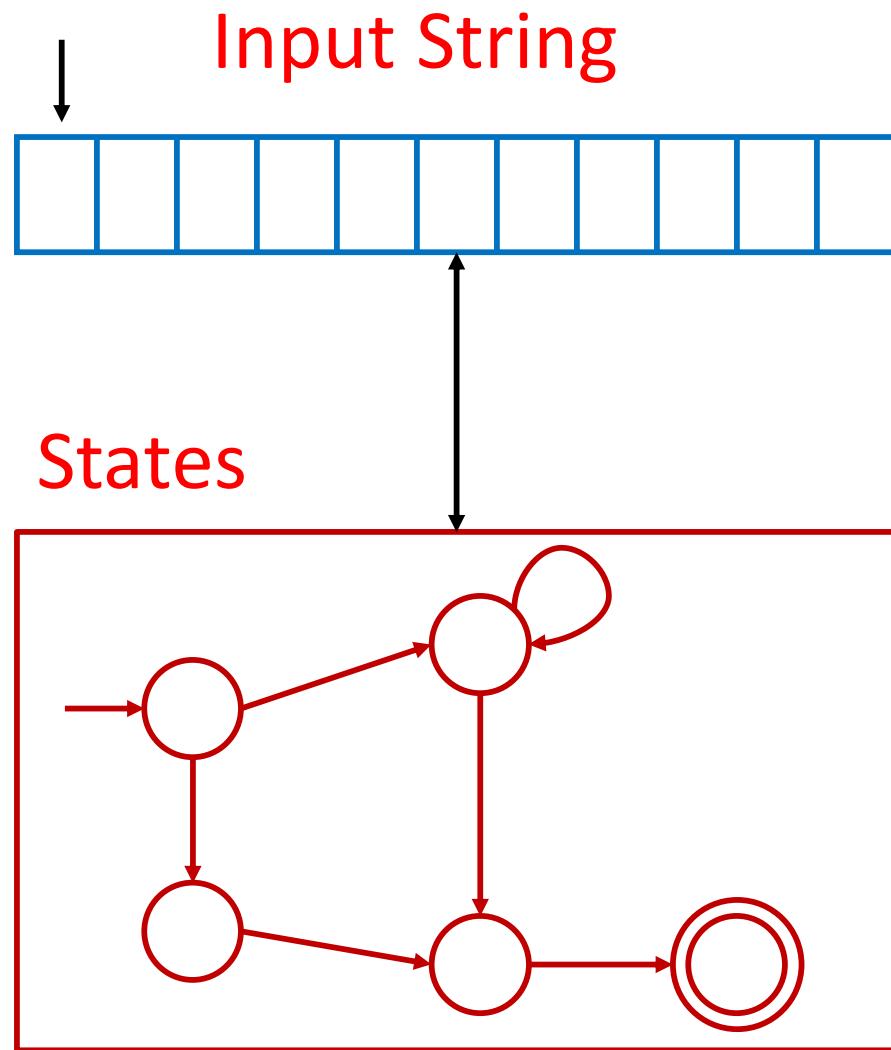
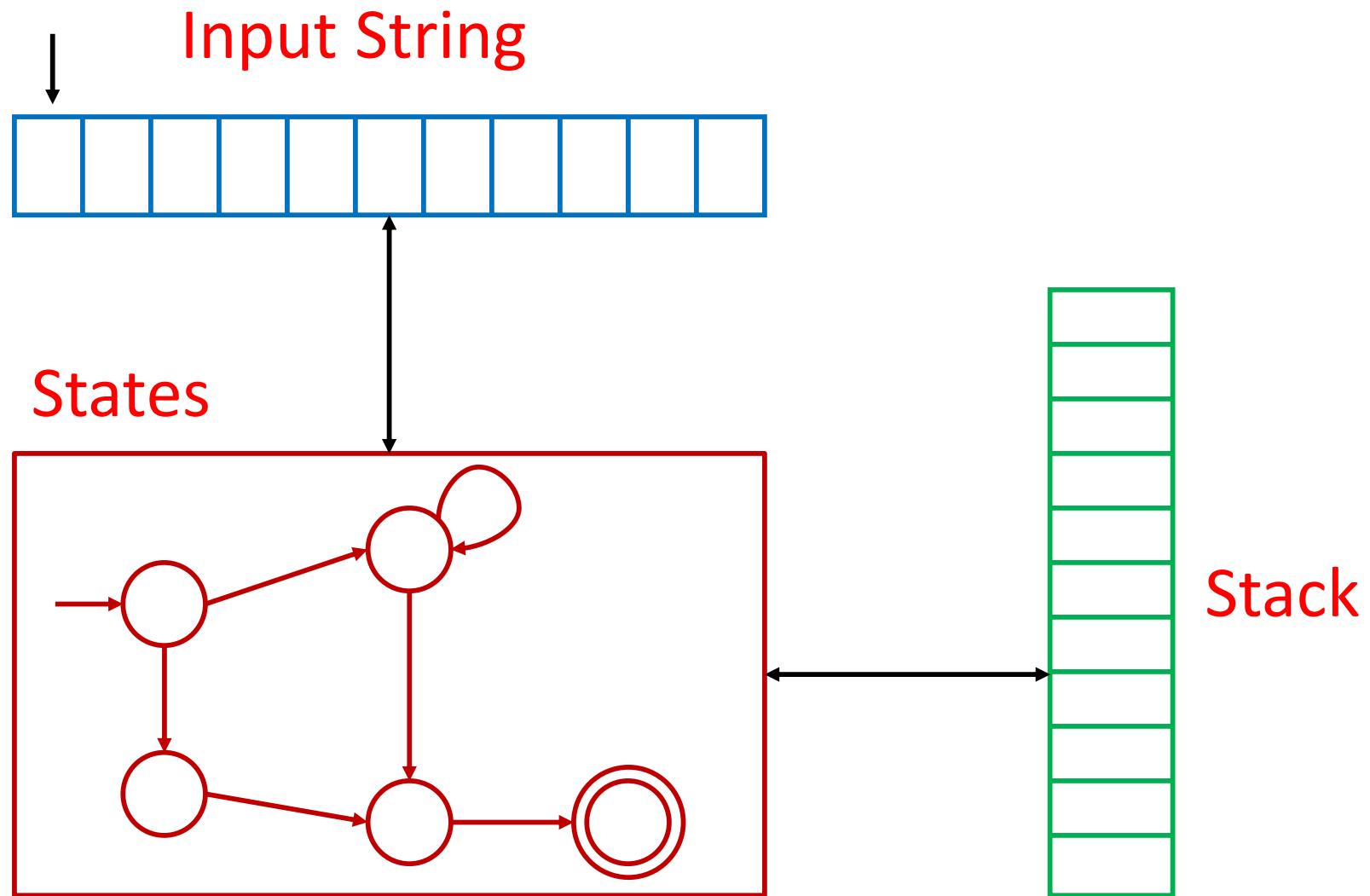


Pushdown Automata

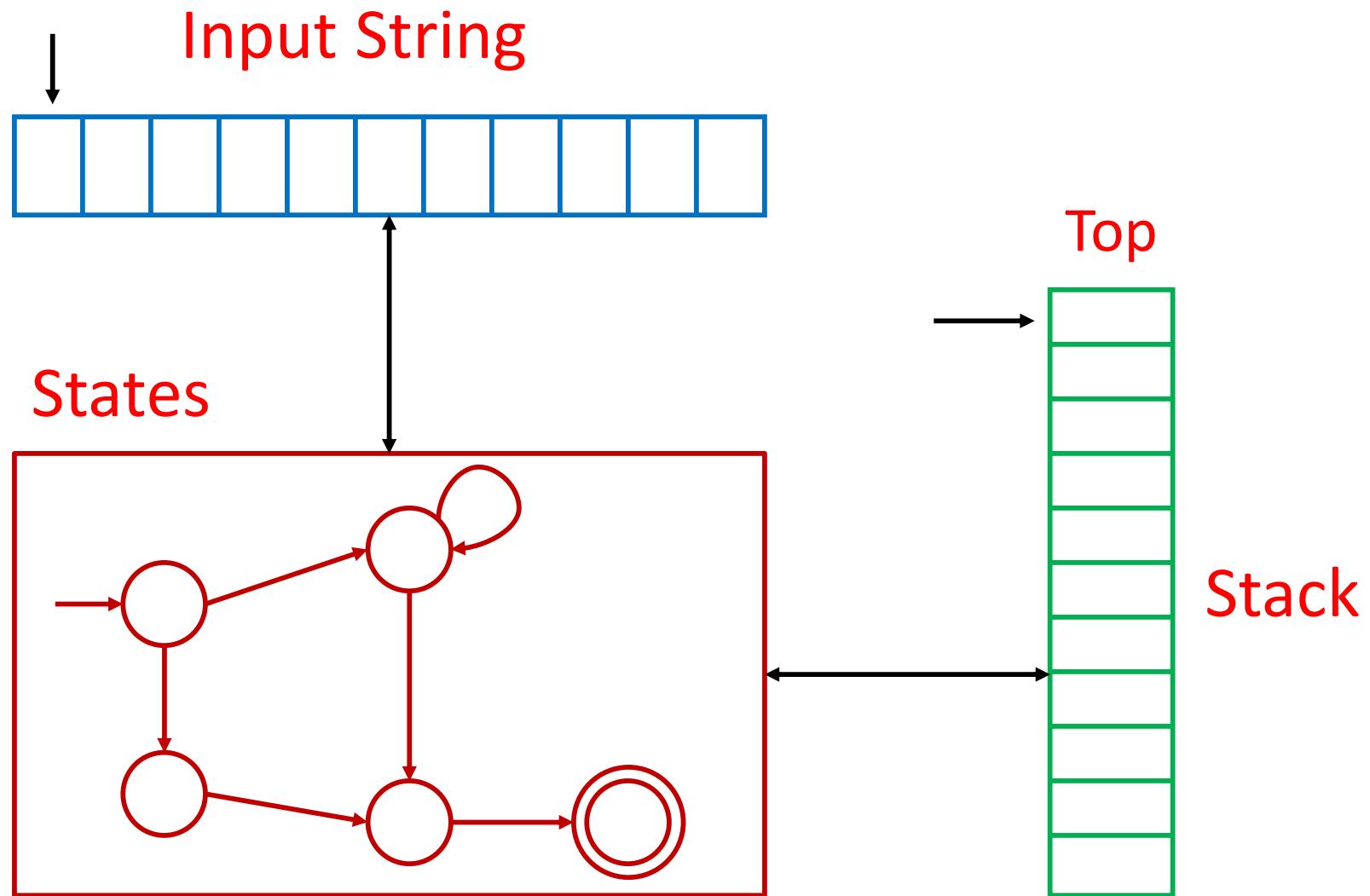
Finite Automaton



Pushdown Automaton - PDA



Pushdown Automaton - PDA



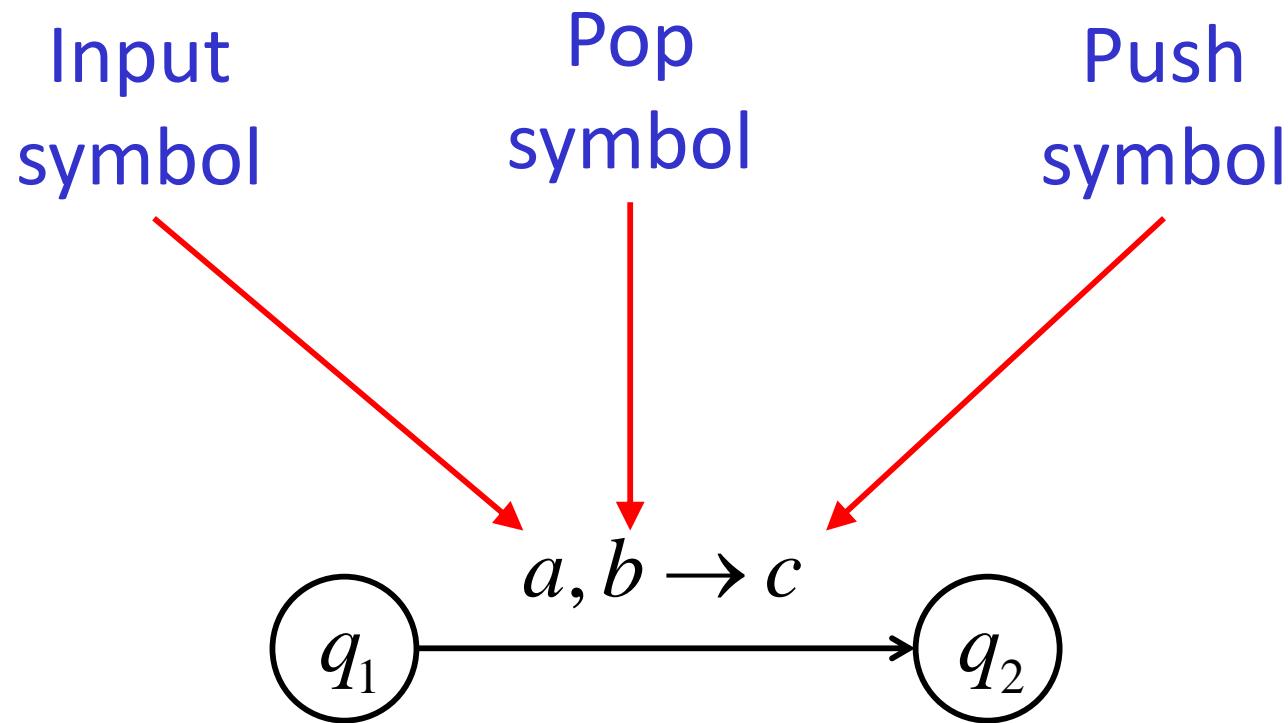
Initial Stack Symbol

Stack

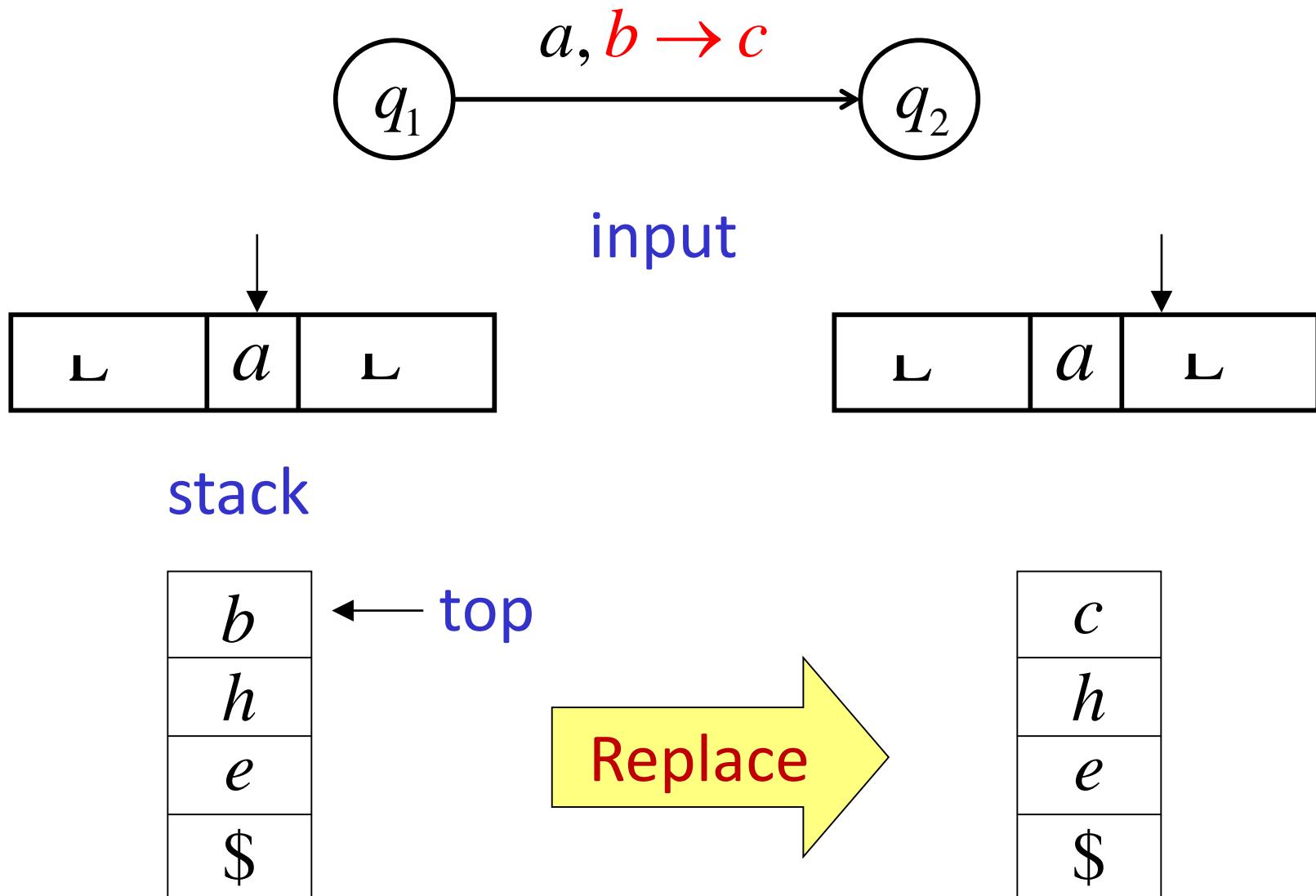


Bottom
special symbol

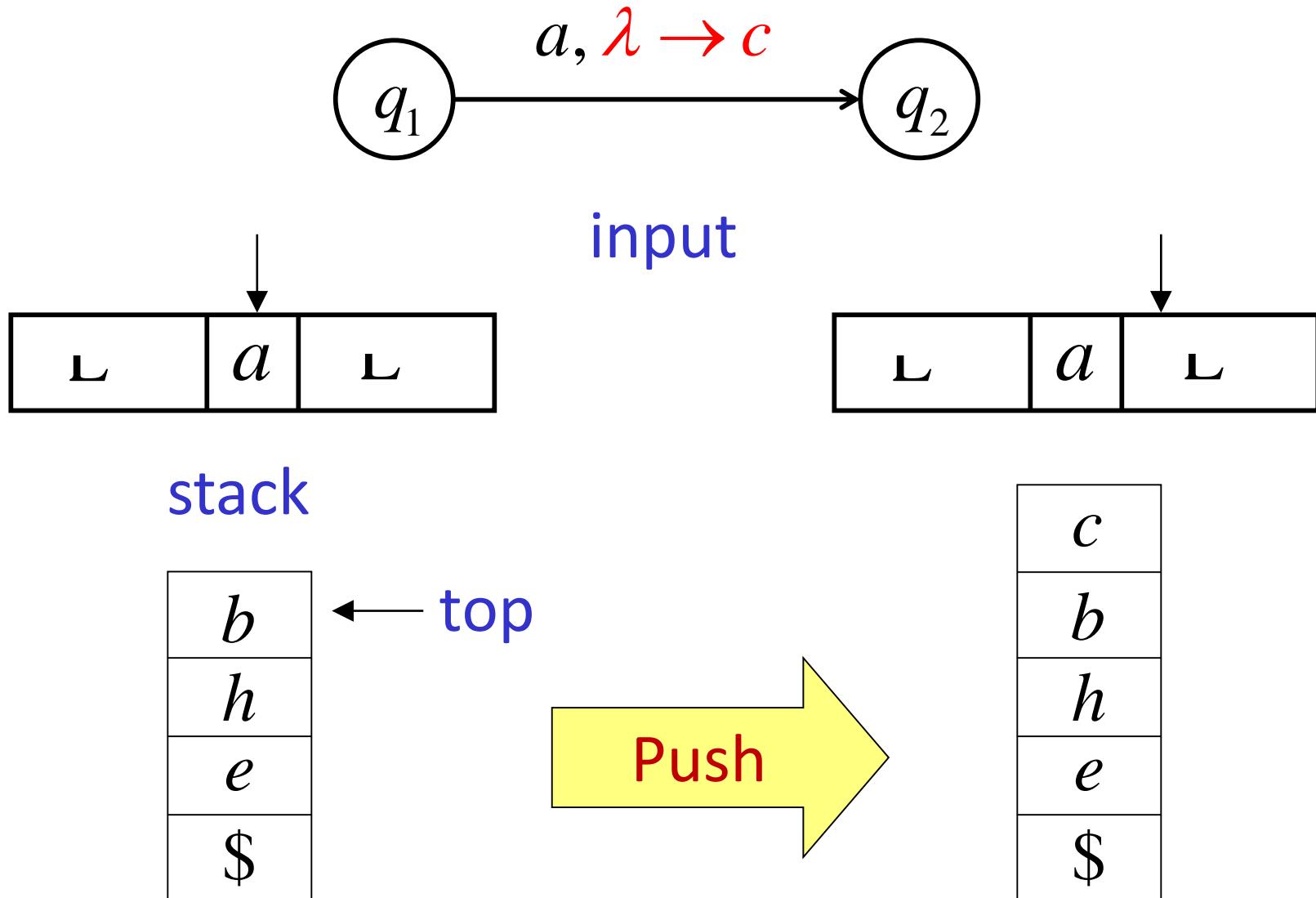
The States



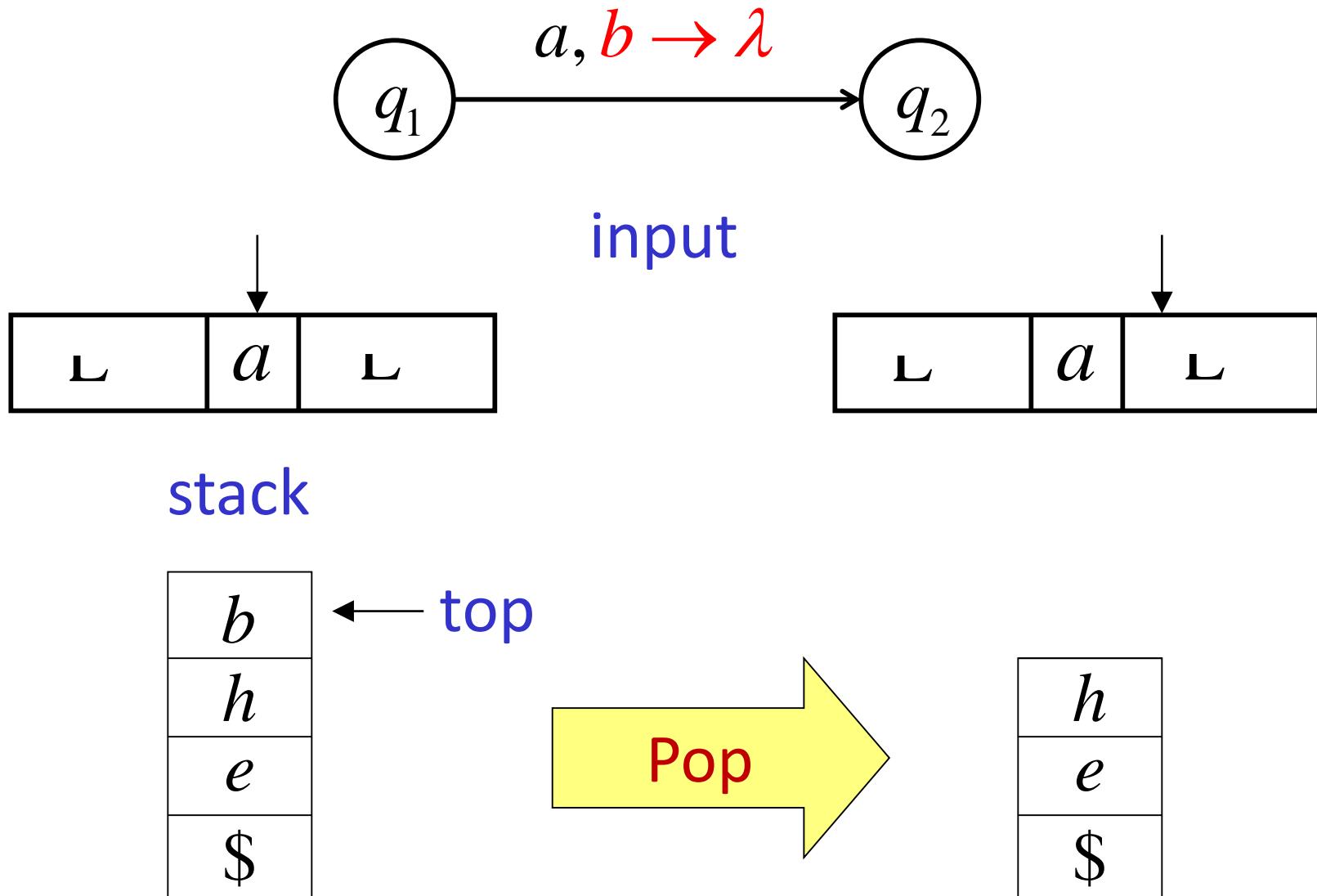
The Stack Operations



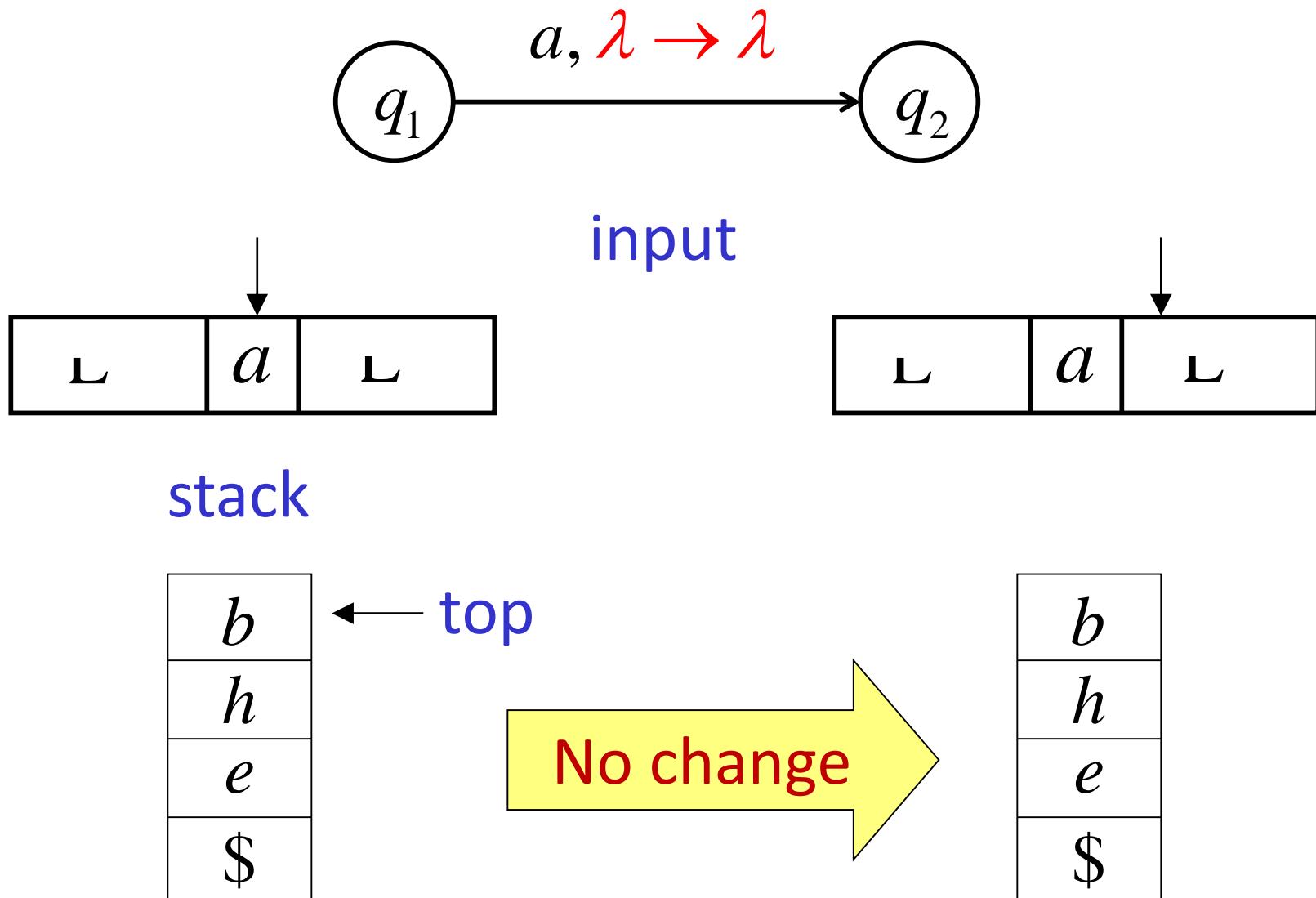
The Stack Operations



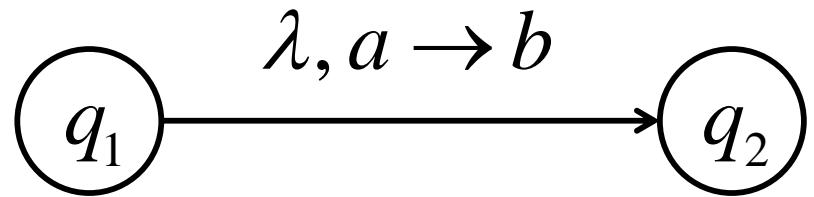
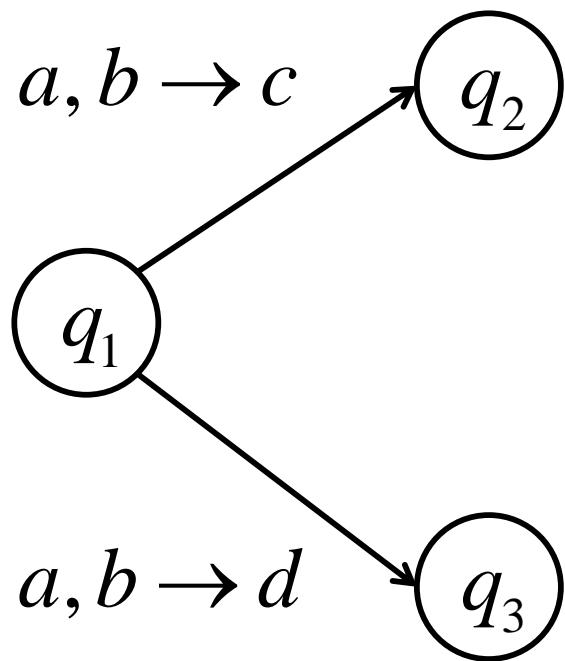
The Stack Operations



The Stack Operations



Non-Determinism

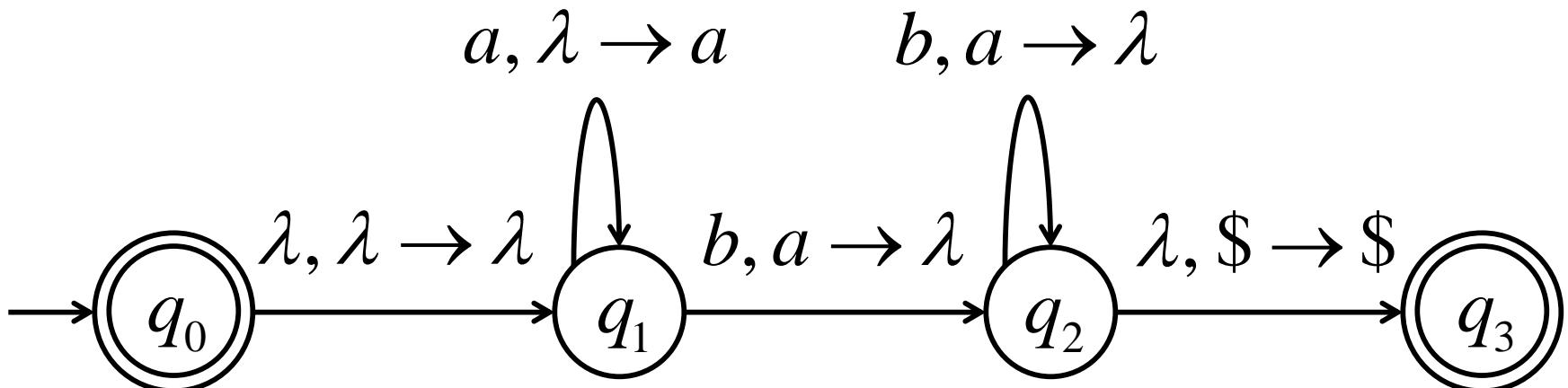


λ -transition

These are allowed transitions in
a non-deterministic PDA (NPDA)

NPDA: Non-Deterministic PDA

- Example:



Execution Example

Input

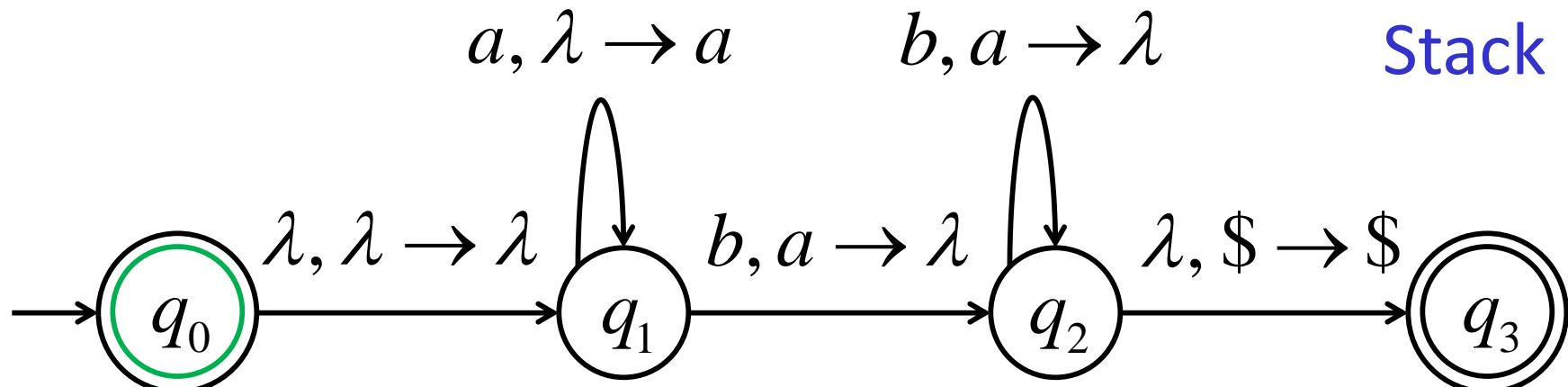
a	a	a	b	b	b
-----	-----	-----	-----	-----	-----

Time 0



\$

Stack



Execution Example

Input

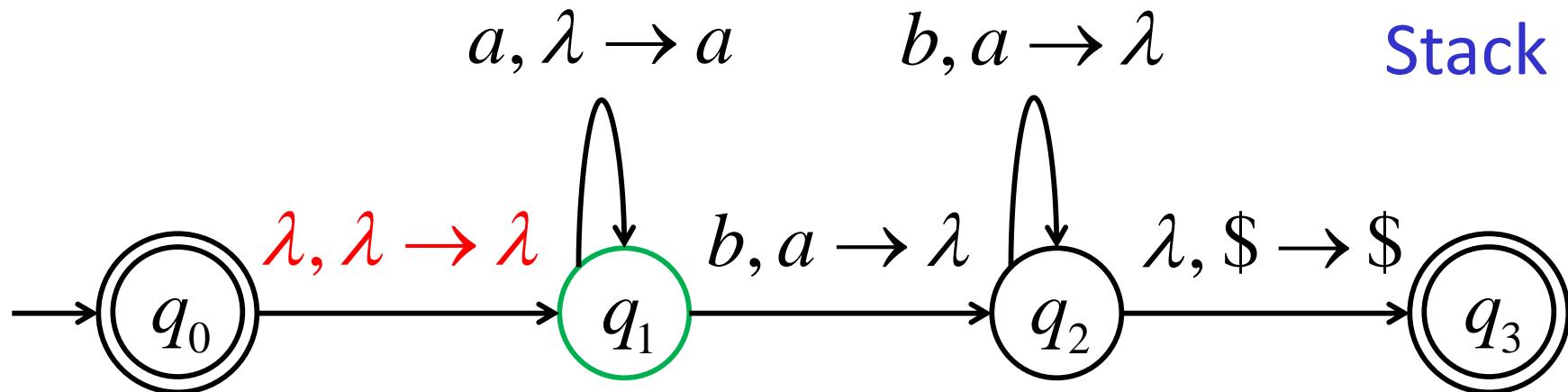
a	a	a	b	b	b
-----	-----	-----	-----	-----	-----

Time 1



\$

Stack



Execution Example

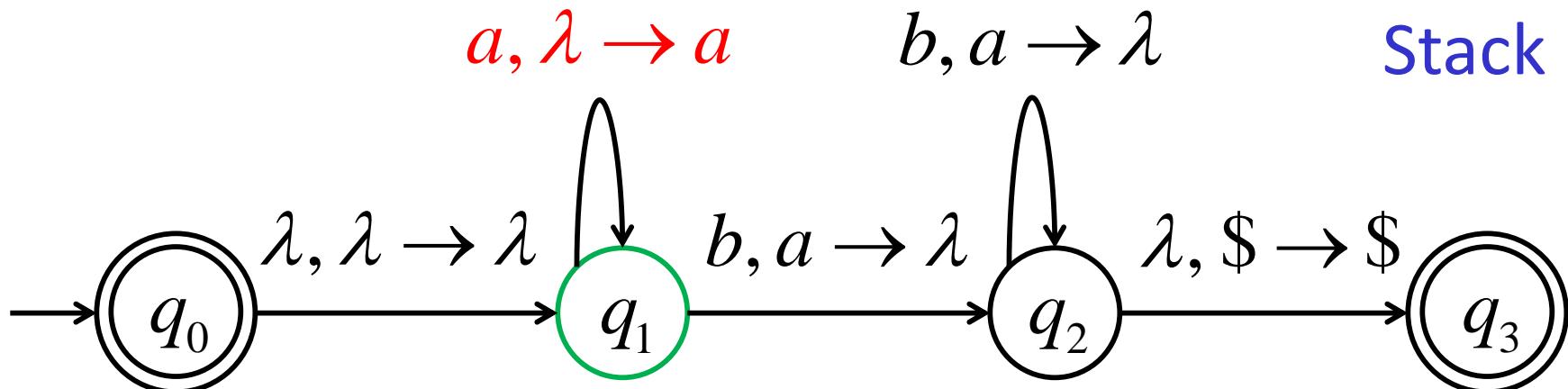
Input

a	a	a	b	b	b
-----	-----	-----	-----	-----	-----

Time 2

a
$\$$

Stack



Execution Example

Input

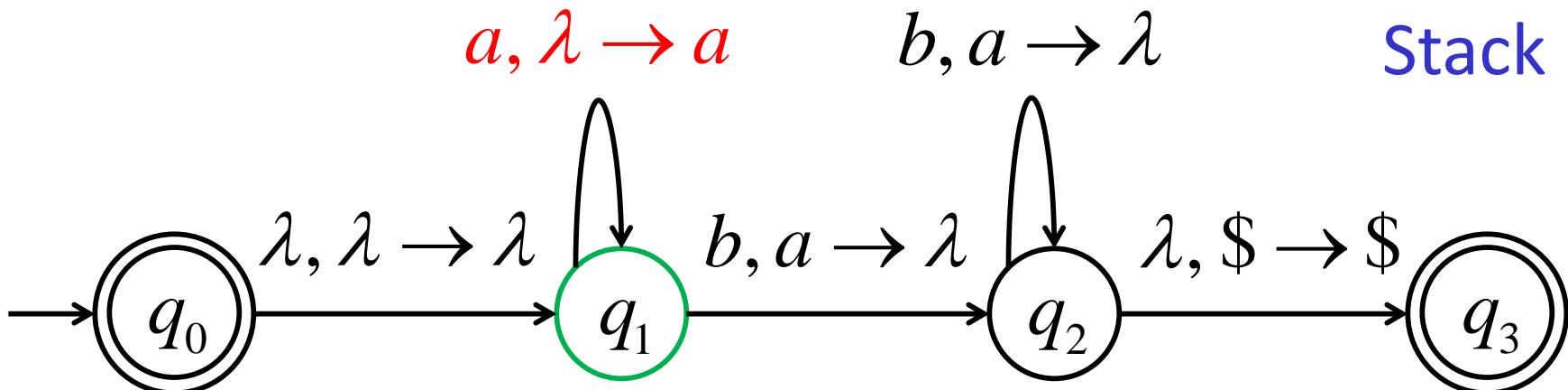
a	a	a	b	b	b
-----	-----	-----	-----	-----	-----



Time 3

a
a
\$

Stack



Execution Example

Input

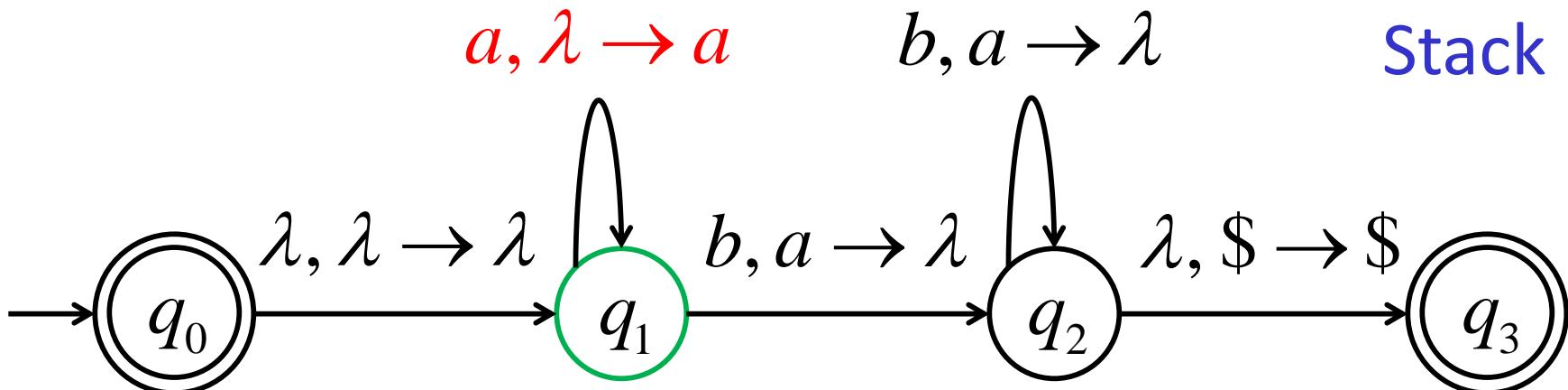
a	a	a	b	b	b
---	---	---	---	---	---



Time 4

a
a
a
\$

Stack



Execution Example

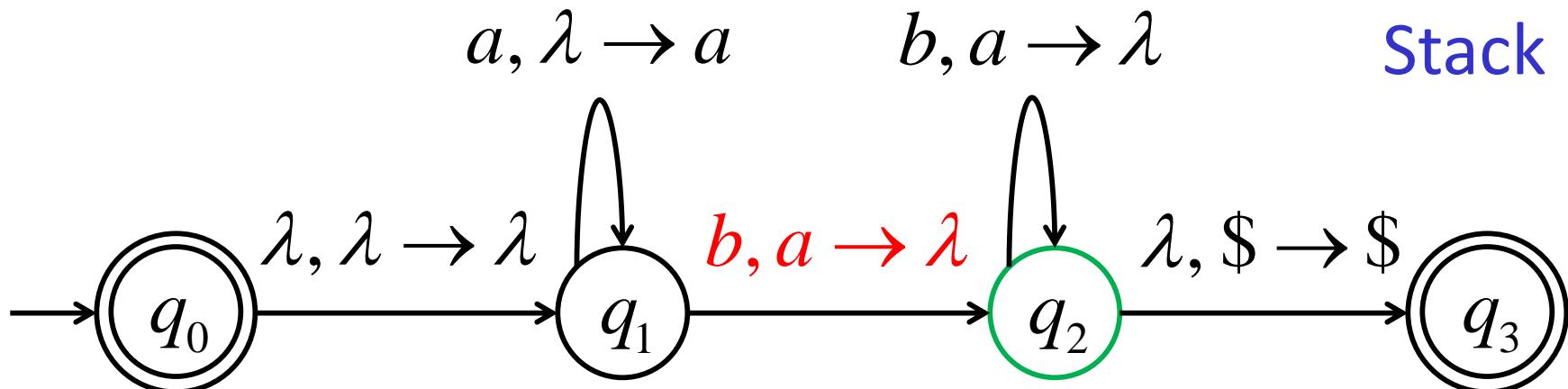
Input

a	a	a	b	b	b
-----	-----	-----	-----	-----	-----

Time 5

a
a
a
\$

Stack



Execution Example

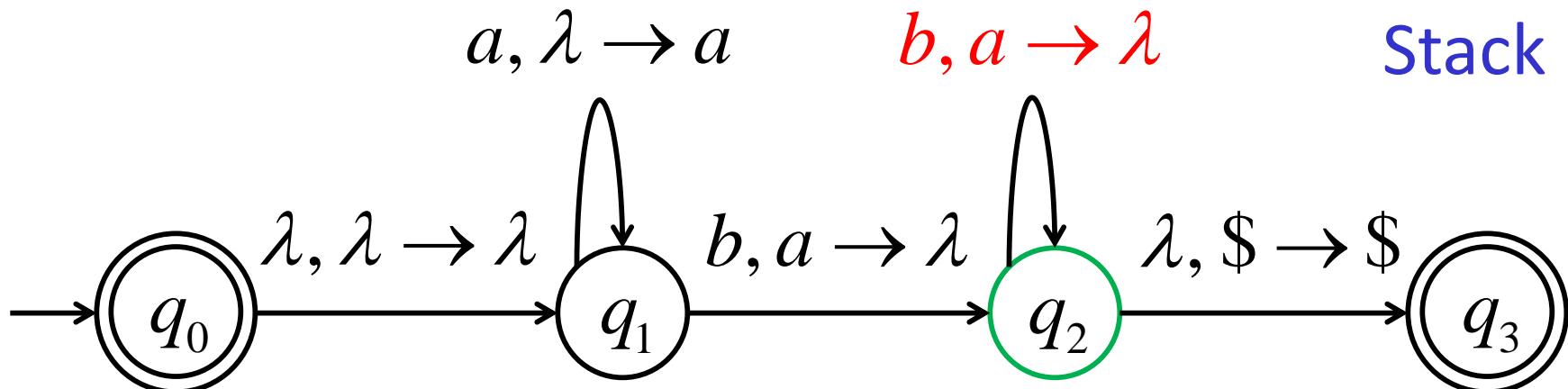
Input

a	a	a	b	b	b
-----	-----	-----	-----	-----	-----

Time 6

a
a
\$

Stack

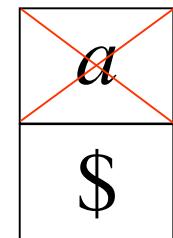


Execution Example

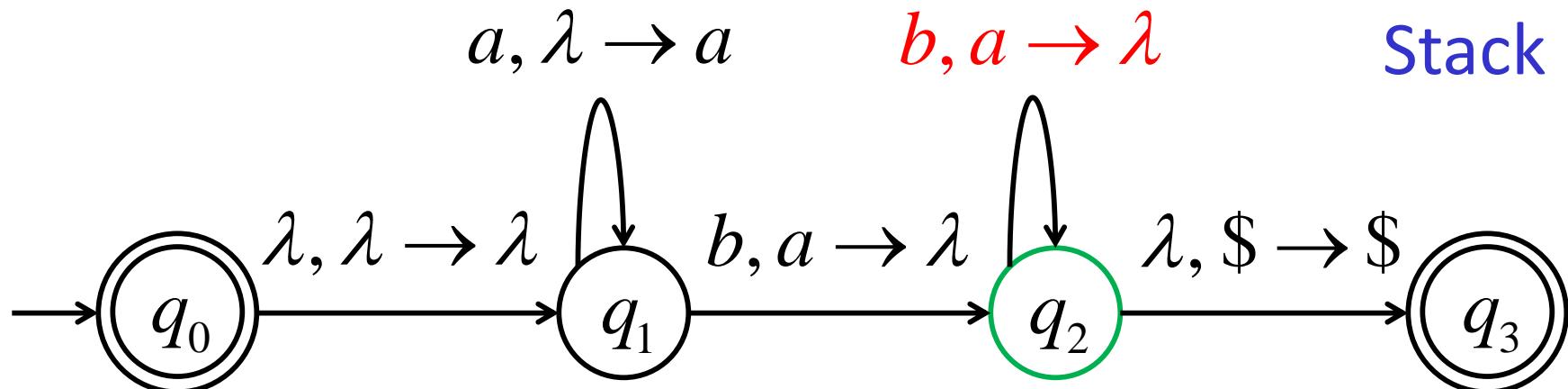
Input

a	a	a	b	b	b
-----	-----	-----	-----	-----	-----

Time 7



Stack



Execution Example

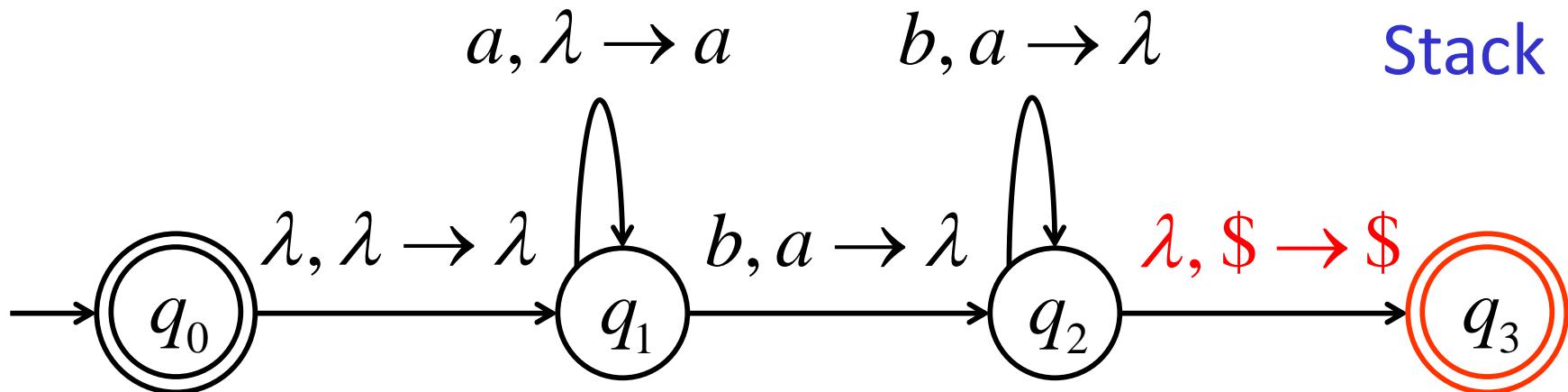
Input

a	a	a	b	b	b
-----	-----	-----	-----	-----	-----

Time 8

\$

Stack



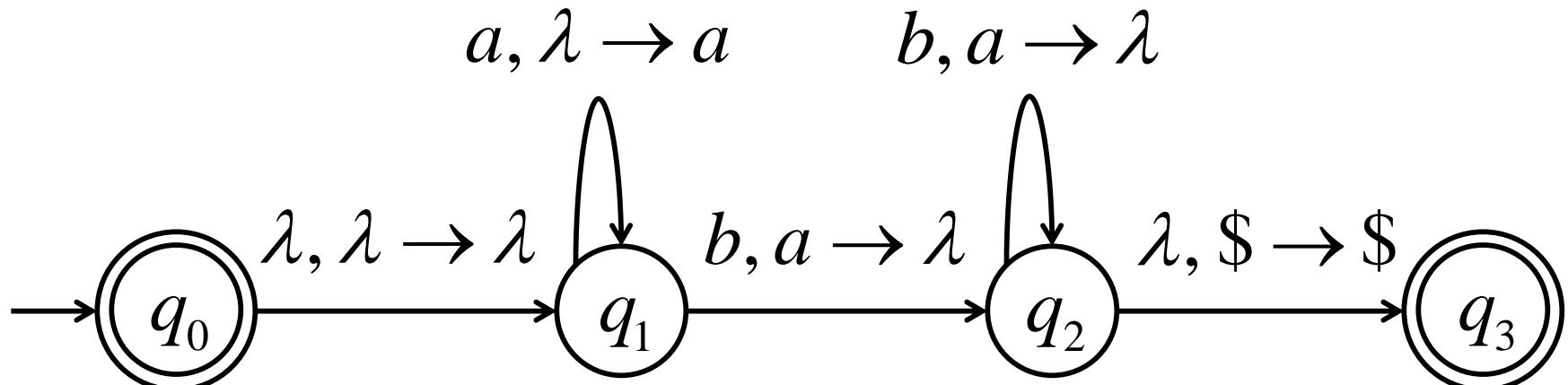
Pushdown Automaton - PDA

- A string is **accepted** if there is a computation such that:
 1. All the input is consumed
 2. The last state is a final state

At the end of the computation, we do not care about the **stack contents**

Pushdown Automaton - PDA

- The input string $aaabbb$ is accepted by the NPDA:

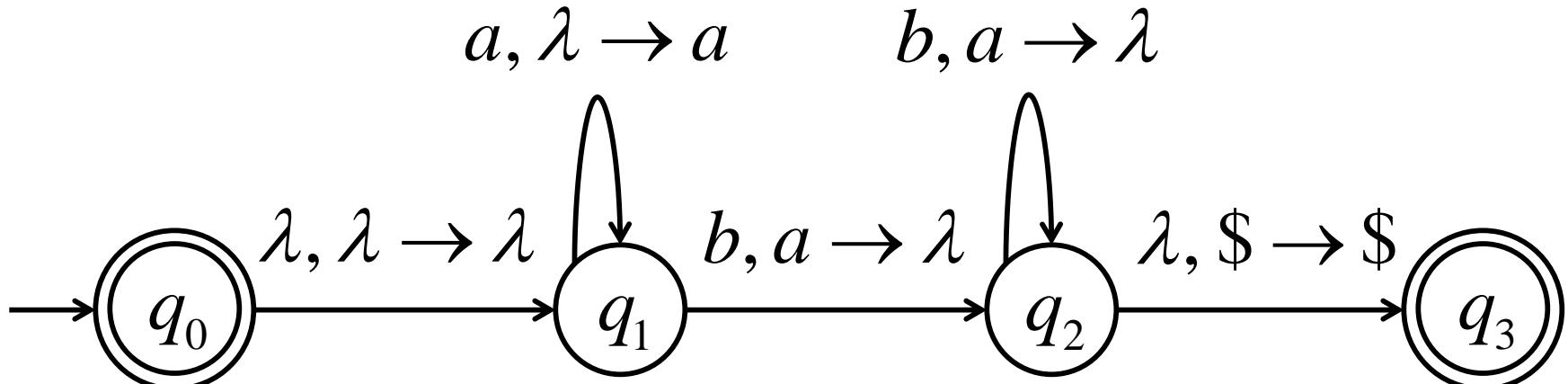


Pushdown Automaton - PDA

- In general,

$$L = \{a^n b^n : n \geq 0\}$$

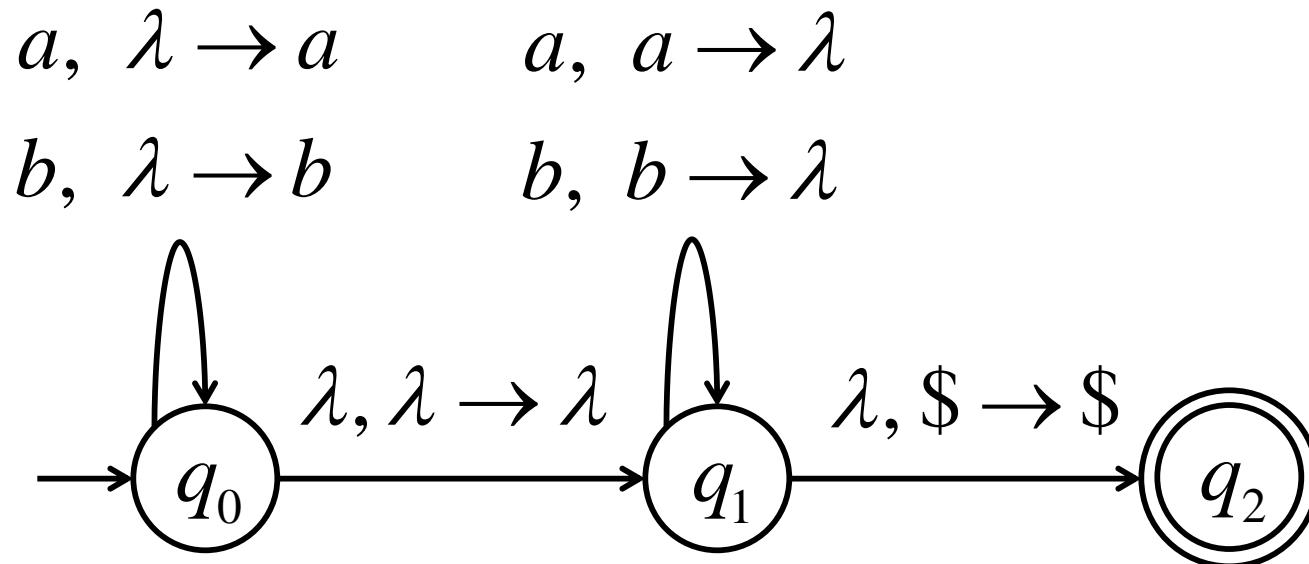
is the language accepted by the NPDA:



Example

- NDPA M

$$L(M) = \{ww^R\}$$



Execution Example

Input

a	b	b	a
-----	-----	-----	-----

Time 0

$$a, \lambda \rightarrow a$$

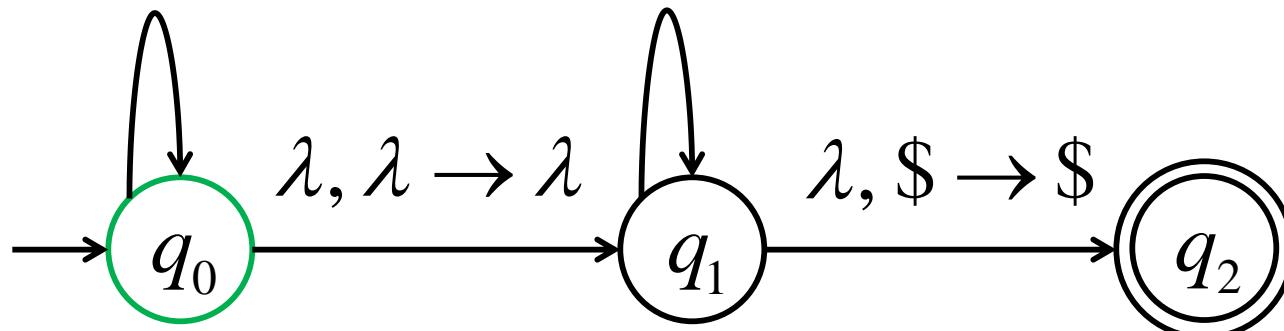
$$b, \lambda \rightarrow b$$

$$a, a \rightarrow \lambda$$

$$b, b \rightarrow \lambda$$

\$

Stack



Execution Example

Input

a	b	b	a
---	---	---	---



Time 1

a
\$

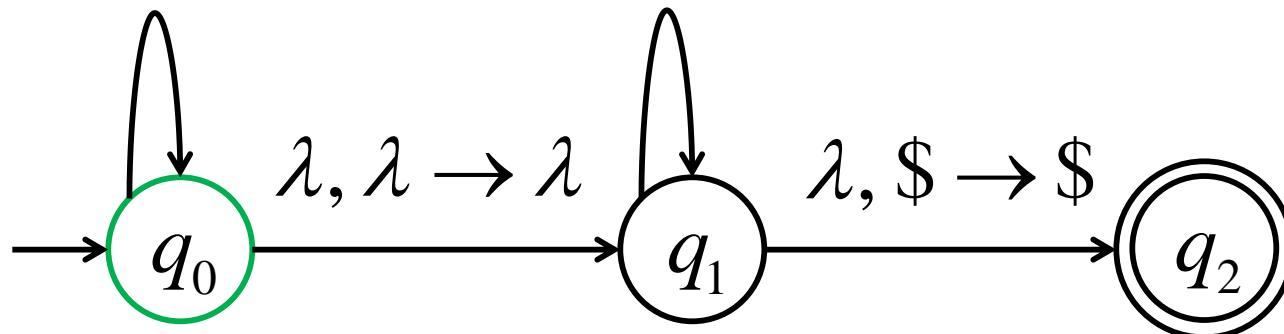
$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$

Stack



Execution Example

Input

a	b	b	a
-----	-----	-----	-----



Time 2

b
a
\$

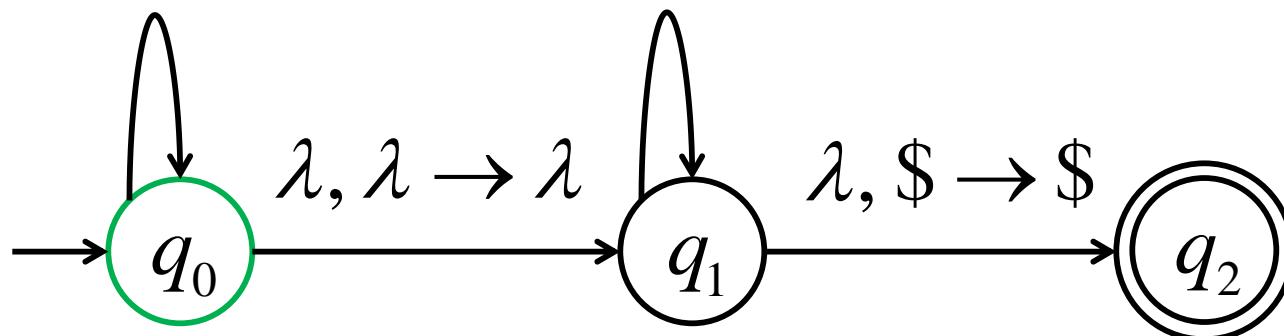
$$a, \lambda \rightarrow a$$

$$a, a \rightarrow \lambda$$

$$b, \lambda \rightarrow b$$

$$b, b \rightarrow \lambda$$

Stack



Execution Example

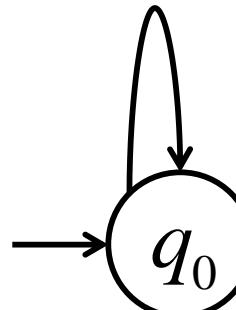
Input

a	b	b	a
---	---	---	---



$$a, \lambda \rightarrow a$$

$$b, \lambda \rightarrow b$$

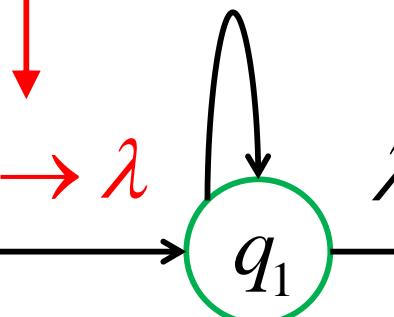


Guess the middle
of string

Time 3

$$a, a \rightarrow \lambda$$

$$b, b \rightarrow \lambda$$



$$\lambda, \$ \rightarrow \$$$



b
a
\$

Stack

Execution Example

Input

a	b	b	a
-----	-----	-----	-----

Time 4

b
a
\$

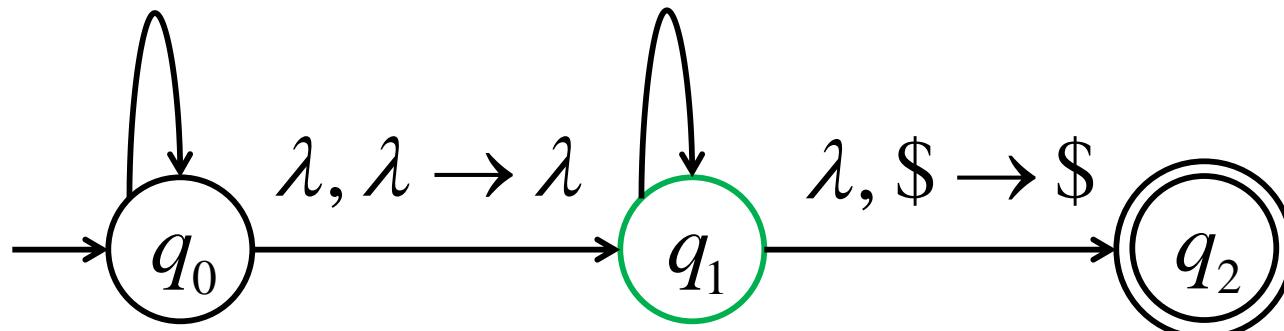
$$a, \lambda \rightarrow a$$

$$a, a \rightarrow \lambda$$

$$b, \lambda \rightarrow b$$

$$b, b \rightarrow \lambda$$

Stack



Execution Example

Input

a	b	b	a
-----	-----	-----	-----

Time 5

$$a, \lambda \rightarrow a$$

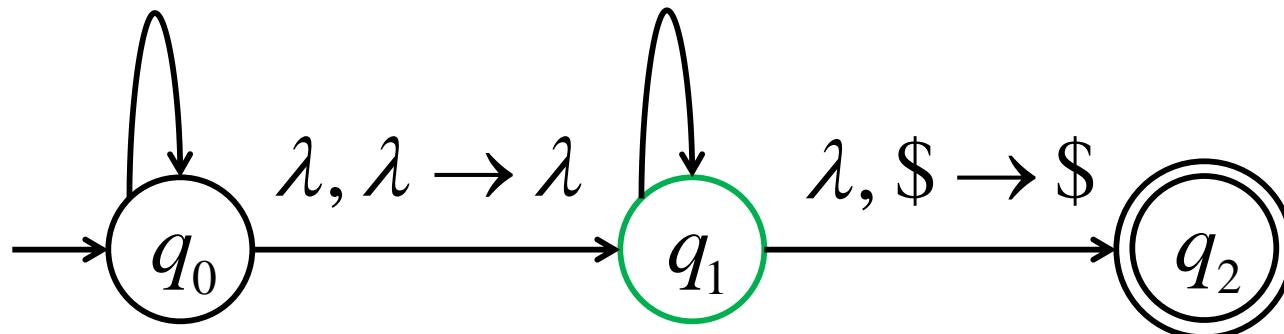
$$b, \lambda \rightarrow b$$

$$a, a \rightarrow \lambda$$

$$b, b \rightarrow \lambda$$

a
\$

Stack



Execution Example

Input

a	b	b	a
-----	-----	-----	-----

Time 6

$$a, \lambda \rightarrow a$$

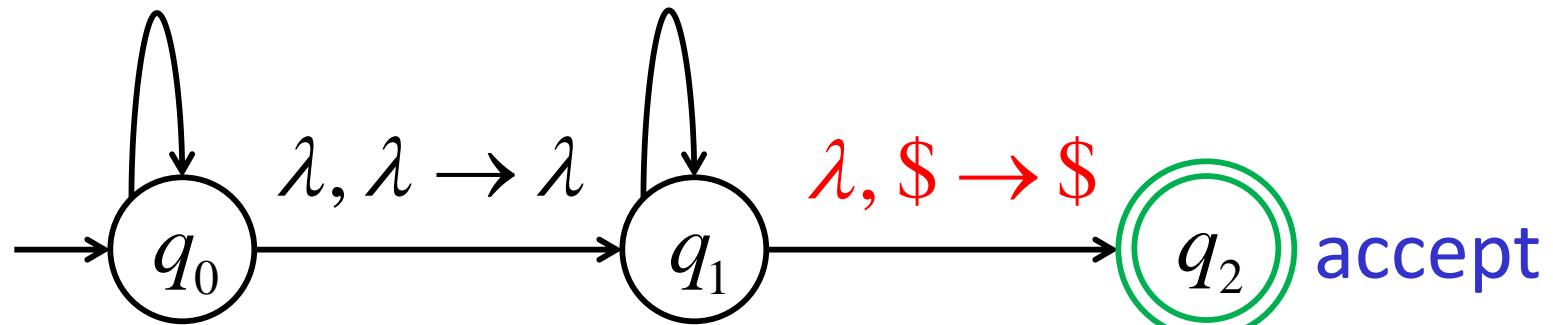
$$b, \lambda \rightarrow b$$

$$a, a \rightarrow \lambda$$

$$b, b \rightarrow \lambda$$

\$

Stack



Rejection Example

Input

a	b	b	b
-----	-----	-----	-----

Time 0

$$a, \lambda \rightarrow a$$

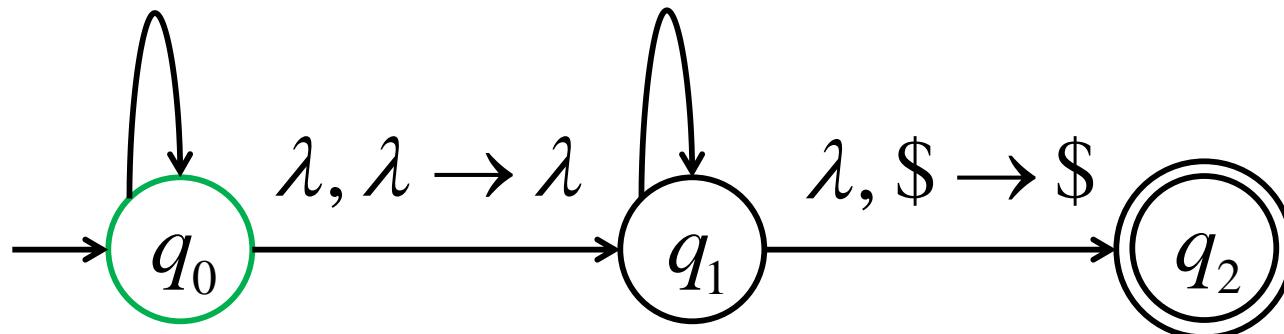
$$b, \lambda \rightarrow b$$

$$a, a \rightarrow \lambda$$

$$b, b \rightarrow \lambda$$

\$

Stack



Rejection Example

Input

a	b	b	b
-----	-----	-----	-----



Time 1

a
\$

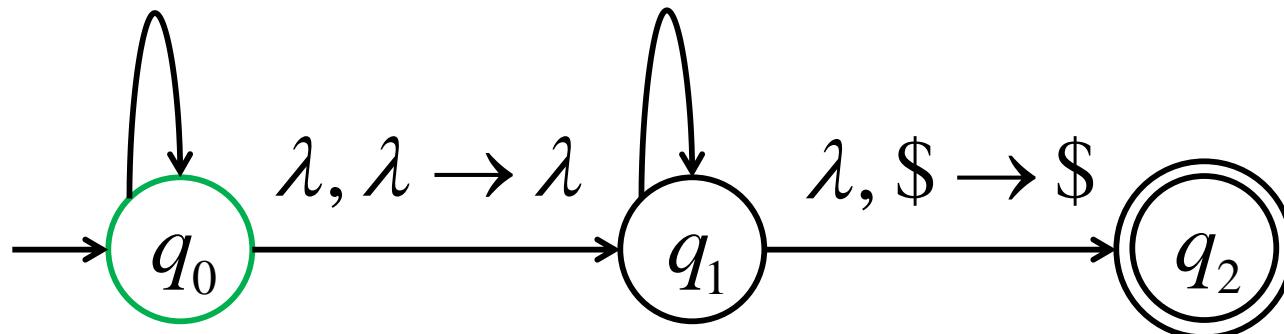
$a, \lambda \rightarrow a$

$b, \lambda \rightarrow b$

$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$

Stack



Rejection Example

Input

a	b	b	b
-----	-----	-----	-----



Time 2

b
a
\$

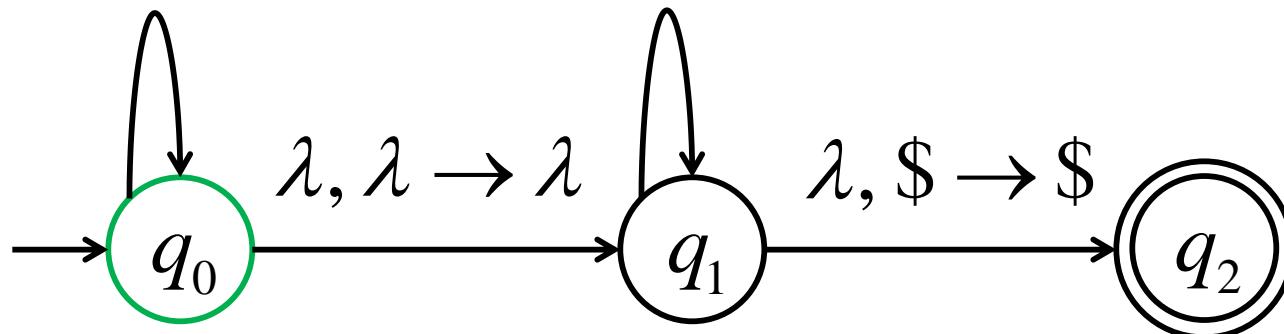
$$a, \lambda \rightarrow a$$

$$a, a \rightarrow \lambda$$

$$b, \lambda \rightarrow b$$

$$b, b \rightarrow \lambda$$

Stack



Rejection Example

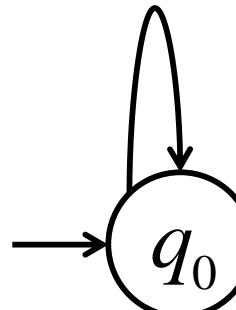
Input

a	b	b	b
-----	-----	-----	-----



$$a, \lambda \rightarrow a$$

$$b, \lambda \rightarrow b$$

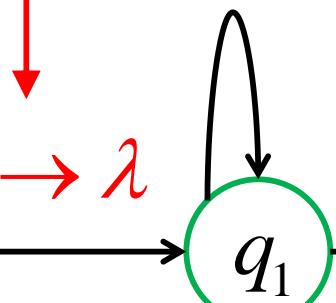


Guess the middle
of string

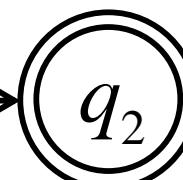
Time 3

$$a, a \rightarrow \lambda$$

$$b, b \rightarrow \lambda$$



$$\lambda, \$ \rightarrow \$$$



b
a
$\$$

Stack

Rejection Example

Input

a	b	b	b
-----	-----	-----	-----

Time 4

b
a
\$

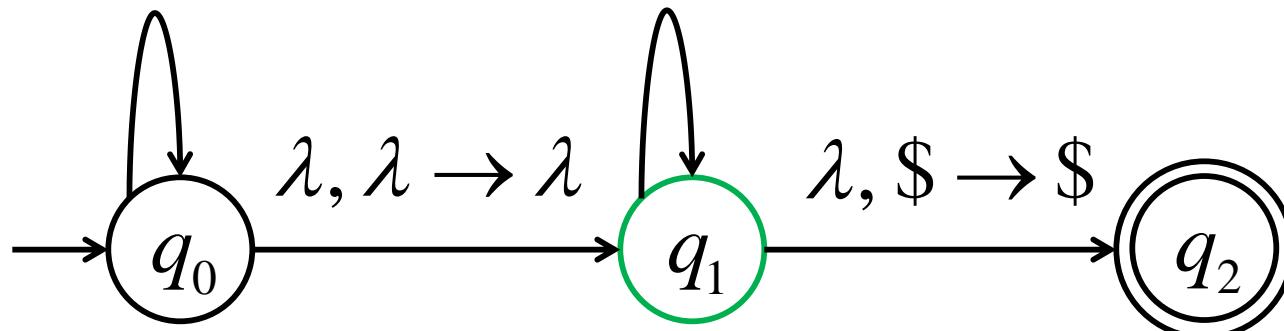
$$a, \lambda \rightarrow a$$

$$a, a \rightarrow \lambda$$

$$b, \lambda \rightarrow b$$

$$b, b \rightarrow \lambda$$

Stack



Rejection Example

Input

a	b	b	b
---	---	---	---

Time 5

There is no possible transition;
input is not consumed

$$a, \lambda \rightarrow a$$

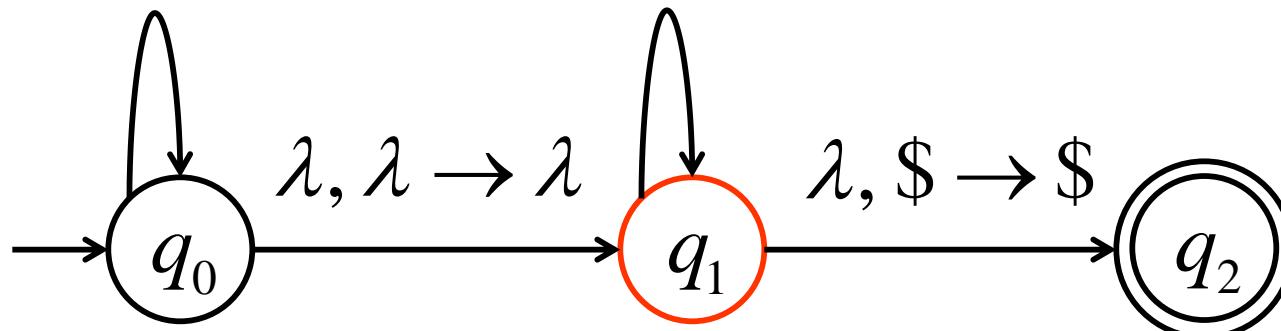
$$b, \lambda \rightarrow b$$

$$a, a \rightarrow \lambda$$

$$b, b \rightarrow \lambda$$

a
\$

Stack



Rejection Example: Another Computation

Input

a	b	b	b
-----	-----	-----	-----

Time 0

$$a, \lambda \rightarrow a$$

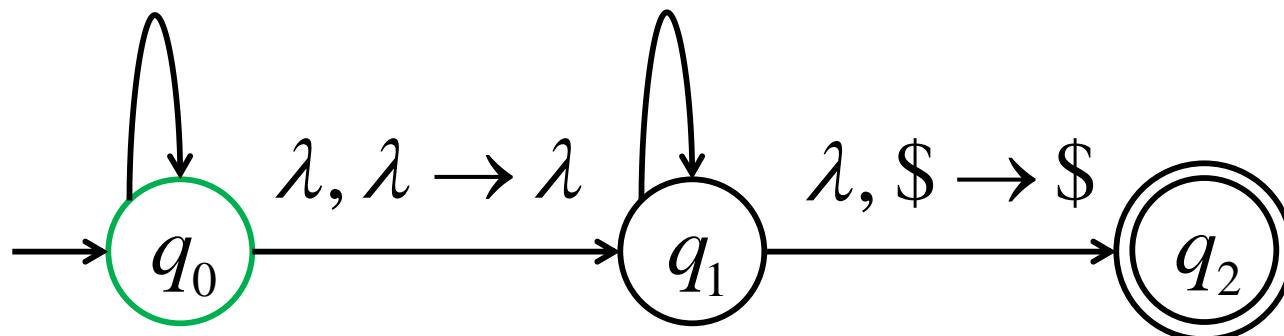
$$b, \lambda \rightarrow b$$

$$a, a \rightarrow \lambda$$

$$b, b \rightarrow \lambda$$

\$

Stack



Rejection Example: Another Computation

Input

a	b	b	b
-----	-----	-----	-----



Time 1

$a, \lambda \rightarrow a$

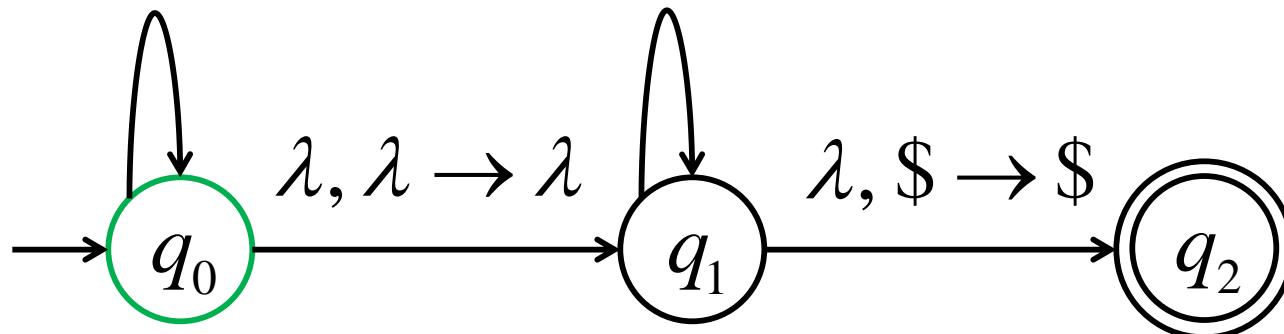
$b, \lambda \rightarrow b$

$a, a \rightarrow \lambda$

$b, b \rightarrow \lambda$

a
\$

Stack



Rejection Example: Another Computation

Input

a	b	b	b
-----	-----	-----	-----



Time 2

b
a
\$

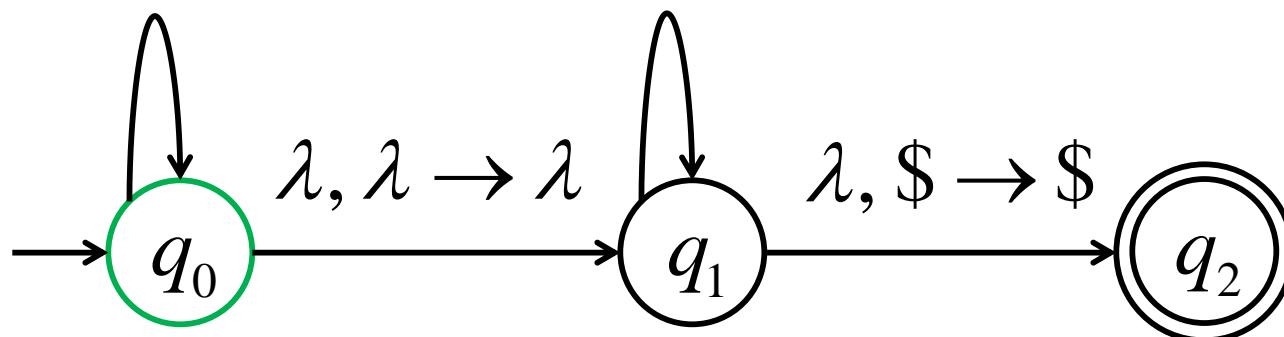
$$a, \lambda \rightarrow a$$

$$b, \lambda \rightarrow b$$

$$a, a \rightarrow \lambda$$

$$b, b \rightarrow \lambda$$

Stack



Rejection Example: Another Computation

Input

a	b	b	b
-----	-----	-----	-----



Time 3

b
b
a
\$

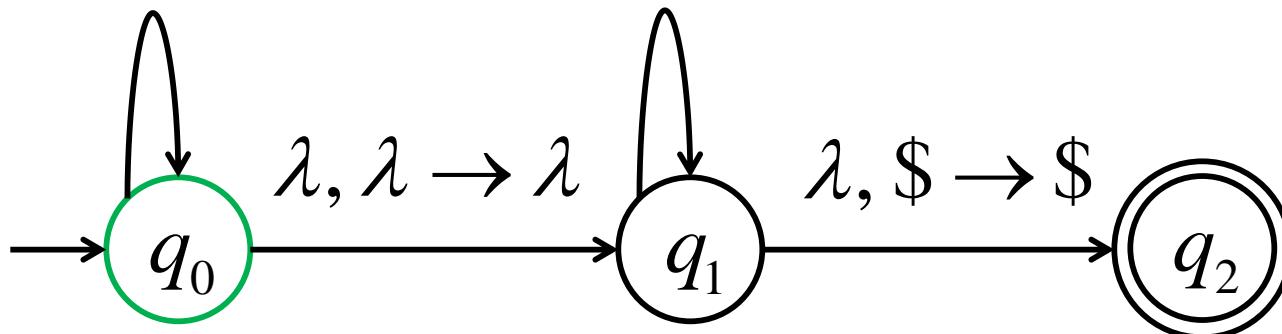
$$a, \lambda \rightarrow a$$

$$a, a \rightarrow \lambda$$

$$b, \lambda \rightarrow b$$

$$b, b \rightarrow \lambda$$

Stack



Rejection Example: Another Computation

Input

a	b	b	b
-----	-----	-----	-----

Time 4

b
b
b
a
\$

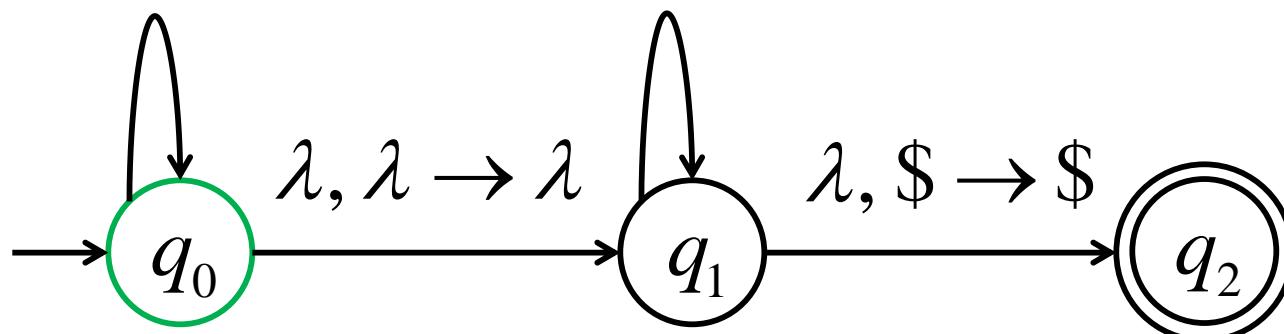
$$a, \lambda \rightarrow a$$

$$b, \lambda \rightarrow b$$

$$a, a \rightarrow \lambda$$

$$b, b \rightarrow \lambda$$

Stack



Rejection Example: Another Computation

Input

a	b	b	b
-----	-----	-----	-----

Time 5

b
b
b
a
\$

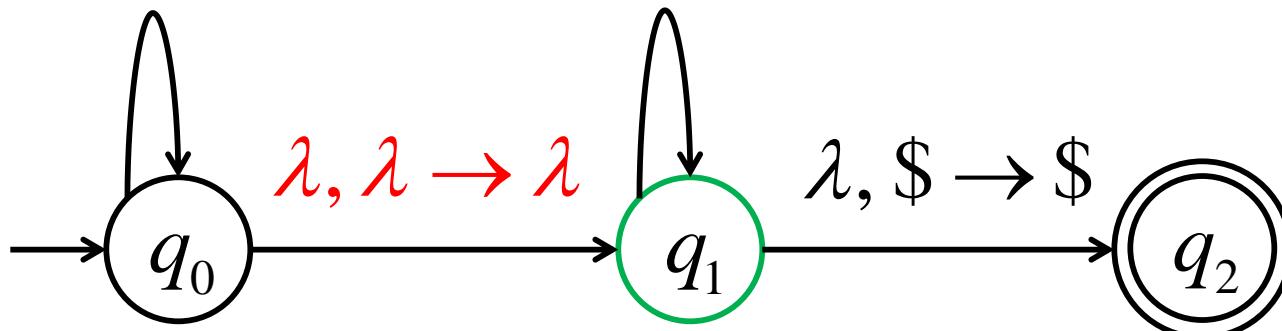
$$a, \lambda \rightarrow a$$

$$b, \lambda \rightarrow b$$

$$a, a \rightarrow \lambda$$

$$b, b \rightarrow \lambda$$

Stack



Rejection Example: Another Computation

Input

a	b	b	b
---	---	---	---

Time 6

No final state
is reached

b
b
b
a
\$

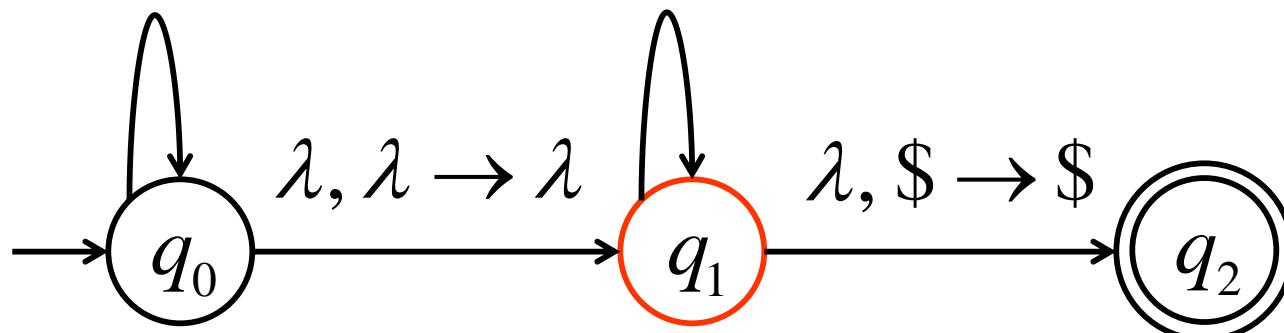
$$a, \lambda \rightarrow a$$

$$a, a \rightarrow \lambda$$

$$b, \lambda \rightarrow b$$

$$b, b \rightarrow \lambda$$

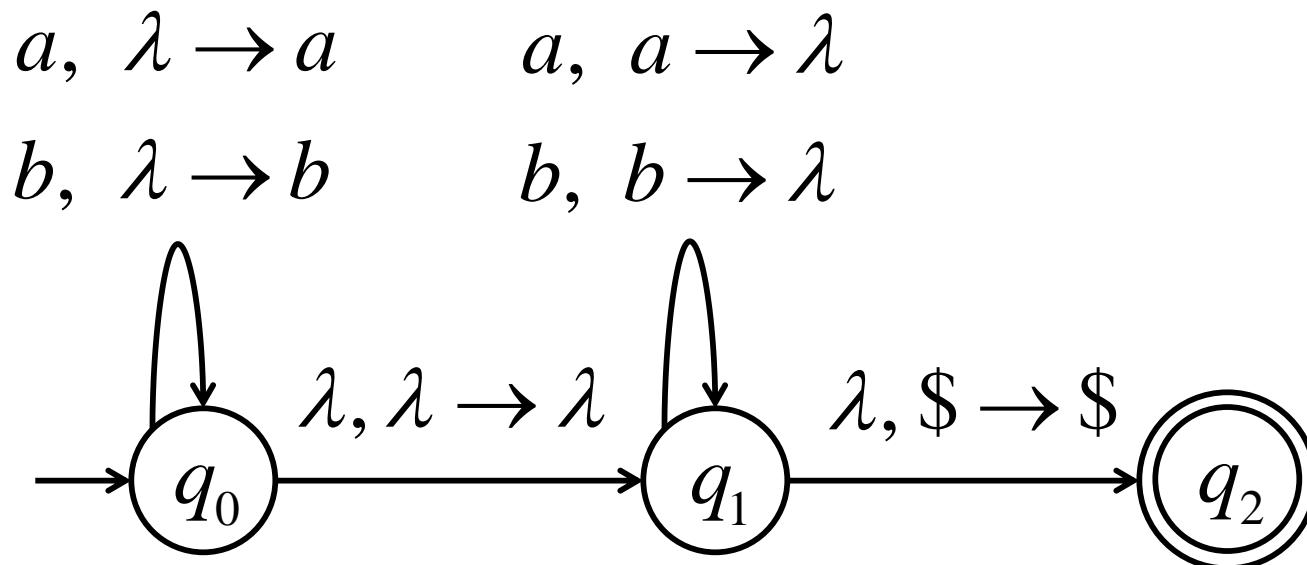
Stack



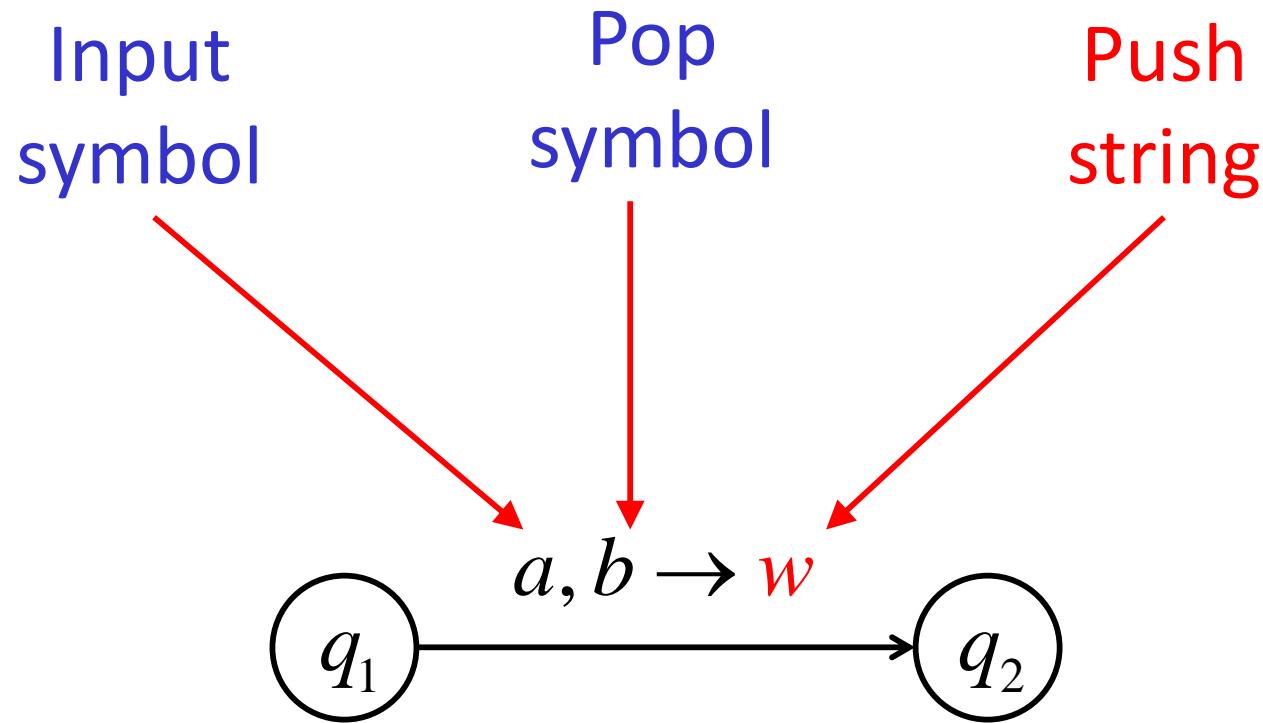
Rejection Example: Another Computation

There is no computation that accepts string $abbb$

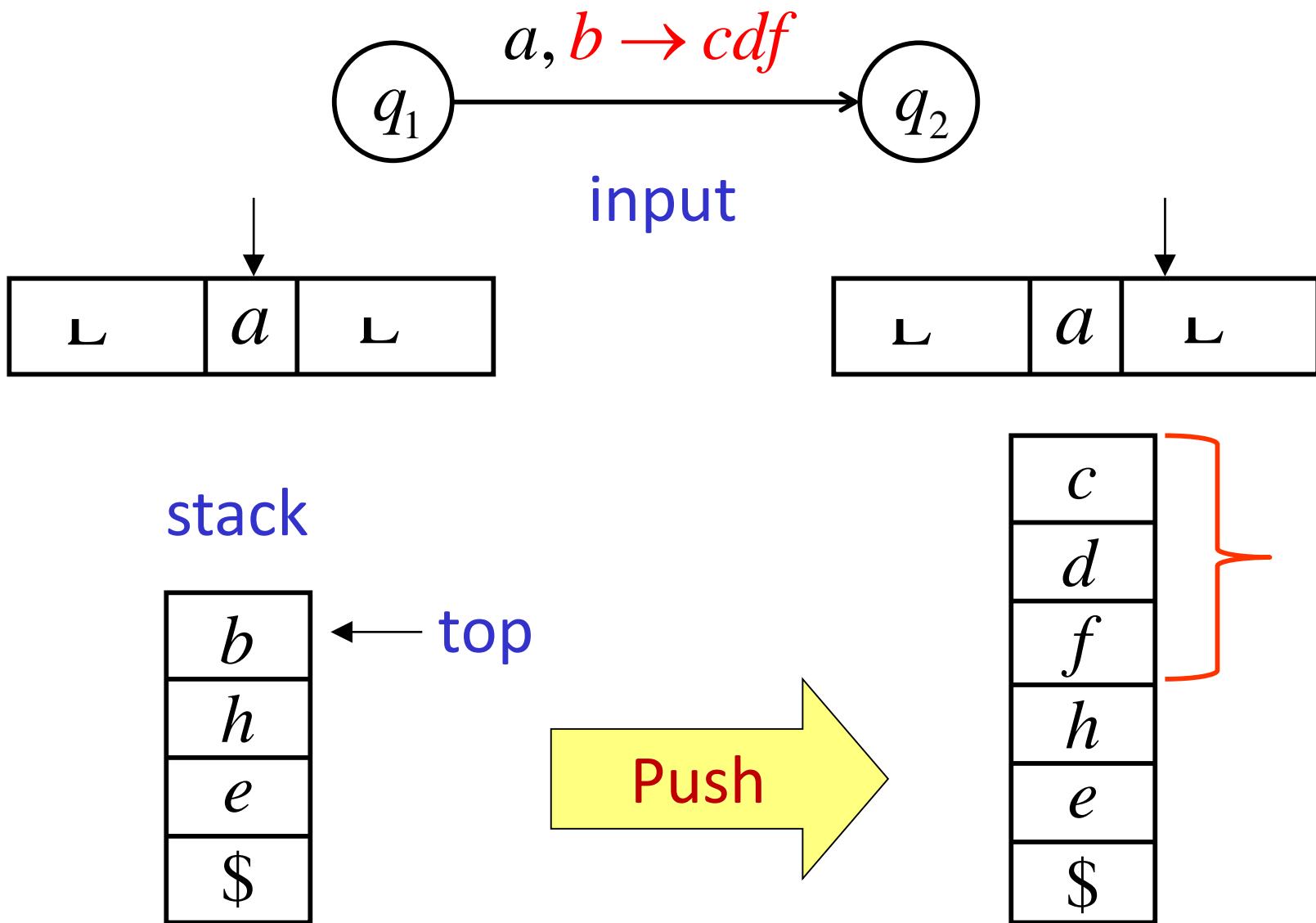
$$abbb \notin L(M)$$



Pushing Strings



Example



Another NPDA example

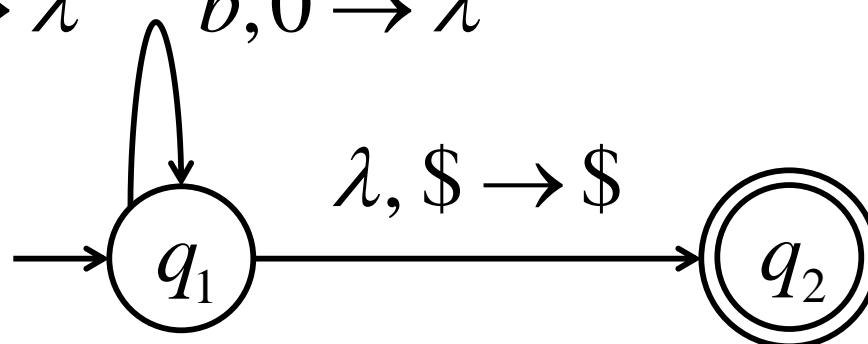
NPDA M

$$L(M) = \{w : n_a(w) = n_b(w)\}$$

$$a, \$ \rightarrow 0\$ \quad b, \$ \rightarrow 1\$$$

$$a, 0 \rightarrow 00 \quad b, 1 \rightarrow 11$$

$$a, 1 \rightarrow \lambda \quad b, 0 \rightarrow \lambda$$



Execution Example

Input

a	b	b	b	a	a
---	---	---	---	---	---

Time 0



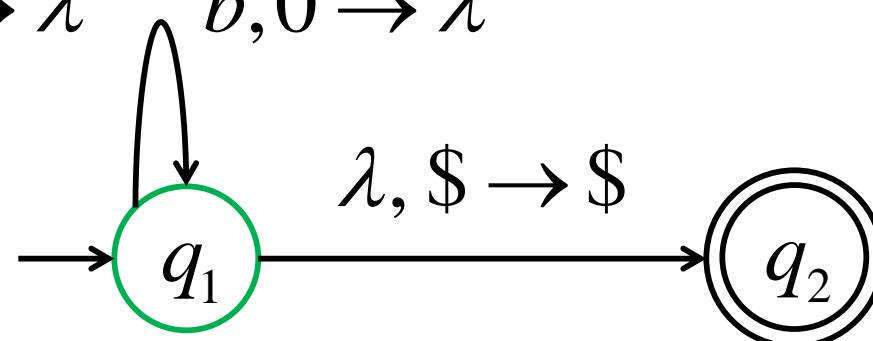
$a, \$ \rightarrow 0\$$ $b, \$ \rightarrow 1\$$

$a, 0 \rightarrow 00$ $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$ $b, 0 \rightarrow \lambda$

\$

Stack



Execution Example

Input

a	b	b	b	a	a
-----	-----	-----	-----	-----	-----

Time 1

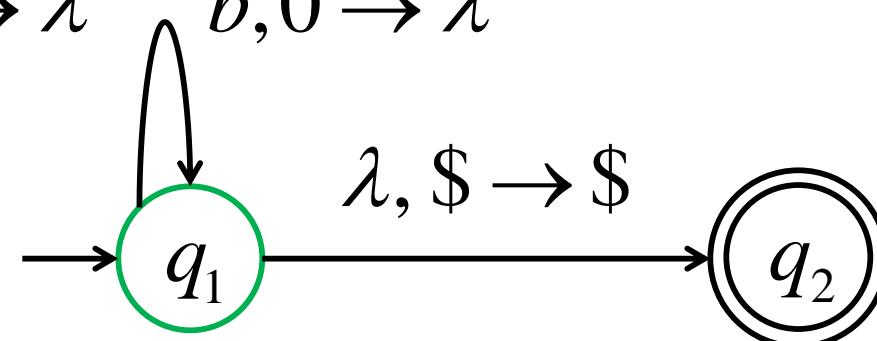
$a, \$ \rightarrow 0\$$ $b, \$ \rightarrow 1\$$

$a, 0 \rightarrow 00$ $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$ $b, 0 \rightarrow \lambda$

0
$\$$

Stack



Execution Example

Input

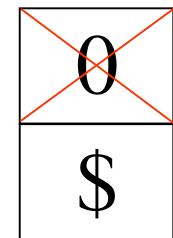
a	b	b	b	a	a
-----	-----	-----	-----	-----	-----

Time 2

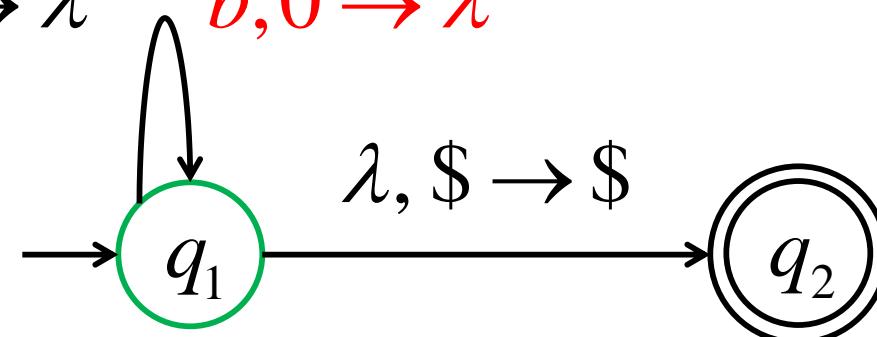
$a, \$ \rightarrow 0\$$ $b, \$ \rightarrow 1\$$

$a, 0 \rightarrow 00$ $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$ $b, 0 \rightarrow \lambda$



Stack



Execution Example

Input

a	b	b	b	a	a
-----	-----	-----	-----	-----	-----

Time 3

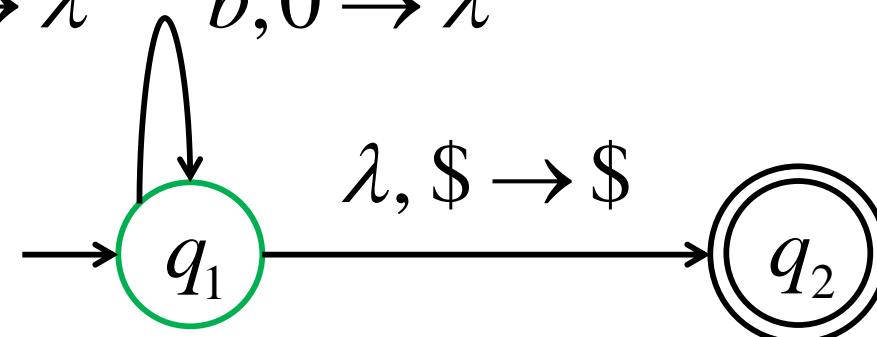
$a, \$ \rightarrow 0\$$ $b, \$ \rightarrow 1\$$

$a, 0 \rightarrow 00$ $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$ $b, 0 \rightarrow \lambda$

1
$\$$

Stack



Execution Example

Input

a	b	b	b	a	a
-----	-----	-----	-----	-----	-----

Time 4

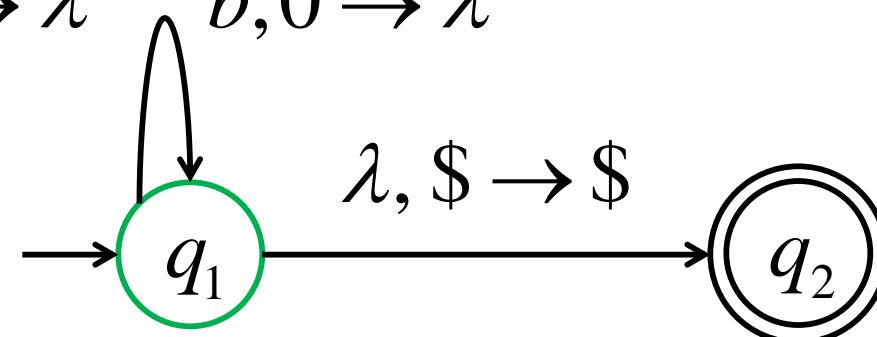
$a, \$ \rightarrow 0\$$ $b, \$ \rightarrow 1\$$

$a, 0 \rightarrow 00$ $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$ $b, 0 \rightarrow \lambda$

1
1
$\$$

Stack



Execution Example

Input

a	b	b	b	a	a
---	---	---	---	---	---

Time 5

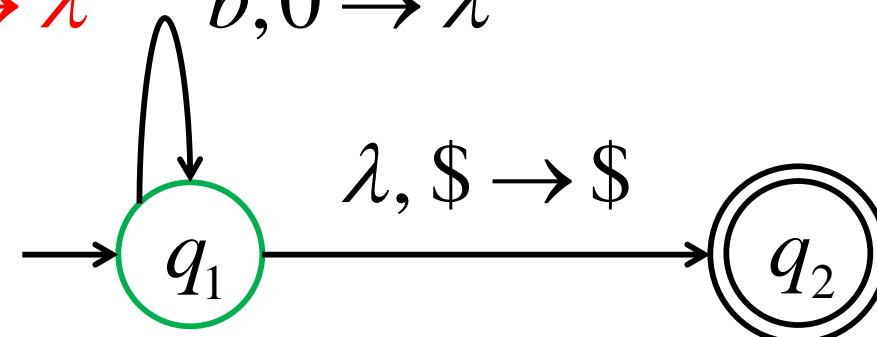
1
1
\$

$a, \$ \rightarrow 0\$$ $b, \$ \rightarrow 1\$$

$a, 0 \rightarrow 00$ $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$ $b, 0 \rightarrow \lambda$

Stack



Execution Example

Input

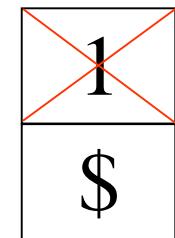
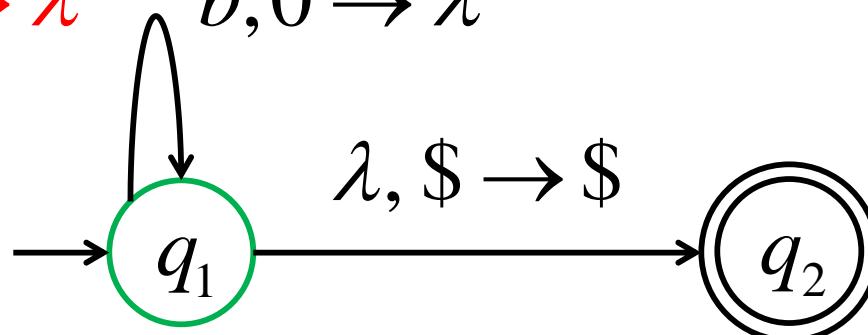
a	b	b	b	a	a
-----	-----	-----	-----	-----	-----

Time 6

$a, \$ \rightarrow 0\$$ $b, \$ \rightarrow 1\$$

$a, 0 \rightarrow 00$ $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$ $b, 0 \rightarrow \lambda$



Stack

Execution Example

Input

a	b	b	b	a	a
-----	-----	-----	-----	-----	-----

Time 7

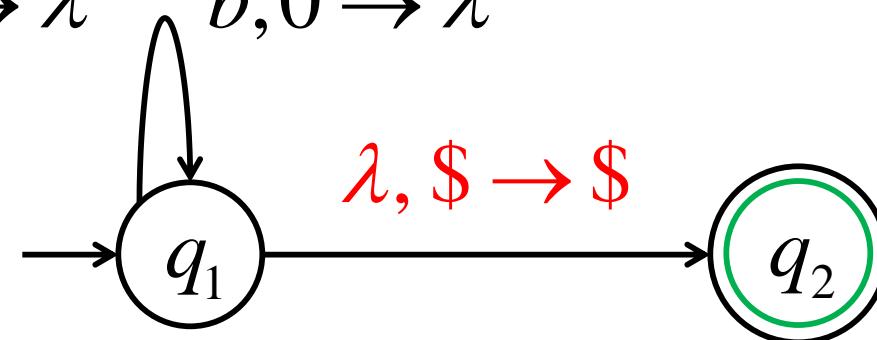
$a, \$ \rightarrow 0\$$ $b, \$ \rightarrow 1\$$

$a, 0 \rightarrow 00$ $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$ $b, 0 \rightarrow \lambda$

$\$$

Stack



Execution Example

Input

a	b	b	b	a	a
-----	-----	-----	-----	-----	-----

Time 8

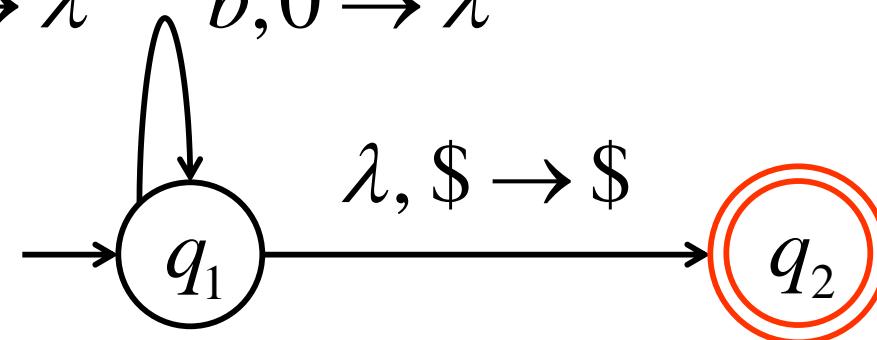
$a, \$ \rightarrow 0\$$ $b, \$ \rightarrow 1\$$

$a, 0 \rightarrow 00$ $b, 1 \rightarrow 11$

$a, 1 \rightarrow \lambda$ $b, 0 \rightarrow \lambda$

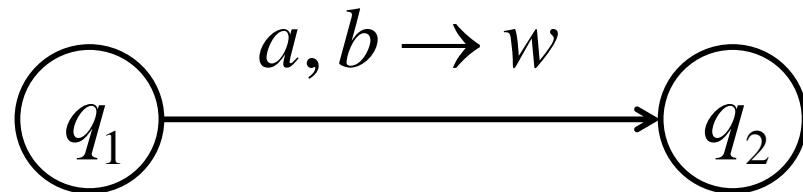
$\$$

Stack



Formalities for NPDAs

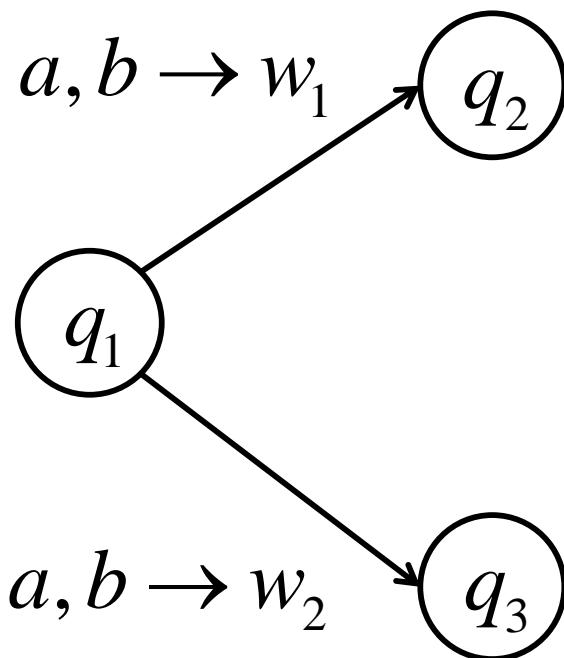
Formal Definition



- Transition function:

$$\delta(q_1, a, b) = \{(q_2, w)\}$$

Formal Definition



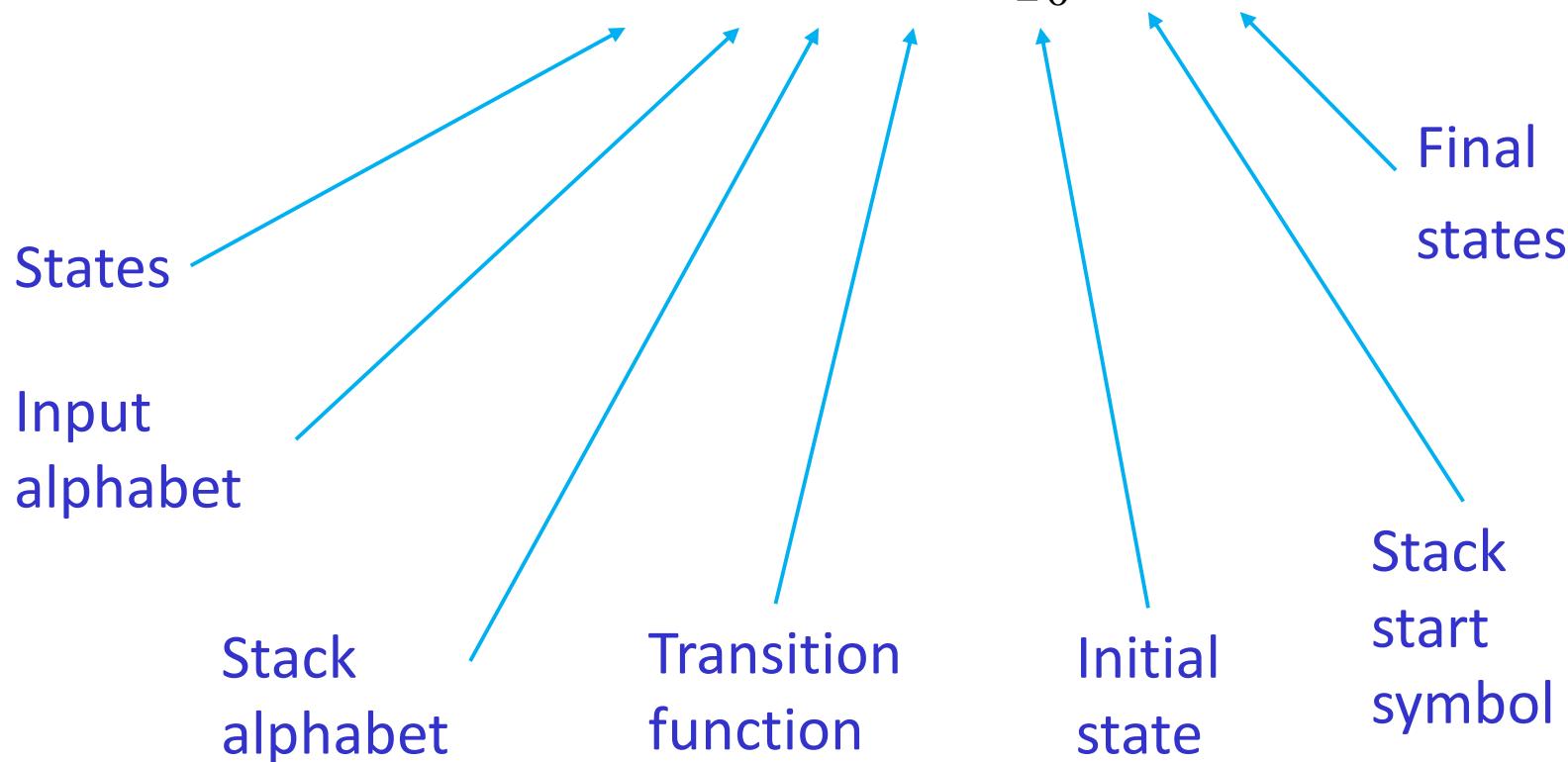
- Transition function:

$$\delta(q_1, a, b) = \{(q_2, w_1), (q_3, w_2)\}$$

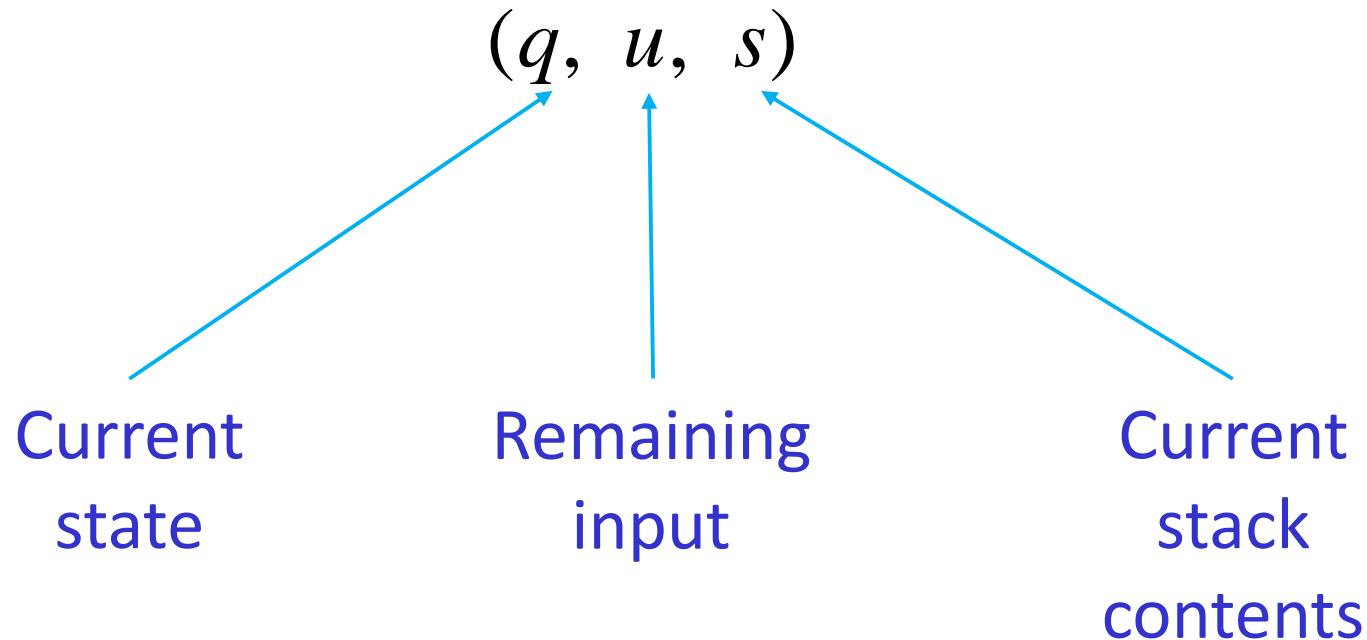
Formal Definition

Non-Deterministic Pushdown Automaton - NPDA

$$M = (Q, \Sigma, \Gamma, \delta, q_0, \$, F)$$



Instantaneous Description



Example: Instantaneous Description

Input

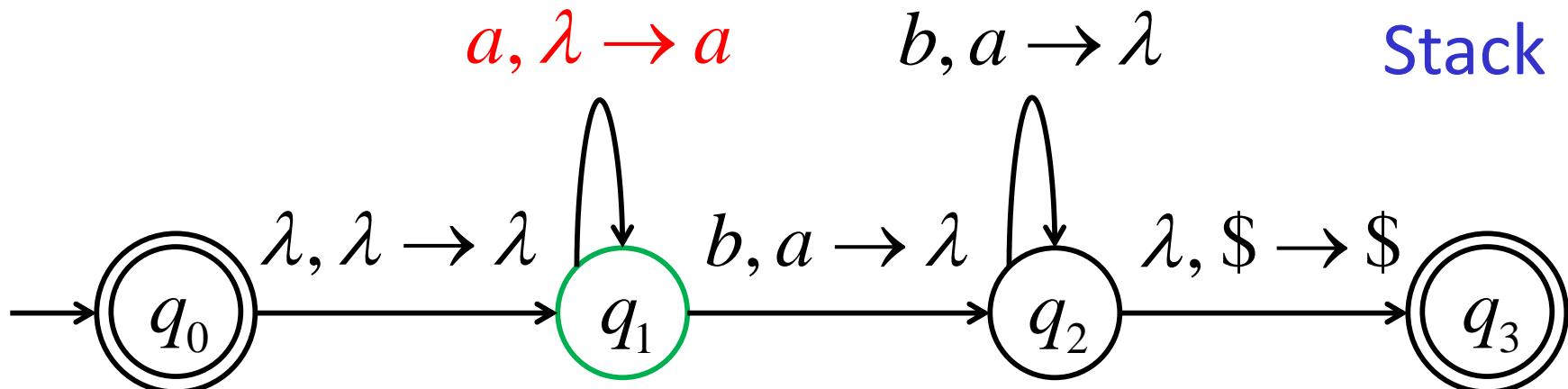
a	a	a	b	b	b
---	---	---	---	---	---



($q_1, bbb, aaa\$$)

a
a
a
\$

Stack



Example: Instantaneous Description

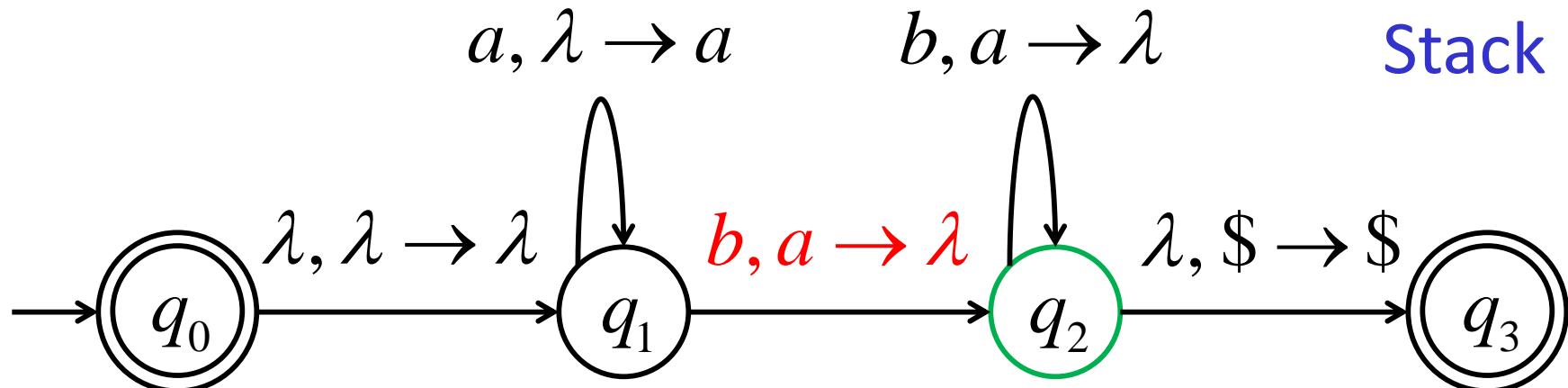
Input

a	a	a	b	b	b
---	---	---	---	---	---

($q_1, bb, aa\$$)

a
a
a
\$

Stack



Computation

We write:

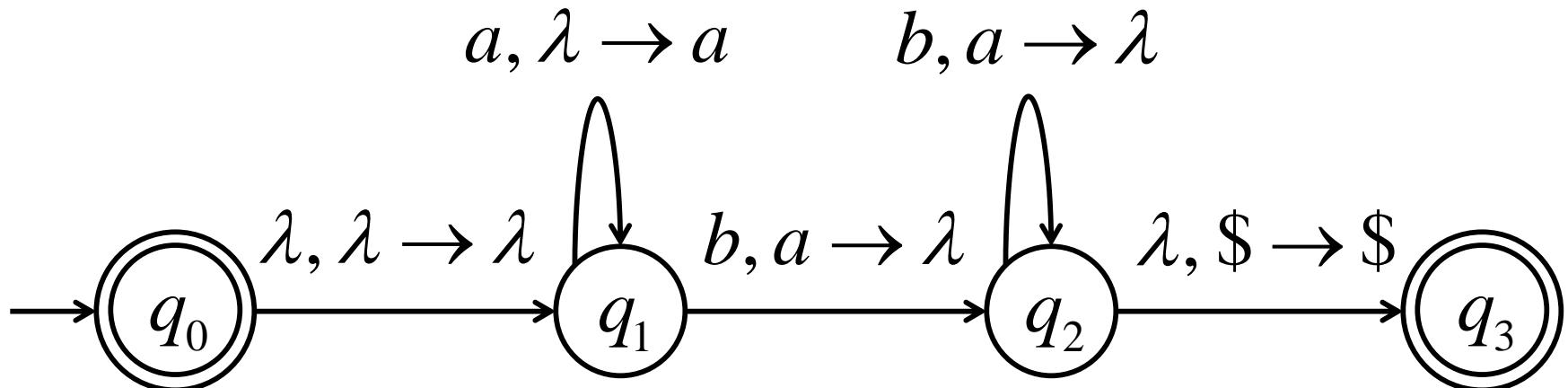
$$(q_1, bbb, aaa\$) \xrightarrow{f} (q_2, bb, aa\$)$$

Computation

$(q_0, aaabbb, \$) \xrightarrow{f} (q_1, aaabbb, \$)$

$(q_1, aabbb, a\$) \xrightarrow{f} (q_1, abbb, aa\$) \xrightarrow{f} (q_1, bbb, aaa\$)$

$(q_2, bb, aa\$) \xrightarrow{f} (q_2, b, a\$) \xrightarrow{f} (q_2, \lambda, \$) \xrightarrow{f} (q_3, \lambda, \$)$



Computation

$(q_0, aaabbb, \$) \xrightarrow{f} (q_1, aaabbb, \$)$

$(q_1, aabbb, a\$) \xrightarrow{f} (q_1, abbb, aa\$) \xrightarrow{f} (q_1, bbb, aaa\$)$

$(q_2, bb, aa\$) \xrightarrow{f} (q_2, b, a\$) \xrightarrow{f} (q_2, \lambda, \$) \xrightarrow{f} (q_3, \lambda, \$)$

For convenience we write:

$(q_0, aaabbb, \$) \xrightarrow{*} (q_3, \lambda, \$)$

Formal Definition

Language of NPDA M :

$$L(M) = \{ w : (q_0, w, s) \xrightarrow{*} (q_f, \lambda, s') \}$$

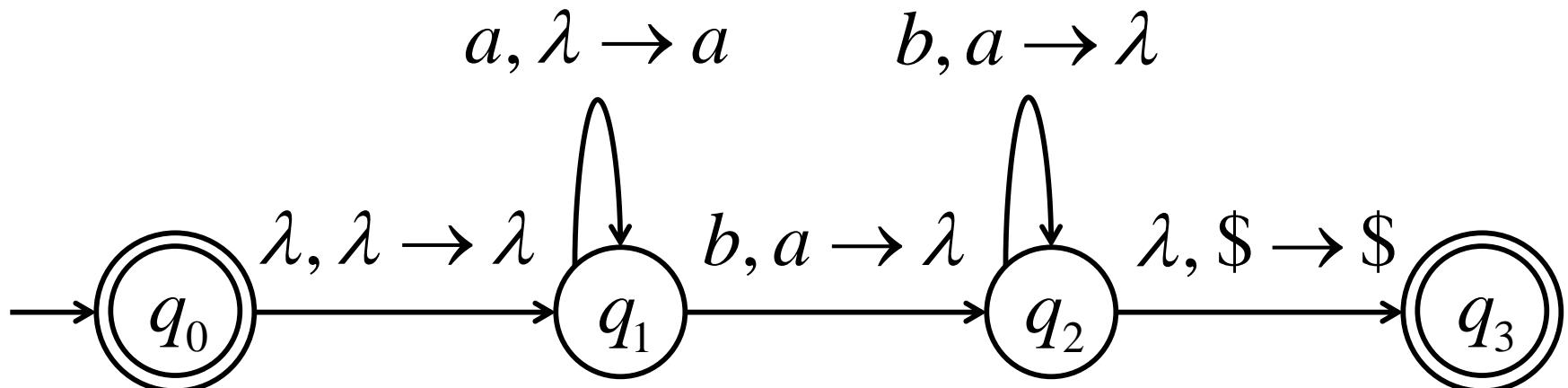
The equation defines the language $L(M)$ as the set of strings w such that there exists a sequence of transitions starting from the initial state q_0 and ending at the final state q_f , with the final configuration being (q_f, λ, s') . The symbol $\xrightarrow{*}$ indicates that the sequence of transitions can be zero or more steps long.

Initial state Final state

Example

$$(q_0, aaabbb, \$) \xrightarrow{*} (q_3, \lambda, \$)$$
$$aaabbb \in L(M)$$

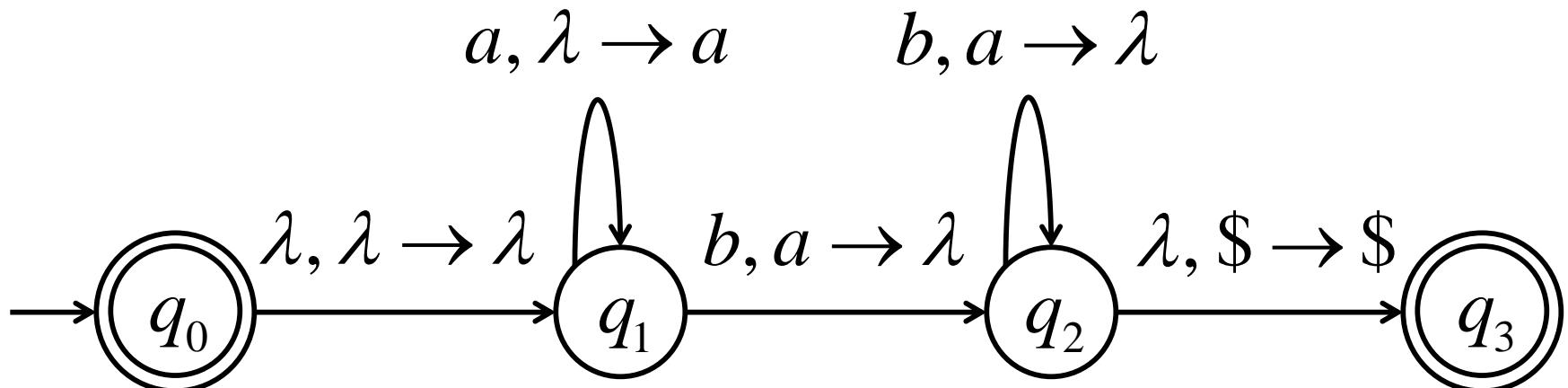
NPDA M :



Example

$$(q_0, a^n b^n, \$) \xrightarrow{*} (q_3, \lambda, \$)$$
$$a^n b^n \in L(M)$$

NPDA M :

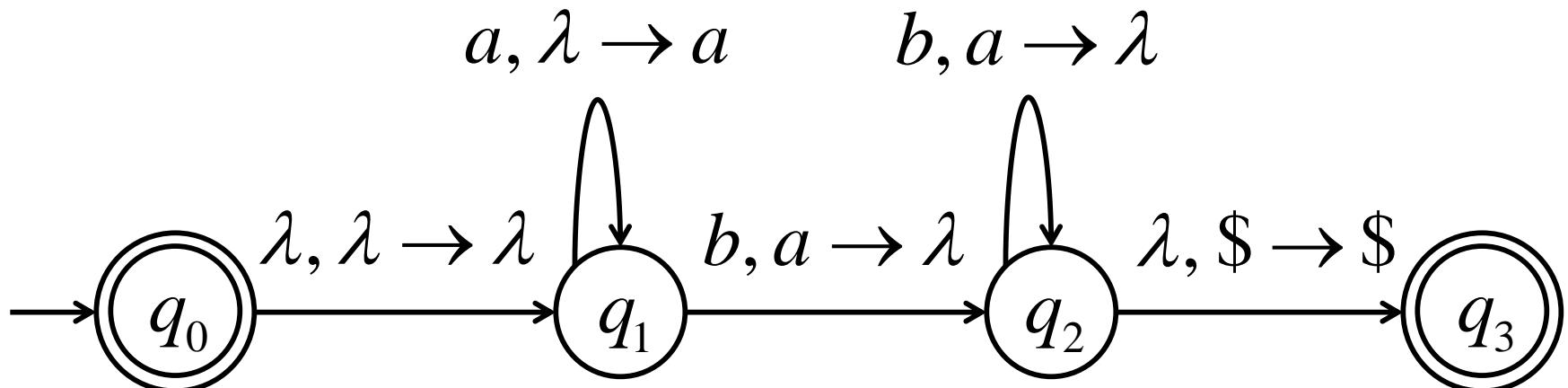


Example

Therefore:

$$L(M) = \{a^n b^n : n \geq 0\}$$

NPDA M :



Exercises

- Construct NPDAs that accept the following languages on $\Sigma = \{a, b, c\}$.
 - a. $L = \{a^n b^{2n} : n \geq 0\}$
 - b. $L = \{wcw^R : w \in \{a, b\}^*\}$
 - c. $L = \{a^n b^{n+m} c^m : n, m \geq 0\}$
 - d. $L = \{a^n b^m c^{n+m} : n, m \geq 0\}$
 - e. $L = \{a^3 b^n c^n : n \geq 0\}$

Exercises

- Construct an NPDA that accepts the language

$$L = \{w : n_a(w) = n_b(w) + 1\}$$

- Construct an NPDA that accepts the language

$$L = \{w : n_a(w) = 2n_b(w)\}$$

- Construct an NPDA that accepts the language

$$L = \{w : n_a(w) + n_b(w) = n_c(w)\}$$