Assignment SMV

CS4211 - Formal Methods for Software Engineering

October, 2021

Instructions



Note:

- This assignment is due before 9:59 PM, Sunday, 14th November, 2021
 No late submissions!
- This is an individual assignment. Acts of plagiarism are subjected to disciplinary action by the university
- Create a folder named your matriculation number YourMatricNumber, e.g. U123456M.
- Create the following files in this folder (name these files exactly as instructed.)
 - Assignment
 - Problem 1: Include your translated NuSMV source file needed
 - Problem 2: Include your implemented NuSMV source file needed
 - Problem 3: Include your implemented NuSMV source file needed
 - Report.PDF
 - ReadME.txt
- Zip (using WinZip) the entire YourMatricNumber folder (including the folder itself and all files in it) into a file YourMatricNumber.zip.
- Submit YourMatricNumber.zip to the Luminus Workbin Folder



Problem 1

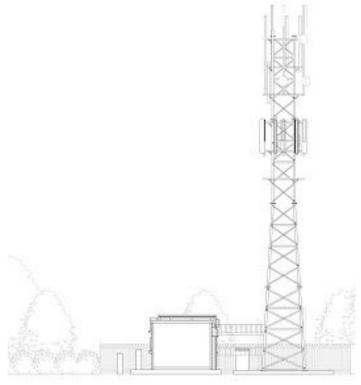


Entities

- 2 Communication Stations
- 4 Operators

Abstracted Information

Data/Information Communicated



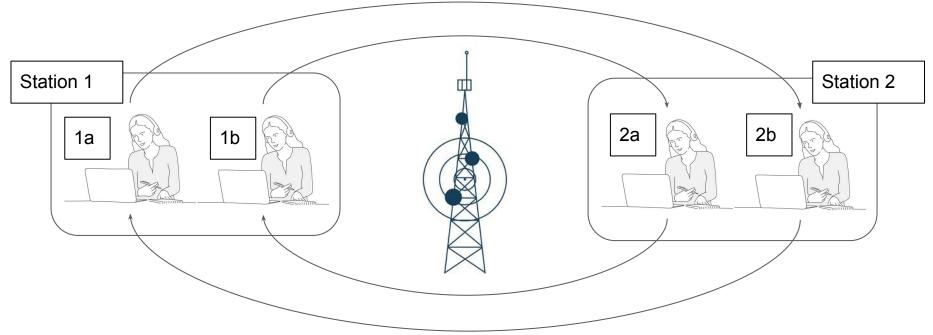


- Consider two stations connected in a ring where each station can both send and receive data. Consequently, we need to assign two operators to each station.
- Let the stations be 1 and 2, and the operators be 1a, 1b, 2a, 2b. Operator 1a handles data-communication from station 1 to station 2, while operator 1b handles data-communication from station 2 to station 1. Similarly for operators 2a and 2b.



- We may not always perform communication from station 1 to station 2 via a single channel (from 1 to 2).
- For example, if the protocol involves transferring data (from 1 to 2) followed by acknowledgment(from 2 to 1), then we might want to have separate channels for the actual data items and the acknowledgment signal.







```
#define true 1
#define false 0
#define limit 3

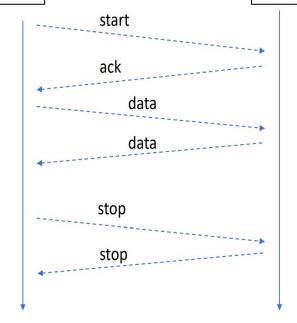
bool busy[2];
int count = 0;

mtype = {start, stop, data, ack};
chan up[2] = [1] of { mtype };
chan down[2] = [1] of { mtype };
```



```
proctype operator(byte id; chan in, out) {
     do
          in?start -> atomic { !busy[id] -> busy[id] = true };
           out!ack;
           do
                 :: in?data -> out!data
                 :: in?stop -> break
           od:
           out!stop;
           busy[id] = false
          atomic { !busy[id] && !busy[1-id] -> busy[id] = true };
           out!start:
           in?ack;
           do
                 :: (count < limit) -> out!data; count++ ; in?data
                 :: (count <= limit) -> out!stop; count = 0; break
           od;
           in?stop;
           busy[id] = false
     od
```

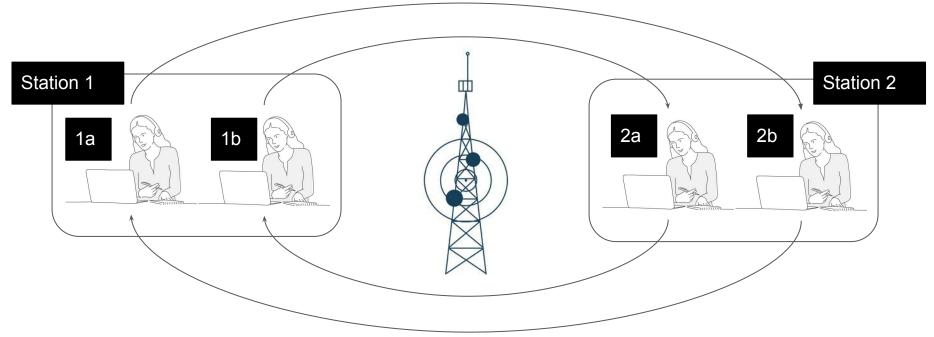
1a 2b



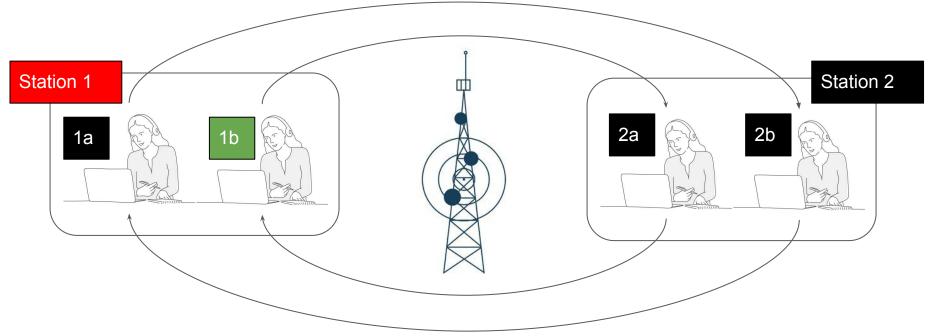


```
init {
    atomic {
        run operator(0, up[1], down[1]);
        run operator(1, up[0], down[0]);
        run operator(0, down[0], up[0]);
        run operator(1, down[1], up[1]);
    }
}
```

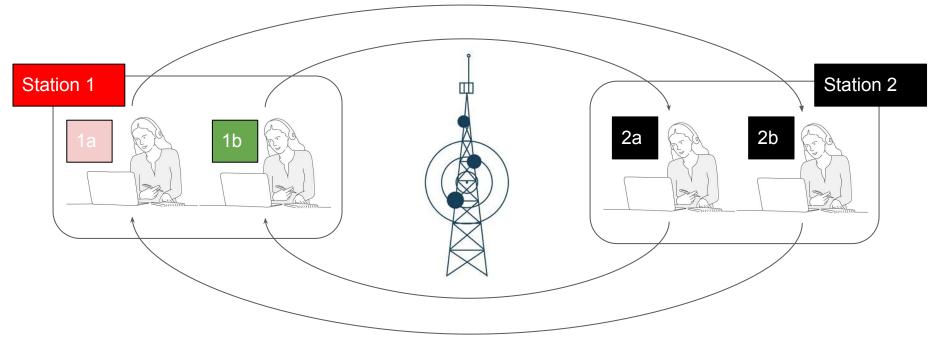




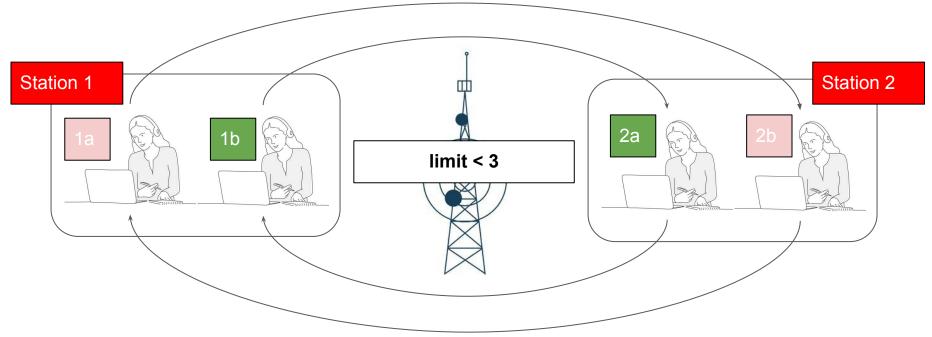




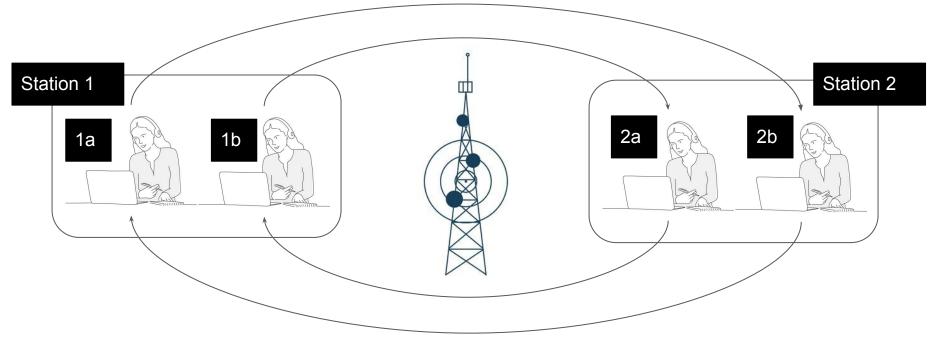














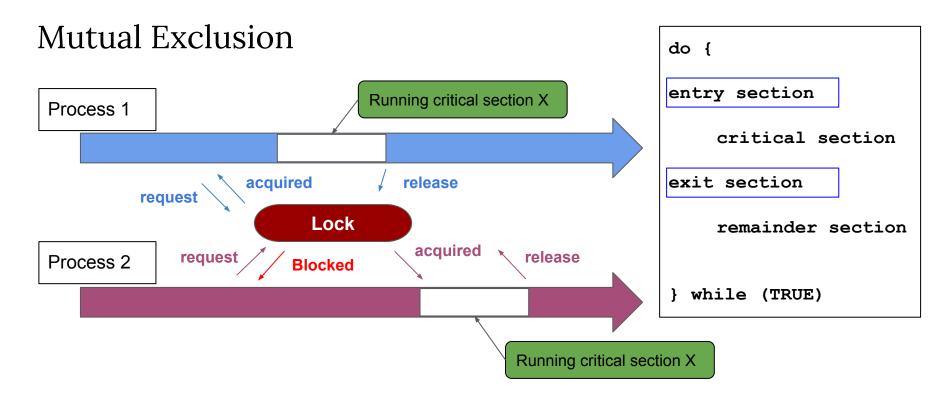
Questions [7 Marks]

- A. Translate above Promela code into SMV code while adhering the protocol defined.
 (5 marks)
- B. Formalize the property that no communication between stations is aborted in the above protocol, i.e. any communication that is started (with a start signal) is finished (with a stop signal). Use Linear time temporal logic (LTL) to express your property description and show whether it holds true using NuSMV. (3 marks)



Problem 2







Algorithm for Mutual Exclusion

```
loop forever {
    /* non-critical section */
    key1 := false;
    while (key1 = false) do
        swap(key1, lock); // atomic operation
    /* critical section */
    lock := true
```



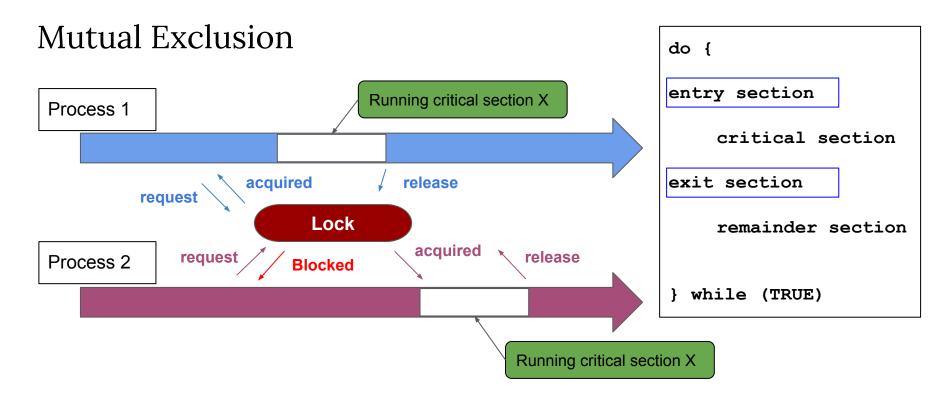
Questions [4 Marks]

- A. Using NuSMV model checker show that the above solution preserves mutual exclusion.
 - a. Write a SMV model to implement above solution(1 mark)
 - b. Using a LTL formula verify whether above solution preserves mutual exclusion (1 marks)
- B. Using a Linear time temporal logic (LTL) property show that this is not a good solution to the mutual exclusion problem. Show your proof using NuSMV and explain your answer (2 marks)



Problem 3







Algorithm for Mutual Exclusion

```
C0: b[i] := true;
    k = 1 - i;
C1: if k != i \text{ and } b[1-i]
                 then go to C1;
    else /*critical section*/;
    b[i] := false;
    /* non-critical section */;
    go to CO;
```



Questions [4 Marks]

- A. Using NuSMV model checker show that the above solution preserves mutual exclusion.
 - a. Write a SMV model to implement above pseudo-code (1 mark)
 - b. Using a LTL formula verify whether above solution preserves mutual exclusion (1 marks)
- B. The above candidate solution is starvation-free.
 - a. Write a LTL formula to show the solution is starvation free (1 mark)
 - b. Show proof using NuSMV and explain your answer (1 marks)



Q&A

