# CS214-system programming

Section 03/08 recitation 04

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### Content

- File operation
- Directory

### File types in Linux system

- A file is a sequence of bytes
- All I/O devices (network, disks, ···) are modeled as files to generalize all input and output operations as uniform ones
- Regular file: .txt, .out, .o, .c, ...
- Directory: a file composing a set of file links
- Another directory can be an element in the set recursively

#### "Everything is a File" and Types of Files in Linux

| Normal           | Ī |   | ${\bf I}_{-}$ | Normal file   |
|------------------|---|---|---------------|---|
| Directories      | ı | d | ı             | Normal directory                                      |
| Hard link        | ı | - | ı             | additional name for existing file Linux Howto's Guide |
| Symbolic link    | ı | 1 | ı             | Shortcut to a file or directory                       |
| Socket           | ı | s | 1             | Pass data between 2 process                           |
| Named pipe       | ı | P | ı             | like sockets, user can't work directly with it        |
| Character device | ı | С | ı             | Processes character hw communication                  |
| Block device     | ı | b | ı             | Major and minor numbers for controlling dev           |

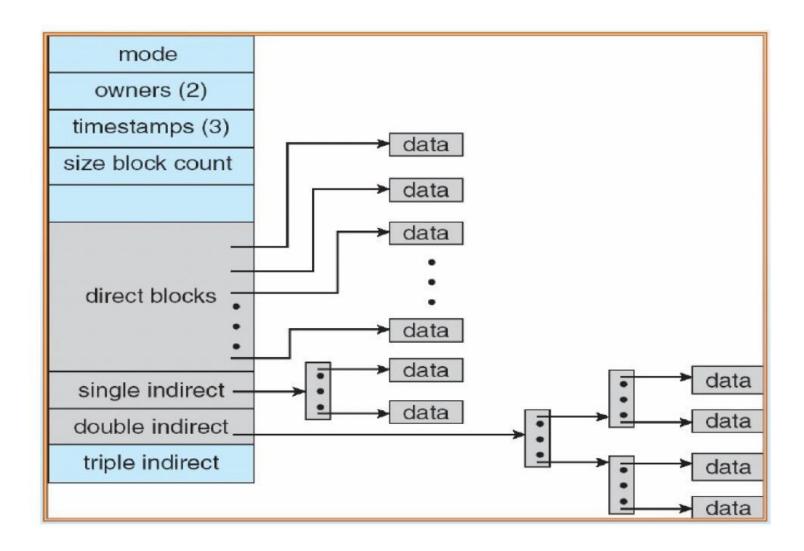
### inode

Each file on the disk has an inode associated with it.

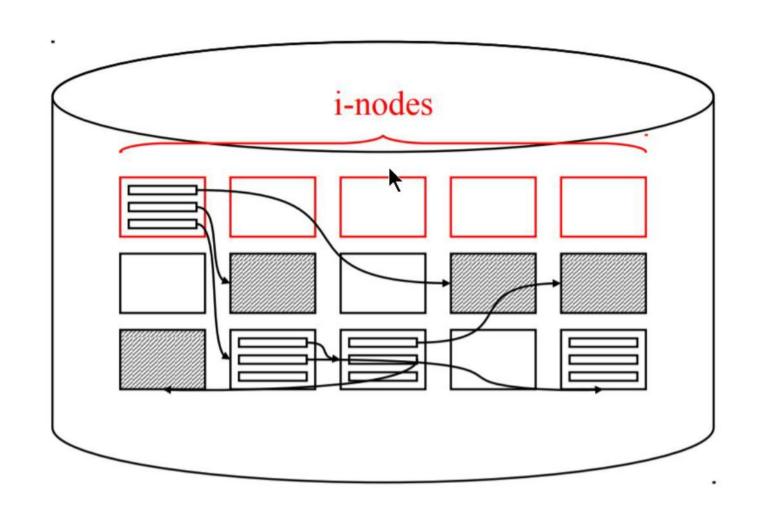
An inode is a data structure on a filesystem that stores meta-information about a file or directory.

#### Consists of:

- File metadata (mode, owner, info, size, etc.)
- Pointers to data blocks
  - Direct mapped pointers to data blocks
  - Indirect pointers
    - Point to blocks that contain pointers



# inode



### File operations in Linux

- Three basic types of operations on a file
  - Read (r)
  - Write (w)
  - Execute (x)
- A permission of an operation is associated with the relationship between current user and the file
- Three types of relationships
  - Owner (u)
  - Users in the group where the owner is in (g)
  - Users in other groups (o)

# File operations

- File descriptor:
  - Unix / Linux I/O functions: access I/O devices via system calls
  - File descriptor is used for file operations, no stream is needed
  - Open(), close(), read(), write(), ···
  - Low-level operations are performed directly using file descriptor
- FILE pointer:
  - Standard I/O functions: access I/O devices in a higher-level
  - A stream (represented with a FILE object) is needed to associate with an opened file
  - Implemented by invoking Linux / Unix I/O functions
  - printf(), scanf(), fopen(), fread(), fscanf(), fgets(), ...
  - Streams interface could provide powerful formatted input and output functions

### An opened file in Linux

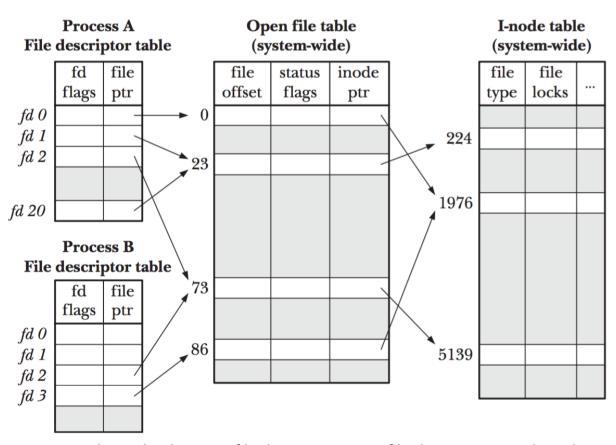


Figure 5-2: Relationship between file descriptors, open file descriptions, and i-nodes

### Open

- Open a file: kernel does the work for the program
- After opening, a file descriptor is returned to the program

- int open(char \*path, int flags, mode\_t mode);
- Returns: new file descriptor if OK, -1 otherwise
- Import library: fcntl.h

#### Close Function and Errno

- Close an opened file represented by a file descriptor
- Prototype:

```
int close(int fd);
```

- Import library: unistd.h
- Returns: 0 on success, -1 on error, errno is set propriately
- Errno: indicate what goes wrong in the event of an error in the system
- Errno is set by system calls and some library functions

#### Read Function

- Read: copies ≤n bytes from the current position of fd to memory location buf
- Import library: unistd.h
- Prototype:

```
ssize_t read(int fd, void *buf, size_t n);
```

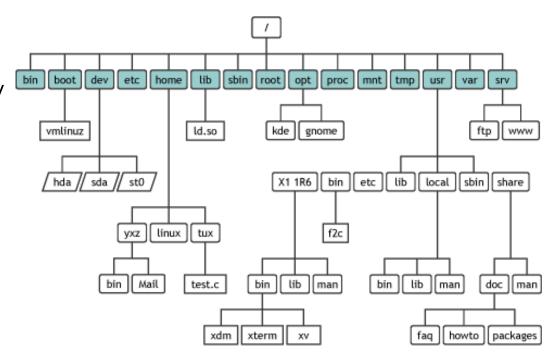
- Return number of bytes read if OK, 0 on EOF, -1 on error
- ssize\_t: signed long in essence, which could be a negative number (-1 on error)
- size\_t: unsigned long in essence, which should satisfy that is bigger or equal to 0

#### Write Function

- Write: copies ≤n bytes from the memory location buf to the current file position of fd
- Import library: unistd.h
- Prototype of writesssize\_t write(int fd, const void \*buf, size\_t n);
- Returns number of bytes written if OK, -1 on error

### Directory

- Directory: special file that contains directory entries
- In Linux, directories are organized in form of a tree and there exists a root directory
- 3 permissions associated with a directory
- Read: get the list of file links
- Write: add/remove a file link
- Execute: enter into the directory (cd command), access files inside it and view the metadata of the files (ls –l command)



http://researchhubs.com/post/computing/linux-cmd/linux-directory.html

# Directory

- It is a special file that contains directory entries
- How to find is a file is a regular file or a directory, and other information?
  - Use stat

```
int stat(const char *pathname, struct stat *statbuf);
```

#### stat

int stat(const char \*pathname, struct stat \*statbuf);

```
The stat structure
   All of these system calls return a stat structure, which contains the following fields:
       struct stat {
                                     /* ID of device containing file */
           dev_t
                     st_dev;
                                     /* Inode number */
                     st_ino;
                     st_mode;
                                     /* File type and mode */
                                     /* Number of hard links */
           nlink_t st_nlink;
                                     /* User ID of owner */
                     st_uid;
                                     /* Group ID of owner */
           gid_t
                     st_gid;
                                     /* Device ID (if special file) */
                     st_rdev;
                                     /* Total size, in bytes */
                     st_size;
                                     /* Block size for filesystem I/O */
           blksize_t st_blksize;
           blkcnt_t st_blocks;
                                     /* Number of 512B blocks allocated */
           /* Since Linux 2.6, the kernel supports nanosecond
              precision for the following timestamp fields.
              For the details before Linux 2.6, see NOTES. */
           struct timespec st_atim; /* Time of last access */
           struct timespec st_mtim; /* Time of last modification */
           struct timespec st_ctim; /* Time of last status change */
       #define st_atime st_atim.tv_sec
                                           /* Backward compatibility */
       #define st_mtime st_mtim.tv_sec
       #define st_ctime st_ctim.tv_sec
```

#### stat

```
The following mask values are defined for the file type:
              0170000 bit mask for the file type bit field
   S_IFMT
   S_IFSOCK
              0140000
                        socket
   S IFLNK
              0120000
                        symbolic link
   S_IFREG
              0100000
                        regular file
   s IFBLK
              0060000
                        block device
   S_IFDIR
              0040000
                        directory
   S_IFCHR
              0020000
                        character device
   s IFIF0
              0010000 FIF0
Thus, to test for a regular file (for example), one could write:
   stat(pathname, &sb);
if ((sb.st_mode & S_IFMT) == S_IFREG) {
       /* Handle regular file */
```

```
Because tests of the above form are common, additional macros are defined by POSIX to allow the test of the file type in <a href="mailto:st_mode">st_mode</a> to be written more concisely:

S_ISREG(m) is it a regular file?

S_ISDIR(m) directory?

S_ISCHR(m) character device?

S_ISBLK(m) block device?

S_ISFIFO(m) FIFO (named pipe)?

S_ISLNK(m) symbolic link? (Not in POSIX. 1-1996.)

S_ISSOCK(m) socket? (Not in POSIX. 1-1996.)

The preceding code snippet could thus be rewritten as:

stat(pathname, &sb);
if (S_ISREG(sb.st_mode)) {
    /* Handle regular file */
}
```

# **Directory Handling**

- DIR \*opendir(char \*path);
  - Returns a pointer to an abstract DIR structure
  - Returns NULL and sets errno on failure
- int closedir(DIR \*directory\_pointer);
  - Cloose the directory & deallocate struct
  - Return -1 and sets errno on failure
- struct dirent \*readdir(DIR \*directory\_pointer);
  - Obtain the next entry from the directory
  - Return NULL if there are no further entries

# **Directory Handling**

```
SYNOPSIS

#include <sys/types.h>
#include <dirent.h>

DIR *opendir(const char *name);

struct dirent *readdir(DIR *dirp);
```

```
DT_BLK This is a block device.

DT_CHR This is a character device.

DT_DIR This is a directory.

DT_FIFO This is a named pipe (FIFO).

DT_LNK This is a symbolic link.

DT_REG This is a regular file.

DT_SOCK This is a UNIX domain socket.

DT_UNKNOWN The file type is unknown.
```

```
1 #include (stdlib.h)
2 #include (unistd.h)
3 #include (stdio.h)
4 #include (dirent.h)
5 #include (string.h)
6
7 int main()
8 {
9    DIR *dirp = opendir(".");
10
11    struct dirent *dp;
12    while ((dp = readdir(dirp))) {
        printf("Got file %llu: %s\n", (unsigned long long) dp->d_ino, dp->d_name);
15
16
17    closedir(dirp);
18
19    return EXIT_SUCCESS;
20 }
21
```

### **Thanks**

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