

# COMP2521 19T0

## Week 4, Thursday: Graphic Content (III)!

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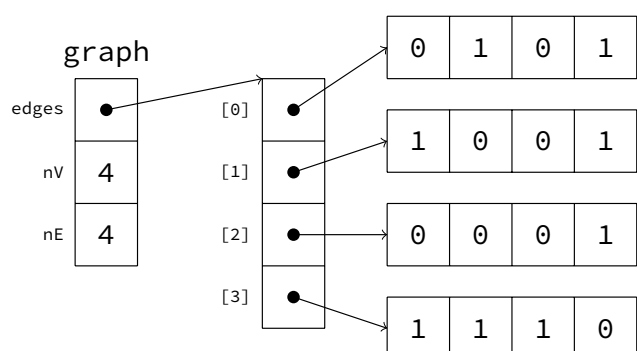
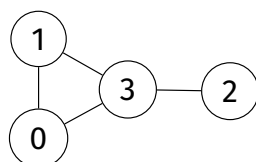
computability  
directed graphs

# Graph Representation

## Recap: Ways of Representing Graphs

### Adjacency Matrices

```
struct graph {  
    size_t nV, nE;  
    bool **matrix;  
};
```



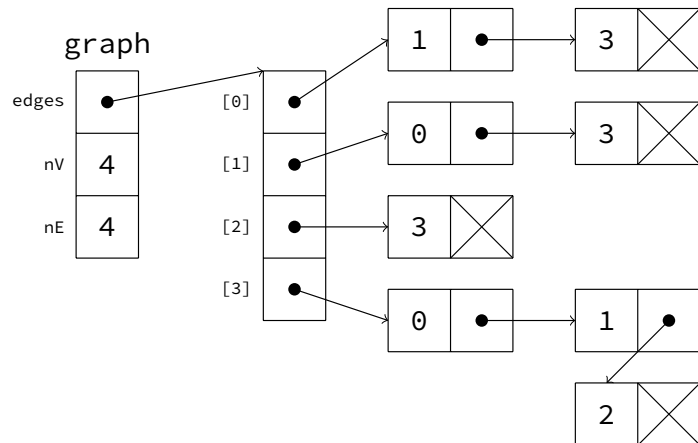
# Recap: Ways of Representing Graphs

Adjacency Lists

Graph Rep.

Computability

```
typedef
    struct adjnode
    adjnode;
struct graph {
    size_t nV, nE;
    adjnode **edges;
};
struct adjnode {
    vertex w;
    adjnode *next;
};
```



# Recap: Ways of Representing Graphs

Edge Lists

Graph Rep.

Computability

```
struct graph {
    size_t nV, nE;
    edge *edges;
};
```

$v$	$w$
0	1
0	3
1	3
2	3

# Recap: Ways of Representing Graphs

Linked Data Structure

Graph Rep.

Computability

```
typedef struct vertex {
    Item it;
    size_t degree;
    vertex *neighbours;
} vertex;

struct graph {
    size_t nV, nE;
    vertex *root;
};
```

	matrix	adj.list	edge list	node links
space	$V^2$	$V + E$	$E$	$V + E$
initialise	$V^2$	$V$	1	$V$
destroy	$V$	$E$	$E$	$V + E$
insert edge	1	$V$	1	$E$
find/remove edge	1	$V$	$E$	$E$
is isolated?	$V$	1	$E$	1
degree	$V$	$E$	$E$	$E$
is adjacent?	1	$V$	$E$	$E$

## Hamilton Paths and Tours

### Hamilton Path:

a simple path connecting two vertices  
that visits each vertex in the graph exactly once

### Hamilton Tour:

a cycle  
that visits each vertex in the graph exactly once

\* \* \*

Given a list of vertices or edges, easy to check.  
Given a graph ... how do we know if one exists?  
Do we have to find one? If so, how do we find one?

## Hamilton Paths and Tours

Brute Force Is Best Force

### IDEA brute force!

enumerate every possible path, and check each one.

hack a BFS or DFS to do it:  
keep a counter of vertices visited in the current path;  
only accept a path only if count is  
equal to the number of vertices.

### PROBLEM how many paths?

given a simple path:  
no path from  $t$  to  $w$  implies no path from  $v$  to  $w$  via  $t...$   
so there's no point visiting a vertex twice on a simple search  
... but that's not true for a Hamilton path!

we must inspect every possible path in the graph.  
in a complete graph, we have  $V!$  different paths ( $\approx (V/e)^V$ )

there are well-known, well-defined subsets of this problem  
which are easy to solve (Dirac, Ore) ... but in general  
this is a **non-deterministic polynomial**, or NP problem

## Euler Paths and Tours

### Euler Path:

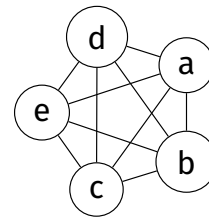
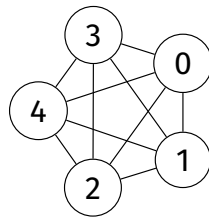
a simple path connecting two vertices  
that visits each edge in the graph exactly once  
... exists iff the graph is connected  
and has exactly two vertices of odd degree

### Euler Tour:

a cycle  
that visits each edge in the graph exactly once ... exists iff the graph is  
connected  
and all vertices are of even degree  
... these can be found in linear time.

## Graph Problems Tractable and Intractable

- tractable: can we find a simple path connecting two vertices in a graph?
- tractable: what's the shortest such path?
- intractable**: what's the longest such path?
- tractable: is there a clique in a given graph?
- intractable**: what's the largest clique?
- tractable: given two colours, can we colour every vertex in a graph  
such that no two adjacent vertices are the same colour?
- intractable**: what about three colours?



Graph isomorphism:

Can we make two given graphs identical by renaming vertices?

No general solution exists.  
We don't know if one can exist.



IN CS, IT CAN BE HARD TO EXPLAIN  
THE DIFFERENCE BETWEEN THE EASY  
AND THE VIRTUALLY IMPOSSIBLE.

xkcd 1425 "Tasks" // CC BY-NC 2.5