

CS211 Computer Architecture Fall 2020

Recitation 12

Today • HW6

- Simulate a L1 write-through cache given input parameters that determine design and replacement policies
- Your program will read trace files
- Given the layout of the cache, you will read each line in the file and determine whether there is a cache hit or miss

HW6 – Cache Simulator

- Your file should be called first.c
- · How the program will be run as

./first.c <cache size> <assoc> <replacement policy> <block size> <trace file> *

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* Always check if you have enough arguments and return error if not (this case argc == 6)

```
./first.c <cache size> <assoc> <replacement policy> <block size> <trace file>
```

- Cache size = number of bytes in cache
 - Check to see if it is a power of 2; return error if not

```
./first.c <cache size> <assoc> <replacement policy> <block size> <trace file>
```

- Associativity layout of cache
 - options are:
 - direct direct mapped
 - assoc fully associative
 - assoc:n n-way associative
 - else, return error

```
./first.c <cache size> <assoc> <replacement policy> <block size> <trace file>
```

- Given associativity, you should be able to calculate
 - Number of sets
 - Associativity (how many cache lines per set)

Associativity	Number of Sets	lines per set
Direct mapped	cache size/block size	1
Fully associative	1	cache size/block size
N-way	cache size/(n*block size)	n



```
./first.c <cache size> <assoc> <replacement policy> <block size> <trace file>
```

- Replacement policy when a set is full and we want to add a new line, which cache line to we replace?
 - Options are
 - 1ru least recently used
 - fifo first in, first out
 - else, return error



```
./first.c <cache size> <assoc> <replacement policy> <block size> <trace file>
```

- Replacement policy when a set is full and we want to add a new line, which cache line to we replace?
 - Least recently used remove the least recently accessed line
 - Must keep track of which lines (determined by tag and set) were accessed last
 - On hit or insertion, that item is now the most recently accessed line
 - First in, first out remove the first line added to the set
 - Can use "queue" like structure or just maintain order

```
./first.c <cache size> <assoc> <replacement policy> <block size> <trace file>
```

- Block size = number of bytes in cache line/block
 - Check to see if it is a power of 2; return error if not



```
    Trace file – what you will be scarning
    Format:
```

- - Ignore address 1 it is the program counter
 - R or W determines if access is read or write
 - Address 2 the address you will be using to determine the set index and tag
- The last line of the file is #eof to help you determine the end of the file
- These files are relatively large so runtime won't be instantaneous (but shouldn't be slow either)



./first.c <cache size> <assoc> <replacement policy> <block size> <trace file>

- Trace file what you will be scanning
 - Address 2
 - NOTE: the addresses are written in HEX, and are 48 bits long
 - · Given the address and its size
 - use the correct datatypes to account for 48 bits
 - use the other parameters to determine the number of bits for the index, tag, and
 offset (though we don't really care about the offset)
 - use bit operators to extract the information

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^{*} Do research about C functions that allow you to do conversions (just do your research in general)

HW6 – Cache Simulator

```
./first.c <cache size> <assoc> <replacement policy> <block size> <trace file>
```

- Trace file what you will be scanning
 - Example:

```
./first.c 128 assoc:4 lru 8 file1.txt
```

```
0x804ae19: R 0x8caf41a057de
```

R -> you are doing a memory read

HW6 – Cache Simulator

```
./first.c <cache size> <assoc> <replacement policy> <block size> <trace file>
```

- Trace file what you will be scanning
 - Example:

```
./first.c 128 assoc:4 lru 8 file1.txt
```

0x804ae19: R 0x8caf41a057de

- For the address, we have the following sizes (as bits):
 - Offset size = log2(block size)
 - Index size = log2(number of sets)
 - Tag size = 48 offset size index size



```
./first.c <cache size> <assoc> <replacement policy> <block size> <trace file>
```

- Trace file what you will be scanning
 - Example:

```
./first.c 128 assoc:4 lru 8 file1.txt
```

--> 2 bits to label 4 sets

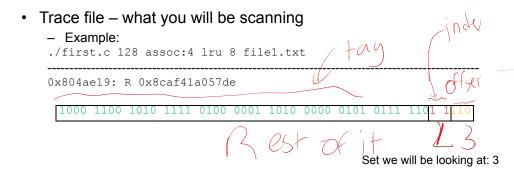
00

```
• For the address, we have the following sizes (as bits):
         - Offset size = log2(8) = 3 assoc mours
```

Index size = $\log 2(128/(4*8)) = \log 2(4) = 2$ we only have 4 sets of 4 lines each » But we need 2 bits to represent 4 sets Tag size = 48 - 3 - 2 = 43



./first.c <cache size> <assoc> <replacement policy> <block size> <trace file>





```
./first.c <cache size> <assoc> <replacement policy> <block size> <trace file>
```

Your job throughout this program is to count and output

```
Memory reads: %d\n
Memory writes: %d\n
Cache hits: %d\n
Cache misses: %d\n
```

 Determine what data structure you will use to design your cache – the replacement policy will really drive your design

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- Hints when do you increment memory reads, memory writes, cache hits, and cache misses?
 - Don't overthink it understand what it means to read from memory
 - Regardless of R or W (w)

 - Check to see if the info is in the cache
 If it is, then you have a cache hit
 - If not, then you have a mins, and you must bring that address into the cache by reading from memory
- Use appropriate data types for the values (they should all be unsigned)
 - In our example a tag of 43 bits exceeds the size of normal int values...

