# Boundary patrol using robotic networks without localization

Topology and Minimalism RSS Workshop, June 28, 2008

#### Francesco Bullo



Center for Control, Dynamical Systems and Computation University of California at Santa Barbara http://motion.mee.ucsb.edu

Ack: Sara Susca (PhD, UCSB), Sonia Martínez (UCSD)

### Incomplete state of the art



AeroVironment Inc, "Raven" small unmanned aerial vehicle



iRobot Inc, "PackBot" unmanned ground vehicle

#### **Distributed algorithms**

automata-theoretic: "Distributed Algorithms" by N. Lynch, D. Peleg

numerical: "Parallel and Distributed Computation" by by Bertsekas and Tsitsiklis

#### **Motion coordination**

"rendezvous" by Suzuki and Yamashita

"consensus, flocking, agreement" by Jadbabaie, Olfati-Saber

"formation control" by Baillieul, Morse, Anderson

### Research directions

**Build:** distributed systems embedded actuator/sensors networks

#### **Develop distributed disciplines:**

- (i) sensor fusion
- (ii) communications
- (iii) coordinated control
- (iv) task allocation and scheduling

#### **Challenges**

- (i) scalability
- (ii) performance
- (iii) robustness
- (iv) models



Environmental monitoring



Building monitoring and evac



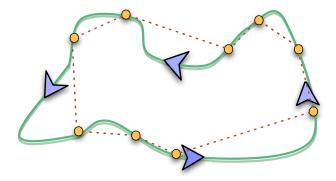
Security systems

# **Scenario 1: Boundary estimation**

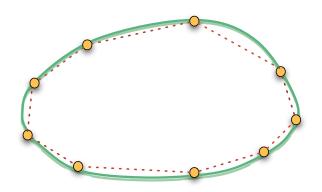
Assumption: local sensing and tracking, interpolation via waypoints

**Objective:** estimate/interpolate moving boundary

#### adaptive polygonal approximation



# **Scenario 1: Interpolation theory**



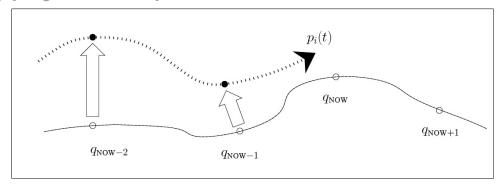
For strictly convex bodies (Gruber '80)

• sufficient condition for optimality: each two consecutive interpolation points  $p_k$ ,  $p_{k+1}$  are separated by same line integral  $\int_{p_k \to p_{k+1}} \kappa(\ell)^{1/3} d\ell$ 

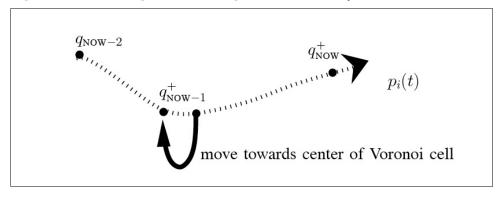
• error estimate 
$$pprox rac{1}{12n^2} \left( \int_{\partial Q} \kappa(\ell)^{1/3} d\ell \right)^3$$

# Scenario 1: Estimate-Update and Pursuit

(i) projection step

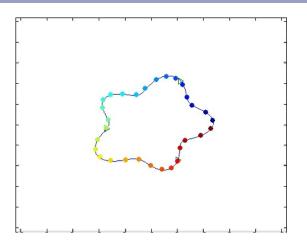


(ii) update interpolation points for "pseudo-uniform" interpolation placement



(iii) accelerate/decelerate for uniform vehicle placement

# Scenario 1: Performance/robustness



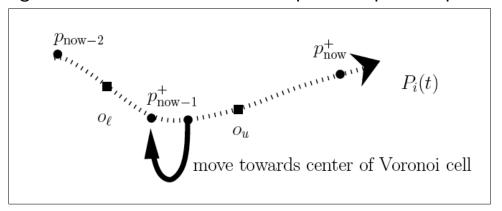
- (i) asynchronous distributed over ring
- (ii) convergence to equally distributed interpolation points and equally spaced vehicles
- (iii) time complexity: worst case  $O(n^2\log(n))$ , where  $n=\frac{\# \text{ interpolation points}}{\# \text{ vehicles}}$
- (iv) ISS robust to: evolving boundary, interpolation, sensor noise

# Scenario 1: translation into average consensus

• pseudo-distance between interpolation points  $(p_k, p_{k+1})$ 

$$d(k) = \lambda \int_{p_k \to p_{k+1}} \kappa(\ell)^{1/3} d\ell + (1 - \lambda) \int_{p_k \to p_{k+1}} d\ell$$

• "go to center of Voronoi cell" update is peer-to-peer averaging rule

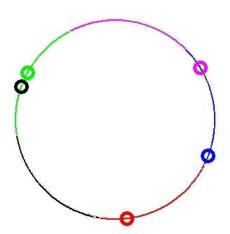


- linear model is:
  - stochastic matrices: switching, symmetric and nondegenerate
  - union of associated graphs over time is a ring (i.e., jointly connected graphs)
  - convergence rate as in Toeplitz tridagonal problem

# Scenario 2: Synchronized boundary patrolling

- (i) some UAVs move along boundary of sensitive territory
- (ii) short-range communication and sensing
- (iii) surveillance objective: minimize service time for appearing events communication network connectivity

#### Example motion:



# Analogy with mechanics and dynamics

- (i) robots with "communication impacts" analogous to beads on a ring
- (ii) classic subject in dynamical systems and geometric mechanics billiards in polygons, iterated impact dynamics, gas theory of hard spheres
- (iii) rich dynamics with even just 3 beads (distinct masses, elastic collisions) dynamics akin billiard flow inside acute triangle dense periodic and nonperiodic modes, chaotic collision sequences

SIAM REVIEW Vol. 47, No. 2, pp. 273-300 © 2005 Society for Industrial and Applied Mathematics

#### Iterated Impact Dynamics of N-Beads on a Ring\*

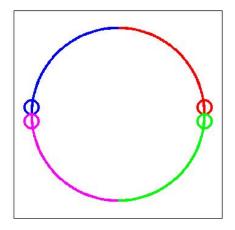
Bryan Cooley<sup>†</sup> Paul K. Newton<sup>†</sup>

Abstract. When N-beads slide along a frictionless hoop, their collision sequence gives rise to a dynamical system that can be studied via matrix products. It is of general interest to understand the distribution of velocities and the corresponding eigenvalue spectrum that a given collision sequence can produce. We formulate the problem for general N and state some basic theorems regarding the eigenvalues of the collision matrices and their products. The

# Boundary patrolling: synchronized bead oscillation

#### Desired synchronized behavior:

- starting from random initial positions and velocities
- every bead impacts its neighbor at the same point
- simultaneous impacts



# Design specification for synchronization algorithm

Achieve: asymptotically stabilize synchronized motion Subject to:

- (i) arbitrary initial positions, velocities and directions of motion
- (ii) beads can measure traveled distance, however
  no absolute localization capability, no knowledge of circle length
- (iii) no knowledge about n, adaptation to changing n (even and odd)
- (iv) anynomous agents with memory and message sizes independent of n
- (v) smooth dependency upon effect of measurement and control noise

Fully-adaptive feedback synchronization



# Slowdown-Impact-Speedup algorithm

Algorithm: (for presentation's sake, beads sense their position)

1st phase: compute average speed v and desired sweeping arcs

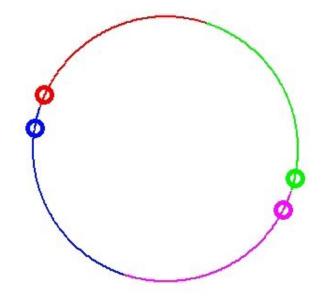
**2nd phase** for  $f \in ]\frac{1}{2}, 1[$ , each bead:

- ullet moves at nominal speed v if inside its desired sweeping arc
- ullet slows down (fv) when moving away of its desired sweeping arc hesitate when early
- when impact, change direction
- speeds up when moving towards its desired sweeping arc

# Simulations results: balanced synchronization

#### Balanced initial condition:

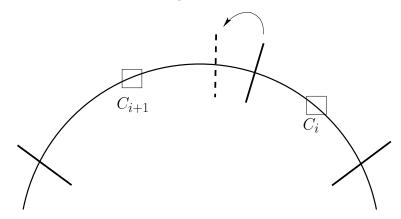
- $\bullet$  n is even
- $d_i$  is direction of motion
- $\sum_{i=1}^{n} d_i(0) = \sum_{i=1}^{n} d_i(t) = 0$
- $\bullet$  n/2 move initially clockwise



### First phase: average speed and sweeping arc

If an impact between bead i and i + 1 occurs:

- beads average nominal speeds:  $v_i^+ = v_{i+1}^+ = 0.5(v_i + v_{i+1})$
- beads change their direction of motion if  $d_i = -d_{i+1}$  (head-head type)
- beads update their desired sweeping arc



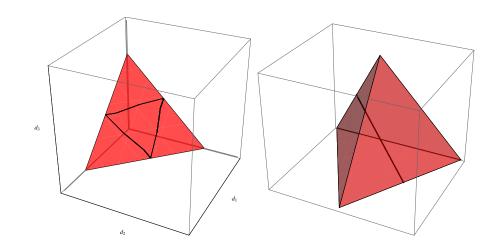
exponential average consensus

# Challenges

- (i) how to prove balanced synchronization?
- (ii) what happens for unbalanced initial conditions  $\sum_{i=1}^{n} d_i(0) \neq 0$ ?
- (iii) what happens for n is odd?
- (iv) how to describe the system with a single variable?

# Modeling detour

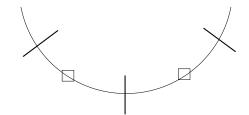
- configuration space
  - (i) order-preserving dynamics  $\theta_i \in \mathsf{Arc}(\theta_{i-1}, \theta_{i+1})$  on  $\mathbb{T}^n$
  - (ii)  $\Delta^n \times \{c, cc\}^n \times (\mathbb{R}_{>0})^n \times (arcs)^n \times \{away, towards\}^n$



- hybrid system with
  - (i) piecewise constant dynamics
  - (ii) event-triggered jumps: impact, cross boundary

### Passage and return times

ullet passage time:  $t_i^k=k$ th time when bead i passes by sweeping arc center



- return time:  $\delta_i(t) =$  duration between last two passage times
- if impact between beads (i, i + 1) at time t, then

$$\begin{bmatrix} \delta_i \\ \delta_{i+1} \end{bmatrix} (t^+) = \underbrace{\begin{bmatrix} \frac{1-f}{1+f} & \frac{2f}{1+f} \\ \frac{2f}{1+f} & \frac{1-f}{1+f} \end{bmatrix}}_{\text{stochastic (irr + aperd)}} \begin{bmatrix} \delta_i \\ \delta_{i+1} \end{bmatrix} (t^-)$$

# Convergence results: balanced synchronization

Balanced Synchronization Theorem: For balanced initial directions, assume

- (i) exact average speed and desired sweeping arcs
- (ii) initial conditions lead to well-defined 1st passage times

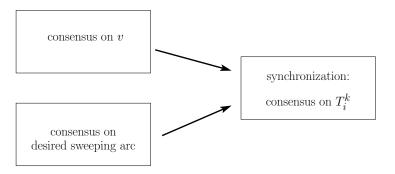
Then balanced synchronization is asymptotically stable

$$\lim_{t \to \infty} \delta(t) = \frac{2\pi}{Nv} \mathbf{1}_n, \qquad \lim_{k \to +\infty} ||T^k - \frac{\mathbf{1}_n \cdot T^k}{n} \mathbf{1}_n|| = 0$$

# Conjectures arising from simulation results

Only assumption required is balanced initial conditions.

(i) analysis of cascade consensus algorithms

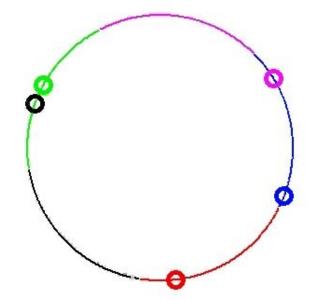


(ii) global attractivity of synchronous behavior

# Simulations results: 1-unbalanced case

#### 1-unbalanced initial condition:

- ullet n is odd
- $\sum_{i=1}^{n} d_i(0) = \sum_{i=1}^{n} d_i(t) = \pm 1$



# 1-unbalanced synchronization

(i) 
$$f \in ]\frac{1}{2}, \frac{n}{1+n}[$$

(ii) 1-unbalanced sync: beads meet at arcs boundaries 
$$\pm \frac{2\pi}{n^2} \frac{f}{1-f}$$

1-unbalanced Synchronization Theorem: For  $\sum_{i=1}^{n} d_{i}(0) = \pm 1$ , assume

- (i) exact average speed and desired sweeping arcs
- (ii) initial conditions lead to well-defined 1st passage times

Then 1-unbalanced synchronization is asymptotically stable

$$\lim_{t \to \infty} \delta(t) = \frac{2\pi}{Nv} \mathbf{1}_n, \qquad \lim_{k \to +\infty} \left( T^{2k} - T^{2(k-1)} \right) = \frac{2\pi}{v} \mathbf{1}_n$$

### General unbalanced case

Conjecture global asy-synchronization in the balanced and unbalanced case

*D*-unbalanced period orbits Theorem:

Let  $\sum_{i=1}^{n} d_i(0) = \pm D$ . If there exists an orbit along which beads i and i+1 meet at boundary  $\pm \frac{2\pi}{n^2} \frac{f}{1-f}$ , then  $f < \frac{n/|D|}{1+n/|D|}$ .

### **Emerging discipline: motion-enabled networks**

- network modeling
   network, ctrl+comm algorithm, task, complexity
- coordination algorithm
   deployment, task allocation, boundary estimation

#### Open problems

- (i) algorithmic design for motion-enabled sensor networks scalable, adaptive, asynchronous, agent arrival/departure tasks: search, exploration, identify and track
- (ii) integration between motion coordination, communication, and estimation tasks
- (iii) Very few results available on:
  - (a) scalability analysis: time/energy/communication/control
  - (b) robotic networks over random geometric graphs
  - (c) complex sensing/actuation scenarios