COMS W3137 - Recitation Notes Week 3: Functional Programming Basics

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1 for loops and expressions

1.1 Iteration with Side Effects

We have previously discussed while loops. While loops require updates to a var, so that the looping condition can eventually become false. Scala also support for loops. At first glance, for loops behave like enhanced for loops in Java: We can iterate over the items in an iterable data structure. This does not require vars, but for loops always have side effects, such as printing output or adding to a mutable data structure.

1.2 for comprehensions

for loops can also be used as expressions, using the yield keyword. This is called a *comprehension*. When we use a for comprehension on a list, the

expression evaluates to a list (of the same generic type), without requiring any explicit change of state.

```
scala> for (x : Int <- li) yield (x * 2)
val res0: List[Int] = List(2, 4, 6, 8, 10)
```

The yield keyword tells Scala to aggregate the result if the expression $\mathbf{x} = \mathbf{2}$ in a new list. We can also include a filter. Only values that pass the filter condition are included in the resulting list.

```
Also works: for (x : Int <- li) yield (if (x%2 == 0) x * 2 else x)
```

Exercise: Without for-comprehensions or some of the other abstractions discussed below, iterating through a list in a side-effect free and expression oriented way would be tedious: It would require recursion. Write a function def double(li: List[Int]): List[Int] that uses only recursion (no while or for loops allowed) to construct a new list in which each integer has been doubled.

NB: `::` creates a new list, whereas `:+` appends to an existing list

2 map, filter

Rather than having to traverse data structures using recursion, or using a for loop, data structures such as lists provide a map method. This method takes a function object as its parameter, applies it to each element in the data structure and returns a new data structure.

```
| scala > def timestwo(x : Int) : Int = x * 2
| scala > li.map(timestwo)
| res0: List[Int] = List(2, 4, 6, 8, 10)
```

```
Also works: li.map((x : Int) => x * 2)
```

A related method, filter, returns a new collection that includes only the elements for which a certain condition is true. Like map, filter takes a function object as a paremeter. This function object must return a boolean. This function is applied to each element in the collection. If the function return true, the element is included in the result, otherwise it is skipped.

```
def isEven(x : Int) = x % 2 == 0
isEven: (x: Int)Boolean

scala> li.filter(isEven)
res1: List[Int] = List(2, 4)
```

```
Also works: li.filter((x : Int) \Rightarrow x % 2 == 0)
```

Can chain these, ie: li.filter((x : Int) => x % 2 == 0).map((x : Int) => x + 2)

Can simplify more by not including type annotation: li.filter(x => x % 2 == 0).map(x => x + 2)

The map and filter method are defined in the trait scala.collection.TraversableLike, which is extended by almost all scala data structures.

map and filter are examples of higher-order functions. A higher-order function is any function that either expects a function object as a parameter, or returns a function object.

3 Function literals / val function

So far we have defined functions with def. Unless def is used in the REPL, such definitions always appear in the context of a class or object. As a result, def is used to define *methods*. A method always has an implicit reference to the class instance or object it is called on.

An alternative way to define a function is as a function literal, using =>.

```
\begin{array}{l} {\rm scala}{\rm >\ val\ incr} = ({\rm x\ :\ Int}) \implies {\rm x\ +\ 1\ :\ Int} \\ {\rm incr:\ Int} \implies {\rm Int} = {\rm s$Lambda}{\rm s$1218/0}{\rm x}0000000800781840@1774c4e2 \end{array}
```

Function literals are also called lambdas in other languages, including more recent versions of Java (≥ 1.8). They are often used to define a small function in-place.

The (x: Int) on the left hand side of the => specifies the parameter list of the function. The: Int at the end of the line specifies the return type. In many cases, we do not have to make the return type explicit and Scala's type inference can infer it automatically. x+1 is the function body, which can be any expression involving the parameters. The type of the function incr is Int=>Int. The function takes an integer as a parameter and returns an integer.

Once the function is defined, we can call it as usual.

```
scala > incr(5)
res4: Int = 6
```

3.1 Partially applied functions

A partially applied function is a function in which some of the arguments have already been provided. Consider for example the following function add.

```
\begin{array}{l} scala > val \ add = (x : Int, \ y : Int) \Rightarrow x + y \\ add: (Int, \ Int) \Rightarrow Int = \$Lambda\$\$1176/0x000000080074f840@5fffb692 \end{array}
```

We can now provide only one of the parameters, leaving the other one to be completed later.

If the parameters of the function appear in the same order in the function body, the following notation can be used as a short-cut in function literals. This simplifies the definition of partially applied functions:

```
scala> val add2 = add(_,2)
add2: Int => Int = $$Lambda$$1240/0x0000000800784840@3ed87b6e
```

We can also use this technique to convert def methods in to val functions. This process is also called *eta expansion*.

```
scala> def add(x : Int , y : Int) = x+y
add: (x: Int , y: Int)Int
scala> val addToo = add(_ , _ )
addToo: (Int , Int) => Int =
    $$Lambda$1242/0x0000000800786040@55e3b64d
```

A single trailing _ can be used to replace the entire parameter list of a function.

```
scala> val addToo = add _
addToo: (Int, Int) => Int =
$$Lambda$1242/0x0000000800786040@55e3b64d
```

3.2 Currying

Typically, when a function is called, all arguments to the function need to be provided. It is also possible to define a function that accepts only some of its parameters at a time. Such a function is called a curried function. ¹ Technically, a curried function is a higher-order function because it returns another function.

¹Currying functions are named after the American mathematician Haskell Curry, but also called Schönfinkeling after the Russian mathematician Moses Schönfinkel

Note that the x in the nested function (y: Int) => x+y is the parameter of the surrounding function. Holding on to the local vairables of a function in a nested function is called a *closure*. Closures appear commonly in functional programming as a way to store program state. The state is preserved in the function object that is returned from the outer function.

For def methods there is special notation to define a curried method by specifying multiple parameter groups. This does not work with function literals.

```
scala> def curriedAdd(x :Int)(y:Int) = x+y
curriedAdd: (x: Int)(y: Int)Int
scala> curriedAdd(1)(2)
res0: Int = 3
```

We can verify that passing only one parameter actually returns a method (which we need to eta convert to a function):

Or we can eta expand the entire method:

```
scala> curriedAdd _
res6: Int ⇒ (Int ⇒ Int) =
$$Lambda$1141/0x0000000800737840@2c08c787
```

Exercise: Write a function curry3(f : (Int,Int,Int) => Int) : (Int=>(Int=>Int))) that returns a curried version of the function f. For an additional challenge, can you make the function generic so that it works for any function with three parameters (regardless of their type)?

3.3 Example: Function Composition and Repeated Application

Given two mathematical functions f and g, the function composition $f \circ g$ ("g following/afer f") produces a function h(x) = f(g(x)). In other words, the function h first computers g(x) and then applies f to the result. Assume we are given functions f: Int=>Int and g: Int=>Int. Write a function combine: (Int=>Int, Int=>Int) => (Int=>Int). Calling combine(f,g) should return a function h: Int=>Int, so that $h = f \circ g$. solution:

```
val combine = (f : Int \Rightarrow Int, g:Int \Rightarrow Int) \Rightarrow (x:Int) \Rightarrow f(g(x))
```

or using a polymorphic (generic) function (note that there are no polymorphic function literals in Scala):

```
def combine [A,B,C](f : A \Rightarrow B, g: B \Rightarrow C) = (x:A) \Rightarrow g(f(x))
```

Now assume you wish to apply the same function k times. Given a function f and a positive integer k, the k-th repeated application of k is the function $h(x) = f(f(\cdots(f(x))\cdots))$. For example, the 2nd repeated application of the function f(x) = x + 1 is h(x) = f(f(x)) and h(2) = 3. We wish to define a Scala function repeat(k: Int, f: Int=>Int, k: Int): Int=>Int that returns thek-th repeated application of f.

solution: We can use recursion and the combine function we just defined.

Interestingly, we can use also define the repeat function as a curried function, like this:

```
\begin{array}{lll} \text{val repeat} & : & ((\text{Int}\Rightarrow \text{Int} \,,\,\, \text{Int}) \Rightarrow (\text{Int}\Rightarrow \text{Int})) & = & (k \,:\,\, \text{Int} \,,\,\, f \,:\,\, \\ & \text{Int}\Rightarrow \text{Int}) & \Rightarrow & k \,\, \text{match} \,\, \{ & \\ & \text{case} \,\, 1 \,\Rightarrow\, f & \\ & \text{case} \,\, - \, \Rightarrow & \text{combine}(f \,,\,\, \text{repeat}(k-1,\,\, f)) \\ \} \end{array}
```

This example illustrates how curried functions can be used as control structures.

```
scala> repeat(3){ (x : Int) \Rightarrow x+1 } (1) res20: Int = 4
```

4 foldLeft, foldRight, fold

Recall the map and filter methods described above. One thing that these methods have in common is that they separate the traversal through the list from the operation performed on the list element. We can generalize this idea even further.

Consider the task of computing the sum of a list. We could do this recursively by first computing the sum of the tail list, adding the head, and returning the value. In the base case, we are calling the recursive function on an empty list Nil and return 0.

We can use the same technique to compute all kinds of things on the list. We only need to specify 1) the return value for the base case and 2) a way to combine the return value of the recursion with the next element.

The foldRight method defined on the list allows you to do exactly that.

foldRight is a curried method with two parameter groups. The first one specifies the base case value. The second one specifies a function object that combines the return value and the next head to produce a new return value, as in the example before.

Here is the signature of that method:

```
foldRight[A,B](base\_case : A)(combiner : (B,A)=>A)
```

Using fold, we can rewrite the sum function as follows.

Also works: li.foldRight(0)((x, y) => x + y)

Interestingly, both map and filter can be defined using fold.

```
li.foldRight(0)((head, tailresult) => head + tailresult)
```

foldLeft starts on the left side—the first item—and iterates to the right; foldRight starts on the right side—the last item—and iterates to the left