# cryptohack RSA 部分wp by crumbling

# 目录

- RSA
  - STARTER
    - 1. RSA Starter 1
    - 2. RSA Starter 2
    - 3. RSA Starter 3
    - 4. RSA Starter 4
    - 5. RSA Starter 5
    - 6. RSA Starter 6
  - PRIMES PART 1
    - 1. Factoring
    - 2. Inferius Prime
    - 3. Monoprime
    - 4. Square Eyes
    - 5. Manyprime
  - PUBLIC EXPONENT
    - 1. Salty
    - 2. Modulus Inutilis
    - 3. Everything is Big

# **RSA**

# **STARTER**

#### **RSA Starter 1**

```
print(pow(101,17,22663))
```

## **RSA Starter 2**

用小数实现加密

```
e=65537
q=23
p=17
N=p*q
m=12
print(pow(m,e,N))
```

#### **RSA Starter 3**

求欧拉函数。

关于公式,直接贴个wiki吧

欧拉函数 - 维基百科, 自由的百科全书 (wikipedia.org)

```
p = 857504083339712752489993810777
q = 1029224947942998075080348647219
print((p-1)*(q-1))
```

## **RSA Starter 4**

在已知q, p, e的情况下求私钥d

```
import gmpy2
p = 857504083339712752489993810777
q = 1029224947942998075080348647219
e = 65537
r=(q-1)*(p-1)
d=gmpy2.invert(e,r)
print(d)
```

#### **RSA Starter 5**

在已知q, p, e的情况下解密RSA

```
N = 882564595536224140639625987659416029426239230804614613279163
e = 65537
c = 77578995801157823671636298847186723593814843845525223303932
p=857504083339712752489993810777
q=1029224947942998075080348647219
r=(q-1)*(p-1)
d=gmpy2.invert(e,r)
m=pow(c,d,N)
print(m)
```

#### **RSA Starter 6**

利用RSA对flag进行签名

```
from hashlib import sha256
flag="crypto{Immut4ble_m3ssag1ng}"
15216583654836731327639981224133918855895948374072384050848479908982286890731769
48660908591885766404607537525316895505874318566439027305807445039023677432490330
56634790465662329672977657316253280298140556353160025912275702712714452260949198
64475407884459980489638001092788574811554149774028950310695112688723853763743238
75334978250812198533874675523781937317869934313509178399229956182738974513288002
22598733875242732988503406487798979093819797140268371720039532210524312179406325
52930880000919436507245150726543040714721553361063311954285289857582079880295199
632757829525723874753306371990452491305564061051059885803
11175901210643014262548222473449533091378848269490518850474399681690547281665059
31715583169230045319733573572845925939236682330240568538958688367004374468399370
91231808051546310885135214569793176280127218815371541072393894660631360073371205
99915456659758559300673444689263854921332185562706707573660658164991098457874495
05485449147406503962192297267158829931584630606984516995945125082104441788663034
62290213054103401004015301461354188065443409083551065820890829805336510955941920
31411679866134256418292249592135441145384466261279428795408721990564658703903787
956958168449841491667690491585550160457893350536334242689
s=sha256(flag.encode())
b=s.hexdigest()
b=int(b,16)
enc=pow(b,d,N)
print(hex(enc))
```

#### **PRIMES PART 1**

## **Factoring**

http://factordb.com/index.php?id=

对于不大的n可以用这个网站分解。

不过大部分时候都没法直接分解,所以只能用来做一些非常基础的RSA分解题目,或是其他用途。

## **Inferius Prime**

n200位,同样网站分解。

```
from Crypto.Util.number import*
n = 742449129124467073921545687640895127535705902454369756401331
e = 3
ct = 39207274348578481322317340648475596807303160111338236677373
p=752708788837165590355094155871
q=986369682585281993933185289261
r=(p-1)*(q-1)
d=inverse(e,r)
flag=pow(ct,d,n)
flag=long_to_bytes(flag)
print(flag)
```

# Monoprime

n不由p、q两个素数相乘得到,而直接是一个素数。

这是一个非常危险的做法,虽然"无法分解",但因为欧拉函数就是n-1,所以反而更加容易求解。

```
from Crypto.Util.number import inverse, long_to_bytes
from gmpy2 import mpz
17173137121806544412548253630224591541560331838028039238529183647229975274793460
72464775085078272840757639102649953260102512684936305019898108554184166433526311
505861203914449338007760990051788980485462592823446469606824421932591
e = 65537
ct =
16136755034673060445145475618902893896494128034766209879877546601946337561070007
48401057768737916050700925546501904860303671210115781715257596007747398904584145
93857709994072516290998135846956596662071379067305011746842247628316996977338024
343628757374524136260758515864509435302781735938531030576289086798942
r=n-1
d=inverse(e,r)
flag=pow(ct,d,n)
flag=long_to_bytes(flag)
print(flag)
```

#### **Square Eyes**

n很大,但是因数是2个p。开根后求欧拉函数即可。

```
from Crypto.Util.number import long_to_bytes
import gmpy2
```

n =

e = 65537

ct =

p=gmpy2.iroot(n,2)[0]

 $\begin{array}{l} \textbf{p=}231486675219980977208571688277907713376624837163484354773605674093550261691659\\ 34446949809664595523770853897203103759106983985113264049057416908191166720008503\\ 27595162573897566601902917237765317060244037357959329257653066777395140764722275\\ 77564378672160951931742013232788960272945177926078818618552646005257724607452594\\ 40301156930943255240915685718552334192230264780355799179037816026330705422484000\\ 08654236208400695815855034639594186238392594203373003000460636030837977625543620\\ 64405294417118592468115866527460284184960201454415130375354753809625621089206999\\ 29022900677901988508936509354385660735694568216631382653107 \end{array}$ 

```
r=p*(p-1)
d=gmpy2.invert(e,r)
m=pow(ct,d,n)
print(long_to_bytes(m))
```

#### Manyprime

n很大,但因为因数较多各个因数较小,所以可以网站可以直接分解。根据公式求出欧拉函数就可以成功求解。

```
from Crypto.Util.number import long_to_bytes
import qmpy2
n =
58064239189884319292956385687089779965088315271876176293229248225215259127987142
15691620371904190364350417977398803895295936744855557922349009694020190556017816
62044515999210032698275981631376651117318677368742867687180140048715627160641771
11804037257357547933083009298980073010557370055771714625186058880250931053479231
07488985043949662638199599632735091197910375255044226066346401732775987748140995
40555569257179715908642917355365791447508751401889724095964924513196281345665480
68802963999947264954916314759954014236757541388572965316651759571999187222301196
9856259344396899748662101941230745601719730556631637
e = 65537
ct =
32072149053462443414999372352732297796055651075062835485626073209810969258133840
99999833761313549183700476251504547287184679988703223449809856351496569777879643
80651868131740312053755501594999166365821315043312308622388016666802478485476059
62588803301719808347297601171999833398553175697867875889747284535816773022150657
38177984671000237547091092742658352017573698297441132336073595264410075778501112
28850004361838028842815813724076511058179239339760639518034583306154826603816927
75723654909633950150331660107889128740868209975016472003297501681418789939927371
9181407940397071512493967454225665490162619270814464
p=
[9282105380008121879,9303850685953812323,9389357739583927789,1033665022087849984
1,10638241655447339831,11282698189561966721,11328768673634243077,114034606390362
43901,11473665579512371723,11492065299277279799,11530534813954192171,11665347949
879312361.
12132158321859677597, 12834461276877415051, 12955403765595949597, 12973972336777979
701,13099895578757581201,13572286589428162097,14100640260554622013,1417886959219
3599187,14278240802299816541,14523070016044624039,14963354250199553339,153645975
61881860737, 15669758663523555763, 15824122791679574573, 15998365463074268941, 16656
402470578844539,16898740504023346457,17138336856793050757,17174065872156629921,1
7281246625998849649]
r=1
for i in p:
    r=r*(i-1)
d=gmpy2.invert(e,r)
m=pow(ct,d,n)
print(long_to_bytes(m))
```

#### **PUBLIC EXPONENT**

#### Salty

e为1, 虽然加密时很快, 但解密时更快。

```
from Crypto.Util.number import long_to_bytes,getPrime
n =
11058179571595856620660039216136021257966963739143709770368515423701735157046476
77253241820511999019203182112904047772597289236149172112915625558647530051793261
01890427669819834642007924406862482343614488768256951616086287044725034412802176
312273081322195866046098595306261781788276570920467840172004530873767
e = 1
ct =
44981230718212183604274785925793145442655465025264554046028251311164494127485
m=pow(ct,e,n)
print(long_to_bytes(m))
```

#### **Modulus Inutilis**

n很大, 但e=3很小, 有m^3=ct<n的可能

```
from Crypto.Util.number import long_to_bytes
import gmpy2
n =
17258212916191948536348548470938004244269544560039009244721959293554822498047075
64341564678119696284734253656080029070612417146882045651701461178697429102730649
09154666642642308154422770994836108669814632309362483307560217924183202838588431
34262255159849974736977129510589035929007314633067738334112124236636830912685009
43715250787494968505200750156367164900874821936035625015773485712562109917320712
82478547626856068209192987351212490642903450263288650415552403935705444809043563
866466823492258216747445926536608548665086042098252335883
e = 3
ct =
24325105361790376030994184483541129237335065597307548026400135291986518015122218
98204733584110377593813286429573248895191923371523553028084006380526205804098132
22660643570085177957
m=gmpy2.iroot(ct,3)[0]
print(long_to_bytes(m))
```

#### **Everything is Big**

e很大时可以考虑wiener's Attack

捡了个脚本

太久了稍微找了下,应该是来自<u>BugkuCTF RSA(wiener's attack) bugkectf rsa ProboxDu的博客-CSDN</u> <u>博客</u>

```
from __future__ import print_function
import libnum

def continued_fractions_expansion(numerator, denominator): # (e,N)
    result = []

    divident = numerator % denominator
    quotient = numerator // denominator
    result.append(quotient)
```

```
while divident != 0:
        numerator = numerator - quotient * denominator
        tmp = denominator
        denominator = numerator
        numerator = tmp
        divident = numerator % denominator
        quotient = numerator // denominator
        result.append(quotient)
    return result
def convergents(expansion):
    convergents = [(expansion[0], 1)]
    for i in range(1, len(expansion)):
        numerator = 1
        denominator = expansion[i]
        for j in range(i - 1, -1, -1):
            numerator += expansion[j] * denominator
            if j == 0:
                break
            tmp = denominator
            denominator = numerator
            numerator = tmp
        convergents.append((numerator, denominator)) # (k,d)
    return convergents
def newtonSqrt(n):
    approx = n // 2
    better = (approx + n // approx) // 2
    while better != approx:
        approx = better
        better = (approx + n // approx) // 2
    return approx
def wiener_attack(cons, e, N):
    for cs in cons:
        k, d = cs
        if k == 0:
            continue
        phi_N = (e * d - 1) // k
        \# x^{**2} - ((N - phi_N) + 1) * x + N = 0
        a = 1
        b = -((N - phi_N) + 1)
        C = N
        delta = b * b - 4 * a * c
        if delta <= 0:</pre>
            continue
        x1 = (newtonSqrt(delta) - b) // (2 * a)
        x2 = -(newtonSqrt(delta) + b) // (2 * a)
```

```
if x1 * x2 == N:
    return [x1, x2, k, d]

if __name__ == "__main__":
    n =
```

0xb8af3d3afb893a602de4afe2a29d7615075d1e570f8bad8ebbe9b5b9076594cf06b6e7b30905b6 420e950043380ea746f0a14dae34469aa723e946e484a58bcd92d1039105871ffd63ffe64534b7d7 f8d84b4a569723f7a833e6daf5e182d658655f739a4e37bd9f4a44aff6ca0255cda5313c3048f56e ed5b21dc8d88bf5a8f8379eac83d8523e484fa6ae8dbcb239e65d3777829a6903d779cd2498b255f cf275e5f49471f35992435ee7cade98c8e82a8beb5ce1749349caa16759afc4e799edb12d299374d 748a9e3c82e1cc983cdf9daec0a2739dadcc0982c1e7e492139cbff18c5d44529407edfd8e75743d 2f51ce2b58573fea6fbd4fe25154b9964d

e =

0x9ab58dbc8049b574c361573955f08ea69f97ecf37400f9626d8f5ac55ca087165ce5e1f459ef6f a5f158cc8e75cb400a7473e89dd38922ead221b33bc33d6d716fb0e4e127b0fc18a197daf856a706 2b49fba7a86e3a138956af04f481b7a7d481994aeebc2672e500f3f6d8c581268c2cfad4845158f7 9c2ef28f242f4fa8f6e573b8723a752d96169c9d885ada59cdeb6dbe932de86a019a7e8fc8aeb077 48cfb272bd36d94fe83351252187c2e0bc58bb7a0a0af154b63397e6c68af4314601e29b07caed30 1b6831cf34caa579eb42a8c8bf69898d04b495174b5d7de0f20cf2b8fc55ed35c6ad157d3e7009f1 6d6b61786ee40583850e67af13e9d25be3

c =

0x3f984ff5244f1836ed69361f29905ca1ae6b3dcf249133c398d7762f5e27791917469429398914
4c9d25e940d2f66058b2289c75d1b8d0729f9a7c4564404a5fd4313675f85f31b47156068878e236
c5635156b0fa21e24346c2041ae42423078577a1413f41375a4d49296ab17910ae214b45155c4570
f95ca874ccae9fa80433a1ab453cbb28d780c2f1f4dc7071c93aff3924d76c5b4068a0371dff8253
1313f281a8acadaa2bd5078d3ddcefcb981f37ff9b8b14c7d9bf1accffe7857160982a2c7d9ee01d
3e82265eec9c7401ecc7f02581fd0d912684f42d1b71df87a1ca51515aab4e58fab4da96e154ea6c
dfb573a71d81b2ea4a080a1066e1bc3474

```
expansion = continued_fractions_expansion(e, n)
cons = convergents(expansion)
p, q, k, d = wiener_attack(cons, e, n)
m = pow(c, d, n)
print(libnum.n2s(m))
```