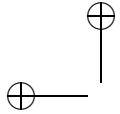
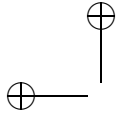


Appendices to Dense Sphere Packings

a blueprint for formal proofs

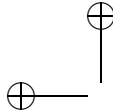
Thomas C. Hales



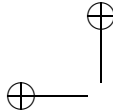


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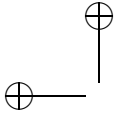


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**Abstract**

These are unpublished appendices for "Dense Sphere Packings" (DSP). These are notes that were useful in the formalization of some lemmas in the proof of the Kepler conjecture but were considered too technical for inclusion in DSP. It also includes material related to some changes that were made to the formalization that occurred after the publication of DSP.

These appendices have not been edited into a form that is suitable for publication. They were primarily written for the benefit of the members of the Flyspeck team who carried out the formalization.



1

Appendix on the main estimate

This appendix gives further details about the proof of the main estimate. It contains numerous improvements over the published text of *Dense Sphere Packings*. The verification of the main estimate can be viewed as a combinatorial proof that $S_{\text{init}} \Rightarrow S_{\text{term}}$.

1.1 statement of results

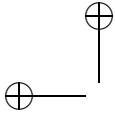
This appendix largely replaces [1][Section 7.4]. For completeness, we repeat a few definitions and results.

Definition 1.1 (h_0, τ) [CUFCNHB] [$\rho_0 \leftrightarrow \text{rho_fun}$] [$\rho_0 \leftrightarrow \text{rho}$] [$\tau \leftrightarrow \text{tau_fun}$] [$\text{sol}_0 \leftrightarrow \text{sol0}$] [$h_0 \leftrightarrow \text{h0}$] [$\text{azim} \leftrightarrow \text{azim_in_fan}$] Let (V, E, F) be a nonreflexive local fan. Recall that $h_0 = 1.26$ and $L(h) = (h_0 - h)/(h_0 - 1)$, when $h \leq h_0$. Set

$$\rho_0(y) = 1 + \frac{\text{sol}_0}{\pi} \cdot \frac{y - 2}{2h_0 - 2} = 1 + \frac{\text{sol}_0}{\pi}(1 - L(y/2))$$

$$\begin{aligned} \tau(V, E, F) &= \sum_{x \in F} \rho_0(\|\text{node}(x)\|) \text{azim}(x) + (\pi + \text{sol}_0)(2 - k(F)) \\ &= \sum_{v \in V} \rho_0(\|\mathbf{v}\|) \angle(\mathbf{v}) + (\pi + \text{sol}_0)(2 - \text{card}(V)), \end{aligned}$$

where $\text{sol}_0 = 3 \arccos(1/3) - \pi \approx 0.551$ is the solid angle of a spherical equilateral triangle of side $\pi/3$, and $k(F)$ is the cardinality of F .



Definition 1.2 ($\tau_{tri}, \tau_3, , \text{dih}_i$) [$\tau_{tri} \leftrightarrow \text{taum}$] We define additional functions in the case $k = 3$. If $\text{card}(V) = 3$, then write $V = \{\mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3\}$. Set

$$\tau_3(V) = \sum_{i=1}^3 \rho_0(\|\mathbf{v}_i\|) \text{dih}_V\{\mathbf{0}, \mathbf{v}_i\}\{\mathbf{v}_{i+1}, \mathbf{v}_{i+2}\} - (\pi + \text{sol}_0).$$

Let

$$\tau_{tri}(y_1, y_2, y_3, y_4, y_5, y_6) = \sum_{i=1}^3 \rho_0(y_i) \text{dih}_i(y_1, \dots, y_6) - (\pi + \text{sol}_0), \quad (1.3)$$

where

$$\begin{aligned} \text{dih}_1(y_1, y_2, y_3, y_4, y_5, y_6) &= \text{dih}(y_1, y_2, y_3, y_4, y_5, y_6), \\ \text{dih}_2(y_1, y_2, y_3, y_4, y_5, y_6) &= \text{dih}(y_2, y_3, y_1, y_5, y_6, y_4), \text{ and} \\ \text{dih}_3(y_1, y_2, y_3, y_4, y_5, y_6) &= \text{dih}(y_3, y_1, y_2, y_6, y_4, y_5). \end{aligned} \quad (1.4)$$

Definition 1.5 (standard, protracted, diagonal) [KRACSCQ] [standard \leftrightarrow standard (deprecated)] [protracted \leftrightarrow protracted (deprecated)] [diagonal \leftrightarrow diagonal0 (deprecated)] Let (V, E) be a fan. We write $\|\varepsilon\|$ for $\|\mathbf{v} - \mathbf{w}\|$, when $\varepsilon = \{\mathbf{v}, \mathbf{w}\} \subset V$. We say that ε is *standard* if

$$2 \leq \|\varepsilon\| \leq 2h_0.$$

We say that ε is *protracted* if

$$2h_0 \leq \|\varepsilon\| \leq \sqrt{8}.$$

If $\mathbf{v}, \mathbf{w} \in V$ are distinct, and $\varepsilon = \{\mathbf{v}, \mathbf{w}\}$ is not an edge in E , then we call ε a *diagonal* of the fan.

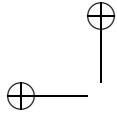
Theorem 1.6 (main estimate) [JEJTVGB] [(ANNULUS) \leftrightarrow ball_annulus] [main estimate \leftrightarrow main_estimate (deprecated)] [(DIAGONAL) \leftrightarrow diagonal1 (deprecated)] [$*$ \leftrightarrow check:quad_std_cs, pent_std_cs, hex_std_cs, pent_diag_cs, pent_pro_cs,] Let (V, E, F) be a nonreflexive local fan ([1][Definition 7.2]). We make the following additional assumptions on (V, E, F) .

1. (PACKING) V is a packing.
2. (ANNULUS) $V \subset \mathcal{B}$.
3. (DIAGONAL) If $\{\mathbf{v}, \mathbf{w}\}$ is a diagonal, then

$$\|\mathbf{v} - \mathbf{w}\| \geq 2h_0.$$

4. (CARD) Let $k = \text{card}(E) = \text{card}(F)$. Then $3 \leq k \leq 6$.

In this context, we have the following conclusions.



1. Assume $k \geq 3$. If every edge of E is standard, then

$$\tau(V, E, F) \geq \begin{cases} 0.0, & \text{if } k = 3, \\ 0.206, & \text{if } k = 4, \\ 0.4819, & \text{if } k = 5, \\ 0.712, & \text{if } k = 6. \end{cases}$$

2. Assume $k = 5$. Assume that every edge of E is standard. Assume that every diagonal ε of the fan satisfies $\|\varepsilon\| \geq \sqrt{8}$. Then

$$\tau(V, E, F) \geq 0.616. \quad (\text{check:pent_diag_cs})$$

3. Assume $k = 5$. Assume there exists some protracted edge in E and that the other four edges are standard. Then

$$\tau(V, E, F) \geq 0.616. \quad (\text{check:pent_pro_cs})$$

4. Assume that $k = 4$. Assume that there exists some protracted edge in E and that the other three edges are standard. Assume that both diagonals ε of the fan satisfy $\|\varepsilon\| \geq \sqrt{8}$. Then

$$\tau(V, E, F) \geq 0.477. \quad (\text{check:quad_pro_cs})$$

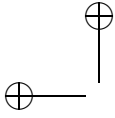
5. Assume $k = 4$. Assume that every edge of E is standard. Assume that both diagonals ε of the fan satisfy $\|\varepsilon\| \geq 3$. Then

$$\tau(V, E, F) \geq 0.467. \quad (\text{check:quad_diag_cs})$$

1.2 definitions

Definition 1.7 (torsor, adjacent) [XCZLSVS] [torsor \leftrightarrow torsor] Let $k > 1$ be an integer. A *torsor* is a set I with a given simply transitive action of $\mathbb{Z}/k\mathbb{Z}$ on I . We write the application of $j \in \mathbb{Z}/k\mathbb{Z}$ to $i \in I$ as $j+i$ or $i+j$. We also write $j+i$ for the application of the image of $j \in \mathbb{Z}$ in $\mathbb{Z}/k\mathbb{Z}$ to $i \in I$. Note that each choice of base point $i_0 \in I$ gives a bijection $i \mapsto i+i_0$ between $\mathbb{Z}/k\mathbb{Z}$ and I . We say that i and j are *adjacent* if $i = j+1$ or $j = i+1$. When i and j are adjacent, we call $\{i, j\}$ an edge of I . A *diagonal* is a pair $\{i, j\} \subset I$ such that $\{i, j\}$ is not a singleton and not an edge.

Definition 1.8 [TEQQCLX] The *opposite* I' of a torsor I is the torsor with the same underlying set and the action is composed with the group automorphism $\mathbb{Z}/k\mathbb{Z} \rightarrow \mathbb{Z}/k\mathbb{Z}$, sending $i \mapsto -i$. An *isomorphism of torsors* is a bijection that respects the action. Two torsors are *equivalent* if they are isomorphic or if one is isomorphic to the opposite of the other.



Example 1.9 If $H = (D, e, n, f)$ is a hypermap with face F , then F is a torsor under the action $x \mapsto fx$. If H is isomorphic to Dih_{2k} and has vertex set V , then V is a torsor under the action of $x \mapsto \rho x$, for some ρ similar to $\rho_{(V,E,F)}$. We may mostly restrict our attention to these instances of torsors.

Recall that $c_{\text{stab}} = 3.01$.

Definition 1.10 (stable constraint system) [ZGFHNKX] [constraint system \leftrightarrow constraint_system] [stable constraint system \leftrightarrow stable_system] $[(k, a, b, d, I, J, +1) \leftrightarrow (k_{\text{sy}}, a_{\text{sy}}, b_{\text{sy}}, d_{\text{sy}}, I_{\text{sy}}, J_{\text{sy}}, f_{\text{sy}})]$ [tri-stable \leftrightarrow tri_stable] $[(a, b, d, k, I, J, +1) \leftrightarrow (a_{\text{ts}}, b_{\text{ts}}, d_{\text{ts}}, k_{\text{ts}}, I_{\text{ts}}, J_{\text{ts}}, f_{\text{ts}})]$ [substandard (stable) \leftrightarrow augmented_constraint_system1] [substandard (tri-stable) \leftrightarrow augmented_constraint_system2] A *stable constraint system* s (or SCS for short) consists of the following data subject to conditions listed below.

1. a natural number $k \in \{3, 4, 5, 6\}$,
2. a $\mathbb{Z}/k\mathbb{Z}$ -torsor I ,
3. a real number d ,
4. real constants $a_{ij}, b_{ij}, \alpha_{ij}, \beta_{ij}$, for $i, j \in I$.
5. a set of edges $J \subset I \setminus \{i, 1+i\} : i \in I$,
6. subsets $I_{lo}, I_{hi}, I_{str} \subset I$.

The data is subject to the following conditions.

1. $d < 0.9$,
2. $a_{ij} = a_{ji}, b_{ij} = b_{ji}, \alpha_{ij} = \alpha_{ji}, \beta_{ij} = \beta_{ji}$, for $i, j \in I$.
3. $a_{ij} \leq \alpha_{ij} \leq \beta_{ij} \leq b_{ij}$, for $i, j \in I$,
- 4.

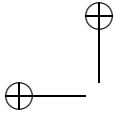
$$0 = a_{ii} \text{ and } 2 \leq a_{ij} \text{ for all } i, j \in I \text{ such that } i \neq j.$$

5. Also,

$$\begin{cases} b_{i,i+1} < 4, & \text{if } k = 3 \\ b_{i,i+1} \leq c_{\text{stab}}, & \text{if } k > 3. \end{cases}$$

6. If $\{i, j\} \in J$, then $[a_{ij}, b_{ij}] = [\sqrt{8}, c_{\text{stab}}]$.
7. $m+k \leq 6$, where m is the number of edges $\{i, i+1\} \subset I$ such that $b_{i,i+1} > 2h_0$ or $a_{i,i+1} > 2$. In particular, $\text{card}(J) + k \leq 6$.

For each SCS s , we write $k(s), d(s), I(s), a_{ij}(s)$, and so forth for the associated parameters. We simply write k, d , and so forth when there is a single SCS s in a given context.



Definition 1.11 (unadorned, basic) [SDJTENL] We say that a SCS s is *unadorned* if the following additional properties hold (with established notation):

1. For all $i, j \in I(s)$, $a_{ij}(s) = \alpha_{ij}(s)$ and $b_{ij}(s) = \beta_{ij}(s)$.
2. $I_{lo}(s) = I_{hi} = I_{str}(s) = \emptyset$.

We say the SCS s is *basic* if it is unadorned and $J(s) = \emptyset$.

Intuitively, we think of the SCS as involving both *hard* and *soft* constraints. The soft constraints are those that involve the adornments $\alpha_{ij}, \beta_{ij}, I_{lo}, I_{hi}, I_{str}$.

Example 1.12 We may always transform a SCS s into another s' that is unadorned by setting $\alpha_{ij}(s') = \alpha_{ij}(s)$, $\beta_{ij}(s') = \beta_{ij}(s)$, $I_{lo}(s') = I_{hi}(s') = I_{str}(s') = \emptyset$, and keeping the rest of the data the same.

Definition 1.13 (ear) [HVFQIBQ] [$a_{ij} \leftrightarrow \mathbf{a_ear0}$] [$b_{ij} \leftrightarrow \mathbf{b_ear0}$] [$\mathbf{ear} \leftrightarrow \mathbf{ear_sy}$] We have an unadorned SCS s given by $k = 3$, $d = 0.11$, J a singleton, and

$$[a_{ij}, b_{ij}] = \begin{cases} [0, 0], & \text{if } i = j, \\ [\sqrt{8}, c_{\text{stab}}], & \text{if } \{i, j\} \in J, \\ [2, 2h_0], & \text{otherwise.} \end{cases}$$

We call s an *ear* (by analogy with an ear in a triangulation of a polygon, which is a triangle that has two of its edges in common with the polygon).

Next we associate a set \mathcal{B}_s with each SCS s .

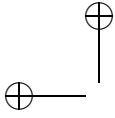
Definition 1.14 (\mathcal{B}_s) [KTFVGXF] For every SCS s , and every function $\mathbf{v} : I(s) \rightarrow \mathbb{R}^3$, let $V_{\mathbf{v}} \subset \mathbb{R}^3$ be the image of \mathbf{v} . Let $E_{\mathbf{v}}$ be the image of $i \mapsto \{\mathbf{v}_i, \mathbf{v}_{i+1}\}$. Let $F_{\mathbf{v}}$ be the image of $i \mapsto (\mathbf{v}_i, \mathbf{v}_{i+1})$. Let \mathcal{B}_s be the set of all functions \mathbf{v} that have the following properties.

1. $V_{\mathbf{v}} \subset \mathcal{B}$.
2. $a_{ij}(s) \leq \|\mathbf{v}_i - \mathbf{v}_j\| \leq b_{ij}(s)$, for all $i, j \in I(s)$.
3. if $k(s) > 3$, then $(V_{\mathbf{v}}, E_{\mathbf{v}}, F_{\mathbf{v}})$ is a nonreflexive local fan.

Note that \mathcal{B}_s does not depend on the data $\alpha_{ij}, \beta_{ij}, J, I_{str}, I_{lo}, I_{hi}$, and d . The set J is used to make a small correction $d(s, \mathbf{v})$ to the constant $d(s)$.

Definition 1.15 ($d(s, \mathbf{v})$) [TPLCZFL] [$\sigma \leftrightarrow \mathbf{sigma_sy}$] Let s be a SCS. Set $\sigma(s) = 1$ when s is an ear; $\sigma = -1$, otherwise. Write

$$d(s, \mathbf{v}) = d(s) + 0.1 \sigma(s) \sum_{\{i, j\} \in J(s)} (c_{\text{stab}} - \|\mathbf{v}_i - \mathbf{v}_j\|). \quad (1.16)$$



This correction to $d(s)$ makes it a bit easier to prove inequalities when $\sigma(s) = -1$, at the cost of slightly more difficult inequalities for ears.

When $k(s) = 3$ and $\mathbf{v} \in \mathcal{B}_s$, the set $V_{\mathbf{v}} = \{\mathbf{v}_i : i \in I(s)\}$ may degenerate to planar configurations, because the local fan constraint in Definition 1.14 is not imposed in this case. Nevertheless, by the constraint $b_{ij}(s) < 4$, the tetrahedron $\{\mathbf{0}, \mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3\}$ has well-defined dihedral angles $\text{dih}_V\{\mathbf{0}, \mathbf{v}_i\}\{\mathbf{v}_{i+1}, \mathbf{v}_{i+2}\}$, so that $\tau_3(V_{\mathbf{v}})$ is defined.

Definition 1.17 (τ^*) [BGCEUKP] Let s be a SCS. Define

$$\tau^* : \{(s, \mathbf{v}) : \mathbf{v} \in \mathcal{B}_s\} \rightarrow \mathbb{R}$$

by

$$\tau^*(s, \mathbf{v}) = \begin{cases} \tau(V_{\mathbf{v}}, E_{\mathbf{v}}, F_{\mathbf{v}}) - d(s, \mathbf{v}), & \text{if } k(s) > 3 \\ \tau_3(V_{\mathbf{v}}) - d(s, \mathbf{v}), & \text{if } k(s) = 3. \end{cases}$$

Definition 1.18 (S_{init}) [XOSFOMP] The constants in the conclusions of the main estimate (Theorem 1.6) can be packaged into unadorned SCSs. For example, the standard main estimate for $k = 6$ gives the SCS data: $d = 0.712$, $J = \emptyset$, I an indexing set of cardinality six, and

$$a_{ij} = \begin{cases} 0, & i = j, \\ 2, & j = i \pm 1, \\ 2h_0, & \text{otherwise,} \end{cases} \quad b_{ij} = \begin{cases} 0, & i = j, \\ 2h_0, & j = i \pm 1, \\ 4h_0^+, & \text{otherwise,} \end{cases}$$

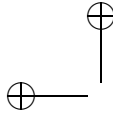
where h_0^+ is any constant greater than h_0 . The upper bound $4h_0$ on any diagonal comes from the triangle inequality: $\|\mathbf{v}_i - \mathbf{v}_j\| \leq \|\mathbf{v}_i\| + \|\mathbf{v}_j\| \leq 4h_0$. We write S_{init} for the set of SCSs s , for all cases of the main estimate.

Lemma 1.19 [ZITHLQN] [formal proof by Hoang Le Truong]. *The main estimate holds if and only if for every $s \in S_{\text{init}}$ and for every $\mathbf{v} \in \mathcal{B}_s$, we have $\tau^*(s, \mathbf{v}) \geq 0$.*

Proof This follows by expanding the definition of S_{init} and \mathcal{B}_s . Note that the set $J(s)$ is empty for $s \in S_{\text{init}}$, so

$$d(s, \mathbf{v}) = d(s), \text{ for all } s \in S_{\text{init}}.$$

□



1.3 minimization

Definition 1.20 (index, ι , \mathcal{M}_s) [FNUEPJW] Let s be a SCS. Let

$$\mathcal{M}_s \subset \mathcal{B}_s'' \subset \mathcal{B}_s' \subset \mathcal{B}_s$$

be defined as follows. Let \mathcal{B}_s' be the set of all $\mathbf{v} \in \mathcal{B}_s$ such that

1. $\tau^*(s, \mathbf{v})$ is equal to the minimum of $\tau^*(s, *)$ over \mathcal{B}_s .
2. $\tau^*(s, \mathbf{v}) < 0$.

Define the *index* $\iota(s, \mathbf{v})$ of $\mathbf{v} \in \mathcal{B}_s$ to be the number of edges $\{i, j\}$ of $I(s)$ for which $a_{ij}(s) = \|\mathbf{v}_i - \mathbf{v}_j\|$. Let $\iota(s)$ be the minimum of the index of \mathbf{v} as \mathbf{v} runs over \mathcal{B}_s' . We let \mathcal{B}_s'' be the set of $\mathbf{v} \in \mathcal{B}_s'$ that attain the smallest possible index $\iota(s)$, and let $\mathcal{M}_s \subset \mathcal{B}_s''$ be the subset of all \mathbf{v} satisfying the additional soft constraints.

1. If $i \in I_{str}$, then \mathbf{v}_i is straight.
2. If $i \in I_{lo}$, then $\|\mathbf{v}_i\| = 2$.
3. If $i \in I_{hi}$, then $\|\mathbf{v}_i\| = 2h_0$.
4. $\alpha_{ij}(s) \leq \|\mathbf{v}_i - \mathbf{v}_j\| \leq \beta_{ij}(s)$ for all i, j .

Note that if s is unadorned, then $\mathcal{M}_s = \mathcal{B}_s''$.

Lemma 1.21 [GKFMJLC] s is a SCS for every $s \in S_{init}$.

Proof This is a direct result of definitions. □

Lemma 1.22 [UXCKFPE] [old:XWITCCN] [formal proof by Hoang Le Truong]. Let s be a SCS, and assume that

$$\tau^*(s, \mathbf{v}) < 0.$$

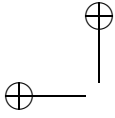
for some $\mathbf{v} \in \mathcal{B}_s$. Then \mathcal{B}_s' is nonempty.

Proof Let s be a SCS. Then by DSP [1], \mathcal{B}_s is compact (as a subset of $\mathcal{B}^k \subset \mathbb{R}^{3k}$).

The function

$$\mathbf{v} \mapsto \tau^*(s, \mathbf{v})$$

is a continuous function on \mathcal{B}_s . Moreover, if \mathcal{B}_s is nonempty, then the function attains a minimum. This follows from Lemma HDPLYGY and Lemma PCRT-TID. The set of minima is then nonempty. □



Lemma 1.23 [SGTRNAF] [old:AYQJTM] [formal proof by Hoang Le Truong]. *Let s be an unadorned SCS, and assume that*

$$\tau^*(s, \mathbf{v}) < 0.$$

for some $\mathbf{v} \in \mathcal{B}_s$. Then \mathcal{M}_s is nonempty.

Proof By Lemma 1.22, the set \mathcal{B}'_s of minimizers is nonempty. The subset \mathcal{B}''_s on which the index is as small as possible is then also nonempty. By assumption, s is unadorned, and $\mathcal{B}''_s = \mathcal{M}_s$. \square

Lemma 1.24 [EAPGLEJ] [formal proof by Hoang Le Truong]. *The main estimate holds if and only if $\mathcal{M}_s = \emptyset$ for all $s \in S_{\text{init}}$.*

Proof By Lemma 1.23 and Lemma 1.19. \square

Lemma 1.25 [JKQEWGV] [formal proof by Hoang Le Truong]. *Let s be a SCS. Let $\mathbf{v} \in \mathcal{B}_s$. Suppose that $\tau^*(s, \mathbf{v}) < 0$ and $k(s) > 3$. Then $\text{sol}(V_{\mathbf{v}}, E_{\mathbf{v}}, F_{\mathbf{v}}) < \pi$. Furthermore, the local fan is not circular, and the local fan can be lunar only when the pole has acute interior angle.*

Proof By the definition of SCS, we have $d(s) < 0.9$. The proof of Lemma GBY-CPXS extends readily to this context. The solid angle of a lune is less than π if and only if the pole has acute interior angle. \square

1.4 operations

This section describes some operations on SCSs. The first of these, a restriction, is a *hardening* of some of the soft constraints.

Definition 1.26 (restriction) [PFE0BSC] Let s be a SCS. We say that s' is a *restriction* of s of the *first type* if $I(s) = I(s')$ and

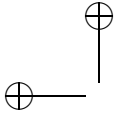
$$b_{ij}(s') = \beta_{ij}(s), \quad i, j \in I(s).$$

and all other parameters $k, d, J, a, \alpha, \beta$, etc. are the same for s and s' .

We say that s' is a *restriction* of s of the *second type* if

1. $I(s) = I(s')$,
2. $\alpha_{ij}(s) = \beta_{ij}(s)$ for all $i, j \in I(s)$,
3. $a_{ij}(s') = b_{ij}(s') = \alpha_{ij}(s)$, and
4. all other parameters k, d , etc. are the same for s and s' .

We say that s' is a restriction of s if it is a restriction of the first or second type.



Definition 1.27 (subdivision) [YYKMEWW] Let s be a SCS, and let $p, q \in I(s)$, with $p \neq q$. Let $c \in \mathbb{R}$. Define constants $\beta_0 = \min(b_{pq}(s), c)$ and $\alpha_0 = \max(a_{pq}(s), c)$. Define s_1 to be the same as s except that

$$\beta_{pq}(s_1) = \beta_0, \quad b_{pq}(s_1) = c.$$

Define s_2 to be the same as s except that

$$a_{pq}(s_2) = c, \quad \alpha_{pq}(s_2) = \alpha_0.$$

Define the subdivision of s to be the following list of one or two SCSs, according to the case.

$$\begin{cases} c \leq a_{pq}(s), & [s] \\ a_{pq}(s) < c \leq \alpha_{pq}(s), & [s_2] \\ \alpha_{pq}(s) < c < \beta_{pq}(s), & [s_1; s_2] \\ \beta_{pq}(s) \leq c < b_{pq}(s), & [s_1] \\ b_{pq}(s) \leq c, & [s] \end{cases}$$

The subdivision thus corresponds to splitting an interval $[a_{pq}, b_{pq}]$ into

$$[a_{pq}, c] \cup [c, b_{pq}].$$

Definition 1.28 [LCTBALA] We say that a SCS s *transfers* to a SCS s' if

1. If s is an ear, then $s = s'$.
2. s' is unadorned.
3. $I(s) = I(s')$.
4. $d(s) \leq d(s')$.
5. For all i, j , we have $a_{ij}(s') \leq a_{ij}(s) \leq b_{ij}(s) \leq b_{ij}(s')$.
6. $J(s') \subset J(s)$.

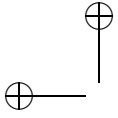
Definition 1.29 (equivalent SCS) If s is a SCS with torsor $I(s)$, and if I' is any equivalent torsor, then we can use the bijection between $I(s)$ and I' to obtain a SCS s' with $I' = I(s')$. An SCS s' related in this way to s is said to be *equivalent* to s . An *equi-transfer* of s to s' is a transfer from s to a SCS that is equivalent to s' .

Definition 1.30 (torsor slice) [ZTBHGM0] Let I be a $\mathbb{Z}/k\mathbb{Z}$ -torsor, with action given by $(j, i) \mapsto j + i$, for $i \in I$. Let $p, q \in I$ be non-adjacent. Set

$$I[p, q] = \{p, 1 + p, 2 + p, \dots, q\} \subset I.$$

Note that the cardinality of $I[p, q]$ is

$$m = 1 + \min\{j \in \mathbb{N} : j + p = q\}.$$



We make $I[p, q]$ into a $\mathbb{Z}/m\mathbb{Z}$ -torsor with action $(j, i) \mapsto j +' i$, given by the iterates of

$$1 +' i = \begin{cases} 1 + i & \text{if } i \neq q, \\ p & \text{if } i = q. \end{cases}$$

The $\mathbb{Z}/m\mathbb{Z}$ -torsor $I[p, q]$ is called the *slice* of I along (p, q) .

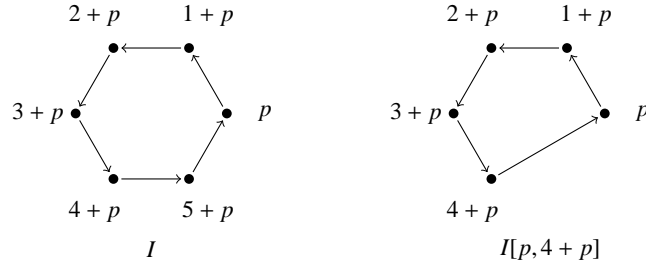


Figure 1.1 [WKUYEXM] Given $p, q \in I$, the slice $I[p, q]$ follows the cyclic order through I from p to q , then returns directly from q to p .

Definition 1.31 (SCS slice) [CJMHFAT] Let s be a SCS and let $\{p, q\} \subset I(s)$ be a diagonal. (In particular, $k(s) > 3$.) We say that a pair $\{s', s''\}$ is a *slice* along the diagonal $\{p, q\}$ of s , if the following conditions hold.

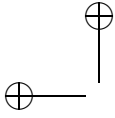
1. s' and s'' are unadorned SCSs.
2. $I(s') = I[p, q]$ and $I(s'') = I[q, p]$.
3. $d(s) \leq d(s') + d(s'')$.
4. $J(t) \subset J(s) \cup \{\{p, q\}\}$, for $t = s', s''$.
5. $\{p, q\} \in J(s')$ iff $\{p, q\} \in J(s'')$.
6. $\{p, q\} \in J(s')$ iff s' or s'' is an ear.
- 7.

$$a_{ij}(t) = a_{ij}(s), \text{ and } b_{ij}(t) = b_{ij}(s),$$

for $t = s', s''$ and all $i, j \in I(t)$ except the given diagonal $\{i, j\} = \{p, q\}$.

8. $a_{pq}(t) = \alpha_{pq}(s)$, $b_{pq}(t) = \beta_{pq}(s)$, $t = s', s''$.
9. $\beta_{pq}(s) < 4$.
10. $k(s) = 4$ or $\beta_{pq}(s) \leq c_{\text{stab}}$.

The word *slice* is used for related operations on the indexing set, the local fan, and the SCS. A slice of a SCS is used in parallel with the slice a fan of cardinality $k(s)$ into two smaller fans with cardinalities $k(s')$ and $k(s'')$. All of the edge length constraints are to be preserved under slicing.



If $\{s', s''\}$ is a slice of a diagonal $\{p, q\}$, then we can use the inclusions of sets $I(s') \subset I(s)$ and $I(s'') \subset I(s)$ to restrict an element $\mathbf{v} : I(s) \rightarrow \mathcal{B}$ to $\mathbf{v}' : I(s') \rightarrow \mathcal{B}$ and $\mathbf{v}'' : I(s'') \rightarrow \mathcal{B}$.

Lemma 1.32 [QKNVMLB] [formal proof by Hoang Le Truong]. *Let s be a SCS with diagonal $\{p, q\}$ and slice $\{s', s''\}$ along $\{p, q\}$. Let $\mathbf{v} \in \mathcal{M}_s$ and let \mathbf{v}' and \mathbf{v}'' be constructed from \mathbf{v} as above. Then $\mathbf{v}' \in \mathcal{B}_{s'}$ and $\mathbf{v}'' \in \mathcal{B}_{s''}$. Moreover,*

$$d(s, \mathbf{v}) \leq d(s', \mathbf{v}') + d(s'', \mathbf{v}'') \quad (1.33)$$

and

$$\tau^*(s, \mathbf{v}) \geq \tau^*(s, \mathbf{v}') + \tau^*(s, \mathbf{v}''). \quad (1.34)$$

Proof See Lemma MTUWLUN. \square

1.5 propagation

The proof of the main lemma consists in showing that the nonemptiness of \mathcal{M}_s propagates in an orderly way under the operations of restriction, slicing, equivalence, subdivision, and deformation.

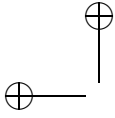
Definition 1.35 (\Rightarrow) [AZGJNZ0] [$\Rightarrow \Leftarrow \rightsquigarrow$ scs_arrow...]. Let **SCS** be the set of SCSs and let $P(\mathbf{SCS})$ be the powerset of **SCS**. We define a binary relation (\Rightarrow) on $P(\mathbf{SCS})$. When $S, T \subset \mathbf{SCS}$, we write $S \Rightarrow T$ to mean either that $\mathcal{M}_s = \emptyset$ for all $s \in S$, or that there exists $t \in T$ such that $\mathcal{M}_t \neq \emptyset$.

Lemma 1.36 [FZIOTEF] [formal proof by Hales]. *The relation (\Rightarrow) is reflexive and transitive.*

Proof Clearly, $\{s\} \Rightarrow \{s\}$, and this implies reflexivity. Transitivity is a simple matter. Assume $S_1 \Rightarrow S_2$ and $S_2 \Rightarrow S_3$. Assume $s_1 \in S_1$, with $\mathcal{M}_{s_1} \neq \emptyset$. Then select $s_2 \in S_2$ such that $\mathcal{M}_{s_2} \neq \emptyset$, and $s_3 \in S_3$ such that $\mathcal{M}_{s_3} \neq \emptyset$. Then $\{s_1\} \Rightarrow \{s_3\}$. So $S_1 \Rightarrow S_3$. \square

Lemma 1.37 (restriction) [EQTTNZI] [formal proof by Hoang Le Truong]. *Let s be a SCS, and let t be a restriction of s . When t has the first type, assume also that for all $\{i, j\} \in J$, we have $\beta_{ij}(s) = b_{ij}(s)$. Assume further for this type that $J(s) = \emptyset$ or $k(s) > 3$. When t has the second type, assume also that $J(s) = \emptyset$ and $m(t) + k(t) \leq 6$, with m as in the definition of SCS. Then $\{s\} \Rightarrow \{t\}$.*

Proof $\mathcal{B}_t \subset \mathcal{B}_s$. Assume $\mathbf{w} \in \mathcal{M}_s$. The assumptions according to the type give $\mathbf{w} \in \mathcal{B}_t$.



Assume the restriction has the first type. Then a global minimizer \mathbf{w} also minimizes on a subset, and the index does not change. Hence $\mathbf{w} \in \mathcal{M}_t$ and $\{s\} \Rightarrow \{t\}$.

Assume that the restriction has the second type. Then $\alpha_{ij}(s) = \beta_{ij}(s)$. We have $\mathcal{B}_t \subset \mathcal{B}_s$. For every $\mathbf{v} \in \mathcal{B}'_t$, the index of \mathbf{v} is k : every edge is fixed at its lower bound. Hence $\mathcal{B}'_t = \mathcal{B}''_t$, and $\mathbf{w} \in \mathcal{M}_t$. \square

Lemma 1.38 (subdivision) [UAGHHBM] [formal proof by Hoang Le Truong].
Let s be a SCS and let $\{s_1, s_2\}$ be a subdivision of s along $\{i, j\} \subset I(s)$. Then $\{s\} \Rightarrow \{s_1, s_2\}$.

Proof Let c denote the constant used for the subdivision. We have $\mathcal{B}_s = \mathcal{B}_{s_1} \cup \mathcal{B}_{s_2}$. If $\mathbf{w} \in \mathcal{M}_s$, then $\mathbf{w} \in \mathcal{B}'_{s_1} \cup \mathcal{B}'_{s_2}$.

Assume first that there exists $\mathbf{w} \in \mathcal{M}_s$ such that $c < \|\mathbf{w}_i - \mathbf{w}_j\|$. We claim that $\iota(s) \leq \iota(s_2)$. In fact,

$$\begin{aligned} \iota(s) &= \min_{\mathbf{v} \in \mathcal{B}'_s} \iota(s, \mathbf{v}) \\ &\leq \min_{\mathbf{v} \in \mathcal{B}'_{s_2}} \iota(s, \mathbf{v}) \\ &\leq \min_{\mathbf{v} \in \mathcal{B}'_{s_2}} \iota(s_2, \mathbf{v}) \\ &= \iota(s_2). \end{aligned}$$

We have $\mathbf{w} \in \mathcal{B}'_{s_2}$ and $\mathcal{B}'_{s_2} = \mathcal{B}'_s \cap \mathcal{B}_{s_2}$. Also,

$$\iota(s_2) \geq \iota(s) = \iota(s, \mathbf{w}) = \iota(s_2, \mathbf{w}) \geq \iota(s_2)$$

so that $\mathbf{w} \in \mathcal{B}''_{s_2}$ and $\mathbf{w} \in \mathcal{M}_{s_2}$.

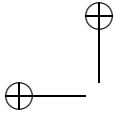
Finally, assume that every $\mathbf{w} \in \mathcal{M}_s$ has $\|\mathbf{w}_i - \mathbf{w}_j\| \leq c$. This means $\mathcal{M}_s \subset \mathcal{B}_{s_1}$. Let $\mathbf{w} \in \mathcal{M}_s$. We claim that $\iota(s) = \iota(s_1)$. In fact,

$$\begin{aligned} \iota(s) &= \min_{\mathbf{v} \in \mathcal{B}'_s} \iota(s, \mathbf{v}) \\ &\leq \min_{\mathbf{v} \in \mathcal{B}'_{s_1}} \iota(s, \mathbf{v}) \\ &= \min_{\mathbf{v} \in \mathcal{B}'_{s_1}} \iota(s_1, \mathbf{v}) \\ &= \iota(s_1). \end{aligned}$$

Also,

$$\iota(s) = \iota(s, \mathbf{w}) = \iota(s_1, \mathbf{w}) \geq \iota(s_1).$$

We have $\mathbf{w} \in \mathcal{M}_s \subset \mathcal{B}'_s \cap \mathcal{B}_{s_1} = \mathcal{B}'_{s_1}$. Since $\iota(s) = \iota(s_1)$, we also have $\mathbf{w} \in \mathcal{B}''_{s_1}$. So $\mathbf{w} \in \mathcal{M}_{s_1}$. \square



Lemma 1.39 (transfer, equivalence) [YXIONXL] [formal proof by Hoang Le Truong]. *Let s be a SCS and let t be a transfer or equivalent of s . Then $\{s\} \Rightarrow \{t\}$.*

Proof We treat transfer first. Assume that $\mathbf{w} \in \mathcal{M}_s$. We have $\mathbf{w} \in \mathcal{B}_s \subset \mathcal{B}_t$, and $d(t, \mathbf{w}) \geq d(s, \mathbf{w})$, as well as $\tau^*(t, \mathbf{w}) \leq \tau^*(s, \mathbf{w}) < 0$. By Lemma 1.22, there exists a global minimizer of smallest index $\mathbf{v} \in \mathcal{B}_t'' \neq \emptyset$. It satisfies $\tau^*(t, \mathbf{v}) \leq \tau^*(t, \mathbf{w}) < 0$. By the definition of transfer, t is unadorned, so that $\mathcal{B}_t'' = \mathcal{M}_t$. (Note that \mathcal{M}_s and \mathcal{M}_t might not be directly related.) This shows that $\mathcal{M}_t \neq \emptyset$. Hence $\{s\} \Rightarrow \{t\}$. Similarly, if t is equivalent to s , then we again have $\{s\} \Rightarrow \{t\}$. \square

Lemma 1.40 (slice) [LKGRQUI] [formal proof by Hoang Le Truong]. *Let $\{s', s''\}$ be a slice of a SCS s along a diagonal $\{p, q\}$. Then $\{s\} \Rightarrow \{s', s''\}$.*

Proof Assume that $\mathbf{w} \in \mathcal{M}_s$. Let $\mathbf{w}' \in \mathcal{B}_{s'}$ and $\mathbf{w}'' \in \mathcal{B}_{s''}$ be obtained by restriction of parameters. From Lemma 1.32, we have $\tau^*(s', \mathbf{w}') < 0$ or $\tau^*(s'', \mathbf{w}'') < 0$. To be concrete, say $\tau^*(s', \mathbf{w}') < 0$. A global minimizer $\mathbf{v}' \in \mathcal{B}_{s'}$ of smallest index then also satisfies $\tau^*(s', \mathbf{v}') < 0$. By the definition of slice, s' is unadorned, so that $\mathcal{B}_{s'}'' = \mathcal{M}_{s'}$ and $\mathbf{v}' \in \mathcal{M}_{s'}$. Hence $\{s\} \Rightarrow \{s', s''\}$. \square

1.6 deformation

This section proves some deformation results and implements them as arrows $\{s\} \Rightarrow T$.

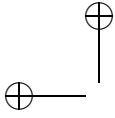
Lemma 1.41 [HXHYTIJ] [formal proof by Hoang Le Truong]. *Let s be a SCS, and let $\mathbf{v} \in \mathcal{B}_s''$. For Let $\mathbf{w} : I(s) \rightarrow \mathbb{R}^3$. Then one of the following holds:*

1. $\mathbf{w} \notin \mathcal{B}_s$.
2. $\tau^*(s, \mathbf{w}) > \tau^*(s, \mathbf{v})$.
3. $\tau^*(s, \mathbf{w}) = \tau^*(s, \mathbf{v})$ and the index of \mathbf{w} is at least that of \mathbf{v} .

Proof If the first two conclusions fail, then $\mathbf{w} \in \mathcal{B}_s'$. Since $\mathbf{v} \in \mathcal{B}_s''$, it must minimize the index over \mathcal{B}_s' . Hence the third conclusion holds. \square

Lemma 1.42 [ODXLSTCv2] [formal proof by Hoang Le Truong]. *Let s be a SCS and let $\mathbf{w} \in \mathcal{M}_s$. Fix $\ell \in I(s)$. Assume that \mathbf{w}_ℓ is not the pole of a lunar local fan $(V_{\mathbf{w}}, E_{\mathbf{w}}, F_{\mathbf{w}})$. Assume that $4h_0 < b_{\ell i}(s)$ for every diagonal $\{\ell, i\}$ at ℓ . Then one of the following hard constraints hold at index ℓ .*

1. $\|\mathbf{w}_\ell - \mathbf{w}_i\|$ attains its lower bound $a_{\ell i}(s)$, for some $i \neq \ell$.



2. $\|\mathbf{w}_\ell\|$ attains its lower bound 2.
3. There exists i adjacent to ℓ such that $\{\ell, i\} \in J(s)$.

Proof For a contradiction, assume that none of the enumerated constraints hold. By Lemma 1.25, the fan is not circular. The hypotheses allow us to use Lemmas MHAHEYJN and ZLZTHIC to deform \mathbf{w} at \mathbf{w}_ℓ .

The function τ^* is decreasing along the curve of the form [1][Equation 7.70] such that $\mathbf{w}_\ell(t) = (1 - t)\mathbf{w}_\ell$. That is, we push the point \mathbf{w}_ℓ radially towards the origin. Explicitly,

$$\tau^*(s, w) = c_1 + c_2 \|\mathbf{w}_\ell(t)\|$$

for some $c_2 > 0$ and c_1 . We have $\mathbf{w}(t) \in \mathcal{B}_s$ for all t positive and sufficiently small. This contradicts the minimality properties of \mathcal{M}_s . \square

Lemma 1.43 [IMJXPHRv2] *Let s be a SCS, $\mathbf{w} \in \mathcal{M}_s$, and \mathbf{w}_ℓ straight. Assume that \mathbf{w}_ℓ is not the pole of a lunar local fan $(V_{\mathbf{w}}, E_{\mathbf{w}}, F_{\mathbf{w}})$. Assume that $4h_0 < b_{\ell i}(s)$ for every diagonal $\{\ell, i\}$ at ℓ . Assume that $\mathbf{w}_\ell \in \text{aff}_+(\mathbf{0}, \{\mathbf{w}_{\ell+1}, \mathbf{w}_{\ell-1}\})$. Then one of the following hard constraints holds at ℓ .*

1. $\|\mathbf{w}_\ell - \mathbf{w}_{\ell+1}\|$ attains its lower bound $a_{\ell, \ell+1}(s)$, and $\|\mathbf{w}_\ell - \mathbf{w}_{\ell-1}\|$ attains its lower bound $a_{\ell, \ell-1}(s)$.
2. $\|\mathbf{w}_\ell\|$ attains its lower bound 2.
3. There exists i adjacent to ℓ such that $\{\ell, i\} \in J(s)$.
4. Some diagonal $\{\ell, i\} \subset I(s)$ at ℓ satisfies $\|\mathbf{w}_\ell - \mathbf{w}_i\| = a_{\ell i}(s)$.

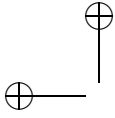
Proof For a contradiction, assume that none of the enumerated constraints hold. By Lemma 1.25, the fan is not circular. The hypotheses allow us to use Lemmas MHAHEYJN and ZLZTHIC to deform \mathbf{w} at \mathbf{w}_ℓ .

The set $\{\mathbf{0}, \mathbf{w}_{\ell-1}, \mathbf{w}_\ell, \mathbf{w}_{\ell+1}\}$ lies in a plane A . By the previous lemma one of the norm constraints is satisfied, say

$$\|\mathbf{w}_\ell - \mathbf{w}_{\ell+1}\| = a_{\ell, \ell+1}(s).$$

We consider a deformation of \mathbf{w} that moves $\mathbf{w}_\ell(t)$ along a circle through \mathbf{w}_ℓ with center $\mathbf{w}_{\ell+1}$ in the plane A . Parameterize the curve so that as t increases, the norm $\|\mathbf{w}_\ell(t)\|$ decreases. The function $\tau^*(s, \mathbf{w})$ is decreasing in t . Explicitly, the function again depends linearly on $\|\mathbf{w}_\ell(t)\|$, because the angle at ℓ remains straight. This contradicts the minimality of \mathcal{M}_s . The result ensues. \square

Lemma 1.44 [NUXCOEA v2] *Let s be a SCS and let $\mathbf{w} \in \mathcal{M}_s$. Fix $\ell \in I_{\text{str}}(s)$. Assume j is an index adjacent to ℓ such that $\|\mathbf{w}_\ell - \mathbf{w}_j\| = a_{\ell j}(s)$. Assume that \mathbf{w}_ℓ is not the pole of a lunar local fan $(V_{\mathbf{w}}, E_{\mathbf{w}}, F_{\mathbf{w}})$. Assume that $4h_0 < b_{\ell i}(s)$ for every diagonal $\{\ell, i\}$ at ℓ . Then one of the following hard conditions holds at index ℓ .*



1. We have $\|\mathbf{w}_\ell - \mathbf{w}_i\| = a_{\ell k}(s)$ for both choices of $i \in I(s)$ adjacent to ℓ .
2. There exists i adjacent to ℓ such that $\{\ell, i\} \in J(s)$.
3. Some diagonal $\{\ell, i\} \subset I(s)$ at ℓ satisfies $\|\mathbf{w}_\ell - \mathbf{w}_i\| = a_{\ell i}(s)$.

Proof The lemma is a special case of the previous lemma, unless $\|\mathbf{w}_\ell\| = 2$, which we assume. Let $i \neq j$ be the other index adjacent to ℓ . Assume that the three enumerated parts of the conclusion fail.

We consider a deformation of \mathbf{w} that only moves \mathbf{w}_ℓ . We take the motion of \mathbf{w}_ℓ to be in a circular arc with center $\mathbf{0}$ through the point \mathbf{w}_ℓ and in the fixed plane determined by $\{\mathbf{0}, \mathbf{w}_\ell, \mathbf{w}_j, \mathbf{w}_i\}$. The function $\tau^*(s, *)$ is constant along this curve. We orient the curve to be increasing in $\|\mathbf{w}_\ell - \mathbf{w}_j\|$. For sufficiently small t , we find that $\mathbf{w}(t) \in \mathcal{B}'_s$ has smaller index than \mathbf{w} . This is contrary to the minimizing properties of $\mathbf{w} \in \mathcal{M}_s$.

As in the proofs of the previous lemmas, the fan is not circular, and the constraints on generic and local fans allow us to use Lemmas MHAHEYJN and ZLZTHIC, showing that fan conditions are preserved. \square

In the preceding three lemmas, we specifically allow the deformations $\mathbf{v}(t)$ to occur within a lunar fan, moving a single node that is not a pole of the lunar fan, as given by Lemma MHAHEYJN. We consider some further deformations on generic \mathbf{v} .

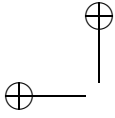
Lemma 1.45 (482) [CUXVZ0Z] *Let s be a basic SCS, and let $p \in I(s)$. Let $\mathbf{v} \in \mathcal{M}_s$ be generic. We make the following assumptions.*

1. For all diagonals except $\{p-1, p+1\}$, we have $a_{i,j}(s) < \|\mathbf{v}_i - \mathbf{v}_j\|$.
2. $3 \leq \|\mathbf{v}_{p-1} - \mathbf{v}_{p+1}\|$.
3. $\angle \mathbf{v}_p < \pi$.
4. $\pi/2 < \angle \mathbf{v}_p$ or $\angle \mathbf{v}_{p+1} < \pi$.
5. The constraints along $\{p, p+1\}$ are standard.
6. $2 \leq \|\mathbf{v}_p - \mathbf{v}_{p-1}\| \leq c_{stab}$.

Then $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| = 2$.

Proof We construct a deformation using Lemma 3.1. Specifically, we take the base triangle to be $\{\mathbf{0}, \mathbf{v}_{p-1}, \mathbf{v}_{p+1}\}$ and variable point to be \mathbf{v}_p with variable edge $\{\mathbf{v}_p, \mathbf{v}_{p+1}\}$. LEMMA_4828966562 states that the condition $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| > 2$ is not a local minimum. Because of the condition $\pi/2 < \angle \mathbf{v}_p < \pi$, decreasing $\|\mathbf{v}_p - \mathbf{v}_{p+1}\|$ along this deformation will preserve the nonreflexivity condition of a nonreflexive local fan. \square

Lemma 1.46 (482 bis) [CJBDXXN] *Let s be a basic SCS, and let $p \in I(s)$. Let $\mathbf{v} \in \mathcal{M}_s$ be generic. Assume that for all diagonals except $\{p-1, p+1\}$, we*



have $a_{i,j}(s) < \|\mathbf{v}_i - \mathbf{v}_j\|$. Assume $3 \leq \|\mathbf{v}_{p-1} - \mathbf{v}_{p+1}\|$. Assume $\pi/2 < \angle \mathbf{v}_p < \pi$ or $\angle \mathbf{v}_{p-1} < \pi$. Assume that the constraints along $\{p-1, p\}$ are standard. Assume $2 \leq \|\mathbf{v}_p - \mathbf{v}_{p+1}\| \leq c_{stab}$. Then $\|\mathbf{v}_p - \mathbf{v}_{p-1}\| = 2$.

Proof We construct a deformation using Lemma 3.1 as in the previous lemma, now using $\{\mathbf{v}_p, \mathbf{v}_{p-1}\}$ as the variable edge. \square

Abbreviate $\text{xrr}_V(\mathbf{u}, \mathbf{v}) = \text{xrr}(\|\mathbf{u}\|, \|\mathbf{v}\|, \|\mathbf{u} - \mathbf{v}\|)$.

Lemma 1.47 684 pent[IUNBUIG] *Let s be a basic SCS with $k(s) = 5$. Let $p \in I(s)$ and $\mathbf{v} \in \mathcal{M}_s$ be generic. Assume that for all diagonals, we have $a_{i,j}(s) \leq c_{stab} < \|\mathbf{v}_i - \mathbf{v}_j\|$ and $4h_0 < b_{ij}(s)$. Assume that $\angle \mathbf{v}_{p+2} < \pi$ and $\angle \mathbf{v}_{p+3} < \pi$. Assume that $2 \leq a_{p+2,p+3} \leq b_{p+2,p+3} \leq c_{stab}$. Assume that $\|\mathbf{v}_i - \mathbf{v}_j\| = 2$ for every edge other than $\{p+2, p+3\}$. Assume that $2 \leq \|\mathbf{v}_i - \mathbf{v}_j\|$ and $\text{xrr}_V(\mathbf{v}_i, \mathbf{v}_j) \leq 15.53$, for $\{i, j\} = \{p, p+2\}, \{p, p+3\}$. Then $\|\mathbf{v}_{p+2} - \mathbf{v}_{p+3}\| = a_{p+2,p+3}$.*

We only need to use the special case $(p+2, p+3) = (0, 1)$.

Proof We construct a deformation using Lemma 3.1. We use $\{\mathbf{0}, \mathbf{v}_p, \mathbf{v}_{p+3}\}$ as the base triangle with constructed point \mathbf{v}_{p+2} and variable edge $\{p+2, p+3\}$. We construct a second time on the base triangle $\{\mathbf{0}, \mathbf{v}_p, \mathbf{v}_{p+2}(t)\}$ with constructed point \mathbf{v}_{p+1} . LEMMA_6843920790 asserts that $\|\mathbf{v}_{p+2} - \mathbf{v}_{p+3}\| > a_{p+2,p+3}$ is incompatible with the minimality of \mathbf{v} . The conditions $\angle \mathbf{v}_{p+2} < \pi$ and $\angle \mathbf{v}_{p+3} < \pi$ are used to maintain nonreflexivity.

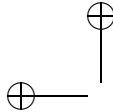
To use LEMMA_6843920790, we must check the precondition that $\{\mathbf{0}, \mathbf{v}_p, \mathbf{v}_{p+2}, \mathbf{v}_{p+3}\}$ is not coplanar. We leave this as an exercise based on the triangle inequality? \square

Lemma 1.48 (1834) [NEHXMWH] *Let s be a basic SCS, with $k \leq 4$, and let $p \in I(s)$. Let $\mathbf{v} \in \mathcal{M}_s$ be generic. Assume that for all diagonals we have $a_{i,j}(s) \leq c_{stab} < \|\mathbf{v}_i - \mathbf{v}_j\|$ and $4h_0 < b_{i,j}(s)$. Assume $\angle \mathbf{v}_i < \pi$ for all $i \neq p+2$. Then $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| = a_{p,p+1}$ or $b_{p,p+1}$.*

Proof As in previous lemmas, use Lemma 3.1 to construct a deformation. Use $\{\mathbf{0}, \mathbf{v}_{p+1}, \mathbf{v}_{p-1}\}$ as the base triangle with constructed point \mathbf{v}_p , where $\|\mathbf{v}_{p+1} - \mathbf{v}_{p+2}\|$ varies. Continue as before with LEMMA_1834976363. \square

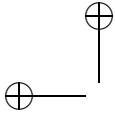
Lemma 1.49 (1834 sym) [BZQNDMM] *Let s be a basic SCS, with $k \leq 4$, and let $p \in I(s)$. Let $\mathbf{v} \in \mathcal{M}_s$ be generic. Assume that for all diagonals we have $a_{i,j}(s) \leq c_{stab} < \|\mathbf{v}_i - \mathbf{v}_j\|$ and $4h_0 < b_{i,j}(s)$. Assume $\angle \mathbf{v}_i < \pi$ for all $i \neq p-1$. Then $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| = a_{p,p+1}$ or $b_{p,p+1}$.*

Proof As in previous lemmas, use Lemma 3.1 to construct a deformation. Use $\{\mathbf{0}, \mathbf{v}_p, \mathbf{v}_{p+2}\}$ as the base triangle with constructed point \mathbf{v}_{p+1} , where $\|\mathbf{v}_p - \mathbf{v}_{p+1}\|$ varies. Continue as before with LEMMA_1834976363. \square



Lemma 1.50 (684) [JOTSWIX] *Let s be a basic SCS, with $k(s) = 4$, and let $p \in I(s)$. Let $\mathbf{v} \in \mathcal{M}_s$ be generic. Assume that for all diagonals we have $a_{i,j}(s) \leq c_{stab} < \|\mathbf{v}_i - \mathbf{v}_j\|$ and $4h_0 < b_{i,j}(s)$. Assume $\angle \mathbf{v}_i < \pi$ for all $i \neq p+2$. Assume $\text{xrr}_V(\mathbf{v}_{p-1}, \mathbf{v}_{p+1}) \leq 15.53$. Assume $a_{p,p+1}(s) < c_{stab}$. Then $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| < c_{stab}$ or $\|\mathbf{v}_p - \mathbf{v}_{p-1}\| < c_{stab}$.*

Proof As in previous lemmas, use Lemma 3.1 to construct a deformation. Use $\{\mathbf{0}, \mathbf{v}_{p-1}, \mathbf{v}_{p+1}\}$ as the base triangle with constructed point \mathbf{v}_p , where $\|\mathbf{v}_{p-1} - \mathbf{v}_p\|$ varies. Continue as before with LEMMA_6843920790. \square



2

Appendix on checking completeness

In this section we prove the implication $S_{\text{init}} \Rightarrow S_{\text{term}}$.

2.1 definitions

We let

$$M(s) = \{(i, i+1) : i \in I(s), \quad (2 < a_{i,i+1}(s) \text{ or } 2h_0 < b_{i,i+1}(s))\}.$$

By the definition of SCS, the cardinality $m(s)$ of $M(s)$ satisfies $k(s) + m(s) \leq 6$.

The stable constraints systems we discuss here are all unadorned with $J = I_{\text{str}} = I_{\text{lo}} = I_{\text{hi}} = \emptyset$; that is, basic stable constraint systems. Little further mention will be made of this sets.

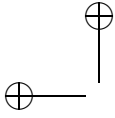
We say that $\mathbf{v} \in \mathcal{B}_s$ is generic if its local fan $(V_{\mathbf{v}}, E_{\mathbf{v}}, F_{\mathbf{v}})$ is generic. We say that $\mathbf{v} \in \mathcal{B}_s$ is lunar if its local fan $(V_{\mathbf{v}}, E_{\mathbf{v}}, F_{\mathbf{v}})$ is generic.

2.2 init list

We name stable constraint systems s with a three-character identifier. The first digit is $k(s)$. The second is the letter I (initial), T (terminal), or M (intermediate). The third character is a distinguishing digit.

We let std denote the range $[2, 2h_0]$ of a standard edge, pro denote the range $[2h_0, \sqrt{8}]$ of a prolonged edge, and pro^+ denote the range $[2h_0, c_{\text{stab}}]$. Let pro_8 denote the range $[\sqrt{8}, c_{\text{stab}}]$. If r is a real number, we write $r+$ for the interval $[r, 6]$, and by abuse of notation r for the interval $[r, r]$.

We have the given table (2.2) of SCSs in the initial list. In each case we give the value of k and d , the range of the diagonals, the default range on edges, and the range of one distinguished edge if there is one.



name	interval
std	$[2, 2h_0]$
pro	$[2h_0, \sqrt{8}]$
pro_8	$[\sqrt{8}, c_{\text{stab}}]$
pro^+	$[2h_0, c_{\text{stab}}]$
pro	$[2h_0, \sqrt{8}]$

Table 2.1 *intervals*

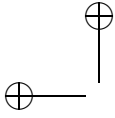
name	k	d	diag	edge	dist.edge
6I1	6	0.712	$2h_0+$	std	*
5I1	5	0.4819	$2h_0+$	std	*
5I2	5	0.616	$\sqrt{8}+$	std	*
5I3	5	0.616	$2h_0+$	std	pro
4I1	4	0.206	$2h_0+$	std	*
4I2	4	0.467	$3+$	std	*
4I3	4	0.477	$\sqrt{8}+$	std	pro
3I1	3	0	*	std	*

Table 2.2 *Initial SCS*

We have the given table (2.2) of SCSs in the terminal list. Two entries in the diagonal column separated by a semicolon indicates that there are two diagonals (in a quadrilateral) with different length constraints. In 3T7, there are three different edge ranges.

name	k	d	diag	edge	dist.edge
6T1	6	0.712	$c_{\text{stab}}+$	2	*
5T1	5	0.616	$c_{\text{stab}}+$	2	*
4T1	4	0.467	$3+$	2	*
4T2	4	0.467	3	std	*
4T3	4	0.513	$c_{\text{stab}}+$	2	c_{stab}
4T4	4	0.477	$\sqrt{8}+; \text{pro}_8$	std	pro
4T5	4	0.513	$c_{\text{stab}}; c_{\text{stab}}+$	std	pro^+
3T1	3	0.11	*	std	pro_8
3T2	3	0	*	2	*
3T3	3	0.476	*	pro^+	*
3T4	3	0.2759	*	pro^+	std
3T5	3	0.103	*	std	pro
3T6	3	0.4348	*	pro_8	std
3T7	3	0.2565	*	2	$c_{\text{stab}}; [c_{\text{stab}}, 3.62]$

Table 2.3 *terminal SCS*



We define basic SCSs with the following parameters. In 4M4, 4M5, 4M7, and 4M8, there are two standard edges and two with range pro^+ , differing according to whether the standard edges are adjacent or opposite. In particular,

name	k	d	diag	edge	dist.edge
6M1	6	0.712	$2h_0+$	std	*
5M1	5	0.616	$2h_0+$	std	pro^+
5M2	5	0.616	$c_{\text{stab}}+$	std	$[2, c_{\text{stab}}]$
4M2	4	0.3789	$2h_0+$	std	pro^+
4M3'	4	0.513	$\sqrt{8}+$	std	pro_8
4M4'	4	0.513	$2h_0+$	std (2 adj)	$\text{pro}^+(2 \text{ adj})$
4M5'	4	0.513	$2h_0+$	std (2 opp)	$\text{pro}^+(2 \text{ opp})$
4M6'	4	0.513	$c_{\text{stab}}+$	std	pro^+
4M7	4	0.513	$c_{\text{stab}}+$	std (2 adj)	$\text{pro}^+(2 \text{ adj})$
4M8	4	0.513	$c_{\text{stab}}+$	std (2 opp)	$\text{pro}^+(2 \text{ opp})$
3M1	3	0.103	*	std	pro^+

2.3 preliminary lemmas

Lemma 2.1 [FEKTYIY]. [formal proof by Hoang Le Truong]. *Let s be an SCS with $k(s) > 3$, and let $\mathbf{v} \in \mathcal{M}_s$. Then $\{\mathbf{0}\} \cup V_{\mathbf{v}}$ is not coplanar.*

Proof If $\{\mathbf{0}\} \cup V_{\mathbf{v}}$ is coplanar, then every interior angle in the local fan $(V_{\mathbf{v}}, E_{\mathbf{v}}, F_{\mathbf{v}})$ is 0 or π . However, the fan conditions prevent angle 0. Hence every interior angle is π . This gives $\text{sol}(V_{\mathbf{v}}, E_{\mathbf{v}}, F_{\mathbf{v}}) = 2\pi$, which is contrary to the estimate of Lemma 1.25 (JKQEWGV). \square

Lemma 2.2 [AURSIPD] [formal proof by Hoang Le Truong]. *Let s be an SCS with $k(s) > 3$, and let $\mathbf{v} \in \mathcal{M}_s$ be generic. Then $3 + \ell \leq k$, where ℓ is the number of straight interior angles of $(V_{\mathbf{v}}, E_{\mathbf{v}}, F_{\mathbf{v}})$.*

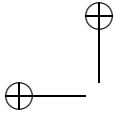
Proof Otherwise $2 + \ell \geq k$, and we can find two distinct i_1, i_2 such that the interior angle is straight at \mathbf{v}_i for $i \neq i_1, i_2$. By the straightness of $i_1 + 1, \dots$ we have that

$$\{\mathbf{0}, \mathbf{v}_{i_1}, \mathbf{v}_{i_1+1}, \dots, \mathbf{v}_{i_2}\}$$

lies in the plane $A = \text{aff}\{\mathbf{0}, \mathbf{v}_{i_1}, \mathbf{v}_{i_2}\}$. Similarly,

$$\{\mathbf{0}, \mathbf{v}_{i_2}, \mathbf{v}_{i_2+1}, \dots, \mathbf{v}_{i_1}\}$$

lies in A . Hence $V_{\mathbf{v}}$ is coplanar, which is contrary to Lemma 2.1. \square



Lemma 2.3 (polar test) [PPBTYDQ] [formal proof by Hoang Le Truong].

Let \mathbf{u} , \mathbf{p} , and $\mathbf{v} \in \mathbb{R}^3$ be vectors such that $\{\mathbf{0}, \mathbf{p}, \mathbf{u}\}$ is not collinear and $\{\mathbf{0}, \mathbf{p}, \mathbf{v}\}$ is not collinear. If

$$\text{arc}_V(\mathbf{0}, \mathbf{u}, \mathbf{p}) + \text{arc}_V(\mathbf{0}, \mathbf{p}, \mathbf{v}) < \pi$$

then $\mathbf{0}$ is not in the convex hull of $\{\mathbf{u}, \mathbf{v}\}$.

Proof Assume for a contradiction that $\mathbf{0}$ does lie in the convex hull. Let $\mathbf{e} = \mathbf{u} \times \mathbf{p}$. By the orientation chosen for \mathbf{e} , and by (DIHV_ARCV), we have

$$\text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}, \mathbf{p}) = \text{dih}_V(\mathbf{0}, \mathbf{e}, \mathbf{u}, \mathbf{p}) = \text{arc}_V(\mathbf{0}, \mathbf{u}, \mathbf{p})$$

$$\text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{p}, \mathbf{v}) = \text{dih}_V(\mathbf{0}, \mathbf{e}, \mathbf{p}, \mathbf{v}) = \text{arc}_V(\mathbf{0}, \mathbf{p}, \mathbf{v}).$$

Hence the assumption of the lemma and additivity of azimuth angles gives

$$\text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}, \mathbf{v}) = \text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}, \mathbf{p}) + \text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{p}, \mathbf{v}) < \pi.$$

If $\mathbf{0}$ does lie in the convex hull, then $\{\mathbf{0}, \mathbf{u}, \mathbf{v}, \mathbf{e}\}$ is coplanar, which implies that $\text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}, \mathbf{v}) \in \{0, \pi\}$. We have eliminated the possibility π . If the azimuth angle is 0, then \mathbf{u} and \mathbf{v} lie in the same half plane $\mathbf{u} \in \text{aff}_+^0(\{\mathbf{0}, \mathbf{e}\}, \{\mathbf{v}\})$. Hence $\mathbf{0}$ is not in the convex hull. \square

Lemma 2.4 [OIQKKEP] Let $\mathbf{u}, \mathbf{v} \in \mathcal{B}$. Assume that $\|\mathbf{u} - \mathbf{v}\| \leq c < 4$, for some c . Then

$$\text{arc}_V(\mathbf{0}, \mathbf{u}, \mathbf{v}) \leq \text{arc}(2, 2, c).$$

Proof

$$\text{arc}_V(\mathbf{0}, \mathbf{u}, \mathbf{v}) = \text{arc}(\|\mathbf{u}\|, \|\mathbf{v}\|, \|\mathbf{u} - \mathbf{v}\|).$$

The function arc is decreasing in the first two arguments and increasing in the third. \square

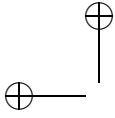
Lemma 2.5 [BGSAXHJ] [formal proof by Hoang Le Truong]. Let (V, E, F) be a lunar fan with pole (\mathbf{v}, \mathbf{w}) . Then $\mathbf{0}$ lies in the convex hull of $\{\mathbf{v}, \mathbf{w}\}$.

Proof This follows directly from the definition of a lunar fan. \square

Lemma 2.6 [MXQTIED] [formal proof by Hoang Le Truong]. Let s be a basic SCS, and let $\mathbf{v} \in \mathcal{M}_s$. Let s' be a basic SCS such that $M(s') = M(s)$, $k(s') = k(s)$, $I(s') = I(s)$, $\mathbf{v} \in \mathcal{B}_{s'}$, $d(s') = d(s)$, $a_{ii}(s') = 0$, $a_{ij}(s') = a_{ij}(s)$ for all edges (i, j) , and

$$a_{ij}(s) \leq a_{ij}(s') \leq \|\mathbf{v}_i - \mathbf{v}_j\| \leq b_{ij}(s') \leq b_{ij}(s).$$

Then $\mathbf{v} \in \mathcal{M}_{s'}$.



Proof Clearly $\mathbf{v} \in \mathcal{B}_{s'}$. The domain shrinks and \mathbf{v} is a global minimizer, so $\mathbf{v} \in \mathcal{B}'_{s'}$. By the hypotheses the constants a_{ij} are the same on edges, so the index remains the same, and $\mathbf{v} \in \mathcal{B}''_{s'}$. The SCS is unadorned, so $\mathbf{v} \in \mathcal{M}_{s'}$. \square

Lemma 2.7 [XWNHLM] [formal proof by Hoang Le Truong]. *Let s and s' be basic SCSs with $d(s') \geq d(s)$, $k(s) = k(s')$, and $I(s) = I(s')$. Suppose that $\mathbf{v} \in \mathcal{M}_s \cap \mathcal{B}_{s'}$. Then $\{s\} \Rightarrow \{s'\}$.*

Proof The condition $\mathbf{v} \in \mathcal{M}_s$ gives $\tau^*(s', \mathbf{v}) \leq \tau^*(s, \mathbf{v}) < 0$. Hence \mathcal{B}'_s is nonempty. The SCS is basic, so $\mathcal{B}''_{s'} = \mathcal{M}_{s'}$ is also nonempty. \square

Lemma 2.8 [SYNQIWN] [formal proof by Hales]. *Let s be a basic SCS and let $\mathbf{v} \in \mathcal{B}_s$. Let $i \in I(s)$. Assume that at least two the following four lengths are equal to 2:*

$$\|\mathbf{v}_i\|, \quad \|\mathbf{v}_i - \mathbf{v}_{i+1}\|, \quad \|\mathbf{v}_{i+2}\|, \quad \|\mathbf{v}_{i+2} - \mathbf{v}_{i+1}\|$$

Assume that $c_{stab} \leq \|\mathbf{v}_i - \mathbf{v}_{i+2}\|$. Then the interior angle at $i+1$ is greater than $\pi/2$.

Proof This is a computer calculation¹ [2]. \square

Lemma 2.9 (lunar prep) [AXJRPNC] [formal proof by Hoang Le Truong]. *Let s be a basic SCS and let $\mathbf{v} \in \mathcal{M}_s$. Assume that $b_{ij}(s) \leq c_{stab}$ for all edges $\{i, j\}$. Assume that $(V_{\mathbf{v}}, E_{\mathbf{v}}, F_{\mathbf{v}})$ is a lune with pole $(\mathbf{v}_p, \mathbf{v}_q)$. Then $k(s) = 6$ and $q = p + 3$.*

Proof We first consider the case that $k = 6$ and $q = p, p + 1$, or $p + 2$. Then $M(s) = \emptyset$ and the polar criterion gives the result, when applied to

$$2 \arccos(2, 2, 2h_0) < \pi.$$

Assume for a contradiction that $k < 6$. In particular, the floor of $k/2$ is at most 2. We have $\|\mathbf{v}_i - \mathbf{v}_j\| < 4$ for every edge, and this implies that (p, q) is not an edge. We have that $k > 3$; for otherwise (p, q) is an edge. Without loss of generality we may assume that $q = p + 2$.

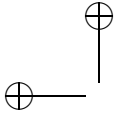
We start a special case by assuming that $M(s)$ meets $\{(p, p+1), (p+1, p+2)\}$ in a set of cardinality at most one. Then by the monotonicity of arc, we apply the polar test

$$\arccos(\mathbf{0}, \mathbf{v}_p, \mathbf{p}) + \arccos(\mathbf{0}, \mathbf{p}, \mathbf{v}_q) \leq \arccos(2, 2, 2h_0) + \arccos(2, 2, c_{stab}) < \pi,$$

to see that (p, q) is not a pole.

In the final case we assume that $M(s)$ meets $\{(p, p+1), (p+1, p+2)\}$ in a set

¹ [1117202051 4559601669 4559601669b]



of cardinality two. By the inequality $m + k \leq 6$, we see that $m = 2$ and $k = 4$. This implies that $M(s)$ is disjoint from $\{(p+2, p+3), (p+3, p)\}$. By the polar test,

$$\text{arc}_V(\mathbf{0}, \mathbf{v}_p, \mathbf{v}_q) \leq 2 \text{arc}(2, 2, 2h_0) < \pi,$$

and again (p, q) is not a pole. \square

A weak version of the following lemma has been formalized as `RRCWNS_WEAK`.

Lemma 2.10 (genericity) [RRCWNSJJ] [formal proof by Hoang Le Truong]. *Let s be a basic SCS and $\mathbf{v} \in \mathcal{M}_s$. Assume that $4h_0 < b_{ij}(s)$ for every diagonal (i, j) . Assume that $a_{ij}(s) \leq c_{stab} < \|\mathbf{v}_i - \mathbf{v}_j\|$ for every diagonal (i, j) . Assume that $b_{ij}(s) \leq c_{stab}$ for every edge (i, j) . Then $(V_{\mathbf{v}}, E_{\mathbf{v}}, F_{\mathbf{v}})$ is generic.*

Proof An earlier lemma shows that $(V_{\mathbf{v}}, E_{\mathbf{v}}, F_{\mathbf{v}})$ is not circular.

Assume for a contradiction that $(V_{\mathbf{v}}, E_{\mathbf{v}}, F_{\mathbf{v}})$ is lunar with pole $(\mathbf{v}_p, \mathbf{v}_q)$. The previous lemma gives $k(s) = 6$ and $q = p + 3$. By the definition of SCS, we have $m = 0$. That is, $b_{ij}(s) \leq 2h_0$ and $a_{ij} = 2$ for all edges.

By the deformation ODX applied at $p + 1$, we have

$$\|\mathbf{v}_{p+1}\| = 2, \text{ or } \|\mathbf{v}_{p+1} - \mathbf{v}_p\| = 2, \text{ or } \|\mathbf{v}_{p+1} - \mathbf{v}_{p+2}\| = 2.$$

Furthermore, if the third possibility holds: $\|\mathbf{v}_{p+1} - \mathbf{v}_{p+2}\| = 2$, then the deformation IMJ at $p + 1$ gives that one of the first two possibilities hold.

By similar reasoning at vertex $p - 1$, we have

$$\|\mathbf{v}_{p-1}\| = 2, \text{ or } \|\mathbf{v}_p - \mathbf{v}_{p-1}\| = 2.$$

By Lemma 2.8, the interior angle at p is greater than $\pi/2$. This is contrary to Lemma 1.25 (JKQEWGV). \square

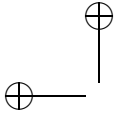
Lemma 2.11 (exists-2) [JCYFMRP] *Let s be a basic SCS, with $3 < k(s)$, and let $\mathbf{v} \in \mathcal{M}_s$ be generic. Assume that $4h_0 < b_{ij}(s)$ for all diagonals, and that $m(s) \leq 1$. Assume that $a_{ij}(s) = 2$ for all edges $\{i, j\}$. Assume that for all diagonals $\{i, j\}$*

$$a_{ij}(s) \leq c_{stab} < \|\mathbf{v}_i - \mathbf{v}_j\| \text{ and } 4h_0 < b_{ij}(s).$$

Then there exists an edge $\{i, j\}$ such that $\|\mathbf{v}_i - \mathbf{v}_j\| = 2$.

Proof For a contradiction, assume $\|\mathbf{v}_i - \mathbf{v}_j\| > 2$ for all edges. Apply the deformation ODX at every vertex ℓ . Then $\|\mathbf{v}_\ell\| = 2$ for all ℓ . Since $m(s) \leq 1$, we have some index p such that for every edge $(i, i + 1)$ other than $(p, p + 1)$, we have $b_{ij}(s) \leq 2h_0$.

By Lemma 2.8, the interior angle is obtuse at every i , such that $i \neq p, p + 1$.



We claim that each of the vertices \mathbf{v}_i is straight $i \neq p, p+1$. For otherwise, we may apply the obtuse case of deformation 482 to conclude that some edge has length 2.

Thus, we have constructed $k-2$ straight angles. This is contrary to Lemma 2.2. \square

Lemma 2.12 (propagate-min-std) [TFITSKC] *Let s be a basic SCS, with $3 < k(s)$, and let $\mathbf{v} \in \mathcal{M}_s$ be generic. Let $p \in I(s)$. Assume that $(p, p+1)$ is a standard edge. Assume that $\|\mathbf{v}_{p-1} - \mathbf{v}_p\| = 2$. Assume that for all diagonals $\{i, j\}$*

$$a_{ij}(s) \leq c_{stab} < \|\mathbf{v}_i - \mathbf{v}_j\| \text{ and } 4h_0 < b_{ij}(s).$$

Assume

$$2 < b_{p-1,p}(s), \quad \leq a_{p+1,p+2}(s) < b_{p+1,p+2}(s)$$

Then $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| = 2$.

Proof Assume for a contradiction that $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| > 2$.

We claim that \mathbf{v}_p is not straight. Otherwise, deformation NUX applied at \mathbf{v}_p gives $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| = 2$.

We claim that $\|\mathbf{v}_{p+1}\| > 2$. Otherwise, by Lemma 2.8, the interior angle at \mathbf{v}_p is obtuse, and the obtuse case of deformation 482 gives the $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| = 2$.

By deformation ODX applied at $p+1$, we have $\|\mathbf{v}_{p+1} - \mathbf{v}_{p+2}\| = a_{p+1,p+2}$.

We claim that \mathbf{v}_{p+1} is not straight. Otherwise, deformation NUX applied at \mathbf{v}_{p+1} gives $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| = 2$.

Finally, we may apply the deformation 482 in the non-obtuse case at $p+1$ to obtain $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| = 2$. \square

Lemma 2.13 (propagate-min-std2) [CQAOQLR] *Let s be a basic SCS, with $3 < k(s)$, and let $\mathbf{v} \in \mathcal{M}_s$ be generic. Assume that for all diagonals $\{i, j\}$, we have $a_{ij}(s) \leq c_{stab} < \|\mathbf{v}_i - \mathbf{v}_j\|$ and $4h_0 < b_{ij}(s)$. Assume*

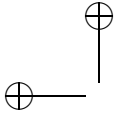
$$2 < b_{p,p+1}, \quad 2 < b_{p,p-1}, \quad a_{p+1,p+2} < b_{p+1,p+2}, \quad a_{p-2,p_1} < b_{p-2,p-1}.$$

Assume

$$a_{ij} < b_{ij}$$

for all edges. Then for every p , such that $(p-1, p)$ and $(p, p+1)$ are both standard, either both of $\|\mathbf{v}_p - \mathbf{v}_{p+1}\|$ and $\|\mathbf{v}_p - \mathbf{v}_{p-1}\|$ equal 2, or neither equals 2.

Proof This follows from Lemma 2.12 and the reversal symmetry of stable constraint systems. \square



Lemma 2.14 (propagate-min-std3) [JLXFDMJ] *Let s be a basic SCS, and let $\mathbf{v} \in \mathcal{M}_s$ be generic. Assume that $m(s) \leq 1$. Assume that for all diagonals $\{i, j\}$, we have $a_{ij}(s) \leq c_{stab} < \|\mathbf{v}_i - \mathbf{v}_j\|$ and $4h_0 < b_{ij}(s)$. Assume*

$$a_{ij} < b_{ij}$$

for all edges. Assume that $\|\mathbf{v}_i - \mathbf{v}_{i+1}\| = 2$ for some edge $\{i, i+1\}$. Then $\|\mathbf{v}_i - \mathbf{v}_j\| = 2$ at every edge not in $M(s)$.

Proof Let $\{j, j+1\}$ be any standard edge. We may create a chain avoiding $M(s)$ of the form

$$\{j, j+1\}, \{j+1, j+2\}, \dots, \{i-1, i\},$$

or

$$\{i+1, i+2\}, \{i+2, i+3\}, \dots, \{j, j+1\}.$$

Induction along the chain using Lemma 2.13 gives the result. \square

If s is a SCS, with diagonal $\{i, j\}$, let $D(s, \{i, j\})$ be the SCS which is identical to s , except that $b_{ij}(D(s, \{i, j\})) = c_{stab}$. In particular, we have a SCS $D(6I1, \{i, j\})$.

Lemma 2.15 [YRTAFYH] *If s is a basic SCS, $\{i, j\}$ is a diagonal, and $a_{ij}(s) \leq c_{stab}$, then $D(s, \{i, j\})$ is a basic SCS.*

Lemma 2.16 (equivalent diagonal) [WKEIDFT] *Let s be a basic SCS with $k(s) > 3$. Assume that the constants $a_{ij}(s)$ and $b_{ij}(s)$ are independent of the subscript for all edges $\{i, j\}$. Assume that the constants $a_{ij}(s)$ and $b_{ij}(s)$ are independent of the subscript for all diagonals. Assume $a_{ij}(s) \leq stab$ for all diagonals. Let $\{p, q\}$ and $\{p', q'\}$ be diagonals with $p' + q = p + q'$. Then*

$$\{D(s, \{p, q\})\} \Rightarrow \{D(s, \{p', q'\})\}.$$

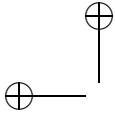
Proof This follows from the equivalence of the two SCSs. \square

2.4 hexagons

In this subsection, we treat the initial SCS 6I1.

Lemma 2.17 [PEDSLGV] *Let $\mathbf{v} \in \mathcal{M}_{6I1}$. If for some diagonal $\{i, j\}$, we have $\|\mathbf{v}_i - \mathbf{v}_j\| \leq c_{stab}$, then $\mathbf{v} \in \mathcal{M}_{D(6I1, \{i, j\})}$. If for every diagonal $\{i, j\}$, we have $c_{stab} < \|\mathbf{v}_i - \mathbf{v}_j\|$, then $\mathbf{v} \in \mathcal{M}_{6M1}$.*

Proof This is a direct consequence of Lemma 2.6. \square



Lemma 2.18 (hex slice) [AQICLXA] *We have*

$$D(6I1, \{0, 2\}) \Rightarrow \{5M1, 3M1\}.$$

Proof Slice $D(6I1, \{0, 2\})$ along the diagonal $\{0, 2\}$. \square

Lemma 2.19 (hex slice2) [FUNOUYH] *We have*

$$D(6I1, \{0, 3\}) \Rightarrow \{4M2\}.$$

Proof Slice $D(6I1, \{0, 3\})$ along the diagonal $\{0, 3\}$. The two halves are equivalent. \square

Lemma 2.20 (hex slice 3) [OEHDBEN] $\{6I1\} \Rightarrow \{6T1, 5M1, 4M2, 3M1\}$.

Proof Let $\mathbf{v} \in \mathcal{M}_{6I1}$.

Assume first that there is a diagonal $\{i, i+2\}$ such that $\|\mathbf{v}_i - \mathbf{v}_{i+2}\| \leq c_{\text{stab}}$. Then by the previous lemmas,

$$\{6I1\} \Rightarrow \{D(6I1, \{i, j\})\} \Rightarrow \{D(6I1, \{0, 2\})\} \Rightarrow \{5M1, 3M1\}.$$

Assume next that there is a diagonal $\{i, i+3\}$ such that $\|\mathbf{v}_i - \mathbf{v}_{i+3}\| \leq c_{\text{stab}}$. Then by the previous lemmas,

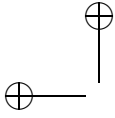
$$\{6I1\} \Rightarrow \{D(6I1, \{i, j\})\} \Rightarrow \{D(6I1, \{0, 3\})\} \Rightarrow \{4M2\}.$$

In a hexagon, every diagonal has the form $\{i, i+2\}$ or $\{i, i+3\}$. We may now assume that for every diagonal, $c_{\text{stab}} < \|\mathbf{v}_i - \mathbf{v}_j\|$. By Lemma 2.10, \mathbf{v} is generic. By Lemma 2.11, there exists an edge $\{i, j\}$ such that $\|\mathbf{v}_i - \mathbf{v}_j\| = 2$. By Lemma 2.14 [propagate-min-std3], all edges have length 2. Hence $\mathbf{v} \in \mathcal{B}_{6T1}$ and $\{6I1\} \Rightarrow \{6T1\}$. \square

In the section on pentagons, we make a further analysis of 5M1; and in the section on quadrilaterals, we analyze 4M2. This completes our analysis of hexagons.

2.5 pentagons

In this section, we treat the pentagonal cases: 5I1, 5I2, 5I3, 5M1.



Lemma 2.21 [OTMTOTJ]

$$\begin{aligned}
 \{5I1\} &\Rightarrow \{D(5I1, \{0, 2\}), 5M2\} \\
 \{5I2\} &\Rightarrow \{D(5I2, \{0, 2\}), 5M2\} \\
 \{5I3\} &\Rightarrow \{D(5M1, \{0, 2\}), D(5M1, \{0, 3\}), D(5M1, \{2, 4\}), 5M2\} \\
 \{5M1\} &\Rightarrow \{D(5M1, \{0, 2\}), D(5M1, \{0, 3\}), D(5M1, \{2, 4\}), 5M2\}
 \end{aligned}$$

Proof Let $s \in \{5I1, 5I2, 5I3, 5M1\}$. Let $\mathbf{v} \in \mathcal{M}_s$. If $\|\mathbf{v}_i - \mathbf{v}_j\| > c_{\text{stab}}$ for every diagonal $\{i, j\}$, then s transfers to $5M2$, and $\{s\} \Rightarrow \{5M2\}$.

For every diagonal, we have

$$\{D(5I3, \{i, j\})\} \Rightarrow \{D(5M1, \{i, j\})\}. \quad (2.22)$$

Now, let $\{i, j\}$ be a diagonal such that $\|\mathbf{v}_i - \mathbf{v}_j\| \leq c_{\text{stab}}$. Then $\{s\} \Rightarrow \{D(s, \{i, j\})\}$ and by equivalence $\{D(s, \{i, j\})\}$ and the arrow (2.22), we may pass to one of the SCSs enumerated in the lemma. \square

Lemma 2.23 [HIJQAHA]

$$\{5M2\} \Rightarrow \{3T1, 3T4, 4M6, 4M7, 4M8, 5T1, D(5I2, \{0, 2\}), D(5M1, \{0, 2\}), D(5M1, \{0, 3\}), D(5M1, \{2, 4\})\}$$

Proof Let $\mathbf{v} \in \mathcal{M}_{5M2}$. Assume first that $\|\mathbf{v}_i - \mathbf{v}_j\| \leq c_{\text{stab}}$ for some diagonal $\{i, j\}$. We consider subcases depending on the length of edges $\|\mathbf{v}_i - \mathbf{v}_j\|$. If $\|\mathbf{v}_0 - \mathbf{v}_1\| \leq 2h_0$, then

$$\{5M2\} \Rightarrow \{D(5I2, \{i, j\})\},$$

and otherwise if $2h_0 < \|\mathbf{v}_0 - \mathbf{v}_1\| \leq \sqrt{8}$, then

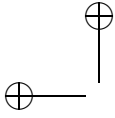
$$\{5M2\} \Rightarrow \{D(5I3, \{i, j\})\},$$

These SCSs transfer to one of the terms on the right-hand side in the lemma. In the final case $\sqrt{8} \leq \|\mathbf{v}_0 - \mathbf{v}_1\| \leq c_{\text{stab}}$. We can take the diagonal to be $\{0, 2\}$, $\{0, 3\}$, or $\{2, 4\}$. For these three diagonals, we have respectively,

$$\{5M2\} \Rightarrow \{D(5M2, \{i, j\})\} \Rightarrow \{4M6, 3T4\}, \quad \text{or } \{3T1, 4M7\}, \quad \text{or } \{3T1, 4M8\}.$$

We now assume $\|\mathbf{v}_i - \mathbf{v}_j\| > c_{\text{stab}}$ for all diagonals. By Lemma 2.10, \mathbf{v} is generic. By Lemma 2.11, there exists an edge $\{i, j\}$ such that $\|\mathbf{v}_i - \mathbf{v}_j\| = 2$. By Lemma 2.14, all standard edges of \mathbf{v} have length 2.

Index the vertices so that $\{0, 1\}$ is the non-standard edge. If either \mathbf{v}_0 or \mathbf{v}_1 is straight, then the deformation NUX implies that the edge $\{0, 1\}$ is as short as possible, length 2. If neither of the two vertices is straight, then we may apply deformation 684, to again conclude the edge $\{0, 1\}$ is as short as possible. In these cases $\{5M2\} \Rightarrow \{5T1\}$. \square



Lemma 2.24 [CNICGSF]

$$\{D(5I1, \{0, 2\})\} \Rightarrow \{4M2, 3M1\}$$

$$\{D(5I2, \{0, 2\})\} \Rightarrow \{4M3, 3T1\}$$

$$\{D(5M1, \{0, 2\})\} \Rightarrow \{4M2, 3T4\}$$

$$\{D(5M1, \{0, 3\})\} \Rightarrow \{4M4, 3M1\}$$

$$\{D(5M1, \{2, 4\})\} \Rightarrow \{4M5, 3M1\}$$

Proof In each case, we slice the given SCS along the given diagonal $\{i, j\}$ to produce a quadrilateral and a triangle. The resulting pieces (up to equivalence and transfer) are shown on the right. \square

2.6 quadrilaterals

Next we turn to the analysis of quadrilaterals.

Lemma 2.25 [GSXRFW] *Let s be an SCS with $k(s) = 4$. Let $\mathbf{v} \in \mathcal{M}_s$. Then \mathbf{v} is generic.*

Proof Assume $k = 4$. Recall $m(s) + k(s) \leq 6$ so that $m(s) \leq 2$. There is no pole, because of the polar test

$$\text{arc}(2, 2, 2h_0) + \text{arc}(2, 2, c_{\text{stab}}) < \pi.$$

\square

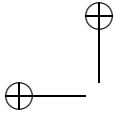
Lemma 2.26 [ARDBZYE]

$$\{4I2\} \Rightarrow \{4T1, 4T2\}.$$

Proof Let $\mathbf{v} \in \mathcal{M}_{4I2}$. By Lemma 2.25, \mathbf{v} is generic.

We consider the case when both diagonals of \mathbf{v} are greater than c_{stab} . By Lemma 2.11, there exists an edge $\{i, j\}$ such that $\|\mathbf{v}_i - \mathbf{v}_j\| = 2$. By Lemma 2.14, all edges have length 2. This gives $\{4I2\} \Rightarrow \{4T1\}$.

We turn to the case when some diagonal, say $\{1, 3\}$ has length in the range $[3, c_{\text{stab}}]$. Without loss of generality, we may assume that $\{1, 3\}$ is the shorter of the two diagonals. By Lemma ear_acute, a quadrilateral with diagonals at least 3 does not have any straight \mathbf{v}_i . If diagonal $\{0, 2\}$ has length 3 then since it is the longest, both diagonals are 3, and $\{4I2\} \Rightarrow \{4T2\}$. If the diagonal $\{0, 2\}$ has length greater than 3, we may triangulate the quadrilateral along the diagonal $\{1, 3\}$ and apply non-obtuse deformation 482 at any edge to show that it has length 2. Hence, $\{4I2\} \Rightarrow \{4T1\}$. \square



Lemma 2.27 [FYSSVEV]

$$\{4I1\} \Rightarrow \{4I2, D(4I1, \{0, 2\})\}$$

Proof Let $\mathbf{v} \in \mathcal{M}_{4I1}$. If \mathbf{v} has some diagonal $\{0, 2\}$ that is less than c_{stab} , then

$$\{4I1\} \Rightarrow \{D(4I1, \{i, j\})\} \Rightarrow \{D(4I1, \{0, 2\})\}.$$

Otherwise, with long diagonals, we have a transfer

$$\{4I1\} \Rightarrow \{4I2\}.$$

□

Lemma 2.28 [AUEAHEH]

$$\{D(4I1, \{0, 2\})\} \Rightarrow \{3M1\}.$$

Proof Slice along the diagonal $\{0, 2\}$.

□

Lemma 2.29 [ZNLLDL]

$$\{D(4I3, \{0, 2\})\} \Rightarrow \{4T4\}.$$

Proof In fact, $4T4$ is equivalent to $D(4I3, \{0, 2\})$.

□

Remark 2.30 The terminal case $4T4$ is treated by slicing along the diagonal $\{0, 2\}$ and introducing a nontrivial J .

Lemma 2.31 [VQFYMZY]

$$\{4I3\} \Rightarrow \{4T4, 4M6\}.$$

Proof Let $\mathbf{v} \in \mathcal{M}_{4I3}$. If there is a diagonal $\{i, j\}$ of length at most c_{stab} , and we have

$$\{4I3\} \Rightarrow \{D(4I3, \{i, j\})\} \Rightarrow \{D(4I3, \{0, 2\})\} \Rightarrow \{4T4\}.$$

If both diagonals are greater than c_{stab} , there is an equi-transfer to $4M6$.

□

Lemma 2.32 [BNAWVNH] $\{4M2\} \Rightarrow \{3M1, 3T4, 4M6\}$.

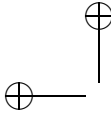
Proof If there exists $\mathbf{v} \in \mathcal{M}_{4M2}$ and a diagonal $\{i, j\}$ such that $\|\mathbf{v}_i - \mathbf{v}_j\| \leq c_{\text{stab}}$, then slice along the diagonal $\{i, j\}$ to obtain $\{4M2\} \Rightarrow \{3T5, 3T4\}$. Otherwise, $\{4M2\} \Rightarrow \{4M6\}$.

□

Lemma 2.33 [RAWZDIB] $\{4M3\} \Rightarrow \{3T1, 3T6, 4M6\}$.

Proof If there exists $\mathbf{v} \in \mathcal{M}_{4M3}$ and a diagonal $\{i, j\}$ such that $\|\mathbf{v}_i - \mathbf{v}_j\| \leq c_{\text{stab}}$, slice along the diagonal $\{i, j\}$ to obtain $\{4M3\} \Rightarrow \{3M1, 3T6\}$. Otherwise, $\{4M3\} \Rightarrow \{4M6\}$.

□



Lemma 2.34 [MFKLVDK] $\{4M4\} \Rightarrow \{3T4, 3M1, 3T3, 4M7\}$.

Proof If there exists $\mathbf{v} \in \mathcal{M}_{4M4}$ and a diagonal $\{i, j\}$ such that $\|\mathbf{v}_i - \mathbf{v}_j\| \leq c_{\text{stab}}$, slice along the diagonal $\{i, j\}$ to obtain $\{4M4\} \Rightarrow \{3T4, 3M1, 3T3\}$. Otherwise, $\{4M4\} \Rightarrow \{4M7\}$. \square

Lemma 2.35 [RYPDIXT] $\{4M5\} \Rightarrow \{3T4, 4M8\}$.

Proof If there exists $\mathbf{v} \in \mathcal{M}_{4M5}$ and a diagonal $\{i, j\}$ such that $\|\mathbf{v}_i - \mathbf{v}_j\| \leq c_{\text{stab}}$, slice along the diagonal $\{i, j\}$ to obtain $\{4M5\} \Rightarrow \{3T4\}$. Otherwise, $\{4M5\} \Rightarrow \{4M8\}$. \square

The quadrilateral cases that remain are 4M6, 4M7, and 4M8.

Lemma 2.36 [WGDHPPI] *Let s be an SCS with $k(s) = 4$. Let $\mathbf{v} \in \mathcal{M}_s$. Then the number of straight vertices \mathbf{v}_i is at most 1.*

Proof Lemma 2.25 implies that \mathbf{v} is generic. Lemma 2.2 gives the result. \square

Lemma 2.37 [ASSWPOW] *Let s be a basic SCS with $k(s) = 4$. Let $\mathbf{v} \in \mathcal{M}_s$. Let $p \in I(s)$. Assume that the constraints at $(p-1, p)$ and $(p, p+1)$ are both standard. Then $\text{rrv}(\mathbf{v}_{p-1}, \mathbf{v}_{p+1}) \leq 15.53$.*

Remark 2.38 The constant 15.53 appears in several of the computer calculations for nonlinear inequalities.

Proof This follows from the triangle inequality.

$$\text{arc}_V(\mathbf{0}, \mathbf{v}_{\ell-1}, \mathbf{v}_{\ell+1}) \leq 2 \text{arc}(2, 2, 2h_0) < \text{arc}(2, 2, \sqrt{15.53}).$$

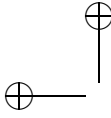
\square

Remark 2.39 (Running assumptions) The following lemmas will have several assumptions in common. We list them here as conditions on (s, v) .

1. s is an SCS
2. $k = 4$
3. $\mathbf{v} \in \mathcal{M}_s$.
4. For all diagonals, $a_{ij} \leq c_{\text{stab}} < \|\mathbf{v}_i - \mathbf{v}_j\|$ and $4h_0 < b_{ij}$.
5. The constraint on each edge is standard or pro^+ . (In particular m is the number of edges with pro^+ constraints.)

Lemma 2.40 [YEBWJNG] *Let (s, v) satisfy the running assumptions. If the constraint on $(p, p+1)$ is standard and if \mathbf{v} is not straight at p and $p+1$, then $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| = 2$.*

Proof This is a direct consequence of deformation 482. \square



Lemma 2.41 [TUAPYYU] *Let (s, v) satisfy the running assumptions. Assume that \mathbf{v} is not straight at p and $p + 1$. Then $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| = a_{p,p+1}$ or $b_{p,p+1}$.*

Proof Either \mathbf{v}_{p-1} or \mathbf{v}_{p+2} is not straight. Use deformation 1834 or deformation (1834 sym), as appropriate. \square

Lemma 2.42 [WKZZEEH] *Let (s, v) satisfy the running assumptions. Assume that \mathbf{v} is not straight at p , $p - 1$, and $p + 1$. Then $\|\mathbf{v}_p - \mathbf{v}_{p-1}\| < c_{stab}$ or $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| < c_{stab}$.*

Proof This is a restatement of deformation 684. \square

Lemma 2.43 [PWEIWBZ] *Let (s, v) satisfy the running assumptions. Then $\|\mathbf{v}_p - \mathbf{v}_{p+1}\| = 2$ for every standard edge.*

Proof Assume $(p, p + 1)$ standard. If \mathbf{v} is not straight at p and $p + 1$, then YEBWJNG gives the result. We may now assume without loss of generality that \mathbf{v} is straight at p or $p + 1$. Shifting indices, we may assume that \mathbf{v} is straight at p and that (p, q) is standard, where $q = p \pm 1$. For a contradiction, we may assume that $\|\mathbf{v}_p - \mathbf{v}_q\|$ is not minimal. In fact, by NUX, we may assume that no edge at p is minimal. By IMJ, $\|\mathbf{v}_p\| = 2$.

We review the conditions at $(i, j) = (p + 1, p + 2)$ and $(p + 2, p + 3)$, the two edges away from the straight angle p . By Lemma TUAPYYU, \mathbf{v} is extremal along edges $(p + 1, p + 2)$ and $(p + 2, p + 3)$. Hence along these two edges $\|\mathbf{v}_i - \mathbf{v}_j\|$ is 2, $2h_0$, or c_{stab} . By Lemma TUAPYYU, one of the two edges $\|\mathbf{v}_i - \mathbf{v}_j\| \leq 2h_0$. By Lemma YEBWJNG, if the constraint at (i, j) is standard, then $\|\mathbf{v}_i - \mathbf{v}_j\| = 2$.

If (p, j) is a standard edge at p , with $j = p \pm 2$, then $(j, p + 2)$ is pro^+ . Otherwise, by earlier remarks $\|\mathbf{v}_j - \mathbf{v}_{p+2}\| = 2$, which propagates to make $\|\mathbf{v}_p - \mathbf{v}_j\| = 2$ by TFITSKC, which is contrary to what we have.

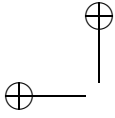
We consider four subcases:

1. $(p - 1, p)$ and $(p, p + 1)$ are both standard. Also, $\|\mathbf{v}_{p+1} - \mathbf{v}_{p+1}\| \leq 2h_0$.
2. $(p - 1, p)$ and $(p, p + 1)$ are both standard. Also, $\|\mathbf{v}_{p+2} - \mathbf{v}_{p-1}\| \leq 2h_0$.
3. $(p, p + 1)$ is standard, $(p - 1, p)$ is pro^+ .
4. $(p, p - 1)$ is standard, $(p, p + 1)$ is pro^+ .

If (p, j) is a standard edge at p , with $j = p \pm 1$, then the angle at j is acute. Otherwise, deformation 482 obtuse makes (p, j) minimal. We use a computer calculation to contradict this fact.

In the first two cases, by *computer calculation*² [2] we may use LEMMA_PWE1 and LEMMA_PWE2, respectively. Note that both edges at $p + 2$ are pro^+ .

² [6184614449, 1348932091, 1348932091 delta, 5557288534, 5557288534 delta]



In the third case, the edge $(p + 1, p + 2)$ is pro^+ , and the edge $(p + 2, p + 3)$ is standard (because $m \leq 2$) of length 2. Use `LEMMA_PWE3`, which is based on a *computer calculation*³ [2].

The fourth case is symmetrical with the third, and uses `LEMMA_PWE4`. \square

Lemma 2.44 [VASYYAU] *Let (s, v) satisfy the running assumptions. Let \mathbf{v} be straight at p . Then both edges at p are minimal.*

Proof If there exists a standard constraint at some edge at p , then by `PWEIBZ`, it is minimal. By `NUX`, the other edge at p is also minimal. In this case we are done.

If, on the other hand, there does not exist a standard constraint at some edge at p , then since $m \leq 2$, both edges at $p + 2$ are standard. The edges are \mathbf{v} -minimal and hence have length 2. The two edges at p are pro^+ and hence have length $\geq 2h_0$. Also, by `ODX`, `IMJ`, since some edge at p is not minimal, we have $\|\mathbf{v}_p\| = 2$.

This violates the triangle inequality:

$$\begin{aligned} \text{arc}(2, \|\mathbf{v}_{p-1}\|, 2) + \text{arc}(2, \|\mathbf{v}_{p+1}\|, 2) &< \text{arc}(2, \|\mathbf{v}_{p-1}\|, 2h_0) + \text{arc}(2, \|\mathbf{v}_{p+1}\|, 2h_0) \\ &\leq \text{arc}_V(\mathbf{0}, \{\mathbf{v}_{p-1}, \mathbf{v}_p\}) + \text{arc}_V(\mathbf{0}, \{\mathbf{v}_{p+1}, \mathbf{v}_p\}) \\ &= \text{arc}_V(\mathbf{0}, \{\mathbf{v}_{p-1}, \mathbf{v}_{p+1}\}) \\ &\leq \text{arc}_V(\mathbf{0}, \{\mathbf{v}_{p-1}, \mathbf{v}_{p+2}\}) + \text{arc}_V(\mathbf{0}, \{\mathbf{v}_{p+2}, \mathbf{v}_{p+1}\}) \\ &\leq \text{arc}(2, \|\mathbf{v}_{p-1}\|, 2) + \text{arc}(2, \|\mathbf{v}_{p+1}\|, 2). \end{aligned}$$

\square

Lemma 2.45 [NWDGKXH]

$$\{4M6\} \Rightarrow \{4T3, 4T5\}.$$

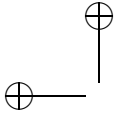
Proof Let $\mathbf{v} \in \mathcal{M}_{4M6}$. By the definition of `4M6`, the diagonals are at least c_{stab} . If some diagonal has $\|\mathbf{v}_i - \mathbf{v}_j\| = c_{\text{stab}}$, then

$$\{4M6\} \Rightarrow \{4T5\}.$$

Otherwise, by the preceding lemmas, all edges are extremal, and all standard edges have length 2.

By `Lemma_quad_nonexist_849`, the nonstandard edge does not have $\|\mathbf{v}_i - \mathbf{v}_j\| = 2h_0$. Since the non-standard edge is extremal, this forces $\|\mathbf{v}_i - \mathbf{v}_j\| = c_{\text{stab}}$. In this case $\{4M6\} \Rightarrow \{4T3\}$. \square

³ [2073661826, 8405387449, 9368433105, 5550839403, 5550839403 delta]



Lemma 2.46 [EFLYGAU] *Let $\mathbf{v} \in \mathcal{M}_{4M7}$. If both diagonals of \mathbf{v} are greater than c_{stab} , then*

$$\{4M7\} \Rightarrow \{4M6\}.$$

Proof By preceding lemmas, the nonstandard edges of \mathbf{v} is extremal.

If one of those edges is minimal $2h_0$, then \mathbf{v} is in the domain of \mathcal{M}_{4M6} and it transfers.

We prove in fact, that if both standard edges are maximal, then we obtain a contradiction. Suppose for a contradiction, the two nonstandard edges are adjacent and have length c_{stab} . By Lemma VASYAU, neither nonstandard edge is adjacent to a vertex that is straight. By WKZZEEH, one of the nonstandard edges is minimal. earlier lemma has ruled out adjacent edges of length c_{stab} . \square

Lemma 2.47 [YOBIMPP] *Let $\mathbf{v} \in \mathcal{M}_{4M7}$. If some diagonal has length c_{stab} , then*

$$\{4M7\} \Rightarrow \{3M1, 3T3, 3T4\}$$

Proof Let $\{i, i+2\}$ be the diagonal such that $\|\mathbf{v}_i - \mathbf{v}_{i+2}\| = c_{stab}$.

Assume first that the diagonal is $\{0, 2\}$. Slice along the diagonal to obtain 3M1 and 3T3. Then $\{4M7\} \Rightarrow \{3M1, 3T3\}$.

In the remaining case, the diagonal is $\{1, 3\}$. Slice along the diagonal to obtain two equivalent pieces that transfer to 3T4. \square

Lemma 2.48 [BJTDWPS] *Let $\mathbf{v} \in \mathcal{M}_{4M8}$. If both diagonals of \mathbf{v} are greater than c_{stab} , then*

$$\{4M8\} \Rightarrow \{4M6, 3T7\}.$$

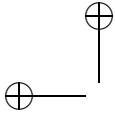
Proof By preceding lemmas, the non-standard edges of \mathbf{v} is extremal. If one of those edges is minimal $2h_0$, then \mathbf{v} is in the domain of \mathcal{M}_{4M6} and it transfers.

Otherwise, both non-standard edges have length c_{stab} . Every non-standard edge is adjacent to a standard edge. By Lemma PWEI WBZ, these standard edges are minimal; that is, they have length 2. By Lemma quad_diag_362, the shortest diagonal has length at most 3.62. Slice along that diagonal, to obtain two pieces like 3T7. Then $\{4M7\} \Rightarrow \{3T7\}$. \square

Lemma 2.49 [MIQMCSN] *Let $\mathbf{v} \in \mathcal{M}_{4M8}$. If some diagonal has length precisely c_{stab} , then*

$$\{4M8\} \Rightarrow \{3T4\}$$

Proof Slice along the diagonal to obtain two equivalent pieces that transfer to 3T4. (Compare the proof of RYPDIXT.) \square



This completes the analysis of the cases 4M6, 4M7, and 4M8. All quad cases have been resolved into terminal cases.

2.7 triangles

Lemma 2.50 [LFLACKU] [formal proof by TCH]. $\{3I1\} \Rightarrow \{3T2, 3T5\}$

Proof By deformation 1834, we may assume that each edge $\{i, j\}$ is extremal, hence of length 2 or $2h_0$. If all edges have length 2, then

$$\{3I1\} \Rightarrow \{3T2\}.$$

Otherwise,

$$\{3I1\} \Rightarrow \{3T5\}.$$

□

Lemma 2.51 [BKOSSGE] [formal proof by TCH]. $\{3M1\} \Rightarrow \{3T1, 3T5\}$.

Lemma 2.52 [OCBICBY] [formal proof by T. Hales].

$$\mathcal{M}_s = \emptyset,$$

for all $s \in S_{\text{term}}$.

Proof When $s \in S_{\text{term}}$ with $k(s) \leq 4$, by a *computer calculation*⁴ [2], we show that $\tau^*(s, \mathbf{v}) \geq 0$ for all $\mathbf{v} \in \mathcal{B}_s$. This implies for such s that $\mathcal{M}_s = \emptyset$. The two cases 5T1 and 6T1 remain. They are treated in [1][Section 7.4.5], where it is shown that $\mathcal{M}_s = \emptyset$ for $s = 5T1$ and $6T1$. This shows that $\mathcal{M}_s = \emptyset$ for all $s \in S_{\text{term}}$. □

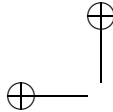
2.8 unsorted lemmas

We dump a few unsorted lemmas at the end of this appendix.

Lemma 2.53 [TECOXBM] [formal proof by Hoang Le Truong]. *Let s be a SCS, and let $\mathbf{v} \in \mathcal{B}_s$. Let $\mathbf{u}, \mathbf{w} \in V_{\mathbf{v}}$ satisfy $2 \leq \|\mathbf{u} - \mathbf{w}\| \leq c_{\text{stab}}$ where $\{\mathbf{u}, \mathbf{w}\} \notin E_{\mathbf{v}}$. Then \mathbf{u} and \mathbf{w} are nonparallel. Moreover, $C^0\{\mathbf{u}, \mathbf{w}\} \subset W_{\text{dart}}^0(x)$ for all $x \in F$.*

Proof This is a repetition of [1][Lemma TECOXBM]. □

⁴ [various inequalities]



Definition 2.54 (XR) [LIHVJNG] We let

$$XR(y_1, y_2, y_6) = 8(1 - \frac{y_1^2 + y_2^2 - y_6^2}{2y_1y_2}).$$

Lemma 2.55 [KPIDBQH] [formal proof by Hales].

$$\text{dih}_y(y_1, y_2, \dots, y_6) = \text{dih}_x(4, 4, 4, XR(y_2, y_3, y_4), XR(y_1, y_2, y_6), XR(y_1, y_3, y_5)).$$

Proof Compute both sides. \square

Lemma 2.56 [DRNDRDV] [formal proof by Hales].

$$\frac{\partial XR(y_1, y_2, y_6)}{\partial y_6} = \frac{8y_6}{y_1y_2}.$$

Lemma 2.57 [TBRMXRZ] [formal proof by Hales]. *Let $h(x) = g(f(x))$ and $f(x) = y$. Assume $f'(x) > 0$. Then under appropriate differentiability conditions:*

1. $g'(y) \sim h'(x)$.
2. $h'(x) = 0$ implies that $h''(x) \sim g''(y)$.

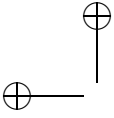
Proof

$$h'(x) = g'(y)f'(x) \sim g'(y).$$

Also, if $h'(x) = 0$, then $g'(y) = 0$ and

$$h''(x) = g''(y)(f'(x))^2 + g'(y)f''(x) = g''(y)(f'(x))^2 \sim g''(y).$$

\square



3

Appendix on explicit deformations

This is an appendix to Section 7.2 of Dense Sphere Packings. This appendix gives the explicit construction of particular deformations.

The first lemma constructs a simplex $\{\mathbf{v}_0, \mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3\} \subset \mathbb{R}^3$ on a given base triangle $\{\mathbf{v}_0, \mathbf{v}_1, \mathbf{v}_2\}$. Our intention is to define a new constant to equal the right-hand side of Equation 3.2. The variable x_5 will run over an interval to define a continuous deformation of a local fan (V, E, F) .

Lemma 3.1 [PQCSXWG] *Let $\{\mathbf{v}_0, \mathbf{v}_1, \mathbf{v}_2\} \subset \mathbb{R}^3$. Assume that $\{\mathbf{v}_0, \mathbf{v}_1, \mathbf{v}_2\}$ is not collinear. Let $x_1, \dots, x_6 \in \mathbb{R}^3$ be given with $x_i > 0$ and*

$$x_1 = \|\mathbf{v}_1 - \mathbf{v}_0\|^2, \quad x_2 = \|\mathbf{v}_2 - \mathbf{v}_0\|^2, \quad x_6 = \|\mathbf{v}_1 - \mathbf{v}_2\|^2.$$

Assume that $\Delta(x_1, \dots, x_6) > 0$ (that is, $\Delta(\mathbf{x}) > 0$). Then there exists \mathbf{v}_3 such that

- $x_3 = \|\mathbf{v}_3 - \mathbf{v}_0\|^2, \quad x_5 = \|\mathbf{v}_3 - \mathbf{v}_1\|^2, \quad x_4 = \|\mathbf{v}_3 - \mathbf{v}_2\|^2$, and
- $(\mathbf{v}_1 - \mathbf{v}_0) \cdot (\mathbf{v}_2 - \mathbf{v}_0) \times (\mathbf{v}_3 - \mathbf{v}_0) > 0$.

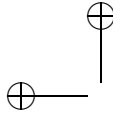
Explicitly, the following vector works:

$$\mathbf{v}_3 = \mathbf{v}_0 + \frac{2\sqrt{\Delta}(\mathbf{v}_1 - \mathbf{v}_0) \times (\mathbf{v}_2 - \mathbf{v}_0) + \Delta_5(\mathbf{v}_1 - \mathbf{v}_0) + \Delta_4(\mathbf{v}_2 - \mathbf{v}_0)}{v(x_1, x_2, x_6)}, \quad (3.2)$$

with subscripts on Δ indicating partial derivatives and $v = \text{ups}_{\mathbf{x}}$, the epsilon function. Moreover, fixing $\mathbf{v}_0, \mathbf{v}_1, \mathbf{v}_2$ and fixing all the variables x_i except x_5 , the vector $\mathbf{v}_3 \in \mathbb{R}^3$ in Equation 3.2 depends continuously on x_5 on the domain

$$\{x_5 : x_5 > 0, \quad \Delta(x_1, \dots, x_6) > 0\}.$$

Remark 3.3 There is a symmetry to the data fixing $\mathbf{v}_0, x_3, x_6, v(x_1, x_2, x_6)$ and swapping $\mathbf{v}_1 \leftrightarrow \mathbf{v}_2, x_1 \leftrightarrow x_2, x_4 \leftrightarrow x_5, \Delta_5 \leftrightarrow \Delta_4$. Under this symmetry, the vector \mathbf{v}_3 is given by the same formula, except that $\sqrt{\Delta}$ is replaced with $-\sqrt{\Delta}$, and the sign of the triple product is reversed.



Proof It can be shown by direct computation that the vector \mathbf{v}_3 works. This proof gives details about how \mathbf{v}_3 is found.

Without loss of generality, we may move \mathbf{v}_0 to the origin. Explicitly, we let $\mathbf{w}_i = \mathbf{v}_i - \mathbf{v}_0$. We will construct a unique \mathbf{w}_3 (for $\mathbf{w}_0 = 0$) and then set $\mathbf{v}_3 = \mathbf{v}_0 + \mathbf{w}_3$.

Note that the non-collinearity condition on $\{\mathbf{v}_0, \mathbf{v}_1, \mathbf{v}_2\}$ gives $\mathbf{w}_1 \times \mathbf{w}_2 \neq 0$. We may write \mathbf{w}_3 in terms of the basis $\mathbf{w}_1 \times \mathbf{w}_2$, \mathbf{w}_1 , and \mathbf{w}_2 :

$$\mathbf{w}_3 = \alpha(\mathbf{w}_1 \times \mathbf{w}_2) + \beta\mathbf{w}_1 + \gamma\mathbf{w}_2.$$

We compute the norms of \mathbf{w}_3 , $\mathbf{w}_3 - \mathbf{w}_1$ and $\mathbf{w}_3 - \mathbf{w}_2$ in this basis:

$$\begin{aligned} \|\mathbf{w}_3\|^2 &= x_3 = \alpha^2 \|\mathbf{w}_1 \times \mathbf{w}_2\|^2 + \|\beta\mathbf{w}_1 + \gamma\mathbf{w}_2\|^2 \\ \|\mathbf{w}_3 - \mathbf{w}_1\|^2 &= x_5 = \alpha^2 \|\mathbf{w}_1 \times \mathbf{w}_2\|^2 + \|(\beta - 1)\mathbf{w}_1 + \gamma\mathbf{w}_2\|^2 \\ \|\mathbf{w}_3 - \mathbf{w}_2\|^2 &= x_4 = \alpha^2 \|\mathbf{w}_1 \times \mathbf{w}_2\|^2 + \|\beta\mathbf{w}_1 + (\gamma - 1)\mathbf{w}_2\|^2. \end{aligned}$$

We eliminate α^2 from the equations and write the resulting equations for β and γ as a linear system:

$$\begin{pmatrix} \|\mathbf{w}_1\|^2 & \mathbf{w}_1 \cdot \mathbf{w}_2 \\ \mathbf{w}_1 \cdot \mathbf{w}_2 & \|\mathbf{w}_2\|^2 \end{pmatrix} \begin{pmatrix} \beta \\ \gamma \end{pmatrix} = \begin{pmatrix} (x_1 + x_3 - x_5)/2 \\ (x_2 + x_3 - x_4)/2 \end{pmatrix}.$$

The determinant of the system is

$$\|\mathbf{w}_1\|^2 \|\mathbf{w}_2\|^2 - (\mathbf{w}_1 \cdot \mathbf{w}_2)^2 = \|\mathbf{w}_1 \times \mathbf{w}_2\|^2 > 0.$$

By the law of sines or cosines (Lemma 2.59), this determinant can also be written in terms of x_i :

$$x_1 x_2 - ((x_6 - x_1 - x_2)/2)^2 = v(x_1, x_2, x_6)/4 > 0.$$

Thus, there are unique solutions β and γ as functions of the variables x_i . Explicitly,

$$\beta = \Delta_5/v(x_1, x_2, x_6), \quad \gamma = \Delta_4/v(x_1, x_2, x_6),$$

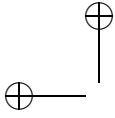
where $\Delta_i = \Delta_i(x_1, \dots, x_6)$ is the i th partial derivative of Δ .

With β and γ in hand, we solve for α using the equation:

$$x_3 - \|\beta\mathbf{w}_1 + \gamma\mathbf{w}_2\|^2 = \alpha^2 \|\mathbf{w}_1 \times \mathbf{w}_2\|^2.$$

The left-hand side of this equation, expressed in terms of the variables x_i is precisely $\Delta/v(x_1, x_2, x_6) > 0$. Hence

$$\alpha = \pm 2 \frac{\sqrt{\Delta(x_1, \dots, x_6)}}{v(x_1, x_2, x_6)}.$$



The triple product is

$$((\mathbf{v}_1 - \mathbf{v}_0) \times (\mathbf{v}_2 - \mathbf{v}_0)) \cdot (\mathbf{v}_3 - \mathbf{v}_0) = (\mathbf{w}_1 \times \mathbf{w}_2) \cdot \mathbf{w}_3 = \alpha |\mathbf{w}_1 \times \mathbf{w}_2|^2.$$

This is positive exactly when α is chosen to be positive. This shows the existence of a unique vector \mathbf{w}_3 subject to the given conditions.

Continuity follows from the continuous dependence of α, β , and γ on x_5 . In fact, β and γ are polynomials in x_5 and α requires the extraction of a square root of a positive polynomial in x_5 . \square

We will need a second lemma for deformations that occur within a plane. Our intention is to define a new constant to equal the right-hand side of Equation 3.5. It will be used in two contexts. Sometimes, the variable x_3 will run over an interval to define a continuous deformation of a local fan (V, E, F) . At other times, the vector \mathbf{v}_2 will vary in a continuous deformation constructed in the previous lemma, and \mathbf{v}_3 will be carried continuously along in the plane of $\{\mathbf{v}_0, \mathbf{v}_1, \mathbf{v}_2(t)\}$ by the construction of this lemma.

Lemma 3.4 [EYYPQDW] [formal proof by Hoang Le Truong]. *Let $\{\mathbf{v}_0, \mathbf{v}_1, \mathbf{v}_2\} \subset \mathbb{R}^3$. Assume that $\{\mathbf{v}_0, \mathbf{v}_1, \mathbf{v}_2\}$ is not collinear. Let $x_1, x_2, x_3, x_5, x_6 \in \mathbb{R}^3$ be given with $x_i > 0$ and*

$$x_1 = \|\mathbf{v}_1 - \mathbf{v}_0\|^2, \quad x_2 = \|\mathbf{v}_2 - \mathbf{v}_0\|^2, \quad x_6 = \|\mathbf{v}_1 - \mathbf{v}_2\|^2.$$

Assume that $v(x_1, x_3, x_5) > 0$. Let $\sigma \in \{\pm 1\}$ be a choice of sign. Then there exists \mathbf{v}_3 such that

- $\{\mathbf{v}_0, \mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3\}$ is coplanar.
- $x_3 = \|\mathbf{v}_3 - \mathbf{v}_0\|^2$, $x_5 = \|\mathbf{v}_3 - \mathbf{v}_1\|^2$, and
- $(\mathbf{v}_3 - \mathbf{v}_0) \times (\mathbf{v}_1 - \mathbf{v}_0)$ is a positive scalar times $\sigma(\mathbf{v}_1 - \mathbf{v}_0) \times (\mathbf{v}_2 - \mathbf{v}_0)$.

Explicitly, the following vector works:

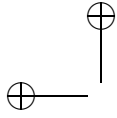
$$\mathbf{v}_3 = \mathbf{v}_0 + \frac{x_1 + x_3 - x_5}{2x_1}(\mathbf{v}_1 - \mathbf{v}_0) + \frac{\sigma}{x_1} \sqrt{\frac{v(x_1, x_3, x_5)}{v(x_1, x_2, x_6)}}(\mathbf{v}_1 - \mathbf{v}_0) \times ((\mathbf{v}_1 - \mathbf{v}_0) \times (\mathbf{v}_2 - \mathbf{v}_0)). \quad (3.5)$$

Moreover, fixing $\mathbf{v}_0, \mathbf{v}_1, \mathbf{v}_2$ and fixing all the variables x_i except x_3 , the vector $\mathbf{v}_3 \in \mathbb{R}^3$ depends continuously on x_3 on the domain

$$\{x_3 : x_3 > 0, \quad v(x_1, x_3, x_5) > 0\}.$$

Moreover, fixing $\mathbf{v}_0, \mathbf{v}_1$ and fixing all the variables x_i , the vector $\mathbf{v}_3 \in \mathbb{R}^3$ depends continuously on \mathbf{v}_2 on the domain

$$\{\mathbf{v}_2 \in \mathbb{R}^3 : \mathbf{v}_2 \text{ is not on the line through } \mathbf{v}_0 \text{ and } \mathbf{v}_1\}.$$



Remark 3.6 There is a symmetry in the data $\mathbf{v}_0 \leftrightarrow \mathbf{v}_1$, $x_3 \leftrightarrow x_5$, $x_2 \leftrightarrow x_6$ fixing \mathbf{v}_2 and x_1 . The symmetry preserves the solution \mathbf{v}_3 .

Proof Without loss of generality, we may move \mathbf{v}_0 to the origin. Explicitly, we let $\mathbf{w}_i = \mathbf{v}_i - \mathbf{v}_0$. We will construct a unique \mathbf{w}_3 (for $\mathbf{w}_0 = 0$) and then set $\mathbf{v}_3 = \mathbf{v}_0 + \mathbf{w}_3$.

Let $\mathbf{n} = \mathbf{w}_1 \times (\mathbf{w}_1 \times \mathbf{w}_2)$. By the non-collinearity assumption and the positivity of x_i , we have $\mathbf{n} \neq 0$. In fact, \mathbf{w}_1 and \mathbf{n} are orthogonal and span $\text{aff}\{\mathbf{0}, \mathbf{w}_1, \mathbf{w}_2\}$. The norm of \mathbf{n} is computed as in the previous lemma by the law of cosines:

$$\|\mathbf{n}\|^2 = \|\mathbf{w}_1\|^2 \|\mathbf{w}_1 \times \mathbf{w}_2\|^2 = x_1 v(x_1, x_2, x_6)/4 > 0.$$

We solve for \mathbf{w}_3 as a combination of \mathbf{w}_1 and \mathbf{n} :

$$\mathbf{w}_3 = \alpha \mathbf{w}_1 + \beta \mathbf{n}.$$

The norms of \mathbf{w}_3 and $\mathbf{w}_3 - \mathbf{w}_1$ are computed as

$$\begin{aligned} \|\mathbf{w}_3\|^2 &= x_3 = \alpha^2 \|\mathbf{w}_1\|^2 + \beta^2 \|\mathbf{n}\|^2 \\ \|\mathbf{w}_3 - \mathbf{w}_1\|^2 &= x_5 = \alpha^2 \|\mathbf{w}_1\|^2 + \beta^2 \|\mathbf{n}\|^2 \end{aligned}$$

Eliminate β and solve the resulting linear equation uniquely for α :

$$\alpha = (x_1 + x_3 - x_5)/(2x_1).$$

The right-hand side of $\beta^2 \|\mathbf{n}\|^2 = x_3 - \alpha^2 x_1$, expressed in terms of the variables x_i is $v(x_1, x_3, x_5)/(4x_1)$. Hence,

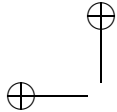
$$\beta = \pm \frac{1}{x_1} \sqrt{\frac{v(x_1, x_3, x_5)}{v(x_1, x_2, x_6)}}.$$

To compute the sign of β , we examine the cross-product condition.

$$\begin{aligned} \mathbf{w}_3 \times \mathbf{w}_1 &= (\alpha \mathbf{w}_1 + \beta \mathbf{n}) \times \mathbf{w}_1 \\ &= \beta (\mathbf{w}_1 \times (\mathbf{w}_1 \times \mathbf{w}_2)) \times \mathbf{w}_1 \\ &= \beta (\mathbf{w}_1 \times \mathbf{w}_2) (\mathbf{w}_1 \cdot \mathbf{w}_1) \\ &= \beta x_1 (\mathbf{w}_1 \times \mathbf{w}_2). \end{aligned}$$

This is a positive multiple of $\sigma \mathbf{w}_1 \times \mathbf{w}_2$ when β has sign σ . This establishes the unique existence of \mathbf{w}_3 .

Continuity follows from the explicit formulas for α and β . \square



4

Appendix on deformations of local fans

This is an appendix to Section 7.2 of Dense Sphere Packings. This appendix gives further details about the proof of the wedge property in Lemma 7.25 (ZLZTHIC).

Lemma 4.1 [WEECNNS] [formal proof by J. Harrison]. *The function dih_V is continuous on the set*

$$\{(\mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3, \mathbf{v}_4) : \text{not collinear } \{\mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3\} \text{ and not collinear } \{\mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_4\}\}.$$

Lemma 4.2 [XBJRPHC] *The function azim is continuous on the set*

$$\{(\mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3, \mathbf{v}_4) : \text{not collinear } \{\mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3\} \text{ and not collinear } \{\mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_4\} \mathbf{v}_4 \notin \text{aff}_+^0(\{\mathbf{v}_1, \mathbf{v}_2\}, \mathbf{v}_3)\}.$$

The following lemma rewrites the generic blade condition as a union of two sets that are manifestly open.

Lemma 4.3 (generic_alt) [IHWVUIZ] [formal proof by T. Hales]. *Assume that $\{0, \mathbf{v}, \mathbf{w}\}$ is not collinear and that $\mathbf{u} \neq \mathbf{0}$. Then the generic blade condition holds*

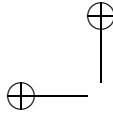
$$\text{aff}_+(\mathbf{0}, \{\mathbf{v}, \mathbf{w}\}) \cap \text{aff}_-^0(\{\mathbf{0}\}, \{\mathbf{u}\}) = \emptyset$$

iff $\{\mathbf{0}, \mathbf{u}, \mathbf{v}, \mathbf{w}\}$ is not coplanar or $-\mathbf{u} \in W^0(\mathbf{0}, \mathbf{v} \times \mathbf{w}, \mathbf{w}, \mathbf{v})$.

Lemma 4.4 [NHCXLRV] *Let (V, E, F) be a generic convex local fan. Let (φ, I, V) be a deformation of the fan. Let $\mathbf{v}, \mathbf{w} \in V$. Assume that $\mathbf{v} \in W_{\text{dart}}^0(\mathbf{w}, \rho\mathbf{w})$. Then for sufficiently small t , we have*

$$\mathbf{v}(t) \in W_{\text{dart}}^0(\mathbf{w}(t), \rho\mathbf{w}(t)).$$

Proof We claim that (\mathbf{v}, \mathbf{w}) is not a pole of the fan. In fact, by WEDGE_ALT, if a pole, then $\mathbf{v} \notin W_{\text{dart}}^0(\mathbf{w}, \rho\mathbf{w})$.



The condition $\mathbf{v} \in W_{\text{dart}}^0(\mathbf{w}, \rho\mathbf{w})$ gives

$$0 < \text{azim}(\mathbf{0}, \mathbf{w}, \rho\mathbf{w}, \mathbf{v}) < \text{azim}(\mathbf{0}, \mathbf{w}, \rho\mathbf{w}, \rho^{-1}\mathbf{w}).$$

Pick c such that

$$0 < \text{azim}(\mathbf{0}, \mathbf{w}, \rho\mathbf{w}, \mathbf{v}) < c < \text{azim}(\mathbf{0}, \mathbf{w}, \rho\mathbf{w}, \rho^{-1}\mathbf{w}). \quad (4.5)$$

The pole claim gives that the domain conditions for the continuity of azim are met. Therefore Equation (4.5) holds under a small deformation. This gives the conclusion. \square

Lemma 4.6 [WNWSHJT] *Let (V, E, F) be generic fan. Let (φ, I, V) be a deformation of the fan. Let $\mathbf{u} \in V$ have interior angle $\angle(\mathbf{u}) < \pi$. Then for sufficiently small t , we have $\angle(\mathbf{u}(t)) < \pi$.*

Proof By INTERIOR_ANGLE1_POS, we have $0 < \angle(\mathbf{u}) < \pi$. The continuity domain conditions for azim are met at $(\mathbf{0}, \mathbf{u}, \rho\mathbf{u}, \rho^{-1}\mathbf{u})$. Hence

$$0 < \text{azim}(\mathbf{u}(t)) < \pi$$

for t sufficiently small. \square

The following case is not needed, because it is a special case of MHAHEYJN.

Lemma 4.7 [CREUXCQ] *Let (V, E, F) be a generic fan. with pole (\mathbf{v}, \mathbf{w}) . Assume that the interior angle at the pole is less than π . Let $\mathbf{u} \neq \mathbf{v}, \mathbf{w}$. Let (φ, I, V) be a deformation of the fan that moves a single $\mathbf{u} \in V \setminus \{\mathbf{v}, \mathbf{w}\}$ and such that $\mathbf{u} \in V$, we have $\mathbf{u}(t) \in \text{aff}\{\mathbf{0}, \mathbf{v}, \mathbf{w}, \mathbf{u}\}$. Fix $\mathbf{u}' \in V$ such that $\angle(\mathbf{u}') = \pi$. Then for sufficiently small t , we have $\angle(\mathbf{u}'(t)) \leq \pi$.*

Proof Let $\mathbf{u}_1 = \mathbf{u}'$, $\mathbf{u}_2 = \rho\mathbf{u}'$, and $\mathbf{u}_0 = \rho^{-1}\mathbf{u}'$. Let $S = \{\mathbf{0}, \mathbf{u}_0, \mathbf{u}_1, \mathbf{u}_2, \mathbf{v}, \mathbf{w}\}$.

We consider the case $\mathbf{u} \notin S$. Then the set S is fixed under the deformation, so that

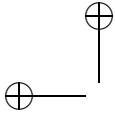
$$\angle(\mathbf{u}_1(t)) = \angle(\mathbf{u}'(t)) = \angle(\mathbf{u}') = \pi.$$

In the remaining case $\mathbf{u} \in S$. By the lunar geometry lemma, $S \subset A$, where $A = \text{aff}\{\mathbf{0}, \mathbf{v}, \mathbf{w}, \mathbf{u}\}$, a plane. The deformation moves \mathbf{u} within this plane, so $S(t) \subset A$. This implies that $S(t)$ is coplanar, which gives

$$\angle(\mathbf{u}'(t)) = \text{azim}(\mathbf{0}, \mathbf{u}_0(t), \mathbf{u}_1(t), \mathbf{u}_2(t)) \in \{0, \pi\}.$$

We claim the the domain conditions for the continuity of azim are met. Hence by continuity, $\angle(\mathbf{u}'(t)) = \pi$ for small π . \square

In the next series of lemmas we introduce an assumption ECAU given as follows.



Remark 4.8 (ECAU) Let $\mathbf{u}_0, \dots, \mathbf{u}_r \in \mathbb{R}^3$, and define $\mathbf{e} = \mathbf{u}_0 \times \mathbf{u}_1$. We make the following assumptions.

1. (C) The set $\{\mathbf{0}, \mathbf{u}_i, \mathbf{u}_j\}$ is not collinear for $i \neq j$.
2. (A) For all $i \geq 1$, we have $\mathbf{u}_i \in \text{aff}_+^0(\{\mathbf{0}, \mathbf{u}_0\}, \mathbf{u}_1)$.
3. (U) The set $U = \{\mathbf{u}_0, \dots, \mathbf{u}_r\}$ is cyclic with respect to $(\mathbf{0}, \mathbf{e})$, with cycle

$$\sigma \mathbf{u}_i = \mathbf{u}_{i+1}, \text{ if } i < r.$$

We let A be the plane $\text{aff}\{\mathbf{0}, \mathbf{u}_0, \mathbf{u}_1\}$.

Lemma 4.9 [VUYCADE] [formal proof by Hales]. *Let $\mathbf{u}_1, \dots, \mathbf{u}_r$ with cross product \mathbf{e} satisfy the conditions ECAU. Then for all $i \leq r$, we have constants t_0 and t_1 , with $t_1 > 0$ such that*

$$\mathbf{u}_i = t_0 \mathbf{u}_0 + t_1 \mathbf{u}_1, \quad \mathbf{e} \cdot \mathbf{u}_i = 0, \quad \text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_0, \mathbf{u}_i) < \pi.$$

Also, $\{\mathbf{0}, \mathbf{e}, \mathbf{u}_i\}$ is not collinear.

Proof Condition A gives $\mathbf{u}_i = t_0 \mathbf{u}_0 + t_1 \mathbf{u}_1$, with $t_1 \geq 0$. We get $t_1 > 0$ from C. Taking the dot product of both sides of this identity gives $\mathbf{e} \cdot \mathbf{u}_i = 0$. We have that $\{\mathbf{0}, \mathbf{e}, \mathbf{u}_i\}$ is not collinear because \mathbf{e} and \mathbf{u}_i are orthogonal and nonzero. The condition $\text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_0, \mathbf{u}_i) < \pi$ is equivalent to $\mathbf{e} \cdot (\mathbf{u}_0 \times \mathbf{u}_i) \neq 0$. The left-hand-side of this equation evaluates to $\|\mathbf{e}\| t_1 > 0$. \square

Lemma 4.10 [YBTASCZ] [formal proof by Hales]. *Let $\mathbf{u}_1, \dots, \mathbf{u}_r$ with cross product \mathbf{e} satisfy the conditions ECAU. For all $i < j \leq r$, we have*

$$\text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_0, \mathbf{u}_i) < \text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_0, \mathbf{u}_j).$$

Furthermore, $\text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_i, \mathbf{u}_j) < \pi$.

Proof Let $j \leq r$. We prove the result by induction on i .

When $i = 0$, we have

$$0 = \text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_0, \mathbf{u}_0) \leq \text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_0, \mathbf{u}_j).$$

We do not have equality because of A and C.

Assume for i and assume for a contradiction $i + 1$ fails. That is,

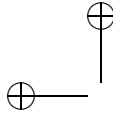
$$\text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_0, \mathbf{u}_i) < \text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_0, \mathbf{u}_j) \leq \text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_0, \mathbf{u}_{i+1}).$$

By AC, the inequality is strict. By U, this contradicts $\sigma \mathbf{u}_i = \mathbf{u}_{i+1}$. This proves the main statement.

The further claim follows from the addition law

$$\text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_0, \mathbf{u}_i) + \text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_i, \mathbf{u}_j) = \text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_0, \mathbf{u}_j).$$

\square



Lemma 4.11 [KCZXLL] [formal proof by Hales]. *Let $\mathbf{u}_1, \dots, \mathbf{u}_r$ with cross product \mathbf{e} satisfy the conditions ECAU. Assume $j < k \leq r$. If $i < j$ or $k < i$ then $\text{azim}(\mathbf{0}, \mathbf{u}_i, \mathbf{u}_j, \mathbf{u}_k) = \mathbf{0}$.*

Proof By A, the set $\{\mathbf{0}, \mathbf{u}_i, \mathbf{u}_j, \mathbf{u}_k\}$ is coplanar. This implies that the azimuth angle is 0 or π . By LDURDPN, using that $\{\mathbf{0}, \mathbf{u}_i, \mathbf{u}_j, \mathbf{u}_k\}$ is coplanar, we get that $\text{azim}(\mathbf{0}, \mathbf{u}_i, \mathbf{u}_j, \mathbf{u}_k) \neq \pi$ if and only iff

$$\text{aff}\{\mathbf{0}, \mathbf{u}_i\} \cap \text{conv}^0\{\mathbf{u}_j, \mathbf{u}_k\} = \emptyset. \quad (4.12)$$

Since $\text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_j, \mathbf{u}_k) < \pi$, we can relate the wedge to a lune:

$$\text{conv}^0\{\mathbf{u}_j, \mathbf{u}_k\} \subset \text{aff}_+^0\{\mathbf{0}, \mathbf{e}\}\{\mathbf{u}_j, \mathbf{u}_k\} = W^0(\mathbf{0}, \mathbf{e}, \mathbf{u}_j, \mathbf{u}_k) \subset \text{aff}_+^0(\{\mathbf{0}, \mathbf{e}, \mathbf{u}_0\}, \mathbf{u}_1)$$

and by the $t_1 > 0$ calculation above, we also have

$$\text{aff}\{\mathbf{0}, \mathbf{u}_i\} \cap \text{aff}_+^0(\{\mathbf{0}, \mathbf{e}, \mathbf{u}_0\}, \mathbf{u}_1) \subset \text{aff}_+^0\{\mathbf{0}\}\{\mathbf{u}_i\}.$$

Hence the intersection (4.12) is contained in $W^0(\mathbf{0}, \mathbf{e}, \mathbf{u}_j, \mathbf{u}_k) \cap \text{aff}_+^0\{\mathbf{0}\}\{\mathbf{u}_i\}$. A point \mathbf{p} in the intersection on the right gives

$$0 < \text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_j, \mathbf{p}) < \text{azim}(\mathbf{0}, \mathbf{e}, \mathbf{u}_j, \mathbf{u}_k) \text{ and } \text{azim}(\mathbf{0}, \mathbf{u}_j, \mathbf{p}) = \text{azim}(\mathbf{0}, \mathbf{u}_j, \mathbf{u}_i).$$

However, when $i < j$ or $k < i$, these inequalities fail. \square

Lemma 4.13 [FSQKWK] [formal proof by JH]. *If $\text{azim}(\mathbf{0}, \mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3) \leq \pi$ then*

$$\text{azim}(\mathbf{0}, \mathbf{v}_2, \mathbf{v}_3, \mathbf{v}_1) \leq \pi.$$

Proof $\text{azim}(\mathbf{0}, \mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3) \leq \pi$ iff $0 \leq \sin(\text{azim}(\mathbf{0}, \mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3))$ iff $0 \leq (\mathbf{v}_1 \times \mathbf{v}_2) \cdot \mathbf{v}_3$. The final condition is invariant under even permutations of the subscripts. \square

Recall that if we have $\text{azim}(\mathbf{0}, \mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3) \leq \pi$ and appropriate noncollinearity constraints, then $\text{azim}(\mathbf{0}, \mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3) = \text{dih}_V(\mathbf{0}, \mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3)$.

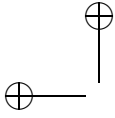
Lemma 4.14 [MKIFWJT] *Let $\mathbf{v}, \mathbf{w}_1, \mathbf{w}_2, \mathbf{w}_3 \in \mathbb{R}^3$. Assume that none of the sets $\{\mathbf{0}, \mathbf{v}, \mathbf{w}_i\}$ are collinear and that $\text{azim}(\mathbf{0}, \mathbf{v}, \mathbf{w}_1, \mathbf{w}_2) + \text{azim}(\mathbf{0}, \mathbf{v}, \mathbf{w}_2, \mathbf{w}_3) < 2\pi$. Then*

$$\text{azim}(\mathbf{0}, \mathbf{v}, \mathbf{w}_1, \mathbf{w}_2) + \text{azim}(\mathbf{0}, \mathbf{v}, \mathbf{w}_2, \mathbf{w}_3) = \text{azim}(\mathbf{0}, \mathbf{v}, \mathbf{w}_1, \mathbf{w}_3).$$

Proof This is (Fan.sum3_azim_fan). \square

Lemma 4.15 (deform-wedge) [XIVPHKS] [formal proof by John Harrison]. *Let $\mathbf{w}_0, \dots, \mathbf{w}_n \in \mathbb{R}^3$, for some $n \geq 1$. Assume that $i \neq j$ implies $\{\mathbf{0}, \mathbf{w}_i, \mathbf{w}_j\}$ is not collinear. Let $r = n - 1$. Set*

$$d(i, j, k) = \text{dih}_V(\mathbf{0}, \mathbf{w}_i, \mathbf{w}_j, \mathbf{w}_k) \quad \text{and} \quad a(i, j, k) = \text{azim}(\mathbf{0}, \mathbf{w}_i, \mathbf{w}_j, \mathbf{w}_k).$$



Introduce abbreviations for wedges $W(i) = W(\mathbf{0}, \mathbf{w}_i, \mathbf{w}_{i+1}, \mathbf{w}_{i-1})$. Let $\epsilon \in \mathbb{R}$ satisfy $0 < 2\epsilon < a(i, i+1, i-1)$, for $i = 1, \dots, r$. Assume $a(i, i+1, i-1) \leq \pi$ for $i = 1, \dots, r$. Assume that $d(p, q, q+1) < \epsilon$ for all sets of three distinct elements $\{p, q, q+1\} \subset \{0, \dots, r\}$. Assume that $d(p, p+1, q) < \epsilon$ for all sets of three distinct elements $\{p, p+1, q\} \subset \{0, \dots, r\}$ with $q > p+1$. Assume that $d(p+1, p, q) < \epsilon$ for all sets of three distinct elements $\{p+1, p, q\} \subset \{0, \dots, r\}$ with $q < p$.

Then for all k we have the statement S_k : for all $j \leq r-k$ we have $\mathbf{w}_j \in W(j+k)$ and $\mathbf{w}_{j+k} \in W(j)$.

Proof We prove S_k by induction on k . The cases $k = 0, 1$ are trivially satisfied. Assume that S_k , for some $k \geq 1$. We show S_{k+1} . Fix $j \leq r - (k+1)$.

By induction, $\mathbf{w}_j \in W(j+k)$, giving

$$a(j+k, j+k+1, j) \leq a(j+k, j+k+1, j+k-1) \leq \pi.$$

By Lemma 4.13, $a(j, j+k, j+k+1) \leq \pi$, and converting to dihedral, $a(j, j+k, j+k+1) < \epsilon$.

By induction, $\mathbf{w}_{j+k} \in W(j)$, giving

$$a(j, j+1, j+k) \leq a(j, j+1, j-1) \leq \pi,$$

and converting to dihedral, $a(j, j+1, j+k) < \epsilon$.

Since $2\epsilon < 2\pi$, we can add the angles (Lemma 4.14),

$$a(j, j+1, j+k) + a(j, j+k, j+k+1) = a(j, j+1, j+k+1) < 2\epsilon < a(j, j+1, j-1).$$

This says that $\mathbf{w}_{j+k+1} \in W(j)$.

The proof that $\mathbf{w}_j \in W(j+k+1)$ is similar.

By induction, $\mathbf{w}_{j+k+1} \in W(j+1)$, giving

$$a(j+1, j+k+1, j) \leq a(j+1, j+2, j) \leq \pi.$$

By Lemma 4.13, $a(j+k+1, j, j+1) \leq \pi$, and converting to dihedral $a(j+k+1, j, j+1) < \epsilon$.

By induction, $\mathbf{w}_{j+1} \in W(j+k+1)$, giving

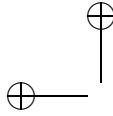
$$a(j+k+1, j+1, j+k) \leq a(j+k+1, j+k+2, j+k) \leq \pi,$$

and converting to dihedral $a(j+k+1, j+1, j+k) < \epsilon$.

Since $2\epsilon < 2\pi$, we can add angles,

$$a(j+k+1, j, j+1) + a(j+k+1, j+1, j+k) = a(j+k+1, j, j+k) < 2\epsilon < a(j+k+1, j+k+2, j+k).$$

This says that $\mathbf{w}_j \in W(j+k+1)$. □



We apply this lemma to prove the wedge property in ZLZTHIC (Dense Sphere Packings Lemma 7.25). We return to the context and notation of that lemma.

Lemma 4.16 [ITNZZRD] *Let (φ, V, I) be a deformation of a generic local fan (V, E, F) over an interval I . Assume that the azimuth angle of $\mathbf{v}_i(t)$ is at most π for all $t \in I$, whenever \mathbf{v}_i is straight. Assume that $(V(t), E(t), F(t))$ is a generic for all $t \in I$. (Or assume that (V, E, F) is a local fan and a deformation that moves a single nonpolar element of V .) Then for all sufficiently small $t \in I$, we have that $(V(t), E(t), F(t))$ satisfies the wedge property of local fans. That is, $V(t) \subset W_{\text{dart}}(x(t))$ for all $x(t) \in F(t)$.*

Proof As in the proof in Dense Sphere Packings, by continuity, the proof of the wedge property reduces to the case $\mathbf{w} \in W_{\text{dart}}(\mathbf{u}, \rho\mathbf{u}) \setminus W_{\text{dart}}^0(\mathbf{u}, \rho\mathbf{u})$. We may prove the more symmetrical statement

$$\mathbf{u}(t) \in W_{\text{dart}}(\mathbf{w}(t), \rho\mathbf{w}(t)) \text{ and } \mathbf{w}(t) \in W_{\text{dart}}(\mathbf{u}(t), \rho\mathbf{u}(t)).$$

Exchanging \mathbf{u} and \mathbf{w} as needed, we have a sequence $\mathbf{u} = \mathbf{v}_0, \dots, \mathbf{v}_r = \mathbf{w}$, with $\mathbf{v}_i = \rho^i \mathbf{v}$, where all of the terms \mathbf{v}_i ($i = 1, \dots, r-1$) are straight. We use Lemmas 7.15 and 7.19 to check the properties ECAU.

The interior angles of the fan are positive. Pick ϵ such that $0 < 2\epsilon < \text{azim}(\mathbf{0}, \mathbf{v}_i, \mathbf{v}_{i+1}, \mathbf{v}_{i-1})$ for $i = 0, \dots, r$. By the continuity of azim , there is some $\epsilon > 0$ such that $2\epsilon < \text{azim}(\mathbf{0}, \mathbf{v}_i(t), \mathbf{v}_{i+1}(t), \mathbf{v}_{i-1}(t))$ for $i = 0, \dots, r$ when $|t| < \epsilon$.

When $t = 0$, the straight conditions give

$$\text{dih}(\mathbf{0}, \mathbf{v}_p, \mathbf{v}_q, \mathbf{v}_{q+1}) = 0,$$

for distinct triples $\{p, q, q+1\} \subset \{0, \dots, r\}$, and

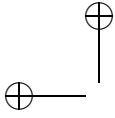
$$\text{dih}(\mathbf{0}, \mathbf{v}_p, \mathbf{v}_{p+1}, \mathbf{v}_q) = 0,$$

for distinct triples $\{p, p+1, q\} \subset \{0, \dots, r\}$ with $p+1 < q$. (This follows from the lemma above, which shows that $\text{azim} = 0$.) By the continuity of dih_V , shrinking ϵ as needed, we may assume the corresponding dihedral angles are less than ϵ , when evaluated at $\mathbf{v}_i(t)$, for $|t| < \epsilon$.

All of the conditions of Lemma 4.15 are satisfied for vectors $\mathbf{v}_i(t)$ for any $|t| < \epsilon$. Hence for all i and all j such that $i, j \leq r$, we have

$$\mathbf{v}_i(t) \in W_{\text{dart}}(\mathbf{v}_j(t), \mathbf{v}_{j+1}(t)) \text{ and } \mathbf{v}_j(t) \in W_{\text{dart}}(\mathbf{v}_i(t), \mathbf{v}_{i+1}(t)).$$

If we take $i = 0$ and $j = r$, this is the desired wedge condition. \square



5

Appendix on the proof of BGMIFTE

This is an appendix to Section 7.3 of Dense Sphere Packings.

This page contains some notes on the verification that the polar (V', E', F') is a local fan. This is an expanded version of the last paragraph of the proof of Lemma 7.34 (BGMIFTE). We check the intersection property of fans, the dihedral property of local fans, and the face property of local fans.

Since we are treating the last paragraph of the proof, we may assume that all the other parts of that lemma dealing with arcs and azimuth angles have been established.

We place ourselves in the context of Lemma BGMIFTE, adopting the notation from that lemma. In particular, $\mathbf{v}_i = \rho^i \mathbf{v}$ and $\mathbf{w}_i = \mathbf{v}_i \times \mathbf{v}_{i+1}$. We take the indices modulo $k = \text{card}(V)$, so that $\mathbf{v}_{i+k} = \mathbf{v}_i$. By earlier parts of the proof, we have

$$\mathbf{w}_j \in W^0(\mathbf{0}, \mathbf{w}_i, \mathbf{w}_{i+1}, \mathbf{w}_{i-1}), \quad j \neq i-1, i, i+1 \quad (5.1)$$

and

$$\{\mathbf{0}, \mathbf{w}_i, \mathbf{w}_j\} \text{ is not collinear, when } i \neq j \quad (5.2)$$

and

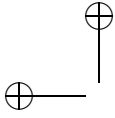
$$\text{arc}_V(\mathbf{0}, \{\mathbf{w}_i, \mathbf{w}_{i+1}\}) = \pi - \angle(\mathbf{v}_{i+1}) > 0 \quad (5.3)$$

and

$$\angle'(\mathbf{w}_{i+1}) = \pi - \text{arc}_V(\mathbf{0}, \{\mathbf{v}_{i+1}, \mathbf{v}_{i+2}\}) \in (0, \pi). \quad (5.4)$$

We verify the fan intersection property of (V', E', F') . Remark GMLWKPK gives some hints about verifying the intersection property, and notes that it comes down to two cases:

1. $\varepsilon \cap \varepsilon' = \emptyset$ implies $C(\varepsilon) \cap C(\varepsilon') = \emptyset$.
2. $\varepsilon \cap \varepsilon' = \{\mathbf{v}\}$ implies $C(\varepsilon) \cap C(\varepsilon') = C\{\mathbf{v}\}$.



If ε and ε' are both singletons then the intersection property follows from (5.2). Without loss of generality assume that ε is not a singleton. By the definition of E' , we have $\varepsilon = \{\mathbf{w}_i, \mathbf{w}_{i+1}\}$, for some i . We may partition $C(\varepsilon)$ as

$$C(\varepsilon) = C^0(\varepsilon) \cup C^0(\mathbf{w}_i) \cup C^0(\mathbf{w}_{i+1}) \cup \{\mathbf{0}\}.$$

Since we know the intersection property for singletons, we are reduced to showing the following.

1. if $\varepsilon' = \{\mathbf{w}_j, \mathbf{w}_{j+1}\} \neq \varepsilon$ is also an edge, then $C^0(\varepsilon) \cap C^0(\varepsilon') = \emptyset$.
2. for every $\mathbf{w}_j \in V'$, we have $C^0(\varepsilon) \cap C^0(\mathbf{w}_j) = \emptyset$.

Consider the first of these two enumerated cases. Exchanging i with j if necessary, we may assume that $j \neq i, i+1$. Set $\alpha(\mathbf{p}) = \text{azim}(\mathbf{0}, \mathbf{w}_i, \mathbf{w}_{i+1}, \mathbf{p})$. For every point \mathbf{p} in $C^0(\varepsilon)$ we have $\alpha(\mathbf{p}) = 0$ and $C^0(\varepsilon) \subset A := \text{aff}\{\mathbf{0}, \mathbf{w}_i, \mathbf{w}_{i+1}\}$. We separate this from $C^0(\varepsilon')$ by showing that $C^0(\varepsilon') \subset A_+^0 := \text{aff}_+^0(\{\mathbf{0}, \mathbf{w}_i, \mathbf{w}_{i+1}\}, \{\mathbf{w}_{i-1}\})$, and using the disjointness of this open half-space A_+^0 from its bounding plane A .

If $j = i-1$, then by (5.4), for every $\mathbf{q} \in C^0(\varepsilon')$, we have

$$\alpha(\mathbf{q}) = \alpha(\mathbf{w}_j) = \alpha(\mathbf{w}_{i-1}) \in (0, \pi).$$

The values of α separate $C^0(\varepsilon)$ from $C^0(\varepsilon')$. Note also that this gives

$$\mathbf{w}_{i-1} \in A_+^0 \tag{5.5}$$

If $\ell \neq i, i+1$, then by (5.1) and (5.5), we have $\mathbf{w}_\ell \in A_+^0$. From this, we obtain $\varepsilon' \subset A_+^0$ and from the conic structure of the halfspace A_+^0 , it follows that $C^0(\varepsilon') \subset A_+^0$.

The second enumerated case is similar, if $j \in \{i, i+1\}$, then the empty intersection property $C^0\{\mathbf{w}_i, \mathbf{w}_{i+1}\} \cap C^0(\mathbf{w}_j) = \emptyset$ follows from the strict inequality in the definition of C^0 and the linear independence of \mathbf{w}_i and \mathbf{w}_{i+1} . For example,

$$t_0 \mathbf{w}_i + t_1 \mathbf{w}_{i+1} = s \mathbf{w}_i,$$

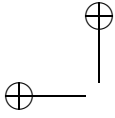
has no solution in positive real numbers t_0, t_1, s . Otherwise, if $j \notin \{i, i+1\}$, we separate $C^0(\varepsilon)$ from $C^0(\mathbf{w}_j)$ by the disjointness of A_+^0 and A , as in the first case.

This completes the proof of the intersection property of fans for the polar. Next we verify the dihedral property of local fans.

For this, we review how a hypermap is attached to the fan (V', E', F') . We have sets (V', E', F') defined to be

$$V' = \{\mathbf{w}_i = \mathbf{v}_i \times \mathbf{v}_{i+1} : i\},$$

$$E' = \{\{\mathbf{w}_i, \mathbf{w}_{i+1}\} : i\},$$



$$F' = \{(\mathbf{w}_i, \mathbf{w}_{i+1}) : i\}.$$

The set of darts of the hypermap consists of all orderings of edges:

$$D = \{(\mathbf{w}_i, \mathbf{w}_{i+1}) : i\} \cup \{(\mathbf{w}_{i+1}, \mathbf{w}_i) : i\} = F' \cup F'' \text{ say .}$$

The set $E'(\mathbf{w}_i)$, with overloaded notation, is defined as the set

$$\{\mathbf{w} \in V' : \{\mathbf{w}, \mathbf{w}_i\} \in E'\},$$

which in this situation reduces to

$$E'(\mathbf{w}_i) = \{\mathbf{w}_{i-1}, \mathbf{w}_{i+1}\}.$$

The permutation $\sigma(\mathbf{w}_i)$ of $E'(\mathbf{w}_i)$ is defined as the azimuth cycle on this set. Since the set has only two elements, the permutation is forced to swap \mathbf{w}_{i-1} and \mathbf{w}_{i+1} . The hypermap is a tuple (D, e, n, f) , where the permutations of D are generally defined as follows.

$$\begin{aligned} n(\mathbf{v}, \mathbf{w}) &= (\mathbf{v}, \sigma(\mathbf{v}, \mathbf{w})), \\ f(\mathbf{v}, \mathbf{w}) &= (\mathbf{w}, \sigma(\mathbf{w})^{-1}\mathbf{v}), \\ e(\mathbf{v}, \mathbf{w}) &= (\mathbf{w}, \mathbf{v}). \end{aligned}$$

In the present context, this reduces to

$$\begin{aligned} n(\mathbf{w}_i, \mathbf{w}_{i\pm 1}) &= (\mathbf{w}_i, \mathbf{w}_{i\mp 1}), \\ f(\mathbf{w}_i, \mathbf{w}_{i\pm 1}) &= (\mathbf{w}_{i\pm 1}, \mathbf{w}_{i\pm 2}), \\ e(\mathbf{w}_i, \mathbf{w}_{i\pm 1}) &= (\mathbf{w}_{i\pm 1}, \mathbf{w}_i). \end{aligned}$$

To prove the dihedral property, by Lemma QQYVCFM, it is enough to prove the following properties

1. the hypermap is connected.
2. the number of darts is $2k$.
3. the orders of f, n, e are $k, 2, 2$, respectively.

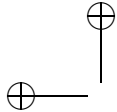
We have

$$f^j(\mathbf{w}_i, \mathbf{w}_{i\pm 1}) = (\mathbf{w}_{i\pm j}, \mathbf{w}_{i\pm(j+1)}).$$

Note that the orbit of f on $(\mathbf{w}_i, \mathbf{w}_{i\pm 1})$ is F' or $F'' \subset D$. (This proves in particular the face property of local fans: F' is a face of the hypermap.) and n exchanges darts in F' and F'' . Hence the hypermap is connected.

The sets F' and F'' are disjoint and each contain k darts. Hence the number of darts is $2k$.

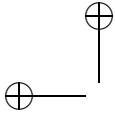
The smallest positive j such that $f^j(\mathbf{w}_i, \mathbf{w}_{i\pm 1}) = (\mathbf{w}_i, \mathbf{w}_{i\pm 1})$ is k . Hence f



has order k . The orders of e and n are 2 by inspection. This completes the verification of the dihedral property of local fans.

While proving the dihedral property, the face property fell out as well.

This completes the verification of properties intersection, fan, and dihedral.



6

Appendix on saturation

Lemma 6.1 [CPNKNXN] *Let $V \subset \mathbb{R}^3$ be any packing. Then there exists a saturated packing V_{sat} that contains V .*

Proof If V is any packing and r is a real number, let $m(V, r)$ be the maximum over the following set of natural numbers:

$$\{\text{card}(V' \cap B(\mathbf{0}, r)) : V \subset V' \text{ and } V' \text{ is a packing}\}.$$

By (the proof of) Lemma 6.2 of DSP (KIUMVTC) and Lemma (WQZISRI), there is a one-to-one map from $V' \cap B(\mathbf{0}, r)$ to $\mathbb{Z}^3 \cap B(\mathbf{0}, 2r+1)$, which gives an upper bound on the cardinality of $V' \cap B(\mathbf{0}, r)$ depending only on r . Thus the maximum exists.

Skolemizing variables, we have a function W of V and r with the property that for any packing V and any real number r , we have that $W(V, r)$ is a packing with $V \subset W(V, r)$ that realizes the maximum cardinality $m(V, r)$.

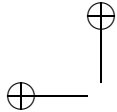
We define a packing by recursion. Let $V_0 = V$ and let $V_{n+1} = W(V_n, n)$. Set $V_{sat} = \bigcup_n V_n$. We will show that V_{sat} is a saturated packing that contains V . It will take a few steps to reach the conclusion.

We show that $V \subset V_{sat}$. In fact, $V = V_0 \subset \bigcup_n V_n = V_{sat}$.

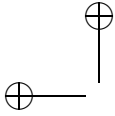
It is easily checked that if $V' \subset V''$ is an inclusion of packings, then for all r , we have $V' \subset W(V'', r)$. From this, it follows that if $V' \subset V_n$ then $V' \subset V_{n+1}$. By induction, it now follows that if $n \leq m$, then $V_n \subset V_m$.

We show that V_{sat} is a packing. If $\mathbf{v}, \mathbf{w} \in V_{sat}$, then $\mathbf{v}, \mathbf{w} \in V_m$ for some m . The set V_m is a packing, so that if $\mathbf{v} \neq \mathbf{w}$, we have $\|\mathbf{v} - \mathbf{w}\| \geq 2$. Hence V_{sat} is a packing.

We claim that V_{sat} is saturated. Otherwise, there exists $\mathbf{w} \in \mathbb{R}^3$, with $\mathbf{w} \notin V_{sat}$ such that $\|\mathbf{v} - \mathbf{w}\| \geq 2$ for all $\mathbf{v} \in V_{sat}$. Choose n such that $\mathbf{w} \in B(\mathbf{0}, n)$. Let $V' = V_{n+1} \cup \{\mathbf{w}\}$. Then V' is a packing containing V_n that achieves a larger



cardinality of intersection with $B(\mathbf{0}, n)$ than does V_{n+1} . This contradicts the maximality of V_{n+1} . \square



7

Appendix on OXLZLEZ

This appendix was written in Nov 2012 to give details of the formalization of Lemma OXLZLEZ.

7.1 formulas for γ and dih

The dihedral angle of a cell is given by a formula $\text{dih}X$, which is defined in terms of $\text{dih}V$. Various lemmas related $\text{dih}V$ to the function giving the dihedral angle of a simplex as a function of its edge lengths.

The function γ can also be expressed as a function of edge lengths. In the following lemmas, we fix a saturated packing V and take cells with respect to V .

Lemma 7.1 [YJBIAOE] *Let X be a 4-cell. Let y_1, \dots, y_6 be its edge lengths. Then $\gamma(X, L)$ is given by*

$$\text{gamma4fgcy } y_1 \ y_2 \ y_3 \ y_4 \ y_5 \ y_6 \ \text{lmfun}$$

defined formally in the module Sphere.

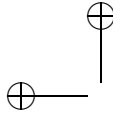
Lemma 7.2 [XKYBPAI] *Let X be a 3-cell. Let y_4, y_5, y_6 be the lengths of the three edges in $E(X)$. Then $\gamma(X, L)$ is given by*

$$\text{gamma3f } y_4 \ y_5 \ y_6 \ \text{sqrt2} \ \text{lmfun}$$

defined formally in the module Sphere.

Lemma 7.3 [KKHWUHM] *Let X be a 2-cell. Let y be the length of the unique edge in $E(X)$. Let α be the dihedral angle along that edge. Then $\gamma(X, L) = \alpha \cdot t$, where t is*

$$\text{gamma2_x_div_azim } (\text{h0cut } y) \ (y * y)$$



defined formally in the module `FunctionalEquation`.

Lemma 7.4 [OWEWPJG] *Let X be a 1-cell. Let \mathbf{v} be the unique element of $V(X)$. Let $s = \text{sol}(X, \mathbf{v})$. Then $\gamma(X, L) = st$, where t is*

$$\frac{8\pi\sqrt{2}}{3} - 8m_1$$

Lemma 7.5 [KPJNKIL] *Let X be a 0-cell. Then $\gamma(X, L)$ is equal to the volume of X .*

7.2 leaf and cell

Definition 7.6 (leaf) [NIPHFIE] *Let V be a saturated packing. A leaf of V is an element $\underline{\mathbf{u}} = [\mathbf{u}_0; \mathbf{u}_1; \mathbf{u}_2] \in \underline{V}(2)$ such that $h(\underline{\mathbf{u}}) < \sqrt{2}$. The stem of the leaf is $\{\mathbf{u}_0, \mathbf{u}_1\}$.*

Lemma 7.7 [GBEWFYFX] *Let V be a saturated packing, and let $\underline{\mathbf{u}} = [\mathbf{u}_0; \mathbf{u}_1; \mathbf{u}_2]$ be a leaf of V . Then $S = \{\mathbf{u}_0, \mathbf{u}_1, \mathbf{u}_2\}$ is not collinear.*

Proof Part 1 of MHFTTZN states that S has affine dimension 2, hence the set is not collinear. \square

Lemma 7.8 [NWVRFMF] *Let V be a saturated packing, and let $\underline{\mathbf{u}}$ be a leaf of V . Let $\mathbf{p} \in \mathbb{R}^3$ be such that $\{\mathbf{p}\}$ is a facet of $\Omega(V, \underline{\mathbf{u}})$. Then there exists $\underline{\mathbf{v}} \in \underline{V}(3)$ such that $d_2 \underline{\mathbf{v}} = \underline{\mathbf{u}}$ and $\omega_3(\mathbf{v}) = \mathbf{p}$.*

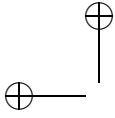
Proof This follows directly from Lemma IDBEZAL and $\underline{\mathbf{u}} \in \underline{V}(2)$. \square

Lemma 7.9 [YBZFUP0] *Let V be a saturated packing with leaf $\underline{\mathbf{u}}$. Then there exist distinct \mathbf{p}_1 and \mathbf{p}_2 such that $\Omega(V, \underline{\mathbf{u}})$ is the convex hull of $\{\mathbf{p}_1, \mathbf{p}_2\}$ and such that F is a facet of $\Omega(V, \underline{\mathbf{u}})$ if and only if $F \in \{\{\mathbf{p}_1\}, \{\mathbf{p}_2\}\}$.*

Proof By the definition of $\underline{V}(2)$, we have that $\underline{\mathbf{u}} \in \underline{V}(2)$ implies that the affine dimension of $\Omega(V, \underline{\mathbf{u}})$ is one. This is a bounded polyhedron of dimension one, hence a segment given as a convex hull of distinct points \mathbf{p}_1 and \mathbf{p}_2 . The facets of a segment are its extreme points as given. \square

Lemma 7.10 [ZASUVOR] *Let V be a saturated packing with leaf $[\mathbf{u}_0; \mathbf{u}_1; \mathbf{u}_2]$. Then $[\mathbf{u}_1; \mathbf{u}_0; \mathbf{u}_2]$ is also a leaf with the same stem.*

Proof The stem is clearly the same, and the circumradius does not change upon reordering of elements. Let $\underline{\mathbf{v}} \in \underline{V}(3)$ be an element constructed in the previous lemma such that $d_2 \underline{\mathbf{v}} = \underline{\mathbf{u}}$. Let $\underline{\mathbf{v}}'$ be obtained by transposing the first



two elements. By YNHYYIT, we have $\underline{v}' \in \underline{V}(3)$. Then $d_2 \underline{v}' = [\underline{u}_1; \underline{u}_0; \underline{u}_2] \in \underline{V}(2)$. The result ensues. \square

Lemma 7.11 [FUZBZGI] *Let V be a saturated packing with leaf \underline{u} . Let \mathbf{q} be the circumcenter of \underline{u} . Then $\mathbf{q} \in \Omega(V, \underline{u})$, but is not an extreme point of $\Omega(V, \underline{u})$.*

Proof The third part of Lemma MHFTTZN gives that $\mathbf{q} \in \Omega(V, \underline{u})$.

Assume for a contradiction that \mathbf{q} is an extreme point, then Lemma 7.8 gives $\underline{v} \in \underline{V}(3)$ such that $d_2 \underline{v} = \underline{u}$ and $\omega_3(\underline{v}) = \mathbf{q}$. The set $\Omega(V, \underline{v})$ is convex of affine dimension 0, and is therefore a singleton $\omega_3(\underline{v})$. By Lemma MHFTTZN applied to \underline{v} , we have that $\mathbf{q} = \omega_3(\underline{v})$ is the circumcenter of \underline{v} . This contradicts the strict inequality of Lemma XYOFCGX. \square

Definition 7.12 (χ) [MSBKFLD] For any list $\underline{u} = [\underline{u}_0; \underline{u}_1; \underline{u}_2]$ of elements in \mathbb{R}^3 , define $\chi(\underline{u}, \mathbf{p}) = ((\underline{u}_1 - \underline{u}_0) \times (\underline{u}_2 - \underline{u}_0)) \cdot (\mathbf{p} - \underline{u}_0)$.

Lemma 7.13 [JDHAWAY] *Let V be a saturated packing with leaf \underline{u} . Let \mathbf{p}_1 and \mathbf{p}_2 be the distinct points constructed in Lemma 7.9. Then $\chi(\underline{u}, \mathbf{p}_i)$ is not zero, and $\chi(\underline{u}, \mathbf{p}_1)$ and $\chi(\underline{u}, \mathbf{p}_2)$ have opposite signs.*

Proof Let \mathbf{q} be the circumcenter of \underline{u} . If $\chi(\underline{u}, \mathbf{p}_i) = 0$, then \mathbf{p}_i lies in the affine hull of \underline{u} . By MHFTTZN, this implies that $\mathbf{q} = \mathbf{p}_i$, which is impossible by the previous lemma. Hence $\chi(\underline{u}, \mathbf{p}_i) \neq 0$.

By the previous lemma, the circumcenter \mathbf{q} of \underline{u} has the form $\mathbf{q} = \mathbf{p}_1 t_1 + \mathbf{p}_2 t_2$, for some t_i such that $t_1 + t_2 = 1$ and $t_i > 0$. Since \mathbf{q} lies in the affine hull of \underline{u} , we have

$$0 = \chi(\underline{u}, \mathbf{q}) = t_1 \chi(\underline{u}, \mathbf{p}_1) + t_2 \chi(\underline{u}, \mathbf{p}_2)$$

Since $t_i > 0$, this implies that $\chi(\underline{u}, \mathbf{p}_1)$ and $\chi(\underline{u}, \mathbf{p}_2)$ have opposite signs. \square

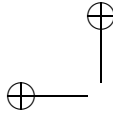
Remark 7.14 [LITLFSC] Recall from Lemma JBDNJJB that $\chi(\underline{u}, \mathbf{p})$ has the same sign as

$$\sin(\text{azim}(\underline{u}_0, \underline{u}_1, \underline{u}_2, \mathbf{p})).$$

Also, $\chi([\underline{u}_0; \underline{u}_1; \underline{u}_2], \mathbf{p}) = -\chi([\underline{u}_1; \underline{u}_0; \underline{u}_2], \mathbf{p})$.

Definition 7.15 ($\mathbf{p}_e, \hat{\mathbf{u}}, k_e, c$) [AQEQEDX] Let V be a saturated packing with leaf \underline{u} . We define $\mathbf{p}_e(\underline{u})$, $\hat{\mathbf{u}}$, and k_e as functions of \underline{u} as follows. By the previous lemma, there exists a unique extreme point \mathbf{p}_e of $\Omega(V, \underline{u})$ such that $\chi(\underline{u}, \mathbf{p}_e) > 0$. By an earlier lemma, there exists $\hat{\mathbf{u}} \in \underline{V}(3)$ (choose one) where $d_2(\hat{\mathbf{u}}) = \underline{u}$ and $\omega_3(\hat{\mathbf{u}}) = \mathbf{p}_e$. Finally, let $k_e \in \{3, 4\}$ be given by

$$k_e = \begin{cases} 4, & h(\hat{\mathbf{u}}) < \sqrt{2} \\ 3, & \text{otherwise.} \end{cases}$$



We abbreviate $c(\underline{\mathbf{u}}) = \text{cell}(\underline{\mathbf{u}}, k_e(\underline{\mathbf{u}}))$.

This construction uses a leaf $\underline{\mathbf{u}} = [\mathbf{u}_0; \mathbf{u}_1; \mathbf{u}_2]$ to select an extreme point of $\Omega(V, \underline{\mathbf{u}})$ and an associated cell. By earlier lemmas, the other extreme point of $\Omega(V, \underline{\mathbf{u}})$ is determined by the other leaf $[\mathbf{u}_1; \mathbf{u}_0; \mathbf{u}_2]$ with the same stem. It reverses the sign of χ . Note that $[\mathbf{u}_1; \mathbf{u}_0; \mathbf{u}_2]$ is the unique nontrivial rearrangement of $\underline{\mathbf{u}}$ with the same stem. If $\underline{\mathbf{u}}$ is a leaf, we write

$$A_+^0(\underline{\mathbf{u}}) = \text{aff}_+^0(\{\mathbf{u}_0; \mathbf{u}_1\}, \mathbf{u}_2), \text{ and } A(\underline{\mathbf{u}}) = \text{aff}\{\mathbf{u}_0, \mathbf{u}_1, \mathbf{u}_2\}.$$

Lemma 7.16 [NUNRRDS] *Let V be a saturated packing with leaf $\underline{\mathbf{u}}$. Then $c(\underline{\mathbf{u}})$ meets $A_+^0(\underline{\mathbf{u}})$. Furthermore, for every $\mathbf{q} \in c(\underline{\mathbf{u}})$, we have $\chi(\underline{\mathbf{u}}, \mathbf{q}) \geq 0$.*

Proof The cell is given as a convex hull of $\{\mathbf{u}_0, \mathbf{u}_1, \mathbf{u}_2, \mathbf{p}\}$, for some point \mathbf{p} . Hence the cell contains \mathbf{u}_2 which lies in $A_+(\underline{\mathbf{u}})$. The convex hull clearly lies in a half space bounded by $A(\underline{\mathbf{u}})$.

We have $\chi(\underline{\mathbf{u}}, \mathbf{p}) > 0$. Since \mathbf{q} is in the convex hull of four points with $\chi \geq 0$, we also have $\chi(\underline{\mathbf{u}}, \mathbf{q}) \geq 0$. \square

7.3 planarity

Lemma 7.17 [RIJRIED] *Let V be a saturated packing. Let X be a cell with an edge. Then X is not coplanar.*

Proof By definition, the vertex set is empty, if the cell is a null set. \square

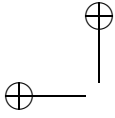
Lemma 7.18 [ZWVCBMN] *Assume that $S = \{\mathbf{u}_0, \dots, \mathbf{u}_3\} \subset \mathbb{R}^3$ is not coplanar. Then the convex hull of S has positive measure.*

Proof The volume of a tetrahedron has the form $\sqrt{\Delta}/12$, and the condition of planarity is $\Delta = 0$. \square

Lemma 7.19 [ASVAYEW] *Let V be a saturated packing and let X be a nonempty 3 or 4-cell. Then X is not coplanar.*

Proof Write the cell as $\text{cell}(\underline{\mathbf{u}}, k)$. A nonempty 3 or 4-cell has the form of a convex hull of four points, $\{\mathbf{u}_0, \mathbf{u}_1, \mathbf{u}_2, \mathbf{p}\}$, where $\mathbf{p} = \mathbf{u}_3$ in the case of a 4-cell ($h(\underline{\mathbf{u}}) < \sqrt{2}$) or $\mathbf{p} = \xi$ in the case of a 3-cell ($h_2 < \sqrt{2} \leq h(\underline{\mathbf{u}})$). In the case of a 4-cell the points are not coplanar by Lemma MHFTTN.

In the case of a 3-cell, we have that $A(\underline{\mathbf{u}})$ has dimension 2 by Lemma MHFTTN. It is enough to show that ξ is not in this affine hull. The point ξ is equidistant



from $\mathbf{u}_0, \mathbf{u}_1$ and \mathbf{u}_2 . If ξ is in the affine hull, then it is the circumcenter, and we arrive at a contradiction

$$\sqrt{2} = \|\xi - \mathbf{u}_0\| = h(d_2 \underline{\mathbf{u}}) < \sqrt{2}.$$

□

7.4 classification of cells

Lemma 7.20 [CFFONNL] *Let V be a saturated packing with leaf $\underline{\mathbf{u}} = [\mathbf{u}_0; \mathbf{u}_1; \mathbf{u}_2]$. Let X be a cell of V , and let $\{\mathbf{u}_0, \mathbf{u}_1\} \in E(X)$ be an edge. Suppose that*

$$X \cap A_+^0(\underline{\mathbf{u}}) \neq \emptyset.$$

Then there exists a 2-rearrangement $\underline{\mathbf{u}}'$ of $\underline{\mathbf{u}}$ such that

$$X = c(\underline{\mathbf{u}}').$$

Proof Since X has an edge, it is at least a 2-cell. The cell X is not a subset of $A(\underline{\mathbf{u}})$. Pick $\mathbf{q} \in X \setminus A$, and choose a 2-rearrangement $\underline{\mathbf{u}}'$ of $\underline{\mathbf{u}}$ such that $\chi(\underline{\mathbf{u}}', \mathbf{q}) > 0$. The cell X contains the convex hull of four points

$$\mathbf{u}_0, \mathbf{u}_1, \mathbf{p}, \mathbf{q}$$

where $\mathbf{p} \in A_+^0$ and $\chi(\underline{\mathbf{u}}', \mathbf{q}) > 0$.

Let $\underline{\mathbf{v}} = \underline{\hat{\mathbf{u}}} \in \underline{V}(3)$ and $k = k_e(\underline{\mathbf{u}}')$. Let $X' = \text{cell}(\underline{\mathbf{v}}, k)$. The cell X' also contains the convex hull of four points

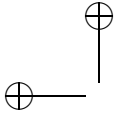
$$\mathbf{u}_0, \mathbf{u}_1, \mathbf{p}', \mathbf{q}'$$

where $\mathbf{p}' \in A_+^0$ and $\chi(\underline{\mathbf{u}}', \mathbf{q}') > 0$. Any two such convex hulls meet in a set of positive measure. So $X \cap X'$ has positive measure. Hence by Lemma AJRIPQN, we have $X = X'$. This gives the result. □

Lemma 7.21 [FUEIMOV] *Suppose that V is a saturated packing with leaves $\underline{\mathbf{u}}$ and $\underline{\mathbf{u}}'$ with the same stem. Let $k = k_e(\underline{\mathbf{u}})$ and $k' = k_e(\underline{\mathbf{u}}')$. Assume that $\underline{\mathbf{u}} \neq \underline{\mathbf{u}}'$. and $c(\underline{\mathbf{u}}) = c(\underline{\mathbf{u}}')$. Then $k = k' = 4$, $\underline{\hat{\mathbf{u}}} = [\mathbf{v}_0; \mathbf{v}_1; \mathbf{v}_2; \mathbf{v}_3]$, and $\underline{\hat{\mathbf{u}}}' = [\mathbf{v}_1; \mathbf{v}_0; \mathbf{v}_3; \mathbf{v}_2]$.*

Proof If the cells are equal, then $k = k' \in \{3, 4\}$ by the definition of k_e and Lemma AJRIPQN.

We consider two cases depending on the value of k . Suppose that $k = 4$. The definition of k_e gives $h(\underline{\hat{\mathbf{u}}}) < \sqrt{2}$. The set $V \cap \text{cell}(\underline{\mathbf{v}}, 4)$ determines the parameter $\underline{\mathbf{v}}$ up to rearrangement. The stem is fixed, giving only two possibilities $[\mathbf{u}_0; \mathbf{u}_1; \dots]$ and $[\mathbf{u}_1; \mathbf{u}_0; \dots]$ for the first two elements. The last two entries



must also be equal or transpositions. To preserve the sign of χ , the permutation must be even. The only nontrivial possibility is the one given in the lemma.

Suppose that $k = 3$. The cell is the convex hull of its extreme points, three of which are elements of the vertex set $V(X)$. The first two entries of $\underline{\mathbf{u}}$ are determined by the stem, up to transposition. The third entry is fixed by membership in $V(X)$. The fourth entry must be the common extreme point ξ . If $\underline{\mathbf{u}} = [\mathbf{u}_0; \mathbf{u}_1; \mathbf{u}_2]$ and $\underline{\mathbf{u}}' = [\mathbf{u}_1; \mathbf{u}_0; \mathbf{u}_2]$, then

$$0 < \chi(\underline{\mathbf{u}}, \xi) = -\chi(\underline{\mathbf{u}}', \xi).$$

This gives incompatible sign constraints. \square

Lemma 7.22 [YSAKKT χ] *Let V be a saturated packing. Let X be a cell with edge $\{\mathbf{u}_0, \mathbf{u}_1\}$. Then $X = \text{cell}(\underline{\mathbf{u}}, k)$, for some $\underline{\mathbf{u}}$ of the form $[\mathbf{u}_0; \mathbf{u}_1; \dots]$.*

Proof This follows from RVFXZBU. \square

Lemma 7.23 [RBUTTC χ] *Let V be a saturated packing. Let X be a 4-cell with edge $\{\mathbf{u}_0, \mathbf{u}_1\}$. Then there exists a leaf $\underline{\mathbf{u}}$ with stem $\{\mathbf{u}_0, \mathbf{u}_1\}$ such that $X = c(\underline{\mathbf{u}})$.*

Proof By YSAKKT χ , we may write $X = \text{cell}(\underline{\mathbf{v}}, 4)$, where $\underline{\mathbf{v}}$ has the form $\underline{\mathbf{v}} = [\mathbf{u}_0; \mathbf{u}_1; \mathbf{v}_2; \mathbf{v}_3]$, for some \mathbf{v}_2 and \mathbf{v}_3 . Set $\underline{\mathbf{u}} = d_2 \underline{\mathbf{v}}$. We have $h(\underline{\mathbf{u}}) < \sqrt{2}$, so that $\underline{\mathbf{u}}$ is a leaf. The cell meets $A_+^0(\underline{\mathbf{u}})$ at \mathbf{v}_2 . Lemma 7.20 gives that $X = c(\underline{\mathbf{u}}')$ for some 2-rearrangement of $\underline{\mathbf{u}}$. \square

Lemma 7.24 [FCHKUGT] *Let V be a saturated packing, and let $\underline{\mathbf{u}} = [\mathbf{u}_0; \mathbf{u}_1; \mathbf{u}_2]$ and $\underline{\mathbf{u}}' = [\mathbf{u}_0; \mathbf{u}_1; \mathbf{u}_2']$ be two leaves with the same first two entries, such that $A_+^0(\underline{\mathbf{u}}) = A_+^0(\underline{\mathbf{u}}')$. Then $\underline{\mathbf{u}} = \underline{\mathbf{u}}'$.*

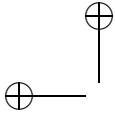
Proof We have that $c(\underline{\mathbf{u}})$ and $c(\underline{\mathbf{u}}')$ both meet A_+^0 . Hence, there is a 2-rearrangement $\underline{\mathbf{u}}''$ of $\underline{\mathbf{u}}'$ such that $c(\underline{\mathbf{u}}) = c(\underline{\mathbf{u}}'')$. By FUEIMOV, $k = k' = 4$, and $\underline{\mathbf{u}}''$ is a rearrangement of $\underline{\mathbf{u}}$. If $\underline{\mathbf{u}} \neq \underline{\mathbf{u}}''$, then $\underline{\mathbf{u}}'' = [\mathbf{u}_1; \mathbf{u}_0; \mathbf{u}_3]$, where $c(\underline{\mathbf{u}}) = \text{conv}\{\mathbf{u}_0, \dots, \mathbf{u}_3\}$. But this final possibility is impossible, since $\{\mathbf{u}_0, \dots, \mathbf{u}_3\}$ is not coplanar. \square

Lemma 7.25 [BDXKHTW] *Let V be a saturated packing. Let X be a cell of V with edge $e = \{\mathbf{u}_0, \mathbf{u}_1\}$. Let $\underline{\mathbf{u}} = [\mathbf{u}_0; \mathbf{u}_1; \mathbf{u}_2]$ and $\underline{\mathbf{u}}' = [\mathbf{u}_0; \mathbf{u}_1; \mathbf{u}_2']$ be distinct leaves with stem e . Assume that the intersection of X with $\text{wedge}^0(\mathbf{u}_0, \mathbf{u}_1, \mathbf{u}_2, \mathbf{u}_2')$ is nonempty. Then X is a subset of $\text{wedge}(\mathbf{u}_0, \mathbf{u}_1, \mathbf{u}_2, \mathbf{u}_2')$.*

Proof We claim that X is not coplanar. Otherwise it is a null set, and by definition a null set has no vertices or edges.

Since the leaves are distinct, by FCHKUGT, we have $A_+^0(\underline{\mathbf{u}}) \neq A_+^0(\underline{\mathbf{u}}')$. This implies that the wedges are nondegenerate.

For a contradiction, pick $\mathbf{p} \in X \cap \text{wedge}^0(\mathbf{u}_0, \mathbf{u}_1, \mathbf{u}_2, \mathbf{u}_2')$ and $\mathbf{q} \in X \cap \text{wedge}^0(\mathbf{u}_0, \mathbf{u}_1, \mathbf{u}_2', \mathbf{u}_2)$. We may assume that $\{\mathbf{u}_0, \mathbf{u}_1, \mathbf{p}, \mathbf{q}\}$ is not coplanar. The



sets wedge^0 are open and disjoint. By the connectedness of the unit interval, the path $t \mapsto (1-t)\mathbf{p} + t\mathbf{q}$ crosses $A_+^0(\underline{\mathbf{u}})$ or $A_+^0(\underline{\mathbf{u}}')$ at \mathbf{q}_t , for some $0 < t < 1$. We then have that $\{\mathbf{u}_0, \mathbf{u}_1, \mathbf{q}_t\}$ is not collinear. This point of crossing, by an earlier lemma gives $X = c(\underline{\mathbf{u}}'')$ for some 2-rearrangement of $\underline{\mathbf{u}}$ or $\underline{\mathbf{u}}'$, and say $A_+^0(\underline{\mathbf{u}}'') = A_+^0(\underline{\mathbf{u}})$. By an earlier lemma,

$$0 \leq (1-t)\chi(\underline{\mathbf{u}}'', \mathbf{p}) + t\chi(\underline{\mathbf{u}}'', \mathbf{q}) = \chi(\underline{\mathbf{u}}'', \mathbf{q}_t) = 0.$$

This forces $\chi(\underline{\mathbf{u}}'', \mathbf{p}) = \chi(\underline{\mathbf{u}}'', \mathbf{q}) = 0$, which is contrary to our assumption that $\{\mathbf{p}, \mathbf{q}, \mathbf{u}_0, \mathbf{u}_1\}$ is not coplanar. \square

Lemma 7.26 [EWYBJUA] *Let V be a saturated packing. Let X be any cell with edge $\{\mathbf{u}_0, \mathbf{u}_1\}$. Let $\underline{\mathbf{u}} = [\mathbf{u}_0; \mathbf{u}_1; \mathbf{u}_2]$ and $\underline{\mathbf{u}}' = [\mathbf{u}_0; \mathbf{u}_1; \mathbf{u}_2']$ be distinct leaves of V . Then*

$$X \subset \text{wedge}(\mathbf{u}_0, \mathbf{u}_1, \mathbf{u}_2, \mathbf{u}_2') \text{ or } X \subset \text{wedge}(\mathbf{u}_0, \mathbf{u}_1, \mathbf{u}_2', \mathbf{u}_2).$$

Proof The cell X is not a null set, so there exists a point in the intersection of X with

$$\text{wedge}^0(\mathbf{u}_0, \mathbf{u}_1, \mathbf{u}_2, \mathbf{u}_2') \text{ or } X \subset \text{wedge}^0(\mathbf{u}_0, \mathbf{u}_1, \mathbf{u}_2', \mathbf{u}_2).$$

Then apply the previous lemma. \square

7.5 angle sums revisited

In the following lemma, the hypothesis (7.28) always holds by preceding lemmas, but we include it to make it independent of these lemmas. The lemma is a variant of Lemma GRUTOTI.

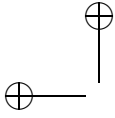
Lemma 7.27 [REUHADY] *Let V be a saturated packing. Assume that $\mathbf{u}_0, \mathbf{u}_1 \in V$ satisfy $\|\mathbf{u}_0 - \mathbf{u}_1\| < 2\sqrt{2}$. Set $\varepsilon = \{\mathbf{u}_0, \mathbf{u}_1\}$. Let $[\mathbf{u}_0; \mathbf{u}_1; \mathbf{w}]$ and $[\mathbf{u}_0; \mathbf{u}_1; \mathbf{w}']$ be two leaves (not necessarily distinct) such that for every cell X with edge $\{\mathbf{u}_0, \mathbf{u}_1\}$ we have*

$$X \subset \text{wedge}(\mathbf{u}_0, \mathbf{u}_1, \mathbf{w}, \mathbf{w}') \quad \text{or} \quad X \subset \text{wedge}(\mathbf{u}_0, \mathbf{u}_1, \mathbf{w}', \mathbf{w}). \quad (7.28)$$

Then

$$\sum_{X \in \mathcal{X}} \text{dih}(X, \varepsilon) = \text{azim}(\mathbf{u}_0, \mathbf{u}_1, \mathbf{w}, \mathbf{w}').$$

The sum runs over the set \mathcal{X} of cells X such that $\varepsilon \in E(X)$ and $X \subset \text{wedge}(\mathbf{u}_0, \mathbf{u}_1, \mathbf{w}, \mathbf{w}')$.



Proof Consider the sets

$$C = B(\mathbf{u}_0, r) \cap \text{rcone}^0(\mathbf{u}_0, \mathbf{u}_1, a), \text{ and } C' = C \cap \text{wedge}(\mathbf{u}_0, \mathbf{u}_1, \mathbf{w}, \mathbf{w}'),$$

where r and a are small positive real numbers. From the definition of k -cells, it follows that we can choose r and a sufficiently small so that if X is a k cell that meets C' in a set of positive measure, then $k \geq 2$ and there exists $\underline{\mathbf{u}} \in \underline{V}(3)$ such that $X = \text{cell}(\underline{\mathbf{u}}, k)$ and $d_1 \underline{\mathbf{u}} = [\mathbf{u}_0; \mathbf{u}_1]$. Moreover,

$$C' \cap X = C \cap X = C \cap A, \quad A = \text{aff}_+(\{\mathbf{u}_0, \mathbf{u}_1\}, \{\mathbf{v}, \mathbf{w}\}),$$

where A is the lune of [1][Definition MVIADQK] and \mathbf{v}, \mathbf{w} are chosen as in [1][Definition RSDYMHV]. By [1][Lemma FMSWMVO] and Definition RSDYMHV, the volume of this intersection is

$$\text{vol}(C \cap A) = \text{vol}(C) \text{dih}_V(\{\mathbf{u}_0, \mathbf{u}_1\}, \{\mathbf{v}, \mathbf{w}\})/(2\pi) = \text{vol}(C) \text{dih}(X, \varepsilon)/(2\pi).$$

The set of cells meeting C' in a set of positive measure gives a partition of C' into finitely many measurable sets. This gives

$$\begin{aligned} \text{vol}(C) \text{azim}(\mathbf{u}_0, \mathbf{u}_1, \mathbf{w}, \mathbf{w}')/(2\pi) &= \text{vol}(C') \\ &= \sum_{X \in \mathcal{X}} \text{vol}(C \cap X) \\ &= \text{vol}(C) \sum_{X \in \mathcal{X}} \text{dih}(X, \varepsilon)/(2\pi). \end{aligned}$$

The calculation of volumes in [1][Chapter 3] gives $\text{vol}(C) > 0$. The conclusion follows by canceling $\text{vol}(C)$ from both sides of the equation. \square

7.6 linear programs

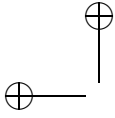
Let V be a saturated packing and let $\{\mathbf{u}_0, \mathbf{u}_1\}$ be a critical edge (of some cell). We may order the set

$$\text{Leaf} = \{\mathbf{v} \in V : [\mathbf{u}_0; \mathbf{u}_1; \mathbf{v}] \text{ is a leaf}\}$$

(and the corresponding leaves) by arbitrarily fixing one element $\mathbf{v}_0 \in L$ and then ordering them by increasing azimuth angle $\text{azim}(\mathbf{u}_0, \mathbf{u}_1, \mathbf{v}_0, \mathbf{v})$, for $\mathbf{v} \in \text{Leaf}$. This partitions \mathbb{R}^3 into finitely many wedges delimited by the leaves. Each cell with edge $\{\mathbf{u}_0, \mathbf{u}_1\}$ lies in one of these wedges.

If $h(\hat{\mathbf{u}}) < \sqrt{2}$, then $c(\underline{\mathbf{u}}) = c(\underline{\mathbf{u}}')$ is a 4-cell in the wedge, and this is the only cell (of positive measure or in fact the only cell at all) contained in the wedge.

If $h(\hat{\mathbf{u}}) \geq \sqrt{2}$, then $c(\underline{\mathbf{u}})$ and $c(\underline{\mathbf{u}}')$ are distinct 3-cells in the wedge, and these are the only 3 cells in the wedge. The wedge may also include a finite number



of 2-cells, but it has no 4-cells. The 2-cells within a given wedge are combined into a total γ and total azimuth angle.

We have a number of nonlinear inequalities bounding the value of

$$\gamma(X, L)\text{wt}(X) + \beta(\{\mathbf{u}_0, \mathbf{u}_1\}, X)$$

as a function of its azimuth angle. These inequalities are based on the partition of cells according to wedges. Based on these inequalities, we run linear programs giving lower bounds for $\sum \gamma(X, L)$, subject to the constraint that the azimuth angles sum to 2π . In every case, we find that $\Gamma \geq 0$ for each cluster. We run a separate linear program depending on the number of leaves. A generic case handles the case of five or more leaves. The following subsection gives details.

7.7 cell cluster inequality

This section shows how to deduce Lemma OXLZLEZ from a collection of nonlinear inequalities. The nonlinear inequalities are specified in an abbreviated way according to the table at the end of this section. Some collections of nonlinear inequalities have been merged into a single inequality. Once the source of a nonlinear inequality has been cited, we assume that the reader has become familiar with it, and it is not repeatedly cited.

Remark 7.29 (preparation) [XSBYGIQ]

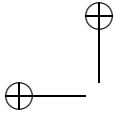
1. If a leaf separates two 3-cells, erase it, scoring it at 0.008azim. This is justified by a *computer calculation*¹ [2].
2. Score all 2-cells at 0.008azim by a *computer calculation*² [2].
3. Score all 3-cells with η^+ at 0.008azim, but keep the leaf, by *computer calculation*³ [2].

Definition 7.30 (small leaf, subcritical, supercritical) [HYTORSD] A *critical* edge has length $2h_- \leq y \leq 2h_+$. A *subcritical* edge has length $2 \leq y < 2h_-$. A *supercritical* edge has length $2h_+ < y \leq \sqrt{8}$. A *small leaf* is a triangle along the critical edge y_1 with one critical edge y_1 and two subcritical edges y_2, y_6 . A β -cell is one with two critical edges y_1, y_4 and the remaining four edges subcritical. Set $\tilde{\gamma} = \gamma(X)\text{wt} + \beta(X)$. A 23-cell is used to refer to the collection of 2 and 3-cells that lie between two consecutive nonerased leaves. The term *cell* will now refer to a 4-cell or a 23-cell, rather than a Marchal cell.

¹ [cell13]

² [grki]

³ [cell13]



A cell along a fixed stem $[\mathbf{u}_0; \mathbf{u}_1]$ is call a quarter, if it is a 4-cell with a critical edge along the stem and five subcritical edges. We write QU for a quarter. We write QX as an abbreviation for any 4-cell that is not a quarter, and QY for any 23-cell. Write $\epsilon = 0.0057$. When we speak of a leaf, we restrict our usage to the two leaves along the common stem. We write η^+ for the condition that the circumradius of a leaf > 1.34 and η^- for the circumradius ≤ 1.34 .

Remark 7.31 (Running Assumptions) In the following lemmas, we work in the context of OXLZLEZ. That is, we have a number of cells sharing a critical edge. For a contradiction, we assume that we are working with a counterexample. That is, we assume as a hypothesis in each lemma that $\sum \check{\gamma}_i < 0$. We have running assumptions `bool_model`, `bool_prep`, `real_model`, and $\sum \check{\gamma}_i < 0$.

Lemma 7.32 [CHQSQEY] *Assume the running assumptions. There are at least three 4-cells (in every counterexample).*

Proof We consider cases, according to the number of 4-cells. If there are no quarters, then $\check{\gamma}_i \geq 0$ for all i by *computer calculation*⁴ [2]. We may in fact assume that there is a quarter with $\check{\gamma} < 0$.

If there is one 4-cell (a quarter), then $\text{dih} < 2\pi$ by *computer calculation*⁵ [2] so there is another cell, a 23-cell. Then $\check{\gamma}_4 + \check{\gamma}_3 > 0$ by *computer calculation*⁶ [2]. If there are two 4-cells. The total angle of the two cells is at most $2(2.8) < 2\pi$, so there is at least one 23-cell, and each 4-cell is flanked by at least one 23-cell. Then use $\check{\gamma}_4 + \check{\gamma}_3 > 0$. \square

Lemma 7.33 [MTMLSRF] *Assume the running assumptions. There exists a quarter with $\check{\gamma} < 0$ and flanked on both sides by another 4-cell.*

Proof Otherwise, we again have $\check{\gamma}_4 + \check{\gamma}_3 > 0$. \square

Lemma 7.34 [LXDEYBO] *Assume the running assumptions. There are at most four 4-cells.*

Proof Otherwise, by a *computer calculation*⁷ [2]

$$\sum_{(5)} \check{\gamma} \geq 5a_5 + b_5\alpha \geq 5a_5 + b_5(2\pi) > 0.$$

\square

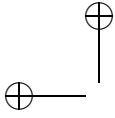
Lemma 7.35 [UNPNFVW] *Assume the running assumptions. The 4-cells are all contiguous. That is, there is at most one 23-cell.*

⁴ [gamma_qx]

⁵ [azim_c4] ,

⁶ [quqy] .

⁷ [ztg4]



Proof The number of 4-cells is three or four. By Lemma 7.33, there is a contiguous block of three or four 4-cells. Thus, if they are not all contiguous, there is a block of three and a block of one, with 23-cells interspersed. That is, there would be at least two 23-cells, each with angle at least 0.606. Then by a computer calculation⁸ [2]

$$\sum \check{\gamma} \geq 2(0.606)(0.008) + 4a_5 + (2\pi - 2(0.606))b_5 > 0.$$

□

Lemma 7.36 [DHCVTVE] *Assume the running assumptions. The configuration must be one of the four possibilities $C_{n,k}$, where n is the number of leaves, k is the number of 4-cells, and $n - k$ is the number of 23-cells:*

$$C_{3,3} \quad C_{4,4} \quad C_{4,3} \quad C_{5,4}.$$

Proof Since there are three or four 4-cells and zero or one 23-cell, there are four cases as given. □

Lemma 7.37 [PMZTATI] *Assume the running assumptions. Let QX be a cell with one small leaf and the other not small. Then $\check{\gamma} > \epsilon$ or QX is a β -cell.*

Proof This is a computer calculation⁹ [2]. □

Lemma 7.38 [RJSZKQX] *Assume the running assumptions. Let QU be a cell with (at least) one leaf η^+ . Then $\check{\gamma} > \epsilon$.*

Proof This is a computer calculation¹⁰ [2]. □

Lemma 7.39 [IXPFBKA] *Assume the running assumptions. If η^- , then the leaf is small.*

Proof This is a computer calculation¹¹ [2]. □

Remark 7.40 (neutralize) [XCLCXWG] $\check{\gamma}_{QU} > -\epsilon$ and $\check{\gamma}_{QX} \geq 0$ by a computer calculation¹² [2] If QX has one leaf that is small and another that is not small, it *neutralizes* a quarter in the sense that it gives ϵ against the $-\epsilon$ of the quarter. Notice that this condition propagates, forcing one cell after another to have small leaves, if it doesn't neutralize. If a 4-cell is next to a 23-cell, then we have by a computer calculation¹³ [2],

$$\check{\gamma}_4 + \check{\gamma}_3 > \epsilon,$$

⁸ [gckb ztg4 cell3 grki]

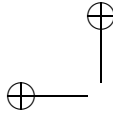
⁹ [gamma8]

¹⁰ [fhvb2]

¹¹ [jsp]

¹² [gamma_qu gamma_gx].

¹³ [gamma10_gamma11 qu_qy]



and again we can neutralize a quarter. Call this 234-neutralization.

Lemma 7.41 [IPVICGW] *Assume the running assumptions. All leaves along the stem are small.*

Proof We consider many cases, breaking it down initially according to the cases $C_{n,k}$.

$C_{3,3}$: If there are at least two quarters, every leaf lies along a quarter and is small. If there is one quarter, we can neutralize it, unless all leaves are small.

$C_{4,4}$, $C_{4,3}$: If there are at least three quarters (or two that are nonadjacent), then every leaf lies along a quarter and is small. If there are two adjacent or one quarter, we can neutralize them unless every leaf is small.

$C_{5,4}$: If there are four quarters, then every leaf is along a quarter. If there is one quarter, we can neutralize it, if it is not small. There remain two cases: two or three quarters.

$C_{5,4}$: two or three quarters. There are various combinatorial placements of the quarters among the four 4-cells. If two of them are not adjacent to the 23-cell, then we can treat it with neutralization arguments. The one case that is not easily neutralized is an arrangement with a single contiguous block of quarters, occupying the two slots that are not adjacent to the 23-cell, and possibly one other slot. Furthermore, there is exactly one leaf (along the 23-cell) that is not small. That nonsmall leaf has η^+ . We break this case into subcases.

Assume $\text{azim}_{23} > 1.074$. By a computer calculation¹⁴ [2], we have

$$\sum \check{\gamma} > 0.008\text{azim}_{23} + (4a_5 + b_5(2\pi - \text{azim}_{23})) > 0.$$

Assume $\text{azim}_{23} \leq 1.074$ and assume η^- on the small leaf along the 23-cell. By a computer calculation¹⁵ [2], the 23-cell can be included in the estimate,

$$\sum \check{\gamma} > 5a_5 + b_5(2\pi) > 0.$$

Assume $\text{azim}_{23} \leq 1.074$ and assume η^+ on the small leaf along the 23-cell. In this case we can split the 4-cells into two groups of two, each with a quarter paired with a neutralizing 4-cell by a computer calculation¹⁶ [2]. \square

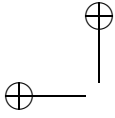
Lemma 7.42 [RSIWAMP] *Assume the running assumptions. There are at most four leaves. That is, $C_{5,4}$ does not occur.*

Proof Otherwise, we have reduced to the case of five leaves, four 4-cells, and one 23-cell, and all leaves small. Let A and B be the two leaves along the 23-cell. We consider various cases.

¹⁴ [cell13 grki ztg4]

¹⁵ [pem ztg4]

¹⁶ [gamma10_gamma11 fhbv2]



Assume $\text{azim}_{23} > 1.074$. This case is the same as in the previous lemma. Without generality we now assume that $\text{azim}_{23} \leq 1.074$.

Assume η_A^- or η_B^- . In this case, by a *computer calculation*¹⁷ [2], we have

$$\sum_{(5)} \check{\gamma} > 5a_5 + b_5(2\pi) > 0.$$

Assume η_A^+ and η_B^+ . If further, some 4-cell is not a quarter, then by a *computer calculation*¹⁸ [2],

$$\sum_{(5)} \check{\gamma} > 0.606(0.008) + 0.21849 + 3(0.161517) - (2\pi - 0.606)0.119482 > 0.$$

And if all 4-cells are quarters, then we can easily neutralizes the two pairs of 4-cells. \square

Lemma 7.43 [BKLETJQ] *Assume the running assumptions. If QX is adjacent to a 23-cell, then $\check{\gamma}_4 + \check{\gamma}_3 > \epsilon$.*

Proof This is a *computer calculation*¹⁹ [2]. \square

Lemma 7.44 [UTE0ITF] *Assume the running assumptions. There is no 23-cell. That is, $C_{4,3}$ does not occur.*

Proof Again, there are several cases. Label the three consecutive 4-cells as A, B, C . All leaves are small, B is a quarter with $\check{\gamma}_B < 0$ and its two leaves have η^- . If any leaf has η^+ , then we can easily neutralize. Assume now that η^- for all leaves. If any 4-cell is not a quarter, we can neutralize.

Assume that all 4-cells are quarters.

Assume $\text{azim}_{23} \geq 2.089$. By a *computer calculation*²⁰ [2],

$$\sum \check{\gamma} > \text{azim}_{23}(0.008) + 3(0.161517) + (2\pi - \text{azim}_{23})(-0.119482) > 0.$$

Assume $\text{azim}_{23} \leq 1.946$. By a *computer calculation*²¹ [2],

$$\sum \check{\gamma} > \text{azim}_{23}(0.008) + 3(-0.0659) + (2\pi - \text{azim}_{23})0.42 > 0.$$

Assume $1.946 \leq \text{azim}_{23} \leq 2.089$. By a *computer calculation*²² [2],

$$\sum \check{\gamma} > 3\epsilon - 3\epsilon \geq 0.$$

\square

¹⁷ [pem tew]

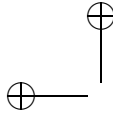
¹⁸ [gaz9 gaz6 gckb]

¹⁹ [gamma10_gamma11]

²⁰ [gaz6]

²¹ [azim1]

²² [txq]



Lemma 7.45 [LUIKGMH] *Assume the running assumptions. The case $C_{3,3}$ does not occur.*

Proof By a *computer calculation*²³ [2], we have $\text{azim} < 2.8$. If QX and $\check{\gamma} < \epsilon$, then $\text{azim} < 2.3$. If QU and $\check{\gamma} < 0$, then $\text{azim} < 1.65$.

If there are at least two negative quarters, then the total angle is

$$2\pi \leq 1.65 + 1.65 + 2.8 < 2\pi.$$

If there is one negative quarter, it will be neutralized unless both others have angle < 2.3 . Then

$$2\pi \leq 1.65 + 2(2.3) < 2\pi.$$

□

Lemma 7.46 [GRHIDFA] *Assume the running assumptions. The case $C_{4,4}$ does not occur.*

Proof Let k be the number of quarters. We have seen that all faces are small.

If $k \leq 2$, then by a *computer calculation*²⁴ [2],

$$kQU + (4 - k)QX : \sum \check{\gamma} > k1.61517 + (4 - k)0.213849 - 2\pi(0.119482) > 0.$$

if $k = 4$, then by a *computer calculation*²⁵ [2],

$$\sum \check{\gamma} > -4(0.0659) + 0.42(2\pi) > 0.$$

Finally assume that $k = 3$. If the fourth cell is a β -cell, we are done by a *computer calculation*²⁶ [2]. Thus, the fourth edge of the fourth cell is supercritical. This gives by a *computer calculation*²⁷ [2],

$$3QU + QX : 3(-0.0142852) + (0.00457511) + 2\pi(0.00609451) > 0.$$

□

7.8 table of inequalities

Here is a table of the inequalities that were used. The first column gives the name referenced in the text. Then it is indicated if the inequality is in the module *Ineq* or *Module* and the name of the inequality in the computer development. The final column gives a shorthand form of the inequality.

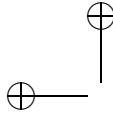
²³ [azim1 azim_c4 g-qxd]

²⁴ [gaz6 gaz9]

²⁵ [azim1]

²⁶ [ox3qlh] .

²⁷ [gaz4 azim2]



general bounds

jsp	ineq	JSPEVYT	η^+
gckb	ineq	GCKBQEA	$\text{azim} > 0.606$
azim_c4	ineq	BIXPCGW 6652007036 a2	$QU \vee QX \Rightarrow \text{azim} < 2.8$
ox3q1h	merge	ox3q1h_merge	$3QU + 1B \Rightarrow \sum_4 \check{\gamma} > 0.$

quarters

gamma_qu	ineq	BIXPCGW 9455898160	$QU \Rightarrow \check{\gamma} > -\epsilon$
fhbv2	ineq	FHBVYXZv2 a	$QU \eta^+ \Rightarrow \check{\gamma} > \epsilon$
quqy	merge	g_quqya-g_quqyb	$QU + QY \Rightarrow \check{\gamma}_4 + \check{\gamma}_3 > 0$
ztg4	merge	ztg4	$QU \vee QX \Rightarrow (a_5, b_5)$
azim1	ineq	QITNPEA 5653573305	$QU \Rightarrow (\check{\gamma}, \text{azim})$
gaz4	ineq	QITNPEA 6206775865	$QU \Rightarrow (\check{\gamma}, \text{azim})$
gaz6	ineq	QITNPEA 3848804089	$QU \Rightarrow (\check{\gamma}, \text{azim})$

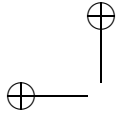
nonquarter 4-cells

gamma_qx	merge	gamma_qx	$QX \Rightarrow \check{\gamma} > 0$
g_qxd	merge	gamma_qxd	$QXD \Rightarrow \check{\gamma} > \epsilon$
gamma10_gamma11	merge	gamma10_gamma11	$QX^{ss} + QY \Rightarrow \check{\gamma}_4 + \check{\gamma}_3 > \epsilon$
gamma8	merge	QITNPEA_9063653052_weak	$QX^{1s} \Rightarrow \check{\gamma} > \epsilon$
gaz9	ineq	QITNPEA 2134082733	$QX^{ss} \Rightarrow (\check{\gamma}, \text{azim})$
azim2	ineq	QITNPEA 9939613598	$QX^{ss,super} \Rightarrow (\check{\gamma}, \text{azim})$

23-cells

cell3	merge	cell3_008_from_ineq	$\check{\gamma}_3 > 0.008\text{azim}$
grki	ineq	GRKIBMP a	$\check{\gamma}_2 > 0.008\text{azim}$
pem	ineq	PEMKWKU	$QY \eta^+ \eta^- \text{azim}^- \Rightarrow (a_5, b_5)$
tew	ineq	TEWNSCJ	$QY, 2\eta^- \Rightarrow (a_5, b_5)$
txq	ineq	TXQTPVC, IXPOTPA	$QY 2\eta^- \text{azim}[] \Rightarrow \check{\gamma} > 3\epsilon$

Table 7.1 inequalities



8

Appendix on connecting with Bauer-Nipkow work in Isabelle

This section builds on the HOL-Light files `hypermap.hl` and `import-tame-classificaton.hl`.

For simplicity of typesetting, we use a dash in this text to represent an underscore in HOL Light.

8.1 basic definitions

(This section has been formalized 2/2014.)

In the Bauer-Nipkow work, there is a type for Isabelle graphs, which we abbreviate to `lgraph`. There is a predicate `planegraph` that expresses `lgraph` planarity. There is also a slightly broader class `planegraph-relaxed`.

Lemma 8.1 [DPZGBYF] [formal proof by tchales]. *If g is plane-graph, then it is also planegraph-relaxed.*

There is a function, `fgraph`, that maps an `lgraph` to the list of list representation of the `lgraph`. For example, the list of list representation of a planar graph consisting of a square triangulated into four squares with common vertex 0 is

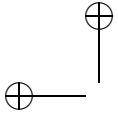
$$[[0; 1; 2]; [0; 2; 3]; [0; 3; 4]; [0; 4; 1]; [4; 3; 2; 1]].$$

We refer to these entries `[0; 1; 2]`, `[0; 2; 3]`, etc. as the faces of the `lgraph`.

There is a function, `hypermap-of-list`, that maps the list of list representation of the planar graph into the corresponding hypermap. The darts of the hypermap are consecutive pairs of elements in the list of list representation. For example, the example above gives the dart set

$$\{(0, 1), (1, 2), (2, 0), (0, 2), (2, 3), (3, 0), (0, 3), (3, 4), (4, 0), (0, 4), (4, 1), (1, 0), (4, 3), (3, 2), (2, 1), (1, 4)\}.$$

The face map follows each pair around the face in which it occurs $f(0, 1) = (1, 2)$, etc. The edge map reverses an ordered pair $e(0, 1) = (1, 0)$, etc. The set



of integers that appear in this list of list is in bijection with the set of vertices of the lgraph. In fact, we can take this as the definition of the set of vertices of the lgraph. When we refer to the hypermap of g , we refer to this construction.

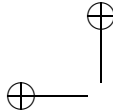
There is a function, `finals`, that returns a sublist of the list of list representaton of the lgraph.

In the Isabelle development, there is a function `facesAt`.

Since our aim is relate hypermaps to the constructions in Isabelle, which are based on lists, we give a series of list-based definitions that run in parallel with hypermap-based definitions.

We make a table of notions about hypermaps and their corresponding list-based notion. This table serves as a dictionary translating between the different data structures.

hypermap	lists	theorem
loop	loop-list	
is-contour	contour-list	contour-list-is-contour
is-normal	normal-list	normal-list-normal
hyp'm, hyp'p, hyp'q	l'm, l'p, l'q	
hyp'y, hyp'z	l'y, l'z	
final-quotient-face	final-list	
quotient	quotient-list	
-	final-dart-list	
-	split-normal-list	
transform	transform-list	
hyp'S	s-list	
s-flagged	s-flag-list	
empty-flagged	flag-list	
is-marked	marked-list	
iso	iso-list	
hyp-iso	isop-list	
f-map, n-map, e-map	f-list, n-list, e-list	
dih2k	dih2k-list	
atom-choice	find-atom	
atom	atom-list	
head-of-atom	TAIL	
tail-of-atom	HD	
is-split-condition	split-condition-list	
index	indexf	index-indexf
loop-of-face	loop-of-list	loop-of-face-list



We also describe the general correspondence of variable names.

hypermap	lists	function
H (hypermap)	L (list of lists)	hypermap-of-list
NF (normal family)	N	loop-family-of-list
L (loop)	r	loop-of-list
x (dart)	x	

We also have conversions from lists to hypermaps. The constant loop-of-list converts a list to a loop, and loop-family-of-list converts a list of lists to a family of loops.

There is a function match-core-list that is true when there is a correspondence between faces of an lgraph g and those of the quotient of L by N , such that the final faces of g are included in those of N . A function match-list also measures the correspondence. The function loop-choice generates a pair (r, x) to be used as part of a marked-list, from a lgraph g .

8.2 properties of planegraph

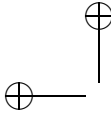
Definition 8.2 (good-list, good-list-nodes, good-graph) Recall that a list of lists L is a good-list if it has three properties:

1. The list of darts of L has no duplicates.
2. Every member of L is non-nil.
3. If (x, y) is a dart of L , then so is (y, x) .

Recall that a good list L is a good-list-nodes if the number of nodes in its hypermap is equal to the number of vertices of L . We say that an lgraph is good-graph if it has the properties:

1. Its fgraph L is a good-list and a good-list-nodes.
2. all faces of the fgraph are final.
3. In each face, each vertex occur at most once: (all uniq L).
4. The vertex set coincides with the set of elements of the fgraph.
5. The facesAt a vertex v is the same as the list of faces that contain v .

(We don't need to formalize JUJUWAT, CGGZYRC, EAHHATZ, ETDLJXT, and HWDMMZDM because we will get them as a corollary of the correspondence between restricted hypermaps and plane-graphs.)



Lemma 8.3 [JUJUWAT] *If g is planegraph-relaxed, then its fgraph is a good-list.*

Lemma 8.4 [CGGZYRC] *Let g be a planegraph. Let F be a face of g that is not final. Let (x, y) be a dart of (the hypermap of) g on face F . Then (y, x) is a dart on a face that is final.*

Lemma 8.5 [EAHHATZ] *If g is a Plane-graph, then it is a good-graph.*

Lemma 8.6 [ETDLJXT] *If g is planegraph-relaxed, then the vertices of g are $0, \dots, n-1$, where n is the number of nodes of g .*

Lemma 8.7 [HWDZDM] *Let g be planegraph-relaxed, and v in the vertex set of g . Then $\text{facesAt } g \ v$ is the set of faces of g that v is a member of.*

In general we want to have the properties of good-list, good-list-nodes, and all-uniq. The following lemmas show how these properties propagate. The following properties can be proved by structural induction for planegraph-relaxed. (We no longer need to use planegraph-relaxed; planegraph will be sufficient for our purposes.)

Lemma 8.8 [PMBRINH] *If L is good-list-nodes and normal-list- L - N , then good-list-nodes holds of quotient-list- L - N .*

Proof It is enough to show that all of the darts with a given first coordinate form a single node of the quotient. This is clear, because by assumption all the darts of L with a given first coordinate form a single node for L . \square

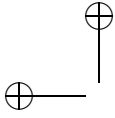
Lemma 8.9 [SNVACWG] [formal proof by tch 3/2014]. *If L is a good-list and good-list-nodes, its hypermap is restricted, and normal-list- L - N , then good-list holds of quotient-list- L - N .*

Proof We examine each of the defining properties of a good-list.

(UNIQ) If a dart (a, b) occurs more than once then there are two darts (x_1, y_1) and (x_2, y_2) in N , that run from nodes a to b . This is impossible by the no-double-joins property of restricted hypermaps.

(ALL-NON-NIL) It is a property of normal that its members are not nil. The same is then true of the quotient.

(SYM). Let (a, b) be a dart in the quotient. Then there is a dart (a, b) in N . In fact, we see that (a, b) is the last member of its atom in N . The family N meets the node b , and (b, a) is a dart at the node b . By properties of normal, all darts at the node, including (b, a) , are on N . We have $n^{-1}(b, a) = f(a, b)$. By unicity, (b, a) is followed by an f -step in N . This implies that (b, a) is a dart in the quotient. \square



Lemma 8.10 [LYNVPSU] [formal proof by tch 3/3/2014]. (L, N, r, s) be a marked list. Then the quotient of L by N is all-uniq.

Proof This is a defining property of a marked hypermap. □

Lemma 8.11 [LEBHIJR] Let L be a good-list and good-list-nodes. Let (L, N, r, x) be a marked list. Let (N', r') be the transform. Then the set of elements of L/N' is the union of the set of elements of L/N , together with the set of nodes that visit N' but not N .

Lemma 8.12 [HOJODCM] Let L be a good-list and good-list-nodes. Let (L, N, r, x) be a marked list with transform (N', r') . Let r'' be the offshoot of the transform. Let V be the set of nodes that visit N' but not N . Let s be a loop in N' , and v a node of N' . Then we have the following rules describing when v is visited by s .

1. If $s \neq r', r''$, then v is visited by s iff v is visited by s in N .
2. If $v \in V$, then v is visited by s iff $s = r'$ or r'' . Now assume $v \notin V$.
3. v is visited by r' iff v is at the node of $l'z, l'y$ or between them.
4. v is visited by r'' iff v is at the node of $l'y, l'z$ or between them.

8.3 main result

This result is imported from Isabelle.

Theorem 8.13 (Import Tame Classification) Let g be a final planegraph that is tame. Then there exists a y in the archive such that the fgraph of g is fgraph congruent to y .

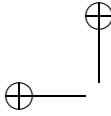
The main result to be proved is the following lemma.

Lemma 8.14 (LSKOKJE) Let H be a restricted hypermap. Then there exists a good-graph g that is a PlaneGraphs, such that H is isomorphic to the hypermap of g .

8.4 isomorphism

We develop the properties of iso-list, which is the list-based version of isomorphism. It is a particularly rigid notion, because it does not allow the rearrangement of the elements of a list. However, it is sufficient for our purposes.

We often need to consider L together with a normal-list N . The definition of iso-list carries with it a second argument that checks that N is sent to N' under the isomorphism.



Lemma 8.15 [GNBEVVU] [formal proof by tchales]. *If L is a good-list and L' is an iso-list, then the hypermaps of L and L' are isomorphic.*

Proof This follows directly from the theorem hypermap-of-list-map. \square

Lemma 8.16 (iso-list-refl, iso-list-sym, iso-list-trans) [formal proof by tchales]. *The relation iso-list is reflexive, symmetric, and transitive.*

The properties of a list L that are invariant under an injective map remain true when L is replaced with an iso-list. This gives the following results.

Lemma 8.17 [PEUTLZH] [formal proof by tchales, 2/2014]. *Let L be a good-list, and L' an iso-list. Then L' is a good-list.*

Lemma 8.18 [OISRWOF] [formal proof by tchales, 2/2014]. *Let L be good-list and good-list-nodes, and L' an iso-list. Then L' is a good-list-nodes.*

Lemma 8.19 [UEYETNI] [formal proof by tchales, 2/2014]. *Let L be all-uniq, and L' an iso-list. Then L' is a all-uniq.*

Lemma 8.20 [XKDZKWV] [formal proof by tch, 2/25/2014]. *Let L be a good-list with normal-list N . If (L', N') is iso-list, then N' is a normal-list of L' .*

Lemma 8.21 [DAKEFCC] [formal proof by tch 3/10/2014]. *If H and K are isomorphic hypermaps, and if H is restricted, then K is restricted.*

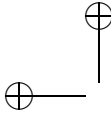
Lemma 8.22 [MEEIXJO] *Let L be a good-list with marked-list (L, N, r, x) . If (L', N', r', x') is iso-list, then (L', N', r', x') is a marked-list.*

8.5 representing hypermap as lists

It is convenient to work with lists of list L rather than hypermaps H . This means we should give a list version of notions such as loop and normal family. We have loop-list as the list version of a loop and norm-list as the list version of a normal family of loops.

Lemma 8.23 [RXOKSKC] [formal proof by HLT, 2/26/2014]. *Let H be a restricted hypermap. Then there exists a list L such that H is its hypermap. In fact, there exists such an L that is good-list, good-list-nodes, all-uniq.*

Proof We may convert a hypermap to list L as follows. From each face F of



H pick a dart x and form the list $[\text{node}(x), \text{node}(fx), \text{node}(f^2x), \dots]$. The list of all these lists for each face is defined to be L .

A restricted hypermap is simple, hence the elements of the list are uniq. This gives the all-uniq property.

By construction the elements of these lists are the nodes. This gives the good-list-nodes property.

We check that it is a good list. For the uniqueness property, note that we show below that the map between darts sets is one-to-one. This means that each dart appears in a face once.

For the non-nil property of good-lists: each list contains $\text{node}(x)$ as the first element. Hence it is not the nil-list.

For the symmetry property of good-lists: if $(\text{node}(x), \text{node}(fx))$ is a dart, then

$$(\text{node}(nfx), \text{node}(fnfx)) = (\text{node}(fx), \text{node}(x))$$

is also a dart.

Finally, we check that H is isomorphic to the hypermap of L . The darts of the hypermap of L are formed by consecutive pairs $(\text{node}(f^i x), \text{node}(f^{i+1} x))$.

We claim the map between dart sets, $y \mapsto (\text{node}(y), \text{node}(fy))$, is one-to-one. This follows from the no-double-joins property of restricted hypermaps. Specifically, if $(\text{node}(y), \text{node}(fy)) = (\text{node}(z), \text{node}(fz))$, then the edges $\{y, ey\}$ and $\{z, ez\}$ run between the same nodes, forcing $y = z$.

The face maps are compatible on the two hypermaps by construction. Finally, the edge map is compatible because

$$(\text{node}(ey), \text{node}(f(ey))) = (\text{node}(fy), \text{node}(y)) = e - \text{list}(\text{node}(y), \text{node}(fy)).$$

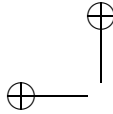
□

Lemma 8.24 [JXBJOAB] [formal proof by tchales, 2/2014]. *If L is any list of lists, then there is a iso-list L' whose elements are natural numbers.*

Proof Pick any injective map from the elements of L to \mathbb{N} and take the resulting iso-list. □

For the correspondence with Isabelle, we will need to use lists L whose elements are natural numbers and whose elements are ordered in a particular way. By taking a permutation on the natural numbers we can obtain any convenient ordering. A particular ordering is given by lemma SHXWKXQ. We refer to the formal statement in the HOL Light files. What follows is an informal summary.

Lemma 8.25 [SHXWKXQ]. *For any L (good-list, good-list-nodes, all-uniq)*



and normal-list N . By passing to an iso-list, we can arrange that the elements appearing in N are the natural numbers $\{0, 1, \dots, n-1\}$, where n is the number of vertices visited by N . Furthermore, if x is any dart, and F is the face in L containing x listed so that x appears at the head of the list. Let F' be the filtered list of F containing all the darts of F that are not in N . Then we may also assume that $F' = [n+1; n+2; n+k]$, where k is the length of F' .

The following lemma gives a reduction of the main result LSKOKJE.

Lemma 8.26 (JCAJYDU) [formal proof by HLT, 2/28/2014]. *It is enough to prove LSKOKJE in the case when H is the hypermap of L , where L is a good-list, good-list-nodes, and having elements in the natural numbers.*

8.6 translating notions between hypermaps and lists

We skip the following proofs. In each case the definitions involved in the correspondence are direct translations of one another. Thus, the proofs are a matter of comparing definitions.

Lemma 8.27 [GZLJIGN] [formal proof by tch 3/9/2014]. *Under the correspondence between L and its hypermap, the constants $l'm, l'p, l'q, l'y, l'z$ correspond respectively to $hyp'm, hyp'p, hyp'q, hyp'y, hyp'z$.*

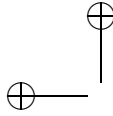
Lemma 8.28 [EVNAPDQ] [formal proof by tch 2/28/2014]. *Let L be a restricted, good-list and good-list-nodes. Let N be a normal list of L . Then the loop family is a normal family of the hypermap of L .*

Lemma 8.29 [ABKCJWD] [formal proof by tch 3/8/2014]. *Let L be a good-list. Let N be a normal list of L . Then the hypermap of the quotient-list is isomorphic to the quotient of the data on the hypermap side.*

Proof Their respective definitions are direct translations of one another. \square

Lemma 8.30 [ODWAFRG] [formal proof by tch 3/8/2014]. *Let L be a good-list and (L, N, r, x) a marked-list. Then the corresponding data in the hypermap is is-marked.*

There are also specific lemmas about hypermaps that we wish to translate back into theorems about lists. Specifically, we have lemmas HQYMRTX1-list, HQYMRTX2-list, HQYMRTX3-list, QRDYXYJ-list, AQIUNPP1-list, and AQIUNPP2-list that are list versions of lemmas given in *Dense Sphere Packings*. We omit the proofs here, because the proofs are direct translations from



hypermap language to list language. They have been deduced from the hypermap versions of the lemmas using transform-assumption-v2.

Lemma 8.31 (transform-assumption-v2) [AQ] *Let L be good-list, good-list-nodes, with marked list (L, N, r, x) and transform (N', r') . Then marked-list (L, N', r', x) .*

8.7 dihedral initialization

The fgraph of Seed p is the list of lists $[[0; 1; \dots; n + 2]; [n + 2; n + 1; \dots; 0]]$.

Lemma 8.32 (good-list-seed, all-uniq-seed) [FOEGZEQ] [formal proof by tch]. *Seed p is a good-list and all-uniq.*

Lemma 8.33 [ENWCUED] [formal proof by tch]. *Seed p is dih2k-list.*

Lemma 8.34 [TAGYMWJ] [formal proof by tch]. *Seed p is a good-list-nodes. More generally, for any L that is dih2k-list, it is good-list-nodes.*

Lemma 8.35 [formal proof by tch, 3/14/2014]. *If L is dih2k-list, then it is all uniq.*

Lemma 8.36 [formal proof by tch, 3/14/2014]. *If L is dih2k-list, then it is good-list.*

The following is not needed:

Lemma 8.37 [DLWOJBB] *If L is dih2k-list, then its hypermap is dih2k.*

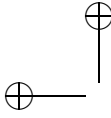
Lemma 8.38 [KUKASGD] [formal proof by tch 3/14/2014]. *Any two dih2k-list's L and L' of the same face size are iso-lists.*

Lemma 8.39 [UNHQYQM] [formal proof by tch, 3/14/2014]. *If L and L' are iso-lists and L is dih2k-list, then so is L' .*

This is a version of AUQTZYZ for lists. It can be deduced from the hypermap version.

Lemma 8.40 (AUQTZYZ-list) *Let L be good-list with restricted hypermap. Let f be a face of L . Then there exists f' such that $[f; f']$ is normal and the quotient is $[f, \text{rev } f]$ (which is dih2k-list).*

Proof Pick f' to be the complementary path (a concatenation of complementary nodes) of f . □



The following is a more precise version of the correspondence of a restricted hypermap with the seed. We will apply this lemma to a face with the largest size of L .

Lemma 8.41 [UYOUIXG] *Let L be a good-list with restricted hypermap. Let f be any face of L . Then there exists data (L', N') and f' such that $(L, [f; f'])$ is an iso-list of (L', N') and such that N' is a normal-list of L' and match-quotient-list between the seed and (L', N') .*

Proof This is obtained by renaming the elements of L to agree exactly with those of the seed. \square

8.8 termination

We will obtain termination after a finite sequence of iterations of `subdivFace`. Each `subdivFace` will increase the number of final faces by one. The number of final faces is no more than the number of faces in the original. Finally, if every face is final, then every loop in N is also final, and the quotient list of L by N is equal to L .

This is not needed. We will count final faces rather than darts.

Lemma 8.42 [ADACDYF] *Let (L, N, r, x) be a marked-list in which r is not final. Let (N', r') be the transform of the marked-list. Then the number of darts in the quotient of L by N is less than the number of darts in the quotient of L by N' .*

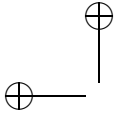
Proof The number of darts in a quotient hypermap is equal to the number of quotient darts (atoms) in the normal family N . The normal family N' replaces a loop r with r_1 and r_2 and keeps the other loops the same. The atoms of r at y and z are split into two atoms to create r_1 and r_2 . Hence the number increases. \square

Lemma 8.43 [ZBHENEI] *The number of final faces in a quotient is no more than the number of faces in the original L .*

Proof Each final face in the quotient is a face in the original L , up to rotation. \square

The following upper bound on the number of elements that are final in an lgraph, combined with a monotonicity result for g , will give termination of the procedure.

Lemma 8.44 [XZAJELF] *If we have a match-quotient-list (g, L, N) , then the number of final faces in g is at most the number of faces in L .*



Proof The set of elements that are final in g injects into the set of darts in the quotient. Then apply ZBHENEI. \square

We can continue to subdivFace until every face is final. Note that match-quotient-list allows for there to be loops that are final in N , whose corresponding face in g is not final. In this case, the iteration of subdivFace has no effect other than to make the face final. In this case, no transform of N occurs.

The following lemma shows that when we terminate, the fgraph of g will equal L .

Lemma 8.45 [XWCNBMA] *If (L, N, r, x) is a marked-list, and if every path of N is a final-list, then L is equal to its quotient-list by N .*

This has been proved at the level of hypermaps in STKBEPH. We can deduce the lemma from that.

Proof This is Example 4.51 (maximal normal family), translated into lists. If f is a loop in the family, it is canonically true. Hence its darts are singletons and the darts in the loop form a face. Thus, the set of darts visited by the family N is a union of faces.

By the third, property of normal family, the set of darts visited by the family N is a union of nodes. Hence the set of such darts is a connected component.

L is restricted, hence connected. Our conditions force N to be a connected set of faces, hence all of L \square

Thus, a series of transforms leads in a finite number of steps to a quotient-list that equals L .

8.9 finals and nonFinals

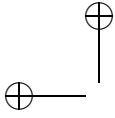
This section shows that showDups and hideDups are inverse operations, enum-of-VertexList and indexToVertexList are inverse operations, mk-triple and dest-triple are inverse operations. It analyzes containsDuplicateEdge and gives an alternative to subdivFace.

We also analyze enumerator and simplify formulas for splitFace.

8.10 enumeration lists

8.11 index calculus of higher transforms

We establish a common notation for the following lemmas.



Let $\iota(y, z, s)$ be the index of the dart z in the list s relative to the base point y (where both y and z are members of s).

Lemma 8.46 [QYHXIVZ] [formal proof by tchales, 2/2014]. [formal note: indexf-refl, indexf1, next-el-indexf, indexf-rotn, indexf-rotn, indexf-add-right, indexf-add-left, indexf-add-sum, indexf-antisym, size-betwn, indexl-betwn, indexf-betwn-eq, next-el-indexf, indexf-n, betwn-cases, indexf-add-betwn]. *The index satisfies the following properties. Assume that a, b, c are members of s , and that s is uniq. Let n be the length of s .*

- $\iota(a, a, s) = 0$.
- ι is invariant under rotations of s .
- $\iota(a, b, s) + \iota(b, c, s) = \iota(a, c, s)$ if $\iota(b, c, s) \leq \iota(a, c, s)$.
- $\iota(a, b, s) + \iota(b, c, s) = \iota(a, c, s)$ if $\iota(a, b, s) \leq \iota(a, c, s)$.
- $\iota(a, b, s) + \iota(b, c, s) = \iota(a, c, s)$ if $\iota(a, b, s) + \iota(b, c, s) < n$.
- $\iota(a, b, s) + \iota(b, a, s) = n$, if $(a = b)$.
- $n^i a = b$, where $i = \iota(a, b, s)$ and n is the next element map.
- $\iota(a, b, s) = 1$ iff b is the next element after a in s .
- b is between a and c on s iff $a \neq c$ and $0 < \iota(a, b, s) < \iota(a, c, s)$.
- We have one of the following: $b = a$, $b = c$, b is between a and c , or b is between c and a .
- if b is not between c and a on s , then $\iota(a, b, s) + \iota(b, c, s) = \iota(a, c, s)$.

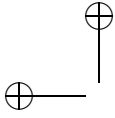
Lemma 8.47 [FWDDPHY] [formal proof by tchales, 2/2014]. [formal note: indexf-prev]. *Assume that b follows a in s and that $a \neq b$. Then $\iota(b, a, s) + 1$ is the length of s .*

Proof This is a special case of the preceeding lemma, since $\iota(a, b, s) = 1$. \square

Remark 8.48 (F, T, w, N_i , r_i , y_i , z_i) Let (L, N, r, x) be a marked list. Let F be the face of x in L . We let $w = fx$ be the dart following x in F . It will serve as a base point.

Let $T = T_{L,x}$ be the transform. Set $(N_i, r_i) = T^i(N, r)$. In particular, $(N_0, r_0) = (N, r)$. Let $y_i = l'y(L, r_i, x)$ and $z_i = l'z(L, N_i, r_i, x)$. By construction y_i and z_i are members of r_i . Also, let z_i^- be the predecessor of z_i in r_i and y_i^+ the successor of y_i in r_i .

Lemma 8.49 [DANGEYJ] [formal proof by HLT 3/2/2014]. *Let (L, N, r, x) be a marked list with notation as established above. Then every dart of N_i is a dart of N_{i+1} .*



Proof The darts of N_i in each face are the same in each face, except the face which is split. Its darts are divided between the two new faces. \square

Lemma 8.50 [PWSSRAT] [formal proof by HLT 3/5/2014]. *Let (L, N, r, x) be a marked list with notation as established above. Then a dart d of r_i appears in r_{i+1} if and only if $\iota(z_i, d, r_i) \leq \iota(z_i, y_i, r_i)$.*

(In particular, z_i appears in r_{i+1} .)

Proof This follows from the explicit ordering of darts on faces and the transform. \square

Lemma 8.51 [OHCGKFU] [formal proof by HLT 3/2014]. *Let (L, N, r, x) be a marked list with notation as established above. Assume that d is a dart in both r_i and r_{i+1} . Then*

$$\iota(z_i, d, r_i) = \iota(z_i, d, r_{i+1}).$$

Proof This follows from the explicit ordering of darts on faces and the transform. \square

Lemma 8.52 [PPLHULJ] [formal proof by HLT 3/2014]. *(L, N, r, x) be a marked list with notation as established above. Then y_i and z_i belong to F , and $\iota(w, y_i, F) < \iota(w, z_i, F)$.*

Proof This follows from the face map power applied to y_i to obtain z_i , and the result that z_i is not in the s-list (HQYMRTX1). \square

Lemma 8.53 [NCVIBWU] [formal proof by HLT 3/2014]. *(L, N, r, x) be a marked list with notation as established above. If $j \leq \iota(w, z_i, F)$, then $f^j w$ belongs to r_{i+1} and $j = \iota(w, f^j w, r_{i+1})$.*

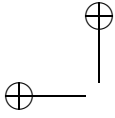
Proof This follows from the explicit ordering of darts on faces and the transform. \square

Lemma 8.54 [QCDVKEA] *(L, N, r, x) be a marked list with notation as established above. Then $\iota(w, z_i, F) \leq \iota(w, y_{i+1}, F)$.*

Proof By construction, $\iota(w, y_{i+1}, F)$ is the largest index j such that the indexing on f and r_{i+1} agree up to j . By the previous lemma, the indexing agrees at least up to $\iota(w, z_i, F)$. \square

By PWSSRAT, y_0 belongs to r_i for all i .

Lemma 8.55 [PBFLHET] *(L, N, r, x) be a marked list with notation as established above. Let d be a dart in r_i such that $\iota(z_i, d, r_i) \leq \iota(z_i, y_0, r_i)$. Then d belongs to r . Moreover, $\iota(z_i, d, r_i) = \iota(z_i, d, r)$.*



Proof By assumption d belongs to r_i . Fix such a d . By PWSSRAT, and the monotonicity of the parameters PPLHULJ and QCDVKEA, by induction, d belongs to r_{i-j} , for $j = 0, \dots, i$. In particular, d belongs to $r_0 = r$.

Taking $j = 0$, we get $d = z_i$ belongs to r .

The index calculation is also backwards induction on indices r_{i-j} , using OHCGKFU. \square

Lemma 8.56 [PNXVWFS] (L, N, r, x) be a marked list with notation as established above. For all i , we have y_i and z_i belong to r .

Proof Apply the previous lemma to $d = y_{i+1}$ and $d = z_i$, which belong to r_i by construction. \square

Lemma 8.57 [DIOWAAS] (L, N, r, x) be a marked list with notation as established above. Then $\iota(w, y_i, r) < \iota(w, z_i, r)$.

Proof This is equivalent to $\iota(z_i, w, r) \leq \iota(z_i, y_i, r)$, or to $\iota(z_i, w, r_i) \leq \iota(z_i, y_i, r_i)$. For this it is enough to show that w is not between y_i and z_i on r_i . This is clear, because the segment between y_i and z_i does not meet r_{i+1} , but w belongs to r_{i+1} . \square

Lemma 8.58 [RYIUUVK] (L, N, r, x) be a marked list with notation as established above. Then $\iota(w, z_i, r) \leq \iota(w, y_{i+1}, r)$.

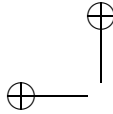
Proof The dart y_{i+1} belongs to r_{i+1} . We claim that it lies on the segment of r_{i+1} between z_i and w . Otherwise, it is between w and z_i . But from w to z_i , the list r_{i+1} coincides with f , by NCVIBWU. This is impossible, since y_{i+1} comes after z_i on f .

On the segment of r_{i+1} between z_i and w , the ordering is the same as on r . So $\iota(z_i, y_{i+1}, r_{i+1}) < \iota(z_i, w, r_{i+1})$ gives $\iota(z_i, y_{i+1}, r) < \iota(z_i, w, r)$ and hence $\iota(w, z_i, r) \leq \iota(w, y_{i+1}, r)$. \square

Lemma 8.59 [KBWPBHQ] (L, N, r, x) be a marked list with notation as established above. The set $\{y_i\}$ consists of darts y such that y on F and N such that y is followed by an n^{-1} -step on r . The set $\{z_i\}$ consists of darts z on F and N such that z is preceded by an n^{-1} -step on r .

Proof We have seen that the darts y_i and z_i lie on F and r . By construction the darts y_i are followed by an n^{-1} -step, and the darts z_i are preceded by an n^{-1} step.

Conversely, let y be any dart on F and N that is followed by an n^{-1} step on N . If y is between y_i and z_i , then these darts are not in N_i , hence not in N . This is contrary to the assumption.



If y is between z_i and y_{i+1} , then y is reached from z_i by f -steps on r_i . So they are also f -steps on r . \square

8.12 match g L and N

8.13 tame hypermaps

The formalization of this section is complete.

Lemma 8.60 [OXAXUCS] [formal proof by tch]. *If a hypermap H is isomorphic to one that has property tame k , for $k \in \{9a, 10, 11a, 11b, 12o, 13a\}$, then H has that property as well.*

Each of the tame properties $I = \{9a, 10, 11a, 11b, 12o, 13a\}$ has a corresponding definition for lgraphs, say I' . This correspondence is defined in such a way that the following holds.

Lemma 8.61 [WMLNYMD] [formal proof by tch]. *Let g be a good graph. Let H be a tame hypermap that is isomorphic to the hypermap of g . Then g is tame.*

Let H be any tame hypermap. It is restricted, so there exists a final plane-graph g and an isomorphism between H and the hypermap of g . It follows that the hypermap of g has all of the tameness properties I . Hence g itself has all of the tameness properties I' .

By the Bauer-Nipkow formalization on tame graphs, g is fgraph congruent to an fgraph y in the archive.

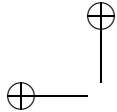
Lemma 8.62 [XRFJND0] [formal proof by Solovyev]. [formal note: See tame/good_list_archive.hl]. *Every member of the archive is a good-list.*

Proof This is by direct enumeration of the archive. \square

Lemma 8.63 [ELLLNYZ] [formal proof by tch]. *Let x and y be two good-lists that are fgraph congruent. Then their hypermaps are isomorphic, or the opposite hypermap of x is isomorphic to the hypermap of y .*

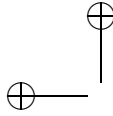
Putting these results together we have that H isomorphic to the hypermap of y or the opposite of H is isomorphic to the hypermap of y .

The main linear programming result, formalized by Solovyev, shows that if H is isomorphic to the hypermap of y , then it is not contravening. We need the opposite as well.



Lemma 8.64 [ASFUTBF][formal proof by tch]. *If the opposite of H is contravening, then H is also contravening.*

Proof If H is contravening, this means there is a finite packing V which is contravening and whose associated hypermap $\text{hyp}(V, E_{\text{std}}(V))$ is H . The finite packing $-V$ obtained by negating all the coordinates is also contravening. Its hypermap is isomorphic to the opposite of H . \square



8.14 notes on AQ

Here we discuss the proof of AQ12 and AQ13. Specifically, we discuss the proof of the second part of those statements (not involving the contour list).

Let $N' = \text{ntrans } L \ N \ r \ x \ 1$, and $r' = \text{rtrans } L \ N \ r \ x \ 1$. Write AQ12' for the goal $AQ(L, N', r', x)$ and AQ12 for the assumption $AQ(L, N, r, x)$. Similarly, for AQ13' and AQ13. We write $S = \text{s-list } L \ r \ x$, and $S' = \text{s-list } L \ r' \ x$. Note that every member of S is a member of S' .

Proof We start with the proof of AQ12'. We may assume (final-list $L \ r'$) and not(final-list $L \ r$). Otherwise, there is nothing to prove.

We let $u = \text{LAST } p$, where p is a part of s in N' . We assume that not(final-list $L \ s$), for otherwise there is nothing to prove. Our goal is to prove that e-list u is a member of final-dart-list $L \ N'$.

(A) We may assume that e-list u is not a member of r' , because otherwise e-list u is in the a final face r' , and we are done. Note that every member of S is in r' . Thus, we may also assume that e-list u is not a member of S' . Since not(final-list $L \ s$) and final-list $L \ r'$, we have $s \neq r'$; so that u is not a member of r' and that u is not a member of S' .

We have a slightly weaker statement: (A') e-list u is not a member S' , and u is not a member of S' .

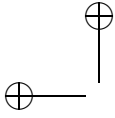
We will see that if we give a proof of AQ12' based on (A'), then exactly the same proof can be used for AQ13'. The following condition also holds:

(B) Every member of final-dart-list $L \ N$ is a member of final-dart-list $L \ N'$.

Case 1. u is a member of N' but not N . Comparing N' with N , we see that u is a new dart that is added at new nodes of degree 2 on the quotient hypermap. (These are the darts at the nodes between $l'y$ and $l'z$ along the face of x .) By (A'), u is not a member of S' . This implies that u is not in r' . It follows that u lies at a node of degree 2 on the quotient hypermap. Hence e-list u is in S' , which is contrary to (A').

Case 2. u is a member of r . Recall that u has the form $\text{LAST } p$, for p a part of s . If p is also a part of r , then using that r is not final, and that e-list u does not depend on N, N' , we see that by AQ13, e-list u is a member of S or a final-dart-list $L \ N$. By (A'), we know that e-list u is not a member of S' . So it is a member of final-dart-list N . The result follows by (B).

If u has the form $\text{LAST } p$, where p is a part of s , but not a part of r , then there are only two possibilities for u . Either it is the dart $l'y$, or the dart node-list ($l'z$). The first case $u = l'y$ is incompatible with (A'), since $l'y$ is a member of S' . The second case $u = \text{node-list } (l'z)$, gives e-list $u = f^{-1}l'z$, which is also a member of S' . This is also incompatible with (A').



Case 3. u is a member of N but not r . For such u , we have that u is final-dart-list for N iff it is in the final-dart-list for N' . Recall that s is not-final.

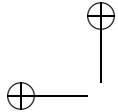
In view of (A'), the assumption AQ13 tells us that since s is a face of both N and N' , and since $\text{not}(\text{final-list } s)$, we have that e-list u is a member of the final-dart-list $L N$. By (B), e-list u is a member of final-dart-list $L N'$ as well. This completes the proof. \square

Proof Now we turn to the proof of AQ13'. We may assume that $\text{not}(\text{final-list } L r')$ and $\text{not}(\text{final-list } L r)$. Otherwise, there is nothing to prove.

We let $u = \text{LAST } p$, where p is a part of s in N' . We assume that $\text{not}(\text{final-list } L s)$, for otherwise there is nothing to prove. To avoid a trivial case, we may also assume that u is not a member of S' .

Our goal is to prove that e-list u is a member of S' or final-dart-list $L N'$. In other words, assuming that e-list u is not a member of S' , we wish to prove it is a member of final-dart-list $L N'$. In this form, we have assumption (A'), as stated in the previous proof. Assumption (B) also holds.

We have written the proof of AQ12' in such a way that the three claims from the previous proof work word-for-word to give the proof of AQ13' as well. \square



References

- [1] T. C. Hales. *Dense Sphere Packings: a blueprint for formal proofs*, volume 400 of *London Math Soc. Lecture Note Series*. Cambridge University Press, 2012.
- [2] Thomas C. Hales, Alexey Solov'yev, and Hoang Le Truong et al. The Flyspeck Project, 2014. <http://code.google.com/p/flyspeck>.