

COMP2521 19T0

Week 3, Thursday: Graphic Content (II)!

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graph representation
graph search

Graph Representation

What do we need to represent?

A graph G is a set of vertices $V := \{v_1, \dots, v_n\}$,
and a set of edges $E := \{(v, w) \mid v, w \in V; (v, w) \in V \times V\}$.

Directed graphs: $(v, w) \neq (w, v)$.

Weighted graphs: $E := \{(v, w, \sigma)\}$.

Multigraphs: E is a list, not a set.

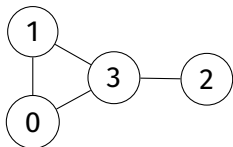
What operations do we need to support?

create/destroy graph;

add/remove vertices, edges;

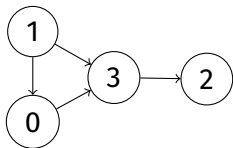
get #vertices, #edges;

A $|V| \times |V|$ matrix; each cell represents an edge.



0	1	0	1
1	0	0	1
0	0	0	1
1	1	1	0

undirected



0	0	0	1
1	0	0	1
0	0	0	0
0	0	1	0

directed

Advantages

- Easy to implement!
two-dimensional array of
`bool/int/float/...`
- Works for:
graphs! digraphs!
weighted graphs!
(unweighted) multigraphs!
- Efficient!
 $O(1)$ edge-insert, edge-delete
 $O(1)$ is-adjacent

Disadvantages

- Huge space overheads!
 V^2 cells of some type
sparse graph \Rightarrow wasted space!
undirected graph \Rightarrow wasted space!
- Inefficient!
 $O(V^2)$ initialisation
 $O(V^2)$ vertex-insert/-delete

Graph Rep.

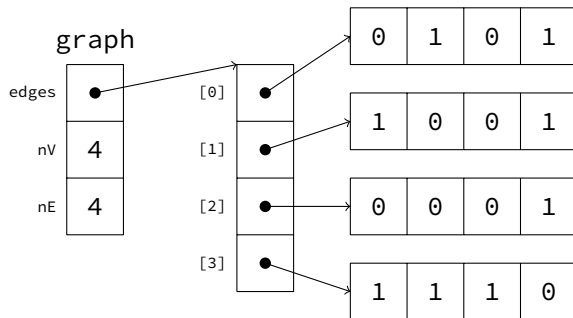
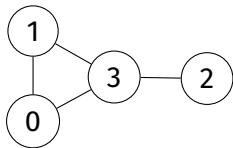
Adj. Matrix

Adj. List

Graph ADT

Graph Search

```
struct graph {  
    size_t nV, nE;  
    bool **matrix;  
};
```

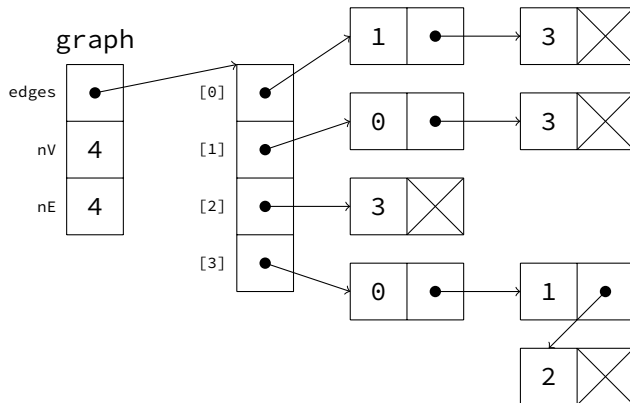


Exercise: Time Complexity

Given an adjacency matrix representation,
find the time complexity, and implement, these functions

- `bool graph_adjacent (Graph g, vertex v, vertex w);`
... returns true if vertices v and w are connected, false otherwise
- `size_t graph_degree (Graph g, vertex v);`
... return the degree of a vertex v

```
typedef
    struct adjnode
    adjnode;
struct graph {
    size_t nV, nE;
    adjnode **edges;
};
struct adjnode {
    vertex w;
    adjnode *next;
};
```



- Space: matrix: V^2 ; adjlist: $V + E$
- Initialise: matrix: V^2 , adjlist: V
- Destroy: matrix: V , adjlist: E
- Insert edge: matrix 1, adjlist: V
- Find/remove edge: matrix: 1, adjlist: V
- is isolated? matrix: V , adjlist: 1
- Degree: matrix: V , adjlist: E
- is adjacent? matrix: 1, adjlist: V

What do we need to represent?
What operations do we need to support?
What behaviours are we trying to model?
How do we interact with other types?

```
typedef struct graph *Graph;
```

Graph Rep.

Adj. Matrix

Adj. List

Graph ADT

Graph Search

```
typedef struct graph *Graph;
```

```
/** A concrete edge type. */
```

```
typedef struct edge { vertex v, w; weight n; } edge;
```

```
typedef struct graph *Graph;

/** A concrete edge type. */
typedef struct edge { vertex v, w; weight n; } edge;

/** Create a new instance of a Graph. */
Graph graph_new (
    size_t max_edges,          /**< maximum value hint */
    size_t max_vertices,      /**< maximum value hint */
    bool directed,             /**< true if a digraph */
    bool weighted              /**< true if edges have weight */
);

/** Deallocate resources used by a Graph. */
void graph_drop (Graph g);
```

Graph Rep.

Adj. Matrix

Adj. List

Graph ADT

Graph Search

```
/** Get the number of vertices in this Graph. */  
size_t graph_num_vertices (Graph g);
```

```
/** Get the number of edges in this Graph. */  
size_t graph_num_edges (Graph g);
```

Graph Rep.

Adj. Matrix

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Graph ADT

Graph Search

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```

```
/** Get the number of edges in this Graph. */  
size_t graph_num_edges (Graph g);
```

```
/** Is this graph directed? */  
bool graph_directed_p (Graph g);
```

```
/** Is this graph weighted? */  
bool graph_weighted_p (Graph g);
```

Graph Rep.

Adj. Matrix

Adj. List

Graph ADT

Graph Search

```
/** Add vertex with index `v' to the Graph.  
 * If the vertex already exists, a no-op returning false. */  
bool graph_vertex_add (Graph g, vertex v);
```


Graph Rep.

Adj. Matrix

Adj. List

Graph ADT

Graph Search

```
/** Add vertex with index `v' to the Graph.  
 * If the vertex already exists, a no-op returning false. */  
bool graph_vertex_add (Graph g, vertex v);  
  
/** Add edge `e', from `v' to `w' with weight `n', to the Graph.  
 * If the edge already exists, a no-op returning false. */  
bool graph_edge_add (Graph g, edge e);
```

Graph Rep.

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void graph_edge_remove (Graph g, edge e);
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/** Remove vertex `v' from the Graph. */  
void graph_vertex_remove (Graph g, vertex v);
```

Graph Rep.

Adj. Matrix

Adj. List

Graph ADT

Graph Search

```
/** Does this Graph have this vertex? */  
bool graph_has_vertex_p (Graph g, vertex v);
```

Graph Rep.

Adj. Matrix

Adj. List

Graph ADT

Graph Search

```
/** Does this Graph have this vertex? */  
bool graph_has_vertex_p (Graph g, vertex v);  
  
/** What is the degree of this vertex on this Graph? */  
size_t graph_vertex_degree (Graph g, vertex v);
```

Graph Rep.

Adj. Matrix

Adj. List

Graph ADT

Graph Search

```
/** Does this Graph have this vertex? */  
bool graph_has_vertex_p (Graph g, vertex v);  
  
/** What is the degree of this vertex on this Graph? */  
size_t graph_vertex_degree (Graph g, vertex v);  
  
/** Does this Graph have this edge? */  
bool graph_has_edge_p (Graph g, edge e);
```

Graph Search

We learn properties of a graph by
systematically examining
each of its edges and vertices —

... to compute the degree of all vertices,
we visit each vertex, and count its edges

... for path-related properties
we move from vertex to vertex along edges
choosing edges as we go

we implement general graph-search algorithms
which can solve a wide range of graph problems

PROBLEM

does a path exist between vertices v and w ?

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- examine vertices adjacent to v ;
- if any of them is w , we're done!
- otherwise, check from all of the adjacent vertices
... rinse and repeat moving away from v

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does a path exist between vertices v and w ?

- examine vertices adjacent to v ;
- if any of them is w , we're done!
- otherwise, check from all of the adjacent vertices
... rinse and repeat moving away from v

What order do we visit nodes in?

'Breadth-first' (**BFS**): adjacent nodes first

'Depth-first' (**DFS**): longest paths first

Dijkstra: lowest-cost paths first

'Greedy Best-First' (**GBFS**): shortest-heuristic-distance

A*: lowest-cost *and* shortest-heuristic-distance

Path searches on graphs tend to follow a simple pattern:

- create a structure that will tell us what next
- add the starting node to that structure
- while that structure isn't empty:
 - get the next vertex from that structure;
 - mark that vertex as visited; and
 - add its neighbours to the structure

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What data structure should we use?

BFS:

DFS:

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BFS: a queue!

DFS:

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What data structure should we use?

BFS: a queue!

DFS: a stack!

`count` number of vertices traversed so far

`pre[]` order in which vertices were visited (for 'pre-order')

`st[]` predecessor of each vertex (for 'spanning tree')

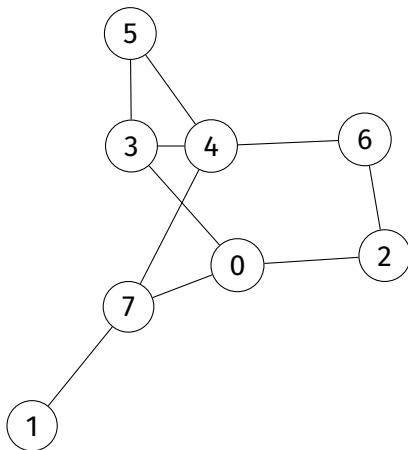
`count` number of vertices traversed so far

`pre[]` order in which vertices were visited (for 'pre-order')

`st[]` predecessor of each vertex (for 'spanning tree')

the edges traversed in all graph walks form a **spanning tree**, which has —

- has edges corresponding to call-tree of recursive function
- is the original graph sans cycles/alternate paths
- (in general, a spanning tree has all vertices and a minimal set of edges to produce a connected graph; no loops, cycles, parallel edges)



If a graph is not connected,
DFS will produce
a spanning forest

An edge connecting a vertex with
an ancestor in the DFS tree
that is not its parent is a back edge

Depth-First Search

... recursively, using the call stack (I)

```
void dfsR (Graph g, edge e) {  
    // ... set up `pre' array of `g->nV' items set to -1  
    // ... set up `st' array of `g->nV' items set to -1  
    // ... set up `count' = 0  
    pre[w] = count++;  
    st[w] = e.v;  
    vertex w = e.w;  
    for (vertex i = 0; i < g->nV; i++)  
        if (g->edges[w][i] && pre[i] == -1)  
            dfsR (g, (edge){.v = w, .w = i});  
}
```

Depth-First Search

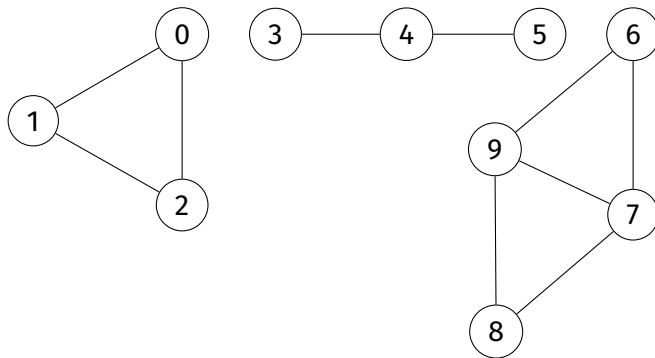
... recursively, using the call stack (II)

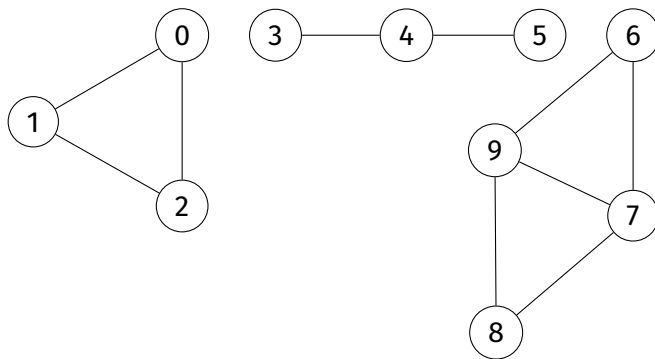
Graph Rep.

Graph Search

DFS

BFS



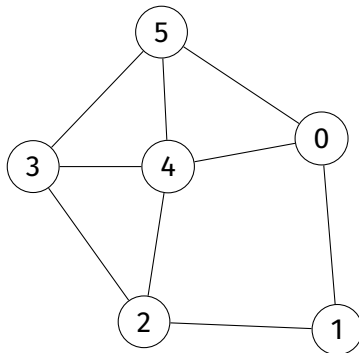


How can we ensure that all vertices are visited?

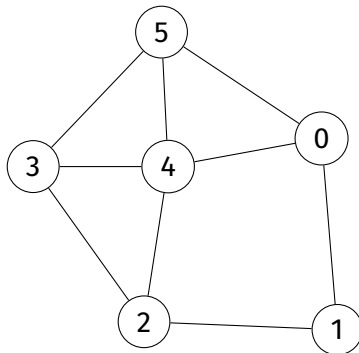
Depth-First Search

... recursively, using the call stack (III)

```
void dfs (Graph g)
{
    count = 0;
    pre = calloc (g->nV * sizeof (int));
    st = calloc (g->nV * sizeof (int));
    for (vertex v = 0; v < g->nV; v++)
        pre[v] = st[v] = 1;
    for (vertex v = 0; v < g->nV; v++)
        if (pre[v] == -1)
            dfsR (g, (edge){.v = v, .w = v});
}
```

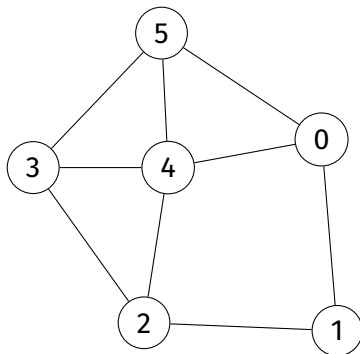


Let's do a DFS!



Let's do a DFS!

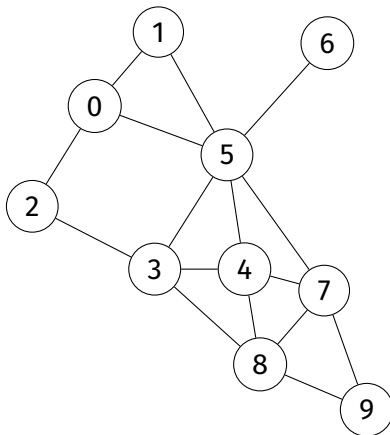
Does a path exist from $0 \dots 5$?



Let's do a DFS!

Does a path exist from 0...5?

Yes: 0, 1, 2, 3, 4, 5.



Let's do a DFS!
What do `pre[]` and `st[]` look like?

```
void dfs (Graph g, edge e)
{
    // ... set up `pre' array of `g->nV' items set to -1
    // ... set up `st' array of `g->nV' items set to -1
    // ... set up `count' = 0
    Stack s = stack_new ();
    stack_push (s, e);
    while (stack_size (s) > 0) {
        e = stack_pop (s);
        if (pre[e.w] != -1) continue;
        pre[e.w] = count++; st[e.w] = e.v;
        for (int i = 0; i < g->nV; i++)
            if (has_edge (e.w, i) && pre[i] == -1)
                stack_push (s, (edge){.v = e.w, .w = i });
    }
}
```

```
void bfs (Graph g, edge e)
{
    // ... set up `pre' array of `g->nV' items set to -1
    // ... set up `st' array of `g->nV' items set to -1
    // ... set up `count' = 0
    Queue q = queue_new ();
    queue_en (q, e);
    while (queue_size (q) > 0) {
        e = queue_de (q);
        if (pre[e.w] != -1) continue;
        pre[e.w] = count++; st[e.w] = e.v;
        for (int i = 0; i < g->nV; i++)
            if (has_edge (e.w, i) && pre[i] == -1)
                queue_en (q, (edge){.v = e.w, .w = i });
    }
}
```