Givens QR Decomposition over Relational Databases

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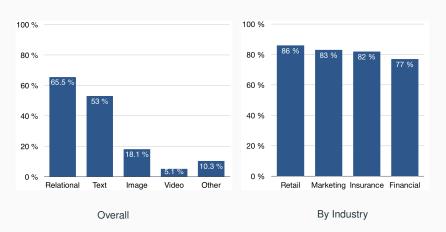




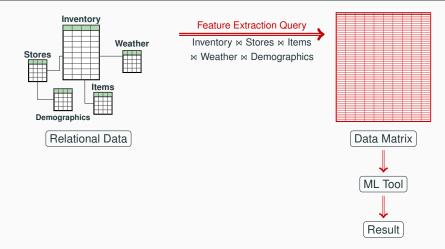
Workshop Factorized Databases, 4.8.2022

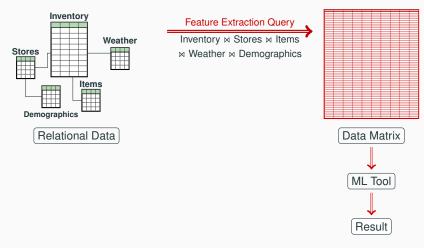
Project Context: Machine Learning for Relational Data

Kaggle Survey: Most Data Scientists use Relational Data at Work!

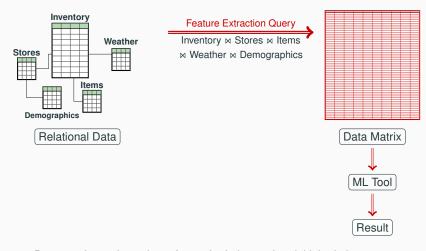


Source: The State of Data Science & Machine Learning 2017, Kaggle, October 2017 (based on 2017 Kaggle survey of 16,000 ML practitioners)

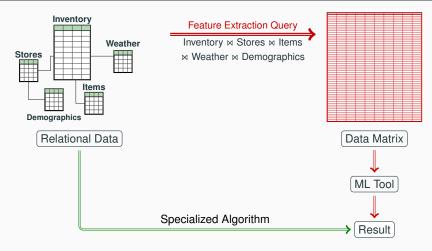




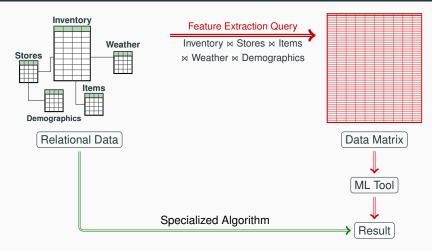
· Data matrix can be orders of magnitude larger than initial relations



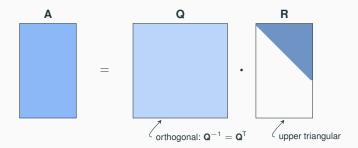
- Data matrix can be orders of magnitude larger than initial relations
- We want to avoid materializing and exporting the data matrix

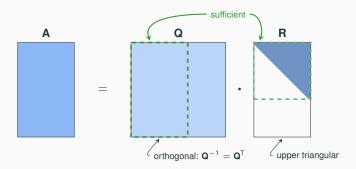


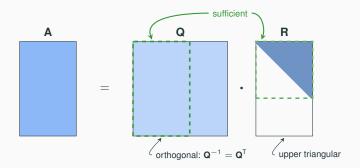
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- Goal: Devise algorithm that runs in linear time wrt. the input database



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- Assumption: Join is acyclic, i.e., admits a join tree



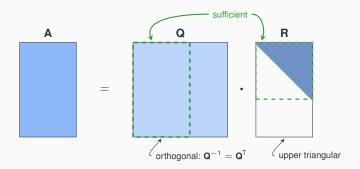




QR decompositions are used for

- computing the least squares estimator $\hat{\mathbf{x}}$ of $\mathbf{A}\mathbf{x} = \mathbf{b}$
- obtaining a Cholesky decomposition of A^TA
- · computing a singular value decomposition of A

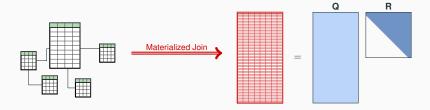
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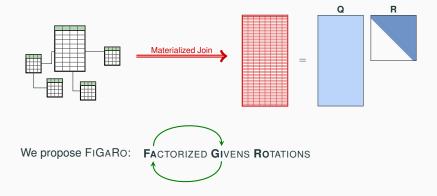


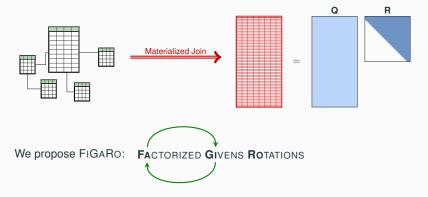
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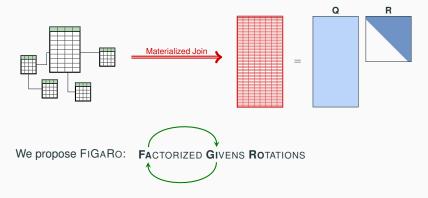
For many applications, only **R** is necessary



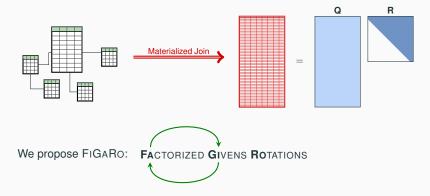




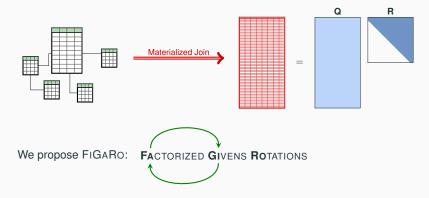
 $+\;$ Pushes the computation of **R** past the joins to the input database



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- + Takes time linear in the number of tuples in the input database
- + Is equivalent to the application of Givens rotations to the materialized join
- $+\,$ If needed, ${f Q}$ can be obtained efficiently from ${f R}$ and the input database

.

 2
 5
 1
 4

 9
 8
 4
 5

 3
 7
 2
 9

 4
 3
 6
 6

 1
 8
 5
 2

 2
 5
 1
 4

 9
 8
 4
 5

 3
 7
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 9

 4
 3
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 1
 8
 5
 2

 2
 5
 1
 4

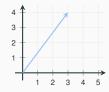
 9
 8
 4
 5

 3
 7
 2
 9

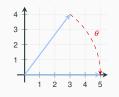
 4
 3
 6
 6

 1
 8
 5
 2

2	5	1	4
9	8	4	5
3	7	2	9
4	3	6	6
1	8	5	2







$$\begin{bmatrix} \cos \theta & \sin \theta \\ -\sin \theta & \cos \theta \end{bmatrix} \begin{bmatrix} 3 \\ 4 \end{bmatrix} = \begin{bmatrix} 5 \\ 0 \end{bmatrix}$$



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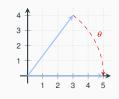




$$\begin{bmatrix} \cos \theta & \sin \theta \\ -\sin \theta & \cos \theta \end{bmatrix} \begin{bmatrix} 3 \\ 4 \end{bmatrix} = \begin{bmatrix} 5 \\ 0 \end{bmatrix}$$



$$\begin{bmatrix} 2 & 5 & 1 & 4 \\ 9 & 8 & 4 & 5 \\ 3 & 7 & 2 & 9 \\ 4 & 3 & 6 & 6 \\ 1 & 8 & 5 & 2 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & \cos\theta & \sin\theta & 0 \\ 0 & 0 & -\sin\theta & \cos\theta & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 2 & 5 & 1 & 4 \\ 9 & 8 & 4 & 5 \\ 3 & 7 & 2 & 9 \\ 4 & 3 & 6 & 6 \\ 1 & 8 & 5 & 2 \end{bmatrix}$$



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$$\begin{array}{c} \\ \\ \\ \end{array}$$

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$$= \begin{bmatrix} 2 & 5 & 1 & 4 \\ 9 & 8 & 4 & 5 \\ 5 & & & & \\ 0 & & & & & \\ \end{bmatrix}$$

Runtime of FIGARO versus Prior Work

Prior work: Givens rotations on the materialized join

N M

- Number of necessary rotations: $\mathcal{O}(\textit{MN})$
- Runtime linear in M, quadratic in N

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FIGARO on the input database with K tuples

- Runtime linear in K, quadratic in N
- If *M* ≫ *K*:
 - FIGARO is faster than the standard algorithms
 - FIGARO is experimentally more accurate than standard implementations
- Even if $M \approx K$, FiGaRo may be faster than the standard algorithms

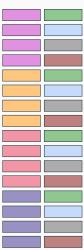
Two relations



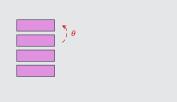
Two relations

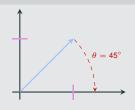


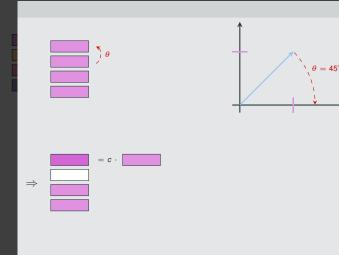
Cartesian product

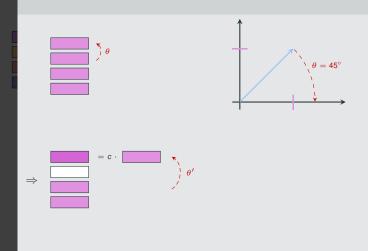


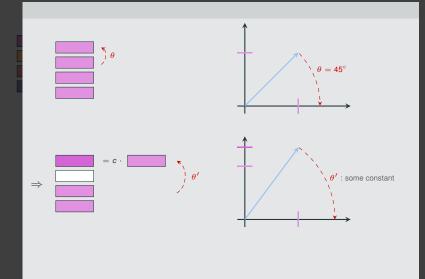


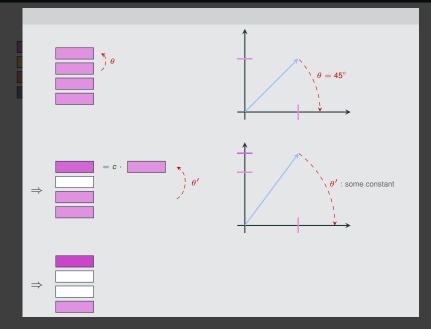


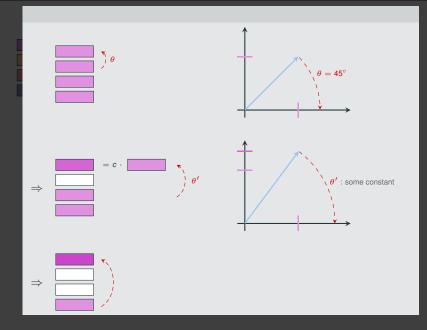


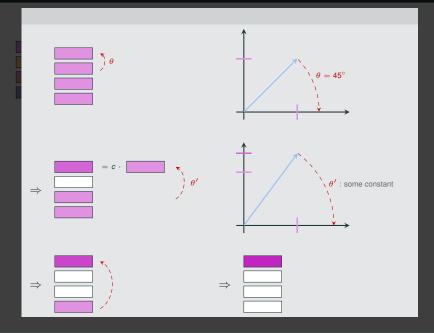








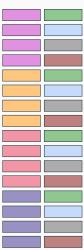


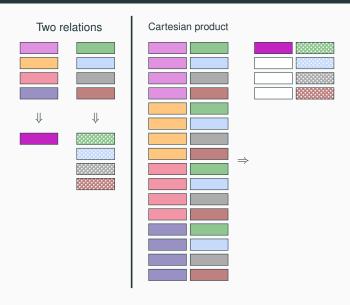


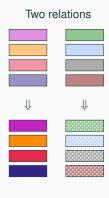
Two relations

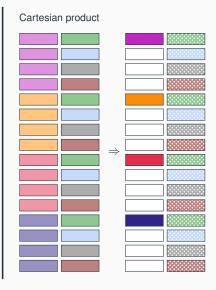


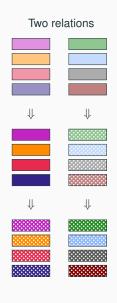
Cartesian product



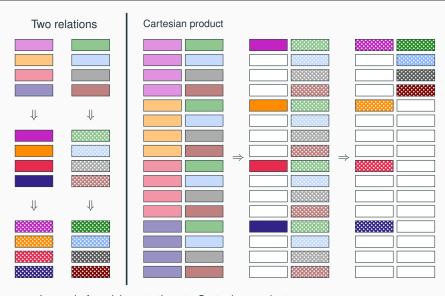












 Instead of applying rotations to Cartesian product: compute non-zero rows from input database in linear time

Combine all rotations in orthogonal matrix M

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For any matrix ${\bf A}$ one can compute ${\bf M}{\bf A}$ in linear time

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For any matrix A one can compute MA in linear time

$$\mathbf{MA} = \begin{bmatrix} \mathcal{H}(\mathbf{A}) \\ \mathcal{T}(\mathbf{A}) \end{bmatrix} \quad \text{Head: first row of } \mathbf{MA}$$

$$\text{Tail: remaining rows of } \mathbf{MA}$$

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QR decomposition for Cartesian Product of matrices **S** and **T** with p resp. q rows:

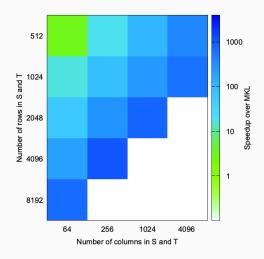
$$\mathbf{S} imes \mathbf{T} = \mathbf{Q} \cdot egin{bmatrix} \sqrt{q} \ \mathcal{H}(\mathbf{S}) & \sqrt{p} \ \mathcal{H}(\mathbf{T}) \\ \sqrt{q} \ \mathcal{T}(\mathbf{S}) & \mathbf{0} \\ \mathbf{0} & \sqrt{p} \ \mathcal{T}(\mathbf{T}) \end{bmatrix}$$

8

Experimental Evaluation

Runtime Performance on Cartesian Products

- Input: Matrices \boldsymbol{S} and \boldsymbol{T} with different numbers of rows and columns
- Intel MKL: on the materialized Cartesian product $\textbf{S} \times \textbf{T}$
- FIGARO: on S and T directly



Accuracy Experiments: Setup

- Usually, accuracy of algorithms for QR decomposition is measured in terms of orthogonality of ${\bf Q}$

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- Usually, accuracy of algorithms for QR decomposition is measured in terms of orthogonality of Q
- · We propose a new accuracy based on R

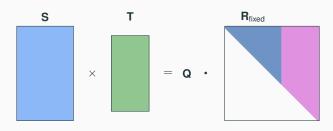
Given:





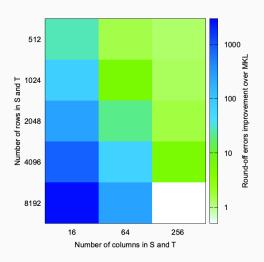
· Dimensions of target matrices S, T

Compute: Matrices S, T such that



Accuracy for QR Decomposition of Cartesian Products

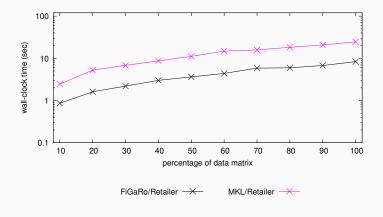
- Input: Matrices \mathbf{S}, \mathbf{T} with varying numbers of rows and columns
- Relative error is measured using Frobenius norm of obtained $\hat{\textbf{R}}_{\text{fixed}}$ vs the ground truth $\textbf{R}_{\text{fixed}}$



Runtime Performance as Function of the Size of the Materialized Join (1/3)

Retailer dataset

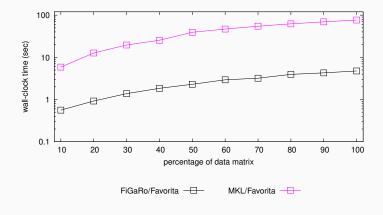
- # tuples in the input database vs materialized join: 86M vs 84M
- # data columns in the materialized join: 43
- # values in the materialized join vs input database (relative): 10x



Runtime Performance as Function of the Size of the Materialized Join (2/3)

Favorita dataset

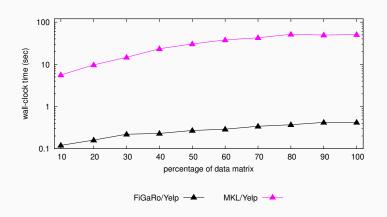
- # tuples in the input database vs materialized join: 125M vs 127M
- # data columns in the materialized join: 30
- # values in the materialized join vs input database (relative): 6x



Runtime Performance as Function of the Size of the Materialized Join (3/3)

Yelp dataset

- # tuples in the input database vs materialized join: 2M vs 150M
- # data columns in the materialized join: 50
- # values in the materialized join vs input database (relative): 332x





Conclusion and Future Work

- FIGARO pushes the computation of the QR decomposition past the join
- It significantly reduces the computation time in experiments on synthetic and real-world relational data
- · It can be numerically more accurate than standard implementations
- The techniques can be transferred to other matrix decompositions, e.g. LU decomposition