

# Phase Transition of the 3-Majority Dynamics with Uniform Communication Noise



Francesco d'Amore  
COATI team

Based on joint work with [I. Ziccardi](#)

Bologna  
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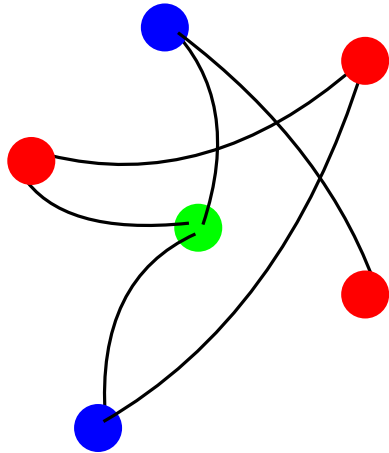
1. The agreement task in distributed computing
2. Opinion dynamics with uniform communication noise
3. 3-MAJORITY dynamics
4. Discussion and techniques

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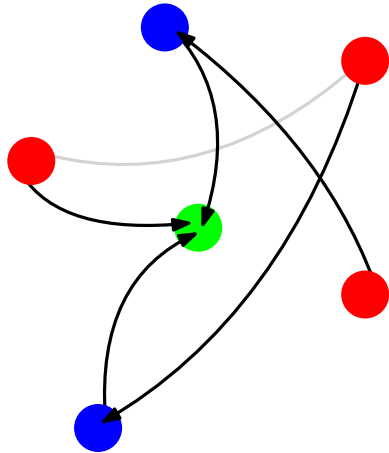
System of agents supporting different opinions



Messages are exchanged upon interaction

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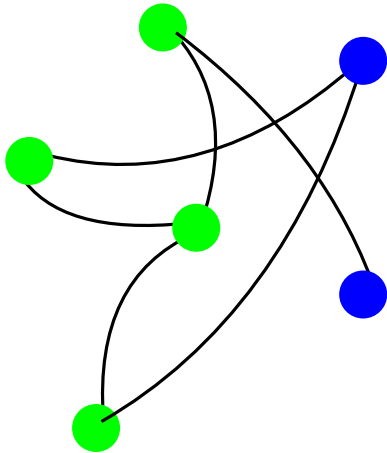
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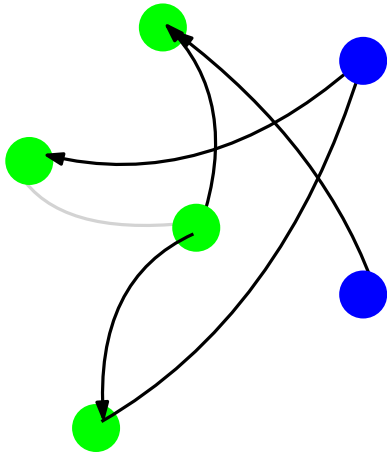
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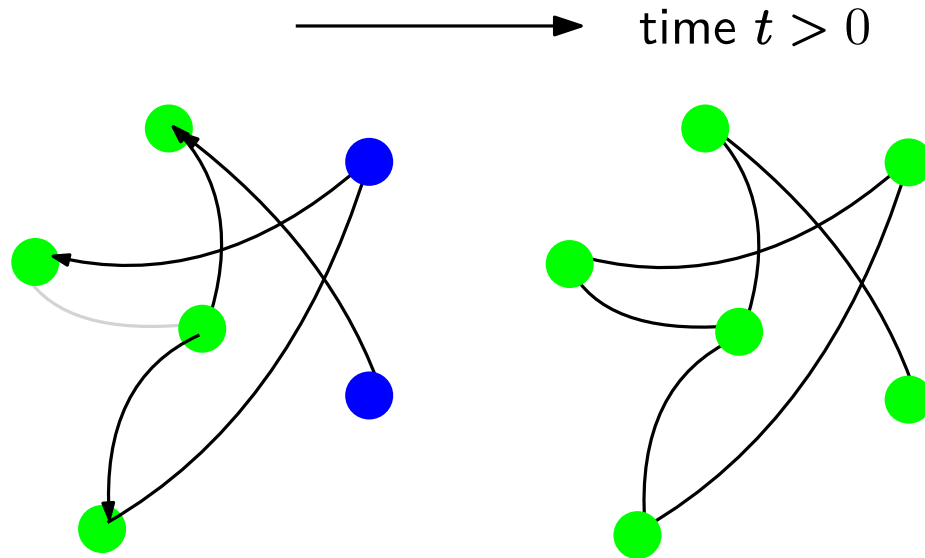
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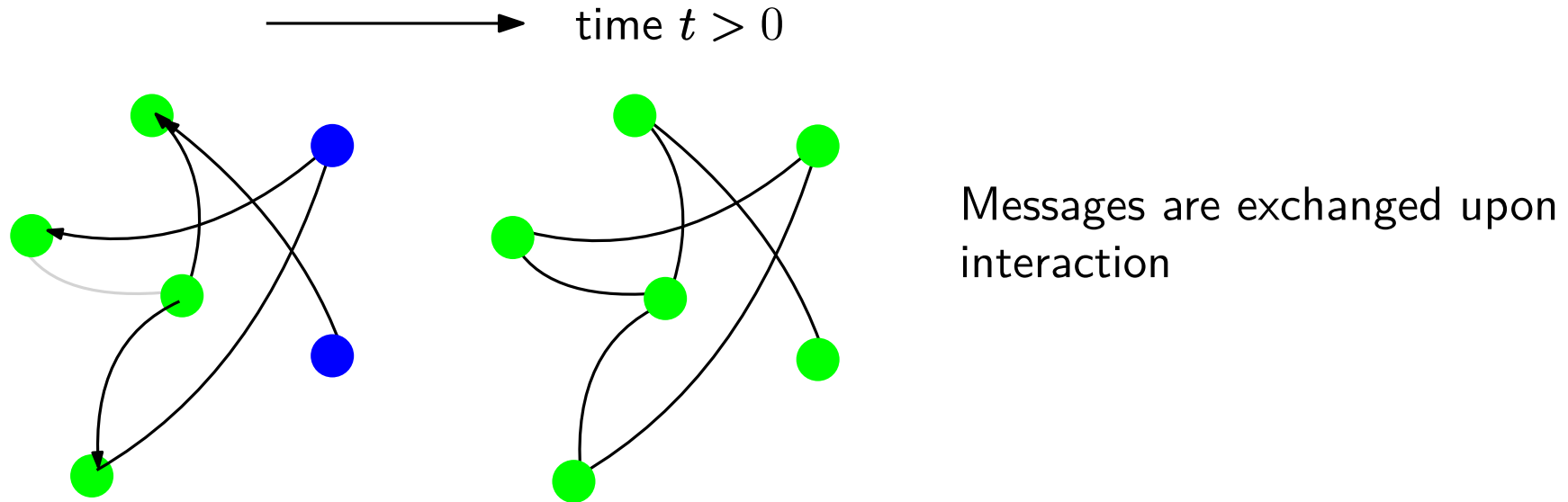
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**Reaching agreement:** fundamental task in distributed computing [Becchetti et al., 2020]



# The agreement task

System of agents supporting different opinions



**Reaching agreement:** fundamental task in distributed computing [Becchetti et al., 2020]

- social networks [Acemoglu et al., Math. Oper. Res. 2013]
- swarm robotics [Bayindir, Neurocomputing 2016]
- communication networks [Ruan et Mostofi, CDC 2008]
- distributed databases [Dietzfelbinger et al., ICALP 2010]
- biological systems [Feinerman et al., Dis. Comput. 2017]

# In biological systems

Often, systems of **interacting agents** performing **collective tasks** (swarm-like)

- Reaching agreement: molecules [Carrol, *Nature Immunology* 2004], bacteria [Bassler, *Cell* 2002], social insects [Franks et al., 2002] (e.g. bees [Reina et al., *Physical Review E* 2017])



honey bees

"Queen bee 1" by quisnovu, CC BY-NC 2.0.

# The consensus problem

**Input:** system of  $n$  agents supporting opinions, with a communication network

**Task:** designing a protocol which brings the system in finite time to a configuration such that

1. all agents support the same opinion (AGREEMENT)
2. the final opinion is among the initial ones (VALIDITY)
3. the agreement keeps on unless external events occur (STABILITY)

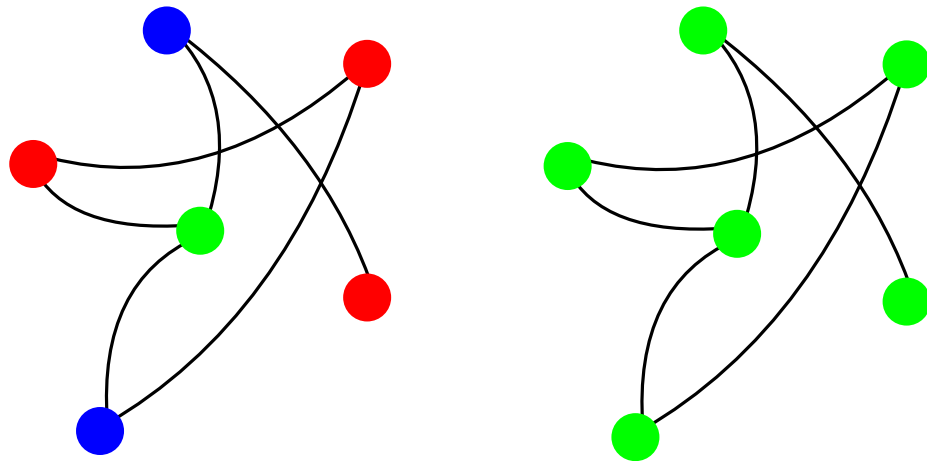
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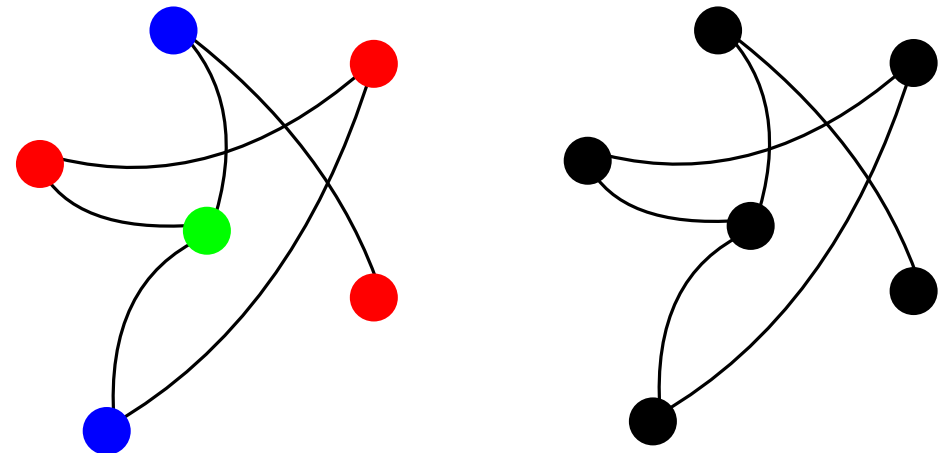
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time  $t = 0$   $\longrightarrow$  time  $t > 0$



valid consensus

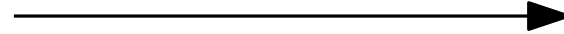
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non valid consensus

# The majority consensus problem

1. AGREEMENT
2. ~~VALIDITY~~
3. ~~STABILITY~~



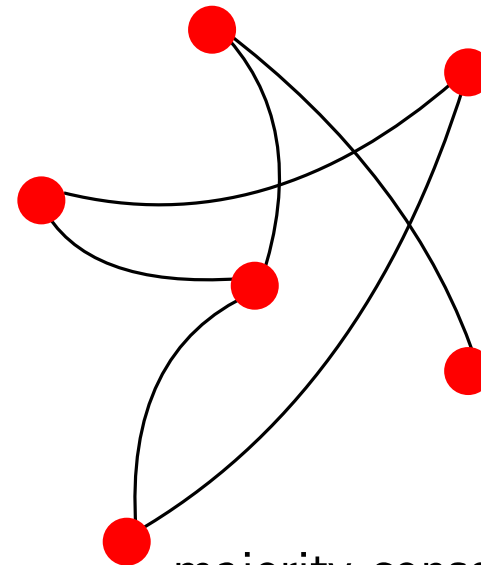
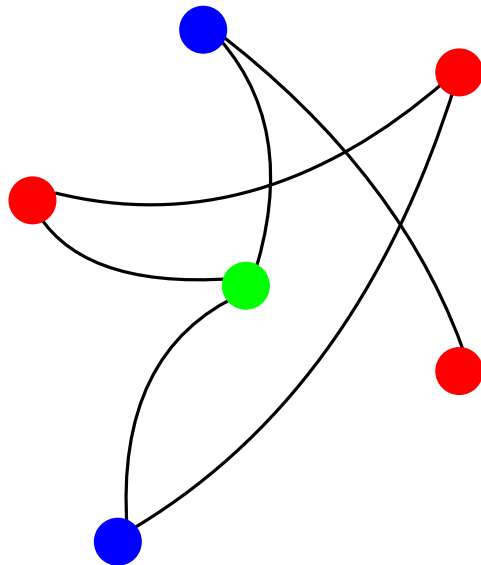
1. AGREEMENT
2. **MAJORITY**
3. STABILITY

2. **MAJORITY** property: the **final opinion** is the **initial majority** one

time  $t = 0$



time  $t > 0$



majority consensus

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# Opinion dynamics for the consensus problem

**Opinion dynamics:** class of simple, lightweight parallel protocols for the consensus problems

Very simple **update-rules**

Many have been investigated, including:

- Voter Model [[Hassin and Peleg, Inf. Comput. 2001](#)]
  - Averaging dynamics [[Becchetti et al., SODA 2017](#)]
  - 3-Majority [[Becchetti et al., SODA 2016](#)]
  - 2-Choices [[Berenbrink et al., PODC 2017](#)]
  - Undecided-State [[Becchetti et al., SODA 2015](#)]
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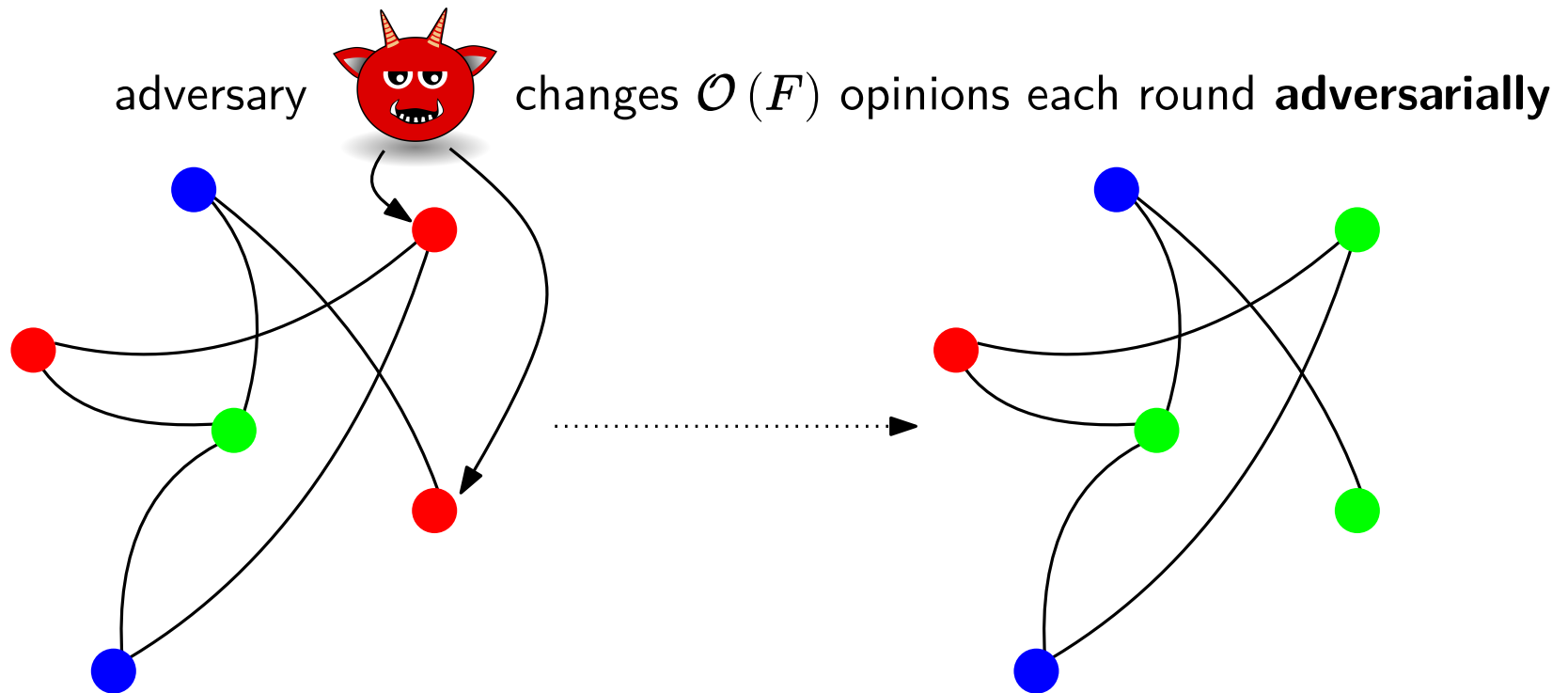
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Majority update-rules and the undecided state dynamics have **biological inspirations** [[Reina et al., Physical Review 2017](#)] [[Condon et al., Nat. Computing 2020](#)] [[Chaouiya et al., PLOS ONE 2013](#)]



# Adversarial failures vs. noise

Often, settings with **adversarial Byzantine failures** are investigated



Not realistic in biological scenarios; rather, **uniform noise** [Feinerman et al., PODC 2014]

# Uniform communication noise

Inspired by [Feinerman et al., PODC 2014], [Freignaud and Natale, PODC 2016]

$\Sigma$  set of  $k$  opinions,  $p \in [0, 1]$  constant

When  $u$  looks at  $v$ 's opinion  $x$

- a) with probability  $1 - p$ ,  $u$  sees  $x$
- b) with probability  $p$ ,  $u$  sees  $y$  where  $y$  is chosen u.a.r. in  $\Sigma$

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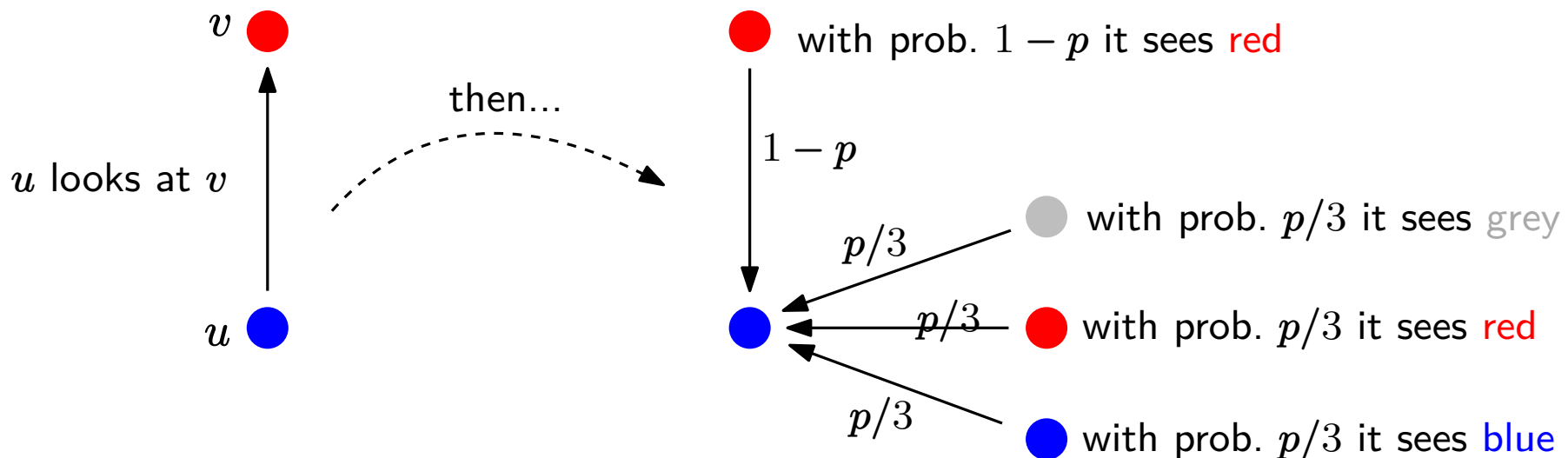
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Example:  $k = 3$



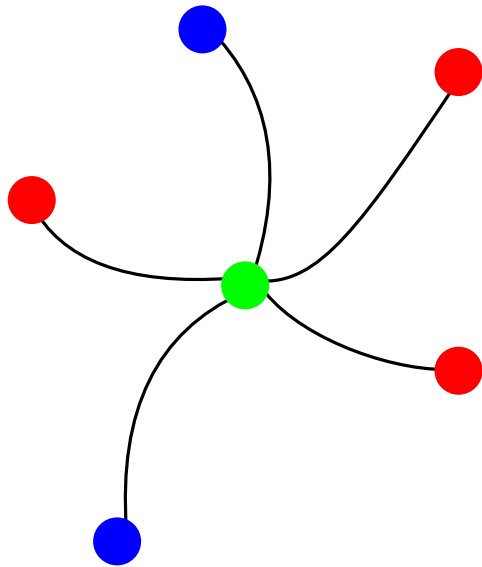
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# The dynamics

**3-Majority** dynamics: each node  $u$

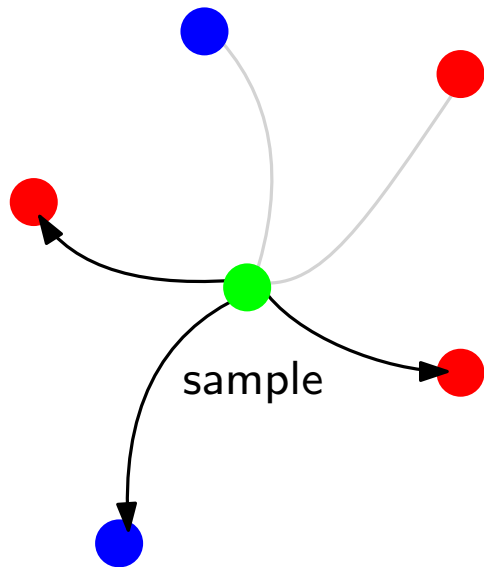
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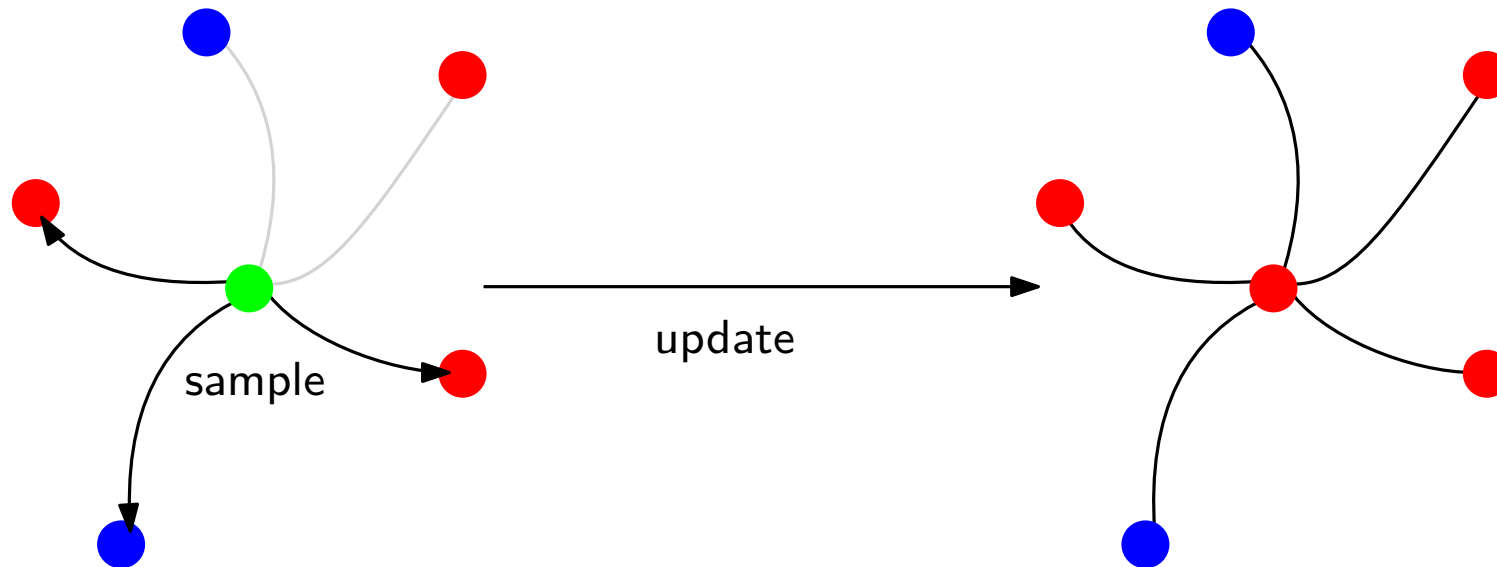
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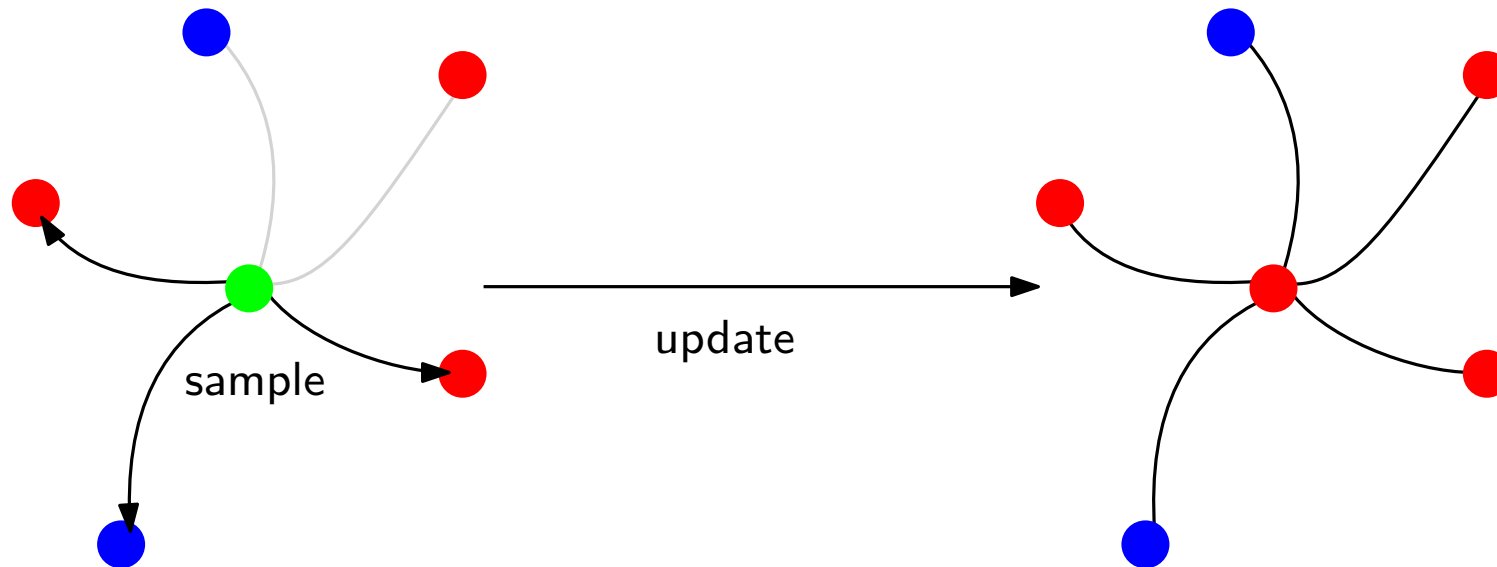


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**Markov chain** with  $k$  absorbing states



# Previous results

Overview of results in **noiseless** settings [Becchetti et al., SIGACT News 2020]

Binary case: (**with high probability**)

- in **time**  $\mathcal{O}(\log n)$  the system reaches **consensus**
- if the initial **bias** is  $\Omega(\sqrt{n \log n})$ , we have **majority consensus**
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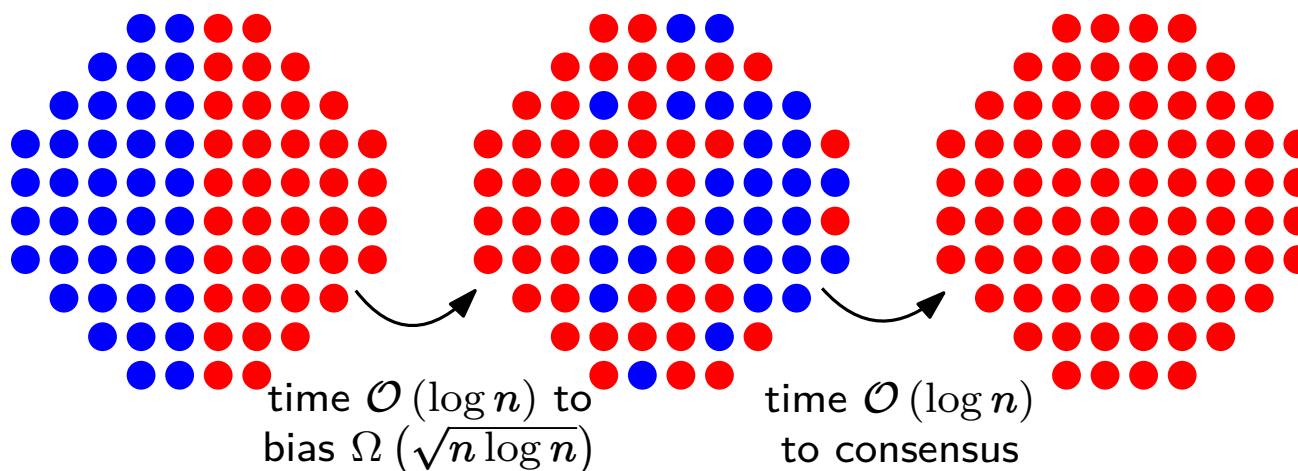
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**Example:**



# Our contribution

[d'Amore et Ziccardi, SIROCCO 2022]

3-MAJORITY dynamics with  $k = 2$  **opinions** in the presence of **uniform noise** in the **complete graph**

**Question:** is there any form of **metastable** consensus?

# Our contribution

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**W.h.p.:** an event  $E$  depending on a parameter  $n$  holds **with high probability** in  $n$  if  $\mathbb{P}(E) \geq 1 - n^{-\gamma}$  for any constant  $\gamma > 0$

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For the 3-MAJORITY dynamics: **phase-transition**

- $p < 1/3$ :
- a value  $\bar{s} = \Theta(n)$  exists such that the **bias** of the system reaches the **interval**  $I_\varepsilon = [(1 - \varepsilon)\bar{s}, (1 + \varepsilon)\bar{s}]$  in **time**  $\mathcal{O}(\log n)$  w.h.p., and **keeps** in  $I_\varepsilon$  for time  $\text{poly}(n)$  w.h.p.  $\longrightarrow$  **almost-consensus**
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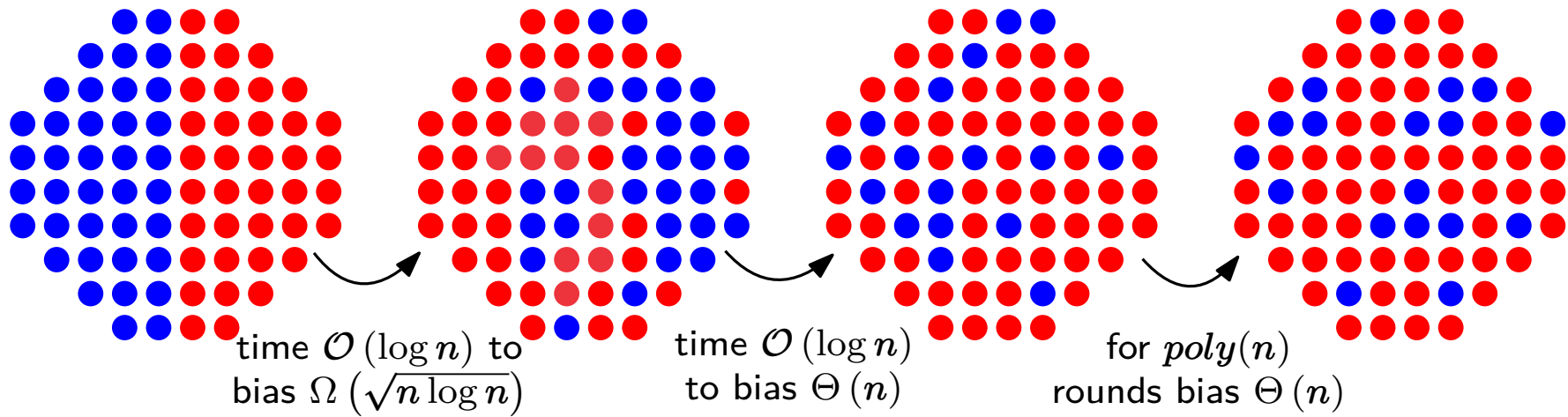
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- $p > 1/3$ :
- in **time**  $\mathcal{O}(\log n)$  the bias becomes **bounded** by  $\mathcal{O}(\sqrt{n \log n})$  and **keeps** bounded for time  $\text{poly}(n)$  wh.p.  $\longrightarrow$  **victory of noise**
  - there is **constant probability** to **switch majority** within **time**  $\mathcal{O}(\log n)$

# Examples

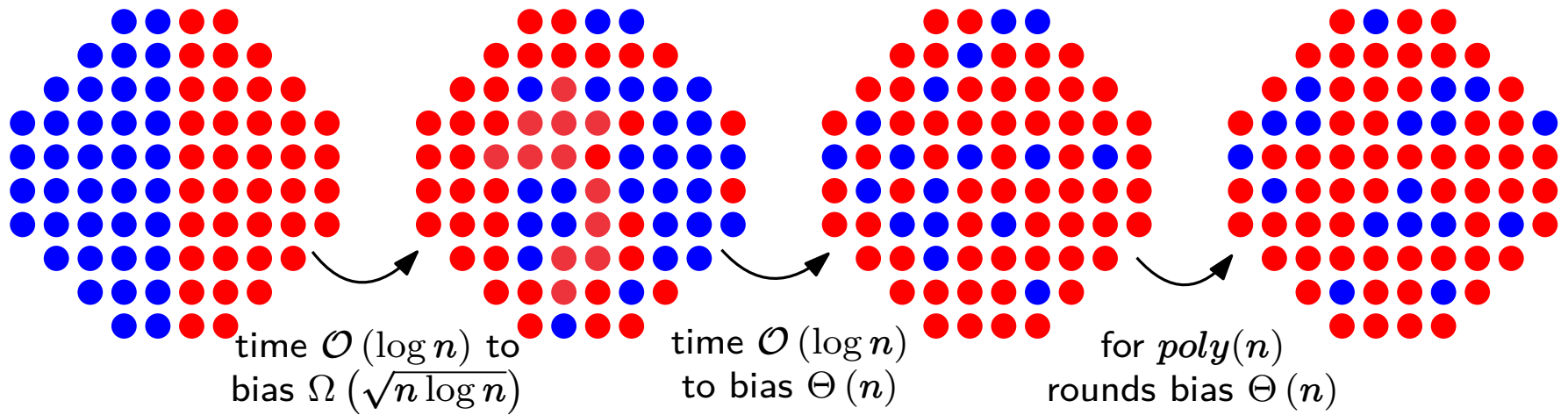
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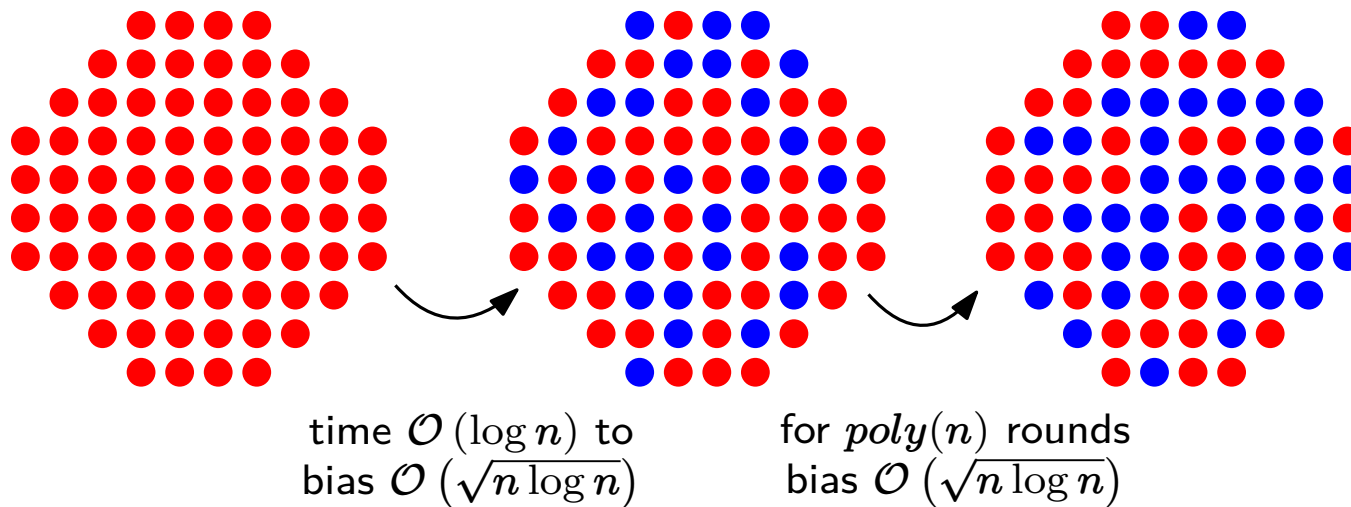


# Examples

## Almost-consensus



## Victory of noise



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$u \backslash v$	opinion $i$	opinion $j$	undecided
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undecided	$i$	$j$	undecided

- phase-transition at  $p = 1/2$ , more resilient to noise
- less characterized
  - no precise equilibrium shown
  - no switch of majority shown

# Techniques

**Equilibrium:**

$$\mathbb{E}[s_{t+1}|s_t] = \frac{s(1-p)}{2} \left( 3 - \frac{s_t^2}{n^2} (1-p)^2 \right) \xRightarrow{\quad \uparrow \quad} \bar{s} = \frac{n}{1-p} \cdot \sqrt{\frac{1-3p}{1-p}}$$
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- conditional on  $s_t$ ,  $s_{t+1}$  sum of Rademacher random variables

**W.h.p.:** an event  $E$  depending on a parameter  $n$  holds with high probability in  $n$  if  $\mathbb{P}(E) \geq 1 - n^{-\gamma}$  for any constant  $\gamma > 0$

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## For majority consensus:

- initial bias  $\Omega(\sqrt{n \log n})$
- $M_t = |s_t - \bar{s}|$ ,  $\bar{s}$  equilibrium value for bias
- Chernoff-Hoeffding inequalities to show  $M_{t+1} \leq (1 - \delta)M_t$  w.h.p. when bias outside  $I_\varepsilon = [(1 - \varepsilon)\bar{s}, (1 + \varepsilon)\bar{s}]$
- chain rule + union bound



# Techniques

For **victory of noise**:

- initial bias  $s = n$
- **super-martingale**  $N_t \sim s_t$  such that  $\mathbb{E}[N_{t+1} | \mathcal{F}_t] \leq (1 - \delta)N_t$
- **concentration arguments** [Lehre and Witt, ISAAC 2014]  $\implies$  **majority disappears** in time  $\mathcal{O}(\log(n))$

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- **concentration arguments** [Lehre and Witt, ISAAC 2014]  $\implies$  **majority disappears** in time  $\mathcal{O}(\log(n))$

For **symmetry breaking**:

- initial bias  $o(\sqrt{n \log n})$
- the bias has enough **standard deviation** to break symmetry (**drift analysis** results)
  - **std** roughly  $\sqrt{n}$
  - $\mathcal{O}(\log n)$  trials suffices

# Discussion

- **3-Majority** dynamics with noise is not implemented by biological systems
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  - despite the bio-inspiration, highly abstract model
  - aiming to capture fundamental phenomena that (**very loosely**) relates to many biological systems
- phase-transition:  $p = 1/3$ 
  - $p = 1/3$  (3-Maj) means 2 out of 3 pulled opinions are non-noisy on average
- Are noise-thresholds independent of  $k$  (num. of opinions)?
- What about a generalized noise with different parameters for the two opinions?
- What about sparser topologies, e.g. expanders?

The End

Thanks

Questions?