

# COMP3231/9201/3891/9283 **Operating Systems 2021/T1**

UNSW

#### Administration

- Notices
- Course Outline
- UNSW Timetable
- Consultations
- Survey Results!!

#### Work

- Lectures
- Tutorials

#### Support

- Ed Forums
- Wiki

#### **Assignments**

- Submission Guide
- Assignment 0 Warm-up

#### Resources **OS/161**

- General
- Man Pages
- Sys161 Pages

#### C coding

- Info Sheet

#### **Debugging**

- Learn Debugging

#### General

- "Hardware" Guide
- R3000 Reference Manual
- Intro. to Prog. Threads

#### **Previous years**

- 2020 T2
- 2020 T1
- 2019 T1
- 2017 S1

## **Tutorial Week 7**

### Questions

## Files and file systems

- 1. Why does Linux pre-allocate up to 8 blocks on a write to a file.
- 2. Linux uses a buffer cache to improve performance. What is the drawback of such a cache? In what scenario is it problematic? What alternative would be more appropriate where a buffer cache is inappropriate?
- 3. What is the structure of the contents of a directory? Does it contain attributes such as creation times of files? If not. where might this information be stored?
- 4. The Unix inode structure contains a reference count. What is the reference count for? Why can't we just remove the inode without checking the reference count when a file is deleted?
- 5. Inode-based filesystems typically divide a file system partition into block groups. Each block group consists of a number of contiguous physical disk blocks. Inodes for a given block group are stored in the same physical location as the block groups. What are the advantages of this scheme? Are they any disadvantages?
- 6. Assume an inode with 10 direct blocks, as well as single, double and triple indirect block pointers. Taking into account creation and accounting of the indirect blocks themselves,

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- TBD (Admin)

#### Grievances

- Student Reps



what is the largest possible number of block reads and writes in order to:

- a. Read 1 byte
- b. Write 1 byte

Assume the inode is cached in memory.

- 7. Assume you have an inode-based filesystem. The filesystem has 512 byte blocks. Each inode has 10 direct, 1 single indirect, 1 double indirect, and 1 triple indirect block pointer. Block pointers are 4 bytes each. Assume the inode and any block free list is always in memory. Blocks are not cached.
  - a. What is the maximum file size that can be stored before
    - 1. the single indirect pointer is needed?
    - 2. the double indirect pointer is needed?
    - 3. the triple indirect pointer is needed?
  - b. What is the maximum file size supported?
  - c. What is the number of disk block reads required to read 1 byte from a file
    - 1. in the best case?
    - 2. in the worst case?
  - d. What is the number of disk block reads and writes required to write 1 byte to a file
    - 1. in the best case?
    - 2. in the worst case?
- 8. A typical UNIX inode stores both the file's size and the number of blocks currently used to store the file. Why store both? Should not blocks = size / block size?
- 9. How can deleting a file leave a inode-based file system (like ext2fs in Linux) inconsistent in the presence of a power failure.
- 10. How does adding journalling to a file system avoid corruption in the presence of unexpected power failures.

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