280 - THE2

Student Information

Let context-free grammar for L_1 be G_1 .

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Q1

a)

I couldn't find a simpler grammar; so, I created the rules such that they are representing all possible combinations.

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G_1 = (V, \sum, R, S) where,
V = \{S, a, b\}
\sum = \{a, b\}
R = \{S \rightarrow e | Sabb | Sbab | Sbba | abbS | babS | babS | babS | babB | bSab | bSab | bbSa | abSb | baSb | bbSa \}
b)
Let context-free grammar for L_2 be G_2.
G_2 = (V, \sum, R, S) where,
V = \{S, a, b\}
\sum = \{a, b\}
R = \{S \to aaSb|aSb|e\}
c)
Let M be the PDA that accepts L_1
M = (\{p, q\}, \sum, V, \Delta, p, \{q\}) where,
\sum = \{a, b\}
V = \{S, a, b\}
\Delta = \{((p, e, e), (q, S)), ((p, e, S), (q, e)),\}
((q, e, S), (q, Sabb)), ((q, e, S), (q, Sbab)), ((q, e, S), (q, Sbba)),
((q, e, S), (q, abbS)), ((q, e, S), (q, babS)), ((q, e, S), (q, bbaS)),
((q, e, S), (q, aSbb)), ((q, e, S), (q, bSab)), ((q, e, S), (q, bSba)),
((q, e, S), (q, abSb)), ((q, e, S), (q, baSb)), ((q, e, S), (q, bbSa))
((q, a, a), (q, e)), ((q, b, b), (q, e))
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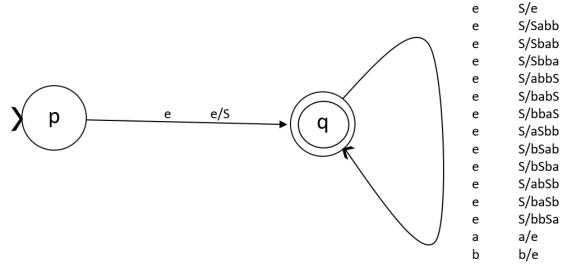


Figure 1: PDA of L_1

d)

I used the construction for union of two context-free grammars.

Changing the names of the elements for G_1 and G_2 , and they become, $G_1 = (V_1, \sum_1, R_1, S_1)$ and $G_2 = (V_2, \sum_2, R_2, S_2)$

Let the context-free grammar for L_3 be G_3 .

 $G_3 = (V_3, \sum_3, R_3, S_3)$ where, $V_3 = V_1 \cup V_2 \cup S_3 = \{S_1, S_2, S_3, a, b\}$

$$\begin{split} &\sum_3 = \{a,b\} \\ &R_3 = \{S_3 \rightarrow S_1 | S_2, \\ &S_1 \rightarrow e | S_1 abb | S_1 bab | S_1 bba | abb S_1 | bab S_1 | bas S_2 | bas S_3 | bas S_3 | bas S_4 |$$

 $S_2 \to aaS_2b|aS_2b|e$

a)

Consider the string 00111. We can create 2 different left-most derivations and the corresponding parse trees for this string as below;

1)
$$S \rightarrow AS \rightarrow A1S \rightarrow 0A11S \rightarrow 00111S \rightarrow 00111$$

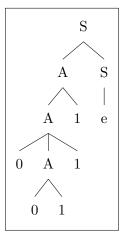


Figure 2: 1^{st} derivation's parse tree

2) $S \rightarrow AS \rightarrow 0A1S \rightarrow 0A11S \rightarrow 00111S \rightarrow 00111$

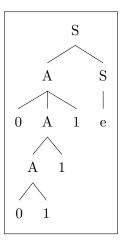


Figure 3: 2^{nd} derivation's parse tree

Since there exist 2 different parse trees for the same string, the grammar is ambiguous.

b)

Let G_2 be the unambiguous grammar for $L(G_1)$ such that;

$$G_2 = (V_2, \sum_2, R_2, S_2) \text{ where;}$$

$$V_2 = \{0, 1, S_2, B, C\}$$

$$\sum_2 = \{0, 1\}$$

$$R_2 = \{S_2 \to CS_2,$$

$$C \to 0C1|01B,$$

$$B \to 1B|e\}$$

c)

Leftmost derivation of the string 00111 is;
$$S_2 \to CS_2 \to 0C1S_2 \to 001B1S_2 \to 0011B1S_2 \to 00111S_2 \to 00111$$

Corresponding parse tree is below;

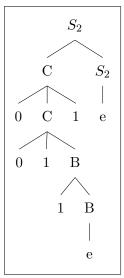


Figure 4: Parse tree of the grammar from part b.