



High Performance Linear Algebra

Lecture 2: Performance Modeling and the Roofline Model

Ph.D. program in High Performance Scientific Computing

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Summary of previous lecture

1 Summary of previous lecture

In the **previous lecture** we have discussed the following topics:

- Introduction to High Performance Scientific Computing
- Overview of Linear Algebra and its importance in scientific computing
- The machines from the TOP500 list
- What tools are we going to use in this course
 - Programming language: Fortran
 - Software Versioning: Git and GitHub



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 - Programming models
 - Programming tools



What is CI/CD?

2 Continuous Integration and Deployment (CI/CD)

Continuous Integration (CI) and **Continuous Deployment (CD)** are practices in software development that automate the process of integrating code changes and deploying applications.

Continuous Integration (CI):

- Developers frequently merge code changes into a shared repository
- Automated builds and tests run to detect issues early

Continuous Deployment (CD):

- Automatically deploys code changes to production after passing tests
- Ensures rapid delivery of new features and bug fixes



Benefits of CI/CD

2 Continuous Integration and Deployment (CI/CD)

Key benefits:

- **Early bug detection:** Automated tests catch issues before they reach production
- **Faster development cycles:** Rapid integration and deployment of changes
- **Improved collaboration:** Teams work together more effectively with shared codebase
- **Higher quality software:** Consistent testing and deployment processes

Popular CI/CD tools:

- GitHub Actions
- GitLab CI/CD
- Jenkins
- Travis CI



An example of CI/CD with GitHub Actions

2 Continuous Integration and Deployment (CI/CD)

Setting up a simple CI workflow: we will use GitHub Actions to automatically build and test our Fortran code whenever we push changes to the repository

First, ensure your Fortran project has a `Makefile` with appropriate build and test targets. This can be as simple as:

```
all:
    gfortran -o hello hello.f90
test:
    ./hello
```

We recall that a **Makefile** is a file that defines a set of tasks to be executed. It is commonly used to automate the build process of software projects.



Makefile basics

2 Continuous Integration and Deployment (CI/CD)

A **Makefile** consists of rules with the following structure:

```
target: dependencies  
    command
```

Example:

```
hello: hello.f90  
    gfortran -o hello hello.f90
```

Here, `hello` is the target, `hello.f90` is the dependency, and the command compiles the Fortran source file into an executable named `hello`.

For **small projects**, a simple Makefile like this is sufficient to automate the build and test process. For **larger projects**, it is better to also have the Makefile programmatically generated using tools like CMake or Autotools.



Creating a workflow on GitHub Actions

2 Continuous Integration and Deployment (CI/CD)

GitHub Actions allows you to automate workflows directly in your GitHub repository.

Key components:

- **Workflows:** Automated processes defined in YAML files
- **Jobs:** A set of steps that execute on the same runner
- **Steps:** Individual tasks within a job (e.g., running commands, setting up environments)

In GitHub, workflows are stored in the `.github/workflows/` directory of your repository as YAML (`.yaml`, Yet Another Markup Language) files.



An example of CI/CD with GitHub Actions

2 Continuous Integration and Deployment (CI/CD)

Create a file `.github/workflows/ci.yml` in your repository:

```
name: CI
on:
  push:
    branches:
      - main
jobs:
  build:
    runs-on: ubuntu-latest
    steps:
      - name: Checkout code
        uses: actions/checkout@v4
      - name: Set up Fortran
```



An example of CI/CD with GitHub Actions

2 Continuous Integration and Deployment (CI/CD)

```
run: |  
    sudo apt-get update  
    sudo apt-get install -y gfortran  
- name: Test  
  run: make test
```

Explanation:

- Triggered on **pushes** to the main branch

```
on:  
  push:  
    branches:  
      - main
```

- Runs on the latest Ubuntu environment



An example of CI/CD with GitHub Actions

2 Continuous Integration and Deployment (CI/CD)

- Defines a job named build with several steps:
 - Checkout code from the repository
 - **name**: Checkout code
 - uses**: actions/checkout@v4
 - Install gfortran
 - **name**: Set up Fortran
 - run**: |
 - sudo apt-get update**
 - sudo apt-get install -y gfortran**
 - Build the project using make
 - **name**: Build
 - run**: make



An example of CI/CD with GitHub Actions

2 Continuous Integration and Deployment (CI/CD)

- Run tests using `make test`
 - `name`: Test
 - `run`: `make test`

You can change the “manual installation” of `gfortran` with a pre-built action from the GitHub Marketplace, such as `setup-fortran`:

- `name`: Setup Fortran
- `uses`: `fortran-lang/setup-fortran@v1.8.0`
- `with`:
 - `compiler`: `gcc`
 - `version`: `'latest'`
 - `update-environment`: `true`



If everything has gone well:

2 Continuous Integration and Deployment (CI/CD)

From the top menu of your GitHub repository, click on the **Actions** tab. You should see your workflow running, and if everything is set up correctly, it should complete successfully, indicating that your Fortran code has been built and tested automatically.

build
succeeded now in 19s

Search logs

| | |
|------------------------|-----|
| > ✓ Set up job | 0s |
| > ✓ Checkout code | 1s |
| > ✓ Set up Fortran | 13s |
| > ✓ Build | 2s |
| > ✓ Test | 0s |
| > ✓ Post Checkout code | 0s |
| > ✓ Complete job | 0s |

Example: github.com/High-Performance-Linear-Algebra/hello-fortran-world



Learning by doing:

2 Continuous Integration and Deployment (CI/CD)

Explore the `setup-fortran` action and modify the provided example workflow to include additional steps, such as:

- Running on different operating systems (e.g., Windows, macOS)
- Running on different Fortran compilers (e.g., Intel Fortran, NVIDIA HPC SDK)

Question: How would you modify the Makefile so that it does not call `gfortran` directly, but uses instead the compiler available in the environment?

Moving to CMake: To automate the search for the compiler, configuration, and building of projects, a good practice is to use CMake.



Further informations on the GitHub Actions

2 Continuous Integration and Deployment (CI/CD)

For more information on GitHub Actions, refer to the official documentation:

- [GitHub Actions Documentation](#)
- [Introduction to GitHub Actions](#)
- [Workflow syntax for GitHub Actions](#)

They can help you explore more advanced features and customize your CI/CD workflows, such as:

- [Running tests on multiple operating systems](#)
- [Integrating with other services](#)
- [Deploying applications automatically](#)
- [Setting up notifications for build status](#)
- [Deploying Documentation, artifacts, and more](#)



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General Parallel Programming Issues

3 General Parallel Programming Issues

In **this lecture** we will discuss some general issues related to parallel programming:

- Performance metrics
- Scalability
- Speedup and Efficiency
- Amdahl's and Gustafson's Laws
- Programming Models
- Parallel Architectures
- Roofline Model



What do we mean by parallelism?

3 General Parallel Programming Issues

Definition

We call *parallelism* the ability to have multiple operations completing their execution at the same time.

- “Operation” may mean a machine instruction, a floating-point operation, or something else depending on context.
- In scientific/engineering applications, the key metric is often the number of floating-point operations (FLOPs), typically the limiting factor for execution speed.



What do we mean by parallelism?

3 General Parallel Programming Issues

Keywords: FLOP and FLOP/s

A *FLOP* is a floating-point operation, typically an addition, subtraction, multiplication, or division between two floating-point numbers. The number of FLOPs required to solve a problem is often used as a measure of the problem's **computational complexity**. The amount of FLOP/s (floating-point operations per second) a computer can perform is a key measure of its performance, and is often used to rank supercomputers (e.g., the TOP500 list).



Levels of parallelism in a computing system

3 General Parallel Programming Issues

- Within a **single machine** instruction specifying multiple operations (e.g., SSE on x86, fused multiply-add on many architectures).
- Within a **single processor** completing more than one instruction per clock cycle; many modern RISC processors have multiple execution units and are *superscalar*.
- Within a **single silicon chip** hosting multiple CPUs; *multicore* processors¹, a core being each complete CPU.
- Within a **single computer** containing multiple processors.
- Using **multiple computers** connected through some sort of communication device.

¹This usage is slightly confusing, since we are calling *processor* both a single CPU and a chip hosting multiple cores; context should avoid confusion.



Programmer's perspective

3 General Parallel Programming Issues

- The first two kinds of parallelism (*single machine* and *single processor*) are mostly handled by the compiler and optimized libraries implementing heavy computational kernels.
- They are almost transparent to application programmers.
- We will concentrate on the **last three kinds of parallel computing systems**:
 - *Single chip* (multicore) systems.
 - *Single computer* with multiple processors.
 - *Multiple computers* connected by a network.



Parallel systems: memory configurations

3 General Parallel Programming Issues

In classifying high-performance parallel computers, the discriminating factor is the memory subsystem configuration. Two main kinds:

1. *Shared memory systems.*
2. *Distributed memory systems.*



Flynn's Taxonomy

3 General Parallel Programming Issues

Overview

Since the early 1970s, the **Flynn taxonomy** [1, 2] has been the standard classification scheme for computer architectures.

Four categories:

- SISD** Single Instruction Single Data
- SIMD** Single Instruction Multiple Data
- MISD** Multiple Instruction Single Data
- MIMD** Multiple Instruction Multiple Data



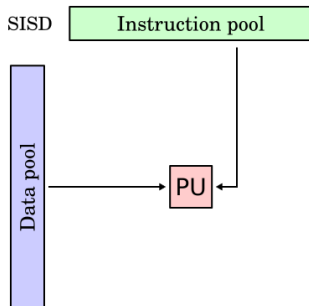
Flynn's Taxonomy: SISD

3 General Parallel Programming Issues

Single Instruction Single Data

Sequential computers where a single stream of data is processed by a single stream of instructions.

- Traditional von Neumann architecture
- One instruction operates on one data element at a time
- No parallelism at the architecture level





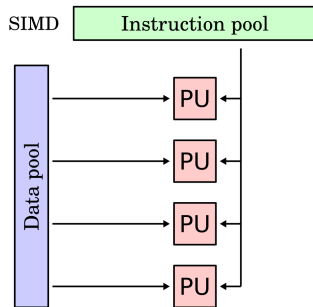
Flynn's Taxonomy: SIMD

3 General Parallel Programming Issues

Single Instruction Multiple Data

Vector processors capable of handling multiple data with a single instruction.

- Data presented in the form of a vector
- **Notable example:** Cray-1 computer with vector registers of length 64
- Modern examples: SSE/AVX instructions on x86, AltiVec on Power processors





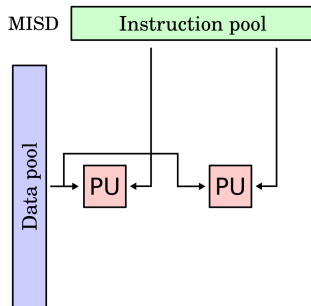
Flynn's Taxonomy: MISD

3 General Parallel Programming Issues

Multiple Instruction Single Data

Multiple instruction streams operating on the same data.

- No significant examples of such architectures have been built
- Flynn classified ancient plug-board machines in this category
- Some embedded devices use this for **fault tolerance**:
 - Same instruction executed redundantly in multiple streams
 - Results verified for accordance





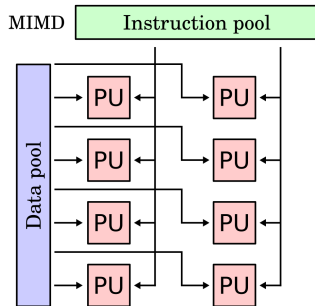
Flynn's Taxonomy: MIMD

3 General Parallel Programming Issues

Multiple Instruction Multiple Data

Multiple instruction streams concurrently operate on different (sub)sets of data.

- **The vast majority** of parallel computers built in the last 30 years
- Includes both:
 - Shared memory multiprocessors
 - Distributed memory multiprocessors
- Most flexible and powerful category





Modern Supercomputers: Beyond Flynn

3 General Parallel Programming Issues

- Flynn's taxonomy still useful for **rough categorization**
- Modern landscape is **much more complex**
- Current supercomputers are **hybrid combinations** of different architectures



The Marenstrum 5 supercomputer at the Barcelona Supercomputing Center (BSC).



Anatomy of a Modern Supercomputer

3 General Parallel Programming Issues

A typical modern supercomputer consists of:

1. **Multiple nodes** connected through a network interconnect
2. **Multiple processors per node**, possibly in NUMA configuration
 - NUMA = Non-Uniform Memory Access
3. **Multicore processors** (multiple cores per processor)
4. **SIMD units per core**
 - SSE/AVX on x86
 - AltiVec on Power processors
 - Vector instructions of size 2 or 4
5. **Accelerators** (e.g., GPUs)
 - Own memory system
 - Connected via PCI or high-speed interconnect (e.g., NVLink)



Challenges in Modern HPC

3 General Parallel Programming Issues

The Challenge

Exploiting the computational power of modern supercomputers demands:

- Programming models capable of matching their **hierarchical structure**
- Algorithms that can handle their **heterogeneity**
- Understanding of performance metrics for parallel programs

Coming up: Basic concepts of parallel programming and performance evaluation



Performance metrics in parallel computing

3 General Parallel Programming Issues

- Many alternatives exist: hardware architectures, programming paradigms, applications.
- It is necessary to **define metrics** to evaluate the performance of a parallel system.



Performance metrics in parallel computing

3 General Parallel Programming Issues

- Many alternatives exist: hardware architectures, programming paradigms, applications.
- It is necessary to **define metrics** to evaluate the performance of a parallel system.

No single criterion fits all

- **Computer scientist:** pure algorithmic speed-up.
- **Computational scientist:** time to completion; maximum problem size that can be analyzed.
- **System administrator:** maximize system utilization.



Beware of benchmarks

3 General Parallel Programming Issues

- No substitute for testing with the workload of interest.
- Benchmarks are indicators only as good as their relation to the intended usage.
- Procurement differs: single critical application vs computing center serving many users.



What do we mean by scalability?

3 General Parallel Programming Issues

- A parallel system = implementation of a parallel algorithm on a given parallel architecture.
- Scalability theory organizes performance evaluation while accounting for usage aspects.
- We must choose appropriate metrics depending on goals and constraints.



Questions a metric should answer

3 General Parallel Programming Issues

- How do we measure the raw performance of a system?
- How do we compare measurements obtained on different machines?
- How does the metric respond to the programming paradigm employed?
- Do we want raw performance or value for money?

Note

Given the variety of questions, only general criteria exist; there is no single, precise measurement procedure.



Definitions of scalability

3 General Parallel Programming Issues

- A parallel system is scalable if it delivers the same performance per processor while increasing the number of processors and/or the problem size.
- A program is scalable if its performance improves when increasing the processors employed from $p - 1$ to p .



Workload, execution times, and I/O

3 General Parallel Programming Issues

- Define problem size W as the number of basic operations for the **best-known sequential algorithm**.

²See e.g. ROOT from CERN: root.cern and CAPIO: github.com/High-Performance-IO/capio.



Workload, execution times, and I/O

3 General Parallel Programming Issues

- Define problem size W as the number of basic operations for the **best-known sequential algorithm**.

How do we know what is the *best* sequential algorithm?

There exist a branch of computer science, called *computational complexity theory*, that studies the **inherent difficulty** of computational problems and classifies them according to the resources needed to solve them. However, in practice, we often rely on empirical performance measurements and established benchmarks to determine the best-known algorithms for specific problems.

²See e.g. ROOT from CERN: root.cern and CAPIO: github.com/High-Performance-IO/capio.



Workload, execution times, and I/O

3 General Parallel Programming Issues

- Define problem size W as the number of basic operations for the **best-known sequential algorithm**.
- **Serial time** T_s : start-to-end time on one processor.
- **Parallel time** T_p on p processors: time until the last processor completes.
- **I/O** may be included or measured separately depending on goals; if I/O is essential, consider parallel I/O.²

²See e.g. ROOT from CERN: root.cern and CAPIO: github.com/High-Performance-IO/capio.



Execution-time models and benchmarking

3 General Parallel Programming Issues

- $T_s = f(W)$
- $T_p = f(W, p, \text{arch})$
- Absolute assessment is application-dependent; for example, the TOP500 rules measure **dense linear algebra factorization performance** (remind the TOP500 list).



Speed-up

3 General Parallel Programming Issues

Definition

$$S(W, p) = T_s(W) / T_p(W, p)$$

- Linear: $S = p$
- Sub-linear: $S < p$
- Super-linear: $S > p$

Linear speed-up is the *ideal target*; in practice communication and overheads limit it.



Why speed-up is often sub-linear

3 General Parallel Programming Issues

- Startup costs
- Communication latency/bandwidth limits
- Synchronization overheads

Example: consider a workload with W operations and communication overhead C per processor. The parallel time can be modeled as:

$$T_p(W, p) = \frac{W}{p} + C \cdot p$$

The speed-up is then:

$$S(W, p) = \frac{W}{\frac{W}{p} + C \cdot p} = \frac{p}{1 + \frac{C \cdot p^2}{W}}$$

With fixed W , the fraction of time spent communicating typically grows with p .



When super-linear speed-up may occur

3 General Parallel Programming Issues

1. Better use of **memory hierarchy**: partitioning reduces working set per process, improving cache and memory behavior (coming back to this in a few slides).
2. **Search problems**: parallel decomposition changes exploration order, finding solutions earlier; sometimes even $S > 1$ on a single core if extra memory helps.

Memory hierarchy: In many computing systems, memory is organized in a hierarchy of levels, each with different speeds and sizes. The fastest level is the CPU cache, followed by main memory (RAM), and then slower storage devices like hard drives or SSDs. When a problem is divided among multiple processors, each processor may work on a smaller subset of the data, which can fit better into the faster levels of the memory hierarchy. This improved data locality can lead to reduced memory access times and increased overall performance, resulting in super-linear speed-up in some cases.



When super-linear speed-up may occur

3 General Parallel Programming Issues

1. Better use of **memory hierarchy**: partitioning reduces working set per process, improving cache and memory behavior (coming back to this in a few slides).
2. **Search problems**: parallel decomposition changes exploration order, finding solutions earlier; sometimes even $S > 1$ on a single core if extra memory helps.

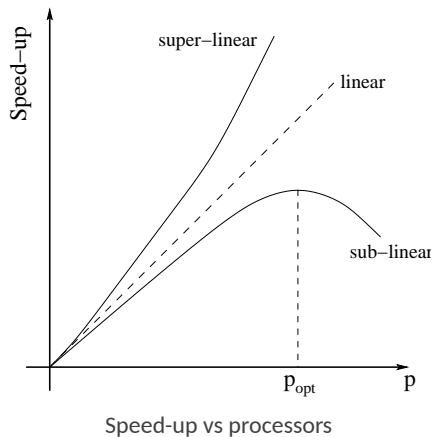
Note: Super-linear speed-up is rare and often specific to certain problem structures or algorithmic techniques. For the linear algebra problems we will consider, it is highly uncommon.



Speed-up vs number of processors

3 General Parallel Programming Issues

- Linear, sub-linear, and occasional super-linear trends.
- For fixed W , the gap from linear typically widens as p increases due to rising overheads.

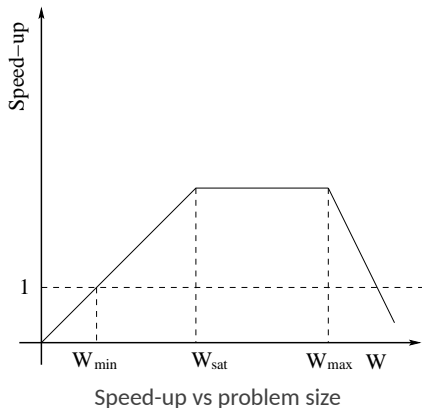




Speed-up vs problem size

3 General Parallel Programming Issues

- Benefit only within a range $[W_{\min}, W_{\max}]$.
- Small W : communication dominates ($S < 1$).
- Peak near W_{sat} ; then saturation.
- For very large W , node memory pressure can degrade performance sharply.





Efficiency

3 General Parallel Programming Issues

Definition

$$E(W, p) = \frac{S(W, p)}{p},$$

or, substituting the definition of speed-up

$$E(W, p) = \frac{S(W, p)}{p} = \frac{T_s(W)}{T_p(W, p) \cdot p}$$

- Typically E in $(1/p, 1]$, excluding rare super-linear cases.



Two models of scalability

3 General Parallel Programming Issues

We can think of two complementary models of scalability:

- **Amdahl's law:** fixed workload W ; increasing p .
- **Gustafson's law:** scale W with p to keep execution time constant.

For the first, the focus is on how speed-up degrades with more processors for a fixed problem size. For the second, the focus is on how larger problems can be solved in the same time as more processors are added.

Strong-scaling vs weak-scaling

- **Strong-scaling:** how T_p changes with p at fixed W (Amdahl).
- **Weak-scaling:** how T_p changes with p when W scales with p (Gustafson).



Amdahl's model

3 General Parallel Programming Issues

- Serial fraction f_s : time inherently serial divided by $T_s(W)$
- Parallel fraction $f_p = 1 - f_s$

Parallel time

$$T_p(W, p) = T_s(W)f_s + T_s(W)f_p/p$$



Amdahl's speed-up and limit

3 General Parallel Programming Issues

Speed-up

$$S(W, p) = \frac{s(W)}{p(W, p)} = \frac{p}{1 + (p - 1)f_s}$$

Asymptotic limit

$$\lim_{p \rightarrow \infty} S(W, p) = 1/f_s$$

- Example: $f_s = 5\% \Rightarrow S \leq 20$ regardless of p .



Implications of Amdahl's law

3 General Parallel Programming Issues

- Reducing the *serial fraction* f_s is critical for absolute performance.
- At fixed W , added sync/comm overhead may cap speed-up even if all code paths are parallelized.
- For many applications in large-scale linear algebra, we scale W with p , making Amdahl less constraining in practice.



Gustafson's law: scaled-speed perspective

3 General Parallel Programming Issues

- With serial fraction α and parallel fraction $(1 - \alpha)$, scale the parallel work with n processors:

$$W(n) = \alpha W + (1 - \alpha) nW, \quad S'_n = \frac{W(n)}{W} = \alpha + (1 - \alpha) n$$

- Better reflects solving larger problems with more processors.
- Caveat: α may change as the problem scales.

Putting it together

- Workload = serial part + parallel part + parallelization/communication overhead.
- Realistic speed-up estimates must include all three and are application dependent.



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Throughput vs parallel applications

4 Paradigms, models and tools for parallel programming

- Parallel machines may run many independent serial jobs to maximize throughput.
- Related to work-pool ideas but outside our main focus on single parallel applications.



Paradigms, models, and tools

4 Paradigms, models and tools for parallel programming

Paradigm Logical structure imposed on a parallel algorithm.

Model How parallelism is expressed in code.

Tool Software instruments (compiler, libraries, etc.).

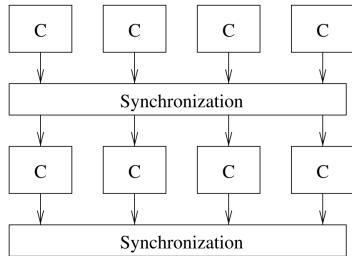
The right choice depends on both software and hardware; shifting abstraction levels is essential.



Algorithmic paradigms (1/2)

4 Paradigms, models and tools for parallel programming

Phase parallel Alternate independent compute phases and synchronization.



Phase parallel

Phase parallel

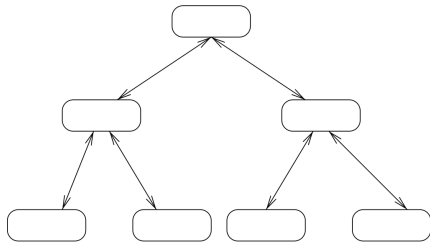


Algorithmic paradigms (1/2)

4 Paradigms, models and tools for parallel programming

Phase parallel Alternate independent compute phases and synchronization.

Divide and conquer Recursively split into subproblems; combine results.



Divide and Conquer

Divide and conquer



Algorithmic paradigms (1/2)

4 Paradigms, models and tools for parallel programming

Phase parallel Alternate independent compute phases and synchronization.

Divide and conquer Recursively split into subproblems; combine results.

Owner computes Partition data; each task processes its own partition.

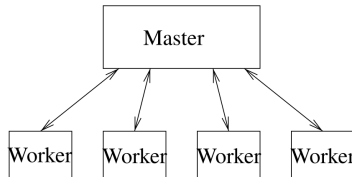


Algorithmic paradigms (2/2)

4 Paradigms, models and tools for parallel programming

Master-worker A controller distributes work and collects results.

PDE-based simulations often use phase-parallel + owner-computes via domain decomposition. Design optimization may mix master-worker with complex simulation kernels.



Master-Worker

Master-worker



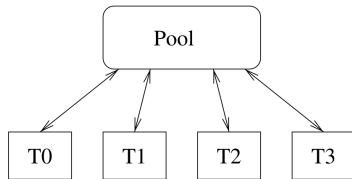
Algorithmic paradigms (2/2)

4 Paradigms, models and tools for parallel programming

Master-worker A controller distributes work and collects results.

Work pool Tasks pull jobs from a shared queue; may generate new jobs.

PDE-based simulations often use phase-parallel + owner-computes via domain decomposition. Design optimization may mix master-worker with complex simulation kernels.



Work pool

Work pool



Programming models

4 Paradigms, models and tools for parallel programming

Implicit parallelism Compiler extracts parallelism; often low efficiency.

Data parallel Single control flow applied to partitioned data; logically shared memory (e.g., HPF constructs like FORALL).

Message passing Processes interact by explicit messages; SPMD common; natural for owner-computes; very flexible but programmer manages communication.



Programming models

4 Paradigms, models and tools for parallel programming

Shared variable Logically shared memory with multiple control flows and private data; attractive with compiler support via directives.

- Data-parallel HPF largely faded in practice.
- Message passing dominates for scalability to many nodes.
- Shared-memory models are effective, especially with compiler support.



Programming tools for parallel computing

4 Paradigms, models and tools for parallel programming

Overview

This set of tools represents the backbone and provides the actual implementations of the parallel programming models we will be using in this course to make Linear Algebra algorithms run on parallel computers.

Four main tools:

MPI Message Passing Interface

OpenMP Open Multi-Processing

OpenACC Open Accelerators

CUDA Compute Unified Device Architecture



MPI: Message Passing Interface

4 Paradigms, models and tools for parallel programming

What is MPI?

A standardized and portable message-passing system for process communication in parallel computing environments.

Key features:

- Widely used in HPC applications
- Point-to-point and collective communication
- Synchronization primitives
- Data distribution mechanisms
- Natural fit for distributed memory systems



MPI implementations

4 Paradigms, models and tools for parallel programming

- Multiple implementations exist, both open-source and commercial.
- Popular open-source implementations:
 - MPICH: www.mpich.org
 - Open MPI: www.open-mpi.org
- Commercial implementations often optimized for specific hardware, e.g., Cray MPI, Intel MPI, IBM Spectrum MPI.

Note

MPI is a specification; different implementations may have varying performance characteristics and features.

Same code different MPI implementations

The same MPI code can be compiled and run with different MPI implementations, allowing flexibility and portability across various HPC systems (at least on paper).



OpenMP: Shared Memory Parallelism

4 Paradigms, models and tools for parallel programming

What is OpenMP?

An API for multi-platform shared memory multiprocessing programming.

Key features:

- Compiler directives for parallel regions
- Library routines and environment variables
- Primary use: parallelizing loops and code sections
- Concurrent execution on multiple threads
- Ideal for intra-node parallelism

Info: www.openmp.org



OpenACC: Directive-Based Accelerator Programming

4 Paradigms, models and tools for parallel programming

What is OpenACC?

A directive-based programming tool for heterogeneous systems (CPUs and GPUs).

Key features:

- High-level directives for parallelism
- Automatic data movement between host and device
- Easier offloading to accelerators
- Portable across different accelerator architectures

Info: www.openacc.org



CUDA: NVIDIA GPU Programming

4 Paradigms, models and tools for parallel programming

What is CUDA?

A parallel computing platform and API developed by NVIDIA for general-purpose computing on GPUs.

Key features:

- Direct access to GPU parallel processing power
- Low-level control over GPU resources
- Extensive libraries and tools ecosystem
- Specific to NVIDIA GPUs

Info: developer.nvidia.com/cuda-zone



The MPI+X framework

4 Paradigms, models and tools for parallel programming

Combining tools

These programming tools are **not mutually exclusive** and can be used together in a single application.

Common pattern: MPI+X

- **MPI** for inter-node communication
- **X** for intra-node parallelism, where X can be:
 - OpenMP (CPU threads)
 - OpenACC (directives for accelerators)
 - CUDA (direct GPU programming)



MPI+X: The current standard

4 Paradigms, models and tools for parallel programming

- Virtually all large-scale HPC libraries and applications are based on the MPI+X framework
- Active research into alternatives to MPI+X exists
- In this course, we focus on MPI+X due to its widespread adoption

Research opportunity

Porting the ideas and algorithms we discuss to alternative frameworks could be an interesting avenue of research.



Summary of Lecture 2

5 Conclusion and summary

Scalability concepts: speed-up, efficiency, Amdahl's and Gustafson's laws

Parallel programming paradigms: phase parallel, divide and conquer, owner computes, master-worker, work pool

Programming models: implicit parallelism, data parallel, message passing, shared variable

Programming tools: MPI, OpenMP, OpenACC, CUDA: the MPI+X framework

Next up: Modern memory hierarchy and Roofline model: for performance analysis

Measuring memory bandwidth with STREAM benchmark, Intra-node parallelism with

OpenMP and starting the implementation of basic linear algebra kernels.



References

6 Bibliography

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- [2] M. J. Flynn. “Some Computer Organizations and Their Effectiveness”. In: *IEEE Transactions on Computers* C-21.9 (1972), pp. 948–960. DOI: [10.1109/TC.1972.5009071](https://doi.org/10.1109/TC.1972.5009071).