

Practical Reverse Engineering Exercises - Write Ups

Chapter 1 – Exercise 2 (15th of July 2014)

TASK

- “1. Given what you learned about CALL and RET, explain how you would read the value of EIP? Why can't you just do MOV EAX, EIP?”
2. Come up with at least two code sequences to set EIP to 0xAABBCCDD.
3. In the example function, addme, what would happen if the stack pointer were not properly restored before executing RET?
4. In all of the calling conventions explained, the return value is stored in a 32-bit register (EAX). What happens when the return value does not fit in a 32-bit register? Write a program to experiment and evaluate your answer. Does the mechanism change from compiler to compiler?”

Excerpt from: “*Practical Reverse Engineering: x86, x64, ARM, Windows Kernel, Reversing Tools, and Obfuscation*”,

Bruce Dang, Alexandre Gazet, Elias Bachaalany, Sebastien Josse, [ISBN: 978-1-118-78731-1](https://www.amazon.com/dp/9781118787311)

MY SHORT ANSWERS

1. According to the “[Intel® 64 and IA-32 Architectures Developer's Manual: Vol. 1](https://www.intel.com/content/www/us/en/development/intel-64-and-ia-32-architectures-developer-s-manual-vol-1.html)” the EIP register is not designed to be accessed directly by the software and could be affected implicitly only by handful of control flow instructions. In order to read the value of the EIP register one needs to execute CALL instruction to a function. CALL saves the EIP value in the stack as the function return address. While in the function body one could read the return address saved in the stack.
2. Here are some possible options for setting EIP to 0xAABBCCDD
Version A
...
MOV EAX, AABBCDDh
JMP EAX

Version B
...
MOV EAX, AABBCDDh
CALL EAX
...
Version C
...
MOV EAX, AABBCDDh
PUSH EAX
RET
3. The execution will NOT continue from the saved return address but from completely different address
4. In case a function return value exceeds 32 bits, a space is allocated to accommodate the function result and EAX is initialised with pointer to that memory.

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To verify this answer I used a C program which defines a function concatenating 2 strings with result string bigger than 32 bit. I used 2 compilers – GCC under Linux and LCC under Windows. Both compilers implemented the same mechanism.

CODE SNIPPETS

READ EIP VALUE

Here is my suggestion for reading EIP value

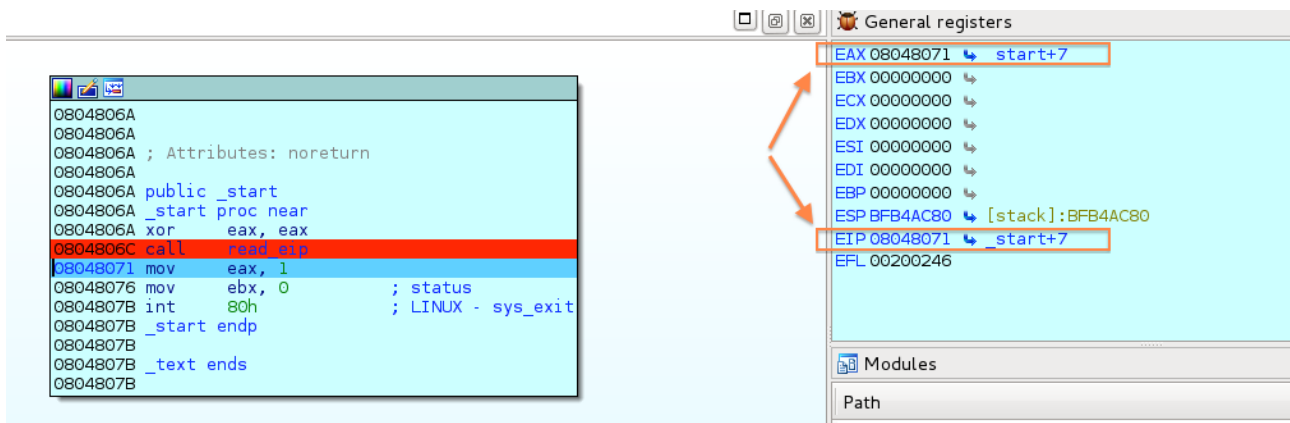
ex2_1-intel.asm listing

```
global _start
section .text
read_eip:
    push ebp
    mov ebp, esp
    mov eax, [ebp+4] ; move the return value in the EAX register
    mov esp, ebp
    pop ebp
    ret
_start:
    xor eax, eax
    call read_eip ; after returning from this function the EAX contains the EIP value
    ; exit gracefully
    mov eax, 1
    mov ebx, 0
    int 080h
```

I used IDA Debugger to trace the program execution and monitor the registers. The following screenshot illustrates a step after the **call read_eip** instruction.

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SET EIP TO 0XAABBCCDD - VERSION A

ex2_2a.asm listing

```
global _start
section .text
_start:
    mov eax, 0AABBCCDDh
    jmp eax
```

ex2_2b.asm listing

```
global _start
section .text
_start:
    mov eax, 0AABBCCDDh
    call eax
```

ex2_2c.asm listing

```
global _start
section .text
_start:
    mov eax, 0AABBCCDDh
```

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```
push eax
```

```
ret
```

HANDLING RETURN VALUES BIGGER THAN EAX SIZE

I used the following C program to verify my theory regarding how function return values bigger than 32b are handled. After compiling it with 2 different compilers I disassembled and traced the executables with IDA. The highlighted lines in the disassembled function clearly indicate that the return value stored in EAX is a pointer to a memory location holding the function result.

ex2_3.c listing

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>

char* stringConcat(char* s1, char* s2) {
    char* result = malloc(strlen(s1)+strlen(s2)+1);
    strcpy(result, s1);
    strcat(result, s2);
    return result;
}

int main() {
    char* r;
    char* a = "Test ";
    char* b = "string";
    r = stringConcat(a, b);
    return 0;
}
```

Disassembled *stringConcat* function (part of the ELF executable generated with GCC)

```
.text:080484BC ; int __cdecl stringConcat(char *s, char *src)
.text:080484BC      public stringConcat
```

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.text:080484BC stringConcat proc near ; CODE XREF: main+28p

.text:080484BC

.text:080484BC dest = dword ptr -0Ch

.text:080484BC s = dword ptr 8

.text:080484BC src = dword ptr 0Ch

.text:080484BC

.text:080484BC push ebp

.text:080484BD mov ebp, esp

.text:080484BF push ebx

.text:080484C0 sub esp, 24h

.text:080484C3 mov eax, [ebp+s]

.text:080484C6 mov [esp], eax ; s

.text:080484C9 call _strlen

.text:080484CE mov ebx, eax

.text:080484D0 mov eax, [ebp+src]

.text:080484D3 mov [esp], eax ; s

.text:080484D6 call _strlen

.text:080484DB add eax, ebx

.text:080484DD add eax, 1

.text:080484E0 mov [esp], eax ; size

.text:080484E3 call _malloc

.text:080484E8 mov [ebp+dest], eax

.text:080484EB mov eax, [ebp+s]

.text:080484EE mov [esp+4], eax ; src

.text:080484F2 mov eax, [ebp+dest]

.text:080484F5 mov [esp], eax ; dest

.text:080484F8 call _strcpy

.text:080484FD mov eax, [ebp+src]

.text:08048500 mov [esp+4], eax ; src

.text:08048504 mov eax, [ebp+dest]

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```
.text:08048507      mov     [esp], eax    ; dest
.text:0804850A      call    _strcat
.text:0804850F      mov     eax, [ebp+dest]
.text:08048512      add     esp, 24h
.text:08048515      pop     ebx
.text:08048516      pop     ebp
.text:08048517      retn
```

Disassembled *stringConcat* function (part of the PE executable generated with LCC)

```
.text:004012D4 ; int __cdecl stringConcat(char *,char *)
.text:004012D4      public _stringConcat
.text:004012D4 _stringConcat  proc near          ; CODE XREF: _main+2Ep
.text:004012D4
.text:004012D4 var_4      = dword ptr -4
.text:004012D4 arg_0      = dword ptr 8
.text:004012D4 arg_4      = dword ptr 0Ch
.text:004012D4
.text:004012D4      push     ebp
.text:004012D4      mov     ebp, esp
.text:004012D5      push     ecx
.text:004012D7      mov     ecx, 1
.text:004012D8
.text:004012DD
.text:004012DD loc_4012DD:                ; CODE XREF: _stringConcat+11j
.text:004012DD      dec     ecx
.text:004012DE      mov     [esp+ecx*4+4+var_4], 0FFFA5A5Ah
.text:004012E5      jnz     short loc_4012DD
.text:004012E7      push     esi
.text:004012E8      push     edi
.text:004012E9      push     [ebp+arg_0]    ; char *
.text:004012EC      call    _strlen
```

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```
.text:004012F1      add     esp, 4
.text:004012F4      mov     edi, eax
.text:004012F6      push    [ebp+arg_4]    ; char *
.text:004012F9      call    _strlen
.text:004012FE      add     esp, 4
.text:00401301      mov     esi, eax
.text:00401303      lea     edi, [edi+esi+1]
.text:00401307      push    edi            ; size_t
.text:00401308      call    _malloc
.text:0040130D      add     esp, 4
.text:00401310      mov     [ebp+var_4], eax
.text:00401313      push    [ebp+arg_0]    ; char *
.text:00401316      push    [ebp+var_4]    ; char *
.text:00401319      call    _strcpy
.text:0040131E      add     esp, 8
.text:00401321      push    [ebp+arg_4]    ; char *
.text:00401324      push    [ebp+var_4]    ; char *
.text:00401327      call    _strcat
.text:0040132C      add     esp, 8
.text:0040132F      mov     eax, [ebp+var_4]
.text:00401332      pop     edi
.text:00401333      pop     esi
.text:00401334      leave
.text:00401335      retn
.text:00401335 _stringConcat endp
```