

Required Algorithms for Heapsort

The Heapsort algorithm involves two primary steps:

1. *Build a Max Heap*: Transform the input array into a max heap.
2. *Sort the Array*: Repeatedly extract the largest element from the heap and move it to the end of the array, reducing the heap size.

Algorithm 1: Heapify

Heapify(A, n, i)

Input:

A is the array, n is the size of the heap, i is the index to heapify.

Output:

Maintains the max-heap property for the subtree rooted at i.

1. Initialize largest = i.
2. Set left = $2i + 1$ (left child index).
3. Set right = $2i + 2$ (right child index).
4. If left < n and $A[\text{left}] > A[\text{largest}]$, set largest = left.
5. If right < n and $A[\text{right}] > A[\text{largest}]$, set largest = right.
6. If largest != i:
 - a. Swap $A[i]$ and $A[\text{largest}]$.
 - b. Recursively call Heapify(A, n, largest).

Algorithm 2: BuildMaxHeap

BuildMaxHeap(A, n)

Input: A is the array, n is the size of the array.

Output: Constructs a max-heap from the input array.

1. For $i = n/2 - 1$ down to 0:
 - a. Call Heapify(A, n, i).

Algorithm 3: HeapSort

HeapSort(A)

Input: A is the array to sort.

Output: Sorted array A.

1. Call BuildMaxHeap(A, n).
2. For $i = n - 1$ down to 1:
 - a. Swap $A[0]$ and $A[i]$.
 - b. Call Heapify(A, i, 0).

b. Analysis of Heapsort Algorithms

Time Complexity:

1. Heapify Operation:

The heapify operation runs in $O(\log n)$, as the height of the heap is proportional to $(\log n)$.

2. BuildMaxHeap Operation:

Building the max heap involves calling Heapify for each non-leaf node, resulting in $O(n)$ time.

3. Sorting the Heap:

Extracting the max element and re-heapifying for n elements takes $O(n \log n)$.

Overall time complexity: $O(n \log n)$.

Space Complexity:

Heapsort operates in-place and requires no additional memory, resulting in $O(1)$ space complexity.