

$$1. \text{ first}(S) = \text{first}(Adg) \cup \text{first}(Bh) = \{a, d, f, g\}$$

$$\text{first}(A) = \text{first}(aBc) \cup \text{first}(\epsilon) = \{a, \epsilon\}$$

$$\text{first}(B) = \text{first}(f) \cup \text{first}(f) = \{f, g\}$$

$$2. \text{ a. } S \rightarrow bL \\ L \rightarrow +aL \mid +bL \mid \epsilon$$

$$\text{ b. } S \rightarrow cL \\ L \rightarrow +M \mid \epsilon \\ M \rightarrow aL \mid bL$$

$$\text{ c. } S \rightarrow aL \\ L \rightarrow bc \mid c$$

$$\text{ d. } S \rightarrow aL \\ L \rightarrow a \mid b \mid \epsilon$$

$$\text{ e. } S \rightarrow aSL \mid b \\ L \rightarrow c \mid b$$

$$3. \text{ a. } \text{first}(\text{true}) = \{\text{true}\} \\ \text{first}(\text{false}) = \{\text{false}\} \\ \text{first}((S)) = \{()\} \\ \text{first}(S \text{ and } S) = \{(\text{, true, false})\} \\ \text{first}(S \text{ or } S) = \{(\text{, true, false})\}$$

b. The first sets of the productions intersect and the grammar is left recursive.

$$4. \text{ a. } \text{first}(abS) = \{a\} \\ \text{first}(acS) = \{a\} \\ \text{first}(c) = \{c\} \\ \text{first}(S) = \{a, c\}$$

b. The first sets of the productions overlap:
 $\text{first}(abS) \cap \text{first}(acS) = \{a\} \cap \{a\} \neq \emptyset$.

$$\text{ c. } S \rightarrow aL \mid c \\ L \rightarrow bS \mid cS$$

```
d. parse_S() {
    if (lookahead == "a") {
        match("a"); // S -> aL
        parse_L();
    } else
        if (lookahead == "c") {
            match("c"); // S -> c
        } else error();
}

parse_L() {
    if (lookahead == "b") {
        match("b"); // L -> bS
        parse_S();
    } else
        if (lookahead == "c") {
            match("c"); // L -> cS
            parse_S();
        } else error();
}
```