WEEK 5 TUTORIAL SOLUTIONS

PART I

Activity 1: Decide whether each of the following sentences is valid, satisfiable, or unsatisfiable. Verify your decisions using truth tables or logical equivalence and inference rules. For those that are satisfiable, list all the models that satisfy them.

a. Smoke ⇒ SmokeValid [implication, excluded middle]

b. Smoke ⇒ FireSatisfiable

Smoke	Fire	Smoke ⇒ Fire
Т	Т	Т
Т	F	F
F	Т	Т
F	F	Т

Models are: {Smoke, Fire}, {Fire}, {}.

c. (Smoke \Rightarrow Fire) \Rightarrow (\neg Smoke $\Rightarrow \neg$ Fire)

Satisfiable

Smoke	Fire	Smoke ⇒ Fire	¬ Smoke ⇒ ¬ Fire	КВ
Т	Т	Т	Т	Т
Т	F	F	Т	Т
F	Т	Т	F	F
F	F	Т	Т	Т

Models are: {Smoke, Fire}, {Smoke}, {}

d. Smoke ∨ Fire ∨ ¬ Fire Valid

e. ((Smoke \land Heat) \Rightarrow Fire) \Leftrightarrow ((Smoke \Rightarrow Fire) \lor (Heat \Rightarrow Fire)) Valid

$$((S \land H) \Rightarrow F) \qquad \Leftrightarrow (F \lor \neg (S \land H)) \qquad [implication] \\ \Leftrightarrow (F \lor \neg S \lor \neg H) \qquad [de Morgan] \\ \Leftrightarrow (F \lor \neg S \lor F \lor \neg H) \qquad [idempotent, commutativity] \\ \Leftrightarrow (S \Rightarrow F) \lor (H \Rightarrow F) \qquad [implication]$$

f. (Smoke \Rightarrow Fire) \Rightarrow ((Smoke \land Heat) \Rightarrow Fire)

Valid

$$(S \Rightarrow F) \qquad \Leftrightarrow (F \lor \neg S) \qquad [implication]$$

$$\Rightarrow (F \lor \neg S \lor \neg H) \qquad [generalization]$$

$$\Rightarrow (F \lor \neg (S \land H)) \qquad [de Morgan]$$

$$\Rightarrow ((S \land H) \Rightarrow F) \qquad [conditional]$$

g. $Big \vee Dumb \vee (Big \Rightarrow Dumb)$

Valid

$$\begin{array}{lll} \text{Big} \vee \text{Dumb} \vee (\text{Big} \Rightarrow \text{Dumb}\,) & \Leftrightarrow \text{Big} \vee \text{Dumb} \vee \neg \text{Big} & [\text{implication}] \\ & \Leftrightarrow \text{Big} \vee \neg \text{Big} \vee \text{Dumb} & [\text{idempotent}] \\ & \Leftrightarrow \text{TRUE} \vee \text{Dumb} & [\text{excluded middle}] \\ & \Leftrightarrow \text{TRUE} \end{array}$$

h. (Big ∧ Dumb) ∨ ¬ Dumb

Satisfiable

Big	Dumb	(Big ∧ Dumb)	(Big ∧ Dumb) ∨ ¬ Dumb
Т	Т	Т	Т
Т	F	F	Т
F	Т	F	F
F	F	F	Т

Models are: {Big, Dumb}, {Big}, {}

Activity 2: Determine whether the following sentences are valid (i.e. tautologies) using truth tables.

(i)
$$((P \lor Q) \land \neg P) \rightarrow Q$$

Р	Q	¬Р	P∨Q	(P∨Q)∧¬P	((P∨Q)∧¬P)→Q
Т	Т	F	Т	F	Т
Т	F	F	Т	F	Т
F	Т	Т	Т	Т	Т
F	F	Т	F	F	Т

Last column is always true no matter what truth assignment to P and Q. Therefore $((P \lor Q) \land \neg P) \rightarrow Q$ is a tautology.

(ii)
$$((P \rightarrow Q) \land \neg (P \rightarrow R)) \rightarrow (P \rightarrow Q)$$

Р	Q	R	P→Q	P→R	¬(P →R)	(P →Q)∧¬(P →R)	$((P \rightarrow Q) \land \neg (P \rightarrow R)) \rightarrow (P \rightarrow Q)$
Т	Т	Т	Т	Т	F	F	Т
Т	Т	F	Т	F	Т	Т	Т
Т	F	Т	F	Т	F	F	Т
Т	F	F	F	F	Т	F	Т
F	Т	Т	Т	Т	F	F	Т
F	Т	F	Т	Т	F	F	Т
F	F	Т	Т	Т	F	F	Т
F	F	F	Т	Т	F	F	Т

Last columnis always true no matter what truth assignment to P, Q and R. Therefore $((P \rightarrow Q) \land \neg (P \rightarrow R)) \rightarrow (P \rightarrow Q)$ is a tautology.

(iii)
$$\neg(\neg P \land P) \land P$$

Р	¬P	¬P ∧P	¬(¬P ∧P)	¬(¬P ∧P)∧P
Т	F	F	Т	Т
F	Т	F	Т	F

Last columnis not always true. Therefore $\neg(\neg P \land P) \land P$ is not a tautology.

(iv)
$$(P \lor Q) \rightarrow \neg (\neg P \land \neg Q)$$

Р	Q	P vQ	¬P	¬Q	¬P ∧¬Q	¬(¬P ∧¬Q)	((P ∨Q)→¬(¬P ∧¬Q)
Т	Т	Т	F	F	F	Т	Т
Т	F	Т	F	Т	F	Т	Т
F	Т	Т	Т	F	F	Т	Т
F	F	F	Т	Т	Т	F	Т

Last column is always true no matter what truth assignment to P and Q. Therefore $(P \lor Q) \rightarrow \neg (\neg P \land \neg Q)$ is a tautology.

Activity 3: Show using the truth table method that the corresponding inferences are valid.

(i)P
$$\rightarrow$$
 Q, \neg Q |= \neg P

P	Q	¬Q	P → Q	¬P
Т	Т	F	Т	F
Т	F	Т	F	F
F	Т	F	Т	Т
F	F	Т	Т	Т

In all rows where both $P \rightarrow Q$ and $\neg Q$ are true $\neg P$ is true. Therefore, valid inference.

(ii)
$$P \rightarrow Q \models \neg Q \rightarrow \neg P$$

Р	Q	¬P	¬Q	$P \rightarrow Q$	$\neg Q \rightarrow \neg P$
Т	Т	F	F	Т	Т
Т	F	F	Т	Т	Т
F	Т	Т	F	F	F
F	F	Т	Т	Т	Т

In all rowswhereboth $P \rightarrow Q$ true, $\neg Q \rightarrow \neg P$ is true. Therefore, valid inference.

(iii) P
$$\rightarrow$$
 Q, Q \rightarrow R |= P \rightarrow R

Р	Q	R	P → Q	$Q \rightarrow R$	$P \rightarrow R$
Т	Т	Т	Т	Т	Т
Т	Т	F	Т	F	F
Т	F	Т	F	Т	Т
Т	F	F	F	Т	F
F	Т	Т	Т	Т	Т
F	Т	F	Т	F	Т
F	F	Т	Т	Т	Т
F	F	Т	Т	Т	Т

In all rows where both $P \rightarrow Q$ and $Q \rightarrow R$ true, P R is true. Therefore, valid inference.

PART II

Activity 1: Consider the following Knowledge Base of facts:

If the unicorn in mythical, then it is immortal, but if it is not mythical, then it is mortal and a mammal. If the unicorn in either immortal or a mammal, then it is horned. The unicorn is magical if it is horned.

1. Translate the above statements into Propositional Logic

Myth	The unicorn is Mythical
Mortal	The unicorn is Mortal
Mammal	The unicorn is a Mammal
Horned	The unicorn is Horned
Magic	The unicorn is Magical

Myth ⇒ ¬ Mortal ¬Myth ⇒ (Mortal ∧ Mammal) ¬Mortal ∨ Mammal ⇒ Horned

Horned ⇒ Magic

2. Convert this Knowledge Base into Conjunctive Normal Form.

 $(\neg Myth \lor \neg Mortal) \land (Myth \lor Mortal) \land (Myth \lor Mammal) \land (Mortal \lor Horned) \land (\neg Mammal \lor Horned) \land (\neg Horned \lor Magic)$

3. Use a series of resolutions to prove that the unicorn is Horned. Using Proof by Contradiction, we add to the database the negative of what we are trying to prove: ¬Horned

We then try to derive the "empty clause" by a series of Resolutions:

¬Horned ∧ (Mortal ∨ Horned)

Mortal

¬Horned ∧ (¬Mammal ∨ Horned)

¬Mammal

Mortal ∧ (¬Myth ∨ ¬Mortal)

¬Myth

¬Myth ∧ (Myth ∨ Mammal)

Mammal

Mammal ∧ ¬Mammal

Having derived the empty clause, the proof (of Horned) is complete.

4. Give all models that satisfy the Knowledge Base. Can you prove that the unicorn is Mythical? How about Magical? Because of the rule (Horned ⇒ Magic), Magic must also be True.

We can construct a truth table for the remaining three variables:

Myth	Mortal	Mammal	Myth ⇒ ¬ Mortal	¬ Myth ⇒ (Mortal ∧ Mammal)	KB
Т	Т	Т	F	Т	F
Т	Т	F	F	Т	F
Т	F	Т	Т	Т	Т
Т	F	F	Т	Т	Т
F	Т	T	Т	Т	Т
F	Т	F	Т	F	F
F	F	Т	Т	F	F
F	F	F	Т	F	F

There are three models which satisfy the entire Knowledge Base: {Horned, Magic, Myth, Mammal}, {Horned, Magic, Myth}, {Horned, Magic, Mortal, Mammal}

(Note: each model is described by listing which variables are True)
We cannot prove that the unicorn is Mythical, because of the third model where Mythical is False.

PART III

Activity 1: Discuss your answers from the activity on the First Order Logic page: Represent the following sentences in first-order logic, using a consistent vocabulary.

- Some students studied French in 2015.
- Only one student studied Greek in 2014.
- The highest score in Greek is always higher than the highest score in French.
- · Every person who buys a policy is smart.
- · No person buys an expensive policy.
- · There is a barber who shaves all men in town who do not shave themselves.
- Politicians can fool some of the people all of the time, and they can fool all of the people some of the time, but they can't fool all of the people all of the time. (Use Fool(p,x,t) to mean that p fools x at time t).
- a. Some students studied French in 2016. $\exists x \ Student(x) \land \ Study(x,French,2016)$
- b. Only one student studied Greek in 2015. $\exists x \; Study(x,Greek,2015) \land \forall y \; (Study(y,Greek,2015) \Rightarrow y=x)$ sometimes written as $\exists !x \; Study(x,Greek,2015)$
- c. The highest score in Greek is always higher than the highest score in French. $\forall t \exists x \forall y \ Score(x,Greek,t) > Score(y,French,t)$
- d. Every person who buys a policy is smart. $\forall x, p \ Person(x) \land Policy(p) \land Buy(x,p) \Rightarrow Smart(x)$
- e. No person buys an expensive policy. $\neg \exists x, p \ Person(x) \land Policy(p) \land Expensive(p) \land Buy(x,p)$
- f. There is a barber who shaves all men in town who do not shave themselves. $\exists b \; Barber(b) \land \forall m \; (Man(m) \land InTown(m) \land \neg Shave(m,m) \Rightarrow Shave(b,m))$
- g. Politicians can fool some of the people all of the time, and they can fool all of the people some of the time, but they can't fool all of the people all of the time. $\forall p \ (\exists x \forall t \ Fool(p,x,t)) \land (\exists t \forall x \ Fool(p,x,t)) \land (\neg \forall x \forall t \ Fool(p,x,t))))$