## An introduction to fractional calculus

**Fundamental ideas and numerics** 

#### Fabio Durastante

Università di Pisa





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We want to find a **numerical solution** of the differential equation written in terms of Caputo Derivatives

$$\alpha > 0, \quad m = \lceil \alpha \rceil, \qquad \begin{cases} c D_{[0,t]}^{\alpha} \mathbf{y}(t) = f(t, \mathbf{y}(t)), & t \in [0, T], \\ \frac{d^k \mathbf{y}(0)}{dt^k} = \mathbf{y}_0^{(k)}, & k = 0, 1, \dots, m-1. \end{cases}$$
 (FODE)

#### Caputo fractional derivative (Caputo 2008)

Let  $\alpha \geq 0$ , and  $m = \lceil \alpha \rceil$ . Then, we define the operator

$$_{C}D_{[a,t]}^{\alpha}y=I_{[a,t]}^{m-\alpha}\frac{d^{m}}{dt^{m}}y,$$

whenever  $\frac{d^m}{dt^m}y \in \mathbb{L}^1([a,b])$ .

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For both *implicit* and *explicit* methods we have also all-at-once formulations, and a middle-ground represented by *Implicit-Explicit* (IMEX) methods.

Our *objective* is to transport what we can for the solution of (FODE).

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$$y(t) = \sum_{k=0}^{m-1} \frac{t^k}{k!} y_0^{(k)} + \frac{1}{\Gamma(\alpha)} \int_0^t (t-\tau)^{\alpha-1} f(\tau, y(\tau)) d\tau, \quad m = \lceil \alpha \rceil.$$

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- Adams-Bashforth-Moulton methods are obtained by applying a quadrature formula to the integral,
- We can use, e.g.,

  - the fractional rectangular formula with nodes  $\{t_j=j\tau\}_{j=1}^{n-1}$ , or the product trapezoidal quadrature formula with nodes  $\{t_j=j\tau\}_{j=1}^n$ .

To obtain a predictor-corrector method.

The main idea behind PI rules is to approximate the integral

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$$y(t) = T_{m-1}(t) + \frac{1}{\Gamma(\alpha)} \sum_{i=0}^{m-1} \int_{t_j}^{t_{j+1}} (t-s)^{\alpha-1} f(\tau, y(\tau)) d\tau, \qquad t \geq t_n.$$

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Replace f in each sub-interval by the first-degree polynomial interpolant

$$p_j( au) = f_{j+1} + rac{s - t_{j+1}}{ au_i}, \quad s \in [t_j, t_j + 1], \quad au_j = t_{j+1} - t_j, \quad f_j = f(t_j, y_j).$$

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- These produce the usual fractional integral that we now how to solve

$$I_{n,j}^{(k)} = \frac{1}{\Gamma(\alpha)} \int_{t_j}^{t_n} (t_n - \tau)^{\alpha - 1} (\tau - t_j)^k d\tau = \frac{(t_n - t_j)^{\alpha + k}}{\Gamma(\alpha + k + 1)}.$$

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- Replace f in each sub-interval by the first-degree polynomial interpolant
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- We plug everything in our expression using that:

$$w_n = I_{n,0}^{(0)} - \frac{I_{n,0}^{(1)}}{\tau_0} + \frac{I_{n,1}^{(1)}}{\tau_0}, \quad b_{n_j} = \frac{I_{n,j-1}^{(1)} - I_{n,j}^{(1)}}{h_{j-1}} - \frac{I_{n,j}^{(1)} - I_{n,j+1}^{(1)}}{\tau_j}, j \leq n-1, \quad b_{n,n} = \frac{I_{n,n-1}^{(1)}}{\tau_{n-1}}.$$

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$$\int_0^t (t-s)^{-lpha} \mathcal{K}(t,s) y(s) \, \mathrm{d}s = f(t), \quad 0 < lpha < 1$$
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If we discretize everything as before we get

$$[B_N \odot K_N]\mathbf{y} = \mathbf{g}, \quad B_N = \tau^{1-\alpha}[b_{i,j}], \quad K_N = [k(t_i,t_j)], \quad \odot \text{ Hadamard product.}$$

where  $\mathbf{y} = (y_0, \dots, y_N)^T$  and g contains the **initial conditions** and the **evaluations** of f.

#### Convergence analysis for (Cameron and McKee 1985)

"[Consistency of order p] demands that  $f(t) \in \mathcal{C}^{1-\alpha}[0,T]$  which is necessary in any case for y(t) to be a smooth function ...  $|y(t_i)-y_i| \leq C\tau^p$ ,  $i=0,1,\ldots,m-1$ ."

The requirements from the standard theory are **far too strong** for what we can reasonably expect from the analysis on the solution regularity we did in the last lecture.

#### Theorem (Dixon 1985)

Let f be Lipschitz continuous with respect to the second variable and  $y_n$  be the numerical approximation obtained by applying the PI trapezoidal rule on the interval  $[t_0, T]$ . There exist a constant  $C = C_1(T - t_0)$ , which does not depend on h, such that

$$||y(t_n)-y_n|| \leq C(t_n^{\alpha-1}\tau^{1+\alpha}+\tau^2), \qquad \tau = \max_{j=0,\dots,n-1}\tau_j.$$

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- The same drop in the convergence order occurs also when higher degree polynomials are employed,
- When  $\alpha > 1$  convergence order 2 is obtained.
- $\blacksquare$  It doesn't make much sense to use higher-degree PI rules if  $0 < \alpha < 1$ .

Let us reduce to the case with  $\alpha \in (0,1)$ , m=1, and a uniform mesh. To build it we need to approximate the integral with the rectangule rule

$$\int_0^t (t-\tau)^{\alpha-1} f(\tau, y(\tau)) \, \mathrm{d}\tau$$

on the grid  $\{t_j = t_0 + j\tau\}_{j=1}^N$  with *uniform* grid spacing  $\tau$ , we denote

$$f^{(j)} = f(t_j, y^{(j)}) \text{ for } y^{(j)} \approx y(t_j),$$

and write it as

$$y^{(n)} = y_0 + \frac{\tau^{\alpha}}{\Gamma(\alpha)} \sum_{i=0}^{n-1} b_{n-j-1} f^{(j)}, \quad b_n = [(n+1)^{\alpha} - n^{\alpha}]/\alpha, \quad n = 1, \dots, N.$$

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$$y^{(n)} = y_0 + \frac{\tau^{\alpha}}{\Gamma(\alpha)} \left( b_0 f^{(n-1)} + \sum_{j=0}^{n-2} b_{n-j-1} f^{(j)} \right) \quad b_n = [(n+1)^{\alpha} - n^{\alpha}]/\alpha, \quad n = 1, \dots, N.$$

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- Using a uniform mesh the evaluation of the weights just involve the computation of real powers of integer numbers, we can simplify also the fractional trapezoidal formula

$$y^{(n)} = T_{m-1}(t_n) + \frac{\tau^{\alpha}}{\Gamma(\alpha+2)} \left( w_n f^{(0)} + \sum_{j=1}^n b_{n-j} f^{(j)} \right),$$

$$w_n = (\alpha+1-n)n^{\alpha} + (n-1)^{\alpha+1},$$

$$b_0 = 1, \ b_n = (n-1)^{\alpha+1} - 2n^{\alpha+1} + (n+1)^{\alpha+1}, \ n \ge 1.$$

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- **\'\>** We have to compute:

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## Some observations

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## Predictor-Corrector algorithms

Now that we have two schemes we can think of using them together to build a **predictor-corrector** algorithm.

## Fractional Predictor-Corrector Scheme (Diethelm 1997)

We are going to write it again for  $0 < \alpha < 1$  on a uniform mesh

1. In the prediction step we use the fractional rectangular formula

$$y_P^{(n+1)} = y^{(0)} + \frac{\tau^{\alpha}}{\Gamma(\alpha)} \sum_{i=0}^n b_{j,n+1} f(t_j, y^{(j)}), \quad b_{j,n+1} = \frac{(n+1-j)^{\alpha} - (n-j)^{\alpha}}{\alpha}$$

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2. In the correction step we use the fractional trapezoidal formula

$$y^{(n+1)} = y^{(0)} + \frac{\tau^{\alpha}}{\Gamma(\alpha)} \left( \sum_{j=0}^{n} a_{j,n+1} f(t_j, y^{(j)}) + a_{n+1,n+1} f(t_{n+1}, y_P^{(n+1)}) \right)$$

where

$$a_{j,n+1} = \begin{cases} (n^{\alpha+1} - (n-\alpha)(n+1)^{\alpha}/\alpha(\alpha+1), & j = 0, \\ (n-j+2)^{\alpha+1} - 2(n-j+1)^{\alpha+1} + (n-j)^{\alpha+1}/\alpha(\alpha+1), & j = 1, 2, \dots, n, \\ 1/\alpha(\alpha+1), & j = n+1. \end{cases}$$

• Predictor-Corrector schemes are of interest because they represent a good compromise between accuracy and ease of implementation.

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- To investigate the convergence we need to look deeper into the convergence results of the two PI integral rules (Diethelm, Ford, and Freed 2004).

### Theorem (Diethelm, Ford, and Freed 2004, Theorem 2.4)

(a) Let  $z \in \mathcal{C}^1([0, T])$ . Then

$$\left| \int_0^{t_{k+1}} (t_{k+1} - t)^{\alpha - 1} z(t) dt - \sum_{j=0}^k b_{j,k+1} z(t_j) \right| \leq \frac{1}{\alpha} \|z'\|_{\infty} t_{k+1}^{\alpha} \tau.$$

(b) Let  $z(t) = t^p$  for some  $p \in (0,1)$ . Then,

$$\left| \int_0^{t_{k+1}} (t_{k+1} - t)^{\alpha - 1} z(t) dt - \sum_{j=0}^k b_{j,k+1} z(t_j) \right| \le C_{\alpha,\rho}^{Re} t_{k+1}^{\alpha + \rho - 1} \tau.$$

And analogously for the product trapezoidal formula.

#### Theorem (Diethelm, Ford, and Freed 2004, Theorem 2.5).

(a) If  $z \in C^2([0,T])$ , then there exist a constant  $C_{\alpha}^{Tr}$  depending only on  $\alpha$  such that

$$\left| \int_0^{t_{k+1}} (t_{k+1} - t)^{\alpha - 1} z(t) dt - \sum_{j=0}^{k+1} a_{j,k+1} z(t_j) \right| \le C_{\alpha}^{Tr} ||z''||_{\infty} t_{k+1}^{\alpha} \tau^2.$$

(b) Let  $z \in \mathcal{C}^1([0,T])$  and assume that z' fulfills a Lipschitz condition of order  $\mu \in (0,1)$ . Then, there exists positive constants  $B_{\alpha,\mu}^{Tr}$  and  $M_{z,\mu}$  such that

$$\left| \int_0^{t_{k+1}} (t_{k+1} - t)^{\alpha - 1} z(t) dt - \sum_{j=0}^{k+1} a_{j,k+1} z(t_j) \right| \leq B_{\alpha,\mu}^{Tr} M_{z,\mu} t_{k+1}^{\alpha} \tau^{1+\mu}.$$

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(b) Let  $z(t) = t^p$  for some  $p \in (0,2)$  and  $\rho = \min(2, p+1)$ . Then

$$\left| \int_0^{t_{k+1}} (t_{k+1} - t)^{\alpha - 1} z(t) dt - \sum_{j=0}^{k+1} a_{j,k+1} z(t_j) \right| \le C_{\alpha,p}^{Tr} t_{k+1}^{\alpha + p - \rho} \tau^{\rho}.$$

Observe that for the fractional rectangular case (b) the bound contains

$$t_{k+1}^{\alpha+p-1},$$

if  $\alpha + p < 1$  then we get that the overall integration error becomes larger if the size of the interval of integration becomes smaller!

Similarly for the case (c) for the fractional trapezoidal rule

$$\alpha < 1, \ \rho < 1, \ \rho = \rho + 1, \quad t_{k+1}^{\alpha + \rho - \rho},$$

has the same explosive behavior.

#### Smaller intervals for harder integrals

By making  $t_{k+1}$  smaller we have two effects

- 1. We reduce the length of the integration interval,
- 2. We change the weight function in a way that makes the integral more difficult.

#### Lemma (Diethelm, Ford, and Freed 2004, Lemma 3.1)

Assume that the solution y of the initial value problem is such that

$$\left| \int_0^{t_{k+1}} (t_{k+1} - t)^{\alpha - 1} {}_{CA} D^{\alpha}_{[0,t]} y(t) dt - \sum_{j=0}^k b_{j,k+1} {}_{CA} D^{\alpha}_{[0,t]} y(t) \right| \leq C_1 t_{k+1}^{\gamma_1} \tau^{\delta_1},$$

and

$$\left| \int_0^{t_{k+1}} (t_{k+1} - t)^{\alpha - 1} {}_{CA} D_{[0,t]}^{\alpha} y(t) dt - \sum_{j=0}^{k+1} a_{j,k+1} {}_{CA} D_{[0,t]}^{\alpha} y(t) \right| \leq C_2 t_{k+1}^{\gamma_2} \tau^{\delta_2},$$

with some  $\gamma_1, \gamma_2 \geq 0$  and  $\delta_1, \delta_2 > 0$ . Then, for some suitably chosen T > 0, we have

$$\max_{0 \leq i \leq N} |y(t_j) - y^{(j)}| = O(\tau^q), \quad q = \min\{\delta_1 + \alpha, \delta_2\}, \quad N = \lceil T/\tau \rceil.$$

### Theorem (Diethelm, Ford, and Freed 2004, Theorem 3.2)

Let 0<lpha and assume  ${}_{C\!A}D^lpha_{[0,t]}y(t)\in\mathcal{C}^2([O,T])$  for some suitable T . Then,

$$\max_{0 \leq j \leq N} |y(t_j) - y^{(j)}| = egin{cases} O( au^2), & ext{if } lpha \geq 1, \ O( au^{1+lpha}), & ext{if } lpha < 1. \end{cases}$$

**Proof.** In view of the two bounds for the Fractional Rectangular and Trapezoidal forms we can apply the previous Lemma with  $\gamma_1=\gamma_2=\alpha>0$ ,  $\delta_1=1$ ,  $\delta_2=2$ . Therefore we find a bound of order  $O(\tau^q)$  where

$$q=\min\{1+lpha,2\}=egin{cases} 2, & ext{if } lpha\geq 1, \ 1+lpha, & ext{if } lpha<1. \end{cases}$$

### Theorem (Diethelm, Ford, and Freed 2004, Theorem 3.2)

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- Can we replace them by similar assumptions on y itself?

### Theorem Diethelm, Ford, and Freed 2004, Theorem 3.3

Let 
$$\alpha>1$$
 and assume  $y\in \mathcal{C}^{1+\lceil\alpha\rceil}([0,T])$  for some suitable  $T$ , then 
$$\max_{0\leq j\leq N}|y(t_j)-y^{(j)}|=O(\tau^{1+\lceil\alpha\rceil-\alpha}).$$

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**Proof.** We need to use the characterization of Caputo's derivative

$$_{CA}D_{[0,t]}^{lpha}y(t) = \sum_{\ell=0}^{m-|lpha|-1}rac{y^{(l+\lceillpha
ceil)}(0)}{\Gamma(\lceillpha
ceil-lpha+\ell+1)}t^{\lceillpha
ceil-lpha+\ell}+g(t), \qquad egin{align*} g\in\mathcal{C}^{m-\lceillpha
ceil}([O,T]), & g^{(m-\lceillpha
ceil)}\in \mathrm{Lip}(\lceillpha
ceil-lpha). & g^{(m-\lceillpha
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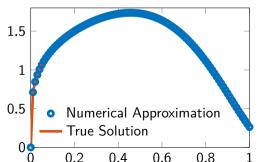
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**Proof.** Then for  $\alpha>1$ , we can apply the Lemma with  $\gamma_1=0$ ,  $\gamma_2=\alpha-1>0$ ,  $\delta_1=1$ ,  $\delta_2=1+\lceil\alpha\rceil-\alpha$  and thus  $\delta_1+\alpha=1+\alpha>2>\delta_2$ ,  $\min\{\delta_1+\alpha,\delta_2\}=\delta_2$ . The overall order is then  $O(\tau^{\delta_2})=O(\tau^{1+\lceil\alpha\rceil-\alpha})$ .

#### Example

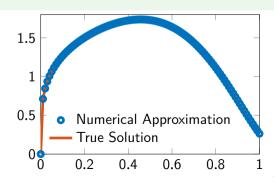
$$\begin{cases} {}_{CA}D^{\alpha}_{[0,t]}y(t) = \frac{40320}{\Gamma(9-\alpha)}t^{8-\alpha} - 3\frac{\Gamma(5+\alpha/2)}{\Gamma(5-\alpha/2)}t^{4-\alpha/2} + \frac{9}{4}\Gamma(\alpha+1) + \left(3t^{\alpha/2}/2 - t^4\right)^3 - y(t)^{3/2}, \\ y(0) = 0. \end{cases}$$



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α	τ	Е	q
0.25	5.00e-01	1.42e+00	
	2.50e-01	4.17e-01	1.77
	1.25e-01	2.13e-01	0.97
	6.25e-02	1.03e-01	1.05
	3.12e-02	5.04e-02	1.03
	1.56e-02	2.44e-02	1.05



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	1.56e-02	2.44e-02	1.05

```
hval = 2.^(-(1:6));

for i=1:length(hval)

h = hval(i);

t0 = 0; T = 1;

[T, Y] = fde_pi12_pc(alpha, f_fun,

t0, T, y0, h, [], 1);

err(i) = norm(Y - ye(T), 'inf');

end
```

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	2.50e-01	4.17e-01	1.77		2.50e-01	1.80e + 00	0.61
	1.25e-01	2.13e-01	0.97		1.25e-01	8.37e-01	1.10
	6.25e-02	1.03e-01	1.05		6.25e-02	2.45e-01	1.77
	3.12e-02	5.04e-02	1.03		3.12e-02	6.57e-02	1.90
	1.56e-02	2.44e-02	1.05		1.56e-02	2.02e-02	1.70

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0.25	5.00e-01	1.42e+00			1.95e-03	9.33e-04	1.4
	2.50e-01	4.17e-01	1.77		9.77e-04	3.58e-04	1.3
	1.25e-01	2.13e-01	0.97		4.88e-04	1.40e-04	1.3
	6.25e-02	1.03e-01	1.05		2.44e-04	5.56e-05	1.3
	3.12e-02	5.04e-02	1.03		1.22e-04	2.23e-05	1.3
	1.56e-02	2.44e-02	1.05		6.10e-05	9.00e-06	1.3

# More than one correction step

One can think of improving convergence by performing **more than one correction step** in the algorithm (Diethelm, Ford, and Freed 2002).

Let us call  $\mu \in \mathbb{N}$  the number of correction steps:

$$\begin{cases} y_{[0]}^{(n+1)} = y^{(0)} + \frac{\tau^{\alpha}}{\Gamma(\alpha)} \sum_{j=0}^{n} b_{j,n+1} f(t_{j}, y^{(j)}), & \text{Prediction step,} \\ y_{[\ell]}^{(n+1)} = y^{(0)} + \frac{\tau^{\alpha}}{\Gamma(\alpha)} \left( \sum_{j=0}^{n} a_{j,n+1} f(t_{j}, y^{(j)}) + a_{n+1,n+1} f(t_{n+1}, y_{[\ell-1]}^{(n+1)}) \right), & \text{Correction steps} \\ y^{(n+1)} \equiv y_{[\mu]}^{(n+1)}. & \ell = 1, \dots, \mu. \end{cases}$$

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lacktriangle The standard predictor corrector method is obtained for  $\mu=1.$ 

The convergence behavior can be described by using repeatedly the result from (Diethelm, Ford, and Freed 2004, Lemma 3.1) that we have used to obtain the other convergence bounds.

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### Corollary

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- In the third case with a single corrector step, and no improvement is possible.
- $\P$  In general we could fix a maximum number of steps  $\mu$  and halt the procedure when the error is under a certain tolerance.

Let us focus on the **test problem** 

$$_{CA}D^{\alpha}_{[t_0,t]}y(t)=\lambda y(t),\quad y(0)=y_0,\quad \lambda\in\mathbb{C},\quad 0<\alpha<1.$$

In the last lecture we have seen that the solution of this problem can be expressed as

$$y(t) = E_{\alpha}(\lambda(t-t_0)^{\alpha})y_0.$$

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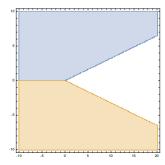
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The solution y(t) asymptotically vanishes as  $t \to +\infty$  for

$$\lambda \in S^* = \{z \in \mathbb{C} \, : \, |\arg(z) - \pi| < (1 - \alpha/2)\pi.\}$$



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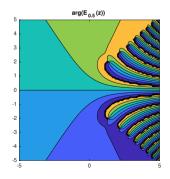
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The application of PI rule leads to a non-homogeneous difference equation ,

$$y^{(n)} = g^{(n)} + \sum_{i=k}^{n} c_{n-j} y^{(j)}, \quad n \ge k,$$



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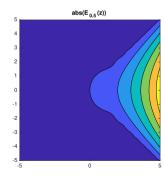
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### Informally

The stability region of the various PI formulas can be described as the set of all  $z=\tau^{\alpha}\lambda$  for which the numerical solution  $\{y^{(n)}\}_n$  behaves as the true solution and tends to 0 as  $n\to +\infty$ .

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As for the other theoretical result we are going to leverage information on the associated Volterra integral equation (Lubich 1986a).

• First we rewrite our non-homogeneous difference equation (in which we simplify the notation assuming to work with scalars) as

$$\begin{cases} y_n = f_n + \tau^{\alpha} \sum_{j=0}^{n} \omega_{n-j} g(y_j), & n \geq 0 \\ f_n = f(t_n) + \tau^{\alpha} \sum_{j=-m}^{-1} w_{n,j} g(y_j), & t_n = t_0 + n\tau, \quad t_0 = mh. \end{cases}$$

• Then we assume that  $h^{\alpha}w_{n,j}g(y_j)=O((n\tau)^{\alpha-1}\tau g(y_j))$ , i.e.,  $w_{n,j}=O(n^{\alpha-1})$  as  $n\to +\infty,\ j=-M,\ldots,-1$ .

#### A connection to the classical theory

In the classical case  $\alpha = 1$ , if we can express the term

$$\sum_{n=0}^{+\infty} \omega_n \zeta^n = \frac{\sigma(\zeta^{-1})}{\rho(\zeta^{-1})}$$

as a rational function, then we have found a standard Linear Multistep Method.

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#### A-stable method

A convolution quadrature  $\{\omega\}_n$  for the Abel equation

$$y(t) = f(t) + \frac{1}{\Gamma(\alpha)} \int_{0}^{t} (t-s)^{\alpha-1} g[y(s)] ds, \quad t \ge 0, \ 0 < \alpha \le 1,$$

is called A-stable if the solution  $\{y_n\}_n$  given by the convolution quadrature satisfies

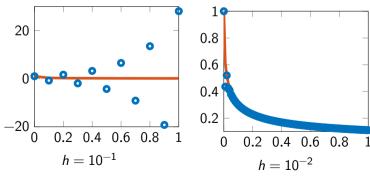
$$y_n \to 0$$
 as  $n \to +\infty$  whenever  $\{f_n\}_n$  has a finite limit  $\forall \tau > 0, \ \forall \ \lambda \in S^*$ .

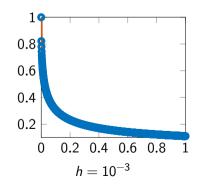
# Stability region

In general we cannot expect to have stability for every  $\lambda \in S^*$ , consider, e.g.

$$_{CA}D^{\alpha}_{[t_0,t]}y(t) = -5y(t), \quad y(0) = 1, \quad T = 1.$$

integrated with the explicit fractional rectangular rule





#### Stability region

The stability region S of a convolution quadrature  $\{\omega_m\}$  is the set of all complex  $z = \tau^{\alpha}\lambda$  for which the numerical solution  $\{y_n\}_n$  satisfies

 $y_n \to 0$  as  $n \to +\infty$  whenever  $\{f_n\}_n$  has a finite limit.

The method is called *strongly stable*, if for any  $\lambda \in S^*$  there exists  $\tau_0(\lambda) > 0$  such that  $\tau^{\alpha}\lambda \in S$  for all  $0 < \tau < \tau_0(\lambda)$ . The method is called  $A(\theta)$ -stable if S contains the sector  $|\arg(z) - \pi| < \theta$ .

#### Stability region

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To obtain the characterization we need, we consider weights

$$\omega_n = (-1)^n {-\alpha \choose n} + \nu_n, \quad n \ge 0, \{\nu_n\}_n \in \ell^1,$$
 (H<sub>1</sub>)

to which corresponds

$$\omega(\zeta) = (1-\zeta)^{-\alpha} + \nu(\zeta) \text{ continuous in } \{\zeta \in \mathbb{C} \,:\, |\zeta| \leq 1, \ \zeta \neq 1\}, \ \lim_{\zeta \to 1^-} \omega(\zeta) = +\infty.$$

#### Theorem (Lubich 1986a, Theorem 2.1)

The stability region of a convolution quadrature under the condition  $(H_1)$  is

$$S = \mathbb{C} \setminus \{1/\omega(\zeta) : |\zeta| \leq 1\}$$
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**Proof.** Let  $z = \tau^{\alpha} \lambda$ . Since 0 is neither contained in  $S^*$  nor in S, we can assume  $z \neq 0$ . We can rewrite our difference equation as

$$y(\zeta) = f(\zeta) + z\omega(\zeta)y(\zeta) \iff y(\zeta) = \frac{f(\zeta)}{1 - z\omega(\zeta)} = \frac{(1 - \zeta)^{\alpha}f(\zeta)}{(1 - \zeta)^{\alpha}[1 - z\omega(\zeta)]}.$$

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We first prove that  $S \subseteq S^*$ .

• The coefficient sequence  $(1-\zeta)^{\alpha}[1-z\omega(\zeta)]$  is in  $\ell^1$ , indeed  $v(\zeta)$  and  $(1-\zeta)^{\alpha}$  are in  $\ell^1$  by using (H<sub>1</sub>) (for the first one with  $-\alpha$  instead of  $\alpha$ ), hence also  $1+(1-\zeta)^{\alpha}v(\zeta)=(1-\zeta)^{\alpha}\omega(\zeta)$ , since for any two sequences in  $\ell^1$  we have  $\sum_n |\sum_i a_{n-i}b_i| \leq \sum_i |a_i||b_i|$ .

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- The coefficient sequence  $(1-\zeta)^{\alpha}[1-z\omega(\zeta)]$  is in  $\ell^1$ ,
- If  $z \in S$  then  $1 z\omega(\zeta) \neq 0$  for  $|\zeta| \leq 1$  with  $\zeta \neq 1$ .

#### Wiener inversion Theorem

$$f(\zeta) = \sum_{n=0}^{+\infty} a_n \zeta^n$$
 with  $||f||_1 < +\infty$ ,  $\zeta = e^{in\theta}$ , then  $1/f(\theta) \in \ell^1$  iff  $f(\theta) \neq 0$  for all  $\theta$ .

**Proof.** Let  $z = \tau^{\alpha} \lambda$ . Since 0 is neither contained in  $S^*$  nor in S, we can assume  $z \neq 0$ . We can rewrite our difference equation as

$$y(\zeta) = f(\zeta) + z\omega(\zeta)y(\zeta) \iff y(\zeta) = \frac{f(\zeta)}{1 - z\omega(\zeta)} = \frac{(1 - \zeta)^{\alpha}f(\zeta)}{(1 - \zeta)^{\alpha}[1 - z\omega(\zeta)]}.$$

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$$(\mathsf{H}_1)\ (1-\zeta)^\alpha[1-z\omega(\zeta)]=(1-\zeta)^\alpha[1-zv(\zeta)]-z$$
 and thus

$$(1-\zeta)^{\alpha}[1-z\omega(\zeta)]\neq 0$$
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 and thus

$$(1-\zeta)^{\alpha}[1-z\omega(\zeta)]\neq 0 \text{ for } |\zeta|\leq 1 \Rightarrow 1/(1-\zeta)^{\alpha}[1-z\omega(\zeta)]\in \ell^1.$$

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The stability region of a convolution quadrature under the condition  $(H_1)$  is

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**Proof.** We first prove that  $S \subseteq S^*$ . Let  $\tilde{f}_n = f_n - f_\infty \stackrel{n \to +\infty}{=} 0$  so that we can write

$$f(\zeta) = \frac{f_{\infty}}{1-\zeta} + \tilde{f}(\zeta) \ \Rightarrow \ (1-\zeta)^{\alpha} f(\zeta) = (1-\zeta)^{\alpha-1} f_{\infty} + (1-\zeta)^{\alpha} \tilde{f}(\zeta) \ \text{has coefficients} \ \to 0.$$

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By (H1) the coefficient sequence of  $(1-\zeta)^{\alpha-1} \to 0$ .

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By (H<sub>1</sub>) the coefficient sequence of  $(1-\zeta)^{\alpha-1}\to 0$ . The coefficient sequence of  $(1-\zeta)^{\alpha}\tilde{f}(\zeta)\to 0$  since  $(1-\zeta)^{\alpha}\in \ell_1$  and  $\ell_1*c_0\subset c_0$  for \* the convolution operator, and  $c_0$  the space of zero sequences

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The stability region of a convolution quadrature under the condition  $(H_1)$  is

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**Proof.** To conclude we need to prove that  $S^*$  is exhausted by S, we assume that

$$1-z\omega(\zeta_0)=0$$
 for some  $|\zeta_0|\leq 1$  and by  $(\mathsf{H}_1)$   $\zeta_0\neq 1$ ,

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$$y(\zeta) = \frac{(1-\zeta)^{\alpha}}{\zeta - \zeta_0} = \frac{(1-\zeta)^{\alpha} - (1-\zeta_0)^{\alpha}}{\zeta - \zeta_0} + (1-\zeta_0)^{\alpha} \frac{1}{\zeta - \zeta_0}.$$

#### Lemma (Lubich 1986a, Lemma 2.1)

Assume that the coefficient sequence of  $a(\zeta)$  is in  $\ell^1$ . Let  $|\zeta_0| \leq 1$ . Then the coefficient sequence of

$$\frac{a(\zeta)-a(\zeta_0)}{\zeta-\zeta_0}$$
 converges to zero.

**Proof.** To conclude we need to prove that  $S^*$  is exhausted by S, we assume that

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and show that then  $z \notin S^*$ . We select

$$y(\zeta) = \frac{(1-\zeta)^{\alpha}}{\zeta - \zeta_0} = \underbrace{\frac{(1-\zeta)^{\alpha} - (1-\zeta_0)^{\alpha}}{\zeta - \zeta_0}}_{=0} + (1-\zeta_0)^{\alpha} \frac{1}{\zeta - \zeta_0}.$$

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and show that then  $z \notin S^*$ . We select

$$y(\zeta) = \frac{(1-\zeta)^{\alpha}}{\zeta - \zeta_0} = +(1-\zeta_0)^{\alpha} \frac{1}{\zeta - \zeta_0}.$$

On the other hand,  $1/\zeta - \zeta_0 = -\sum_{n=0}^{+\infty} \zeta_0^{-n-1} \zeta^n$  diverges! Hence also the sequence associated to  $y(\zeta)$  diverges.

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**Proof.** We can now collect the various parts together

$$f(\zeta) = [1 - z\omega(\zeta)]y(\zeta) = (1 - \zeta)^{\alpha}[1 - z\omega(\zeta)](1 - \zeta)^{-\alpha}y(\zeta)$$

$$= \frac{(1 - \zeta)^{\alpha}(1 - z\omega(\zeta)) - (1 - \zeta_0)(1 - z\omega(\zeta_0))}{\zeta - \zeta_0}$$

using again the lemma we get that  $\{f_n\}_n$  goes to zero, but,  $\{y_n\}_n$  does not, hence  $z \notin S^*$ .

#### Theorem (Lubich 1986a, Theorem 2.1)

The stability region of a convolution quadrature under the condition  $(H_1)$  is

$$S = \mathbb{C} \setminus \{1/\omega(\zeta) : |\zeta| \leq 1\}$$
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#### Corollary

If a convolution quadrature satisfying  $(H_1)$  is applied to the Volterra equation and if  $\tau^{\alpha}\lambda \in S$ , then  $\{y_n\}_n$  is bounded whenever  $\{f_n\}_n$  is bounded. Conversely, if  $\{y_n\}_n$  is bounded whenever  $\{f_n\}_n$  is bounded then  $\tau^{\alpha}\lambda \in \overline{S}$ .

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#### Corollary

The stability region of an explicit convolution quadrature ( $\omega_0=0$ ) satisfying (H<sub>1</sub>) is bounded.

**Proof.** By the open mapping theorem  $\omega(\zeta)$  maps neighborhood of 0 into neighborhood of 0. Hence  $S^*$  is a neighborhood of  $\infty$ , and the result follows from the Theorem.

### Theorem (Lubich 1986a, Theorem 2.1)

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The stability region of an explicit convolution quadrature ( $\omega_0 = 0$ ) satisfying (H<sub>1</sub>) is bounded.

#### Corollary

Every convolution quadrature satisfying  $(H_1)$  is strongly stable.

Using these results we can recover the stability regions for the different methods,

For the Predictor-Corrector method we have

$$\begin{cases} y_P^{(n+1)} = y^{(0)} + \frac{\tau^{\alpha}}{\Gamma(\alpha)} \sum_{j=0}^n b_{j,n+1} f(t_j, y^{(j)}), \\ y^{(n+1)} = y^{(0)} + \frac{\tau^{\alpha}}{\Gamma(\alpha)} \left( \sum_{j=0}^n a_{j,n+1} f(t_j, y^{(j)}) + a_{n+1,n+1} f(t_{n+1}, y_P^{(n+1)}) \right) \end{cases}$$

where

$$b_{j,n+1} = \frac{(n+1-j)^{\alpha} - (n-j)^{\alpha}}{\alpha}$$

$$a_{j,n+1} = \begin{cases} (n^{\alpha+1} - (n-\alpha)(n+1)^{\alpha}/\alpha(\alpha+1), & j = 0, \\ (n-j+2)^{\alpha+1} - 2(n-j+1)^{\alpha+1} + (n-j)^{\alpha+1}/\alpha(\alpha+1), & j = 1, 2, \dots, n, \\ 1/\alpha(\alpha+1), & j = n+1. \end{cases}$$

For the Predictor-Corrector method we have

$$\begin{cases} y_P^{(n+1)} = y^{(0)} + \tau^{\alpha} \sum_{j=0}^n b_{n-j-1} f(t_j, y^{(j)}), \\ y^{(n+1)} = y^{(0)} + \tau^{\alpha} a_{n,0} f^{(0)} + \tau^{\alpha} \sum_{j=1}^n a_{n-j} f(t_n, y_P^{(n+1)}) \end{cases}$$

where

$$\begin{split} b_n = & \frac{(n+1)^{\alpha} - n^{\alpha}}{\Gamma(\alpha+1)} \\ a_{n,0} = & \frac{(n-1)^{\alpha+1} - n^{\alpha}(n-\alpha-1)}{\Gamma(\alpha+2)}, \\ a_n = & \begin{cases} 1/\Gamma(\alpha+2), & n = 0, \\ (n-1)^{\alpha+1} - 2n^{\alpha+1} + (n+1)^{\alpha+1}/\Gamma(\alpha+2), & n \geq 1. \end{cases} \end{split}$$

For the Predictor-Corrector method we have

$$y^{(n)} = g^{(n)} + \sum_{j=k}^{n} c_{n-j} y^{(j)}, \quad n \geq k,$$

where

$$\begin{cases} g^{(n)} = (1 + za_{n,0} + za_0 + z^2 a_0 b_{n-1}) y^{(0)}, \\ c_0 = 0, \ c_n = za_n + z^2 a_0 b_{n-1}, & n \ge 1. \end{cases}$$

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To apply the stability region Theorem we have then to investigate the quantity  $1 - c(\zeta)$  for  $|\zeta| \leq 1$ , and  $c(\zeta) = \sum_{n=0}^{+\infty} c_n \zeta^n$ .

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#### Proposition

The stability region of the Predictor-Corrector method is

$$S = \{ z \in \mathbb{C} \mid 1 - z(\alpha(\zeta) - a_0) - z^2 a_0 \zeta b(\zeta) \neq 0 : |\zeta| \leq 1 \}.$$

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**Proof.** To apply the Theorem we need to prove  $(H_1)$ , we use the binomial series to write

$$(n-1)^p = n^p - pn^{p-1} + \frac{p(p-1)}{2}n^{p-2} + \frac{p(p-1)(p-2)}{6}n^{p-3} + O(n^{p-4}),$$

and similarly for  $(n+1)^p$ , from which we obtain

$$b_n = \frac{1}{\Gamma(\alpha)} n^{\alpha-1} + O(n^{\alpha-2}), \ a_{n,0} = \frac{1}{2\Gamma(\alpha)} n^{\alpha-1} + O(n^{\alpha-2}), \ \alpha_n = \frac{1}{\Gamma(\alpha)} n^{\alpha-1} + O(n^{\alpha-3}),$$

and the expression we need for  $c(\zeta)$  as

$$c(\zeta) = z(\alpha(\zeta) - \alpha_0) + z^2 \alpha_0 \zeta b(\zeta).$$

#### **Proposition**

The stability region of the Predictor-Corrector method is

$$S = \{ z \in \mathbb{C} \mid 1 - z(\alpha(\zeta) - a_0) - z^2 a_0 \zeta b(\zeta) \neq 0 : |\zeta| < 1 \}.$$

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and the expression we need for  $c(\zeta)$  as

$$c(\zeta) = z(\alpha(\zeta) - \alpha_0) + z^2 \alpha_0 \zeta b(\zeta).$$

The expression can be evaluated only numerically.

We have written a predictor-method in an explicit form, we can write and analyze in a similar way also a predictor-corrector made of two *implicit methods*.

- We have now to solve a (possibly) non-linear problem at each step, thus things don't seem to good...
- But we can expect better stability and convergence properties.

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Multiprecision algorithms on specialized hardware can give both an acceleration and maintain the overall accuracy. This idea has already been partially explored for the ODE case, but not yet for FODEs:

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#### Further analyses

One can investigate also stability regions, effects of multiple correction steps, tolerances and step-size selections...

To obtain methods that can be analyzed we can move to Linear Multistep Methods.

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• For an ODE a FLMM with *k* step is a method of the form:

$$\sum_{j=0}^{k} a_j y_{n+j} = \tau \sum_{j=0}^{k} b_j f_{n+j}, \quad n = 0, \dots, s.$$
 (1)

for 
$$t_j = t_0 + j\tau$$
, for  $j = 0, ..., N$ ,  $\tau = (T - t_0)/N$ ,

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- They are associated with the polynomials  $\rho(z) = \sum_{i=0}^k a_i z^i$ ,  $\sigma(z) = \sum_{i=0}^k b_i z^i$ ,
- The fractional version has been introduced in the pioneering work (Lubich 1986b)

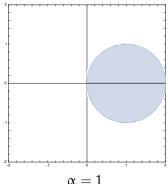
#### Theorem (Lubich 1986b, Theorem 2.6)

Let  $(\rho,\sigma)$  denote an implicit linear multistep method which is stable and consistent of order p. Assume that the zeros of  $\sigma(\zeta)$  have absolute values less than 1. Let  $w(\zeta) = \frac{\sigma(\zeta^{-1})}{\rho(\zeta^{-1})}$  denote the generating power series of the corresponding convolution quadrature  $\omega$ . We define  $\omega^{\alpha} = \{\omega_n^{(\alpha)}\}_{n=0}^{+\infty}$  by  $\omega^{\alpha}(\zeta) = \omega(\zeta)^{\alpha}$ , then the convolution quadrature  $\omega^{\alpha}$  is convergent of order p.

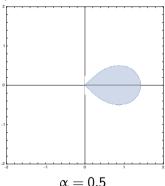
An example is represented by Backward Differentiation Formulas, for which we have

p	$\omega^{\alpha}(\zeta)$
1	$(1-\zeta)^{-\alpha}$
2	$(3/2 - 2\zeta + 1/2\zeta^2)^{-\alpha}$
	$(11/6 - 3\zeta + 3/2\zeta^2 - 1/3\zeta^3)^{-\alpha}$
4	$(25/12 - 4\zeta + 4\zeta^2 - 4/3\zeta^3 + 1/4\zeta^4)^{-\alpha}$
5	$(137/60 - 5\zeta + 5\zeta^2 - 10/3\zeta^3 + 5/4\zeta^4 - 1/5\zeta^5)^{-\alpha}$
6	$(147/60 - 6\zeta + 15/2\zeta^2 - 20/3\zeta^3 + 15/4\zeta^4 - 6/5\zeta^5 + 1/6\zeta^6)^{-\alpha}$

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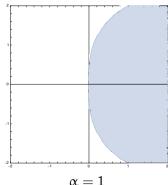


An example is represented by Backward Differentiation Formulas, for which we have

lpha=0.3  $|\zeta|\leq 1\}$  and plot the part

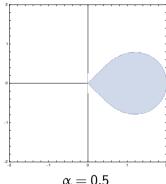
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р	$\omega^{\alpha}(\zeta)$	
	$(1-\zeta)^{-\alpha}$	
2	$(3/2 - 2\zeta + 1/2\zeta^2)^{-\alpha}$	c
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6	$(147/60 - 6\zeta + 15/2\zeta^2 - 20/3\zeta^3 + 15/4\zeta^4 - 6/5\zeta^5 + 1/6\zeta^6)^{-\alpha}$	-2



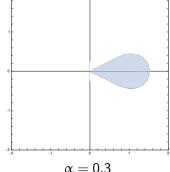
An example is represented by Backward Differentiation Formulas, for which we have

$$\begin{array}{lll} p & \omega^{\alpha}(\zeta) \\ 1 & (1-\zeta)^{-\alpha} \\ 2 & (3/2-2\zeta+1/2\zeta^2)^{-\alpha} \\ 3 & (11/6-3\zeta+3/2\zeta^2-1/3\zeta^3)^{-\alpha} \\ 4 & (25/12-4\zeta+4\zeta^2-4/3\zeta^3+1/4\zeta^4)^{-\alpha} \\ 5 & (137/60-5\zeta+5\zeta^2-10/3\zeta^3+5/4\zeta^4-1/5\zeta^5)^{-\alpha} \\ 6 & (147/60-6\zeta+15/2\zeta^2-20/3\zeta^3+15/4\zeta^4-6/5\zeta^5+1/6\zeta^6)^{-\alpha} \end{array}$$



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	$\omega^{\alpha}(\zeta)$		
1	$(1-\zeta)^{-\alpha}$		
2	$(3/2 - 2\zeta + 1/2\zeta^2)^{-\alpha}$	0	
3	$(11/6 - 3\zeta + 3/2\zeta^2 - 1/3\zeta^3)^{-\alpha}$		
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		-2	2 -1



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### $p \quad \omega^{\alpha}(\zeta)$

$$1 (1-\zeta)^{-\alpha}$$

2 
$$(3/2 - 2\zeta + 1/2\zeta^2)^{-\alpha}$$

3 
$$(11/6 - 3\zeta + 3/2\zeta^2 - 1/3\zeta^3)^{-\alpha}$$

4 
$$(25/12 - 47 + 47^2 - 4/37^3 + 1/47^4)^{-\alpha}$$

5 
$$(137/60 - 5\zeta + 5\zeta^2 - 10/3\zeta^3 + 5/4\zeta^4 - 1/5\zeta^5)^{-\alpha}$$

$$6 \quad (147/60 - 6\zeta + 15/2\zeta^2 - 20/3\zeta^3 + 15/4\zeta^4 - 6/5\zeta^5 + 1/6\zeta^6)^{-\alpha}$$

? How do we obtain the coefficients?

How can we obtain the coefficient describing the method?

$$I_{\tau}^{\alpha}g(t_n) = \tau^{\alpha}\sum_{j=0}^{n}\omega_{n-j}g(t_j) + \tau^{\beta}\sum_{j=0}^{s}w_{n,j}g(t_j),$$

- $\{\omega_j\}_{j=0}^n$  convolution coefficients from  $\omega(\zeta)$ ,
- $\{w_{n,j}\}_{j=0}^k$  starting quadrature weights.

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  - Fast Fourier Transform (FFT) techniques for formal power series,
  - A recursion technique for complex binomial series.
- Solving a small  $k \times k$  Vandermonde system.

Let us suppose that  $\alpha = 1/2$  and that we have a power series of the form

$$\omega(\zeta) = \sum_{j=0}^{+\infty} \omega_j \zeta^j,$$

for which we want to compute for a generic pth degree BDF

$$\omega(\zeta)^{-2} = q(\zeta)$$
 with  $q(\zeta) = \sum_{k=1}^{p} \frac{1}{k} (1-\zeta)^k$ ,

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$$F(\omega(\zeta)) = 0$$
 with  $F(w) = w^{-2} - q(\zeta)$ .

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$$\begin{cases} \omega^{(0)}(\zeta) = \omega_0, \\ \omega^{(m+1)}(\zeta) = \left[ \omega^{(m)}(\zeta) - F'(\omega^{(m)}(\zeta))^{-1} F(\omega^{(m)}(\zeta)) \right]_{2^{m+1}}, \end{cases}$$

for  $[\cdot]_k$  the truncation operator for a power series, i.e.,  $\left[\sum_{j=0}^{+\infty}a_j\zeta^j\right]_k=\sum_{j=0}^ka_j\zeta^j$ , and  $\omega_0$  the solution of  $[F(\omega_0)]_1=0$ .

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$$\begin{cases} \omega^{(0)}(\zeta) = \omega_0 = q(0)^{-1/2}, \\ \omega^{(m+1)}(\zeta) = \left[ \frac{3}{2}\omega^{(m)}(\zeta) - \frac{1}{2}\left(\omega^{(m)}(\zeta)\right)^3 q(\zeta) \right]_{2^{m+1}}, \end{cases}$$

for  $[\cdot]_k$  the truncation operator for a power series, i.e.,  $\left[\sum_{j=0}^{+\infty}a_j\zeta^j\right]_k=\sum_{j=0}^ka_j\zeta^j$ . After m step we have that

$$\omega^{(m)}(\zeta) = [\omega(\zeta)]_{2^m} = \sum_{j=0}^{2^m-1} \omega_j \zeta^j \quad \forall m \geq 0 \ \text{ and cost } O(2^m \log(2^m)).$$

### Recurrence relation

#### Theorem Henrici 1974, Theorem 1.6c, p. 42

Let  $\phi(\zeta) = 1 + \sum_{n=1}^{+\infty} a_n \zeta^n$  be a formal power series. Then for any  $\alpha \in \mathcal{C}$ , we have

$$(\phi(\zeta))^{\alpha} = \sum_{n=0}^{+\infty} v_n^{(\alpha)} \zeta^n,$$

where coefficients  $v_n^{(\alpha)}$  can be evaluated recursively as

$$v_0^{(lpha)}=1, \qquad v_n^{(lpha)}=\sum_{i=1}^n\left(rac{(lpha+1)j}{n}-1
ight)a_jv_{n-j}^{(lpha)}$$

This approach costs an  $O(N^2)$  in general, but can be simplified, e.g., when  $a_1 = \pm 1$ , and  $a_i > 0$  for i > 1 it involves only 2N multiplications and N additions.

### Computing the starting weights

The starting weights  $w_{n,j}$  in

$$I_{\tau}^{\alpha}g(t_n) = \tau^{\alpha}\sum_{j=0}^{n}\omega_{n-j}g(t_j) + \tau^{\beta}\sum_{j=0}^{s}w_{n,j}g(t_j),$$

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#### Starting weight selection

We fix them by imposing that  $I_{\tau}^{\alpha}t^{\gamma}$  is exact for  $\nu \in \mathcal{A} = \mathcal{A}_{p-1} \cup \{p-1\}$  with p the order of convergence of the FLMM, and  $\mathcal{A}_{p-1} = \{\nu \in \mathbb{R} \mid \nu = i+j\alpha, \quad i,j \in \mathbb{N}, \nu < p-1\}.$ 

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$$\tau^{\alpha} \sum_{i=0}^{s} w_{n,j} (jh)^{\gamma} = \frac{1}{\Gamma(\alpha)} \int_{0}^{n\tau} (n\tau - \xi)^{\alpha-1} \chi^{\gamma} d\chi - \tau^{\alpha} \sum_{i=0}^{n} \omega_{n-j} (jh)^{\gamma}, \quad \gamma \in \mathcal{A}.$$

The resulting linear system is of "real" Vandermonde type, i.e.,

$$(A)_{j,\nu_i=1}^s=(jh)^{\nu_i}, \qquad \nu_i\in A, \quad s=|\mathcal{A}|.$$

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- The right-hand side

$$\frac{1}{\Gamma(\alpha)} \int_0^{n\tau} (n\tau - \xi)^{\alpha - 1} \chi^{\nu} d\chi - \tau^{\alpha} \sum_{j=0}^n \omega_{n-j}(jh)^{\nu}$$

can suffer from cancellation of digits!

We know a general way to obtain FLMM methods of the form

$$y^{(n)} = T_{m-1}(t_n) + \tau^{\beta} \sum_{j=0}^{s} w_{n,j} f(t_j, y^{(j)}) + \tau^{\alpha} \sum_{j=0}^{n} \omega_{n-j} f(t_j, y^{(j)}),$$

 $\bigcirc$  starting from the polynomials  $(\rho, \sigma)$  of an implicit order p method,

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- ightharpoons we make the starting values for a multi-step method,
- ightharpoons we still need to discuss how we can efficiently treat the memory term.

To initialize the computation we need the values  $y^{(0)}, \ldots, y^{(s)}, s+1, s=|\mathcal{A}|$ .

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- To avoid mixing methods we evaluate all the approximations at the same time by solving

$$\begin{bmatrix} y^{(1)} \\ y^{(2)} \\ \vdots \\ y^{(s)} \end{bmatrix} = \begin{bmatrix} T_{m-1}(t_1) \\ T_{m-1}(t_2) \\ \vdots \\ T_{m-1}(t_s) \end{bmatrix} + \tau^{\alpha} \begin{bmatrix} (\omega_1 + w_{1,0})f_0 \\ (\omega_2 + w_{2,0})f_0 \\ \vdots \\ (\omega_s + w_{s,0})f_0 \end{bmatrix} + \tau^{\alpha} (\Omega \otimes I + W \otimes I) \begin{bmatrix} f(t_1, y^{(1)}) \\ f(t_2, y^{(2)}) \\ \vdots \\ f(t_s, y^{(s)}) \end{bmatrix}$$

where

$$\Omega = \begin{bmatrix} \omega_0 & & & & \\ \omega_1 & \omega_0 & & & \\ \vdots & \vdots & \ddots & \\ \omega_{s-1} & \omega_{s-2} & \cdots & \omega_0 \end{bmatrix}, \quad W = \begin{bmatrix} w_{1,1} & w_{1,2} & \cdots & w_{1,s} \\ w_{2,1} & w_{2,2} & \cdots & w_{2,s} \\ \vdots & \vdots & \ddots & \vdots \\ w_{s,1} & w_{s,2} & \cdots & w_{s,s} \end{bmatrix}$$

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- This will be in general an  $s \times \dim(y^{(j)})$  nonlinear system that we need to solve before starting the iteration.
- If the value of  $\alpha$  is not very small, viz s is moderate, and the system of ODEs is moderate this is manageable.

### Treating the memory term

If we compute the sum on the coefficients  $\omega_i$  naively for

$$y^{(n)} = T_{m-1}(t_n) + \tau^{\beta} \sum_{j=0}^{s} w_{n,j} f(t_j, y^{(j)}) + \tau^{\alpha} \sum_{j=0}^{n} \omega_{n-j} f(t_j, y^{(j)}),$$

we end up having a  $O(N^2)$  cost! If we do not perform this task efficiently the numerical solution degenerates in an unworkable task as we either refine our grid or enlarge our computational domain.

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- $\blacksquare$  We can consider using a stretched grid towards  $t_0$  to reduce N,
- We can try an approach with nested meshes to reduce the load,
- We can exploit the fact that this is a convolution and adopt some FFT tricks.

The treatment remains the same indifferently for both PI and FLMM method, let us focus here on the generic formulation

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• Let r be a moderate number of step, e.g.,  $r = 2^k$  for a small k, we compute the first step directly

$$y^{(n)} = \phi_n + \sum_{j=0}^n c_{n-j} f_j, \quad n = 0, 1, \dots, r-1.$$

The treatment remains the same indifferently for both PI and FLMM method, let us focus here on the generic formulation

$$y^{(n)} = \phi_n + \sum_{j=0}^n c_{n-j} f_j.$$

• Let r be a moderate number of step, e.g.,  $r = 2^k$  for a small k, we compute the first step directly

$$y^{(n)} = \phi_n + \sum_{i=0}^n c_{n-j} f_j, \quad n = 0, 1, \dots, r-1.$$

• If we now want to compute the next r approximations we can separate the lag term as

$$y^{(n)} = \phi_n + \sum_{j=0}^{r-1} c_{n-j} f_j + \sum_{j=r}^n c_{n-j} f_j, \quad n \in \{r, r+1, \dots, 2r-1\}.$$

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#### We can use FFT!

If we call  $S_r(n,0,r-1) = \sum_{j=0}^{r-1} c_{n-j} f_j$ ,  $n \in \{r,r+1,\ldots,2r-1\}$ , the set of partial sums each of length r we can evaluate them with FFT in  $O(2r \log_2(2r))$ .

• If we now want to compute the next r approximations we can separate the lag term as

$$y^{(n)} = \phi_n + \sum_{j=0}^{r-1} c_{n-j} f_j + \sum_{j=r}^n c_{n-j} f_j, \quad n \in \{r, r+1, \dots, 2r-1\}.$$

 We can apply the same process recursively if we double every time-interval under consideration

$$y^{(n)} = \phi_n + \sum_{j=0}^{2r-1} c_{n-j} f_j + \sum_{j=2r}^{n} c_{n-j} f_j, \quad n \in \{2r, 2r+1, \dots, 3r-1\},$$

$$y^{(n)} = \phi_n + \sum_{j=0}^{2r-1} c_{n-j} f_j + \sum_{j=2r}^{3r-1} c_{n-j} f_j + \sum_{j=3r}^{n} c_{n-j} f_j, \quad n \in \{3r, 3r+1, \dots, 4r-1\},$$

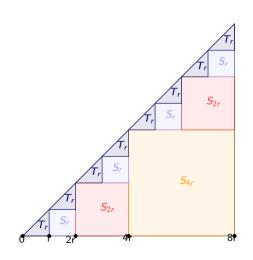
#### 

If we call  $S_{2r}(n,0,2r-1)=\sum_{j=0}^{2r-1}c_{n-j}f_j$ , and  $S_r(n,2r,3r-1)=\sum_{j=2r}^{3r-1}c_{n-j}f_j$  the set of partial sums of lengths 2r and r we can evaluate them with FFT in  $O(4r\log_2(4r))$  and  $O(2r\log_2(2r))$  respectively.

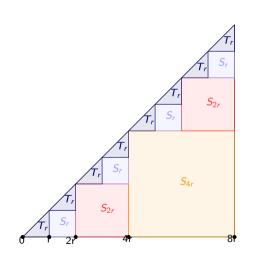
 We can apply the same process recursively if we double every time-interval under consideration

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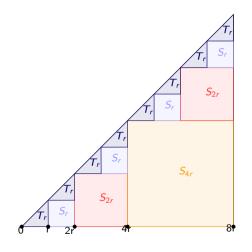


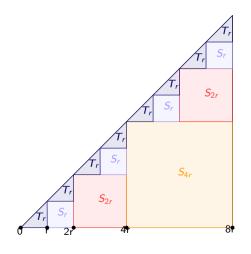
• We can iterate the process for the 4r approximations in the interval  $n \in \{4r, \ldots, 8r-1\}$ , together with the partial sums  $S_{4r}(n, 0, 4r-1)$ ,  $S_{2r}(n, 4r, 6r-1)$ ,  $S_r(n, 6r, 7r-1)$  that can be evaluated in  $O(8r \log_2(8r))$ ,  $O(4r \log_2(4r))$  and  $O(2r \log_2(2r))$  respectively.



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- At each level we have to complete the recursion by computing

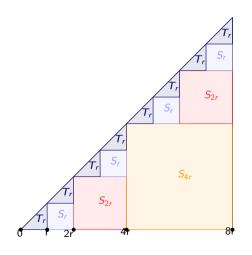
$$T_r(p,n) = \sum_{j=p}^n c_{n-j} f_j, \ p = \ell r,$$
  $n \in \{\ell r, \ell r + 1, \dots, (\ell+1)r - 1\},$   $\ell = 0, 1, 2, \dots$ 



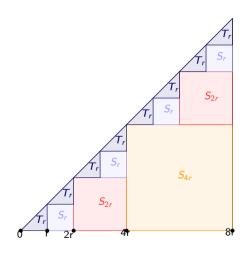


To determine the whole cost we just have to sum the various components

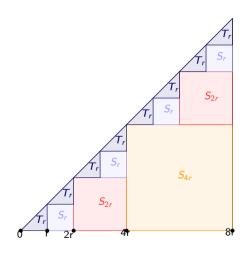
• Assume that  $N = 2^{n_t}$ 



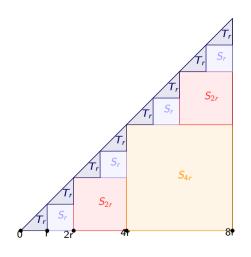
- Assume that  $N = 2^{n_t}$
- $O(N \log_2 N)$  for  $S_{4r}$ ,



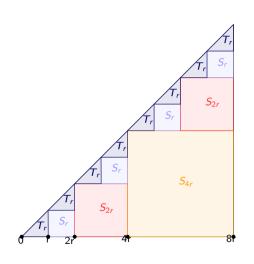
- Assume that  $N = 2^{n_t}$
- $O(N \log_2 N)$  for  $S_{4r}$ ,
- +  $O(N/2 \log_2 N/2)$  for 2  $S_{2r}$



- Assume that  $N = 2^{n_t}$
- $O(N \log_2 N)$  for  $S_{4r}$ ,
- +  $O(N/2 \log_2 N/2)$  for 2  $S_{2r}$
- $+ O(N/4 \log_2 N/4)$  for 4  $S_r$

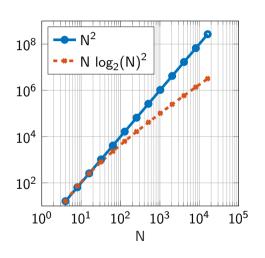


- Assume that  $N = 2^{n_t}$
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- + r(r+1)/2 for the N/r convolutions  $T_r$



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- + r(r+1)/2 for the N/r convolutions  $T_r$
- In general:

$$N \log_2 N + 2 \frac{N}{2} \log_2 \frac{N}{2} + 4 \frac{N}{4} \log_2 \frac{N}{4} + \cdots$$
  
  $+ p \frac{N}{p} \log_2 \frac{N}{p} + \frac{N}{r} \frac{r(r+1)}{2}, \quad p = \frac{N}{2r}$ 



- Assume that  $N = 2^{n_t}$
- $O(N \log_2 N)$  for  $S_{4r}$ ,
- +  $O(N/2 \log_2 N/2)$  for 2  $S_{2r}$
- $+ O(N/4 \log_2 N/4)$  for 4  $S_r$
- + r(r+1)/2 for the N/r convolutions  $T_r$
- In general:

$$\sum_{i=0}^{\log_2 p} N \log_2 \frac{N}{2^j} + N \frac{r+1}{2} = O(N(\log_2 N)^2).$$

We can try to use a "fixed memory length" to reduce the computational (and memory) load.

$$y(t_{n+1}) = y(t_n) + \frac{1}{\Gamma(\alpha)} \int_{t_n}^{t_{n+1}} (t_{n+1} - \tau)^{\alpha - 1} f(\tau, x(\tau)) d\tau + \frac{1}{\Gamma(\alpha)} \int_{0}^{t_n} ((t_{n+1} - \tau)^{\alpha - 1} - (t_n - \tau)^{\alpha - 1}) f(\tau, y(\tau)) d\tau, \quad \alpha \in (0, 1).$$

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$$E = \left| \frac{1}{\Gamma(\alpha)} \int_0^{t_n - T_M} ((t_{n+1} - \tau)^{\alpha - 1} - (t_n - \tau)^{\alpha - 1}) f(\tau, y(\tau)) d\tau \right|$$

$$\leq \frac{M}{\Gamma(\alpha)} \left| \int_0^{t_n - T_M} ((t_{n+1} - \tau)^{\alpha - 1} - (t_n - \tau)^{\alpha - 1}) d\tau \right|$$

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$$= \frac{M}{\alpha \Gamma(\alpha)} \left| (\tau + T_M)^{\alpha} - t_{n+1}^{\alpha} - T^{\alpha} + t_n^{\alpha} \right|$$

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$$= \frac{M}{\Gamma(\alpha)} \left| \int_{T_{M}}^{T_{M} + \tau} z^{\alpha - 1} dz - \int_{t_{n}}^{t_{n+1}} z^{\alpha - 1} dz \right|$$

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$$= \frac{M}{\Gamma(\alpha)} \left| \int_{T_{M}}^{T_{M} + \tau} z^{\alpha - 1} dz - \int_{t_{n}}^{t_{n+1}} z^{\alpha - 1} dz \right| (MeanValueTheorem)$$

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$$\begin{split} E &= \left| \frac{1}{\Gamma(\alpha)} \int_{0}^{t_{n} - T_{M}} ((t_{n+1} - \tau)^{\alpha - 1} - (t_{n} - \tau)^{\alpha - 1}) f(\tau, y(\tau)) d\tau \right| \\ &\leq \frac{M}{\Gamma(\alpha)} \left| \int_{0}^{t_{n} - T_{M}} ((t_{n+1} - \tau)^{\alpha - 1} - (t_{n} - \tau)^{\alpha - 1}) d\tau \right| \\ &= \frac{M}{\Gamma(\alpha)} \left| (z_{1}^{*})^{\alpha - 1} \tau - (z_{2}^{*})^{\alpha - 1} \tau \right| \quad z_{1}^{*} \in [T_{M}, T_{M} + \tau], \ z_{2}^{*} \in [t_{n}, t_{n+1}] \end{split}$$

We can try to use a "fixed memory length" to reduce the computational (and memory) load.

$$y(t_{n+1}) = y(t_n) + \frac{1}{\Gamma(\alpha)} \int_{t_n}^{t_{n+1}} (t_{n+1} - \tau)^{\alpha - 1} f(\tau, x(\tau)) d\tau + \frac{1}{\Gamma(\alpha)} \int_{0}^{t_n} ((t_{n+1} - \tau)^{\alpha - 1} - (t_n - \tau)^{\alpha - 1}) f(\tau, y(\tau)) d\tau, \quad \alpha \in (0, 1).$$

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$$< \frac{M}{\Gamma(\alpha)} T_M^{\alpha - 1} \tau, \qquad \alpha \in (0, 1).$$

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Let us introduce now a fixed window  $T_M$  of memory, then If we have a global error bound  $E_{\rm global}$  with step-length  $\tau$  we just need to choose

$$T_M > \left(\frac{M}{\Gamma(\alpha)E_{\mathsf{global}}}\right)^{1/1-\alpha}, \quad \alpha \in (0,1),$$

while if we have a local error bound  $E_{local}$ 

$$T_M > \left(\frac{M\tau}{\Gamma(\alpha)E_{local}}\right)^{1/1-\alpha}, \quad \alpha \in (0,1).$$

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- $oxed{B}$  Similar bounds can be written for the case  $\alpha > 1$ , that is

$$E < rac{M}{\Gamma(lpha)}(t_{n+1}^{lpha} - T_M^{lpha-1}) au, \quad lpha > 1.$$

But now to preserve the order of accuracy, we must choose

$$T_M^{\alpha-1} > t_{n+1}^{\alpha-1} - \frac{E_{\mathsf{global}}\Gamma(\alpha)}{M}, \quad \alpha > 1,$$

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The idea can be refined by using nested meshes.

Zeroing out the memory term is too drastic, we may want to relax this.

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#### Scaling properties

$$I_{[0,t]}^{\alpha}f(t)=\int_0^t rac{f( au)}{(t- au)^{1-lpha}}\,\mathrm{d} au$$

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#### Scaling properties

$$I_{[0,t]}^{\alpha}f(wt) = \int_0^t \frac{f(\tau)}{(wt-\tau)^{1-\alpha}} d\tau$$

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#### Scaling properties

$$J_{[0,t]}^{\alpha}f(wt) = w^{\alpha} \int_0^t \frac{f(w\tau)}{(t-\tau)^{1-\alpha}} d\tau$$

## **Nested meshes**

Zeroing out the memory term is too drastic, we may want to relax this.

#### Scaling properties

Given  $p \in \mathbb{N}$  we then have

$$I_{[0,t]}^{\alpha} f(w^p t) = w^{p\alpha} \int_0^t \frac{f(w^p \tau)}{(t-\tau)^{1-\alpha}} d\tau$$

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Given  $p \in \mathbb{N}$  we then have

$$I_{[0,t]}^{\alpha}f(w^pt)=w^{p\alpha}\int_0^t \frac{f(w^p au)}{(t- au)^{1-lpha}}\,\mathrm{d} au$$

• We can use the weight on the mesh

$$\Omega_{\tau}^{\alpha}f(\emph{n}\tau) pprox \emph{I}_{[0,t]}^{\alpha}f(\emph{n}\tau),$$
 step length  $\tau$ 

to compute

$$\Omega^{\alpha}_{w^{\rho}\tau}f(w^{\rho}n\tau)\approx \textit{I}^{\alpha}_{[0,t]}f(nw^{\rho}\tau), \text{ step length } w^{\rho}\tau$$

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• We can use the weight on the mesh

$$\Omega_{\tau}^{\alpha} f(n\tau) \approx I_{[0,t]}^{\alpha} f(n\tau)$$
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to compute

$$\Omega_{w^p\tau}^{\alpha}f(w^pn\tau)\approx I_{[0,t]}^{\alpha}f(nw^p\tau), \text{ step length } w^p\tau$$

ullet In summary for any  $p\in\mathbb{N}$  we get

$$\Omega^{\alpha}_{\tau}f(n\tau) = \sum_{i=0}^{n} \omega_{n-j}f(j\tau) \iff \Omega^{\alpha}_{w^{p}\tau}f(nw^{p}\tau) = w^{p\alpha}\sum_{i=0}^{n} \omega_{n-j}f(jw^{p}\tau).$$

#### Nested mesh

Given  $\tau \in \mathbb{R}^+$ , the mesh  $M_{\tau} = \{\tau n, \ n \in \mathbb{N}\}$ . Selected  $w, r, p \in \mathbb{N}$ , w > 0, r > p, we have  $M_{w^p\tau} \supset M_{w^r\tau}$  and we decompose the interval as

$$[0,t] = [0,t-w^mT] \cup [t-w^mT,t-w^{m-1}T] \cup \cdots \cup [t-wT,t-T] \cup [t-T,t]$$

for  $m \in \mathbb{N}$  the smallest integer such that  $t < w^{m+1}T$ .

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for  $m \in \mathbb{N}$  the smallest integer such that  $t < w^{m+1}T$ .

 $\P$  This links the **scaling property** with the singularity of the type  $1/(t-\tau)^{1-\alpha}$  suggesting that we should distribute the computational effort logarithmically, and not uniformly.

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- $\P$  This links the **scaling property** with the singularity of the type  $1/(t-\tau)^{1-\alpha}$  suggesting that we should distribute the computational effort logarithmically, and not uniformly.
- We rewrite our integral as

$$I_{[0,t]}^{\alpha}f(t) = I_{[t-T,t]}^{\alpha}f(t) + \sum_{i=0}^{m-1} I_{[t-w^{i+1}T,t-w^{i}T]}^{\alpha}f(t) + I_{[0,t-w^{m}T]}^{\alpha}f(t)$$

#### Nested mesh

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$$[0,t] = [0,t-w^mT] \cup [t-w^mT,t-w^{m-1}T] \cup \cdots \cup [t-wT,t-T] \cup [t-T,t]$$

for  $m \in \mathbb{N}$  the smallest integer such that  $t < w^{m+1}T$ .

- $\P$  This links the **scaling property** with the singularity of the type  $1/(t-\tau)^{1-\alpha}$  suggesting that we should distribute the computational effort logarithmically, and not uniformly.
- We rewrite our integral using the scaling property as

$$I_{[0,t]}^{\alpha}f(t)=I_{[t-T,t]}^{\alpha}f(t)+\sum_{i=0}^{m-1}w^{i\alpha}I_{[t-wT,t-T]}^{\alpha}f(w^{i}t)+w^{m\alpha}I_{[0,t-T]}^{\alpha}f(w^{m}t).$$

In the discrete approximation of

$$I_{[0,t]}^{\alpha}f(t)=I_{[t-T,t]}^{\alpha}f(t)+\sum_{i=0}^{m-1}w^{i\alpha}I_{[t-wT,t-T]}^{\alpha}f(w^{i}t)+w^{m\alpha}I_{[0,t-T]}^{\alpha}f(w^{m}t).$$

we approximate

$$\Omega^{\alpha}_{\tau,[t-w^{i+1}T,t-w^{i}T]}f(t)\approx\Omega^{\alpha}_{w^{i}\tau,[t-w^{i+1}T,t-w^{i}T]}f(t)$$

and substitute

$$w^{i\alpha}\Omega^{\alpha}_{\tau,[t-wT,t-T]}f(t)=\Omega^{\alpha}_{w^i\tau,[t-w^{i+1}T,t-w^iT]}f(t).$$

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## Theorem (Ford and Simpson 2001, Theorem 1)

The nested mesh scheme preserves the order of the underlying quadrature rule on which it is based.

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### Theorem (Ford and Simpson 2001, Theorem 1)

The nested mesh scheme preserves the order of the underlying quadrature rule on which it is based.

**Proof.** For integration over a fixed interval [0, t] the choice of T fixes (independent of h) the number of subranges over which the integral is evaluated, on each of them we have en error  $O(h^p)$ .

• The first benefit is that we evaluate a fixed number of quadrature coefficients and then re-use them on all successive intervals,

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- The first benefit is that we evaluate a fixed number of quadrature coefficients and then re-use them on all successive intervals.
- This approach cost  $O(w^m)$  with respect to  $O(w^{2m})$  of the full method,
- We could use linear extrapolation techniques to improve the results.
- Selecting the various parameter may need a bit of tuning.

# Available codes

With respect to the ordinary case for which there exists many reliable and high-performance codes, the choices for computing the solution of fractional differential equation is much more *sparse*.

- From (Garrappa 2018)
  - FDE\_PI1\_Ex.m Explicit Product-Integration of rectanguar type
  - FDE\_PI1\_Im.m Implicit Product-Integration of rectanguar type
  - FDE\_PI2\_Im.m Implicit Product-Integration of trapezoidal type
  - FDE\_PI12\_PC.m Product-Integration with predictor-corrector
- From (Garrappa 2015)
  - FLMM2 Matlab code Three implicit second order Fractional Linear Multistep Methods.

### A remark

All these methods use direct-solver for the Newton method inside them, there is space to make improvement on the solution strategies. Furthermore, a challenge that yet remains: can we find a strategy that combines the convolution features and savings on the memory?

## **Conclusions**

We know a general way to obtain FLMM methods of the form

$$y^{(n)} = T_{m-1}(t_n) + \tau^{\beta} \sum_{j=0}^{s} w_{n,j} f(t_j, y^{(j)}) + \tau^{\alpha} \sum_{j=0}^{n} \omega_{n-j} f(t_j, y^{(j)}),$$

- $\bigcirc$  starting from the polynomials  $(\rho, \sigma)$  of an implicit order  $\rho$  method,
- $\bigcirc$  we have seen how to compute the convolution coefficients  $\omega_n$ ,
- $\bigcirc$  we have seen how to compute the starting nodes  $w_{n,j}$ ,
- we know how we can compute the starting values for a multi-step method by solving a nonlinear system with Newton,
- we have some hints on how we can efficiently treat the memory term.

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- $\bigcirc$  starting from the polynomials  $(\rho, \sigma)$  of an implicit order p method,
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- we know how we can compute the starting values for a multi-step method by solving a nonlinear system with Newton,
- we have some hints on how we can efficiently treat the memory term.

## Things we didn't pack...

There have been some effort in devising methods of Runge-Kutta type (Fischer 2019), collocation methods, and exponential type integrators.

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