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AIUB Eclipse

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Contents

```
1 Setup
   1.1 Sublime Build . . . . . . . . .
2 Stress Testing
   2.1 Input Gen . . . . . . . . . . . . . . .
   2.2 Bash Script . . . . . . . . . . . . . . . .
3 Number Theory
   3.1 Euler Totient Function . . . . .
        Phi 1 to N . . . . . . . . . . . . .
        Segmented Sieve . . . . . . . . .
        Extended GCD . . . . . . . . .
        Linear Diophantine Equation . .
        Modular Inverse using EGCD . .
        Exclusion DP . . . . . . . . . .
        Legendres Formula . . . . . . .
        Binary Expo . . . . . . . . . . . .
        Digit Sum of 1 to N . . . . . . .
   3.11 [Problem] How Many Bases - UVa
1 Setup
1.1 Sublime Build
"shell_cmd": "g++ -std=c++17 -o
    \"$file_base_name\" \"$file\" &&
    timeout 2.5s ./\"$file_base_name\" <</pre>
    input.txt > output.txt",
"file regex":
    "^(..[^:]*):([0-9]+):?([0-9]+)?:?
"working_dir": "${file_path}",
"selector": "source.c, source.c++"
2 Stress Testing
2.1 Input Gen
mt19937_64 rnd(chrono::steady_clock::now_
    ().time_since_epoch().count());
11 get_rand(11 1, 11 r) {
  assert(1 <= r);
  return 1 + rnd() \% (r - 1 + 1);
2.2 Bash Script
// run -> bash script.sh
g++ code.cpp -o code
g++ gen.cpp -o gen
g++ brute.cpp -o brute
for((i = 1; ; ++i)); do
    ./gen $i > input_file
    ./code < input_file > myAnswer
    ./brute < input_file > correctAnswer
    diff -Z myAnswer correctAnswer >
        /dev/null || break
    echo "Passed test: " $i
echo "WA on the following test:"
cat input_file
echo "Your answer is:"
cat myAnswer
echo "Correct answer is:"
cat correctAnswer
```

```
3 Number Theory
3.1 Euler Totient Function
// Time: O(\sqrt{N})
map<int, int> dp; // memo
int phi(int n) {
  if (dp.count(n)) return dp[n];
  int ans = n, m = n;
  for (int i = 2; i * i <= m; i++) {
    if (m \% i == 0) {
      while (m \% i == 0) m /= i;
      ans = ans / i * (i - 1);
  if (m > 1) ans = ans / m * (m - 1);
  return dp[n] = ans;
3.2 Phi 1 to N
void phi_1_to_n(int n) {
  vector<int> phi(n + 1);
  for (int i = 0; i <= n; i++)
    phi[i] = i;
  for (int i = 2; i <= n; i++) {
    if (phi[i] == i) {
      for (int j = i; j <= n; j += i)
phi[j] -= phi[j] / i;
 }
3.3 Segmented Sieve
vector<char> segmentedSieve(ll L, ll R) {
  // generate all primes up to \sqrt{R}
  ll lim = sqrt(R);
  vector<char> mark(lim + 1, false);
  vector<ll> primes;
  for (11 i = 2; i \le 1im; ++i) {
    if (!mark[i]) {
      primes.emplace_back(i);
      for (ll j = i * i; j <= lim; j +=
          i) mark[j] = true;
  vector<char> isPrime(R - L + 1, true);
 for (11 i : primes)
for (11 j = max(i * i, (L + i - 1) /
       i * i); j <= R; j += i)
      isPrime[j - L] = false;
  if (L == 1) isPrime[0] = false;
  return isPrime;
3.4 Extended GCD
// ax + by = \gcd(a, b)
int egcd(int a, int b, int& x, int& y) {
  if (b == 0) {
    x = 1, y = 0;
```

return a;

int $d = \operatorname{egcd}(b, a \% b, x1, y1);$

y = x1 - y1 * (a / b);

int x1, y1;

return d;

```
3.5 Linear Diophantine Equation
// ax + by = c, find any x and y
bool find_any_solution(int a, int b, int
    c, int &x0, int &y0, int &g) {
  g = egcd(abs(a), abs(b), x0, y0);
  if (c % g) return false;
  y0 *= c / g;
  if (a < 0) x0 = -x0;
  if (b < 0) y0 = -y0;
  return true;
void shift_solution(int & x, int & y,
    int a, int b, int cnt) {
  x += cnt * b:
y -= cnt * b;
}
int find_all_solutions(int a, int b, int
    c, int minx, int maxx, int miny, int
    maxy) {
  int x, y, g;
  if (!find_any_solution(a, b, c, x, y,
      g)) return 0;
  a /= g, b /= g;
  int sign_a = a > 0 ? +1 : -1;
  int sign_b = b > 0 ? +1 : -1;
  shift_solution(x, y, a, b, (minx - x)
  if (x < minx) shift_solution(x, y, a,
      b, sign_b);
  if (x > maxx) return 0;
  int lx1 = x:
  shift_solution(x, y, a, b, (maxx - x)
  if (x > maxx) shift_solution(x, y, a,
      b, -sign_b);
  int rx1 = x:
  shift_solution(x, y, a, b, -(miny - y)
  if (y < miny) shift_solution(x, y, a,
      b, -sign_a);
  if (y > maxy) return 0;
  int 1x2 = x;
  shift_solution(x, y, a, b, -(maxy - y)
  if (y > maxy) shift_solution(x, y, a,
      b, sign_a);
  int rx2 = x;
  if (1x2 > rx2) swap(1x2, rx2);
  int lx = max(lx1, lx2);
  int rx = min(rx1, rx2);
  if (lx > rx) return 0;
  return (rx - lx) / abs(b) + 1;
```

```
3.6 Modular Inverse using EGCD
// finding inverse(a) modulo m
int g = extended_euclidean(a, m, x, y);
if (g != 1) cout << "No solution!";</pre>
else {
 x = (x \% m + m) \% m;
 cout << x << endl;</pre>
```

```
3.7 Exclusion DP
11 f[N], g[N];
for (int i = N - 1; i >= 1; i --) {
  f[i] = nC4(div_cnt[i]);
  g[i] = f[i];
  for (int j = i + i; j < N; j += i) {
    g[i] -= g[j];
   Here,
f[i] = \text{how many pairs/k-tuple s.t. their gcd}
is i or it's multiple (count of pairs those are
```

divisible by i). q[i] = how many pairs/k-tuple such that their

 \gcd is i. $g[i] = f[i] - \sum_{i|j} g[j].$

Sum of all pair gcd:

We know, how many pairs are there such that their gcd is i for every i (1 to n). So now, $\sum_{i=1}^{n} g[i] * i.$

Sum of all pair lcm (i = 1, j = 1): We know, $lcm(a, b) = \frac{a * b}{gcd(a, b)}$.

Now, f[i] = All pair product sum of those,whose gcd is i or it's multiple. q[i] = All pair product sum of those, whose

Ans $=\sum_{i=1}^n \frac{g[i]}{i}$.

All pair product sum = $(a_1 + a_2 + \cdots +$ $(a_n) * (a_1 + a_2 + \cdots + a_n)$

3.8 Legendres Formula

```
\frac{n!}{n^x} - you will get the largest x
int legendre(int n, int p) {
  int ex = 0:
  while(n) {
    ex += (n / p);
    n /= p;
  return ex:
```

3.9 Binary Expo

```
int power(int x, long long n, int mod) {
 int ans = 1 % mod;
 while (n > 0) {
   if (n & 1) {
     ans = 1LL * ans * x \% mod;
   x = 1LL * x * x % mod;
   n >>= 1;
 return ans;
```

```
3.10 Digit Sum of 1 to N
// for n=10, ans = 1+2+...+9+1+0
ll solve(ll n) {
  ll res = 0, p = 1;

while (n / p > 0) {

ll left = n / (p * 10);

ll cur = (n / p) % 10;
     ll right = n % p;
     res += left * 45 * p;
     res += (cur * (cur - 1) / 2) * p;
     res += cur * (right + 1);
     p *= 10;
  } return res;
3.11 [Problem] How Many Bases - UVa // Given a number N^{M} , find out the
     number of integer bases in which it
has exactly T trailing zeroes.
int solve_greater_or_equal(vector<int>
     e, int t) {
   int ans = 1;
  for (auto i : e) {
     ans = 1LL * ans * (i / t + 1) % mod;
  return ans;
mod) % mod;
```