

# Masters Theorem Proof

Francisco Diaz

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$$T(n) = aT(n) + O(n^d)$$

$$T(n) = \begin{cases} O(n^{d \log(n)}) , & \text{if } a = b^d \\ O(n^d) , & \text{if } a < b^d \\ O(n^{\log_b(a)}) , & \text{if } a > b^d \end{cases}$$

$$\text{number of problems} = a^j$$

$$\text{work} = a^j c(n/b^j)$$

$$\text{total work} = \sum_{j=0}^{\log_b(n)} a^j \cdot c\left(\frac{n}{b^j}\right)^d = \sum_{j=0}^{\log_2(n)} cn^d \left(\frac{a}{b^d}\right)^j = cn^d \sum_{j=0}^{\log_2(n)} \left(\frac{a}{b^d}\right)^j$$

### 0.1 Case: $a < b^d$

$$\text{Suppose } \frac{a}{b^d} < 1$$

$$\text{then } cn^d \sum_{j=0}^{\log_2(n)} \left(\frac{a}{b^d}\right)^j = cn^d \left( \frac{1}{1 - \frac{a}{b^d}} \right) = cn^d \left( \frac{b^d}{b^d - a} \right) = O(n^d)$$

Given that  $a$  and  $b$  are constant

### 0.2 Case: $a > b^d$

$$\text{Suppose } \frac{a}{b^d} > 1$$

$$\sum_{j=0}^x r^j = \frac{r^{x+1} - 1}{r - 1}$$

$$T(n) = cn^d \left( \frac{\frac{a}{b^d}^{\log_b(n)+1} - 1}{\frac{a}{b^d} - 1} \right) = O\left(n^d \frac{a^{\log_b(n)}}{b^{d \cdot \log(n^d)}}\right)$$

$$O\left(n^d \frac{a^{\log_b(n)}}{b^{\log_b(n^d)}}\right) = O\left(n^d \frac{a^{\log_b(n)}}{n^d}\right) = O\left(a^{\log_b(n)}\right) = O\left(n^{\log_b(a)}\right)$$

### 0.3 Case $a = b^d$

$$\text{Suppose } a = b^d$$

$$\text{then } cn^d \sum_{j=0}^{\log_b(n)} \left(\frac{a}{b^d}\right)^j = cn^d (\log_b(n) + 1) = cn^d \log_b(n) + cn^d$$

$$O\left(n^d \log_b(n)\right)$$