CMSC 330: Organization of Programming Languages

Functional Programming in Object-Oriented Languages

Iteration

· Goal: Loop through all objects in an aggregate

```
class Node { Element elt; Node next; }
Node n = ...;
while (n != null) { ...; n = n.next; }
```

- Problems:
 - Depends on implementation details
 - Varies from one aggregate to another

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Iterators in Java

```
public interface Iterator<A> {
    // returns true if the iteration has more elements
    public boolean hasNext();

    // returns the next element in the iteration
    public A next() throws NoSuchElementException;
}
```

Advantages

- The implementation of iterators is not exposed
- Generic for many different aggregates
- Supports multiple traversal strategies
- In Java, iterators can be used explicitly, or implicitly via the enhanced for (also called foreach) loop

Writing an "Iterator" in OCaml

```
let iterate 1 =
  let cur_list = ref 1 in
  ((fun () -> (!cur_list) != []),
    (fun () ->
    let temp = List.hd (!cur_list) in
        cur_list := List.tl (!cur_list);
        temp))
```

```
# let iterate 1 = ...;;
val iterate : 'a list -> (unit -> bool) * (unit -> 'a) = <fun>
# let (has_next, next) = iterate [1; 2; 3];;
val has_next : unit -> bool = <fun>
val next : unit -> int = <fun>
```

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Relating Objects and Closures

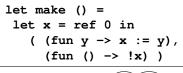
- An object...
 - Is a collection of fields (data)
 - ...and methods (code)
 - When a method is invoked, it is passed an implicit this parameter it can use to access fields
- A closure...

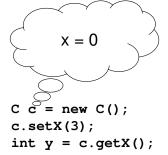
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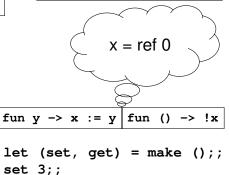
- Is a pointer to an environment (data)
- ...and a function body (code)
- When a closure is invoked, it is passed its environment it can use to access variables

Relating Objects and Closures (cont'd)

```
class C {
  int x = 0;
  void setX(int y) { x = y; }
  int getX() { return x; }
}
```







let y = get();

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Encoding Objects in General

We can apply this transformation in general

```
class C { f_1 \ldots f_m; m_1 \ldots m_n;  }
```

- becomes

```
let make () =
  let f<sub>1</sub> = ... in
  ...
  let f<sub>m</sub> = ... in
  ( fun ... , (* body of m<sub>1</sub> *)
   ...
  fun ..., (* body of m<sub>n</sub> *)
  )
```

- make () is like the constructor
- the closure environment contains the fields

Recall a Useful Higher-Order Function

```
let rec map f = function
[] -> []
| (h::t) -> (f h)::(map f t)
```

Can we encode these in Java?

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A Map Method for Stack

- To write a map method, we need some way of passing a function into another function
 - We can do that with an object with a known method

```
public interface Function {
  Object eval(Object arg);
}
```

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A Map Method for Stack, con't.

 Here are two classes which both implement this Function interface:

```
class AddOne implements Function {
  Object eval(Object arg) {
    return new Integer(((Integer) arg) + 1);
  }
}
```

```
class MultTwo implements Function {
  Object eval(Object arg) {
    return new Integer(((Integer) arg) * 2);
  }
}
```

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A Map Method for Stack, con't.

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```
class Stack { /* ignore type parameter for now... */
    ...
    private Entry theStack;
    Stack map(Function f) {
        Stack s = new Stack();
        map_helper(f, s, theStack);
        return s;
    }
    void map_helper(Function f, Stack s, Entry e) {
        if (e != null) {
            map_helper(f, s, e.next);
            s.push(f.eval(e.elt));
        }
}
```

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A Map Method for Stack, con't.

• Then to apply the function, we just do

```
Stack s = ...;
Stack t = s.map(new AddOne());
Stack u = s.map(new MultTwo());
```

- We make a new object that has a method that performs the function we want
- This is sometimes called a *callback*, because map "calls back" to the object passed into it
- But it's really just a higher-order function, written more awkwardly

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Relating Closures and Objects

```
a = 3

fun b -> a + b
```

```
let add a b = a + b;;
let f = add 3;;
app f 4;;
```

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```
interface F {
   Object eval(Object y);
}
class C {
   static Object app(F f, Object x) {
     return f.eval(x);
   }
}
class G implements F {
   int a;
   G(int a) { this.a = a; }
   Object eval(Object y) {
     return a + (Integer) y;
   }
}
F adder = new G(3);
C.app(adder, 4);
   a = 3
```

Encoding Functions with Objects

We can apply this transformation in general

```
...(fun x -> (* body of f<sub>n</sub> *)) ...

let h f ... = ...f y...

- becomes
```

```
interface F { Object eval(Object x); }
class G implements F {
   Object eval(Object x) { /* body of f<sub>n</sub> */ }
}
class C {
   Typ h(F f, ...) {
        ...f.eval(y)...
   }
}
```

- F is the interface to the callback
- G represents the particular function

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Code as Data

- The key insight in all of these examples is to treat code as if it were data
 - Higher-order functions allow code to be passed around the program
 - As does object-oriented programming
- This is a powerful programming technique
 - And it can solve a number of problems quite elegantly
- Closures and objects are related
 - Both of them allow data to be associated with higherorder code as its passed around (but we can even get by without this)

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