

Homework 2

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1 Assignment 1

1.1 SVD

For every matrix $A \in \mathbb{R}^{m \times n}$, having rank r , there always exist its *Singular Value Decomposition*:

$$A = U \Sigma V^T \quad (1)$$

Where:

- The left singular vectors are the columns of the matrix U . They form an orthonormal basis in \mathbb{R}^n .
- The right singular vectors are the columns of the matrix V . They form an orthonormal basis in \mathbb{R}^m .
- $\Sigma = \text{diag}(\sigma_1, \sigma_2, \dots, \sigma_r)$, where $\{\sigma_i\}$ are the singular values.

1.2 Eigenvectors of AA^T and $A^T A$

For any rectangular matrix A , we can always say that:

- the left singular vectors are the eigenvectors of the matrix AA^T associated to the all non zero eigenvalues of AA^T .
- the right singular vectors are the eigenvectors of the matrix $A^T A$ associated to the all non zero eigenvalues of $A^T A$.

We used those values as columns of U and V to prove the existence and uniqueness of the *SVD* of A .

Alternatively, those claims can be also derived starting from the existence of the *SVD* of A . That is what we're going to do now.

Notice that we can write A^T as:

$$A^T = ((U \Sigma) V^T)^T = V(U \Sigma)^T = V \Sigma^T U^T = V \Sigma U^T \quad (2)$$

The last equality stands because Σ is diagonal and then symmetric.

$$AA^T = U \Sigma V^T V \Sigma U^T = U \Sigma^2 U^T \quad (3)$$

$$A^T A = V \Sigma U^T U \Sigma V^T = V \Sigma^2 V^T \quad (4)$$

Multiplying both sides of the equation (3) by U and the (4) by V , we get:

$$AA^T \cdot U = \Sigma^2 U^T \implies AA^T \mathbf{u}_i = \sigma^2 \cdot \mathbf{u}_i, \forall i \in 1..r \quad (5)$$

$$A^T A \cdot V = \Sigma^2 V^T \implies A^T A \mathbf{v}_i = \sigma^2 \cdot \mathbf{v}_i, \forall i \in 1..r \quad (6)$$

That proves the claims.

1.3 Solution to the point 1(b)

Since A is symmetric, we know that:

$$AA^T = A^T A = A^2 \quad (7)$$

Using the result of the last chapter we can say that:

- the left singular vectors are the eigenvectors of the matrix A^2 associated to the all non zero eigenvalues of A^2 .
- the right singular vectors are the eigenvectors of the matrix A^2 associated to the all non zero eigenvalues of A^2 .

That implied that left and right singular vectors are the same.

1.4 Eigenvectors of A

Let's denote \mathbf{x}_i the i -th eigenvector of A . By definition, we know that:

$$\exists \lambda_i : A \cdot \mathbf{x}_i = \lambda_i \cdot \mathbf{x}_i \quad (8)$$

Multiplying both sides by A we have:

$$A^2 \cdot \mathbf{x}_i = \lambda_i \cdot A \cdot \mathbf{x}_i \quad (9)$$

Applying the (8) again to the right side of (9) we have:

$$A^2 \cdot \mathbf{x}_i = \lambda_i^2 \cdot \mathbf{x}_i \quad (10)$$

So the eigenvectors \mathbf{x}_i are the same as singular vectors, in case of course the matrix A is symmetric.

1.5 Singular values and eigenvalues

The i -th singular value σ_i of the matrix A can be computed as the square root of the corresponding eigenvalue of AA^T , in our case also of the corresponding eigenvalue of A^2 .

Looking at the (10), we can notice that i -th eigenvalue of A^2 is λ^2 , so the corresponding (i -th) singular value is λ , that is the (i -th) eigenvalue of A .

So we can conclude that also eigenvalue and singular values of A are the same.

2 Assignment 2

To calculate $E[R_{ij}]$ we can apply the definition of expected value of a random variable.

$$E[R_{ij}] = \sum_x x \cdot P[R_{ij} = x] \quad (11)$$

where x varies over all possible values in the range of x . According to the definition of R_{ij} :

$$E[R_{ij}] = -\sqrt{\frac{s}{k}} \frac{1}{2s} + 0(1 - \frac{1}{s}) + \sqrt{\frac{s}{k}} \frac{1}{2s} = 0 \quad (12)$$

To say something about $E[\|f(u)\|^2]$, we need first to compute $E[R_{ij}^2]$, we will reuse the result later. Since:

$$R_{ij}^2 = \begin{cases} \frac{s}{k}, & \text{with prob. } \frac{1}{s}; \\ 0, & \text{with prob. } 1 - \frac{1}{s}; \end{cases} \quad (13)$$

We can say that:

$$E[R_{ij}^2] = \frac{s}{k} \frac{1}{s} + 0(1 - \frac{1}{s}) = \frac{1}{k} \quad (14)$$

$$E[\|f(u)\|^2] = E[\|R^T \mathbf{u}\|^2] = \sum_{i=1}^k E[(R(i, *) \mathbf{u})^2] \quad (15)$$

where $R(i, *)$ is the i -th row of the matrix R^T . $R(i, *)$ has dimension $(1 \times d)$, while \mathbf{u} has dimension $(d \times 1)$, so the result of their product $R(i, *) \mathbf{u}$ is a scalar value. Since the random variable R_{ij} does not depend on the row number, we can say that:

$$E[\|f(u)\|^2] = k \cdot E[(\mathbf{r}^T \cdot \mathbf{u})^2] \quad (16)$$

Where we have replaced $R(i, *)$ with \mathbf{r}^T , to indicate a generic vector of size d having as values the R_{ij} .

$$\begin{aligned}
E[(\mathbf{r}^T \cdot \mathbf{u})^2] &= E \left[\left(\sum_{i=1}^d r_i u_i \right)^2 \right] \\
&= E \left[\sum_{i=1}^d (r_i u_i)^2 + \sum_{l=1}^d \sum_{m=1}^d 2r_l r_m u_l u_m \right] \\
&= \sum_{i=1}^d E[r_i^2] u_i^2 + \sum_{l=1}^d \sum_{m=1}^d 2E[r_l] E[r_m] u_l u_m \\
&= \frac{1}{k} \cdot \|u\|^2
\end{aligned} \tag{17}$$

In which we applied that $E[r_i^2] = \frac{1}{k}$, $\sum_{i=1}^d u_i^2 = \|u\|^2$ and $E[r_i] = 0$.
Finally, we can use the result and compute:

$$E[\|f(u)\|^2] = k \frac{1}{k} \|u\|^2 = \|u\|^2 \tag{18}$$

3 Assignment 3

After we've processed n pairs of the stream $[(a_1, w_1) \dots (a_n, w_n)]$, we want that the probability of any pair (a_i, w_i) , with $i \leq n$, is in the sample is exactly equal to $\frac{w_i}{W_n}$, where $W_n = \sum_{j=1}^n w_j$.

3.1 Define the algorithm

The algorithm can be inductively defined. The base case is when the first pair of the stream (a_1, w_1) arrives, in this case we initialize the variables:

$$sample = a_1; W = w_1 \tag{19}$$

The other case is when we already seen n pairs of the stream, with n greater than 0, and the new pair (a_{n+1}, w_{n+1}) pair arrives. The central idea is to choose as sample the incoming (a_{n+1}, w_{n+1}) pair with probability exactly equals to $\frac{w_{n+1}}{W + w_{n+1}}$.

If we choose it, we update the variable *sample* to a_{n+1} . Otherwise, we keep the *sample* as it is. Either way we update the variable W to $W + w_{n+1}$.

3.2 Prove that the algorithm is correct

We can prove by induction on n .

In the base case, $n=1$, the probability that a_1 is the sample variable is 1. Then it is equal to w_1/W_1 , so we achieves the desired probabilistic guarantees stated above for the only element corresponding to $i = 1$.

Now let's consider the inductive step in which we process the $(n+1)$ -st pair of the stream. The probability that the $(n+1)$ -st pair is selected is equal to:

$$\frac{w_{n+1}}{W_n + w_{n+1}} = \frac{w_{n+1}}{W_{n+1}} \quad (20)$$

Then the algorithm has the desired probabilistic guarantees also in this case.

About the other pairs in the first n positions, applying the induction hypothesis, we know that they all have probability equals to $\frac{w_i}{W_n}$ to be in the sample just before the $(n+1)$ -st pair is processed.

If the $(n+1)$ -st pair is chosen, the probability that the i -th pair, with $1 \leq i \leq n$, is in the sample is simply zero. In this setting ($s = 1$).

While the probability of having the i -th pair at the end of the step if the $(n+1)$ -st pair is **not** chosen is equals to $\frac{w_i}{W_n}$.

To compute the probability of having the i -th pair in the sample at the end of the step (e.g., just after the $(n+1)$ -st pair is processed) we can use the law of total probability conditioning the result on the fact that $(n+1)$ -st pair may be chosen or not.

Putting all together we have:

$$\begin{aligned} & \left(1 - \frac{w_{n+1}}{W_n + w_{n+1}}\right) \left(\frac{w_i}{W_n}\right) + \left(\frac{w_{n+1}}{W_n + w_{n+1}}\right) \cdot 0 \\ &= \frac{W_n + w_{n+1} - w_{n+1}}{W_n + w_{n+1}} \left(\frac{w_i}{W_n}\right) = \frac{w_i}{W_n + w_{n+1}} = \frac{w_i}{W_{n+1}} \end{aligned} \quad (21)$$

That should complete the proof.