

01

Recursion and Stack



Fatima Khan
Omer Shamsi
Areeba Zehra Jaffri
Faiq Ahmed

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Part 1

Recursion



```
sum(5) = 5 + sum(4)
↓
sum(4) = 4 + sum(3)
↓
sum(3) = 3 + sum(2)
↓
sum(2) = 2 + sum(1)
```

Definition

The process in which a function calls itself directly or indirectly is called recursion and the corresponding function is called a recursive function.

Recursion breaks down problems into simpler sub-problems and hence, finds optimal solution for a problem.



Simplicity

Recursive solutions are often shorter and easier to understand (e.g., factorial, Fibonacci, tree traversals).



Reduced Code Size

Recursion can replace loops and complex logic, making code more concise.



Higher memory usage

Takes up more memory, hence, its not space efficient.

You need to identify two cases to solve a recursive problem

Base Case: the case where your recursive call is stopped.

Recursive case: invokes the recursive call.

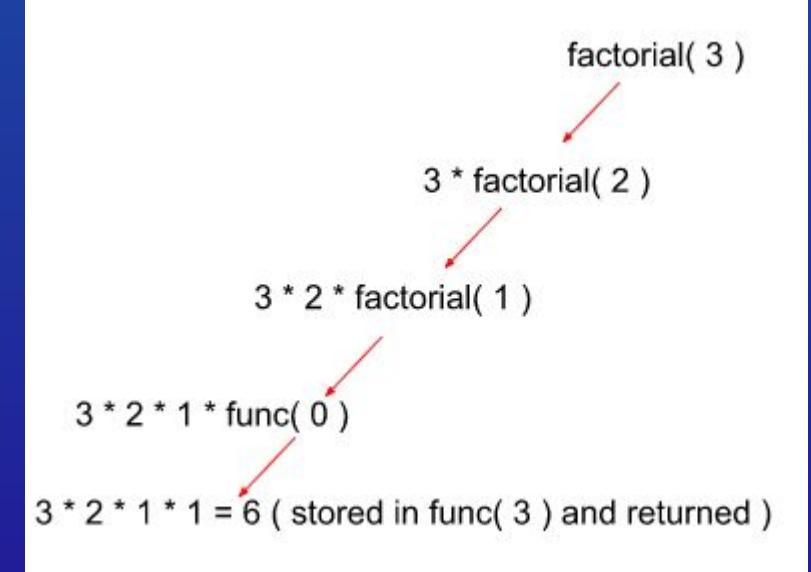
Algorithm (finding factorial of a number)

1. If the number is equal to one or zero, stop computing

2. Otherwise calculate factorial num*factorial(num-1)

Pseudocode (finding factorial of a number)

```
if num==0 or num==1  
    return 1  
else  
    return num*factorial(num-1)
```



In this example, base case will be applied to (number==0 or number==1)

Recursive case will be applied to all values except 0 and 1.

Non tail recursion

Tail recursion is defined by having the recursive call as the last operation in the function before returning

```
// Tail recursive GCD function
int gcd(int a, int b) {
    if (b == 0)
        return a;          // base case
    return gcd(b, a % b); // recursive call is the last operation
}
```

Tail recursion

Tail recursion is when other operations are invoked after the recursive call.

```
// Non-tail recursive factorial
int factorial(int n) {
    if (n == 0) return 1;
    return n * factorial(n - 1); // recursive call is not the last
                                operation
}
```

Stack

Stack is a linear data structure that can be accessed only at one of its ends for storing and retrieving data. It follows LIFO. Example a stack of plates.



Operations of stack

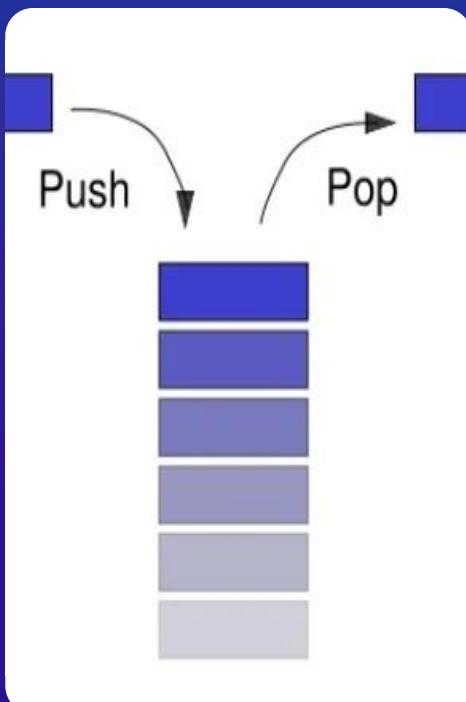
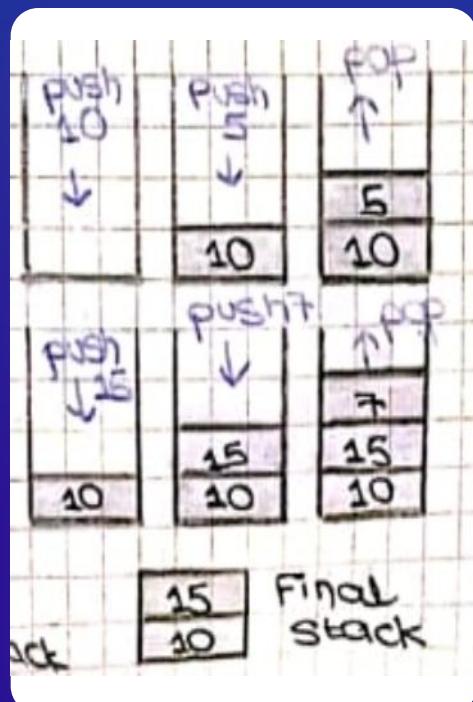
Clear() – To clear the stack

IsEmpty() – Check If Stack Is Empty

Push(EI) – Pushing element on top of stack

Pop()- Remove from stack

TopEI() – Peek the top element



```
void Clear() {  
    top = -1;  
    cout << "Stack cleared!" << endl;  
}  
  
// IsEmpty() - Check If Stack Is Empty  
bool IsEmpty() {  
    return (top == -1);  
}  
  
// Push(EI) - Pushing element on top of stack  
void Push(int EI) {  
    if (top == capacity - 1) {  
        cout << "Stack Overflow!" << endl;  
    } else {  
        arr[++top] = EI;  
        cout << EI << " pushed into stack." << endl;  
    }  
}  
  
// Pop() - Remove from stack  
void Pop() {  
    if (IsEmpty()) {  
        cout << "Stack Underflow!" << endl;  
    } else {  
        cout << arr[top--] << " popped from stack." << endl;  
    }  
}  
  
// TopEl() - Peek the top element  
int TopEl() {  
    if (IsEmpty()) {  
        cout << "Stack is empty!" << endl; return -1;  
    } else { return arr[top]; } }
```

Applications Of stack

1. Delimiter Matching
2. Postfix Expression
3. Tower Of Hanoi



Delimiter Matching

- Stack helps solve complex mathematical equations

```
( a + b ) * { c + [ d - e ] }  
[ { ( x + y ) * z } ]  
a * ( b + c )  
( a + b ] * { c + d )      // mismatched  
[ ( a + b )           // missing closing ]  
a + b ) * ( c + d       // extra closing )
```

Delimiter Matching PseudoCode

```
function isvalid delimiterstring(s):  
stack = empty stack  
for each character in s:  
if character is '(', '{', '[':  
push character onto stack  
else if character is ')', '}', ']':  
if stack is empty:  
return FALSE  
top = pop from stack  
if character doesn't match the type of top:  
return FALSE  
return stack is empty
```

Postfix Expression

The **use of postfix expression** is mainly in **computer systems and compilers** because it makes evaluation easier and faster.

$(5 + 3) * 2$ will turn to $5\ 3\ +\ 2\ *$

Postfix Expression Pseudocode

```
function evaluatePostfix(expression):
```

```
    stack = empty stack
```

```
    for each token in expression:
```

```
        if token is an operand:
```

```
            push token onto the stack
```

```
        else if token is an operator:
```

```
            Operand2 = pop from stack
```

```
            Operand1 = pop from stack
```

```
            RESULT = Apply operator on Operand1 and
```

```
            Operand2
```

```
            Push result onto stack
```

```
    return Pop from stack
```

pop(op2)
pop(op1)

3		
2		6
1		1

push

switch (expression[i])
case '+':

Operand1 =

stack.pop() +
 stack.pop();

Operand: The values or variables on which an operation is performed. Example values: numbers (2, 5, 10) or variables (A, B, X).

Operator: The symbol that tells what operation to perform on operands. Example operators:

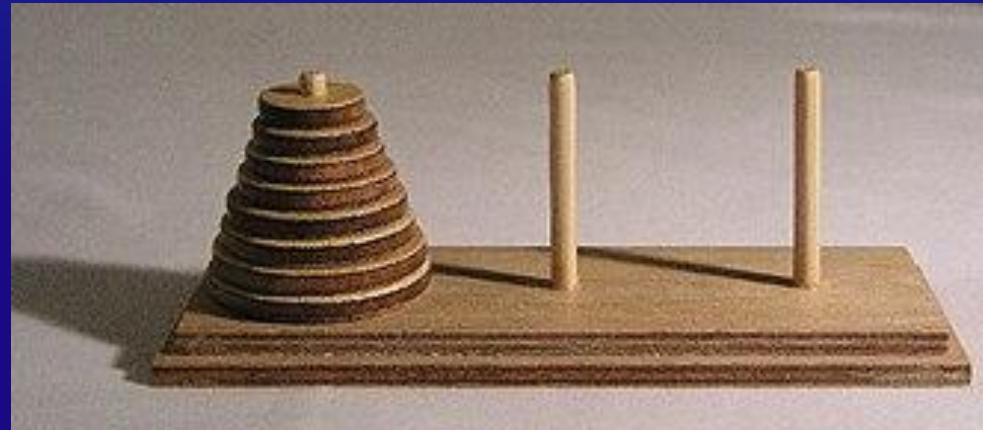
Arithmetic: + , - , * , /

Relational: > , < , == , !=

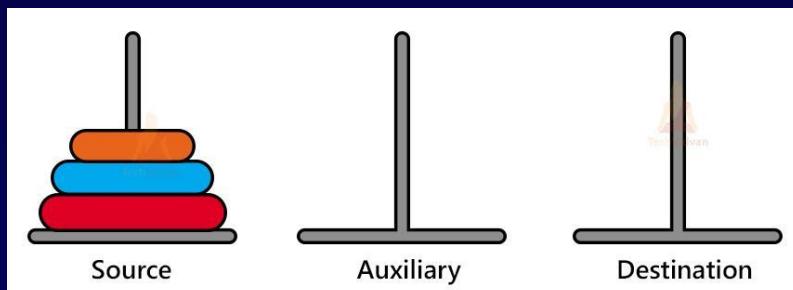
Logical: && , || , !

Tower Of Hanoi

The tower of Hanoi is a classic problem in computer science that involves moving a set of sticks from one rod to another, following specific rules.

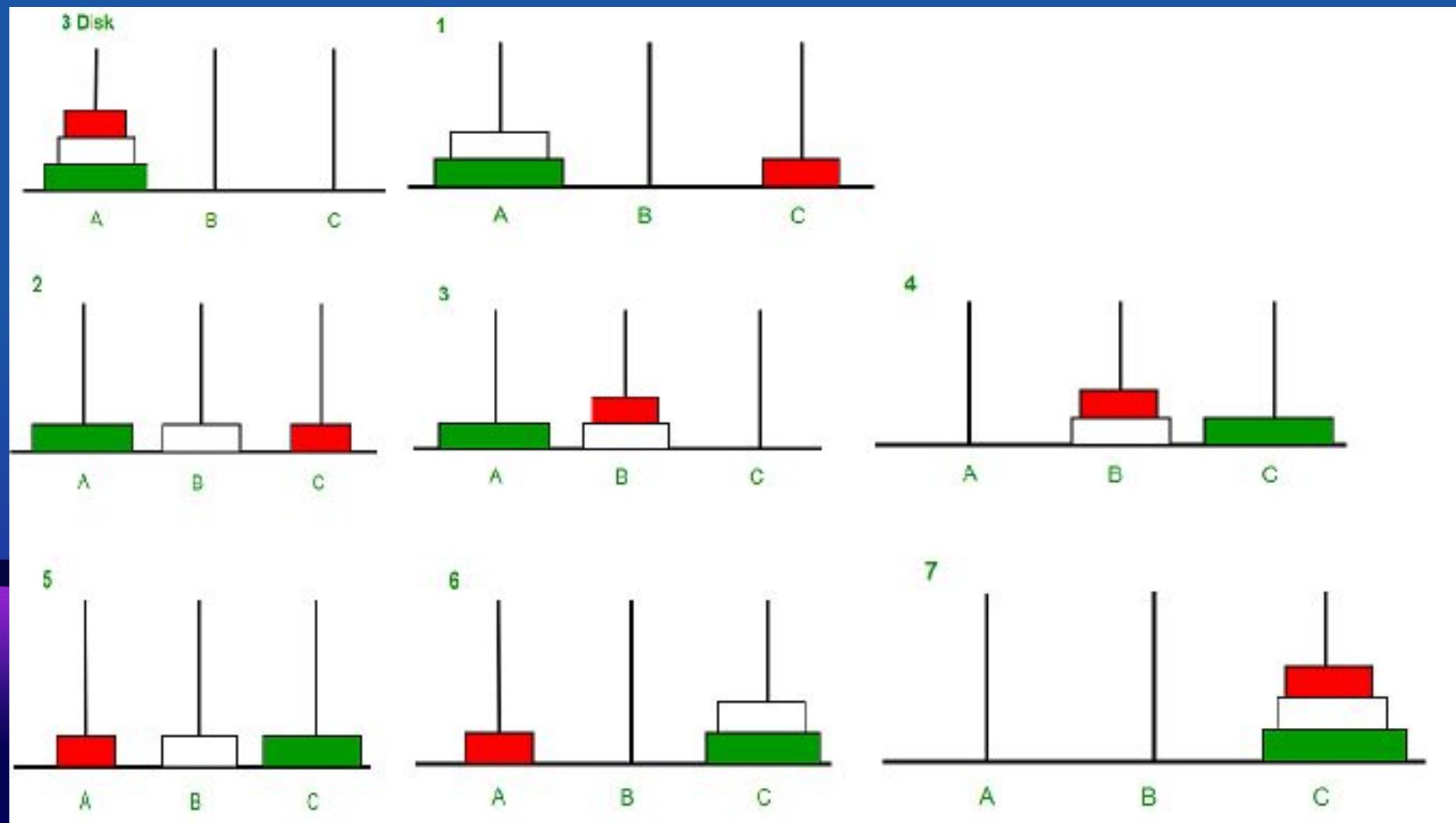


Rules Of Tower Of Hanoi



Constraints

1. Only one disk can be moved at a time
2. A disk can only be placed on top of a larger disk
3. All disks must be moved from source rod to destination rod



Leetcode Questions



921. Minimum Add to Make Parentheses Valid

A parentheses string is valid if and only if:

- It is the empty string,
- It can be written as `AB` (`A` concatenated with `B`), where `A` and `B` are valid strings, or
- It can be written as `(A)`, where `A` is a valid string.

You are given a parentheses string `s`. In one move, you can insert a parenthesis at any position of the string.

- For example, if `s = "())"`, you can insert an opening parenthesis to be `"((()))"` or a closing parenthesis to be `"(())())"`.

Return the minimum number of moves required to make `s` valid.

Example 1:

Input: `s = "()"`
Output: 1

Example 2:

Input: `s = "((("`
Output: 3

Constraints:

- `1 <= s.length <= 1000`
- `s[i]` is either `'('` or `')'`.

1021. Remove Outermost Parentheses

A valid parentheses string is either empty "", "(" + A + ")", or A + B, where A and B are valid parentheses strings, and + represents string concatenation.

- For example, "", "()", "(()())", and "(((())))" are all valid parentheses strings.

A valid parentheses string s is primitive if it is nonempty, and there does not exist a way to split it into $s = A + B$, with A and B nonempty valid parentheses strings.

Given a valid parentheses string s , consider its primitive decomposition: $s = P_1 + P_2 + \dots + P_k$, where P_i are primitive valid parentheses strings.

Return s after removing the outermost parentheses of every primitive string in the primitive decomposition of s .

Example 1:

Input: $s = "((())())()$

Output: $"()()()$

Explanation:

The input string is $"((())())()$, with primitive decomposition $"((())") + "(()())"$.

After removing outer parentheses of each part, this is $"()()" + "()" = "()()()$.

Example 2:

Input: $s = "((())())((())())()$

Output: $"()()()()((())())()$

Explanation:

The input string is $"((())())((())())()$, with primitive decomposition $"((())") + "(())" + "((())())"$.

After removing outer parentheses of each part, this is $"()()" + "()" + "()((())") = "()()()((())())()$.

Constraints:

- $1 \leq s.length \leq 10^5$
- $s[i]$ is either '(' or ')'.

- s is a valid parentheses string.

02

Thank you

