

Solving Linear Systems of the Form $(A + \gamma UU^T)x = b$

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Iterative Methods for Large-Scale Saddle-Point Problems
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Palazzo della Carovana (1567), seat of the Scuola Normale



Basic facts about the Scuola Normale

- Established by Napoleon in 1813 as a branch of the ENS in Paris
- Public institution (independent from U. of Pisa)
- Three Faculties: Sciences, Letters and Philosophy, Political and Social Sciences (approximately 65 permanent faculty members)
- Most selective educational institution in Italy (5% admission rate)
- Student body is nearly 600 students (half of them PhD students), all on full scholarships
- Some notable alumni: Giosuè Carducci (Nobel prize 1906), Enrico Fermi (Nobel prize 1938), Carlo Rubbia (Nobel prize 1984), Alessio Figalli (Fields medal 2018)
- Two of Italy's presidents and several prime ministers were SNS alumni

Our conference room, the "Sala Azzurra"



Outline

- 1 Motivation
- 2 The proposed preconditioner
- 3 Eigenvalue bounds
- 4 Numerical experiments
- 5 Conclusions and future work

Joint work with Chiara Faccio (SNS).

Please note that this is **work in progress!**

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The problem

We are interested in finding efficient solvers for large systems of the form

$$(A + \gamma UU^T)x = b, \quad (1)$$

where $A \in \mathbb{R}^{n \times n}$, $U \in \mathbb{R}^{n \times k}$, $\gamma > 0$ and $b \in \mathbb{R}^n$. Here $1 \ll k \ll n$.

We target problems with the following characteristics:

- A is possibly singular, but $A + \gamma UU^T$ is nonsingular for $\gamma > 0$.
- A has one or more desirable property (sparsity, structure, etc.) which is lost if we form $A + \gamma UU^T$ explicitly.
- Mat-vec products with $A + \gamma UU^T$ can be computed efficiently.
- k may not be “small”, but $k \times k$ systems can be solved accurately.
- Problem (1) must be solved repeatedly within a given application.
Often, either A or U remains constant.

The problem (cont.)

Problems of the form (1) which such characteristics arise frequently in scientific computing.

Examples include:

- Augmented Lagrangian methods for saddle point problems;
- Solution of KKT systems in constrained optimization;
- Solution of sparse-dense least squares problems;
- Certain types of integro-differential equations;
- Solution of PDEs describing almost incompressible materials;
- Numerical solution of PDEs with nonlocal BC's;
- ...

Example 1: Augmented Lagrangian methods

Consider the saddle point problem

$$\mathcal{A}\mathbf{x} = \begin{bmatrix} A & B^T \\ B & 0 \end{bmatrix} \begin{bmatrix} u \\ p \end{bmatrix} = \begin{bmatrix} f \\ g \end{bmatrix} = \mathbf{f}.$$

Such systems arise frequently from the [finite element discretization](#) of systems of PDEs, such as for example the Stokes equations, the Oseen problem (obtained from the steady Navier-Stokes equations via Picard linearization), or the coupled Stokes-Darcy system.

A powerful approach to solve such systems is the one based on the [augmented Lagrangian](#).

M. Benzi and M. Olshanskii, *An augmented Lagrangian-based approach to the Oseen problem*, SIAM J. Sci. Comput., 28 (2006), pp. 2095–2113.

P. E. Farrell, L. Mitchell, and F. Wechsung, *An augmented Lagrangian preconditioner for the 3D stationary incompressible Navier-Stokes equations at high Reynolds numbers*, SIAM J. Sci. Comput., 41 (2019), pp. A3075–A3096.

Example 1: Augmented Lagrangian methods (cont.)

The idea is to replace the original saddle point problem with an equivalent one of the form:

$$\mathcal{A}_\gamma \mathbf{x} = \begin{bmatrix} A + \gamma B^T W^{-1} B & B^T \\ B & 0 \end{bmatrix} \begin{bmatrix} u \\ p \end{bmatrix} = \begin{bmatrix} \hat{f} \\ g \end{bmatrix} = \hat{\mathbf{f}},$$

where $\hat{f} = f + \gamma B^T W^{-1} g$. Here W is usually diagonal and positive definite. In the finite element setting, W is often the diagonal of the (pressure) mass matrix.

This new, augmented system is then solved by a Krylov subspace method like (F)GMRES with preconditioner

$$\mathcal{P}_\gamma = \begin{bmatrix} A + \gamma B^T W^{-1} B & B^T \\ 0 & -\gamma^{-1} W \end{bmatrix}.$$

In practice, the preconditioner is applied **inexactly**.

Example 1: Augmented Lagrangian methods (cont.)

The convergence of the preconditioned iteration is usually **very fast** and independent of parameters like the mesh size and viscosity, especially in the “large γ ” limit.

However, at each iteration of the Krylov subspace method a linear system with coefficient matrix $A + \gamma B^T W^{-1} B$ must be solved (inexactly).

This linear system is of the form (1) with $U = B^T W^{-1/2}$. Here A is **sparse**, often **block diagonal**, and **positive definite** (or $A + A^T$ is).

Forming $A + \gamma B^T W^{-1} B$ explicitly would lead to **loss of sparsity and structure**. This system can be quite ill-conditioned (esp. for large γ) and its solution is the **main challenge** associated with the augmented Lagrangian approach.

It is therefore necessary to develop efficient iterative methods for it.

Ideally, we would like such solvers to be robust with respect to $\gamma > 0$.

Example 2: KKT systems in constrained optimization

The solution of (smooth) constrained minimization problems by interior point (IP) methods leads to sequences of linear systems of the form

$$\mathcal{A}\mathbf{x} = \begin{bmatrix} H & -C^T & 0 \\ C & 0 & -I \\ 0 & Z & \Lambda \end{bmatrix} \begin{bmatrix} \delta x \\ \delta \lambda \\ \delta z \end{bmatrix} = \begin{bmatrix} -r_1 \\ -r_2 \\ -r_3 \end{bmatrix} = \mathbf{f}.$$

Here $H = H^T$ is the Hessian of the objective function at the current point x_k , C is the Jacobian of the constraints at the same point, and Z and Λ are diagonal, positive definite matrices associated with the current values of the Lagrange multipliers λ_k and slack variables z_k , respectively.

The variable δz can easily be obtained using the last equation:

$$\delta z = -\Lambda^{-1}(r_3 + Z\delta\lambda)$$

and substituted into the second (block) equation.

Example 2: KKT systems in constrained optimization (cont.)

This yields the reduced system

$$\begin{bmatrix} H & -C^T \\ C & \Lambda^{-1}Z \end{bmatrix} \begin{bmatrix} \delta x \\ \delta \lambda \end{bmatrix} = \begin{bmatrix} -r_1 \\ -r_2 - \Lambda^{-1}r_3 \end{bmatrix}.$$

Eliminating $\delta \lambda$ leads to the **fully reduced** (Schur complement) system

$$(H + C^T Z^{-1} \Lambda C) \delta x = -r_1 - C^T Z^{-1} (r_3 + \Lambda r_2) =: b.$$

After solving for δx , the other unknowns $\delta \lambda$ and δz are readily obtained.

This system is of the form (1) with $A = H$, $U = C^T(Z^{-1}\Lambda)^{1/2}$ and $\gamma = 1$.

The Hessian is usually **positive semidefinite, sparse and possibly structured**. Again, forming $H + C^T Z^{-1} \Lambda C$ explicitly is generally undesirable. Instead, we propose to solve the fully reduced system iteratively, using a suitable (algebraic) preconditioner.

Example 3: sparse-dense LS problems

Consider a large LS problem of the form

$$\|Bx - c\|_2 = \min,$$

where $B \in \mathbb{R}^{m \times n}$ and $c \in \mathbb{R}^m$. Assume that B has the following structure:

$$B = \begin{bmatrix} B_1 \\ B_2 \end{bmatrix}, \quad B_1 \in \mathbb{R}^{(m-k) \times n}, \quad B_2 \in \mathbb{R}^{k \times n},$$

where B_1 is **sparse** and B_2 is **dense**. Then the LS problem is equivalent to the $n \times n$ system of normal equations:

$$(B_1^T B_1 + B_2^T B_2)x = B^T c,$$

which is of the form (1) with $A = B_1^T B_1$, $U = B_2^T B_2$, $\gamma = 1$ and $b = B^T c$.

Once again, we would like to solve this system by an iterative method. The main challenge is again constructing an effective preconditioner.

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The preconditioner

Consider again the linear system (1): $(A + \gamma UU^T)x = b$.

As we have seen, in applications the matrix A (or $A + A^T$) is usually at least positive semidefinite, and we will make this assumption.

Also, although this is not strictly necessary, we will assume that

$$\text{Ker}(A) \cap \text{Ker}(U^T) = \{0\},$$

so that $A + \gamma UU^T$ is nonsingular (and positive definite) for all $\gamma > 0$.

When A is **nonsingular**, we could use the Sherman-Morrison-Woodbury (SMW) formula to solve (1), but this is only applicable to problems of **moderate size**. Recall that SMW states that

$$(A + \gamma UU^T)^{-1} = A^{-1} - \gamma A^{-1}U(I_k + \gamma U^T A^{-1}U)^{-1}U^T A^{-1}.$$

Another possibility would be to build preconditioners based on the SMW formula, where the action of A^{-1} is replaced by some approximation, but our attempts were unsuccessful. Also, often A is **singular**.

The preconditioner (cont.)

When k is small (say, $k = 10$ or less) then any good preconditioner for A (or $A + \alpha I_n$, $\alpha > 0$, if A is singular) tends to give good results. In fact, using CG preconditioned with A^{-1} yields convergence in at most $k + 1$ steps. However, if k is in the hundreds (or larger), this approach is not appealing.

Hence, we need to take into account both A and γUU^T when building the preconditioner. We do this by forming a suitable **product preconditioner**, as follows.

Let $\alpha > 0$ and consider the two splittings

$$A + \gamma UU^T = (A + \alpha I_n) - (\alpha I_n - \gamma UU^T)$$

and

$$A + \gamma UU^T = (\alpha I_n + \gamma UU^T) - (\alpha I_n - A).$$

Note that both $A + \alpha I_n$ and $\alpha I_n + \gamma UU^T$ are invertible.

The preconditioner (cont.)

Let $x^{(0)} \in \mathbb{R}^n$ and consider the alternating iteration

$$(A + \alpha I_n)x^{(k+1/2)} = (\alpha I_n - \gamma UU^T)x^{(k)} + b,$$

$$(\alpha I_n + \gamma UU^T)x^{(k+1)} = (\alpha I_n - A)x^{(k+1/2)} + b,$$

with $k = 0, 1, \dots$. This alternating scheme is analogous to that of other well-known iterative methods like ADI, HSS, etc.

Theorem 1: Assume $A + A^T$ is positive definite. Then the sequence $\{x^{(k)}\}$ converges, as $k \rightarrow \infty$, to the unique solution of equation (1), for any choice of $x^{(0)}$ and for all $\alpha > 0$.

To turn this into a practical method, we will use it as a preconditioner for a Krylov-type method rather than as a stationary iterative scheme. This will also allow inexact solves.

The preconditioner (cont.)

To derive the preconditioner we eliminate $x^{(k+1/2)}$ and write the iterative scheme as the fixed-point iteration

$$x^{(k+1)} = T_\alpha x^{(k)} + c = (I_n - P_\alpha^{-1} A_\gamma) x^{(k)} + P_\alpha^{-1} b,$$

where we have set $A_\gamma = A + \gamma UU^T$. An easy calculation reveals that the preconditioner P_α is given, in factored form, by

$$P_\alpha = \frac{1}{2\alpha} (A + \alpha I_n)(\alpha I_n + \gamma UU^T).$$

The scalar factor $\frac{1}{2\alpha}$ is immaterial for preconditioning, and can be ignored.

Applying this preconditioner requires **two solves** involving $A + \alpha I_n$ and $\alpha I_n + \gamma UU^T$ at each Krylov iteration.

Generally speaking, each of these should be considerably **simpler** than solving systems involving the matrix $A_\gamma = A + \gamma UU^T$.

The preconditioner (cont.)

Consider first solves involving $A + \alpha I_n$. If A is sparse, and/or structured (e.g., block diagonal, Toeplitz, etc.) then so is $A + \alpha I_n$.

Exact solves with $A + \alpha I_n$ can be replaced, if necessary, with inexact solves using either a good preconditioner for $A + \alpha I_n$ or a few steps of an inner iteration (PCG, AMG, or other).

Note the usual trade-off: larger values of α make solves with $A + \alpha I_n$ easier, but may degrade the performance of the preconditioner P_α .

Numerical experiments suggest that the solution of linear systems involving $\alpha I_n + \gamma UU^T$ is **more critical**. Note that this matrix is SPD for all $\alpha > 0$, but ill-conditioned for small α (or very large γ).

The preconditioner (cont.)

The Sherman-Morrison-Woodbury formula yields

$$(\alpha I_n + \gamma UU^T)^{-1} = \alpha^{-1} I_n - \alpha^{-1} \gamma U (\alpha I_k + \gamma U^T U)^{-1} U^T.$$

The main cost is the solution at each step of a $k \times k$ linear system with matrix $\alpha I_k + \gamma U^T U$, which can be performed by Cholesky factorization (computed once and for all at the outset) or possibly by a suitable inner PCG iteration or maybe an (algebraic) MG method.

Note that for incompressible flow problems, $\alpha I_k + \gamma U^T U$ is essentially a (shifted) discrete pressure Laplacian.

In the numerical solution of the Navier–Stokes equations using (say) Picard iteration, this matrix remains constant, whereas the matrix A changes.

Hence, the cost of a Cholesky factorization can be **amortized** over many nonlinear (or time) steps.

Two variants

Building on the main idea, some variants of the preconditioner can be envisioned.

If A happens to be nonsingular and linear systems with A are not too difficult to solve (e.g., well-conditioned), then it may not be necessary to shift A , leading to a preconditioner of the form

$$P_{\alpha,0} = A(\alpha I_n + \gamma UU^T).$$

Note that Theorem 1, however, is no longer applicable.

When A is symmetric positive semidefinite, we'd like the preconditioner to be SPD so that it can be used with the CG method. In this case we can consider a symmetrized version of the preconditioner, for example

$$P_\alpha^S = L(\alpha I + \gamma UU^T)L^T$$

where L is the Cholesky (or incomplete Cholesky) factor of $A + \alpha I$ (or of A itself if A is SPD and not very ill-conditioned).

Possible variants (cont.)

In some cases (but not always) the performance of the method improves if A_γ is diagonally scaled so that it has unit diagonal prior to forming the preconditioner.

Note that the matrix

$$D_\gamma = \text{diag}(A + \gamma UU^T)$$

can be easily computed:

$$(D_\gamma)_{ii} = a_{ii} + \gamma \|u_i^T\|_2^2,$$

where u_i^T is the i th row of U .

It is easy to see that applying the preconditioner to the diagonally scaled matrix $D_\gamma^{-1/2} A_\gamma D_\gamma^{-1/2}$ is equivalent to using the modified preconditioner

$$\hat{P} = (A + \alpha D_\gamma)(\alpha D_\gamma + \gamma UU^T)$$

on the original matrix. Other choices of D_γ are also possible.

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Bounds on the eigenvalues

Let $A_\gamma := A + \gamma UU^T$ and $P_\alpha := \frac{1}{2\alpha}(A + \alpha I)(\alpha I + \gamma UU^T)$.

WLOG we can assume that $\|A\|_2 = 1$ and $\|U\|_2 = 1$. We also assume that A_γ is nonsingular (that is, $\text{Ker}(A) \cap \text{Ker}(U^T) = \{0\}$).

Theorem 2. Let $A + A^T$ be positive semidefinite. If (λ, x) is an eigenpair of the preconditioned matrix $P_\alpha^{-1}A_\gamma$, with $\|x\|_2 = 1$, then

$$\mu < \operatorname{Re}(\lambda) < 2, \quad |\operatorname{Im}(\lambda)| < 1 \tag{2}$$

where

$$\mu = \frac{\alpha \lambda_{\min}(A + A^T)}{(1 + \alpha)(\alpha + \gamma)}.$$

If (λ, x) is an eigenpair with $x \in \text{Ker}(U^T)$, then

$$\lambda = \frac{2x^*Ax}{x^*Ax + \alpha}$$

(independent of γ).

Bounds on the eigenvalues (cont.)

Some comments on this result are in order:

- We see from (2) that the lower bound is uninformative if $A + A^T$ is singular ($\mu = 0$).
- The lower bound (if $\neq 0$) is maximized for $\alpha = \sqrt{\gamma}$.
- Choosing $\alpha = \sqrt{\gamma}$ to maximize the lower bound may not be optimal.
- The lower bound approaches 0 if $\gamma \rightarrow \infty$, indicating that the case of large γ *may* be challenging.
- The result assumes the preconditioner is applied **exactly** (often not true in practice).
- Eigenvalues alone may not be descriptive of GMRES convergence.

Bounds on the eigenvalues (cont.)

γ	α	$\max(Re(\lambda))$	$\min(Re(\lambda))$	lower bound
0.1	0.1	1.818e+00	1.700e-02	5.709e-04
	0.3162	1.519e+00	5.409e-03	7.250e-04
	5.0	3.333e-01	3.430e-04	2.052e-04
1.0	0.5	1.333e+00	6.590e-03	2.791e-04
	1.0	1.000e+00	3.300e-03	3.140e-04
	5.0	4.683e-01	6.609e-04	1.744e-04
50.0	1.0	1.532e+00	3.323e-03	1.231e-05
	7.0711	1.658e+00	4.707e-04	1.928e-05
	10.0	1.606e+00	3.328e-04	1.903e-05

Table: Stokes problem from IFISS, 64×64 mesh and Q2-Q1 discretization. A and U normalized so that $\|A\|_2 = 1 = \|U\|_2$. In boldface the value $\alpha = \sqrt{\gamma}$.

Bounds on the eigenvalues (cont.)

α	$\max(Re(\lambda))$	$\min(Re(\lambda))$	lower bound
0.001	1.998e+00	6.508e-03	7.343e-04
0.01	1.980e+00	6.321e-02	7.213e-03
0.1	1.818e+00	4.834e-01	6.081e-02
0.5	1.333e+00	8.484e-01	1.635e-01
1.0	1.000e+00	5.384e-01	1.839e-01
5.0	4.335e-01	1.372e-01	1.022e-01
10.0	2.457e-01	7.106e-02	6.081e-02
20.0	1.313e-01	3.617e-02	3.337e-02
50.0	5.470e-02	1.463e-02	1.414e-02

Table: Problem `mosarqp1` from Maros and Mészáros collection (Schur complement of KKT system form constrained optimization), $\gamma = 1$. A and U normalized so that $\|A\|_2 = 1 = \|U\|_2$.

Bounds on the eigenvalues (cont.)

α	$\max(Re(\lambda))$	$\min(Re(\lambda))$	lower bound
0.001	1.998e+00	4.247e-03	1.597e-03
0.01	1.980e+00	4.167e-02	1.568e-02
0.1	1.818e+00	3.477e-01	1.322e-01
0.5	1.333e+00	9.087e-01	3.556e-01
1.0	1.000e+00	8.889e-01	4.000e-01
5.0	5.383e-01	2.759e-01	2.222e-01
10.0	3.180e-01	1.481e-01	1.322e-01
20.0	1.738e-01	7.692e-02	7.256e-02
50.0	7.350e-02	3.150e-02	3.076e-02

Table: Problem `1p_fit2p` from SuiteSparse Matrix Collection (sparse-dense least-squares problem), $\gamma = 1$. A and U normalized so that $\|A\|_2 = 1 = \|U\|_2$.

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Numerical experiments with matrices from Stokes and Oseen problems (leaky-lid driven cavity)

We tested **inexact** variants of the proposed preconditioner

$$P_\alpha = \frac{1}{2\alpha} (A + \alpha I)(\alpha I + \gamma UU^T)$$

on a number of linear systems of the form

$$(A + \gamma B^T W^{-1} B)x = b$$

associated with 2D steady Stokes and Oseen problems, varying γ , α , the mesh size h , the viscosity ν and the type of discretization used. Note that $U = B^T W^{-1/2}$ and that W is diagonal. Also, A is **block diagonal**.

For efficiency, we replace the factor $(A + \alpha I)$ with its no-fill Cholesky or ILU factorization, denoted by M_α . The factor $(\alpha I + \gamma B^T W^{-1} B)$ is inverted exactly via the SMW formula and the Cholesky factorization of the $k \times k$ matrix $\alpha I_k + \gamma W^{-1/2} B B^T W^{-1/2}$.

We also describe a simple but effective heuristic for the choice of α .

Numerical experiments with steady 2D Oseen problem

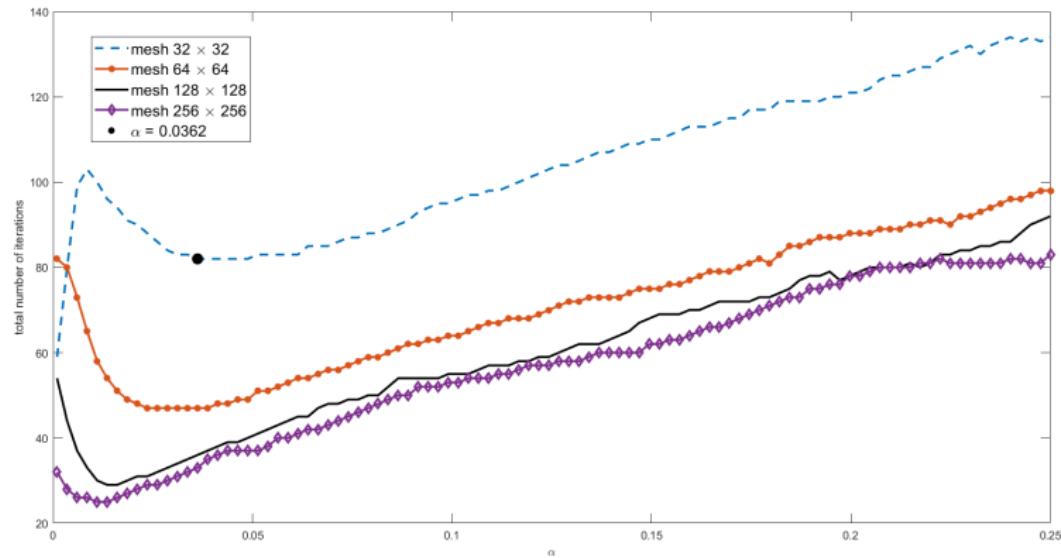


Figure: Number of PGMRES iterations versus α for the 2D Oseen problem (leaky-lid driven cavity from IFISS) with $\nu = 0.01$, $\gamma = 100$, Q2-Q1 finite element discretization and different mesh sizes. GMRES restart $m = 20$, convergence residual tolerance = 1e-06. Diagonal scaling is applied.

A simple heuristic for the choice of α

mesh	α^*	iterations with α^*	optimal α	iterations with optimal α
64×64	0.0256	47	0.0236	47
128×128	0.0181	30	0.0136	29
256×256	0.0128	25	0.0111	25

Table: PGMRES iteration counts for 2D Oseen problem with $\nu = 0.01$, $\gamma = 100$, Q2-Q1 finite element discretization. For 32×32 mesh we find $\alpha^* = 0.0362$, then for the $2^{5+k} \times 2^{5+k}$ mesh we set $\alpha^* = \frac{0.0362}{2^{k/2}}$. That is, we divide α^* by $\sqrt{2}$ each time h is halved.

Numerical experiments for steady 2D Oseen problem

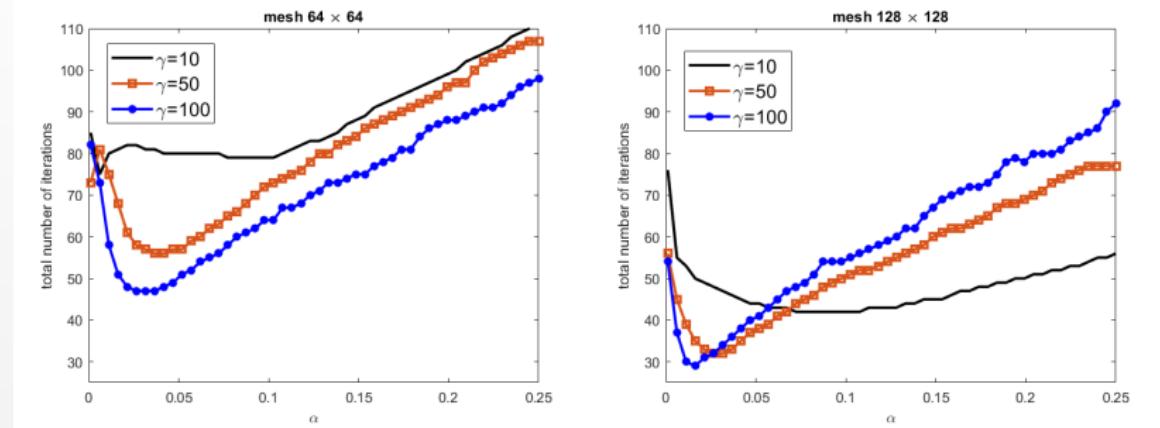


Figure: Number of iterations versus α for the 2D Oseen problem from IFISS with $\nu = 0.01$, Q2-Q1 finite element discretization on 64×64 mesh (LEFT) and on 128×128 mesh (RIGHT) for different values of γ .

Numerical experiments: timings for Oseen problem

mesh	M-Time	P-Time	M_α		P_α	
			Sol-Time	Its	Sol-Time	Its
32×32	8.91e-04	1.55e-03	4.21e-01	1476	3.88e-02	82
64×64	1.79e-03	4.13e-03	1.82e+00	1408	9.11e-02	47
128×128	6.58e-03	2.02e-02	9.50e+00	1516	3.26e-01	30
256×256	2.52e-02	1.16e-01	4.90e+01	1285	1.61e+00	25

Table: Linear system from Oseen problem with $\gamma = 100$, $\nu = 0.01$, Q2-Q1 finite element discretization. Diagonal scaling is used. M_α is ILU(0) of $A + \alpha I$, $P_\alpha = M_\alpha(\alpha I_k + \gamma U^T U)$ with Cholesky factorization of $k \times k$ matrix in SMW formula. M-Time and P-Time are the preconditioner construction times. Sol-Time is the time for the preconditioned iteration to converge.

Note: In this application, the Cholesky factorization needed for inverting $\alpha I + \gamma U^T U$ can be reused many times since U is constant.

Numerical experiments for steady 2D Stokes problem

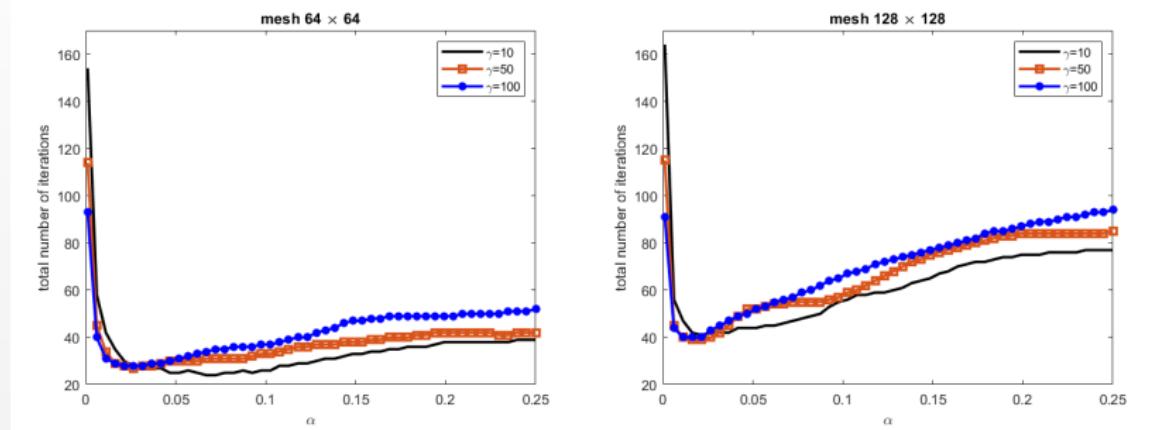


Figure: Number of iterations versus α for the 2D Stokes problem from IFISS with Q2-Q1 finite element discretization on 64×64 mesh (LEFT) and on 128×128 mesh (RIGHT) for different values of γ .

Numerical experiments on reduced KKT systems, I

Problem primal4 (Maros and Mészáros collection), $n = 1489$, $k = 75$.
The reduced system matrix $H + C^T(Z^{-1}\Lambda)C$ has condition number 3.41×10^5 . **Note:** H is singular.

α	M_α	P_α	no prec.
0.001	2000*	13	2000*
0.01	1458	13	
0.1	1379	11	
0.5	1358	9	
1.0	2000*	2	
10.0	2000*	15	
20.0	2000*	18	

Table: Total number of iterations for primal4. M_α is IC(0) of $H + \alpha I$. No diagonal scaling is applied.

Numerical experiments on reduced KKT systems, II

Problem `mosarqp1` problem (Maros ans Mészáros collection), $n = 2500$, $k = 700$. The reduced system matrix $H + C^T(Z^{-1}\Lambda)C$ has condition number 3.35×10^4 .

α	M_α	P_α	no prec.	α	P_α^S	no prec. CG
0.001	2000*	309	2000*	0.001	293	246
0.01	2000*	66		0.01	125	
0.1	2000*	20		0.1	45	
0.5	2000*	9		0.5	20	
1.0	2000*	6		1.0	14	
10.0	2000*	11		10.0	15	
20.0	2000*	13		20.0	17	

Table: Total number of iterations for `mosarqp1`. No diagonal scaling is applied. (LEFT) PGMRES. (RIGHT) PCG. We consider the symmetrized version of the preconditioner: $P_\alpha^S = L(\alpha I_n + \gamma UU^T)L^T$, where L is the no-fill incomplete Cholesky factor of $H + \alpha I_n$.

Numerical experiments on sparse-dense LS problems, I

Problem `1p_fit2p` (SuiteSparse Collection). Here B_1 is 13500×3000 , B_2 is 25×3000 (hence $n = 3000$, $k = 25$), $\kappa(B_1^T B_1 + B_2^T B_2) = 2.52 \times 10^9$.

Note: the norm of $B_2^T B_2$ is 7 orders of magnitude larger than that of $B_1^T B_1$.

α	M_α	P_α	no prec.
0.001	109	10	174
0.01	109	10	
0.1	109	10	
0.5	110	10	
1.0	110	8	
10.0	139	7	
20.0	149	9	

Table: Total number of (P)GMRES iterations for `1p_fit2p`. M_α is IC(0) approximation of $B_1^T B_1 + \alpha I$. No diagonal scaling is used.

We remark that the cost for P_α is only slightly larger than for M_α .

Numerical experiments on sparse-dense LS problems, II

Problem `stormg2-1000`.

Here B_1 is 1377185×528185 , B_2 is 121×528185 (hence we have $n = 528185$, $k = 121$).

α	M_α	P_α	no prec.
0.001	2000*	2000*	2000*
0.01	2000*	334	
0.1	2000*	98	
0.5	2000*	43	
1.0	2000*	50	
5.0	2000*	89	
10.0	2000*	118	
20.0	2000*	176	

Table: Total number of (P)GMRES iterations for `stormg2-1000`. M_α is as before. No diagonal scaling is used.

Numerical experiments: timings for sparse-dense LS problem

α	M_α			P_α		
	M-Time	P-Time	Sol-Time	Its	Sol-Time	Its
0.5	7.05e-02	7.09e-02	1.01e+02	2000*	2.28e+00	43
1.0	6.90e-02	6.93e-02	1.01e+02	2000*	2.61e+00	50
1.5	6.93e-02	6.96e-02	1.01e+02	2000*	3.10e+00	59

Table: stormg2_ 1000 problem. M_α is IC(0) of $A = B_1^T B_1 + \alpha I$. No diagonal scaling is needed. In PGMRES, we do not form the matrix $A = B_1^T B_1$, but we compute mat-vecs as $B_1^T(B_1 x)$. Here B_1 is 1377185×528185 , B_2 is 121×528185 (hence we have $n = 528185$, $k = 121$).

Numerical experiments on sparse-dense LS problems, III

Problem `scfxml1-2r` (SuiteSparse Collection): B_1 is 65886×37980 , B_2 is 57×37980 (so $n = 37980$, $k = 27$), $\kappa(B_1^T B_1 + B_2^T B_2) = 9.32 \times 10^6$.

Note: $A = B_1^T B_1$ is singular.

α	M_α	P_α	no prec.
0.001	1572	555	240
0.01	693	91	
0.1	183	36	
0.5	154	39	
1.0	155	50	
10.0	213	141	

α	P_α^S	no prec.	CG
0.001	1331		184
0.01	415		
0.1	105		
0.5	58		
1.0	65		
10.0	109		

Table: Total number of iterations for `scfxml1-2r` problem. Diagonal scaling is applied. M_α is IC(0) of $B_1^T B_1 + \alpha I$. Left: PGMRES. Right: PCG. We consider the symmetrized version of the preconditioner: $P_\alpha^S = L(\alpha I_n + \gamma UU^T)L^T$, where L is the no-fill incomplete Cholesky factor of $B_1^T B_1 + \alpha I$.

Outline

- 1 Motivation
- 2 The proposed preconditioner
- 3 Eigenvalue bounds
- 4 Numerical experiments
- 5 Conclusions and future work

Conclusions

- Introduced a new solver for a wide class of tough linear systems
- The proposed preconditioner seems to work well in practice
- Solves with $A + \alpha I_n$ can be (very) inexact...
- ... but exact solves with $\alpha I + \gamma UU^T$ seem to be necessary
- SMW formula \Rightarrow only a $k \times k$ solve needed
- Often, much of this work can be reused
- For some PDE problems we found a simple heuristic for choosing α
- Some theory available for ideal case

Future work

- Further investigate the spectrum of $P_\alpha^{-1} A_\gamma$. Clustering?
- Try to find eigenvalue bounds for singular A
- More work to be done on the choice of α
- Investigate other approaches to solving systems with $\alpha I + \gamma UU^T$ (iterative?)
- Are there better ways to symmetrize the preconditioner for use with CG?
- Extension to matrices of the form $A + \gamma UV^T$, with $V \neq U$.