

INTRODUCTION TO ARTIFICIAL INTELLIGENCE

LECTURE 2: UNINFORMED SEARCH

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PROBLEM SOLVING

Problem solving is an essential part of intelligence.

WOLF, GOAT AND CABBAGE PROBLEM

Once upon a time a farmer went to a market and purchased a wolf, a goat, and a cabbage. On his way home, the farmer came to the bank of a river and rented a boat. But crossing the river by boat, the farmer could carry only himself and a single one of his purchases: the wolf, the goat, or the cabbage. If left unattended together, the wolf would eat the goat, or the goat would eat the cabbage.

The farmer's challenge was to carry himself and his purchases to the far bank of the river, leaving each purchase intact. How did he do it?



SEARCH PROBLEMS FORMALLY

Formally, a **search problem** has 5 components:

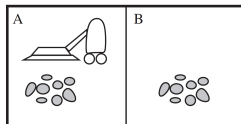
- ▶ s_0 : Initial state
- ▶ $\text{ACTIONS}(s)$: Returns the set of actions **applicable** in state s . (E.g. the action of opening a door might only be applicable if the door is unlocked.)
- ▶ $\text{RESULTS}(s, a)$: Returns the state s' reached from s by executing action a .
- ▶ $\text{GOAL-TEST}(s)$: Returns true if s is goal state, otherwise false.
- ▶ $\text{STEP-COST}(s, a)$: The cost of executing action a in s . Most often we will assume $\text{STEP-COST}(s, a) = 1$ for all s and a .

A state g is called a **goal state** if $\text{GOAL-TEST}(g) = \text{true}$.

A **solution** to a search problem is a **sequence of actions** (a **path**) from s_0 to a goal state. It is **optimal** if it has minimum sum of step costs.

A TOY PROBLEM: VACUUM WORLD

Vacuum World consists of two locations, each of which may or may not contain dirt and the vacuum is in one of the locations.



A TOY PROBLEM: VACUUM WORLD

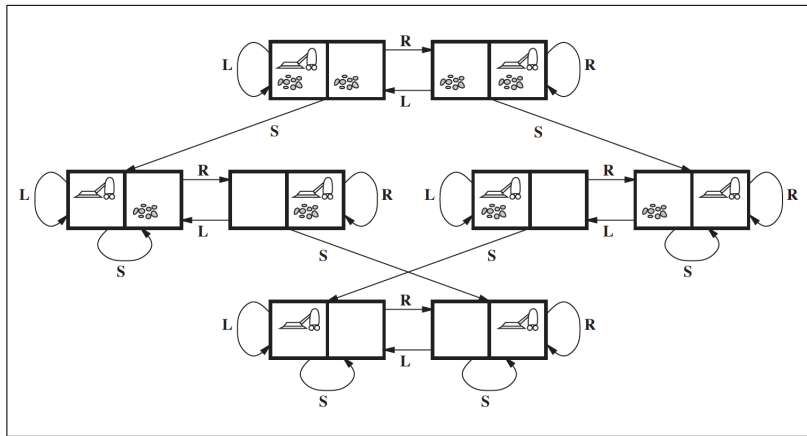
States space consists of each possible configuration (2×2^2 possible states).

- ▶ s_0 : Initial state
- ▶ $ACTIONS(s)$: for each state three possible actions: L, R, S.
- ▶ $RESULTS(s, a)$: actions have their expected results.
- ▶ $GOAL-TEST(s)$: *are all squares clean?*
- ▶ $STEP-COST(s, a)$: each step costs 1.

A TOY PROBLEM: VACUUM WORLD

TRANSITION MODEL

The transitions in the Vacuum World can be represented as a graph.



ANOTHER TOY PROBLEM: 8-PUZZLE

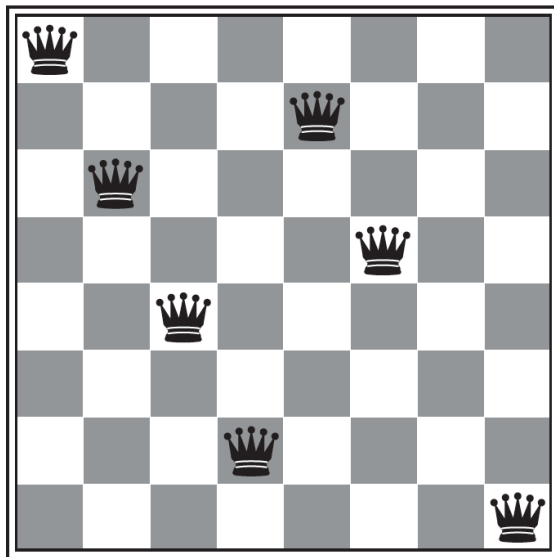
7	2	4
5		6
8	3	1

Start State

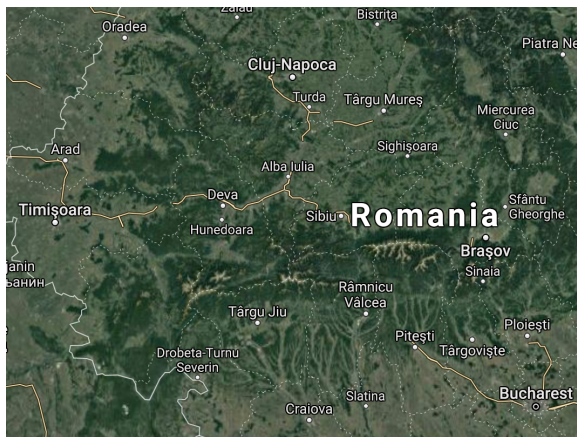
	1	2
3	4	5
6	7	8

Goal State

YET ANOTHER TOY PROBLEM: 8-QUEENS



LEVELS OF ABSTRACTION

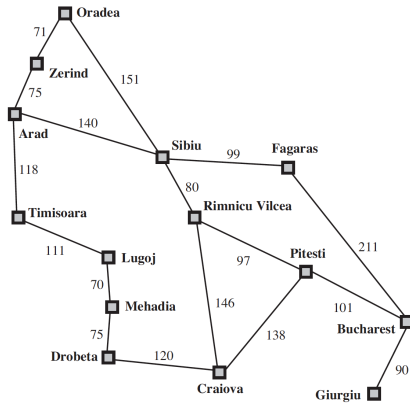


LEVELS OF ABSTRACTION



A REAL WORLD PROBLEM

ROUTE-FINDING



A REAL WORLD PROBLEM

ROUTE-FINDING

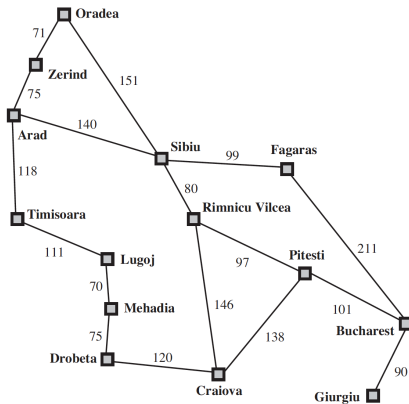
States include a location and the current time, etc.

- ▶ s_0 : specified by the user's query.
- ▶ $ACTIONS(s)$: take any flight from the current location, in any class, etc.
- ▶ $RESULTS(s, a)$: the state resulting from taking a flight.
flight destination as current location and flight arrival time as current time.
- ▶ $GOAL-TEST(s)$: *Are we at the final destination specified by the user?*
- ▶ $STEP-COST(s, a)$: monetary cost, waiting time, flight time, customs and immigration procedures, seat quality, time of day, type of airplane, frequent-flyer mileage awards, etc.

ANOTHER REAL WORLD PROBLEM

TOURING PROBLEM

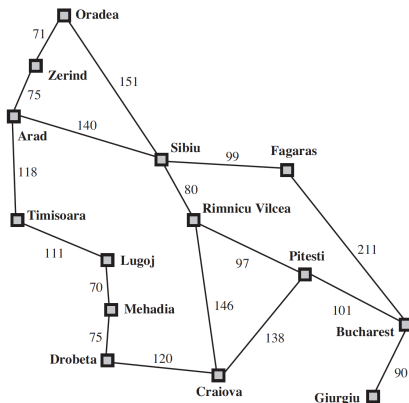
Visit every city in the figure **at least once**, starting and ending in Bucharest.



YET ANOTHER REAL WORLD PROBLEM

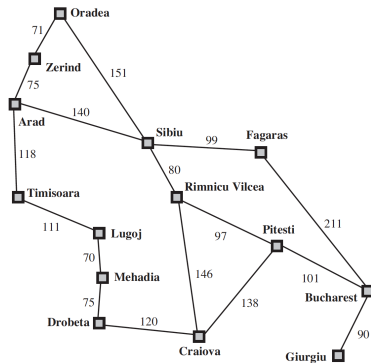
TRAVELLING SALESPERSON PROBLEM

Find a travel plan that allows visiting every city **exactly once**.



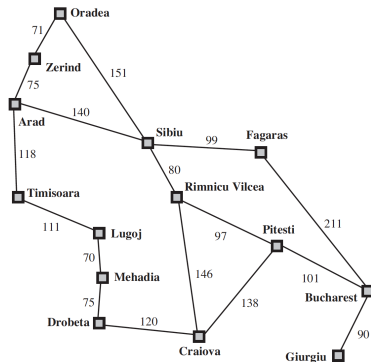
BACK TO ROUTE-FINDING PROBLEM: FROM ARAD TO BUKAREST

PROBLEM SOLVING AS TREE-SEARCH



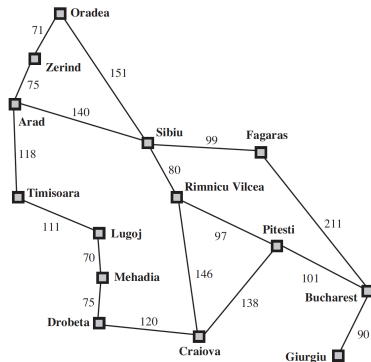
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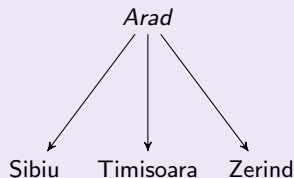
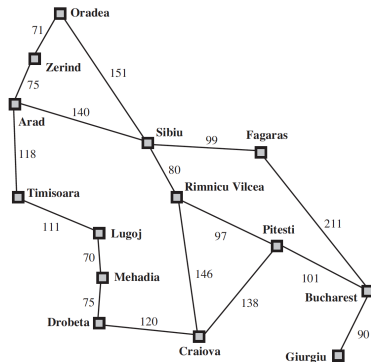
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Arad

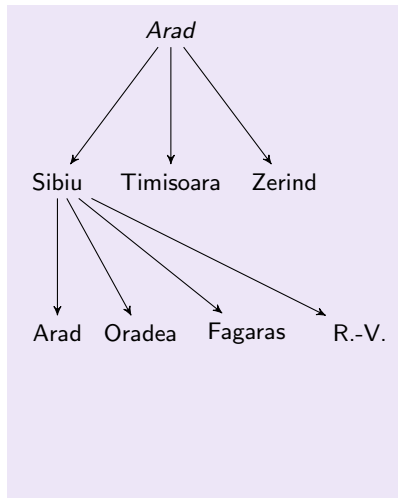
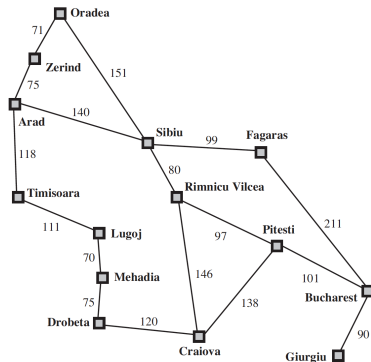
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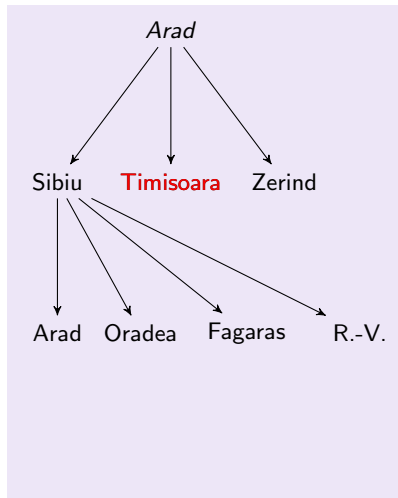
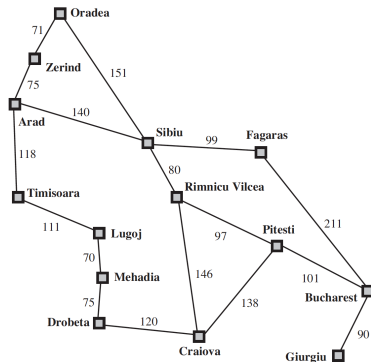
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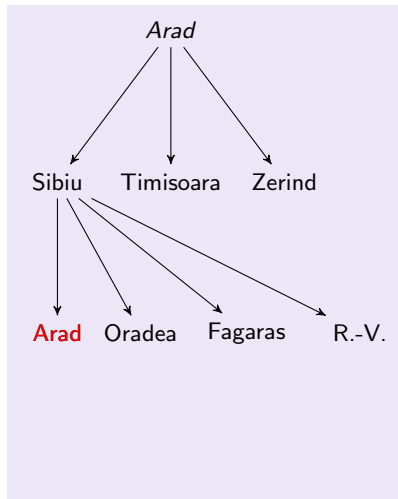
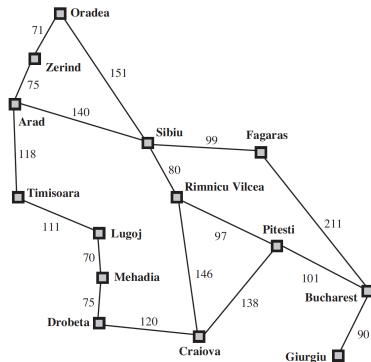
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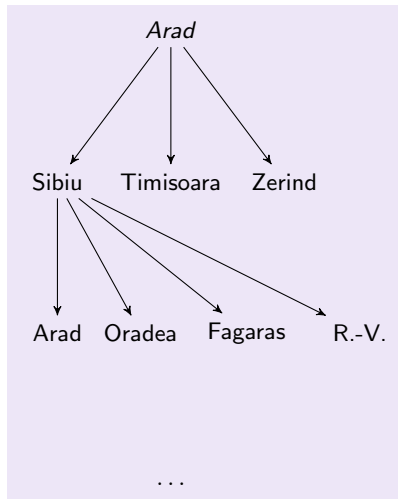
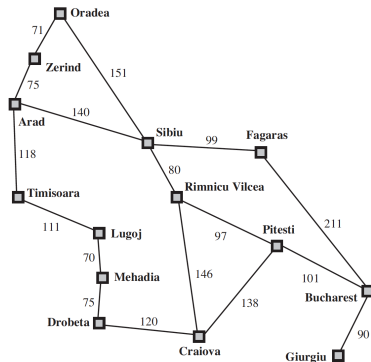
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BACK TO ROUTE-FINDING PROBLEM: FROM ARAD TO BUKAREST

PROBLEM SOLVING AS TREE-SEARCH



EXPANDED NODES & FRONTIER

In tree and graph search, we always generate **all** children of a chosen state s :
We compute $\text{RESULT}(s, a)$ for *all* applicable actions a .
We call this process **expanding** s .

In tree search states are called **nodes**: In tree search we don't keep track of repeated states, and hence distinct tree nodes might represent the same state.

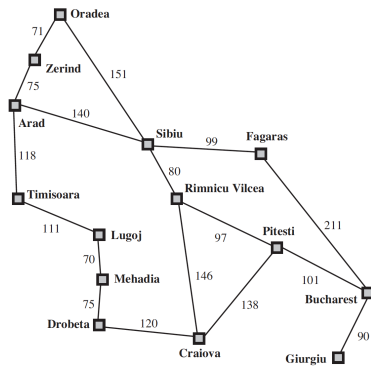
The search has two types of nodes:

- ▶ **Expanded nodes**: Nodes for which all children have been generated.
- ▶ **Frontier**: Nodes that we generated, but not expanded.

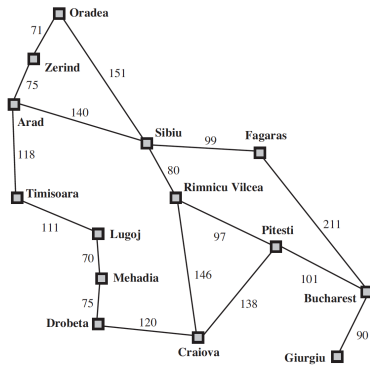
THE TREE-SEARCH ALGORITHM

```
function TREE-SEARCH (problem) returns a solution, or failure
  frontier := { $s_0$ } (initial state)    // we initialise the frontier
  loop do
    if frontier =  $\emptyset$  then return failure
    choose a node n from frontier
    remove n from frontier
    if n is a goal state then return solution
    for each child m of n    // we expand n
      add child m to frontier
```

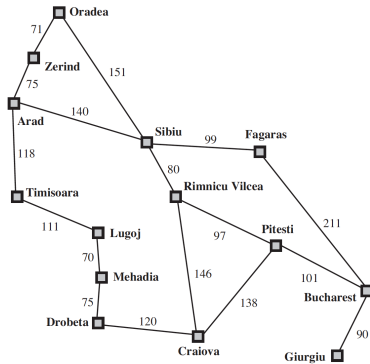
FROM TREE-SEARCH TO GRAPH-SEARCH



FROM TREE-SEARCH TO GRAPH-SEARCH

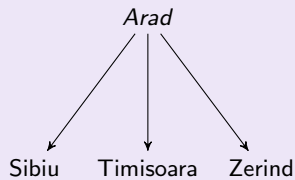
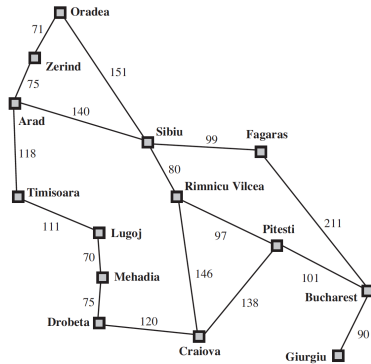


FROM TREE-SEARCH TO GRAPH-SEARCH

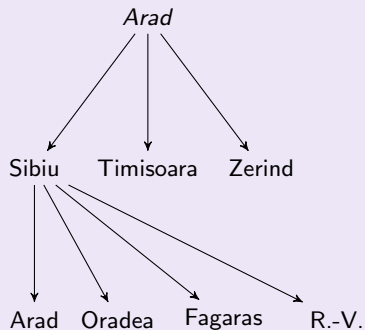
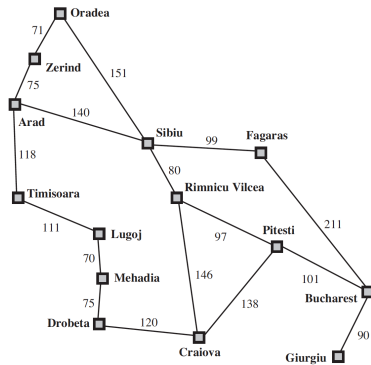


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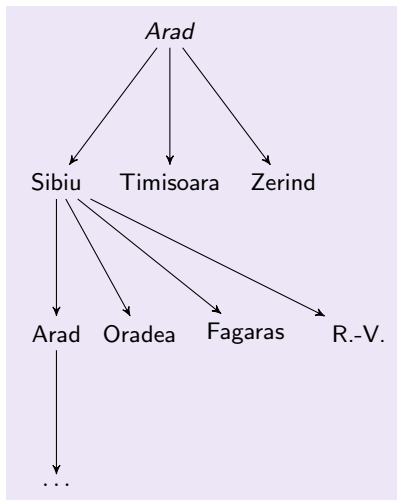
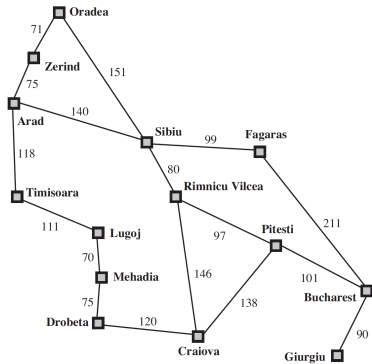
FROM TREE-SEARCH TO GRAPH-SEARCH



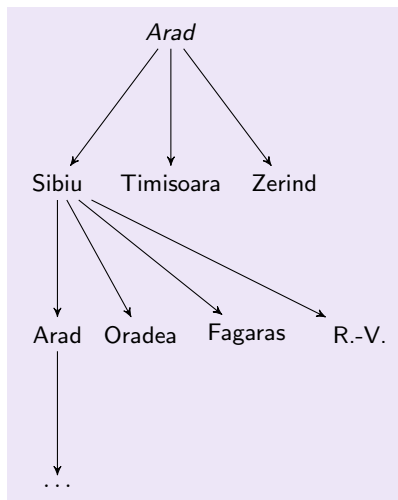
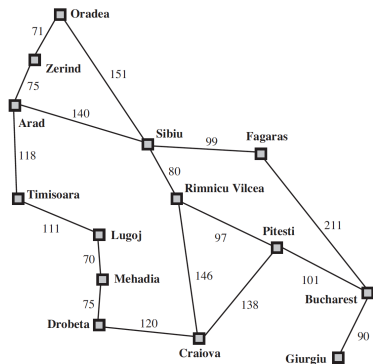
FROM TREE-SEARCH TO GRAPH-SEARCH



FROM TREE-SEARCH TO GRAPH-SEARCH



FROM TREE-SEARCH TO GRAPH-SEARCH



Failure to detect repeated states can lead to infinite loops in TREE-SEARCH.

Solution: Keep track of already generated states.

Adding this check to TREE-SEARCH gives us GRAPH-SEARCH.

THE GRAPH-SEARCH ALGORITHM

The **red lines** are those added to the TREE-SEARCH algorithm.

function GRAPH-SEARCH (problem) returns a solution, or failure

expandedNodes := \emptyset

frontier := $\{s_0\}$ (initial state) // we initialise the frontier

loop do

if *frontier* = \emptyset **then return** failure

 choose a node *n* from *frontier*

 remove *n* from *frontier*

 add *n* to *expandedNodes*

if *n* is a goal state **then return** solution

for each child *m* of *n* // we expand *n*

if *m* \notin *frontier* and *m* \notin *expandedNodes* **then**

 add child *m* to *frontier*

SOLVING SEARCH PROBLEMS

Solutions to search problems split into:

- ▶ **Tree search**
- ▶ **Graph search**

The only difference is: graph search keeps track of repeated states.

Along a different dimension they split into:

- ▶ **Uninformed search:** No sense of closeness to the goal.
- ▶ **Informed search:** Heuristics used to give a sense of closeness to the goal (e.g. straight-line distance to goal in case of route finding on a map).

For now, we only consider uninformed search.

Along yet another dimension:

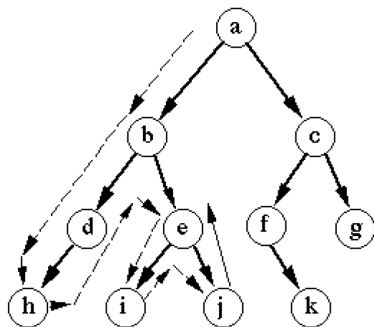
- ▶ **Breadth-based search:** Exploring shallow nodes first (close to s_0).
- ▶ **Depth-based search:** Exploring as deeply as possible, sacrificing breadth (might miss the goal).

TREE SEARCH: DFS vs BFS

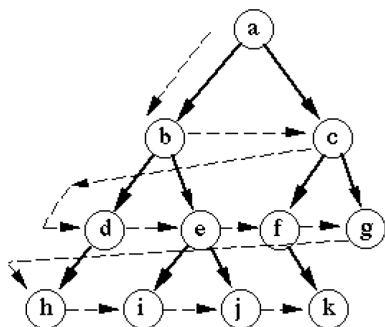
DFS: Depth-first search.

BFS: Breadth-first search.

Order of node expansion:



Depth-first search



Breadth-first search

SEARCH STRATEGIES OF TREE-SEARCH AND GRAPH-SEARCH

Different search strategies can be achieved by simply changing how **choose node from frontier** and **add child to frontier** work.

Breadth-first search (BFS):

- ▶ Frontier is **queue** (FIFO).
- ▶ **choose node from frontier**: dequeue node from frontier.
- ▶ **add child to frontier**: enqueue node to frontier.

Depth-first search (DFS):

- ▶ Frontier is **stack** (LIFO).
- ▶ **choose node from frontier**: pop node from frontier.
- ▶ **add child to frontier**: push node to frontier.

EVALUATION OF SEARCH ALGORITHMS

Which search algorithm would you rather choose: BFS or DFS?

Go to www.menti.com and use the code 72 26 24

EVALUATION OF SEARCH ALGORITHMS

We can measure problem-solving performance by:

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EVALUATION OF SEARCH ALGORITHMS

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Optimal solution has the lowest path cost among solutions.

EVALUATION OF SEARCH ALGORITHMS

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Optimal solution has the lowest path cost among solutions.
- ▶ **Time complexity** How long does it take to find a solution?

EVALUATION OF SEARCH ALGORITHMS

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- ▶ **Completeness** Is the algorithm guaranteed to find a solution if one exists?
- ▶ **Optimality** Does the strategy find the optimal solution?
Optimal solution has the lowest path cost among solutions.
- ▶ **Time complexity** How long does it take to find a solution?
- ▶ **Space complexity** How much memory is needed to perform the search?

EVALUATION OF SEARCH ALGORITHMS

FOR TREE-SEARCH

Criterion	Breadth-First	Uniform-Cost	Depth-First	Depth-Limited	Iterative Deepening	Bidirectional
Complete?	Yes ^a	Yes ^{a,b}	No	No	Yes ^a	Yes ^{a,d}
Time	$O(b^d)$	$O(b^{1+\lceil C^*/\epsilon \rceil})$	$O(b^m)$	$O(b^\ell)$	$O(b^d)$	$O(b^{d/2})$
Space	$O(b^d)$	$O(b^{1+\lceil C^*/\epsilon \rceil})$	$O(bm)$	$O(b\ell)$	$O(bd)$	$O(b^{d/2})$
Optimal?	Yes ^c	Yes	No	No	Yes ^c	Yes ^{c,d}

Figure 3.21 Evaluation of tree-search strategies. b is the branching factor; d is the depth of the shallowest solution; m is the maximum depth of the search tree; ℓ is the depth limit. Superscript caveats are as follows: ^a complete if b is finite; ^b complete if step costs $\geq \epsilon$ for positive ϵ ; ^c optimal if step costs are all identical; ^d if both directions use breadth-first search.

THE END OF LECTURE 2