



American International University-Bangladesh

Aiub_Eclipse

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1 Setup	
1.1 Sublime Build	
<pre>"shell_cmd": "g++ -std=c++17 -o\n\"\$file_base_name\" \"\$file\" &&\ntimeout 2.5s ./\"\$file_base_name\" <\ninput.txt > output.txt",\n"file_regex":\n"^(\\.\\.\\.\\[\\^:]*):([0-9]+)?(:([0-9]+)??:\n(.*)\$",\n"working_dir": "\$file_path",\n"selector": "source.c, source.c++"</pre>	
1.2 Template	
<pre>#include <bits/stdc++.h>\n#include <ext/pb_ds/assoc_container.hpp>\n#include <ext/pb_ds/tree_policy.hpp>\nusing namespace __gnu_pbds;\nusing namespace std;\ntemplate <typename T> using o_set =\ntree<T, null_type, less<T>,\nrb_tree_tag,\ntree_order_statistics_node_update>;\ntypedef long long ll;\nint32_t main() {\n ios_base::sync_with_stdio(0);\n cin.tie(0);\n}</pre>	
2 Stress Testing	
2.1 Input Gen	
<pre>mt19937_64 rnd(chrono::steady_clock::now)\n().time_since_epoch().count());\nll get_rand(ll l, ll r) {\n assert(l <= r);\n return 1 + rnd() % (r - l + 1);\n}</pre>	
2.2 Bash Script	
<pre>// run -> bash script.sh</pre>	

<pre>set -e\ng++ code.cpp -o code\ng++ gen.cpp -o gen\ng++ brute.cpp -o brute\nfor((i = 1; ; ++i)); do\n ./gen \$i > input_file\n ./code < input_file > myAnswer\n ./brute < input_file > correctAnswer\n diff -Z myAnswer correctAnswer >\n /dev/null break\n echo "Passed test: " \$i\ndone\necho "WA on the following test:"\ncat input_file\necho "Your answer is:"\ncat myAnswer\necho "Correct answer is:"\ncat correctAnswer</pre>	
3 Number Theory	
3.1 Euler Totient Function	
<pre>// Time: O(√N)\nmap<int, int> dp; // memo\nint phi(int n) {\n if (dp.count(n)) return dp[n];\n int ans = n, m = n;\n for (int i = 2; i * i <= n; i++) {\n if (m % i == 0) {\n while (m % i == 0) m /= i;\n ans = ans / i * (i - 1);\n }\n }\n if (m > 1) ans = ans / m * (m - 1);\n return dp[n] = ans;\n}</pre>	
3.2 Phi 1 to N	
<pre>void phi_1_to_n(int n) {\n vector<int> phi(n + 1);\n for (int i = 0; i <= n; i++)\n phi[i] = i;\n for (int i = 2; i <= n; i++) {\n if (phi[i] == i) {\n for (int j = i; j <= n; j += i)\n phi[j] -= phi[j] / i;\n }\n }\n}</pre>	
3.3 Segmented Sieve	
<pre>vector<char> segmentedSieve(ll L, ll R) {\n // generate all primes up to √R\n ll lim = sqrt(R);\n vector<char> mark(lim + 1, false);\n vector<ll> primes;\n for (ll i = 2; i <= lim; ++i) {\n if (!mark[i]) {\n primes.emplace_back(i);\n for (ll j = i * i; j <= lim; j +=\n i) mark[j] = true;\n }\n }\n vector<char> isPrime(R - L + 1, true);\n for (ll i : primes)</pre>	

<pre>for (ll j = max(i * i, (L + i - 1) /\ni * i); j <= R; j += i)\n isPrime[j - L] = false;\nif (L == 1) isPrime[0] = false;\nreturn isPrime;\n}</pre>	
3.4 Extended GCD	
<pre>// ax + by = gcd(a,b)\nint egcd(int a, int b, int& x, int& y) {\n if (b == 0) {\n x = 1, y = 0;\n return a;\n }\n int x1, y1;\n int d = egcd(b, a % b, x1, y1);\n x = y1;\n y = x1 - y1 * (a / b);\n return d;\n}</pre>	
3.5 Linear Diophantine Equation	
<pre>// ax + by = c, find any x and y\nbool find_any_solution(int a, int b, int\nc, int &x0, int &y0, int &g) {\n g = egcd(abs(a), abs(b), x0, y0);\n if (c % g) return false;\n x0 *= c / g;\n y0 *= c / g;\n if (a < 0) x0 = -x0;\n if (b < 0) y0 = -y0;\n return true;\n}\nvoid shift_solution(int &x, int &y,\n int a, int b, int cnt) {\n x += cnt * b;\n y -= cnt * a;\n}\nint find_all_solutions(int a, int b, int\nc, int minx, int maxx, int miny, int\nmaxy) {\n int x, y, g;\n if (!find_any_solution(a, b, c, x, y,\n g)) return 0;\n a /= g, b /= g;\n int sign_a = a > 0 ? +1 : -1;\n int sign_b = b > 0 ? +1 : -1;\n shift_solution(x, y, a, b, (minx - x)\n / b);\n if (x < minx) shift_solution(x, y, a,\n b, sign_b);\n if (x > maxx) return 0;\n int lx1 = x;\n shift_solution(x, y, a, b, (maxx - x)\n / b);\n if (x > maxx) shift_solution(x, y, a,\n b, -sign_b);\n int rx1 = x;\n shift_solution(x, y, a, b, -(miny - y)\n / a);\n if (y < miny) shift_solution(x, y, a,\n b, -sign_a);\n if (y > maxy) return 0;\n int lx2 = x;\n shift_solution(x, y, a, b, -(maxy - y)\n / a);\n}</pre>	

```

if (y > maxy) shift_solution(x, y, a,
    b, sign_a);
int rx2 = x;
if (lx2 > rx2) swap(lx2, rx2);
int lx = max(lx1, lx2);
int rx = min(rx1, rx2);
if (lx > rx) return 0;
return (rx - lx) / abs(b) + 1;
}

```

3.6 Modular Inverse using EGCD

```

// finding inverse(a) modulo m
int x, y;
int g = extended_euclidean(a, m, x, y);
if (g != 1) cout << "No solution!";
else {
    x = (x % m + m) % m;
    cout << x << endl;
}

```

3.7 Exclusion DP

```

ll f[N], g[N];
for (int i = N - 1; i >= 1; i--) {
    f[i] = nC4(div_cnt[i]);
    g[i] = f[i];
    for (int j = i + 1; j < N; j += i) {
        g[i] -= g[j];
    }
}

```

Here, $f[i]$ = how many pairs/k-tuple such that their gcd is i or it's multiple (count of pairs those are divisible by i).

$g[i]$ = how many pairs/k-tuple such that their gcd is i .

$$g[i] = f[i] - \sum_{i|j} g[j].$$

Sum of all pair gcd:

We know, how many pairs are there such that their gcd is i for every i (1 to n). So now, $\sum_{i=1}^n g[i] * i$.

Sum of all pair lcm (i = 1, j = 1):

We know, $\text{lcm}(a, b) = \frac{a*b}{\text{gcd}(a, b)}$.

Now, $f[i]$ = All pair product sum of those, whose gcd is i or it's multiple.

$g[i]$ = All pair product sum of those, whose gcd is i .
 $\text{Ans} = \sum_{i=1}^n \frac{g[i]}{i}$.

All pair product sum = $(a_1 + a_2 + \dots + a_n) * (a_1 + a_2 + \dots + a_n)$

3.8 Legendres Formula

```

//  $\frac{n!}{p^x}$  - you will get the largest x
int legendre(int n, int p) {
    int ex = 0;
    while(n) {
        ex += (n / p);
        n /= p;
    }
    return ex;
}

```

3.9 Binary Expo

```

int power(int x, long long n, int mod) {
    int ans = 1 % mod;
    while (n > 0) {
        if (n & 1) {
            ans = 1LL * ans * x % mod;
        }
        x = 1LL * x * x % mod;
        n >>= 1;
    }
    return ans;
}

```

3.10 Digit Sum of 1 to N

```

// for n=10, ans = 1+2+...+9+1+0
ll solve(ll n) {
    ll res = 0, p = 1;
    while (n / p > 0) {
        ll left = n / (p * 10);
        ll cur = (n / p) % 10;
        ll right = n % p;
        res += left * 45 * p;
        res += (cur * (cur - 1) / 2) * p;
        res += cur * (right + 1);
        p *= 10;
    }
    return res;
}

```

3.11 Pollard Rho

```

namespace PollardRho {
    mt19937 rnd(chrono::steady_clock::now().
        time_since_epoch().count());
    const int P = 1e6 + 9;
    ll seq[P];
    int primes[P], spf[P];
    inline ll add_mod(ll x, ll y, ll m) {
        return (x + y) < m ? x + y : x + y - m;
    }
    inline ll mul_mod(ll x, ll y, ll m) {
        ll res = __int128(x) * y % m;
        return res;
        // ll res = x * y - (ll)((long double)x
        // * y / m + 0.5) * m;
        // return res < 0 ? res + m : res;
    }
    inline ll pow_mod(ll x, ll n, ll m) {
        ll res = 1 % m;
        for (; n >= 1; n >>= 1) {
            if (n & 1) res = mul_mod(res, x, m);
            x = mul_mod(x, x, m);
        }
        return res;
    }
    // O(it * (logn)^3), it = number of rounds performed
    inline bool miller_rabin(ll n) {
        if (n <= 2 || (n & 1 ^ 1)) return (n == 2);
        if (n < P) return spf[n] == n;
        ll c, d, s = 0, r = n - 1;
        for (; !(r & 1); r >>= 1, s++) {}
        // each iteration is a round
        for (int i = 0; primes[i] < n &&
            primes[i] < 32; i++) {
            c = pow_mod(primes[i], r, n);

```

```

        for (int j = 0; j < s; j++) {
            d = mul_mod(c, c, n);
            if (d == 1 && c != 1 && c != (n - 1)) return false;
            c = d;
        }
        if (c != 1) return false;
        return true;
    }
    void init() {
        int cnt = 0;
        for (int i = 2; i < P; i++) {
            if (!spf[i]) primes[cnt++] = spf[i] = i;
            for (int j = 0, k; (k = i * primes[j]) < P; j++) {
                spf[k] = primes[j];
                if (spf[i] == spf[k]) break;
            }
        }
    }
    // returns O(n^(1/4))
    ll pollard_rho(ll n) {
        while (1) {
            ll x = rnd() % n, y = x, c = rnd() % n, u = 1, v, t = 0;
            ll *px = seq, *py = seq;
            while (1) {
                *py++ = y = add_mod(mul_mod(y, y, n), c, n);
                *py++ = y = add_mod(mul_mod(y, y, n), c, n);
                if ((x = *px++) == y) break;
                v = u;
                u = mul_mod(u, abs(y - x), n);
                if (!u) return __gcd(v, n);
                if (++t == 32) {
                    t = 0;
                    if ((u = __gcd(u, n)) > 1 && u < n) return u;
                }
            }
            if (t && (u = __gcd(u, n)) > 1 && u < n) return u;
        }
    }
    vector<ll> factorize(ll n) {
        if (n == 1) return vector<ll>();
        if (miller_rabin(n)) return vector<ll>{n};
        vector<ll> v, w;
        while (n > 1 && n < P) {
            v.push_back(spf[n]);
            n /= spf[n];
        }
        if (n >= P) {
            ll x = pollard_rho(n);
            v = factorize(x);
            w = factorize(n / x);
            v.insert(v.end(), w.begin(), w.end());
        }
        return v;
    }
}

```

```

}
3.12 [Problem] How Many Bases - UVa
// Given a number  $N^M$ , find out the number of integer bases in which it has exactly T trailing zeroes.
int solve_greater_or_equal(vector<int> e, int t) {
    int ans = 1;
    for (auto i : e) {
        ans = 1LL * ans * (i / t + 1) % mod;
    }
    return ans;
}
// e contains  $e_1, e_2 \rightarrow p_1^{e_1}, p_2^{e_2}$ 
int solve_equal(vector<int> e, int t) {
    return (solve_greater_or_equal(e, t) - solve_greater_or_equal(e, t + 1) + mod) % mod;
}

```

3.13 [Problem] Power Tower - CF

```

// A sequence  $w_1, w_2, \dots, w_n$  and Q queries, l and r will be given.
Calculate  $w_l^{(w_{l+1}^{(w_{l+2}^{(w_r)}))})}$ 
//  $n^x \bmod m = n^{\phi(m) + x \bmod \phi(m)} \bmod m$ 
inline int MOD(int x, int m) {
    if (x < m) return x;
    return x % m + m;
}
int power(int n, int k, int mod) {
    int ans = MOD(1, mod);
    while (k) {
        if (k & 1) ans = MOD(ans * n, mod);
        n = MOD(n * n, mod);
        k >>= 1;
    }
    return ans;
}
int f(int l, int r, int m) {
    if (l == r) return MOD(a[l], m);
    if (m == 1) return 1;
    return power(a[l], f(l + 1, r, phi(m)), m);
}

```

3.14 Formula and Properties

- $\phi(n) = n \cdot \frac{p_1-1}{p_1} \cdot \frac{p_2-1}{p_2} \dots$
- $\phi(p^e) = p^e - \frac{p^e}{p} = p^e \cdot \frac{p-1}{p}$
- For $n > 2$, $\phi(n)$ is always even.
- $\sum_{d|n} \phi(d) = n$
- NOD: $(e_1 + 1) \cdot (e_2 + 1) \dots$
- SOD: $\frac{p_1^{e_1+1}-1}{p_1-1} \cdot \frac{p_2^{e_2+1}-1}{p_2-1} \dots$
- $\log(a \cdot b) = \log(a) + \log(b)$
- $\log(a^x) = x \cdot \log(a)$
- $\log_a(x) = \frac{\log_b(x)}{\log_b(a)}$
- Digit Count of n: $\lfloor \log_{10}(n) \rfloor + 1$

- **Arithmetic Progression Sum:** $\frac{n}{2} \cdot (a + p), \quad \frac{n}{2} \cdot (2a + (n - 1)d)$
- **Geometric Sum:** $S_n = a \cdot \frac{r^n - 1}{r - 1}$
- $(1^2 + 2^2 + \dots + n^2) = \frac{n(n+1)(2n+1)}{6}$
- $(1^3 + 2^3 + \dots + n^3) = \frac{n^2(n+1)^2}{4}$
- $(2^2 + 4^2 + \dots + (2n)^2) = \frac{2n(n+1)(2n+1)}{3}$
- $(1^2 + 3^2 + \dots + (2n-1)^2) = \frac{n(2n-1)(2n+1)}{3}$
- $(2^3 + 4^3 + \dots + (2n)^3) = 2n^2(n+1)^2$
- $(1^3 + 3^3 + \dots + (2n-1)^3) = n^2(2n^2 - 1)$
- For any number n and bases $> \sqrt{n}$, there will be no representation where the number contains 0 at its second least significant digit. So it is enough to check for bases $\leq \sqrt{n}$.
- For some x and y , let's try to find all m such that $x \bmod m \equiv y \bmod m$. We can rearrange the equation into $(x - y) \equiv 0 \pmod{m}$. Thus, if m is a factor of $|x - y|$, then x and y will be equal modulo m .

4 Combinatorics and Probability

4.1 Combinations

```
// Prime Mod in O(n)
void prec() {
    fact[0] = 1;
    for (int i = 1; i < N; i++) {
        fact[i] = 1ll * fact[i - 1] * i % mod;
    }
    ifact[N - 1] = inverse(fact[N - 1]);
    for (int i = N - 2; i >= 0; i--) {
        ifact[i] = 1ll * ifact[i + 1] * (i + 1) % mod;
    }
}
int nCr(int n, int r) {
    if (r > n) return 0;
    return 1ll * fact[n] * ifact[r] % mod
        * ifact[n - r] % mod;
}
int nPr(int n, int r) {
    if (r > n) return 0;
    return 1ll * fact[n] * ifact[n - r] % mod;
}
```

4.2 nCr for any mod

```
// Time: O(n^2)
// nCr = (n-1)C(r-1) + (n-1)C(r)
for (int i = 0; i < N; i++) {
    C[i][i] = 1;
    for (int j = 0; j < i; j++) {
        C[i][j] = (C[i - 1][j] + C[i - 1][j + 1]) % mod;
    }
}
```

4.3 nCk without mod in O(r)

```
ll nCk(ll n, ll k) {
    double res = 1;
    for (ll i = 1; i <= k; ++i)
        res = res * (n - k + i) / i;
    return (ll)(res + 0.01);
}
```

4.4 Lucas Theorem

```
// returns nCr modulo mod where mod is a
// prime
// Complexity: ?
ll Lucas(ll n, ll r) {
    if (r < 0 || r > n) return 0;
    if (r == 0 || r == n) return 1;
    if (n >= MOD) {
        return (Lucas(n / MOD, r / MOD) %
                MOD * Lucas(n % MOD, r % MOD) %
                MOD) % MOD;
    }
    return (((fact[n] * invFact[r]) % MOD)
            * invFact[n - r]) % MOD;
}
```

4.5 Catalan Number

```
const int MOD = 1e9 + 7, int MAX = 1e7;
int catalan[MAX];
void init(ll n) {
    catalan[0] = catalan[1] = 1;
    for (ll i = 2; i <= n; i++) {
        catalan[i] = 0;
        for (ll j = 0; j < i; j++) {
            catalan[i] += (catalan[j] *
                           catalan[i - j - 1]) % MOD;
            if (catalan[i] >= MOD) {
                catalan[i] -= MOD;
            }
        }
    }
}
```

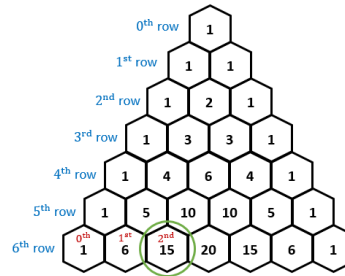
4.6 Derangement

```
// number of combinations such that
// a_i != i of a permutation a
const int N = 1e6 + 100, int p = 1e9 + 7;
ll der[N];
void countDer() {
    der[1] = 0;
    der[2] = 1;
    for (ll i = 3; i <= N; ++i) {
        der[i] = (i - 1) % p * (der[i - 1] %
                                p + der[i - 2] % p);
        der[i] %= p;
    }
}
```

4.7 Stars and Bars Theorem

- Find the number of k -tuples of non-negative integers whose sum is n . $\binom{n+k-1}{n}$
- Find the number of k -tuples of non-negative integers whose sum is $\leq n$. $\binom{n+k}{k}$
- Combination with Repetition (choose k elements from n objects, same element can be chosen multiple times). $\binom{n+k-1}{k}$
- How many ways to go from $(0,0)$ to (n,m) . $\binom{n+m}{m}$

Pascals Triangle is equivalent to nCr:



4.8 Properties of Pascal's Triangle

- $(a + b)^n = \sum_{k=0}^n \binom{n}{k} a^k b^{n-k}$
- $(k + 1)^n = \sum_{i=0}^n k^i \cdot \binom{n}{i}$
- $\sum_{i=0}^n \binom{n}{i} = 2^n$
- $\binom{k}{n} = \frac{k}{n} \binom{k-1}{n-1}$
- $\sum_{k=0}^m \binom{n+k}{k} = \binom{n+m+1}{m}$
- $\binom{n}{0}^2 + \binom{n}{1}^2 + \dots + \binom{n}{n}^2 = \binom{2n}{n}$
- $1 \binom{n}{1} + 2 \binom{n}{2} + \dots + n \binom{n}{n} = n 2^{n-1}$

4.9 Contribution Technique

- Sum of all pair sums: $\sum_{i=1}^n \sum_{j=1}^n (a_i + a_j)$
Every element will be added $2n$ times. $\sum_{i=1}^n (2 \times n \times a_i) = 2 \times n \times \sum_{i=1}^n a_i$.
- Sum of all subarray sums — $\sum_{i=1}^n (a_i \times i \times (n - i + 1))$.
- Sum of all Subsets sums — $\sum_{i=1}^n (2^{n-1} \times a_i)$.
- Product of all pair product — $\prod_{i=1}^n (a_i^{2 \times n})$.
- XOR of subarray XORs — How many subarrays does an element have? $(i \cdot (n - i + 1))$ times. If subarray count is odd then this element can contribute in total XORs.
- Sum of max minus min over all subset — Sort the array. $Min = 2^{n-i}, Max = 2^{i-1}$. $\sum_{i=1}^n (a_i \cdot 2^{i-1} - a_i \cdot 2^{n-i})$
- Sum using bits — $\sum_{k=0}^{30} (cnt_k[1] \times 2^k)$.
- Sum of Pair XORs — XOR will 1 if two bits are different $\sum_{k=0}^{30} (cnt_k[0] \times cnt_k[1] \times 2^k)$.
- Sum of Pair ANDs — $\sum_{k=0}^{30} (cnt_k[1]^2 \times 2^k)$.
- Sum of Pair ORs — $\sum_{k=0}^{30} ((cnt_k[1]^2 + 2 \times cnt_k[1] \times cnt_k[0]) \times 2^k)$.
- Sum of Subset XORs — where $cnt0! = 0$ $\sum_{k=0}^{30} (2^{cnt_k[1] + cnt_k[0] - 1} \times 2^k)$.
- Sum of Subset ANDs — $\sum_{k=0}^{30} ((2^{cnt_k[1]} - 1) \times 2^k)$.
- Sum of Subset ORs — $\sum_{k=0}^{30} ((2^n - 2^{cnt_k[0]}) \times 2^k)$.
- Sum of subarray XORs — Convert to prefix xor, then solve for pairs.
- Sum of product of all subsequence — $\prod_{i=1}^n (a_i + 1) - 1$. Example array — $[a, b]$ the subsequences are $\{a\}, \{b\}, \{a, b\}$ so ans is $a + b + (a \cdot b)$

4.10 Probability and Expected Value

- Expected Value: $E = \frac{\text{Sum of all possible values}}{\text{Total number of outcomes}}$
- Expected Value: $E = \sum_{i=1}^n P_i \cdot i$
- Variance: $V(x) = E(x^2) - \{E(x)\}^2$
- Expected Value with DP: $E(i) = \sum P(i \rightarrow j) \times (R(i \rightarrow j) + E(j))$ Where $R()$ is Immediate Reward(cost/count)
- Linearity of Expectation:
 - $E[X + Y] = E[X] + E[Y]$
 - $E[\text{Total}] = E[I_1] + E[I_2] + \dots = \sum P(I_i = 1)$

We need to define the Indicator Random Variable (I) correctly. Some Examples below —

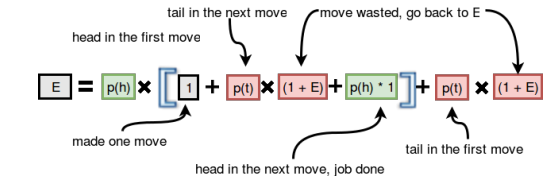
Problem-1: Find $E[\text{correct hats}]$ in a random permutation. Indicator I_i : "Does person i get their own hat?"

Problem-2: Find $E[\text{total inversions}]$ in a random permutation. Indicator I_{ij} : "Is pair (i, j) an inversion?"

Problem-3: Given a string S , delete a random index until it becomes empty. Find the expected count of palindromes seen (exclude S , include empty string). Indicator I_L : "Is the string of length L a palindrome?"

Problem-4: Pick N random integers from $[1, k]$. Merge consecutive equal values (e.g., $1, 1, 2, 2, 1, 1 \rightarrow 1, 2, 1$). Find the expected final length. Indicator I_i : "Is item i different from item $i - 1$?"

- To get two consecutive heads, what is the expected number of tosses?



- To get n heads, what is the expected number of tosses? Let's define: to get n heads, we need to toss $E(n)$ times. Now — I can get a head; I need to toss $E(n - 1)$ more times, or if I get a tail; I need to toss $E(n)$ times. So, the recurrence is: $E(n) = 0.5 \cdot (1 + E(n - 1)) + 0.5 \cdot (1 + E(n))$
- You have n bulbs, all of which are initially off. In each move, you randomly select one bulb. If the selected bulb is **off**, you toss a coin:
 - If you get head, you turn it on.
 - If you get tail, you do nothing.

If the bulb is already **on**, you skip that move (nothing happens).

Now, what is the expected number of moves required to turn all bulbs on?

The coin is not fair — the probability of getting tail is p . This problem can also be solved recursively.

Let's assume at some moment, x bulbs are already on, and the expected number of moves needed from here is $e(x)$.

The probability of picking an already on bulb is $\frac{x}{n}$. In that case, the expected number of moves is $\frac{x}{n} \times (1 + e(x))$.

The probability of picking an off bulb is $\frac{n-x}{n}$.

Now two things can happen:

- With probability p , you get tail, so you stay at the same state ($e(x)$ more moves).
- With probability $(1 - p)$, you get head, so one more bulb turns on ($e(x + 1)$ moves from there).

So, the recurrence relation is:

$$e(x) = \frac{x}{n}(1 + e(x)) + \frac{n-x}{n}(p(1 + e(x)) + (1 - p)(1 + e(x + 1)))$$

5 Data Structure

5.1 Trie for bit

```
struct Trie {
    static const int B = 31;
    struct node {
        node* nxt[2];
        int sz;
        node() {
            nxt[0] = nxt[1] = NULL;
            sz = 0;
        }
    }*root;
    Trie() {
        root = new node();
    }
    void insert(int val) {
        node* cur = root;
        cur -> sz++;
        for (int i = B - 1; i >= 0; i--) {
            int b = val >> i & 1;
            if (cur -> nxt[b] == NULL) cur ->
                nxt[b] = new node();
            cur = cur -> nxt[b];
            cur -> sz++;
        }
    }
    int query(int x, int k) { // number
        // of values s.t. val ⊕ x < k
        node* cur = root;
        int ans = 0;
        for (int i = B - 1; i >= 0; i--) {
            if (cur == NULL) break;
            int b1 = x >> i & 1, b2 = k >> i &
                1;
            if (b2 == 1) {
                if (cur -> nxt[b1]) ans += cur
                    -> sz;
                cur = cur -> nxt[!b1];
            } else cur = cur -> nxt[b1];
        }
        return ans;
    }
    int get_max(int x) { // returns
        // maximum of val ⊕ x
        node* cur = root;
        int ans = 0;
```

```
    for (int i = B - 1; i >= 0; i--) {
        int k = x >> i & 1;
        if (cur -> nxt[!k]) cur = cur ->
            nxt[!k], ans <= 1, ans++;
        else cur = cur -> nxt[k], ans <= 1;
    }
    return ans;
}
int get_min(int x) { // returns
    // minimum of val ⊕ x
    node* cur = root;
    int ans = 0;
    for (int i = B - 1; i >= 0; i--) {
        int k = x >> i & 1;
        if (cur -> nxt[k]) cur = cur ->
            nxt[k], ans <= 1;
        else cur = cur -> nxt[!k], ans <=
            1, ans++;
    }
    return ans;
}
void del(node* cur) {
    for (int i = 0; i < 2; i++) if (cur
        -> nxt[i]) del(cur -> nxt[i]);
    delete(cur);
}
} t;
```

6 Dynamic Programming

6.1 Knapsack 2

Constraints: $N \leq 100$, $W \leq 1e9$, $val[i] \leq 1000$

$dp[i][cur_val]$ = min weight needed to achieve cur_val from i to n . if $dp[1][val] \leq W$ then val is a candidate ans. where val is all_possible_val (1 to 100×1000)

6.2 LIS using Segment Tree

```
int32_t main() {
    int n; cin >> n;
    for (int i = 1; i <= n; i++) {
        cin >> a[i]; // a[i] must be >= 2
    }
    // dp[i] = LIS ending at pos i
    st.build(1, 1, M); // range max query,
        // and upd idx with max(cur_val,
        // new_val)
    for (int i = 1; i <= n; i++) {
        dp[i] = 1;
        if (a[i] != 1) {
            int mx = st.query(1, 1, M, 1, a[i]
                - 1);
            mx++;
            dp[i] = max(dp[i], mx);
        }
        st.upd(1, 1, M, a[i], dp[i]);
    }
    int ans = 0;
    for (int i = 1; i <= n; i++) ans =
        max(ans, dp[i]);
    cout << ans << '\n';
}
```

6.3 Digit DP

```
// Problem: How Many Zeroes? - LightOJ
// How many zeroes between n to m
// (n <= m). 100 to 102 ans 4
ll fun2(int i, bool is_small) {
    if (i == sz) return 1;
    int l = 0, r = s[i] - '0';
    if (is_small) r = 9;
    ll &ans = dp2[i][is_small];
    if (ans != -1) return ans;
    ans = 0;
    for (int x = 1; x <= r; x++) {
        ans += fun2(i + 1, (is_small | (x <
            r)));
    }
    return ans;
}
ll fun(int i, bool is_small, bool
    has_started) {
    if (i == sz) return 0;
    int l = 0, r = s[i] - '0';
    if (is_small) r = 9;
    ll &ans = dp1[i][is_small][has_started];
    if (ans != -1) return ans;
    ans = 0;
    for (int x = 1; x <= r; x++) {
        bool new_has_started = has_started |
            (x != 0);
        ans += fun(i + 1, (is_small | (x <
            r)), new_has_started);
    }
    if (x == 0 and has_started) {
        ans += fun2(i + 1, (is_small | (x
            < r)));
    }
    return ans;
}
void get(long long x) {
    if (x < 0) return; s = "";
    while (x > 0) {
        char c = (x % 10) + '0';
        s += c; x /= 10;
    }
    reverse(s.begin(), s.end());
    sz = s.size();
    memset(dp1, -1, sizeof(dp1));
    memset(dp2, -1, sizeof(dp2));
}
void solve() {
    ll n, m; cin >> n >> m;
    get(n - 1);
    memset(dp1, -1, sizeof(dp1));
    memset(dp2, -1, sizeof(dp2));
    ll ans1 = (n == 0) ? 0 : fun(0, false,
        false);
    get(m);
    ll ans2 = fun(0, false, false);
    cout << ans2 - ans1 + (n == 0) << '\n';
}
```

6.4 Bitmask DP

// Problem: DNA Sequence - LOJ

```
// Given n strings, find the shortest
// string that contains all the given
// strings as substrings and
// lexicographically smallest.
const int N = 17;
int n, tail[N][N], dp[(1 << N) + 2][N +
    2];
string s[N + 2];
bool cmp(string a, string b) {
    if (a.size() < b.size()) {
        return true;
    }
    return false;
}
int fun(int mask, int last) {
    if (__builtin_popcount(mask) >= n)
        return 0;
    int &ans = dp[mask][last];
    if (ans != -1) return ans;
    ans = 1e9;
    for (int j = 1; j <= n; j++) {
        if (!(mask & (1 << j))) {
            if (last == n + 1) {
                int x = s[j].length() + fun(mask
                    | (1 << j), j);
                ans = min(ans, x);
            }
            else {
                int x = (s[j].length() -
                    tail[last][j]) + fun(mask |
                        (1 << j), j);
                ans = min(ans, x);
            }
        }
    }
    return ans;
}
void print(int mask, int last) {
    if (__builtin_popcount(mask) >= n)
        return;
    int ans = fun(mask, last), idx = -1;
    string str = "{";
    for (int j = 1; j <= n; j++) {
        if (!(mask & (1 << j))) {
            if (last == n + 1) {
                int x = s[j].length() + fun(mask
                    | (1 << j), j);
                if (x == ans) {
                    string d = s[j];
                    if (d <= str) str = d; idx = j;
                }
            }
            else {
                int x = (s[j].length() -
                    tail[last][j]) + fun(mask |
                        (1 << j), j);
                if (x == ans) {
                    string d =
                        s[j].substr(tail[last][j]);
                    if (d <= str) str = d; idx = j;
                }
            }
        }
    }
}
```

```

cout << s[idx].substr(tail[last][idx]);
print(mask | (1 << idx), idx);
}

void solve() {
    cin >> n;
    string str[n + 1];
    for (int i = 1; i <= n; i++) {
        cin >> str[i]; // len <= 100
    }
    // remove duplicates and strings which
    // is a subarray of others
    sort(str + 1, str + n + 1, cmp);
    vector<string> vec;
    for (int i = 1; i <= n; i++) {
        bool ok = true;
        for (int j = i + 1; j <= n; j++) {
            if (str[j].find(str[i]) !=
                string::npos) {
                ok = false;
                break;
            }
        }
        if (ok) vec.push_back(str[i]);
    }
    n = vec.size();
    int idx = 1;
    for (auto x : vec) s[idx++] = x;
    // maximum length that suffix of a[i] =
    // prefix of a[j]
    memset(tail, 0, sizeof(tail));
    for (int i = 1; i <= n; i++) {
        string x = s[i];
        for (int j = 1; j <= n; j++) {
            string y = s[j];
            int cnt = 0;
            for (int k = min(x.length(),
                            y.length()); k >= 1; k--) {
                if (x.substr(x.length() - k) ==
                    y.substr(0, k)) {
                    cnt = k;
                    break;
                }
            }
            tail[i][j] = cnt;
        }
    }
    memset(dp, -1, sizeof(dp));
    fun(0, n + 1);
    print(0, n + 1);
}

```

6.5 MCM DP

// Problem: Slimes - Atcoder DP Contest
// Given n slimes. Choose two adjacent
slimes, and combine them into a new
slime. The new slime has a size of
x+y, where x and y are the sizes of
the slimes before combining them.
Here, a cost of x+y is incurred.
Example-
// (10, 20, 30, 40) + (30, 30, 40)
// (30, 30, 40) + (60, 40)
// (60, 40) + (100) ans = 190

// Solution: Think reverse. We
are given the final sum, from i to j.
Now we will cut any point between i to
j and calculate the cost
// Time: O(n^3)
 ll fun(int i, int j) {
 if (i == j) return 0;
 ll &ans = dp[i][j];
 if (ans != -1) return ans;
 ll cur = 0;
 for (int x = i; x <= j; x++) {
 cur += a[x];
 }
 ans = inf;
 for (int x = i; x < j; x++) {
 ans = min(ans, cur + fun(i, x) +
 fun(x + 1, j));
 }
 return ans;
 }
 cout << fun(1, n) << '\n';

7 Graph Theory

7.1 Binary Lifting and LCA

```

const int N = 2e5 + 9, LOG = 20;
vector<int> g[N];
int par[N][LOG], depth[N];
void dfs(int u, int p) {
    par[u][0] = p;
    depth[u] = depth[p] + 1;
    for (int i = 1; i < LOG; i++) {
        par[u][i] = par[par[u][i - 1]][i - 1];
    }
    for (auto v : g[u]) {
        if (v != p) {
            dfs(v, u);
        }
    }
}

int lca(int u, int v) {
    if (depth[u] < depth[v]) {
        swap(u, v);
    }
    int k = depth[u] - depth[v];
    for (int i = 0; i < LOG; i++) {
        if (CHECK(k, i)) u = par[u][i];
    }
    if (u == v) return u;
    for (int i = LOG - 1; i >= 0; i--) {
        if (par[u][i] != par[v][i]) {
            u = par[u][i];
            v = par[v][i];
        }
    }
    return par[u][0];
}

int kth(int u, int k) { // kth parent of
    u
    assert(k >= 0);
    for (int i = 0; i < LOG; i++) {
        if (CHECK(k, i)) u = par[u][i];
    }
    return u;
}

```

```

int dist(int u, int v) { // distance from
    u to v
    int l = lca(u, v);
    return (depth[u] - depth[l]) +
        (depth[v] - depth[l]);
}

// kth node from u to v, 0th node is u
int kth(int u, int v, int k) {
    int l = lca(u, v);
    int d = dist(u, v);
    assert(k <= d);
    if (depth[l] + k <= depth[u]) {
        return kth(u, k);
    }
    k -= depth[u] - depth[l];
    return kth(v, depth[v] - depth[l] - k);
}

```

7.2 LCA and Sparse Table on Tree

```

// max and min weights of a path
const int N = 1e5 + 9, LOG = 20, inf =
    1e9; // change here
vector<array<int, 2>> g[N];
int par[N][LOG], tree_mx[N][LOG],
    depth[N];
void dfs(int u, int p, int dis) {
    par[u][0] = p;
    tree_mx[u][0] = dis;
    depth[u] = depth[p] + 1;
    for (int i = 1; i < LOG; i++) {
        par[u][i] = par[par[u][i - 1]][i - 1];
        tree_mx[u][i] = max(tree_mx[u][i -
            1], tree_mx[par[u][i - 1]][i -
            1]);
    }
    for (auto [v, w] : g[u]) {
        if (v != p) {
            dfs(v, u, w);
        }
    }
}

int query_max(int u, int v) { // max
    weight on path u to v
    int l = lca(u, v);
    int d = dist(l, u);
    int ans = 0;
    for (int i = 0; i < LOG; i++) {
        if (CHECK(d, i)) {
            ans = max(ans, tree_mx[u][i]);
            u = par[u][i];
        }
    }
    d = dist(l, v);
    for (int i = 0; i < LOG; i++) {
        if (CHECK(d, i)) {
            ans = max(ans, tree_mx[v][i]);
            v = par[v][i];
        }
    }
    return ans;
}

```

7.3 Dijkstra

```

vector<int> dijkstra(int s) {
    vector<int> dis(n + 1, inf);

```

```

    vector<bool> vis(n + 1, false);
    dis[s] = 0;
    priority_queue<array<int, 2>,
        vector<array<int, 2>>,
        greater<array<int, 2>>> pq;
    pq.push({0, s});
    while (!pq.empty()) {
        auto [d, u] = pq.top(); pq.pop();
        if (vis[u]) continue;
        vis[u] = true;
        for (auto [v, w] : g[u]) {
            if (dis[v] > d + w) {
                dis[v] = d + w;
                pq.push({dis[v], v});
            }
        }
    }
    return dis;
}

7.4 Bellman Ford
// works for neg edge, can detect neg
// cycle
// Time: O(n^2)
const ll inf = 1e18;
vector<ll> dis(N, inf);
bool bellman_ford(int s) {
    dis[s] = 0;
    bool has_cycle = false;
    for (int i = 1; i <= n; i++) {
        for (int u = 1; u <= n; u++) {
            for (auto [v, w] : g[u]) {
                if (dis[v] > dis[u] + w) {
                    if (i == n) has_cycle = true;
                    dis[v] = dis[u] + w;
                }
            }
        }
    }
    return has_cycle;
}

```

7.5 Floyd Warshall

```

// dis[i][j] = min distance to reach i to
// j, works for neg edge (no neg cycle)
// Time: O(n^3)
vector<int> construct_path(int u, int v)
{ // O(n)
    if (nxt[u][v] == -1) return {};
    vector<int> path = { u };
    while (u != v) {
        u = nxt[u][v];
        path.push_back(u);
    }
    return path;
}

void floyd_warshall() {
    for (int i = 1; i <= n; i++) {
        for (int j = 1; j <= n; j++) {
            if (i == j) dis[i][j] = 0;
            else if (g[i][j] == 0) dis[i][j] =
                inf;
            else dis[i][j] = g[i][j];
        }
    }
}

```

```

for (int k = 1; k <= n; ++k) {
    for (int i = 1; i <= n; ++i) {
        for (int j = 1; j <= n; ++j) {
            if (dis[i][k] < inf and
                dis[k][j] < inf)
                dis[i][j] = min(dis[i][j],
                                dis[i][k] + dis[k][j]);
            nxt[i][j] = nxt[i][k];
        }
    }
}
int32_t main() {
    int q; cin >> n >> m >> q;
    memset(nxt, -1, sizeof nxt);
    while (m--) {
        int u, v, w; cin >> u >> v >> w;
        g[u][v] = (g[u][v] != 0 ?
                    min(g[u][v], w) : w);
        g[v][u] = (g[v][u] != 0 ?
                    min(g[v][u], w) : w);
        nxt[u][v] = v;
        nxt[v][u] = u;
    }
    floyd_warshall();
    while (q--) {
        int u, v; cin >> u >> v;
        cout << (dis[u][v] == inf ? -1 :
                  dis[u][v]) << '\n';
    }
    return 0;
}

```

7.6 Strongly Connected Components

```

// Time:  $O(n + m)$ 
const int N = 1e5 + 9;
vector<int> g[N], gT[N], G[N];
vector<bool> vis(N, false);
vector<vector<int>> components;
vector<int> order;
int n, roots[N], sz[N];
void dfs(int u) {
    vis[u] = true;
    for (auto v : g[u]) {
        if (!vis[v]) dfs(v);
    }
    order.push_back(u);
}
void dfs2(int u, vector<int> &component) {
    vis[u] = true;
    component.push_back(u);
    for (auto v : gT[u]) {
        if (!vis[v]) dfs2(v, component);
    }
}
void scc() {
    // get order sorted by end time
    order.clear();
    for (int u = 1; u <= n; u++) {
        if (!vis[u]) dfs(u);
    }
    reverse(order.begin(), order.end());
    // transpose the graph
    for (int u = 1; u <= n; u++) {
        for (auto v : g[u]) {

```

```

            gT[v].push_back(u);
        }
    }
    // get all components
    components.clear();
    for (int i = 1; i <= n; i++) vis[i] =
        false;
    for (auto u : order) {
        if (!vis[u]) {
            vector<int> component;
            dfs2(u, component);
            sort(component.begin(),
                 component.end());
            components.push_back(component);
            for (auto v : component) {
                roots[v] = component.front();
                sz[v] = component.size();
            }
        }
    }
    // add edges to condensation graph
    for (int u = 1; u <= n; u++) {
        for (auto v : g[u]) {
            if (roots[u] != roots[v]) {
                G[roots[u]].push_back(roots[v]);
            }
        }
    }
    // when you need to use condensed graph,
    use it carefully (Specially  $g \rightarrow G$ ,
     $i \rightarrow roots[i]$ )
}

```

7.7 Articulation Points

```

int disc[N], low[N], timer, n;
vector<bool> vis(N, false), is_ap(N,
    false);
void ap_dfs(int u, int p) {
    disc[u] = low[u] = ++timer;
    vis[u] = true;
    int children_cnt = 0;
    for (auto v : g[u]) {
        if (v == p) continue;
        if (vis[v]) low[u] = min(low[u],
                                disc[v]);
        else {
            ap_dfs(v, u);
            low[u] = min(low[u], low[v]);
            if (disc[u] <= low[v] and p != -1)
                is_ap[u] = true;
            children_cnt++;
        }
    }
    if (p == -1 and children_cnt > 1)
        is_ap[u] = true;
}
void find_articulation_points() {
    for (int u = 1; u <= n; u++) {
        if (!vis[u]) {
            timer = 0;
            ap_dfs(u, -1);
        }
    }
}

```

7.8 Find Bridges

```

map<pair<int, int>, int> bridges;
void bridges_dfs(int u, int p) { // find
    bridges
    disc[u] = low[u] = ++timer;
    vis[u] = true;
    for (auto v : g[u]) {
        if (v == p) continue;
        if (vis[v]) low[u] = min(low[u],
                                disc[v]);
        else {
            bridges_dfs(v, u);
            low[u] = min(low[u], low[v]);
            if (disc[u] < low[v]) {
                bridges[(make_pair(min(u, v),
                                    max(u, v)))]++;
            }
        }
    }
}
void find_bridges() {
    for (int u = 1; u <= n; u++) {
        if (!vis[u]) {
            timer = 0;
            bridges_dfs(u, -1);
        }
    }
}

```

7.9 DSU on Tree

```

// Problem: Distinct Colors CSES (Number
// of distinct color in a subtree)
// Complexity:  $O(n(\log n)^2)$ 
set<int> se[N];
int col[N], ans[N], par[N]; // par[i] = i
// initially
int find(int i) {
    return (i == par[i] ? i : par[i] =
        find(par[i]));
}
void merge(int u, int v) {
    if ((u = find(u)) == (v = find(v)))
        return;
    if (se[u].size() > se[v].size())
        swap(u, v);
    for (auto x : se[u]) {
        se[v].insert(x);
    }
    se[u].clear();
    par[u] = v;
}
void dfs(int u, int p) {
    se[find(u)].insert(col[u]);
    for (auto v : g[u]) {
        if (v != p) {
            dfs(v, u);
            merge(u, v);
        }
    }
    ans[u] = se[find(u)].size();
}

```

7.10 Heavy-Light Decomposition

```

// Per Query Complexity:  $O(\log n^2)$ 
// Path and subtree updates and queries.

```

```

const int N = 2e5 + 9, LOG = 20, inf =
    1e9; // change here
vector<int> g[N];
int par[N][LOG], depth[N], sz[N];
int disc[N], finish[N], timer, head[N];
int n;
void dfs(int u, int p = 0) {
    par[u][0] = p;
    depth[u] = depth[p] + 1;
    sz[u] = 1;
    for (int i = 1; i < LOG; i++) {
        par[u][i] = par[par[u][i - 1]][i - 1];
    }
    if (p) g[u].erase(find(g[u].begin(),
                           g[u].end(), p));
    for (auto &v : g[u]) {
        if (v != p) {
            dfs(v, u);
            sz[u] += sz[v];
            if (sz[v] > sz[g[u][0]]) swap(v,
                g[u][0]);
        }
    }
}
void dfs_hld(int u) {
    disc[u] = ++timer;
    for (auto v : g[u]) {
        head[v] = (v == g[u][0] ? head[u] :
            v);
        dfs_hld(v);
    }
    finish[u] = timer;
}
int kth(int u, int k) {
    assert(k >= 0);
    for (int i = 0; i < LOG; i++) {
        if (CHECK(k, i)) u = par[u][i];
    }
    return u;
}
int query_up(int u, int v) {
    int ans = -inf; // change here
    while (head[u] != head[v]) {
        ans = max(ans, st.query(1, 1, n,
            disc[head[u]], disc[u])); //
        // change here
        u = par[head[u]][0];
    }
    ans = max(ans, st.query(1, 1, n,
        disc[v], disc[u])); // change here
    return ans;
}
int query(int u, int v) {
    int lc = lca(u, v);
    int ans = query_up(u, lc);
    if (v != lc) {
        ans = max(ans, query_up(v, kth(v,
            depth[v] - depth[lc] - 1))); //
        // change here
    }
    return ans;
}
void solve() {

```

```

cin >> n;
for (int i = 2; i <= n; i++) {
    int u, v; cin >> u >> v;
    g[u].push_back(v);
    g[v].push_back(u);
}
dfs(1);
head[1] = 1;
dfs_hld(1);
st.build(1, 1, n);
}

```

7.11 Dinic

```

const int N = 105, inf = 1e9;
int n, m, st, en;
// Edges should be added in both
// direction separately if the graph is
// undirected
const int INF = 2000000000;
struct Edge {
    int from, to, cap, flow, index;
    Edge(int from, int to, int cap, int
        flow, int index) :
        from(from), to(to), cap(cap),
        flow(flow), index(index) {}
};
struct Dinic {
    int N;
    vector<vector<Edge>> > G;
    vector<Edge*> dad;
    vector<int> Q;
    Dinic(int N) : N(N), G(N), dad(N),
        Q(N) {}
    void AddEdge(int from, int to, int
        cap) {
        G[from].emplace_back(from, to, cap,
            0, G[to].size());
        if (from == to)
            G[from].back().index++;
        G[to].emplace_back(to, from, 0, 0,
            G[from].size() - 1);
    }
    long long BlockingFlow(int s, int t) {
        fill(dad.begin(), dad.end(), (Edge
            *) NULL);
        dad[s] = &G[0][0] - 1;
        int head = 0, tail = 0;
        Q[tail++] = s;
        while (head < tail) {
            int x = Q[head++];
            for (int i = 0; i < G[x].size();
                i++) {
                Edge &e = G[x][i];
                if (!dad[e.to] && e.cap - e.flow
                    > 0) {
                    dad[e.to] = &G[x][i];
                    Q[tail++] = e.to;
                }
            }
        }
        if (!dad[t]) return 0;
        long long totflow = 0;
        for (int i = 0; i < G[t].size();
            i++) {

```

```

Edge *start =
    &G[G[t][i].to][G[t][i].index];
    int amt = INF;
    for (Edge *e = start; amt && e !=
        dad[s]; e = dad[e->from]) {
        if (!e) { amt = 0; break; }
        amt = min(amt, e->cap - e->flow);
    }
    if (amt == 0) continue;
    for (Edge *e = start; amt && e !=
        dad[s]; e = dad[e->from]) {
        e->flow += amt;
        G[e->to][e->index].flow -= amt;
    }
    totflow += amt;
}
return totflow;
}
long long GetMaxFlow(int s, int t) {
    long long totflow = 0;
    while (long long flow =
        BlockingFlow(s, t))
        totflow += flow;
    return totflow;
}
};
void solve() {
    cin >> n >> st >> en >> m;
    Dinic *dinic = new Dinic(n + 1);
    while (m--) {
        int u, v, w; cin >> u >> v >> w;
        dinic->AddEdge(u, v, w);
        dinic->AddEdge(v, u, w);
    }
    cout << dinic->GetMaxFlow(st, en) <<
        '\n';
    delete(dinic);
}

```

8 String

8.1 Hashing

```

const int N = 1e6 + 9; // change here
const int MOD1 = 127657753, MOD2 =
    987654319;
const int p1 = 137, p2 = 277; // change
    here
int ip1, ip2;
pair<int, int> pw[N], ipw[N];
void prec() {
    pw[0] = {1, 1};
    for (int i = 1; i < N; i++) {
        pw[i].first = 111 * pw[i - 1].first
            * p1 % MOD1;
        pw[i].second = 111 * pw[i -
            1].second * p2 % MOD2;
    }
    ip1 = power(p1, MOD1 - 2, MOD1);
    ip2 = power(p2, MOD2 - 2, MOD2);
    ipw[0] = {1, 1};
    for (int i = 1; i < N; i++) {
        ipw[i].first = 111 * ipw[i -
            1].first * ip1 % MOD1;
        ipw[i].second = 111 * ipw[i -
            1].second * ip2 % MOD2;
    }
}

```

```

}
}
struct Hashing {
    int n;
    string s;
    vector<pair<int, int>> hash_val;
    vector<pair<int, int>> rev_hash_val;
    Hashing() {}
    Hashing(string _s) {
        s = _s;
        n = s.size();
        hash_val.emplace_back(0, 0);
        for (int i = 0; i < n; i++) {
            pair<int, int> p;
            p.first = (hash_val[i].first + 111
                * s[i] * pw[i].first % MOD1) %
                MOD1;
            p.second = (hash_val[i].second +
                111 * s[i] * pw[i].second %
                MOD2) % MOD2;
            hash_val.push_back(p);
        }
        rev_hash_val.emplace_back(0, 0);
        for (int i = 0, j = n - 1; i < n;
            i++, j--) {
            pair<int, int> p;
            p.first = (rev_hash_val[i].first +
                111 * s[i] * pw[j].first %
                MOD1) % MOD1;
            p.second = (rev_hash_val[i].second
                + 111 * s[i] * pw[j].second %
                MOD2) % MOD2;
            rev_hash_val.push_back(p);
        }
    }
    pair<int, int> get_hash(int l, int r)
    { // 1 indexed
        pair<int, int> ans;
        ans.first = (hash_val[r].first -
            hash_val[l - 1].first + MOD1) *
            111 * ipw[l - 1].first % MOD1;
        ans.second = (hash_val[r].second -
            hash_val[l - 1].second + MOD2) *
            111 * ipw[l - 1].second % MOD2;
        return ans;
    }
    pair<int, int> rev_hash(int l, int r)
    { // 1 indexed
        pair<int, int> ans;
        ans.first = (rev_hash_val[r].first -
            rev_hash_val[l - 1].first +
            MOD1) * 111 * ipw[n - r].first %
            MOD1;
        ans.second = (rev_hash_val[r].second
            - rev_hash_val[l - 1].second +
            MOD2) * 111 * ipw[n - r].second
            % MOD2;
        return ans;
    }
    pair<int, int> get_hash() { // 1
        indexed
        return get_hash(1, n);
    }
    bool is_palindrome(int l, int r) {

```

```

        return get_hash(l, r) == rev_hash(l,
            r);
    }
};

```

8.2 Hashing with Updates

```

using T = array<int, 2>;
const T MOD = {127657753, 987654319};
const T p = {137, 277}; // change here
T operator + (T a, int x) {return {(a[0]
    + x) % MOD[0], (a[1] + x) % MOD[1]};}
T operator - (T a, int x) {return {(a[0]
    - x + MOD[0]) % MOD[0], (a[1] - x +
    MOD[1]) % MOD[1]};}
T operator * (T a, int x) {return
    {(int)((long long) a[0] * x %
    MOD[0]), (int)((long long) a[1] * x
    % MOD[1])};}
T operator + (T a, T x) {return {(a[0] +
    x[0]) % MOD[0], (a[1] + x[1]) %
    MOD[1]};}
T operator - (T a, T x) {return {(a[0] -
    x[0] + MOD[0]) % MOD[0], (a[1] -
    x[1] + MOD[1]) % MOD[1]};}
T operator * (T a, T x) {return
    {(int)((long long) a[0] * x[0] %
    MOD[0]), (int)((long long) a[1] *
    x[1] % MOD[1])};}
ostream& operator << (ostream& os, T
    hash) {return os << "(" << hash[0]
    << ", " << hash[1] << ")";}
T pw[N], ipw[N];
void prec() {
    pw[0] = {1, 1};
    for (int i = 1; i < N; i++) {
        pw[i] = pw[i - 1] * p;
    }
    ipw[0] = {1, 1};
    T ip = {power(p[0], MOD[0] - 2,
        MOD[0]), power(p[1], MOD[1] - 2,
        MOD[1])};
    for (int i = 1; i < N; i++) {
        ipw[i] = ipw[i - 1] * ip;
    }
}
struct Hashing {
    int n;
    string s; // 1 - indexed
    vector<array<T, 2>> t; // (normal, rev)
        hash
    array<T, 2> merge(array<T, 2> l,
        array<T, 2> r) {
        l[0] = l[0] + r[0];
        l[1] = l[1] + r[1];
        return l;
    }
    void build(int node, int b, int e) {
        if (b == e) {
            t[node][0] = pw[b] * s[b];
            t[node][1] = pw[n - b + 1] * s[b];
            return;
        }
        int mid = (b + e) >> 1, l = node <<
            1, r = l | 1;

```

```

    build(l, b, mid);
    build(r, mid + 1, e);
    t[node] = merge(t[l], t[r]);
}

void upd(int node, int b, int e, int i, char x) {
    if (b > i || e < i) return;
    if (b == e && b == i) {
        t[node][0] = pw[b] * x;
        t[node][1] = pw[n - b + 1] * x;
        return;
    }
    int mid = (b + e) >> 1, l = node << 1, r = l | 1;
    upd(l, b, mid, i, x);
    upd(r, mid + 1, e, i, x);
    t[node] = merge(t[l], t[r]);
}

array<T, 2> query(int node, int b, int e, int i, int j) {
    if (b > j || e < i) return {T{0, 0}};
    if (b == e && b == i) return t[node];
    int mid = (b + e) >> 1, l = node << 1, r = l | 1;
    return merge(query(l, b, mid, i, j), query(r, mid + 1, e, i, j));
}

Hashing() {}
Hashing(string _s) {
    n = _s.size();
    s = "." + _s;
    t.resize(4 * n + 1);
    build(1, 1, n);
}

void upd(int i, char c) {
    upd(1, 1, n, i, c);
    s[i] = c;
}

T get_hash(int l, int r) { // 1 - indexed
    return query(1, 1, n, l, r)[0] * ipw[l - 1];
}

T rev_hash(int l, int r) { // 1 - indexed
    return query(1, 1, n, l, r)[1] * ipw[n - r];
}

T get_hash() {
    return get_hash(1, n);
}

bool is_palindrome(int l, int r) {
    return get_hash(l, r) == rev_hash(l, r);
}
};

```

8.3 Hashing with Upd and Deletes

// update or delete a char in the string or check whether a range [l,r] is a palindrome or not (Palindromic Query I - Toph)

#define int long long

const int N = 1e5 + 9;

```

int en;
struct ST {
    pair<int, int> tree[4 * (N + N)];
    void build(int n, int b, int e) {
        if (b == e) {
            tree[n].first = b;
            tree[n].second = 1;
            return;
        }
        int mid = (b + e) >> 1, l = n << 1, r = l + 1;
        build(l, b, mid);
        build(r, mid + 1, e);
        tree[n].second = tree[l].second + tree[r].second;
    }
    void upd(int n, int b, int e, int i, int x1, int x2) {
        if (b > i || e < i) return;
        if (b == e && b == i) {
            tree[n].first = x1;
            tree[n].second = x2;
            return;
        }
        int mid = (b + e) >> 1, l = n << 1, r = l + 1;
        upd(l, b, mid, i, x1, x2);
        upd(r, mid + 1, e, i, x1, x2);
        tree[n].second = tree[l].second + tree[r].second;
    }
    pair<int, int> query(int n, int b, int e, int x) {
        if (b > e) return {-1, -1};
        if (tree[n].second < x) return {tree[n].second, -1};
        if (b == e) return tree[n];
        int mid = (b + e) >> 1, l = n << 1, r = l + 1;
        pair<int, int> L = query(l, b, mid, x);
        if (L.second != -1) return L;
        pair<int, int> R = query(r, mid + 1, e, x - L.first);
        return R;
    }
} st, st2;
using T = array<int, 2>;
const T MOD = {127657753, 987654319};
const T p = {137, 277};
// add operators overloading of T (from only upd) + prec()
int get(int i, int n) {
    return n - i + 1;
}

struct Hashing {
    int n; string s;
    vector<T> tree, lazy;
    void push(int node, int b, int e) {
        if (lazy[node][0] == 1) return;
        tree[node] = tree[node] * lazy[node];
        if (b != e) {
            int l = node << 1, r = l + 1;
            lazy[l] = lazy[l] * lazy[node];

```

```

            lazy[r] = lazy[r] * lazy[node];
        }
        lazy[node] = T{1, 1};
    }
    void build(int node, int b, int e) {
        lazy[node] = T{1, 1};
        if (b == e) {
            tree[node] = pw[b] * s[b];
            return;
        }
        int mid = (b + e) >> 1, l = node << 1, r = l | 1;
        build(l, b, mid);
        build(r, mid + 1, e);
        tree[node] = tree[l] + tree[r];
    }
    void upd(int node, int b, int e, int i, T x) {
        push(node, b, e);
        if (b > i || e < i) return;
        if (b == e && b == i) {
            tree[node] = x;
            return;
        }
        int mid = (b + e) >> 1, l = node << 1, r = l + 1;
        upd(l, b, mid, i, x);
        upd(r, mid + 1, e, i, x);
        tree[node] = tree[l] + tree[r];
    }
    void del(int node, int b, int e, int i, int j) {
        push(node, b, e);
        if (b > j || e < i) return;
        if (b >= i && e <= j) {
            lazy[node] = lazy[node] * ipw[1];
            push(node, b, e);
            return;
        }
        int mid = (b + e) >> 1, l = node << 1, r = l + 1;
        del(l, b, mid, i, j);
        del(r, mid + 1, e, i, j);
        tree[node] = tree[l] + tree[r];
    }
    T query(int node, int b, int e, int i, int j) {
        push(node, b, e);
        if (b > j || e < i) return {0, 0};
        if (b >= i && e <= j) return tree[node];
        int mid = (b + e) >> 1, l = node << 1, r = l + 1;
        T L = query(l, b, mid, i, j);
        T R = query(r, mid + 1, e, i, j);
        return L + R;
    }
}

Hashing() {}
Hashing(string _s) {
    s = _s;
    n = s.size();
    s = '.' + s;
    tree.resize(4 * n + 1);
    lazy.resize(4 * n + 1);
    build(1, 1, n);
}

```

```

}

void upd(int i, char c, int cur) {
    T x = pw[i] * c;
    if (cur == 1) i = st.query(1, 1, en, i).first;
    else i = st2.query(1, 1, en, i).first;
    upd(1, 1, n, i, x);
}

void del(int i, int cur) {
    int orgi = i;
    T x = pw[i] * 011;
    if (cur == 1) i = st.query(1, 1, en, i).first;
    else i = st2.query(1, 1, en, i).first;
    upd(1, 1, n, i, x);
    del(1, 1, n, i + 1, n);
    if (cur == 1) st.upd(1, 1, en, i, i, 0);
    else st2.upd(1, 1, en, i, i, 0);
}

T get_hash(int l, int r, int cur) { // 1 - indexed
    int ll = st.query(1, 1, en, l).first;
    int rr = st2.query(1, 1, en, r).first;
    if (cur == 2) {
        ll = st2.query(1, 1, en, l).first;
        rr = st2.query(1, 1, en, r).first;
    }
    return query(1, 1, n, ll, rr) * ipw[l - 1];
}

};

int32_t main() {
    prec(); // must include
    string s; cin >> s;
    int n = s.size();
    int q; cin >> q;
    string t = s;
    reverse(t.begin(), t.end());
    Hashing hs(s), hs2(t);
    en = n + q + 5;
    st.build(1, 1, en);
    st2.build(1, 1, en);
    while (q--) {
        char c; cin >> c;
        if (c == 'C') {
            int l, r; cin >> l >> r;
            int l2 = get(l, n);
            int r2 = get(r, n);
            if (hs.get_hash(l, r, 1) == hs2.get_hash(r2, l2, 2)) cout << "Yes!\n";
            else cout << "No!\n";
        }
        else if (c == 'U') {
            int i; char x; cin >> i >> x;
            int i2 = get(i, n);
            hs.upd(i, x, 1);
            hs2.upd(i2, x, 2);
        }
        else {
            int i; cin >> i;
            int i2 = get(i, n);

```

```

        hs.del(i, 1);
        hs2.del(i2, 2);
        --n;
    }
}

```

8.4 Hashing on Tree

// Given a tree, Check whether it is symmetrical or not. Problem - CF G. Symmetree

// The value for each node is it's subtree size and position is the level (ordered). But the order of childs doesn't matter (unordered)

```

const int N = 2e5 + 9;
vector<int> g[N];
vector<array<int, 3>> hassh[N]; // hash1, hash2, node
int n, sz[N];
const int MOD1 = 1e9 + 9, MOD2 = 1e9 + 21;
const int p1 = 1e5 + 19, p2 = 1e5 + 43;
void dfs2(int u, int p, int lvl) {
    array<int, 3> my_hash;
    my_hash[0] = 1ll * sz[u] * pw[lvl].first % MOD1;
    my_hash[1] = 1ll * sz[u] * pw[lvl].second % MOD2;
    my_hash[2] = u;
    bool leaf = true;
    for (auto v : g[u]) {
        if (v != p) {
            dfs2(v, u, lvl + 1);
            leaf = false;
        }
    }
    if (!leaf) {
        int sum1 = 1, sum2 = 1;
        for (auto here : hassh[u]) {
            auto [x, y, _] = here;
            sum1 = (sum1 * x) % MOD1;
            sum2 = (sum2 * y) % MOD2;
        }
        my_hash[0] = power(my_hash[0], sum1, MOD1);
        my_hash[1] = power(my_hash[1], sum2, MOD2);
    }
    hassh[p].push_back(my_hash);
}
bool ok(int u) {
    map<pair<int, int>, int> mp;
    for (auto [x, y, who] : hassh[u]) {
        mp[{x, y}]++;
    }
    int odd = 0;
    pair<int, int> val;
    for (auto [here, cnt] : mp) {
        odd += cnt & 1;
        if (cnt & 1) val = here;
    }
    if (odd == 0) return true;
    if (odd > 1) return false;
    int node;
    for (auto [x, y, who] : hassh[u]) {

```

```

        pair<int, int> here = {x, y};
        if (here == val) node = who;
    }
    return ok(node);
}
void solve() {
    cin >> n; clr(n);
    for (int i = 2; i <= n; i++) {
        int u, v; cin >> u >> v;
        g[u].push_back(v);
        g[v].push_back(u);
    }
    dfs(1, 0); // calc. subtree size
    dfs2(1, 0, 1);
    if (ok(0)) cout << "YES\n";
    else cout << "NO\n";
}

```

8.5 Compare 2 strings Lexicographically

// Time: O(logn)

```

string s;
Hashing hs;
// return 0 if both equal
// return 1 if first substring greater
// return -1 if second substring greater
// here lcp() provides the len of longest common prefix
int compare(int i, int j, int x, int y) {
    int common_prefix = lcp(i, j, x, y);
    int len1 = j - i + 1, len2 = y - x + 1;
    if (common_prefix == len1 and len1 == len2) return 0;
    else if (common_prefix == len1) return -1;
    else if (common_prefix == len2) return 1;
    else return (s[i + common_prefix - 1] < s[x + common_prefix - 1] ? -1 : 1);
}

```

8.6 KMP

```

vector<int> build_lps(string &pat) {
    int n = pat.size();
    vector<int> lps(n, 0);
    for (int i = 1; i < n; i++) {
        int j = lps[i - 1];
        while (j > 0 and pat[i] != pat[j]) {
            j = lps[j - 1];
        }
        if (pat[i] == pat[j]) j++;
        lps[i] = j;
    }
    return lps;
}
int kmp(string &txt, string &pat) {
    string s = pat + '#' + txt;
    vector<int> lps = build_lps(s);
    int ans = 0;
    for (auto x : lps) {
        if (x == pat.size()) ans++;
    }
    return ans;
}
int kmp(string &txt, string &pat) {

```

```

vector<int> lps = build_lps(pat);
int n = txt.size(), m = pat.size();
int ans = 0;
int j = 0;
for (int i = 0; i < n; i++) {
    while (j > 0 and txt[i] != pat[j]) {
        j = lps[j - 1];
    }
    if (txt[i] == pat[j]) j++;
    if (j == m) {
        ans++;
        j = lps[j - 1];
    }
}
return ans;
}

```

8.7 KMP Automata

// like DFA. if string is "abcdeabg", aut[7]['c'] = 3. Means 7th index e 'c' bosaile LPS koto, aut[7]['g'] = 8

```

void compute_automaton(string s, vector<vector<int>>& aut) {
    s += '#';
    int n = s.size();
    vector<int> pi = build_lps(s);
    aut.assign(n, vector<int>(26));
    for (int i = 0; i < n; i++) {
        for (int c = 0; c < 26; c++) {
            if (i > 0 && 'a' + c != s[i])
                aut[i][c] = aut[pi[i - 1]][c];
            else
                aut[i][c] = i + ('a' + c == s[i]);
        }
    }
}

```

8.8 Prefix Occurance Count

// Count the number of occurrences of each prefix

```

vector<int> ans(n + 1);
for (int i = 0; i < n; i++) ans[lps[i]]++;
for (int i = n - 1; i > 0; i--)
    ans[lps[i - 1]] += ans[i];
for (int i = 0; i <= n; i++) ans[i]++;

```

8.9 Number of palindromic substring in L to R using Wavelet Tree

// Problem - Kattis palindromes

```

ll f(int x) {
    return (1ll * x * (x + 1)) / 2;
}
ll f(int l, int r) {
    if (l > r) return 0;
    return f(r) - f(l - 1);
}
bool ok(int l, int r) {
    return hash_s.is_palindrome(l, r);
}

```

```

int32_t main() {
    cin >> s;
    n = s.size();
    hash_s = Hashing(s);
    for (int i = 1; i <= n; i++) {
        int l = 0, r = min(n - i, i - 1),
            cnt = 1;

```

```

        while (l <= r) {
            int mid = (l + r) >> 1;
            if (ok(i - mid, i + mid)) {
                cnt = mid;
                l = mid + 1;
            }
            else r = mid - 1;
        }
        pi1[i] = cnt + 1;
        pi1_left[i] = pi1[i] - i;
        pi1_right[i] = i + pi1[i];
    }
    for (int i = 2; i <= n; i++) {
        if (s[i - 1] == s[i - 2]) {
            int l = 0, r = min(n - i, i - 1),
                cnt = 2;
            while (l <= r) {
                int mid = (l + r) >> 1;
                if (ok(i - 1 - mid, i + mid)) {
                    cnt = mid;
                    l = mid + 1;
                }
                else r = mid - 1;
            }
            pi2[i] = cnt + 1;
        }
        else pi2[i] = 0;

        pi2_left[i] = pi2[i] - i;
        pi2_right[i] = i + pi2[i];
    }
    // wavelet trees (odd_len_left, odd_len_right, even_len_left, even_len_right)
    t1.init(pi1_left + 1, pi1_left + n + 1, -N, N);
    t2.init(pi1_right + 1, pi1_right + n + 1, -N, N);
    t3.init(pi2_left + 1, pi2_left + n + 1, -N, N);
    t4.init(pi2_right + 1, pi2_right + n + 1, -N, N);

    int q; cin >> q;
    while (q--) {
        int l, r; cin >> l >> r;
        // define k, find cnt > k and summation whose are <= k;
        int m = (l + r) / 2;
        int k = 1 - l;
        ll ans = f(l, m);
        ans += t1.sum(l, m, k);
        int cnt = t1.GT(l, m, k);
        ans += 1ll * k * cnt;
        k = 1 + r;
        ans += -f(m + 1, r);
        ans += t2.sum(m + 1, r, k);
        cnt = t2.GT(m + 1, r, k);
        ans += 1ll * k * cnt;
        if (l + 1 <= m) { // a bit different than others
            k = -1;

```

```

    ans += f(l + 1, m);
    ans += t3.sum(l + 1, m, k);
    cnt = t3.GT(l + 1, m, k);
    ans += 1ll * k * cnt;
}
k = 1 + r;
ans += -f(m + 1, r);
ans += t4.sum(m + 1, r, k);
cnt = t4.GT(m + 1, r, k);
ans += 1ll * k * cnt;
cout << ans << '\n';
}
}

```

It is easier to explain by considering only palindromes centered at indices (so, odd length), the idea is the same anyway. For each index i , r_i will be the longest radius of a palindrome centered there (in other words, the amount of palindromes centered at index i). Directly from manacher, this takes $\mathcal{O}(n)$ to calculate.

For a query $[l, r]$, we first compute $m = \frac{l+r}{2}$. Now we want to calculate

$$\sum_{i=l}^m \min(i - l + 1, r_i) + \sum_{i=m+1}^r \min(r - i + 1, r_i)$$

$$\sum_{i=l}^m \min(i - l + 1, r_i) = \sum_{i=l}^m i + \min(1 - l, r_i - i).$$

The sum over i can be found in constant time. As for the other term, if we create some array $r'_i = r_i - i$ during the preprocessing, then the queries are asking for some over range of $\min(C, r'_i)$ where C is constant. You can solve this in $\mathcal{O}(\log n)$ per query using wavelet tree.

8.10 Trie

```

const int N = 10; // change here
const char base_char = '0'; // change here
struct TrieNode {
    int cnt;
    TrieNode * nxt[N];
    TrieNode() {
        cnt = 0;
        for (int i = 0; i < N; i++) nxt[i] = NULL;
    }
} *root;
void insert(const string &s) {
    TrieNode *cur = root;
    int n = (int)s.size();
    for (int i = 0; i < n; i++) {
        int idx = s[i] - base_char;
        if (cur -> nxt[idx] == NULL) cur ->
            nxt[idx] = new TrieNode();
        cur = cur -> nxt[idx];
        cur -> cnt++;
    }
}
void rem(TrieNode *cur, string &s, int pos) { // free :: De Alloactes Memory
    if (pos == s.size()) return;

```

```

    int idx = s[pos] - base_char;
    rem(cur -> nxt[idx], s, pos + 1);
    cur -> nxt[idx] -> cnt--;
    if (cur -> nxt[idx] -> cnt == 0) {
        free(cur -> nxt[idx]);
        cur -> nxt[idx] = NULL;
    }
}
int query(const string &s) { // "s" is a
    // prefix of some element or not
    int n = (int)s.size();
    TrieNode *cur = root;
    for (int i = 0; i < n; i++) {
        int idx = s[i] - base_char;
        if (cur -> nxt[idx] == NULL) return 0;
        cur = cur -> nxt[idx];
    }
    return cur -> cnt;
}
void del(TrieNode *cur) {
    for (int i = 0; i < N; i++) if (cur ->
        nxt[i]) del(cur -> nxt[i]);
    delete(cur);
}
int32_t main() {
    root = new TrieNode(); // init new trie
    del(root); // clear trie
}

```