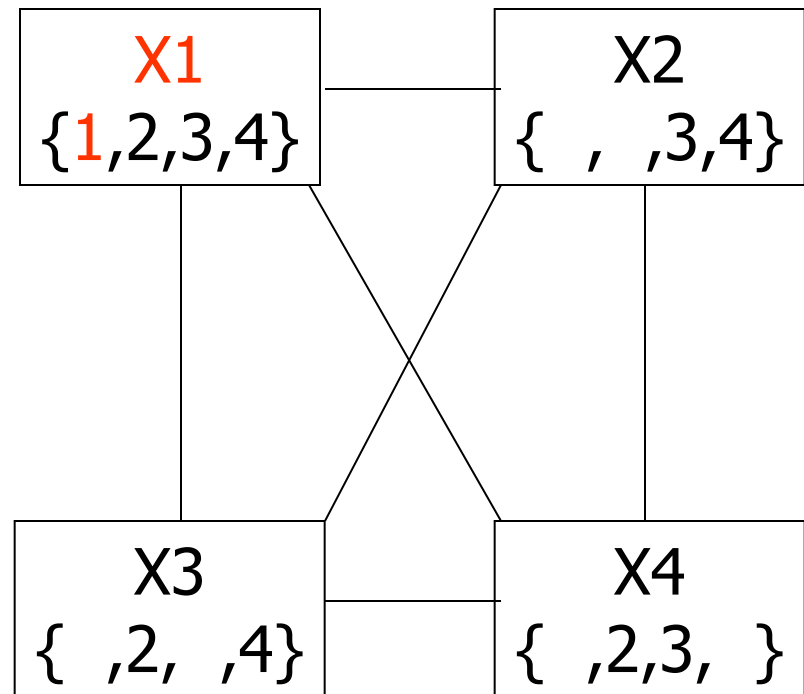
Example: 4-Queens Problem

	1	2	3	4
1				
2				
3				
4				














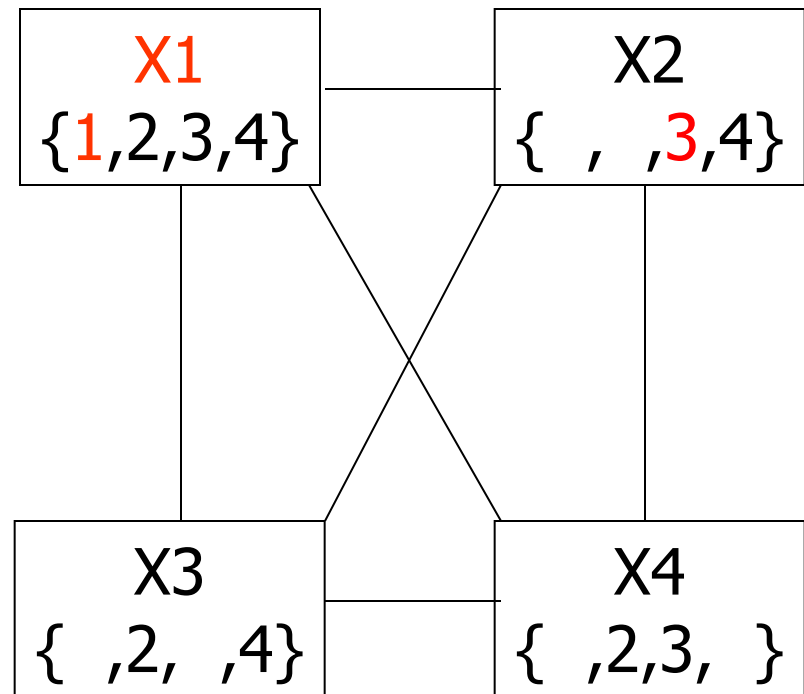
Example: 4-Queens Problem

	1	2	3	4
1				
2				
3				
4				



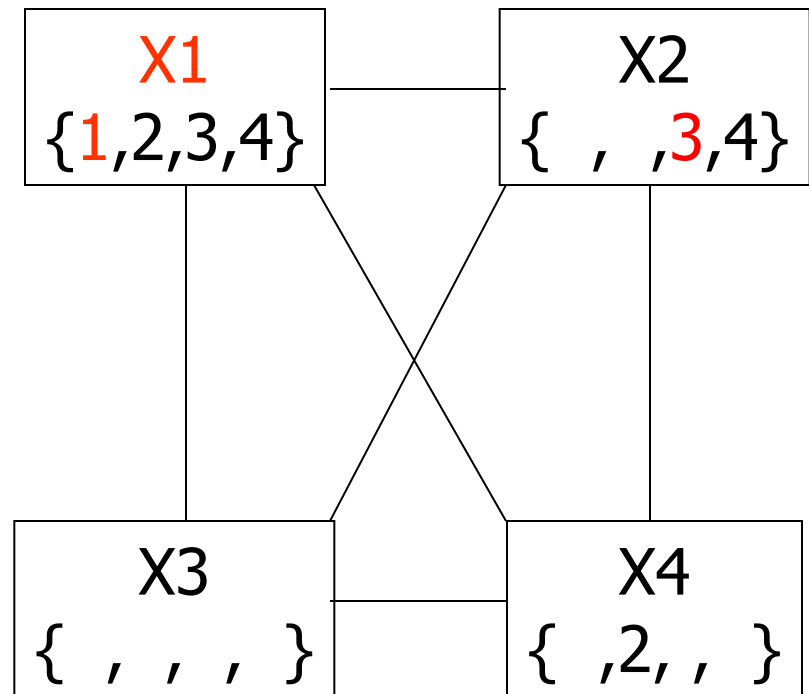
Example: 4-Queens Problem

	1	2	3	4
1				
2				
3				
4				



Example: 4-Queens Problem

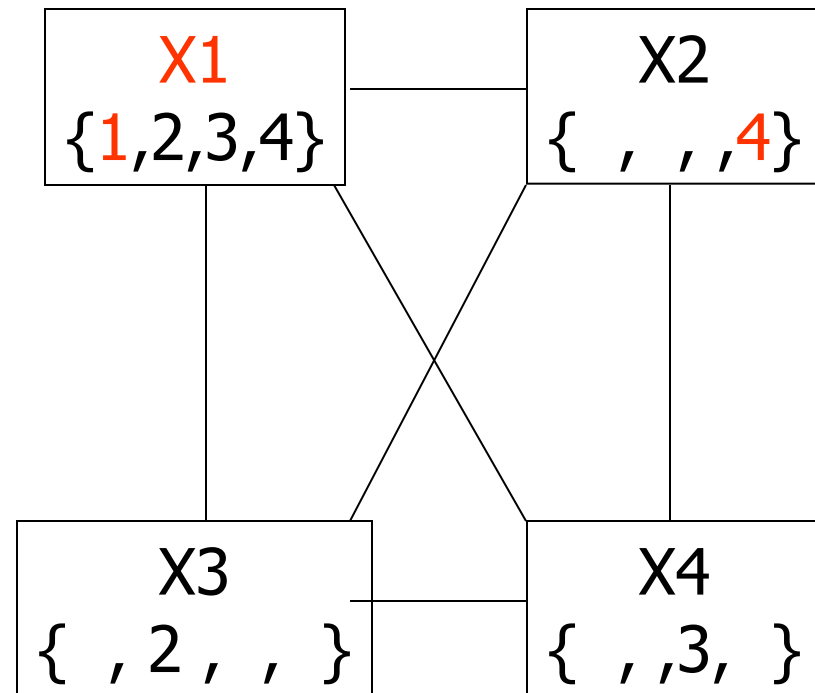
	1	2	3	4
1	★	●	●	●
2		●	●	
3		★	●	●
4			●	●



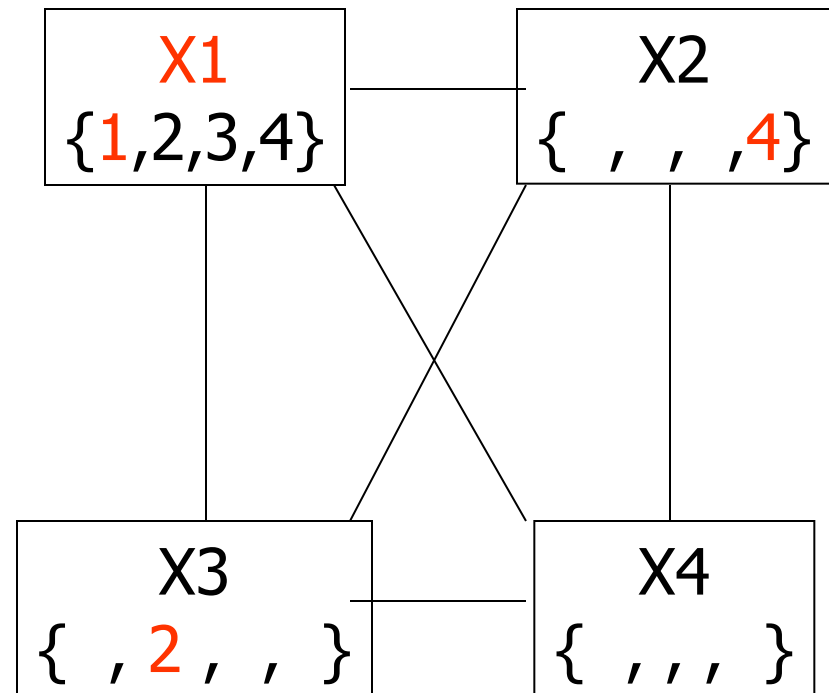
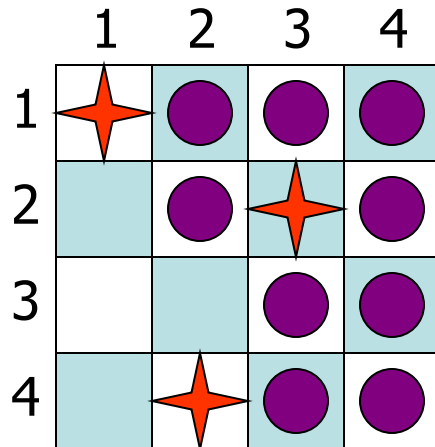
Dead End → Backtrack

Example: 4-Queens Problem

	1	2	3	4
1	★	●	●	●
2		●		●
3			●	
4		★	●	●



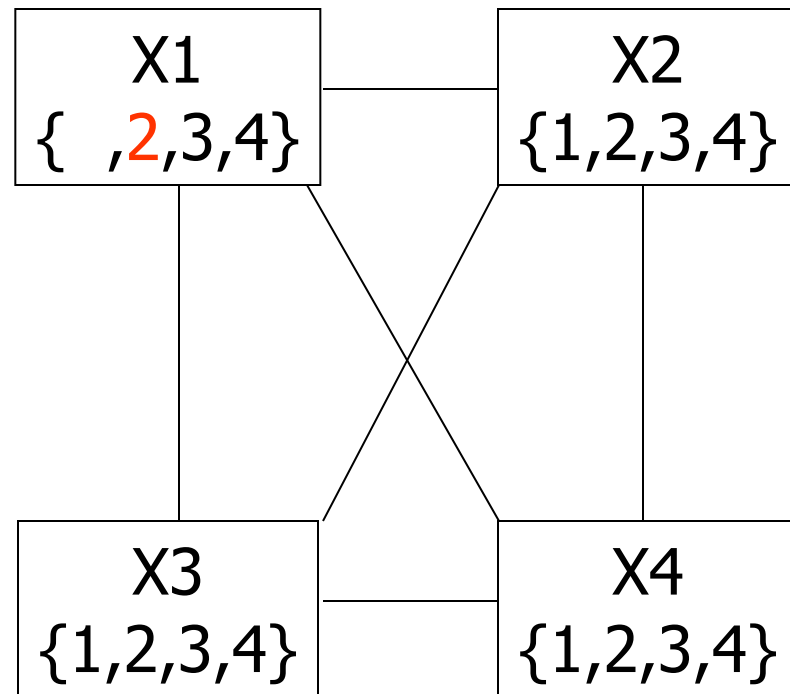
Example: 4-Queens Problem



Dead End → Backtrack

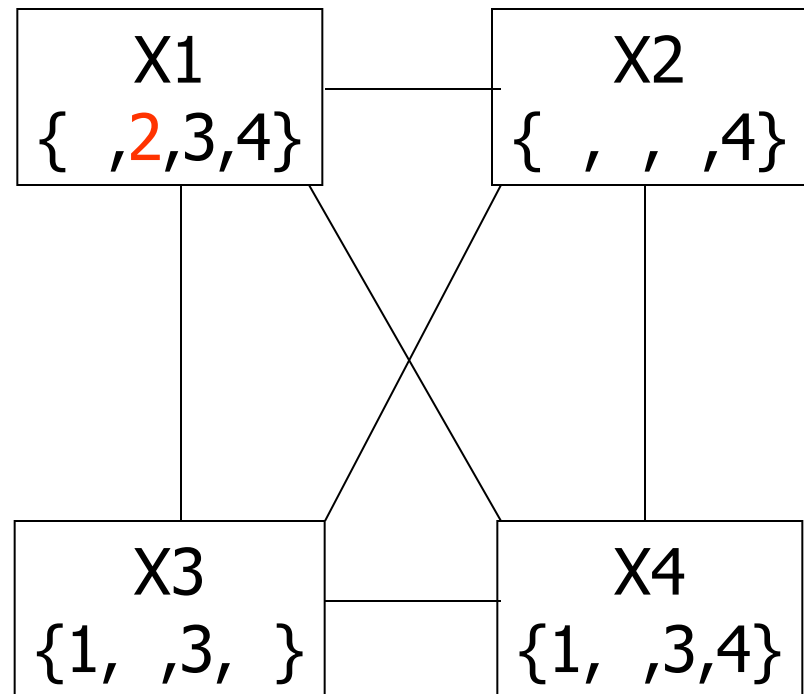
Example: 4-Queens Problem

	1	2	3	4
1		●		
2	★	●	●	●
3		●		
4			●	



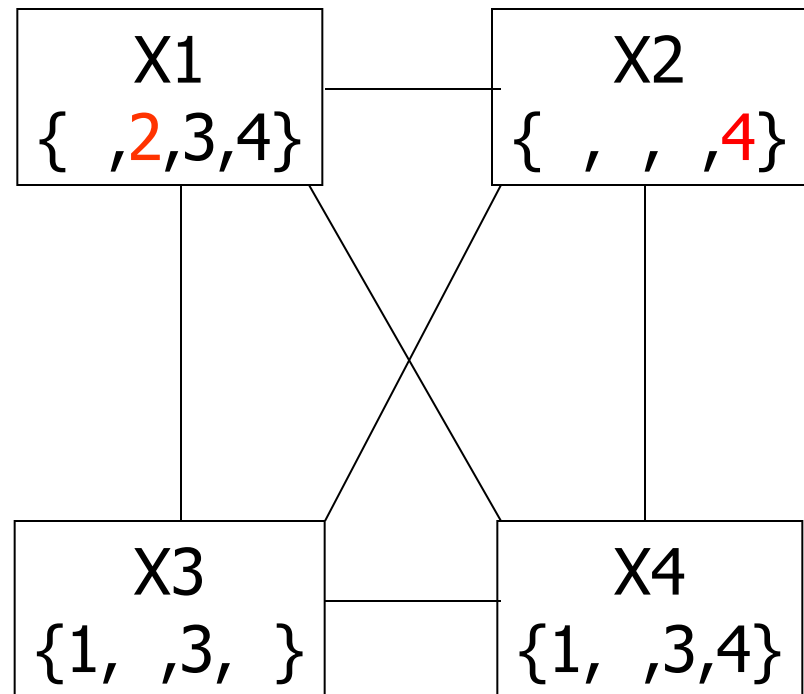
Example: 4-Queens Problem

	1	2	3	4
1		●		
2	★	●	●	●
3		●		
4			●	



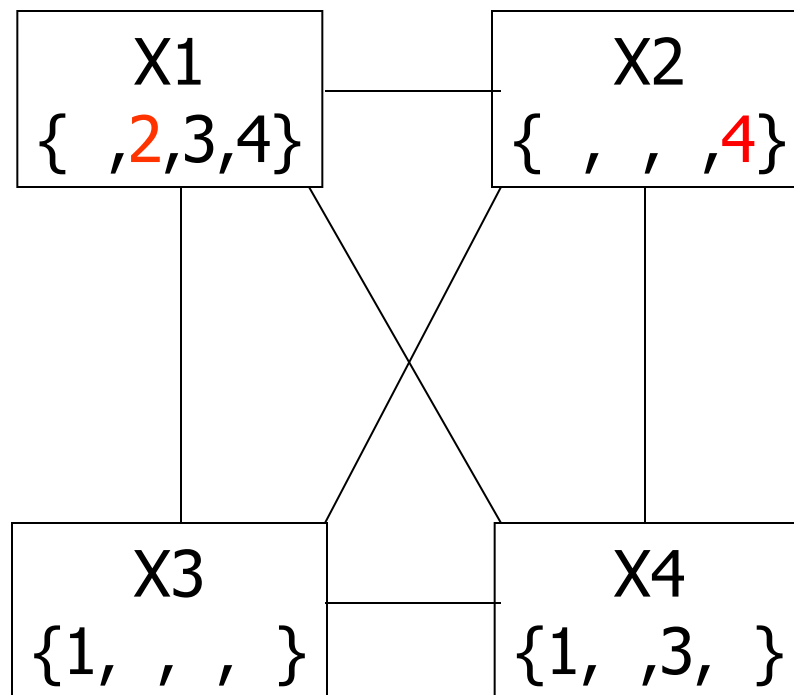
Example: 4-Queens Problem

	1	2	3	4
1		●		
2	★	●	●	●
3		●	●	
4		★	●	●



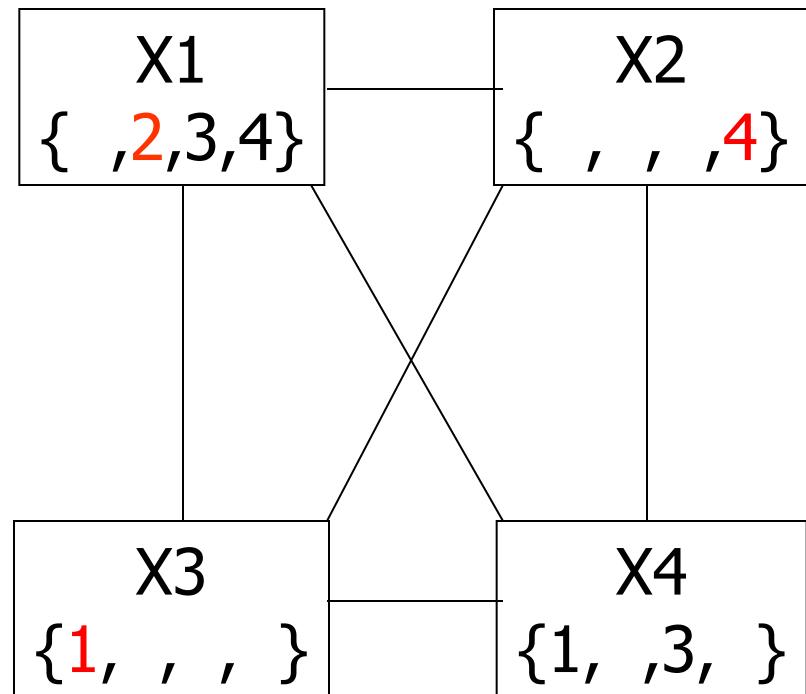
Example: 4-Queens Problem

	1	2	3	4
1		●		
2	★	●	●	●
3		●	●	
4		★	●	●



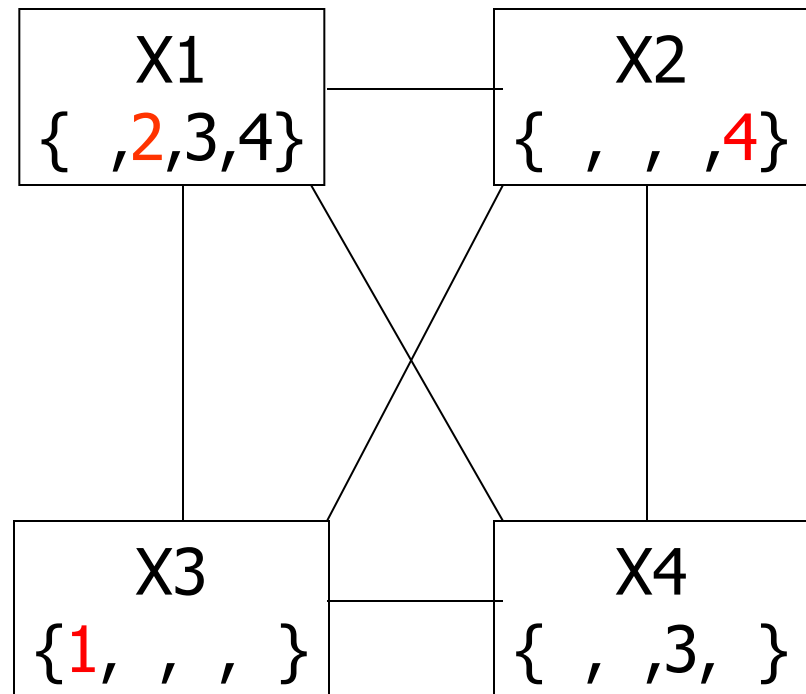
Example: 4-Queens Problem

	1	2	3	4
1		●	★	●
2	★	●	●	●
3		●	●	
4		★	●	●



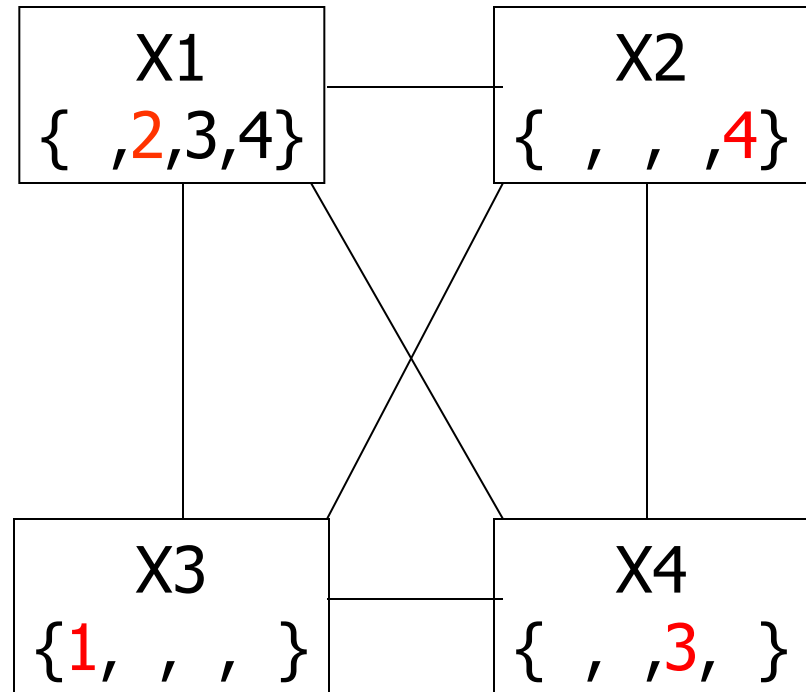
Example: 4-Queens Problem

	1	2	3	4
1		●	★	●
2	★	●	●	●
3		●	●	
4		★	●	●



Example: 4-Queens Problem

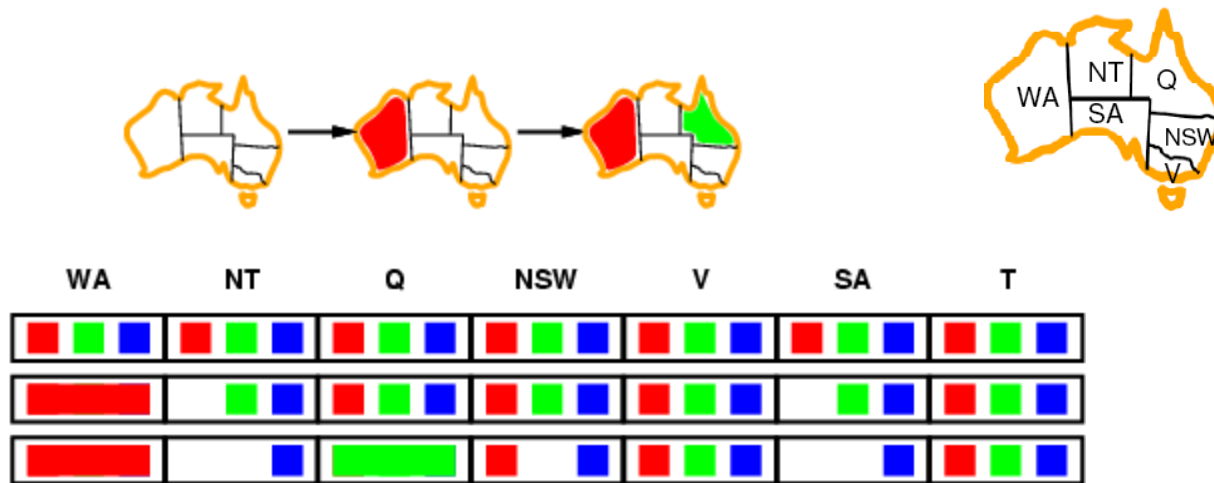
	1	2	3	4
1		●	★	●
2	★	●	●	●
3		●	●	★
4		★	●	●



Solution !!!!

Constraint propagation

- Forward checking propagates information from assigned to unassigned variables, but doesn't provide early detection for all failures:

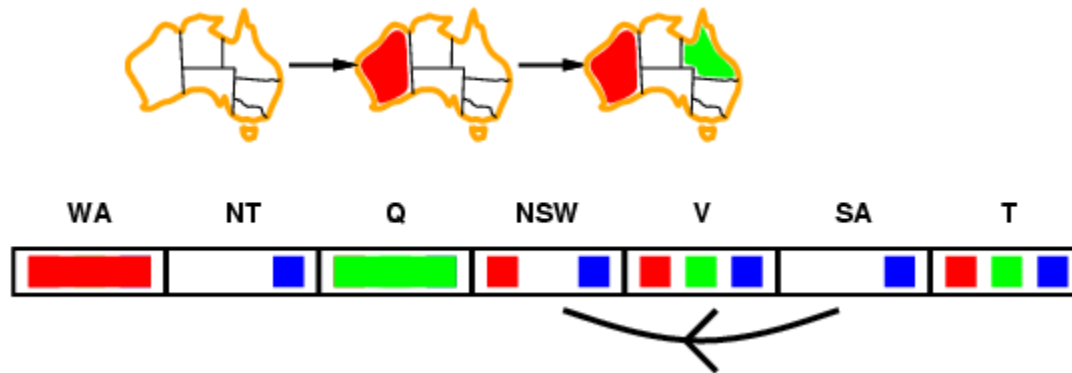


- NT and SA cannot both be blue!
- Constraint propagation repeatedly enforces constraints locally

Arc consistency

- Simplest form of propagation makes each arc **consistent**
- $X \rightarrow Y$ is consistent iff

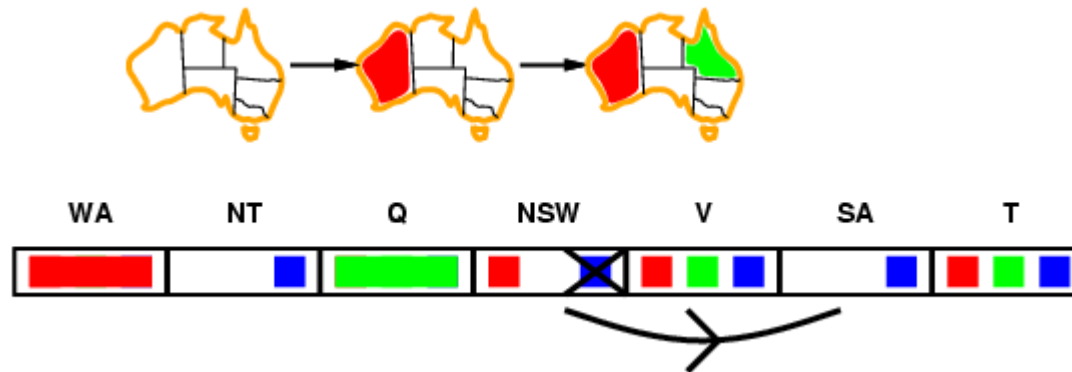
for **every** value x of X there is **some** allowed y



Arc consistency

- Simplest form of propagation makes each arc **consistent**
- $X \rightarrow Y$ is consistent iff

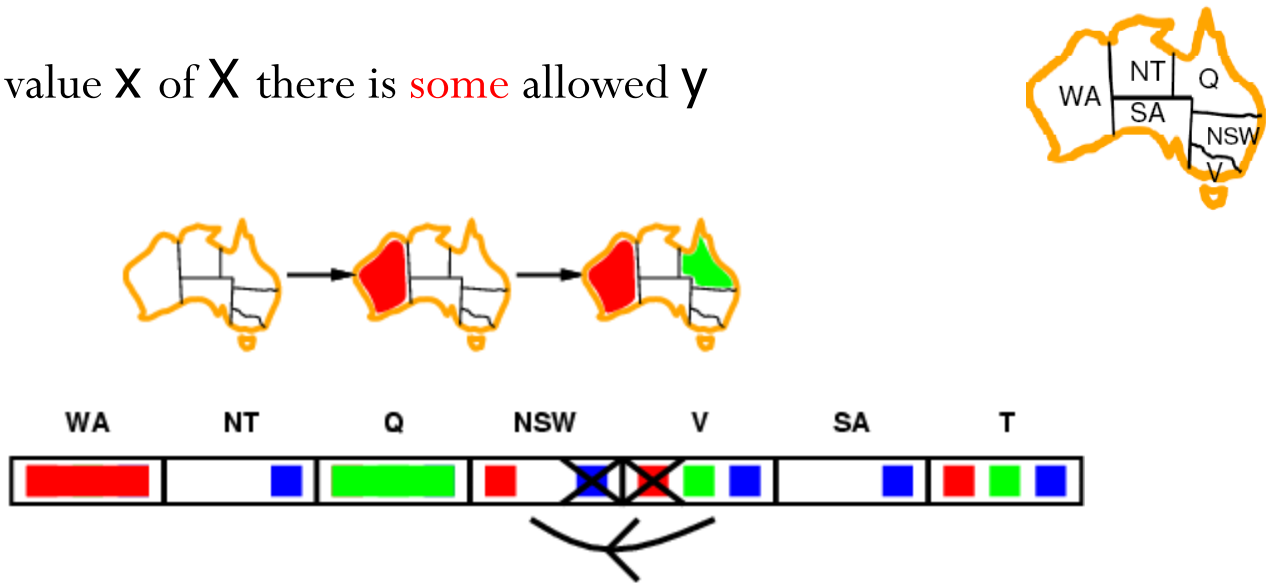
for **every** value x of X there is **some** allowed y



Arc consistency

- Simplest form of propagation makes each arc **consistent**
- $X \rightarrow Y$ is consistent iff

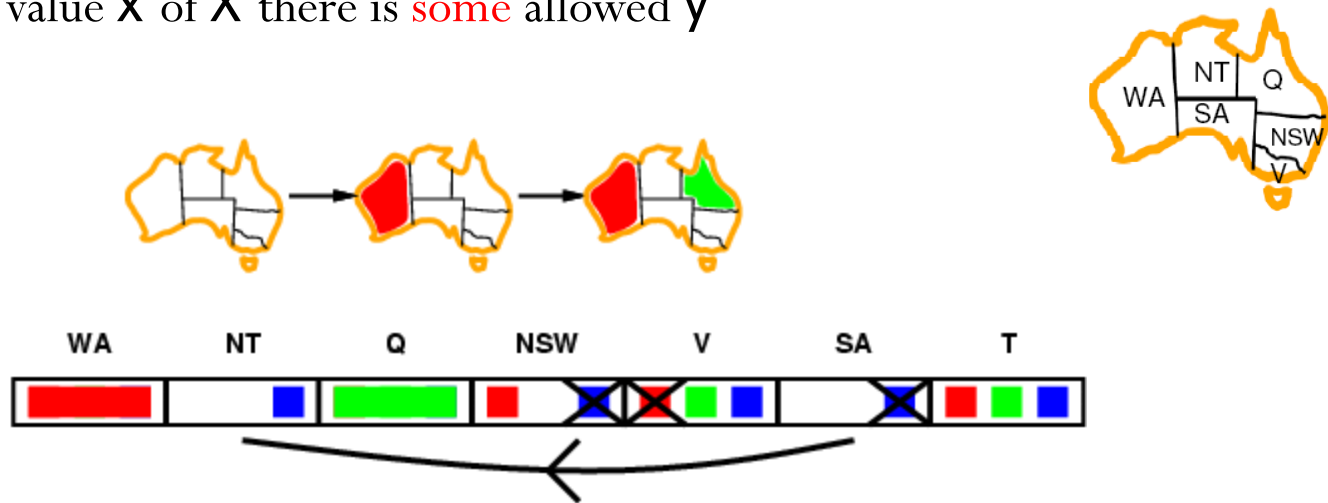
for **every** value **x** of **X** there is **some** allowed **y**



- If **X** loses a value, neighbors of **X** need to be rechecked

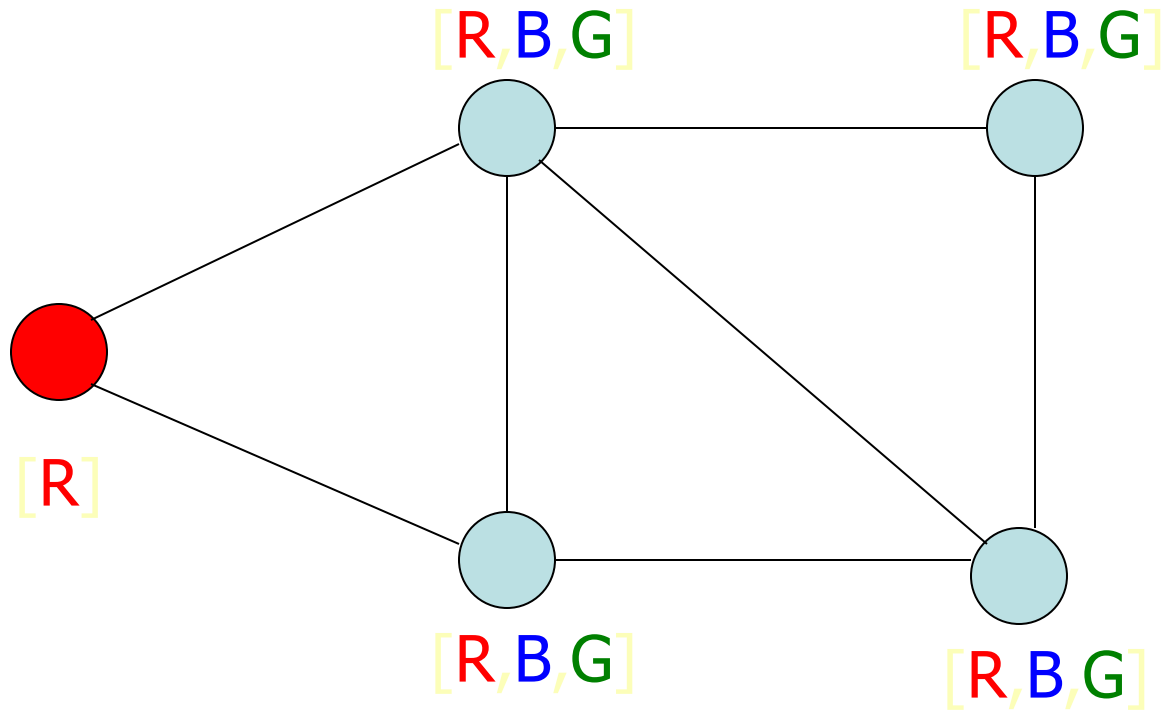
Arc consistency

- Simplest form of propagation makes each arc **consistent**
- $X \rightarrow Y$ is consistent iff
for **every** value x of X there is **some** allowed y

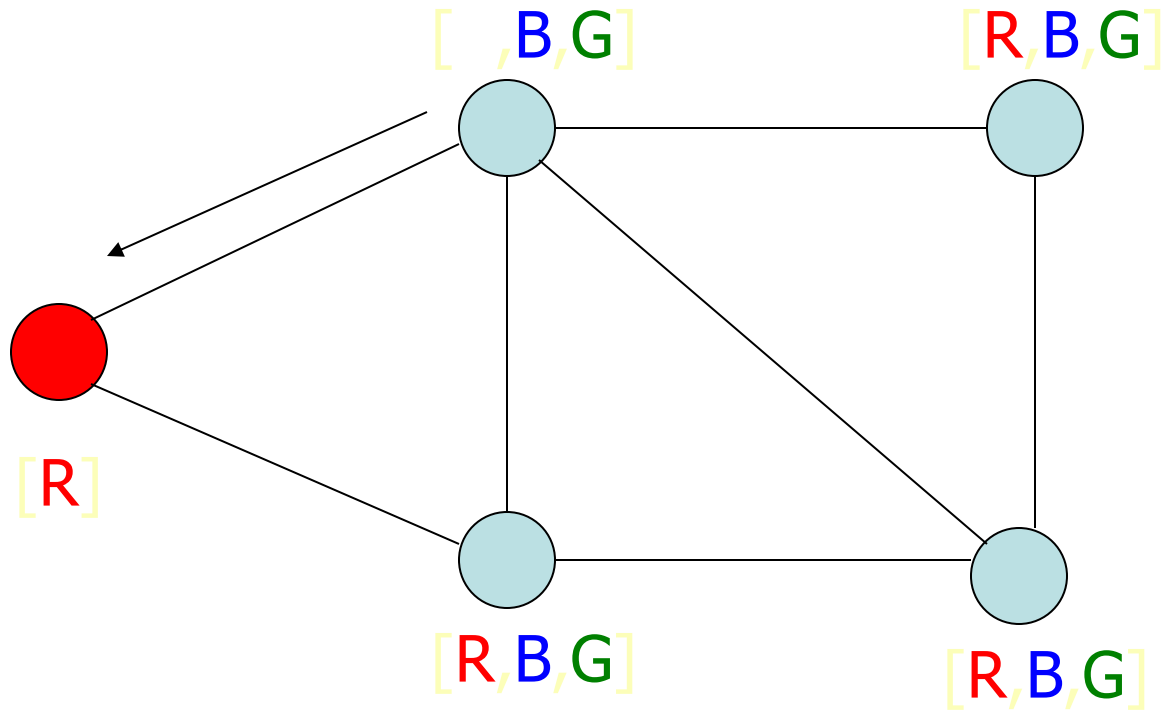


- If X loses a value, neighbors of X need to be rechecked
- Arc consistency detects failure earlier than forward checking
- Can be run as a preprocessor or after each assignment

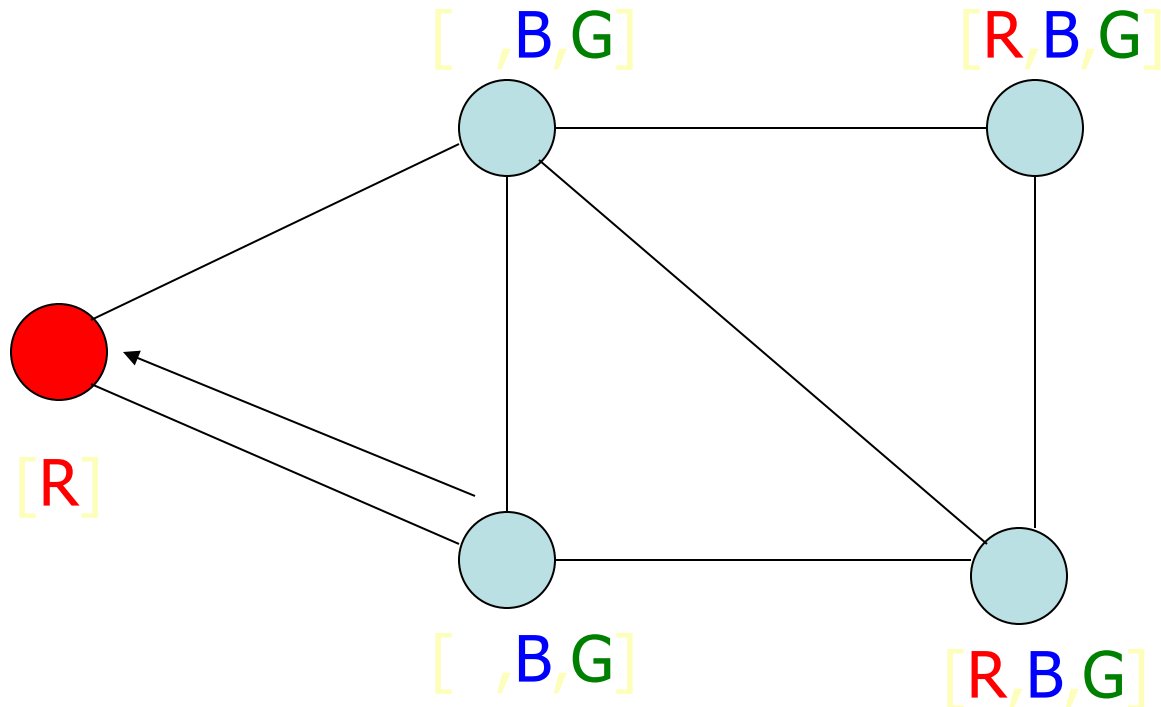
Arc Consistency: AC3



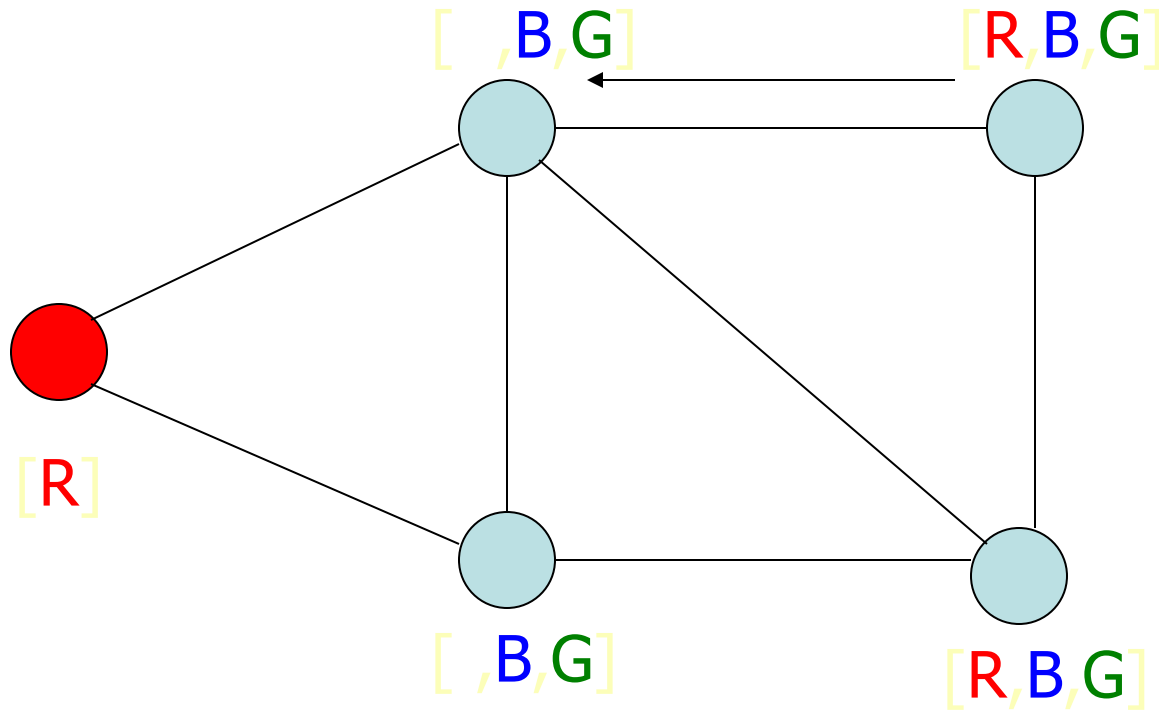
Arc Consistency: AC3



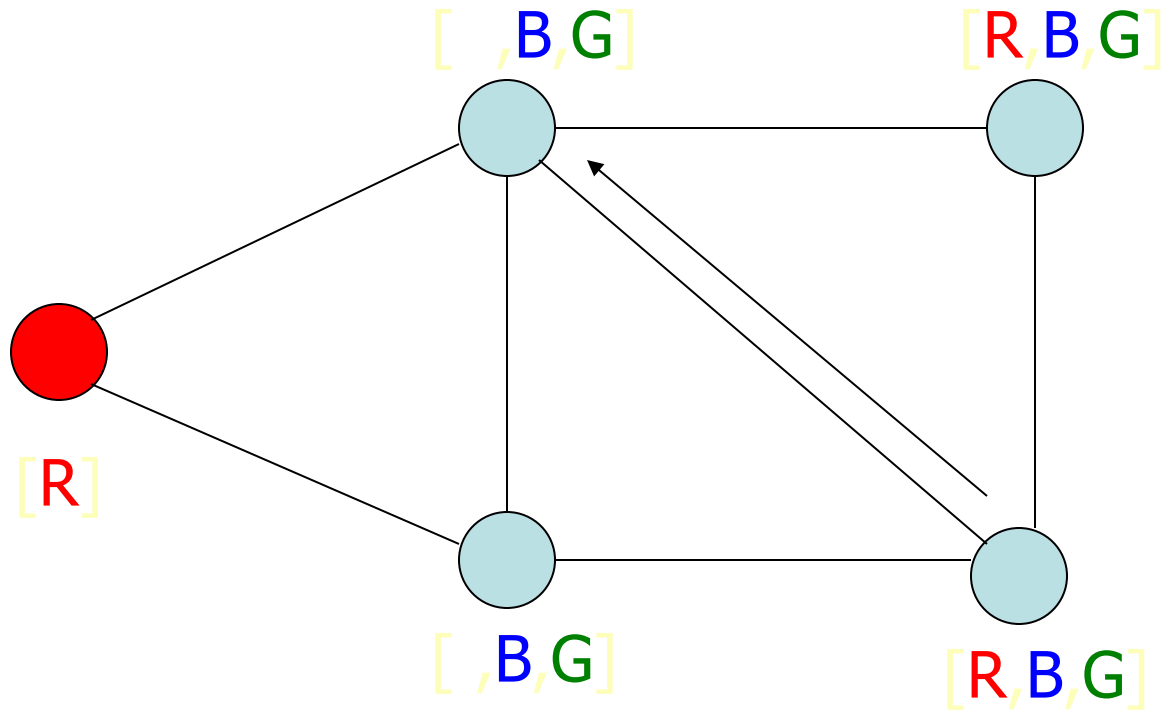
Arc Consistency: AC3



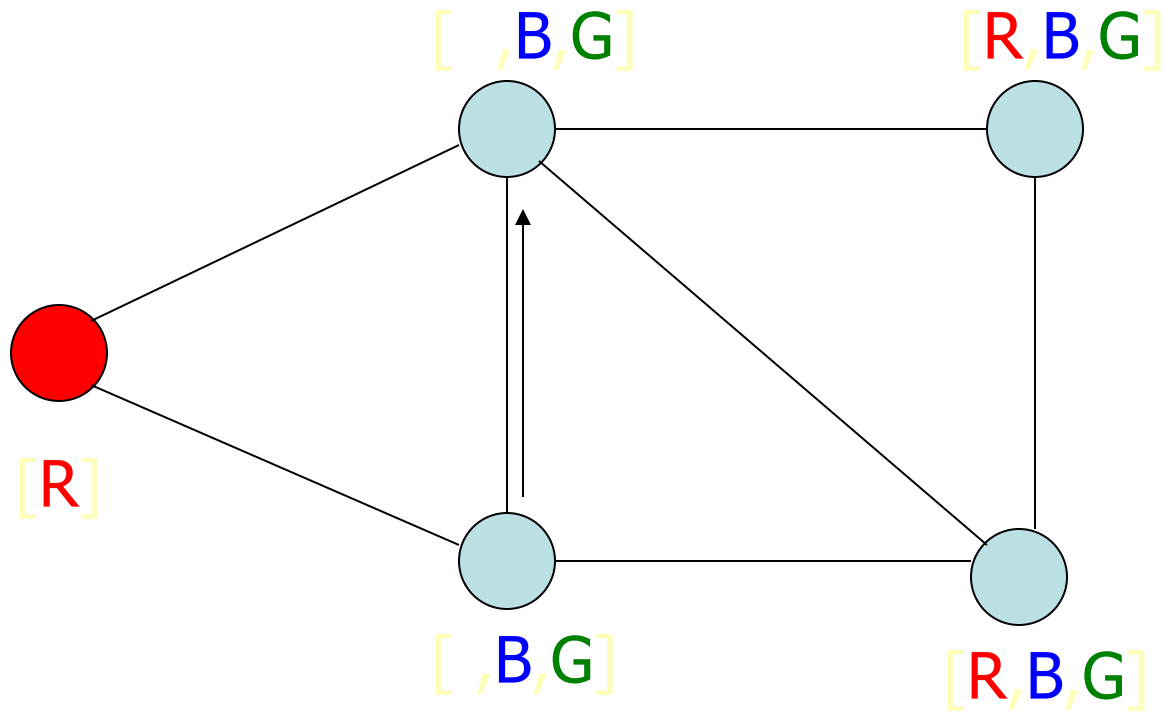
Arc Consistency: AC3



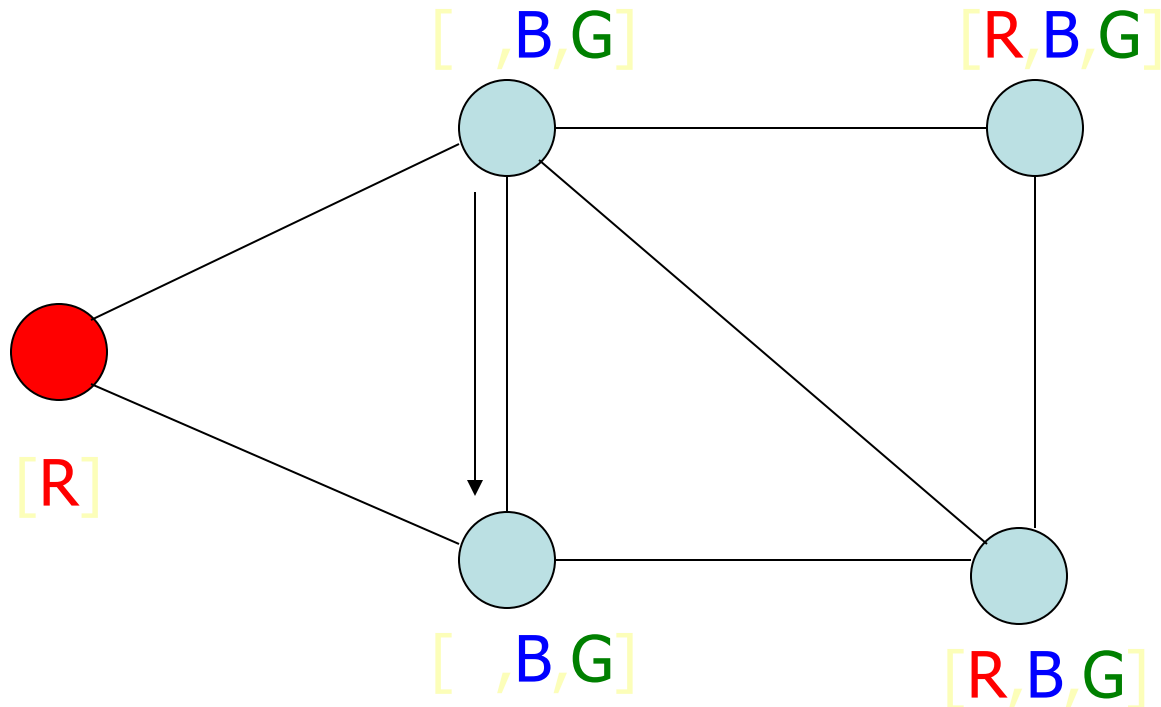
Arc Consistency: AC3



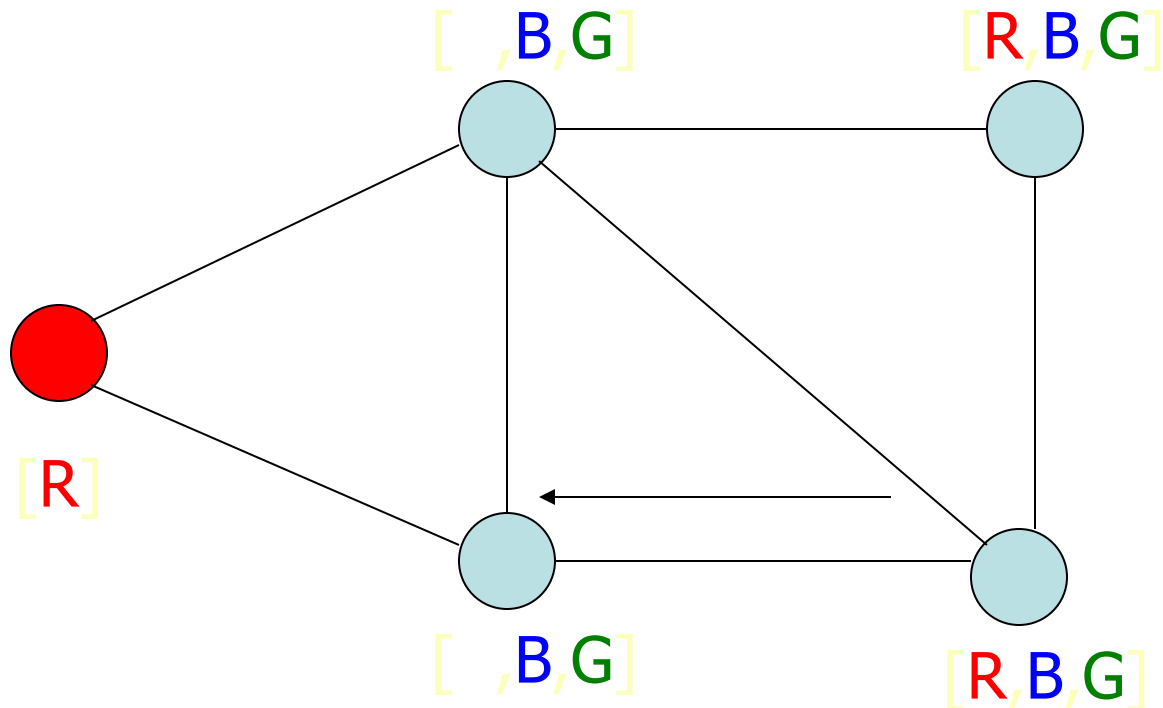
Arc Consistency: AC3



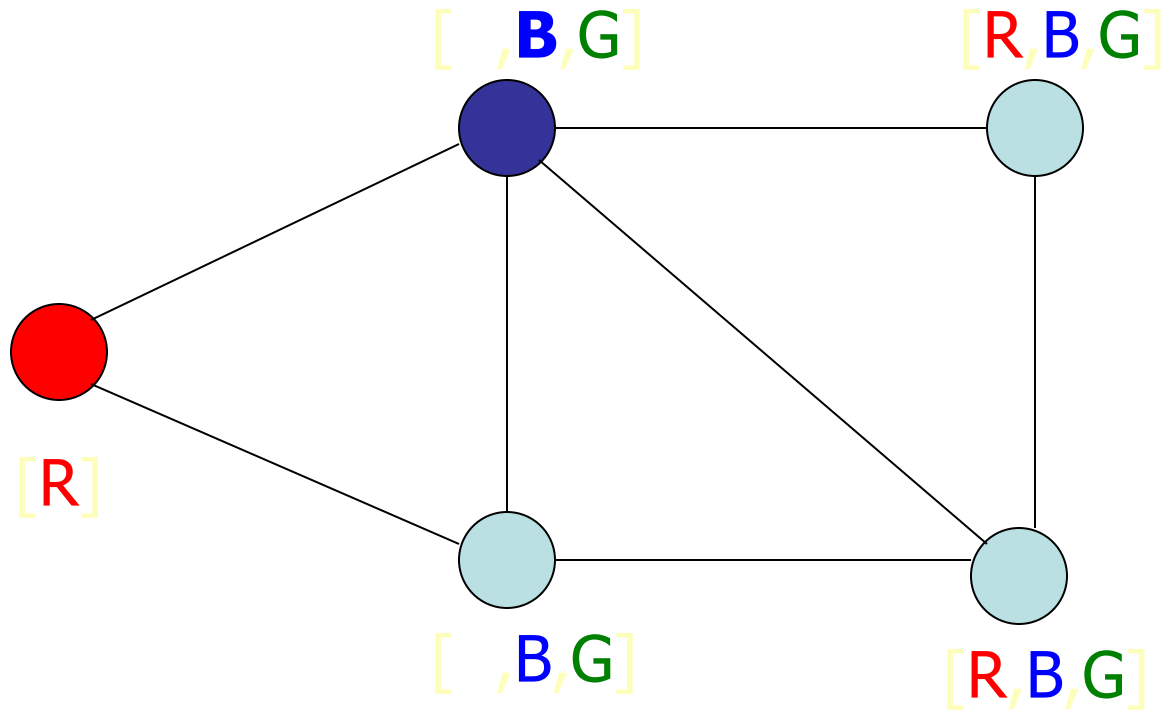
Arc Consistency: AC3



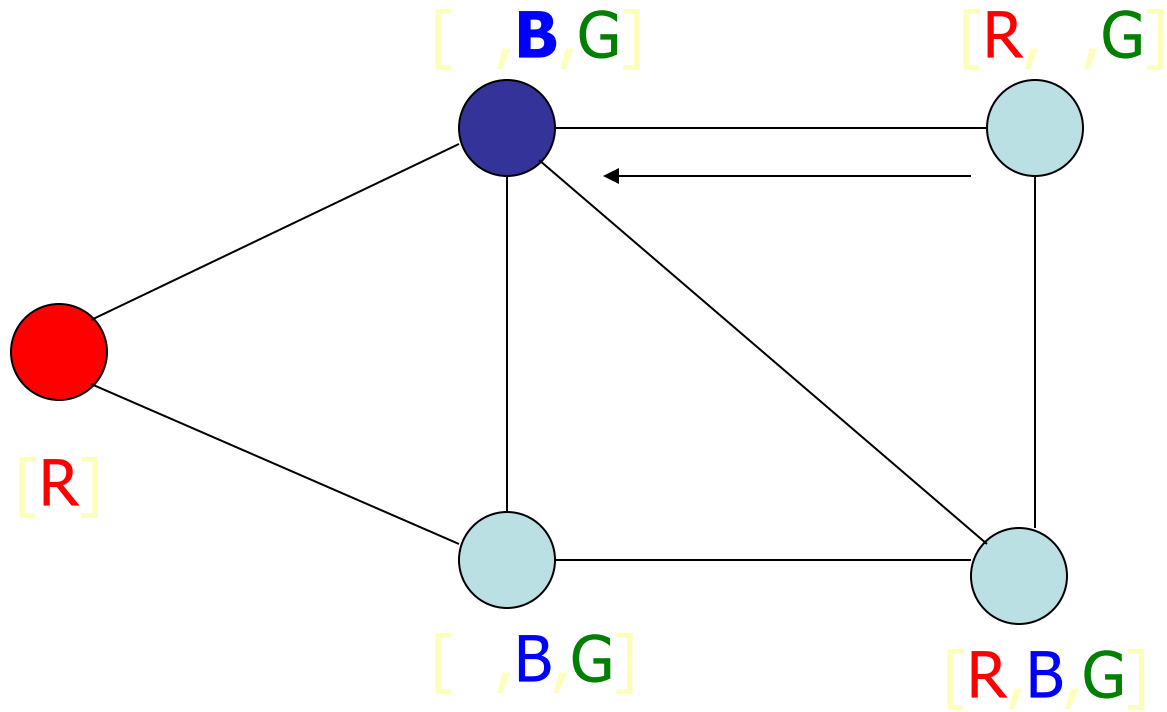
Arc Consistency: AC3



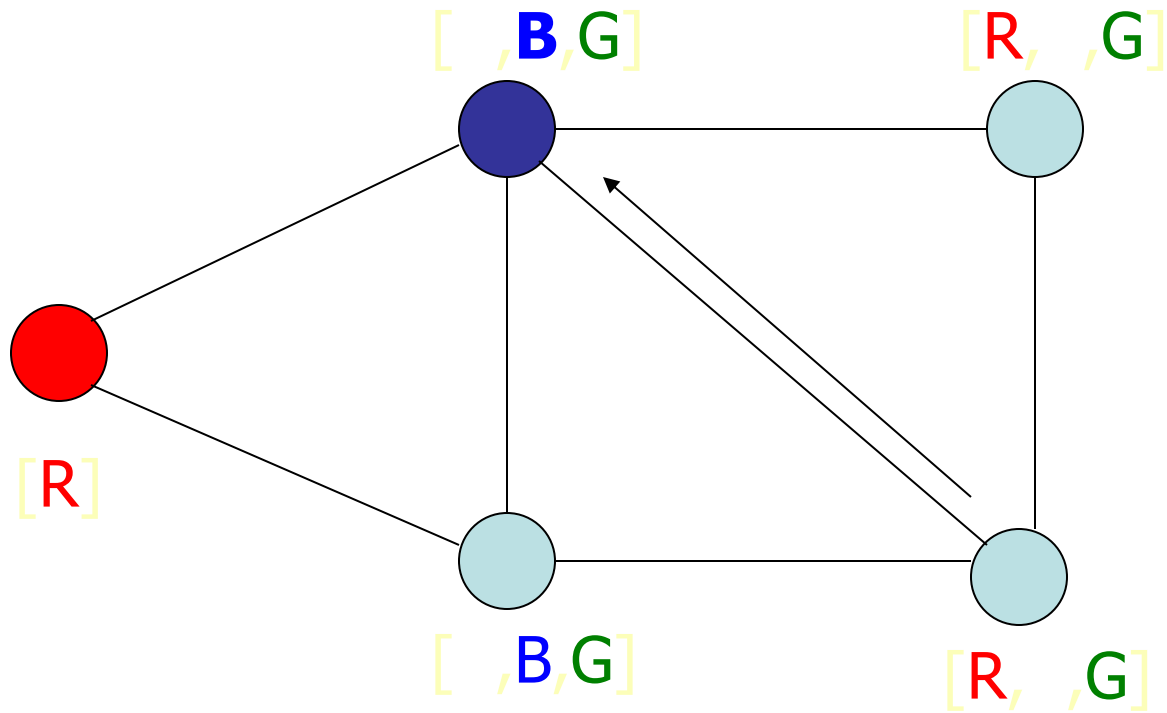
Arc Consistency: AC3



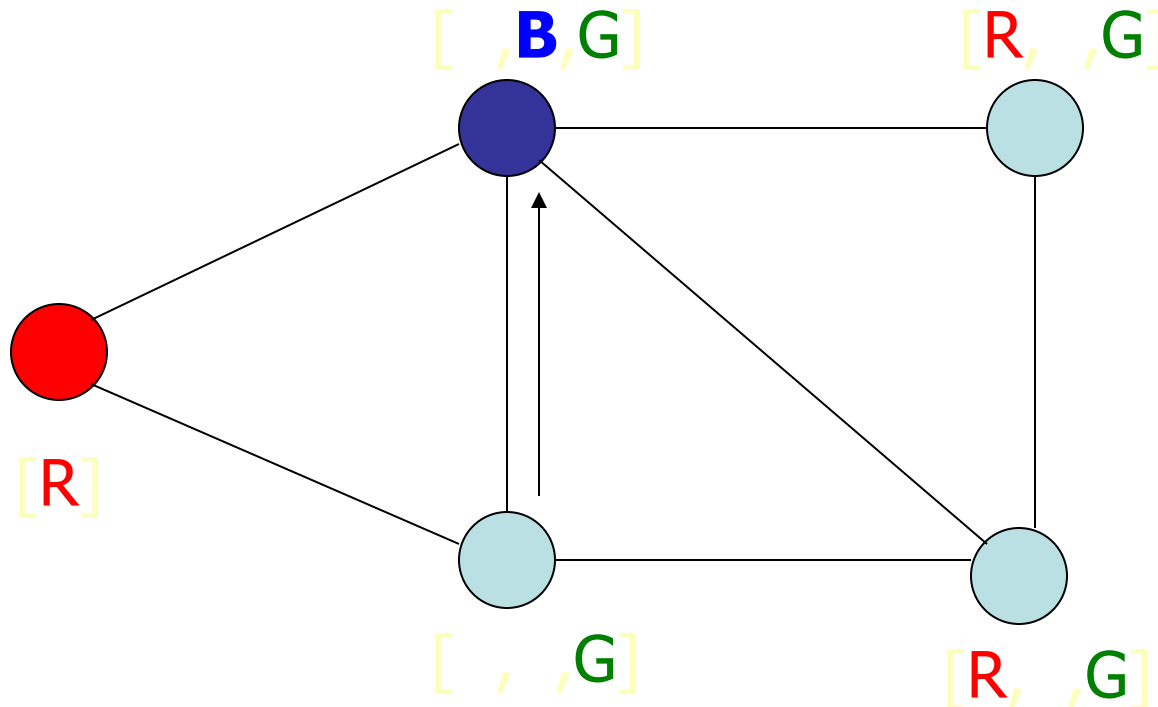
Arc Consistency: AC3



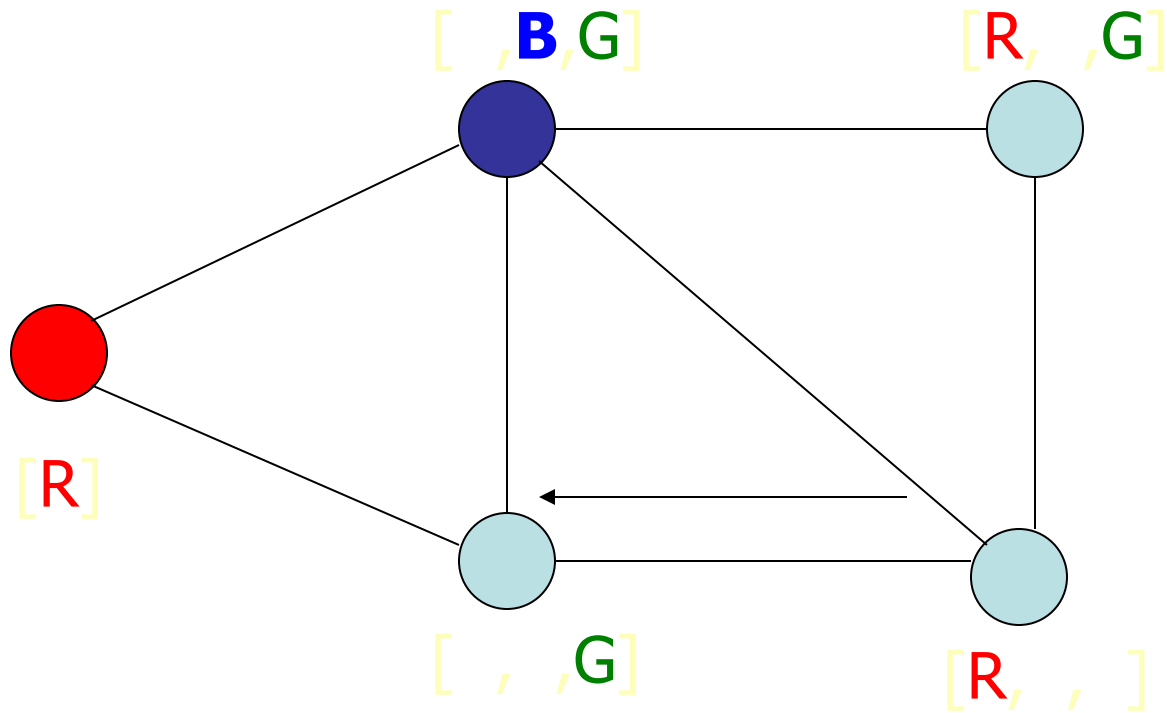
Arc Consistency: AC3



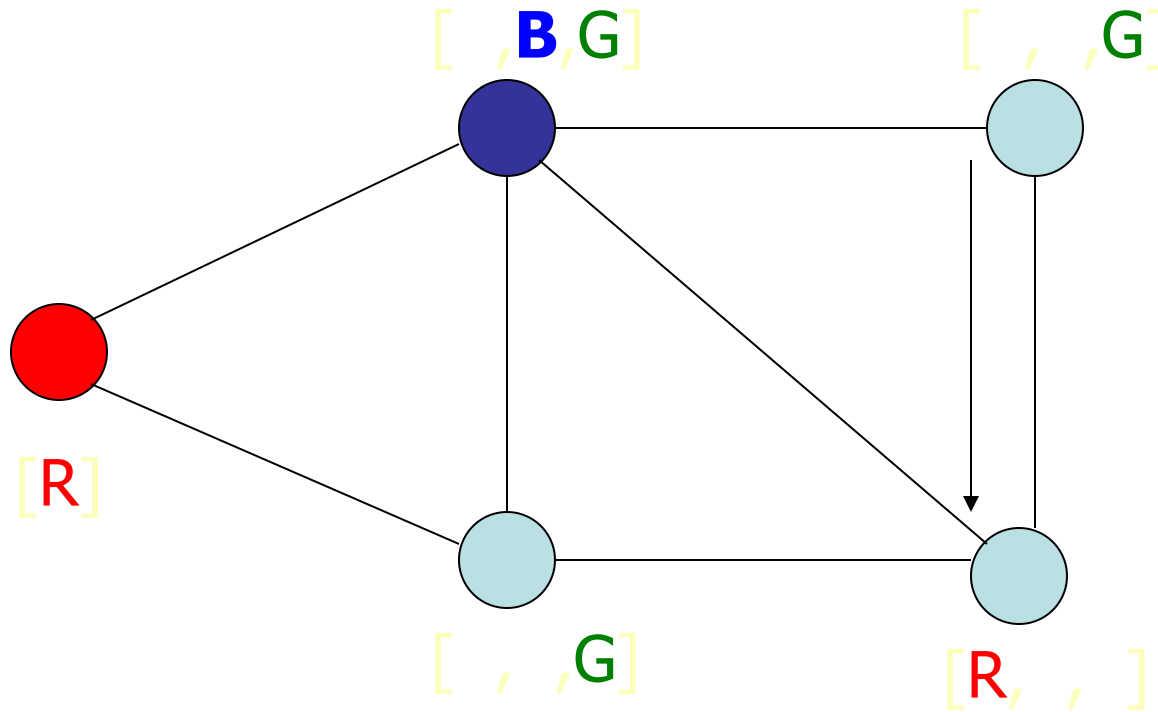
Arc Consistency: AC3



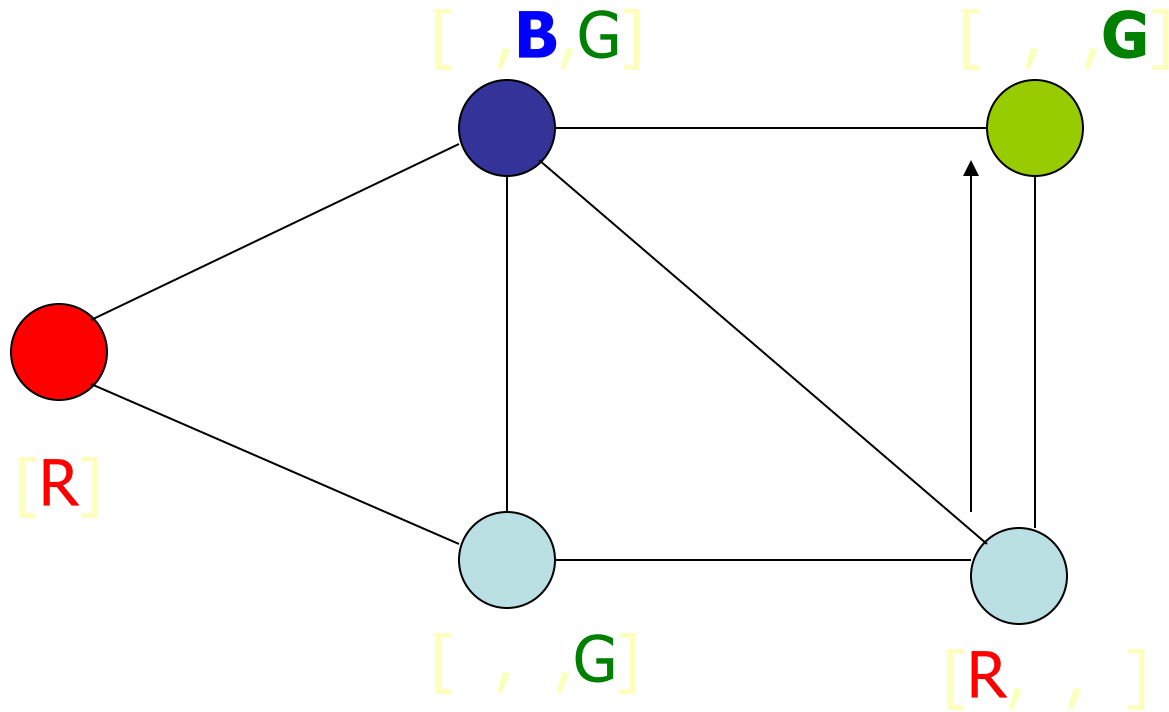
Arc Consistency: AC3



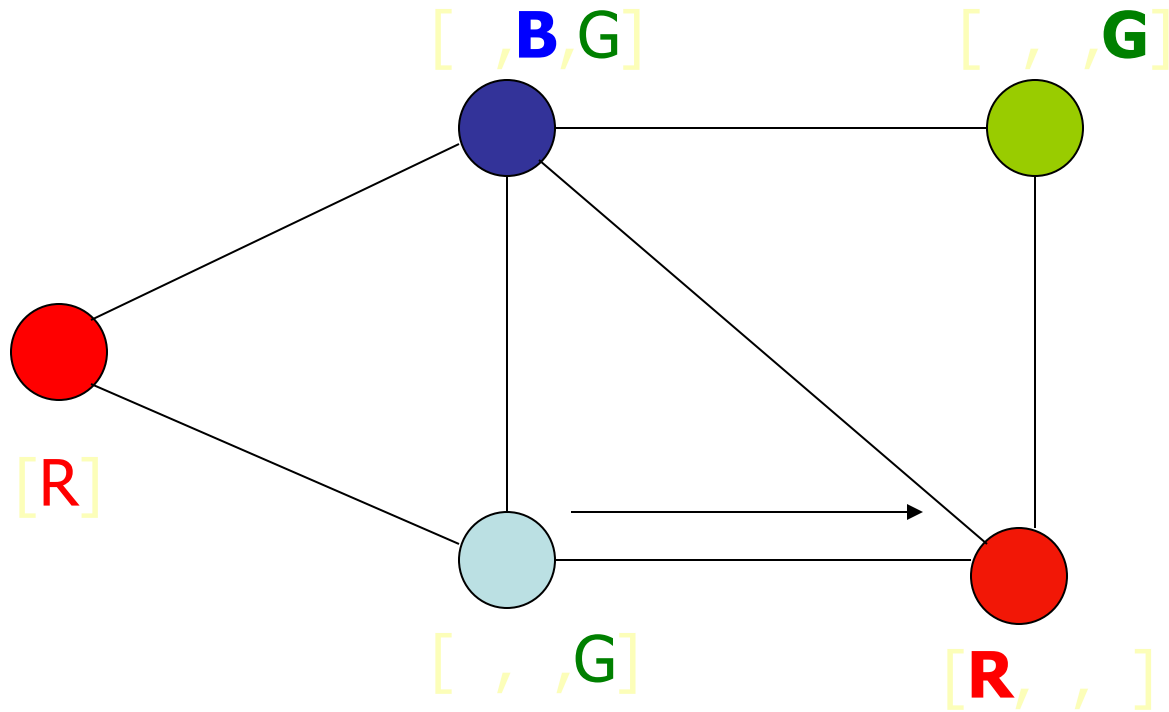
Arc Consistency: AC3



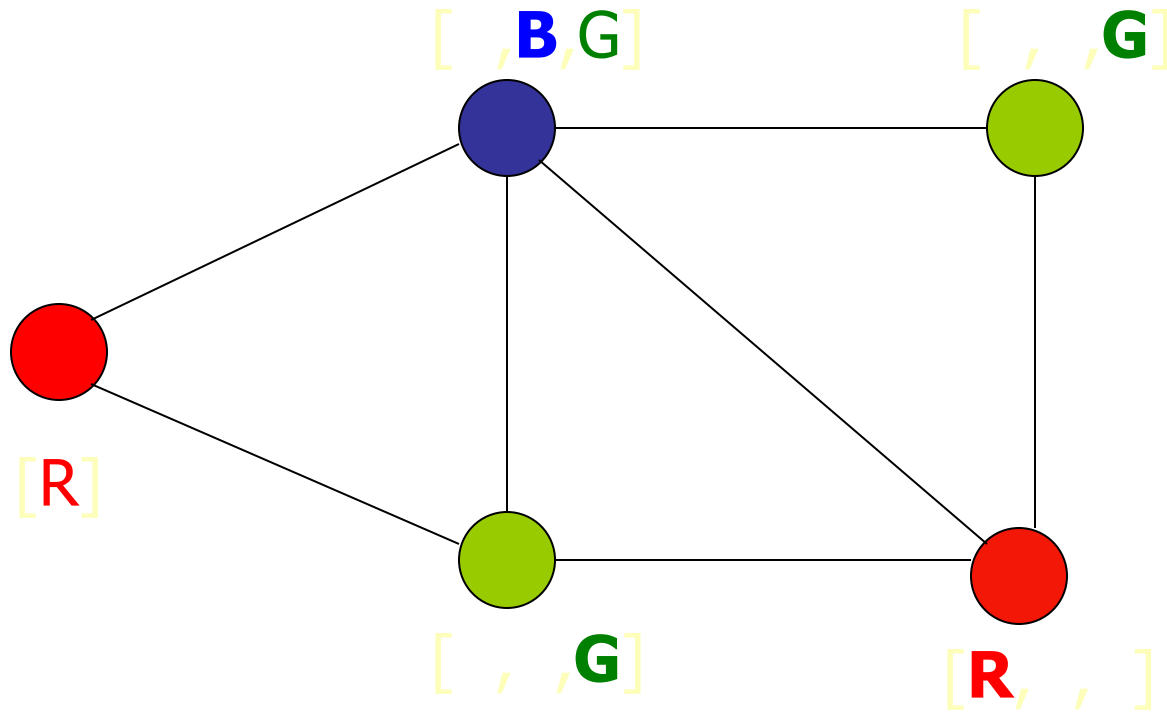
Arc Consistency: AC3



Arc Consistency: AC3



Arc Consistency: AC3



Solution !!!

Local Search and CSP

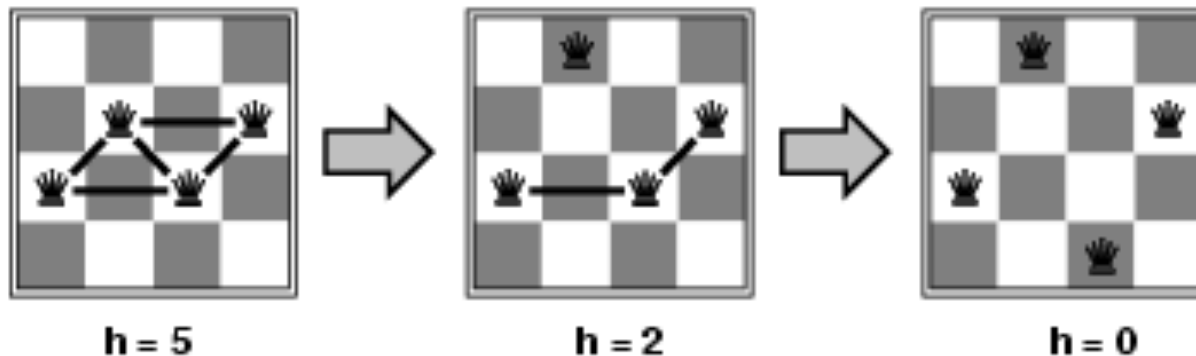
- local search (iterative improvement) is frequently used for constraint satisfaction problems
 - values are assigned to all variables
 - modification operators move the configuration towards a solution
- often called heuristic repair methods
 - repair inconsistencies in the current configuration
- simple strategy: min-conflicts
 - minimizes the number of conflicts with other variables
 - solves many problems very quickly
 - million-queens problem in less than 50 steps
- can be run as **online** algorithm
 - use the current state as new initial state

Local search for CSPs

- Hill-climbing, simulated annealing typically work with "complete" states, i.e., all variables assigned
- To apply to CSPs:
 - allow states with unsatisfied constraints
 - operators **reassign** variable values
- Variable selection: randomly select any conflicted variable
- Value selection by **min-conflicts** heuristic:
 - choose value that violates the fewest constraints
 - i.e., hill-climb with $h(n)$ = total number of violated constraints

Example: 4-Queens

- **States:** 4 queens in 4 columns ($4^4 = 256$ states)
- **Actions:** move queen in column
- **Goal test:** no attacks
- **Evaluation:** $h(n)$ = number of attacks



- Given random initial state, can solve n -queens in almost constant time for arbitrary n with high probability (e.g., $n = 10,000,000$)